

IGPSS

A General Purpose Simulation

Language for the Interdata

by

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A large, stylized handwritten signature in black ink, appearing to read 'Fred Maryanski', is written over a horizontal line.

Fred Maryanski - Major Professor

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INTRODUCTION

The purpose of this project is to develop a simulation language which is portable to different kinds of hardware systems. The language developed is called IGPSS and resembles IBM's General Purpose Simulation System. The program is written in Fortran since Fortran is available on most systems and is fairly well standardized. Although developed for use primarily on the Interdata 8/32, IGPSS is also implemented on the IBM/370.

This report will provide a description of the IGPSS internal organization and its user requirements. The user is assumed to have a knowledge of simulation.

The parser for IGPSS remains to be written, so the current input required for IGPSS is the anticipated output from the parser. The special requirements for this input are discussed in the section entitled "Coding IGPSS."

General Structure

There are three methods of building a simulation model. The event scheduling approach emphasizes a detailed description of the steps that occur when an individual event takes place. Each type of event has a distinct set of steps associated with it. The activity scanning approach emphasizes a review of all activities in a simulation to determine which can be begun or terminated each time an event occurs. The process interactive approach emphasizes the progress of an entity through a system from its arrival event to its departure event. (1)

IGPSS is a simulation language based on GPSS. It uses the process interactive approach to organizing behavior in a simulation. In IGPSS active objects called transactions act on a system of passive objects which include facilities and storages. For example, a job requiring service is a transaction, whereas a single server is a facility and a group of servers is a storage unit. The transaction passes through the system, encountering a facility or storage, depending on whether the simulation is a single or multiserver representation.

The program maintains a record of when each transaction in the system is due to move. The transactions move through the system on a first-come, first-served basis within a priority class. The program moves a transaction through the system until a blocking condition arises. When a transaction is moved as far as it can go, the program moves any other transaction due to move at the same time. If all movements are completed, the program advances the clock to the time of the next most imminent event and repeats the movement process.

A system to be simulated in IGPSS can be described as a block diagram in which the blocks represent activities and the lines joining the blocks indicate the sequence in which the events are to occur. If a choice of

activities can occur depending on a condition, more than one line leaves a block and the condition for the choice is stated at the block. The sequence of events in real time is reflected in the movement of transactions from block to block in simulated time.

Transactions are created at one or more GENERATE blocks and are removed from the simulation at TERMINATE blocks. There can be many transactions moving simultaneously through the system. Each transaction is always positioned at a block and most blocks can hold many transactions simultaneously. The movement of a transaction from one block to another occurs instantaneously at a specific time or when some change of system condition occurs. A block number is given to each block, and the movement of transactions is usually from one block to the block with the next highest location. The TRANSFER block allows a transaction to move to a block other than the next sequential block.

A block type called ADVANCE is used to represent a time delay. IGPSS computes an action time for each transaction as it enters an ADVANCE block, and the transaction remains at that block for the calculated period of simulated time before it attempts to proceed to the next block. Queues may be established to gather statistics on delays transactions may encounter in the system.

The first example in Appendix G is a simulation of jobs as they progress through an office. Jobs arrive at the office at the rate of one every 10 to 20 minutes. Normally they go to clerk A, who takes 7 to 13 minutes and then go to clerk B, who takes 3 to 7 minutes. However, if clerk B is busy at the time the job is brought to the office, the entire job is given to clerk C, who takes 10 to 30 minutes. Assuming that each clerk handles one job at a time, the simulation runs until 1000 jobs have been completed. The example shows the model run on GPSS followed by the

same model run on IGPSS.

In the GPSS program, the GENERATE block creates jobs every 10 to 20 minutes. The TEST block determines if clerk B is busy. If so, the job is sent to clerk C. The QUEUE block indicates entering into clerk A's stack of jobs to be processed. The stack is maintained in a FIFO arrangement. The SEIZE block indicates clerk A has begun processing the job; thus, the job leaves clerk A's stack of work in the DEPART block. The ADVANCE block indicated clerk A's processing time. The job leaves clerk A in the RELEASE block and enters clerk B's stack of work in the next QUEUE block. The job is processed by clerk B in the same manner as with clerk A. It then leaves the office in the TERMINATE block.

The output produced by both GPSS and IGPSS contains a section labeled 'BLOCK COUNTS'. This section indicates the total number of jobs that are currently in each block and the total number that has passed through each block. The facilities section shows how many jobs were processed by each clerk, the average processing time of each job, and the utilization for each clerk. The queues section indicates the maximum contents in each clerk's stack of jobs waiting to be processed was one. Most of the jobs had a zero waiting time, and the waiting time for those jobs that had to wait was 1.588 minutes for clerk A and 2.8 minutes for clerk C. No jobs had to wait for clerk B for processing.

The second example in Appendix G is a simulation of a grocery store checkout line. Thirty-two percent of the customers will use the express checkout line and the other sixty-eight percent will use one of the two regular checkouts. Customers arrive uniformly from two to fourteen minutes apart. Checkout time at the express counter requires 2 minutes and the regular lines require from 2 to 14 minutes. Queues are maintained to gather statistics on the waiting lines which form at each checkout

counter. The simulation is run for 100 customers. The example shows the model run on GPSS followed by the same model run on IGPSS.

In the GPSS program, the STORAGE block defines 2 regular checkout lines. The GENERATE block creates a customer uniformly every 2 to 14 minutes. The TRANSFER block sends 32% of the customers to EXPRS, the rest to REGCK. The express customer is first entered into the line at the QUEUE block, is serviced at the counter in the SEIZE block, thus leaving the line in the DEPART block. The ADVANCE block indicates a two-minute checkout time. The customer then leaves the counter in the RELEASE block and is recorded as leaving the store at the TERMINATE block. The customer is serviced at a regular checkout counter in a similar manner. The START block indicates the simulation will halt when the 100th customer enters the terminate block.

The output of the second example is similar to that of the first example. The additional section entitled STORAGES shows the statistics on the two regular checkout counters.

IGPSS Entities

IGPSS is built around a set of simple entities which are divided into four classes: dynamic, equipment, statistical, and operational.

Dynamic Entities. The dynamic entities are called transactions. Transactions are the items being simulated, such as customers in a grocery store. They are created and destroyed during a simulation run. Associated with each transaction are ten parameters which can be assigned values by the user to represent characteristics of the transaction. For example, a transaction representing a customer in a grocery store might have a parameter indicating the number of grocers at checkout time. This

number could be used in the simulation to determine how long the checkout process would take.

Equipment Entities. The second type of entity, the equipment entity, can be acted on by transactions. Facilities and storages are equipment entities. A facility can handle only one transaction at a time and represents a potential bottleneck in the system. A facility could be used to represent a grocery store checkout counter. A record of which transaction is using each facility within the system is maintained. If a transaction attempts to use a facility that is already in use, that transaction is delayed until the facility is available. IGPSS also maintains a record of the total number of time units for which each facility was in use. At the end of the simulation, this figure is used to compute the utilization of each facility. A count is also maintained of the total number of times each facility was used and the average time duration of each use.

Storages are used to represent parallel processing equipment which can handle more than one transaction at a time. A storage entity could be used to represent a parking lot. The user defines the capacity of each storage that he uses and IGPSS maintains a record of the number of storage units that are occupied at any given point in time. If a transaction attempts to occupy more space than is available in a given storage, the transaction is delayed until sufficient space becomes available. IGPSS maintains a record of the total number of transactions that enter each storage. The average number of clock units each transaction remained in each storage is reported at the simulation's end.

Statistical Entities. Queues are the third type of entity. In any system, transactions may be delayed by the unavailability of an equipment entity. Transactions that are delayed in their attempt to obtain a facility or a storage constitute such a queue. Queues are used to gather statistics on system behavior. Each queue maintains a list of transactions delayed at one or more points in the system, and keeps a record of average number of transactions delayed and the length of these delays. A grocery store checkout line could be represented by a queue. The user must specify selected points in the model where queue statistics are desired. IGPSS will then gather all queue statistics; a count of total number of transactions that enter the queue and also a count of the transactions that were delayed in each queue, the average time the transactions were delayed in each queue, and the maximum number of transactions in the queue at any time.

Operational Entities. The operational entities are called blocks. Like the blocks of a flowchart, they provide the logic flow of a system, instructing the transactions where to go and what to do next.

Transactions move from block to block in the simulation model in a manner similar to the way the items they represent would progress in the real system. Each movement of a transaction is considered as an event that is due to occur at some particular point in time. The program maintains a record of the times at which these events are due to occur in their correct time sequence (first come, first served, except when different priorities exist). In those cases where the action called for cannot be performed at the time originally scheduled (such as attempting to seize a facility which is already in use), processing the transaction ceases. When the blocking condition changes, the transaction is again

moved into its proper processing sequence; that is, transactions with the earliest event time are serviced first.

IGPSS maintains a clock that records the instance of real time that has been reached in the simulation. It is referred to as the clock time. All times are represented as integer numbers. The unit of system represented by a unit of clock time is implied by the user. The time unit chosen should be the smallest time difference between event occurrences and this same unit must be maintained consistently throughout the model.

PROGRAM STRUCTURE

An IGPSS program consists of a set of block definition cards, entity definition cards, and control cards.

Each defined block is entered on a separate 'block definition card'. Each block is assigned a sequential number as the program is read. The block type is entered in the card followed by the subfields specifying the block arguments. The meaning associated with each subfield (A-E) is a function of the particular block type and will be discussed later.

The entity definition cards define storage capacities or functions. The basic format is the same as the block definition card. The function definition card is followed by cards specifying the X and Y coordinate values for the function.

The control cards tell IGPSS the length of the run desired and when the end of the input deck is reached.

STANDARD NUMERICAL ATTRIBUTES

Quantities such as function (FN) and clock time (C1) are generally termed standard numerical attributes (SNA). They represent properties or states of the simulation model and are provided to give the user dynamic

access to the results of the simulation. For example, whenever a reference is made to Q12, the contents of queue 12 are obtained. By making such a reference, a transfer can be made by a transaction if, for example, more than 20 transactions are currently in a certain queue. Thus, a user may simulate a system whose logic depends dynamically on any standard numerical attributes.

All standard numerical attributes consist of a one- or two-letter mnemonic followed by a number. The mnemonic indicates the type of attribute desired and the number indicates the specific attribute desired.

A list of all SNA's available in IGPSS is shown in Appendix C. Standard numerical attributes may be used as arguments of almost any block types.

EQUIPMENT ORIENTED BLOCK TYPES

As discussed earlier, a facility can handle only one transaction at a time and storages can handle multiple transactions. This section discusses the various block types used to deal with these equipment entities.

Two block types, SEIZE and RELEASE, are provided so that transactions may make use of a facility. When a transaction enters a SEIZE block it is recorded as using the facility. The facility then remains in use until the seizing transaction enters a RELEASE block. Upon entering the SEIZE block, the usage count of the facility is incremented. The facility to be seized is specified in field A of the block definition card. When the seizing transaction enters a RELEASE block with the same facility specified in field A, the facility is returned to the available state, and the in-use time for the facility is incremented by the number of clock

units that the transaction remained in possession of that facility. A transaction may seize any number of facilities.

Two block types, ENTER and LEAVE, provide for usage of storages. The STORAGE definition card is used to define the capacity of a storage. Field A specifies the maximum storage capacity and the location field specifies the storage number. Each ENTER and LEAVE block specifies a storage number in field A and the number of units to be occupied or returned in field B. When a transaction encounters an ENTER block, the record of the number of units occupying the storage is incremented by field B. If field B is blank it is assumed to be one. If there is not enough space available for a particular transaction, it will be refused entry until sufficient storage becomes available. If there is room for a subsequent transaction, that transaction will be admitted before the rejected one. When a transaction encounters a LEAVE block, the number specified in field B is returned to the available storage. A transaction need not remove the same number of units that it added; in fact, the transaction entering a LEAVE block need not have entered the storage previously. The contents of the storage cannot be allowed to be made negative, however.

TRANSACTION-ORIENTED BLOCK TYPES

IGPSS provides five block types which modify transaction attributes.

Advance

The first of these, the ADVANCE block, provides for the explicit specification of a time delay. This block requests a computation of the number of clock units which a transaction is to spend in the block. This time may have any integer value, including zero. If the computed time is

zero, the simulator continues processing the entering transaction and immediately attempts to move it to the next block in its path. To specify the delay time at an ADVANCE block, the user states the mean in field B. If the modifier is omitted, a constant time delay is produced.

Two types of ADVANCE block modifiers exist: a spread and a function. A spread modifier is used when the delay times are equiprobable within a given range. For example, a spread of 5, used with a mean of 10, produces action times in the range of 10 ± 5 . The delay time then becomes one of eleven possible integers since clock time only assumes integer values. All values are computed with equal probability. If delay times in a system follow a more complex distribution pattern than the simple uniform distribution produced by the spread modifier, a function modifier must be used. A value for the specified function is computed and multiplied by the mean. If the result is noninteger, the decimal portion is truncated and the resulting integer is the computed delay time.

Generate

The second transaction oriented block type is the GENERATE block. This block is used to create transactions for entry into the system. Field A is the mean time between generations. This mean can be modified by specifying either a spread or a function modifier in field B, as in the ADVANCE block. Each transaction contains ten parameters which are set equal to zero when the transaction is created. Field C is an offset interval. If specified it is used without modification as the time interval preceding the creation of the first transaction by that block. Field D is a limit count on the number of transactions that can be created by the GENERATE block. If this field is omitted, the block will continue to create transactions indefinitely. Field E is the priority of the

transactions to be created by this block. Newly created transactions have a priority of zero unless another value is specified in this field.

Terminate

The TERMINATE block removes transactions from the system. It represents the completion of a path of flow, such as the customer leaving the grocery store. Field A of the TERMINATE block specifies whether or not this block contributes to the total termination count and, if so, how many units it will contribute. The termination count is specified on the START card which will be discussed later. If the field is not specified it is interpreted to be 0, and transactions entering the block do not decrement the termination count. There must be at least one TERMINATE block with a value of 1 or more in field A of every model to terminate the simulation. If this does not happen, the model runs indefinitely.

Assign

Each transaction has 10 parameters associated with it. Contents of parameters are signed integers ranging from $-(2^{31})$ to $(2^{31}-1)$. At the time of creation of a transaction, these parameters are initialized to zero. The ASSIGN block is the principal means of entering values into the parameter fields of a transaction. The value of field A specifies which parameter is to be modified. Field B specifies the integer value the parameter is to become.

Priority

The PRIORITY block is used to set the priority of a transaction to a specified value. Field A specifies the priority the transaction will assume.

FLOW MODIFICATION

In IGPSS, the flow of transactions through blocks is generally sequential. That is, it is assumed that a transaction goes from its current block to the next sequentially numbered block. However, in an actual system, the flow may be diverted.

Transfer

The TRANSFER block is generally used to direct the transaction to a nonsequential next block. This may be done in a variety of ways. The choice is specified by means of a mnemonic selection mode specified in field A of the block. If the selection mode at the TRANSFER block is blank, every transaction that enters the block will be sent to the block specified in field B. If the selection mode is a fraction, a random choice between the blocks specified in fields B and C is made. The probability of selecting the next block in field C is given by the fraction. For example, if a selection mode of .370 were specified in field A, 37% of the transactions would proceed to the next block specified in field C and 63% would proceed to the next block specified in field B.

Test

The TEST block also provides a possible modification of the sequential block flow. The TEST block specifies a condition specified in the form of an algebraic comparison between two arguments. If the desired relation is satisfied, the transaction proceeds to the next sequential block. If not satisfied, the transaction proceeds to the block specified in field C. The arguments are specified in fields A and B. The following six mnemonic relations may be specified:

- L Less than. If the first argument is greater than or equal to the second argument, the relation is not satisfied.
- LE Less than or equal to. The relation is satisfied unless the first argument is greater than the second.
- E Equal to. The relation is satisfied only if the two arguments are equal.
- NE Not equal. The relation is satisfied unless the two system variables are equal.
- G Greater than. If the first argument is less than or equal to the second argument, the relation is not satisfied.
- GE Greater than or equal to. The relation is satisfied unless the first argument is less than the second argument.

COMPUTATIONAL AND REFERENCE ENTITIES

Very often, the complex logical and mathematical interrelations which exist among items within a model may be best expressed in terms of a mathematical function.

Functions

A FUNCTION returns a numerical value that is computed by a rule defined by the user. IGPSS provides for two types of functions: continuous and discrete. A unique relationship is defined between the independent and dependent variables. The independent variable can be any of the standard numerical attributes and is specified in field A. Field B contains the number of pairs of points desired, preceded by an

indication of whether a discrete or continuous function is desired. The X and Y values are entered in cards immediately following the FUNCTION block card. These values are entered by specifying $X_1, Y_1, X_2, Y_2, \dots, X_n, Y_n$, until the desired number of pairs is given.

Random Number Generation

There are eight random number seeds in IGPSS. The user may specify any one of these seeds (RN1-RN8). In cases where reference to a random number is implied (such as TRANSFER .5, a1b1, b1b1), RN1 is used.

STATISTICAL BLOCK TYPES

The IGPSS user will primarily be interested in observing the values of numerical attributes which characterize a system in order to determine how the system is performing in simulated operation. Some statistics are gathered automatically by the program, but others which might be of interest can be gathered by inserting various queues into the system.

Queues are measured by inserting QUEUE and DEPART blocks into the model. A QUEUE block may be placed before any block that might cause a delay. The transactions increment the contents of a queue when they enter a QUEUE block and decrement the contents of that queue when they succeed in entering a DEPART block.

Queue Block

The QUEUE block is analogous to the ENTER block and signals the program that statistics should be gathered for a queue at the specified point in the model. The number of the queue the transaction is to enter is specified in field A. The length of time for which the queue specified has been at its present content is then computed. The product of that interval and the present contents of the queue are added to a sum that

represents the total unit occupancy (cumulative time integral) of the queue. At the end of the simulation run, this sum is divided by the length of the run to form the average contents of the queue. The number of units to be added to the queue is specified at field B in the same manner as in the ENTER block and LEAVE block. If field B is left blank it is assumed to be 1. When a transaction enters the block, the number of units specified in field B is added to the current contents of the queue. IGPSS then compares the new contents with the previous maximum length. If the new contents is larger, the new queue length is recorded as the maximum. The count of the total number of units that have passed through the queue is also incremented by the amount in field B. This total is used to compute the weighted average time per transaction in the queue.

Depart

The DEPART block is analogous to the LEAVE block. Field A specifies the queue number and field B, the number of units to leave the queue. When a transaction enters a DEPART block, the number of units specified in field B is subtracted from the contents of the queue. IGPSS also computes the length of time which transactions remained in the queue. If the time interval is 0, the specified number of units is added to a count which records the total number of units to pass through the queue with no delay.

CONTROL CARDS

Start Card

This card indicates to the simulator that all input data has been received and that the run may now proceed. Field A specifies the number of terminations to be executed before printing the final summary statistics. This quantity is called the run termination count and when it has been

decremented to 0 or less, the run will terminate. The count is decremented by the amount specified in field A of a TERMINATE block each time a transaction enters that block.

End Block

This card specifies the end of the input deck.

PROGRAM OUTPUT

Program output occurs when an error condition arises or upon normal completion of a run.

The output produced upon normal completion is as follows:

1) Clock time - This is the actual clock time at the end of the simulation run.

2) Block counts - The entry and wait counts of all defined blocks is printed out. This printout includes block number, the number of transactions currently waiting at that block and the total number of transactions which have entered that block in the course of the simulation.

3) Facility statistics - The statistics which are gathered for each facility referred to in the simulation are listed. These statistics include utilization (the fraction of time that each facility was seized), the total number of entries, the average number of time units that a transaction held each facility, and the number of the transaction seizing the facility at the clock instance the run ended.

4) Storage statistics - A listing is provided of the statistics gathered for each storage that was referenced during a run. These include the capacity of each storage as defined by the user, the average contents of each storage, the average utilization (average contents divided by maximum capacity), the total number of entries, the average length of

time transactions remained in each storage, the current contents at the end of the run, and the maximum occupancy recorded for each storage during the run.

5) Queue statistics - The statistics gathered for each queue referenced in the simulation. These statistics include the maximum contents of each queue, the average contents, the total number of entries, the number of transactions which entered each queue but were not delayed, the average time per transaction, the average time per transaction for those transactions which were delayed (excluding zero delay entries), and the current contents of each queue at the time of the simulation run termination.

INTERNAL ORGANIZATION

Transactions in IGPSS are maintained in chains. There are two such chains and a transaction may be in either one at a given time. The current event chain (CEC) contains those transactions whose scheduled event time is less than or equal to the current value of the clock. The future events chain (FEC) contains those transactions whose scheduled event times are greater than the current clock time; that is, transactions whose event time is in the future.

Current Events Chain

The current events chain is organized in descending order of priority. Each transaction in the CEC is either in the active status or in a delay status. The scan status indicator specifies the status condition. If the transaction is in the active status IGPSS attempts to move it to its next block. If the transaction has been blocked for some reason it is put into delay status. Transactions in this status are

skipped by the overall scan. When the blocking situation is changed, the transaction is returned to the active status.

Future Events Chain

The future events chain transactions are strictly in ascending block departure time order. There are no distinctions made by priority in this chain.

DATA STRUCTURES

The data structures in IGPSS are implemented through the use of arrays in a common block. A maximum of 200 transactions are allowed at any one time. Each transaction word contains a pointer to the next block to be executed, a scan status indicator, the priority of the transaction, the time of the last status change, the current block, and 10 parameters which are initially zero. A maximum of 120 blocks are permitted. Each block word contains the block type, field A, B, C, D, and E if applicable, and a current and total block entry count. Thirty-five facilities can be used. Each facility word contains the number of the transaction seizing the facility, the cumulative time usage, the clock time of the last status change, an indicator as to whether a delay chain exists, and an entry count. Each storage word of the 35 possible storages contains the current contents, the available contents, the cumulative time usage, the clock time of the last status change, the entry count, the maximum contents, and a field indicating the existence of a delay chain. Each of the 70 queue words contains the clock time of the last status change, the entry count, the zero delay entries, the cumulative time usage, the current and the maximum contents. Twenty functions are allowed with a maximum of twenty-five X and Y values. Each function word also indicates

the number of X and Y values present, whether it is continuous or discrete, and the SNA associated with function. Also contained in the common block are the future and current events chains. Each chain may contain fifty transactions at any one time. Each entry in the chains contains pointers to the previous and next entry in the chain, the priority of the transaction and the transaction number. The future events chain entries also contain the block departure time. A pointer is maintained to the first and last entry in each chain. A pointer also exists which indicates the next available array location into which a new entry is to be placed when a transaction is added to a chain. The common block also contains the clock time and the status change flags. Eight random number seeds are available and are located in the common area.

INTERNAL FLOW

The main routine initializes all necessary common block data then calls the input routine. The input routine reads all the simulation input data. If a generate block type is read, the GENERATE routine is called and a transaction is entered into the future events chain. If a START block type is encountered, the start count is initialized. If a FUNCTION definition is encountered, the function initialize routine is called and the function X, Y cards are read and the function common block is initialized. If a STORAGE definition is read, the appropriate storage available area is initialized. When an END block is read, the input routine returns control to the main program. If all input is read before an END block is encountered, the program encounters an end of file error and execution is halted. The main routine then calls the SCAN routine which controls the overall transaction scanning. A check is first

made to determine if the start count is greater than zero. If not, an error is written and the program is stopped.

The overall scan is divided into three phases: clock updating, current event chain scan, and transaction move. The overall scan begins by increasing the clock time to the block departure time of the first transaction in the future events chain. Next the first transaction is moved from the future events chain to the current events chain. Succeeding future events chain transactions are moved to the current events chain as long as their block departure time equals the new clock time. Transactions are inserted into the current events chain maintaining the priority ordering. New transactions are added behind existing ones with the same priority. When all transactions have been moved, the status change flag is turned off and the current events chain scan begins. The first current events chain transaction is obtained. If the scan status indicator (SSI) is on, the transaction is not processed since it is waiting on an unavailable entity. The next current events chain transaction is then obtained. When all transactions are processed the clock time is again updated to the block departure time of the first future events chain transaction and the process is repeated. When a current events chain transaction is obtained whose SSI is not on, IGPSS attempts to move the transaction through the next block. The block number is obtained from the transaction word and the appropriate block type routine is called. If the transaction did not succeed in entering the next block, the status change flag is tested. If on, a facility or storage was released which had a transaction waiting on it. IGPSS then returns to the first of the current events chain and repeats processing of the entire chain. If the status change flag is not on, the next transaction in the CEC is obtained for processing. If the transaction did move through the current block, a test

is made to determine if the block terminated the transaction or if it entered a positive time advance block. If the test is true, IGPSS then checks the status change flag and proceeds as above. If the test is false, an attempt is made to move the transaction through the next block by calling the appropriate block routine. Thus each CEC transaction is in turn moved through as many blocks as possible. Then movement is stopped; control passes to the next CEC transaction. If a transaction frees a facility or storage which has waiting transactions, control returns to the first of the CEC after the freeing transactions movement is stopped.

The flow of the routines which process each individual block type is illustrated by the flowcharts in Appendix F.

CODING IGPSS

The parser of IGPSS remains to be written. Currently IGPSS uses as its input the anticipated output from the parser. A coding sheet to simplify data entry is provided in Appendix A. The NUM field, columns 1-3, is an integer variable which contains the storage or function number. The remaining eight fields correspond to the fields 1-8 of the block word. They are all real numbers. The block type entered in columns 4-5 is shown in Appendix B. The SNA for field A is entered in columns 6-7. A table of SNA's is shown in Appendix C. Field A is contained in columns 8-17. The SNA for field B is placed in 18-19. Fields B, C, D, and E are placed in columns 20-29, 30-39, 40-49, 50-59, respectively. If the SNA's for fields A and B are constants, the SNA field may be left blank.

RUNNING IGPSS

To run IGPSS on the Interdata 8/32 the user must sign on, load IGPSS, select the logical units desired for input and output, then start the program. Following is an example of some possible methods of running IGPSS.

To use card input and the terminal as output:

```
SIGNON yourname,5,GRAF
V USR6
LO IGPSS
AS 5,CR:
AS 6,.ME
ST
SIGNOFF
```

To use terminal as input and printer as output:

```
SIGNON yourname,5,GRAF
V USR6
LO IGPSS
AS 5,.ME
AS 6,pr:
ST
SIGNOFF
```

The terminal will return:

```
STOP
*
```

when execution is completed. If another run is desired the program must be reloaded and the units reassigned.

RESTRICTIONS

IGPSS is a subset of GPSS. The available block types in IGPSS are identical to those described in the GPSS Introductory User's Guide (3) with the following exceptions. The number of parameters set in field F of the GENERATE block in GPSS is set to 10 in IGPSS. The function modifier in field C of the ASSIGN block is not implemented. The TRANSFER block of IGPSS does not recognize the BOTH, ALL, or SIM options. Functions

are limited to 25 X,Y coordinates and the standard numerical attributes are limited to parameters, random numbers, storage contents, storage remainder, and queue length. The quantity limitations for all IGPSS entities are shown in Appendix D.

TESTING

IGPSS was first implemented on the IBM/370. A test case was first made to test correctness of the simplest overall scan using only a GENERATE, TERMINATE, START, and END blocks. All options of the GENERATE were then tested. Next a facility was added followed by simple advance block. The modifier of the ADVANCE block was then tested. Queues were added, followed by storages and functions. The transfer was then tested and the resulting simulation was compared to the same model using GPSS. The results were very similar and any variations were attributed to the difference in the random number generators. Next the TEST block, the ASSIGN, and the PRIORITY blocks were tested. When the program appeared to be working, several test cases were run to compare results between GPSS and IGPSS. All results were very similar.

In the meantime, IGPSS was compiled on the Interdata 8/32. Minor differences were discovered between compilers on the two machines. The Interdata would not allow comparison of integers and reals or the use of a real as a subscript.

When the program was successfully working on the 370, a test run was run on the Interdata. A problem existed in passing an argument to a subroutine. The implicit real statement apparently had no effect on changing the mode of the argument. When this problem was corrected, the models tested on the 370 produced identical results to those run on the Interdata.

The run times for all test models are less than a minute. There is no means of determining CPU time on the 8/32 so no direct time comparisons could be made. The compile time on the Interdata, however, was over fifteen minutes while it was less than one minute on the 370. No apparent reason was found for this long a compile time required on the Interdata.

CONCLUSION

IGPSS is currently implemented on the IBM/370 and the Interdata 8/32. Minor compiler differences were encountered between the two hardware configurations during development. Simulation results from the test models were identical on both systems.

Some IGPSS standard numerical attributes such as variable and function can contain subfields which are also standard numerical attributes. This would require a recursive call to the routine which determines the SNA. Since recursion is not possible in Fortran, the variable block type was not implemented and the standard numerical attributes available with functions are limited. The variable block type is very useful in GPSS and could possibly be added as a future project through the use of an assembler subroutine. The SNA's for use with functions could be changed fairly easily if the need for additional SNA's arises.

Input to IGPSS is somewhat clumsy, but this was anticipated since the parser remains to be written. The parser will allow the user to enter his data in a format similar to that used for GPSS. This format is more free form than that necessary for IGPSS. The parser will also accept variable names for entities and labels and resolve them to the absolute numbers required by the IGPSS program.

Most of the restrictions on internal limitations shown in Appendix D could be increased with little difficulty. The dimensions of the arrays in common would have to be increased and some minor program changes might have to be made, depending on the entity that is to be increased. The limitations set for IGPSS are the same as the 64K version of GPSS.

REFERENCES

1. Fishman, George S. *Concepts and Methods in Discrete Event Digital Simulation*. New York: John Wiley & Sons, 1973.
2. Geoffrey, Gordon. *System Simulation*. Englewood Cliffs, New Jersey: Prentice-Hall, 1969.
3. IBM, GPSS/360 Introductory User's Manual. IBM Corp., Form No. H20-0304.
4. IBM, GPSS/360 User's Manual. IBM Corp., Form No. GH20-0326.

**THIS BOOK
CONTAINS
NUMEROUS PAGES
WITH DIAGRAMS
THAT ARE CROOKED
COMPARED TO THE
REST OF THE
INFORMATION ON
THE PAGE.**

**THIS IS AS
RECEIVED FROM
CUSTOMER.**

IGPSS BLOCK TYPES

BLOCK TYPE	A	B	C	D	E
1 - Generate	mean	(spread)	(offset)	(count)	(priority)
2 - seize	facility no.				
3 - release	facility no.				
4 - storage	capacity				
5 - enter	storage	(units)			
6 - leave	storage	(units)			
7 - advance	mean	(spread)			
8 - terminate	(units)				
9 - assign	parameter	source			
10 - priority	priority				
11 - transfer	selection mode	next block a	next block b		
12 - test	arg1	arg1	next block b		
13 - function	ind. variable	no points			
14 - variable	not implemented				
15 - queue	queue no.	(units)			
16 - depart	queue no.	(units)			
17 - start	units				
18 - end					

STANDARD NUMERICAL ATTRIBUTES (SNA)

1	Pj	parameters
2	PR	priority
3	Kj	constant
4	RN(x)	1 x 8 random numbers
5	C1	clock time
6	FNj	function
7	Vj	variable
8	Sj	contents of storage
9	Rj	remaining storage
10	SRj	utilization of storage
11	SAj	average of storage
12	SMj	maximum of storage
13	SCj	number of entries
14	STj	average time in storage
15	Fj	status of facility
16	FRj	utilization of facility
17	FCj	number of entries in facility
18	FTj	average time in facility
19	Qj	queue length
20	QAJ	average length of queue
21	QMj	maximum length of queue
22	QCj	number of entries in queue
23	QZj	number of zero entries in queue
24	QTj	average time with zero entries in queue
25	QXj	average time with no zero entries in queue

QUANTITY LIMITATIONS ON IGPSS ENTITIES

Internal Limitations

Parameters	10
Transactions	200
Block	120
Facilities	35
Storage	35
Queues	70
Functions	20
Variables	not implemented

Currents Events Chain (CEC)

CEC1	previous # in chain
CEC2	next # in chain
CEC3	priority
CEC4	tran #

Future Events Chain (FEC)

FEC1	previous # in chain
FEC2	next # in chain
FEC3	block departure time (BDT)
FEC4	priority
FEC5	tran #

DATA STRUCTURES

Facility word

F1	tran # seizing
F2	cumulative time usage
F3	clock time of last status change
F4	delay chain
F5	entry count

Storage word

S1	current contents
S2	available
S3	cumulative time usage
S4	last clock status change
S5	entry count
S6	maximum contents
S7	delay chain

Queue word

Q1	time of last status change
Q2	entry count
Q3	zero-delay entries
Q4	cumulative time
Q5	current contents
Q6	maximum contents

Transaction word

T1	next block for transaction
T2	scan status indicator
T3	priority
T4	time of last status change
T5	current block
T6 - T 15	parameters

Block word

B1	block type
B2	sna for field A
B3	field A
B4	sna for field B
B5	field B
B6	field C
B7	field D
B8	field E
B9	current block count
B10	total block count

Function word

Func1-Func25	x values
Func26-Func50	y values
Func51	number of points
Func52	continuous(C) or discrete(D)
Func53	sna for field C
Func54	field C

APPENDIX F

PROGRAM LISTINGS

[illegible]

```

50 FEC(1,2) = I + 1
   CECA = 1
   FECA = 1
   FECF = 0
   FECL = 0
   CECF = 0
   CFC(50,2) = 0
   CECL = 0
   FEC(50,2) = 0
   STARTC = 0
   DO 60 I = 1,120
     BLOCK(1,9) = 0
     BLOCK(1,10) = 0
     BLOCK(1,11) = 0
60  CALL INPUT
    CALL SCAN
    CALL OUTPUT
    STOP
    END

```

SUBROUTINE ADVANC(CECT, ID)

ADVANCE

SUBROUTINE ADVANCE DETERMINES THE BDT OF THE TRANSACTION IF THE CALCULATED TIME IS GREATER THAN CI IT, IS PLACED IN THE FEC. IF TIME IS = CI, MOVE IS SET TO ZERO INDICATING THE TRANSACTION HAS MOVED THROUGH THE BLOCK.

```

IMPLICIT INTEGER(C,F,Q,S,O)
REAL FUNC
INTEGER CECT
INTEGER TYPE,TRANO,BLOCKN,BLOCKT
COMMON /A/ TRAN(200,15),BLOCK(120,10),FACIL(35,5),STOR(35,7),
1QUEUE(70,6),FEC(50,5),FECF,FECI,FECA,CECF,CECI,CECA,
2STATUS,MOVE,C1,TRANO,BLOCKN,STARTC,RN(8),TRANA,FUNC(20,54)
C   UPDATE BLOCK COUNTS
   ILST = TRAN(TRANO,5)
   TRAN(TRANO,5) = BLOCKN
   BLOCK(BLOCKN,9) = BLOCK(BLOCKN,9) + 1
   BLOCK(ILST,9) = BLOCK(ILST,9) - 1
   BLOCK(BLOCKN,10) = BLOCK(BLOCKN,10) + 1
   TRAN(TRANO,1) = TRAN(TRANO,1) + 1
C   DETERMINE FIELD A
   SNAC = BLOCK(BLOCKN,2)
   SNAN = BLOCK(BLOCKN,3)
   CALL SNA(SNAC,SNAN,A)
   IF(BLOCK(BLOCKN,5).EQ.0.0) GO TO 200
C   DETERMINE FIELD B
   SNAC = BLOCK(BLOCKN,4)
   SNAN = BLOCK(BLOCKN,5)
   CALL SNA(SNAC,SNAN,B)
   IF(BLOCK(BLOCKN,4).EQ.6.0) GO TO 100
   CALL RANDU(1,IY,YFL)
   IR = 2*8*YFL+A-B+C1+.5
   GO TO 300
C   FIELD B MODIFIER IS FUNCTION
100  IR = A * B + C1
   GO TO 300
200  IR = A + C1
300  IF(IR.EQ.C1) GO TO 700
C   PLACE IN FEC IF POSITIVE ADVANCE BLOCK
   IF(FECF.NE.0) GO TO 350
   FECT = FEC(FECA,2)
   FEC(FECA,1) = 0
   FEC(FECA,2) = 0
   FEC(FECA,3) = IR

```

```

FEC(FECA,4) = CEC(CECT,3)
FEC(FECA,5) = CEC(CECT,4)
FECF = FECA
FECL = FECA
FECA = FECT
GO TO 600
      350 FECT = FECF
      400 IF(FEC(FECT,3).GT.IR) GC TO 500
      FECT = FEC(FECT,2)
      IF(FECT.NE.0) GO TO 400
      ADD TO END OF FEC
      FECT = FEC(FECA,2)
      FEC(FECA,1) = FECL
      FFC(FECA,2) = 0
      FEC(FECA,3) = IR
      FEC(FECA,4) = CEC(CECT,3)
      FEC(FECA,5) = CEC(CECT,4)
      FEC(FECL,2) = FECA
      FECL = FECA
      FECA = FECT
      GO TO 600
C
      ADD INTO FEC
      500 FECAT = FEC(FECA,2)
      FEC(FECA,2) = FECT
      FEC(FECA,1) = FEC(FECT,1)
      FEC(FECA,3) = IR
      FEC(FECA,4) = CEC(CECT,3)
      FEC(FECA,5) = CEC(CECT,4)
      IF(FECT
        .EQ.FECF) FECF = FECA
      OLDPRE = FEC(FECT,1)
      IF(OLCPRE.NE.0) FEC(OLDPRE,2) = FECA
      FEC(FECT,1) = FECA
      FECA = FECAT
      UNLINK FROM CEC
C
      600 OLD1 = CEC(CECT,1)
      OLD2 = CEC(CECT,2)
      IF(OLD1.NE.0) CEC(OLD1,2) = OLD2
      IF(OLD2.NE.0) CEC(OLD2,1) = OLD1

```

```

      CEC(CECT,2) = CECA
      CECA = CECT
      IF(CECT.EQ.CECF) CECF = OLD2
      IF(CECT.EQ.CECL) CECL = OLD1
      IF(CECL.NE.0) GO TO 650
      CECT = 0
      ID = 1
      RETURN
650  ID = 1
      CECT = OLD2
      RETURN
700  MOVE = 1
      RETURN
      END
      SUBROUTINE ASSIGN
C*****
C      ASSIGN
C
C      SUBROUTINE PLACES THE RESULT DETERMINED FROM FIELD B IN THE
C      PARAMETER SPECIFIED IN FIELD A.
C*****
C      IMPLICIT INTEGER(C,F,Q,S,O)
      REAL FUNC
      INTEGER TYPE,TRANO,BLOCKN,BLOCKT
      COMMON /A/ TRAN(200,15),BLOCK(120,10),FACIL(35,5),STOR(35,7),
      IQUEUE(70,6),FEC(50,5),FECF,FECI,FECA,CEC(50,4),CECF,CECL,CECA,
      2STATUS,MOVE,C1,TRANO,BLOCKN,STARTC,RN(8),TRANA,FUNC(20,54)
      MOVE = 1
      DETERMINE SNA
      SNAC = BLOCK(BLOCKN,4)
      SNAN = BLOCK(BLOCKN,5)
      CALL SNA(SNAC,SNAN,RESULT)
      MOVE VALUE TO PARAMETER
      IPAPM = BLOCK(BLOCKN,3) + 5
      TRAN(TRANC,IPARM) = RESULT
C*****

```

```

C      UPDATE BLOCK COUNTS
      ILST = TRAN(TRANQ,5)
      TRAN(TRANQ,5) = BLCCKN
      BLOCK(BLOCKN,9) = BLCCK(BLOCKN,9) + 1
      BLOCK(ILST,9) = BLOCK(ILST,9) - 1
      BLOCK(BLCCKN,10) = BLOCK(BLOCKN,10) + 1
      TRAN(TRANQ,1) = TRAN(TRANQ,1) + 1
      RETURN
      END
      SUBROUTINE DEPART
C*****
C      DEPART
C
C      SUBROUTINE DEPART REMOVES THE NO OF UNITS SPECIFIED IN
C      FIELD B FROM THE QUEUE SPECIFIED IN FIELD A.
C*****
C      IMPLICIT INTEGER(C,F,Q,S,O)
      REAL FUNC
      INTEGER TYPE,TRANQ,BLOCKN,BLOCKT
      COMMON /A/ TRAN(200,15),BLOCK(120,10),FACIL(35,5),STOR(35,7),
1      IQUEUE(70,6),FEC(50,5),FECF,FECCL,CECA,CECF,CECL,CECA,
2      STATUS,MOVE,C1,TRANQ,BLOCKN,STARTC,RN(8),TRANA,FUNC(20,54)
      DETERMINE Q
      SNAC = BLOCK(BLOCKN,2)
      SNAN = BLCCK(BLOCKN,3)
      CALL SNA(SNAC,SNAN,RES)
      Q = RES
C
C      ARE YOU TRYING TO REMOVE MORE UNITS FROM QUEUE THEN IN QUEUE
C
      DETERMINE NO TO RELEASE
      SNAC = BLOCK(BLOCKN,4)
      SNAN = BLCCK(BLOCKN,5)
      CALL SNA(SNAC,SNAN,RESULT)
      IR = RESULT
      IF(IR.EQ.0) IR = 1
      IF(IR.GT.QUEUE(Q,5)) GO TO 1000

```



```

C      DETERMINE STORAGE
10  SNAC = BLOCK(BLOCKN,2)
   SNAN = BLOCK(BLOCKN,3)
   CALL SNA(SNAC,SNAN,RES)
   SNO = RES
C      HAS STORAGE BEEN DEFINED
   IF(STOR(SNO,2).NE.-1) GO TO 100
   WRITE(6,50) SNO
50  FORMAT('0 STORAGE NO.,13,' HAS NCT BEEN DEFINED.')
   STOP
100 SNAC = BLOCK(BLOCKN,4)
C      DETERMINE NO TO ENTER
C      CAN YOU GET DESIRED NC OF STORAGE UNITS
   SNAN = BLOCK(BLOCKN,5)
   CALL SNA(SNAC,SNAN,RESULT)
   IR = RESULT
   IF(IR.EQ.0) IR = 1
   IF(STOR(SNO,2).GE.IR) GO TO 1000
C      SET UP DELAY CHAIN
   STOR(SNO,7) = 1
   GO TO 2000
C      CALCULATE NEW CUMULATIVE TIME
1000 STOR(SNO,3) = STOR(SNO,3) + (C1 - STOR(SNO,4))* STOR(SNO,1)
   STOR(SNO,4) = C1
   STOR(SNO,2) = STOR(SNO,2) - IR
   STOR(SNO,1) = STOR(SNO,1) + IR
   STOR(SNO,5) = STOR(SNO,5) + IR
   IF(STOR(SNO,1).GT.STOR(SNO,6)) STOR(SNO,6) = STOR(SNO,1)
C      XACT WAS MOVED THRU BLOCK
C      UPDATE BLCKCK COUNTS
   ILST = TRAN(TRANO,5)
   TRAN(TRANO,5) = BLOCKN
   BLOCK(BLOCKN,9) = BLOCK(BLOCKN,9) + 1
   BLOCK(ILST,9) = BLOCK(ILST,9) - 1
   BLOCK(BLOCKN,10) = BLOCK(BLOCKN,10) + 1
   TRAN(TRANO,2) = 1

```



```

200 IR      = A * B + C1
   GO TO 400
300 IR      = A + C1
400 IF(BLOCK(IN,6).EQ.0.0) GO TO 450
   IR      = IR      + BLCKCK(IN,6)
      BLOCK(IN,6) = 0
   PLACE IN FEC
C
450 IF(FECF.NE.0) GO TO 500
   FECTA = FEC(FECA,2)
   FEC(FECA,1) = 0
   FEC(FECA,2) = 0
   FEC(FECA,3) = IR
   FEC(FECA,4) = TRAN(TRANO,3)
   FEC(FECA,5) = TRANO
   FECF=FECA
   FECL=FECA
   FECA = FECTA
   GO TO 800
500 FECT = FECF
600 IF(FEC(FECT,3).GT.IR) GO TO 700
   FECT = FEC(FECT,2)
   IF(FECT.NE.0) GO TO 600
   TRAN(TRANC,I)= 0
100 DETERMINE FIELD A
   SNAC = BLOCK(IN,2)
   SNAN = BLOCK(IN,3)
   CALL SNA(SNAC,SNAN,A)
   IF(BLCKCK(IN,5).EQ.0.0) GO TO 300
C
   DETERMINE FIELD B
   SNAC = BLOCK(IN,4)
   SNAN = BLCKCK(IN,5)
   CALL SNA(SNAC,SNAN,B)
   IF(BLOCK(IN,4).EQ.6.0) GO TO 200
   CALL RANDU(1,IY,YFL)
   IR = 2*B*YFL+A-B+C1+.5
   GO TO 400
C
   FIELD B MODIFIER IS FUNCTION

```

```

C      ADD TO END OF FEC
      FECT = FEC(FECA,2)
      FEC(FECA,1) = FECT
      FEC(FECA,2) = 0
      FEC(FECA,3) = IR
      FEC(FECA,4) = TRAN(TRANO,3)
      FEC(FECA,5) = TRANO
      FEC(FECL,2) = FECA
      FECL = FECA
      FECA = FECT
      GO TO 800

C      ADD INTO FEC
700    FECAT = FEC(FECA,2)
      FEC(FECA,1) = FEC(FECT,1)
      FEC(FECA,2) = FECT
      FEC(FECA,3) = IR
      FEC(FECA,4) = TRAN(TRANO,3)
      FEC(FECA,5) = TRANO
      OLDPRE=FEC(FECT,1)
      IF(OLDPRE.NE.0) FEC(CLOPRE,2) = FECA
      FEC(FECT,1) = FECA
      IF(FECT .EQ.FECF) FECF = FECA
      FECA = FECAT
800    IF(ID.EQ.2)      GO TO 900
      IF(BLOCK(IN,7).EQ.0.0) BLOCK(IN,7) = 999999999
      RETURN
900    BLOCK(IN,7) = BLCKK(IN,7) -1
      TRANO = OLDT
1000   TRAN(TRANO,1) = TRAN(TRANO,1) + 1
      TRAN(TRANO,5) = IN
      RETURN
      END

```



```

SUBROUTINE LEAVE
C*****
C
C
C
C
C
C
C
C*****
      LEAVE
C
C
C
C
C
C
C
C*****
      SUBROUTINE LEAVE RELEASES THE NO OF UNITS SPECIFIED IN FIELD
      B FROM THE STORAGE SPECIFIED IN FIELD A.
C*****
      IMPLICIT INTEGER(C,F,Q,S,O)
      REAL FUNC
      INTEGER TYPE,TRANQ,BLOCKN,BLOCKT
      COMMON /A/ TRAN(200,15),BLOCK(120,10),FACIL(35,5),STOR(35,7),
      IQUEUE(70,6),FEC(50,5),FECF,FECCL,CECA,CECF,CECL,CECA,
      2STATUS,MOVE,C1,TRANQ,BLOCKN,SIARTC,RN(8),TRANA,FUNC(20,54)
C
      DETERMINE STORAGE
      SNAC = BLOCK(BLOCKN,2)
      SNAN = BLOCK(BLOCKN,3)
      CALL SNA(SNAC,SNAN,RES)
      SNO = RES
C
C ARE YOU TRYING TO RELEASE MORE STORAGE THAN IN USE
C
      DETERMINE NO TO RELEASE
      SNAC = BLOCK(BLOCKN,4)
      SNAN = BLOCK(BLOCKN,5)
      CALL SNA(SNAC,SNAN,RESULT)
      IR = RESULT
      IF(IR.EQ.0) IR = 1
      IF(IR.GT.STOR(SNO,1)) GC TO 1000
C
      CALCULATE NEW CUMULATIVE TIME
      STOR(SNO,3) = STOR(SNO,3) + (C1 - STOR(SNO,4)) * STOR(SNO,1)
C
      DETERMINE NEW CURRENT AND REMAINING STORAGE UNITS
      STOR(SNO,2) = STOR(SNO,2) + IR
      STOR(SNO,1) = STOR(SNO,1) - IR

```

```

      STOR(SNO,4) = C1
C DID A DELAY CHAIN EXIST
      IF(STOR(SNO,7).EQ.0) GO TO 2000
      STOR(SNO,7) = 0
C SET STATUS CHANGE FLAG
      STATUS = 1
      GO TO 2000
1000 WRITE(6,1001)
1001 FORMAT('OTRANSACTION TRYING TO RELEASE MORE STORAGE THEN IN USE.')
      STOP
2000 MOVE = 1
C UPDATE BLOCK COUNTS
      ILST = TRAN(TRANC,5)
      TRAN(TRANC,5) = BLCKCN
      BLOCK(BLOCKN,9) = BLOCK(BLOCKN,9) + 1
      BLOCK(ILST,9) = BLCKCN(ILST,9) - 1
      BLOCK(BLOCKN,10) = BLOCK(BLOCKN,10) + 1
      TRAN(TRANC,1) = TRAN(TRANC,1) + 1
      RETURN
END
SUBROUTINE OUTPUT
C*****
C                                     OUTPUT
C
C                                     SUBROUTINE OUTPUT PRINTS THE SIMULATION STATISTICS.
C
C*****
      IMPLICIT INTEGER(C,F,Q,S,O)
      REAL FUNC
      INTEGER TYPE,TRANO,BLOCKN,BLOCKT
      COMMON /A/ TRAN(200,15),BLOCK(120,10),FACIL(35,5),STOR(35,7),
1        IQUEUE(70,6),FEC(50,5),FECF,FECCL,CECA,CECL,CECA,
2        STATUS,MOVE,C1,TRANO,BLOCKN,STARTC,RN(8),TRANA,FUNC(20,54)

```

```

C      OUTPUT BLOCK COUNTS
      XC1 = C1
      WRITE(6,10) C1
10     FORMAT('1CLOCK',I10)
      WRITE(6,20)
20     FORMAT(' BLOCK COUNTS')
25     WRITE(6,30)
30     FORMAT(' BLOCK CURRENT    TOTAL ')
      DO 50 I = 1,120
      IF(BLOCK(I,1).EQ.18.0) GO TO 155
      I2 = BLOCK(I,9)
      I3 = BLOCK(I,10)
      WRITE(6,70) I,I2,I3
70     FORMAT(' ',I5,I7,I8)
50     CONTINUE
155    DO 156 I=1,35
      IF(FACIL(I,5).NE.0) GO TO 160
156    CONTINUE
      GO TO 205
160    WRITE(6,170)
170    FORMAT('OFACILITY    AVERAGE',10X,'NUMBER',7X,' AVERAGE',7X,
1' SEIZING')
      WRITE(6,180)
180    FORMAT(' ',12X,'UTILIZATION',6X,'ENTRIES',6X,'TIME/TRAN',5X,
1' TRAN NO. ')
      DO 200 I = 1,35
      IF(FACIL(I,5).EQ.0) GO TO 200
      XF2 = FACIL(I,2)
      XF5 = FACIL(I,5)
      UTIL = XF2/XC1
      AVG = XF2/XF5
      WRITE(6,190) I,UTIL,FACIL(I,5),AVG,FACIL(I,1)
190    FORMAT('15,F12.3,I16,F16.3,I12)
200    CONTINUE
      STORAGE OUTPUT
C      205 CONTINUE

```

```

DO 206 I=1,35
  IF(STOR(I,5).NE.0) GC TO 207
  206 CONTINUE
  GO TO 255
  207 WRITE(6,210)
  210 FORMAT('0',23X,'AVERAGE',19X,'AVERAGE',
    16X,'CURRENT',19X,'MAXIMUM')
  WRITE(6,220)
  220 FORMAT(' STORAGE CAPACITY CONTENTS UTILIZATION',
    14X,'ENTRIES TIME/TRAN CONTENTS CCNTENTS')
  DO 240 I=1,35
    IF(STOR(I,5).EQ.0) GO TO 240
    XS3 = STOR(I,3)
    XS5 = STOR(I,5)
    XCAP = STOR(I,1) + STOR(I,2)
    CAP = XCAP
    AVGCCN = XS3/XC1
    UTIL = XS3/(XC1*XCAP)
    AVGTM = XS3/XS5
    WRITE(6,230) I,CAP,AVGCCN,UTIL,STOR(I,5),AVGTM,STOR(I,1),STOR(I,6)
  230 FORMAT('14,I13,F12.3,F12.3,I14,F13.3,I10,I12)
  240 CONTINUE
  C
  255 QUEUE OUTPUT
  CONTINUE
  DO 256 I=1,70
    IF(QUEUE(I,2).NE.0) GO TO 257
  256 CONTINUE
  GO TO 300
  257 WRITE(6,250)
  250 FORMAT('0',9X,'MAXIMUM',5X,'AVERAGE',4X,'TOTAL',8X,'ZERO',7X,
    1,'PERCENT',5X,'AVERAGE',6X,'$ AVERAGE',5X,'CURRENT')
  WRITE(6,260)
  260 FORMAT(' QUEUE CONTENTS ENTRIES ENTRIES ENTRIES',5X,
    1,' ZEROS TIME/TRAN TIME/TRAN',5X,'CONTENTS')
  DO 280 I = 1,70
    IF(QUEUE(I,2).EQ.0) GO TO 280

```



```

C      UPDATE BLOCK COUNTS
      ILST = TRAN(TRANQ,5)
      TRAN(TRANQ,5) = BLOCKN
      BLOCK(BLOCKN,9) = BLOCK(BLOCKN,9) + 1
      BLOCK(ILST,9) = BLOCK(ILST,9) - 1
      BLOCK(BLOCKN,10) = BLOCK(BLOCKN,10) + 1
      TRAN(TRANQ,1) = TRAN(TRANQ,1) + 1
      RESULT = BLOCK(BLOCKN,3)
      IF(RESULT.EQ.TRAN(TRANQ,3)) RETURN
      CEC(CECP,3) = RESULT
      TRAN(TRANQ,3) = RESULT
C      MOVE TRAN LAST IN NEW PRIORITY CLASS IN CEC
      CECT = CECF
      IR = RESULT
      IF(CECF.EQ.CECL) RETURN
      UNLINK FROM OLD LOCATION
      OLD1 = CEC(CECP,1)
      OLD2 = CEC(CECP,2)
      IF(OLD1.EQ.0) CECF = CEC(CECP,2)
      IF(OLD2.EQ.0) CECL=CEC(CECP,1)
      IF(OLD1.NE.0) CEC(OLD1,2) = OLD2
      IF(OLD2.NE.0) CEC(OLD2,1) = OLD1
      CECT = CECF
50    IF(CEC(CECT,3).GE.IR) GO TO 100
      OLD1 = CFC(CECT,1)
      CEC(CECP,1) = CEC(CECT,1)
      CEC(CECP,2) = CECT
      CEC(CECT,1) = CECF
      IF(OLD1.NE.0) CEC(OLD1,2) = CECF
      IF(OLD1.EQ.0) CECF = CECF
      STATUS = 1
      RETURN
100   IF(CEC(CECT,2).EQ.0) GO TO 150
      CECT = CEC(CECT,2)
      GO TO 50
150   OLD2 = CEC(CECT,2)

```



```

C INCREMENT ENTRY COUNT
  IF (QUEUE(Q,5).GT.QUEUE(Q,6)) QUEUE(Q,6) = QUEUE(Q,5)
C IS NEW CURRENT > OLD MAX
  QUEUE(Q,2) = QUEUE(Q,2) + 1
  QUEUE(Q,1) = C1
  MOVE = 1
  TRAN(TRANQ,1) = TRAN(TRANQ,1) + 1
  UPDATE BLCKCK COUNTS
  ILST = TRAN(TRANQ,5)
  TRAN(TRANQ,5) = BLOCKN
  BLOCK(BLOCKN,9) = BLCKCK(BLOCKN,9) + 1
  BLOCK(ILST,9) = BLOCK(ILST,9) - 1
  BLOCK(BLOCKN,10) = BLOCK(BLOCKN,10) + 1
  RETURN
END
SUBROUTINE RANDU(IN,IY,YFL)
C*****
C
C
C
C
C
C
C
C*****
      RANDU
      SUBROUTINE RANDU CALCULATES A RANDOM NUMBER USING THE
      SPECIFIED RANDOM NUMBER SEED.
C*****
      IMPLICIT INTEGER(C,F,Q,S,O)
      REAL FUNC
      INTEGER TYPE,TRANQ,BLOCKN,BLOCKT
      COMMON /A/ TRAN(200,15),BLOCK(120,10),FACIL(35,5),STOR(35,7),
      1QUEUE(70,6),FEC(50,5),FECF,FECCL,FECA,CEC(50,4),CECF,CECL,CECA,
      2STATUS,MOVE,C1,TRANQ,BLCKCKN,STARTC,RN(8),TRANA,FUNC(20,54)
      IX = RN(IN)
      IY = IX * 1220703125
      IF(IY) 1,2,2
      1 IY = IY + 2147483647 + 1
      2 YFL = IY
      YFL = YFL * 0.4656613E-9
      RN(IN) = IY

```



```

      RETURN
      END
      SUBROUTINE RELEAS
C*****
C
C
C
C
C
C
C*****
      SUBROUTINE RELEAS RELEASES THE FACILITY SPECIFIED IN FIELD A.
C*****
      IMPLICIT INTEGER(C,F,Q,S,O)
      REAL FUNC
      INTEGER TYPE,TRANQ,BLOCKN,BLOCKT
      COMMON /A/ TRAN(200,15),BLOCK(120,10),FACIL(35,5),STOR(35,7),
      1QUEUE(70,6),FEC(50,5),FECF,FECCL,CECA,CECF,CECL,CECA,
      2STATUS,MOVE,C1,TRANO,BLOCKN,STARIC,RN(8),TRANA,FUNC(20,54)
C
      DETERMINE FACILITY
      SNAC = BLOCK(BLOCKN,2)
      SNAN = BLOCK(BLOCKN,3)
      CALL SNA(SNAC,SNAN,RES)
      FAC = RES
C
      C TRAN# SAME AS SEIZING TRAN
      IF(FACIL(FAC,1).NE.TRANO)GO TO 1000
C
      C CALCULATE NEW CUMULATIVE TIME
      FACIL(FAC,2) = FACIL(FAC,2) + C1 - FACIL(FAC,3)
C
      C SET FACILITY NOT IN USE
      FACIL(FAC,1) = 0
      FACIL(FAC,3) = C1
      IF(FACIL(FAC,4).EQ.0) GC TO 100
C
      C ELIMINATE DELAY CHAIN
      FACIL(FAC,4) = 0
C
      C SET STATUS CHANGE FLAG IF OTHER TRANS WAITING TO SEIZE FACILITY
      STATUS = 1
      100 MOVE = 1
C
      C UPDATE BLOCK COUNTS

```



```

C      IS CURRENT EVENT CHAIN NULL
      IF(CECT.NE.0) GO TO 100
C      SET UP CEC LINKAGE
      CECF = CECA
      CECL = CECA
      OLDNXT = CEC(CECA,2)
      CEC(CECF,1) = 0
      CEC(CECF,2) = 0
      GO TO 200
50      IF(CEC(CECT,2).EQ.0) GO TO 75
      CECT = CEC(CECT,2)
      GO TO 100
C      WAS TRAN ADDED AT END OF CHAIN
75      OLDNXT = CEC(CECA,2)
      OLD2 = CEC(CECT,2)
      CEC(CECA,1) = CECT
      CEC(CECA,2) = CEC(CECT,2)
      CEC(CECT,2) = CECA
      IF(OLD2.NE.0) CEC(CECT,1) = CECA
      IF(OLD2.EQ.0) CECL = CECA
      GO TO 200
C      INSERT TRAN IN PROPER PRIORITY IN CEC
100     IF(CEC(CECT,3).GE.FEC(FECF,4)) GO TO 50
      OLDNXT = CEC(CECA,2)
C      SET UP LINKAGE OF NEW TRAN
      OLD1 = CEC(CECT,1)
      CEC(CECA,1) = CEC(CECT,1)
      CEC(CECA,2) = CECT
      CEC(CECT,1) = CECA
      IF(OLD1.NE.0) CEC(OLD1,2) = CECA
      IF(OLD1.EQ.0) CECF = CECA
C      MOVE LP AVAILABLE PTR
200     CEC(CECA,3) = FEC(FECF,4)
      CEC(CECA,4) = FEC(FECF,5)
      CECA = OLDNXT
C      UNLINK FROM FEC
      FECF = FEC(FECF,2)
      FECF(FECF,2) = FECA

```

```

FECA = FECF
FECF = FECT
IF(FECF.NE.0) FEC(FECF,1) = 0
IF(FECF.EQ.0) FECL = 0
CECT = FECF
C   IS BDT OF NEXT TRAN EQUAL CLOCK
IF(FECF.EQ.0) GO TO 300
IF(FEC(FECF,3).EQ.C1) GO TO 100
C   SET STATUS CHANGE OFF
300 STATUS = 0
C   GET FIRST CEC TRAN
CECT = FECF
310 TRANO = CEC(CECT,4)
SSI ON
C   IF(TRAN(TRANO,2).EQ.1) GO TO 400
GO TO 400
320 CECT = CEC(CECT,2)
LAST CEC TRAN
IF(CECT.EQ.0) GO TO 10
GO TO 310
C   MOVE TRAN - DETERMINE BLOCK TYPE
BLOCKN = TRAN(TRANO,1)
BLOCKT = BLOCK(BLOCKN,1)
IF(BLOCKT.EQ.1) CALL GENERA(2,BLOCKN)
IF(BLOCKT.EQ.2) CALL SEIZE
IF(BLOCKT.EQ.3) CALL RELEAS
IF(BLOCKT.EQ.5) CALL ENTER
IF(BLOCKT.EQ.6) CALL LEAVE
IF(BLOCKT.EQ.7) CALL ADVANC(CECT,CD)
IF(BLOCKT.EQ.8) CALL TERMIN(CECT,CD)
IF(BLOCKT.EQ.9) CALL ASSIGN
IF(BLOCKT.EQ.10) CALL PRIORI(CECT)
IF(BLOCKT.EQ.11) CALL TRANSF
IF(BLOCKT.EQ.12) CALL TEST
IF(BLOCKT.EQ.15) CALL QUEUES
IF(BLOCKT.EQ.16) CALL DEPART

```

```

C      DID TRAN ENTER NEXT BLOCK
      IF(MOVE.NE.1) GO TO 500
      MOVE = 0
C      TERMINATE BLOCK
      IF(BLOCKT.NE.8) GO TO 450
      IF(STARTC.LE.0) RETURN
      GO TO 500
C      ADVANCE BLOCK
      450 IF(BLOCKT.NE.7) GO TO 400
      IS STATUS CHANGE CN
      500 IF(STATUS.EQ.1) GO TO 300
      IS END OF CEC REACHED
      IF(CD.NE.1) CECT = CEC(CECT,2)
      IF(CECT.EQ.0 ) GO TO 10
      CD = 0
      GO TO 310
      END
      SUBROUTINE SEIZE
C*****
C      SEIZE
C      SUBROUTINE SEIZE ATTEMPTS TO SEIZE THE FACILITY SPECIFIED
C      IN FIELD A. IF THE FACILITY IS CURRENTLY IN USE THE
C      TRANSACTION IS BLOCKED.
C*****
      IMPLICIT INTEGER(C,F,Q,S,O)
      REAL FUNC
      INTEGER TYPE,TRANC,BLOCKN,BLOCKT
      COMMON /A/ TRAN(200,15),BLOCK(120,10),FACIL(35,5),STOR(35,7),
      1QUEUE(70,6),FEC(50,5),FECF,FECF,FECA,CEC(50,4),CECF,CECL,CECA,
      2STATUS,MOVE,C1,TRANO,BLOCKCN,STARTC,RN(8),TRANA,FUNC(20,54)
      DETERMINE FACILITY
C      10 SNAC = BLOCK(BLOCKN,2)
      SNAN = BLOCK(BLOCKN,3)
      CALL SNA(SNAC,SNAN,RES)
      FAC = RES

```

```

C      IS FACILITY IN USE
      IF(FACIL(FAC,1).EQ.0) GC TO 1000
C SET UP DELAY CHAIN
      FACIL(FAC,4) = 1
      GO TO 2000
C      NOT IN USE
C PLACE TRAN # IN F1
      1000 FACIL(FAC,1) = TRANO
      TRAN(TRANO,2) = 1
C PLACE CLOCK TIME IN F3
      FACIL(FAC,3) = C1
C INCREMENT ENTRY COUNT
      FACIL(FAC,5) = FACIL(FAC,5) + 1
      TRAN(TRANO,1) = TRAN(TRANO,1) + 1
      MOVE = 1
C      UPDATE BLOCK COUNTS
      ILST = TRAN(TRANO,5)
      TRAN(TRANO,5) = BLCCKN
      BLOCK(BLOCKN,9) = BLOCK(BLOCKN,9) + 1
      BLOCK(ILST,9) = BLOCK(ILST,9) - 1
      BLOCK(BLOCKN,10) = BLOCK(BLOCKN,10) + 1
      TRAN(TRANO,2) = 1
2000 RETURN
      END
      SUBROUTINE SNA(SNAC,SNAN,RESULT)
C*****
C      SNA
C
C      SUBROUTINE DETERMINES THE VALUE CF THE SNA.
C*****
      IMPLICIT INTEGER(C,F,Q,S,D)
      REAL FUNC
      INTEGER SNAC,SNAN
      INTEGER TYPE,TRANO,BLOCKN,BLOCKT

```

**THIS BOOK
CONTAINS
NUMEROUS
PAGES WITH
THE ORIGINAL
PRINTING ON
THE PAGE BEING
CROOKED.**

**THIS IS THE
BEST IMAGE
AVAILABLE.**

```

COMMON /A/ TRAN(200,15),BLOCK(120,10),FACIL(35,5),STOR(35,7),
1QUEUE(70,6),FEC(50,5),FECF,FECF,FECF,CEC(50,4),CECF,CECL,CECA,
2STATUS,MOVE,C1,TRANQ,BLOCKN,STARTC,RN(8),TRANA,FUNC(20,54)

C
PARAMETER
IF(SNAC.NE.1) GO TO 50
RESULT = TRAN(TRANQ,SNAN+5)
RETURN

C
RANDOM NUMBER
50 IF(SNAC.NE.4) GO TC 100
I = SNAN
CALL RANDU(I,IY,YFL)
RESULT = YFL
RETURN

100 IF(SNAC.NE.3.AND.SNAC.NE.0) GO TO 200
RESULT = SNAN
RETURN

C
PRIORITY
200 IF(SNAC.NE.2) GO TO 300
RESULT = TRAN(TRANQ,3)
RETURN

C
CLOCK
300 IF(SNAC.NE.5) GO TO 400
RESULT = C1
RETURN

C
FUNCTION
400 IF(SNAC.NE.6) GO TO 500
CALL FUNC(RESULT,SNAN)
RETURN

C
VARIABLE
500 IF(SNAC.NE.7) GO TO 600
CALL VARIAB(RESULT,SNAN)
RETURN

C
CONTENTS OF STORAGE
600 IF(SNAC.NE.8) GO TO 700
RESULT = STOR(SNAN,1)
RETURN

```



```

C      REMAINING STORAGE
700  IF(SNAC.NE.9) GO TO 800
    RESULT = STOR(SNAN,2)
    RETURN
C      UTILIZATION OF STORAGE
800  IF(SNAC.NE.10) GO TO 900
    RESULT = STOR(SNAN,3)/(C1*(STOR(SNAN,1)+STOR(SNAN,2)))
    RETURN
C      AVERAGE OF STORAGE
900  IF(SNAC.NE.11) GO TO 1000
    RESULT = STOR(SNAN,3)/C1
    RETURN
C      MAXIMUM OF STORAGE
1000 IF(SNAC.NE.12) GO TO 1100
    RESULT = STOR(SNAN,6)
    RETURN
C      NO OF ENTRIES IN STORAGE
1100 IF(SNAC.NE.13) GO TO 1200
    RESULT = STOR(SNAN,5)
    RETURN
C      AVERAGE TIME IN STORAGE
1200 IF(SNAC.NE.14) GO TO 1300
    RESULT = STOR(SNAN,3)/STOR(SNAN,5)
    RETURN
C      STATUS OF FACILITY
1300 IF(SNAC.NE.15) GO TO 1400
    RESULT = 1
    IF(FACIL(SNAN,1).EQ.0) RESULT = 0
    RETURN
C      FACILITY UTILIZATION
1400 IF(SNAC.NE.16) GO TO 1500
    RESULT = FACIL(SNAN,2)/C1
    RETURN
C      NO OF ENTRIES IN FACILITY

```

```

1500 IF(SNAC.NE.17) GO TO 1600
    RESULT = FACIL(SNAN,5)
    RETURN
C
    AVERAGE TIME IN FACILITY
1600 IF(SNAC.NE.18) GO TO 1700
    RESULT = FACIL(SNAN,2)/FACIL(SNAN,5)
    RETURN
C
    QUEUE LENGTH
1700 IF(SNAC.NE.19)GO TO 1800
    RESULT = QUEUE(SNAN,5)
    RETURN
C
    AVE QUEUE LENGTH
1800 IF(SNAC.NE.20) GO TO 1900
    RESULT = QUEUE(SNAN,4)/C1
    RETURN
C
    MAXIMUM QUEUE LENGTH
1900 IF(SNAC.NE.21) GO TO 2000
    RESULT = QUEUE(SNAN,6)
    RETURN
C
    NO OF ENTRIES IN QUEUE
2000 IF(SNAC.NE.22) GO TO 2100
    RESULT = QUEUE(SNAN,2)
    RETURN
C
    NO OF ZERC ENTRIES IN QUEUE
2100 IF(SNAC.NE.23) GO TO 2200
    RESULT = QUEUE(SNAN,3)
    RETURN
C
    AVERAGE QUEUE TIME WITH ZERO
2200 IF(SNAC.NE.24) GO TO 2300
    RESULT = QUEUE(SNAN,4)/QUEUE(SNAN,2)
    RETURN
C
    AVERAGE QUEUE TIME WITH NO ZEROS
2300 IF(SNAC.NE.25) GO TO 2400
    RESULT = QUEUE(SNAN,4)/((QUEUE(SNAN,2)-QUEUE(SNAN,3))
    RETURN

```

```

C      ERROR
      2400 WRITE(6,2500) SNAC
      2500 FORMAT('OINVALID SNA ',I4)
      STOP
      END
      SUBROUTINE TERMIN(CECT, ID)
C*****
C      TERMIN
C
C      SUBROUTINE TERMIN REMOVES A TRANSACTION FROM THE CEC AND
C      DECREMENTS THE START COUNT.
C*****
C      IMPLICIT INTEGER(C,F,Q,S,O)
      REAL FUNC
      INTEGER CECT
      INTEGER TYPE, TRANO, BLOCKN, BLOCKT
      COMMON /A/ TRAN(200,15), BLOCK(120,10), FACIL(35,5), STOR(35,7),
      1 QUEUE(70,6), FEC(50,5), FECF, FECL, FECA, CEC(50,4), CECF, CECL, CECA,
      2 STATUS, MOVE, CI, TRANO, BLOCKN, STARTC, RN(8), TRANA, FUNC(20,54)
C      UPDATE BLOCK COUNTS
      ILST = TRAN(TRANO,5)
      BLOCK(ILST,9) = BLOCK(ILST,9) - 1
      BLOCK(BLOCKN,10) = BLOCK(BLOCKN,10) + 1
      MOVE = 1
C      UNLINK BLOCK FROM CEC
      OLD1 = CEC(CECT,1)
      OLD2 = CEC(CECT,2)
      IF(OLD1.NE.0) CEC(OLD1,2) = OLD2
      IF(OLD2.NE.0) CEC(OLD2,1) = OLD1
      CEC(CECT,2) = CECA
      CECA = CECT
      IF(CECT.EQ.CECF) CECF = OLD2
      IF(CECT.EQ.CECL) CECL = OLD1
      IF(CECL.NE.0) GO TO 100
      CECT = 0
      ID = 1

```

```

GO TO 200
100 ID = 1
    CECT = OLD2
    DECREMENT START COUNT
200 SNAC = BLOCK(BLOCKN,2)
    SNAN = BLOCK(BLOCKN,3)
    CALL SNA(SNAC,SNAN,RESULT)
    STARTC = STARTC - RESULT
    REMOVE TRANSACTION
    TRAN(TRANC,1) = TRANA
    TRANA = TRANO
    RETURN
END

```

SUBROUTINE TEST

TEST

SUBROUTINE TEST TESTS THE VALUE IN FIELD A AGAINST THE VALUE IN FIELD B. IF THE CONDITION IS MET THE TRANSACTION NEXT BLOCK IS CHANGED TO THE BLOCK SPECIFIED IN FIELD C OTHERWISE IT IS INCREMENTED BY ONE.

IMPLICIT INTEGER(C,F,Q,S,O)

REAL FUNC

INTEGTR TYPE,TRANO,BLOCKN,BLOCKT

```
COMMON /A/ TRAN(200,15),BLOCK(120,10),FACIL(35,5),STOR(35,7),
```

IQUEUE(70,6),FEC(50,5),FECF,FECI,FECA,CEC(50,4),CECF,CECL,CECA,

2STATUS,MCVE,C1,TRANO,BLOCKN,STARTC,RN(8),TRANA,FUNC(20,54)

$$\text{TRAN}(\text{TRANO},1) = \text{TRAN}(\text{TRANO},1) + 1$$

MOVE = 1

```
ILST = TRAN(TRANQ,5)
```

TRAN(TRANC,5) = BLOCKN

$$\text{BLOCK}(\text{BLOCKN},9) = \text{BLOCK}(\text{BLOCKN},9) + 1$$

```

BLOCK(ILST,9) = BLOCK(ILST,9) - 1

```

```

BLOCKN(BLOCKN,10) = BLOCK(BLOCKN,10) + 1

```

```

C      DETERMINE FIELD A
      SNAC = BLOCK(BLOCKN,2)
      SNAN = BLOCK(BLOCKN,3)
      CALL SNA(SNAC,SNAN,A)
      DETERMINE FIELD B
      SNAC = BLOCK(BLOCKN,4)
      SNAN = BLOCK(BLOCKN,5)
      CALL SNA(SNAC,SNAN,B)
      IF(BLOCK(BLOCKN,8).NE.1.0) GO TO 100
      TEST L
      IF (A.GE.B) TRAN(TRANQ,1) = BLOCK(BLOCKN,6)
      RETURN
100  IF(BLOCK(BLOCKN,8).NE.2.0) GO TO 200
      TEST LE
      IF(A.GT.B) TRAN(TRANC,1) = BLOCK(BLOCKN,6)
      RETURN
200  IF(BLOCK(BLOCKN,8).NE.3.0) GO TO 300
      TEST E
      IF(A.NE.B) TRAN(TRANC,1) = BLOCK(BLOCKN,6)
      RETURN
300  IF(BLOCK(BLOCKN,8).NE.4.0) GO TO 400
      TEST NE
      IF(A.EQ.B) TRAN(TRANQ,1) = BLOCK(BLOCKN,6)
      RETURN
400  IF(BLOCK(BLOCKN,8).NE.5.0) GO TO 500
      TEST G
      IF(A.LE.B) TRAN(TRANC,1) = BLOCK(BLOCKN,6)
      RETURN
500  IF(BLOCK(BLOCKN,8).NE.6.0) GO TO 600
      TEST GF
      IF(A.LT.B) TRAN(TRANC,1) = BLOCK(BLOCKN,6)
      RETURN
      EPROR
C      600 WRITE(6,700) BLOCK(BLOCKN,8)
      700 FORMAT('OINVALID TEST RELATION',I3)
      STOP
      END

```



```
      RESULT = BLOCK(BLOCKN,3)  
      IF(YFL .LE.RESULT) GO TO 200  
      TRAN(TRANC,1) = BLOCK(BLOCKN,5)  
      RETURN  
200 TRAN(TRANC,1) = BLOCK(BLOCKN,6)  
      RETURN  
      END
```

NUMBER	*LCC	OPERATION	A,B,C,D,E,F,G,H,I	COMMENTS	NUMBER
1		SIMULATE			1
2		GENERATE	15,5	JOBS ARRIVE EVERY 10 TO 20 MINUTES	2
3		TEST E	F2,0,CLKC	IS CLERK B BUSY	3
4		QUEUE	ACLK	ENTER LINE FOR CLERK A	4
5		SEIZE	ACLK	CLERK A BEGINS PROCESSING	5
6		DEPART	ACLK	LEAVE LINE FOR CLERK A	6
7		ADVANCE	10,3	CLERK TAKES 7 TO 13 MIN TO PROCESS JOB	7
8		RELEASE	ACLK	CLERK A COMPLETES JOB	8
9		QUEUE	BCLK	ENTER LINE FOR CLERK B	9
10		SEIZE	BCLK	CLERK B BEGINS PROCESSING	10
11		DEPART	BCLK	LEAVE LINE FOR CLERK B	11
12		ADVANCE	5,2	CLERK B PROCESSES JOBS IN 3 TO 7 MIN	12
13		RELEASE	BCLK	CLERK B COMPLETES JOB	13
14		TERMINATE	1	JOB LEAVES OFFICE	14
15	CLKC	QUEUE	CCLK	ENTER LINE FOR CLERK C	15
16		SEIZE	CCLK	CLERK C BEGINS PROCESSING	16
17		DEPART	CCLK	LEAVE LINE FOR CLERK C	17
18		ADVANCE	20,10	CLERK C TAKES 10 TO 30 MIN TO PROCESS	18
19		RELEASE	CCLK	CLERK C COMPLETES JOB	19
20		TERMINATE	1	JOB LEAVES OFFICE	20
21		START	1000	RUN SIMULATION FOR 1000 JOBS	21
22		END			22

NUMBER	*LCC	OPERATION	A,B,C,D,E,F,G,H,I	COMMENTS	NUMBER
1		SIMULATE			1
2		GENERATE	15,5	JOBS ARRIVE EVERY 10 TO 20 MINUTES	2
3		TEST E	F2,0,14	IS CLERK B BUSY	3
4		QUEUE	1	ENTER LINE FOR CLERK A	4
5		SEIZE	1	CLERK A BEGINS PROCESSING	5
6		DEPART	1	LEAVE LINE FOR CLERK A	6
7		ADVANCE	10,3	CLERK TAKES 7 TO 13 MIN TO PROCESS JOB	7
8		RELEASE	1	CLERK A COMPLETES JOB	8
9		QUEUE	2	ENTER LINE FOR CLERK B	9
10		SEIZE	2	CLERK B BEGINS PROCESSING	10
11		DEPART	2	LEAVE LINE FOR CLERK B	11
12		ADVANCE	5,2	CLERK B PROCESSES JOBS IN 3 TO 7 MIN	12
13		RELEASE	2	CLERK B COMPLETES JOB	13
14		TERMINATE	1	JOB LEAVES OFFICE	14
15		QUEUE	3	ENTER LINE FOR CLERK C	15
16		SEIZE	3	CLERK C BEGINS PROCESSING	16
17		DEPART	3	LEAVE LINE FOR CLERK C	17
18		ADVANCE	20,10	CLERK C TAKES 10 TO 30 MIN TO PROCESS	18
19		RELEASE	3	CLERK C COMPLETES JOB	19
20		TERMINATE	1	JOB LEAVES OFFICE	20
21		START	1000	RUN SIMULATION FOR 1000 JOBS	21
22		END			22

BLOCK COUNTS		TOTAL	BLOCK CURRENT	TOTAL	BLOCK CURRENT	TOTAL	BLOCK CURRENT	TOTAL
1	0	1000	11	0	658			
2	0	1000	12	0	658			
3	0	658	13	0	658			
4	0	658	14	0	302			
5	0	658	15	0	302			
6	0	658	16	0	302			
7	0	658	17	0	202			
8	0	658	18	0	202			
9	0	658	19	0	302			
10	0	658						

 *
 *
 *
 *
 *

 FACILITIES

 *
 *
 *
 *

FACILITY	NUMBER ENTRIES	-AVERAGE UTILIZATION DURING-		CURRENT STATUS	PERCENT AVAILABILITY	TRANSACTION NUMBER SEIZING PREEMPTING
		AVERAGE TIME/TRAN	TOTAL TIME			
ACLK	698	10.024	-469		100.0	
BCLK	658	5.017	-225		100.0	
CCLK	302	15.553	-405		100.0	

 *
 *
 *
 *

 QUEUES

 *
 *
 *

QUEUE	MAXIMUM CCNTENTS	AVERAGE CONTENTS	TOTAL ENTRIES	ZERO ENTRIES	PERCENT ZEROS	AVERAGE TIME/TRANS	\$AVERAGE TIME/TRANS	TABLE NUMBER	CURRENT CONTENTS
BCLK	1	.030	698	698	100.0	-000	.000		
CCLK	1	.003	302	287	95.0	-182	3.666		

\$AVERAGE TIME/TRANS = AVERAGE TIME/TRANS EXCLUDING ZERO ENTRIES

***** TOTAL RUN TIME (INCLUDING ASSEMBLY) = .05 MINUTES *****

ENC

BLOCK	TYPE	A	B	C	D	E
0	1.000	0.0	0.0	0.0	0.0	0.0
0	12.000	15.000	0.0	0.0	0.0	3.000
0	15.000	0.0	0.0	0.0	0.0	0.0
0	2.000	0.0	0.0	0.0	0.0	0.0
0	16.000	0.0	0.0	0.0	0.0	0.0
0	7.000	0.0	0.0	0.0	0.0	0.0
0	3.000	0.0	0.0	0.0	0.0	0.0
0	15.000	0.0	0.0	0.0	0.0	0.0
0	2.000	0.0	0.0	0.0	0.0	0.0
0	16.000	0.0	0.0	0.0	0.0	0.0
0	7.000	0.0	0.0	0.0	0.0	0.0
0	3.000	0.0	0.0	0.0	0.0	0.0
0	8.000	0.0	0.0	0.0	0.0	0.0
0	15.000	0.0	0.0	0.0	0.0	0.0
0	2.000	0.0	0.0	0.0	0.0	0.0
0	16.000	0.0	0.0	0.0	0.0	0.0
0	7.000	0.0	0.0	0.0	0.0	0.0
0	3.000	0.0	0.0	0.0	0.0	0.0
0	8.000	0.0	0.0	0.0	0.0	0.0
0	17.000	0.0	0.0	0.0	0.0	0.0
0	18.000	0.0	0.0	0.0	0.0	0.0

CLOCK COUNTS
BLOCK CURRENT TOTAL

1	1	1002
2	0	1001
3	0	677
4	0	677
5	0	677
6	0	677
7	0	677
8	0	677
9	0	677
10	0	677
11	0	677
12	0	677
13	0	677
14	0	324
15	0	324
16	0	324
17	1	324
18	0	323
19	0	323

FACILITY AVERAGE UTILIZATION NUMBER ENTRIES SEIZING TRAN NO.

1	0.454	677	10.057	0
2	0.226	677	5.019	0
3	0.451	324	20.951	4

QUEUE MAXIMUM CONTENTS AVERAGE ENTRIES TOTAL ENTRIES ZERO ENTRIES PERCENT ZEROS

1	1	0.004	677	643	94.978
2	1	0.0	677	677	100.000
3	1	0.005	324	299	92.284

AVERAGE TIME/TRAN AVERAGE TIME/TRAN CURRENT CONTENTS

0.080	1.588	0
0.0	0.0	0
0.216	2.800	0

\$ AVERAGE TIME/TRAN EXCLUDES ZERO DELAY ENTRIES

BLCK NUMBER	*LOC	OPERATION	A,B,C,D,E,F,G,H,I	COMMENTS	STATEMENT NUMBER
1	REGCK	SIMULATE			1
2		STORAGE	2	TWO REGULAR CHECKOUT COUNTERS	2
3		GENERATE	3,2	CUSTOMERS ARRIVE FROM 1 TO 5 MINUTES APART	3
4	EXPRS	TRANSFER	.32,REG,EXPRS		4
5		QUEUE	EXPRS	ENTER EXPRESS QUEUE	5
6		SEIZE	EXPRS	ENTER EXPRESS CHECKOUT	6
7		DEPART	EXPRS	LEAVE EXPRESS QUEUE	7
8		ADVANCE	2	TWO MINUTES REQUIRED FOR EXPRESS CHECKOUT	8
9		RELEASE	EXPRS	LEAVE EXPRESS CHECKOUT	9
10	REG	TRANSFER	*TERM		10
11		QUEUE	REGCK	ENTER REGULAR QUEUE	11
12		ENTER	REGCK	ENTER REGULAR CHECKOUT	12
13		DEPART	REGCK	DEPART REGULAR QUEUE	13
14		ADVANCE	8,6	2 TO 14 MIN REQUIRED IN REGULAR CHECKOUT	14
15		LEAVE	REGCK	LEAVE REGULAR CHECKOUT	15
16	TERM	TERMINATE	1	CUSTOMER LEAVES STORE	16
17		START	100	RUN 100 CUSTOMERS	17
18		END			18

BLCK NUMBER	*LOC	OPERATION	A,B,C,D,E,F,G,H,I	COMMENTS	STATEMENT NUMBER
1	REGCK	SIMULATE			1
2		STORAGE	2	TWO REGULAR CHECKOUT COUNTERS	2
3		GENERATE	3,2	CUSTOMERS ARRIVE FROM 1 TO 5 MINUTES APART	3
4	EXPRS	TRANSFER	.32,REG,EXPRS		4
5		QUEUE	EXPRS	ENTER EXPRESS QUEUE	5
6		SEIZE	EXPRS	ENTER EXPRESS CHECKOUT	6
7		DEPART	EXPRS	LEAVE EXPRESS QUEUE	7
8		ADVANCE	2	TWO MINUTES REQUIRED FOR EXPRESS CHECKOUT	8
9		RELEASE	EXPRS	LEAVE EXPRESS CHECKOUT	9
10	REG	TRANSFER	*TERM		10
11		QUEUE	REGCK	ENTER REGULAR QUEUE	11
12		ENTER	REGCK	ENTER REGULAR CHECKOUT	12
13		DEPART	REGCK	DEPART REGULAR QUEUE	13
14		ADVANCE	8,6	2 TO 14 MIN REQUIRED IN REGULAR CHECKOUT	14
15		LEAVE	REGCK	LEAVE REGULAR CHECKOUT	15
16	TERM	TERMINATE	1	CUSTOMER LEAVES STORE	16
17		START	100	RUN 100 CUSTOMERS	17
18		END			18

BLOCK TYPE	A	B	C	D	E
1	4.000	0.0			0.0
0	1.000	0.0			0.0
0	11.000	0.0			0.0
0	15.000	0.0			0.0
C	2.000	0.0			0.0
0	16.000	0.0			0.0
0	7.000	0.0			0.0
0	3.000	0.0			0.0
0	11.000	0.0			0.0
0	15.000	0.0			0.0
0	5.000	0.0			0.0
0	16.000	0.0			0.0
0	7.000	0.0			0.0
0	6.000	0.0			0.0
0	8.000	0.0			0.0
0	17.000	0.0			0.0
0	18.000	0.0			0.0

CLCCK BLOCK	COUNTS CURRENT	278 TOTAL
1	1	102
2	0	101
3	0	37
4	0	37
5	0	37
6	0	37
7	0	37
8	0	37
9	0	64
10	0	64
11	0	64
12	1	64
13	0	63
14	0	100

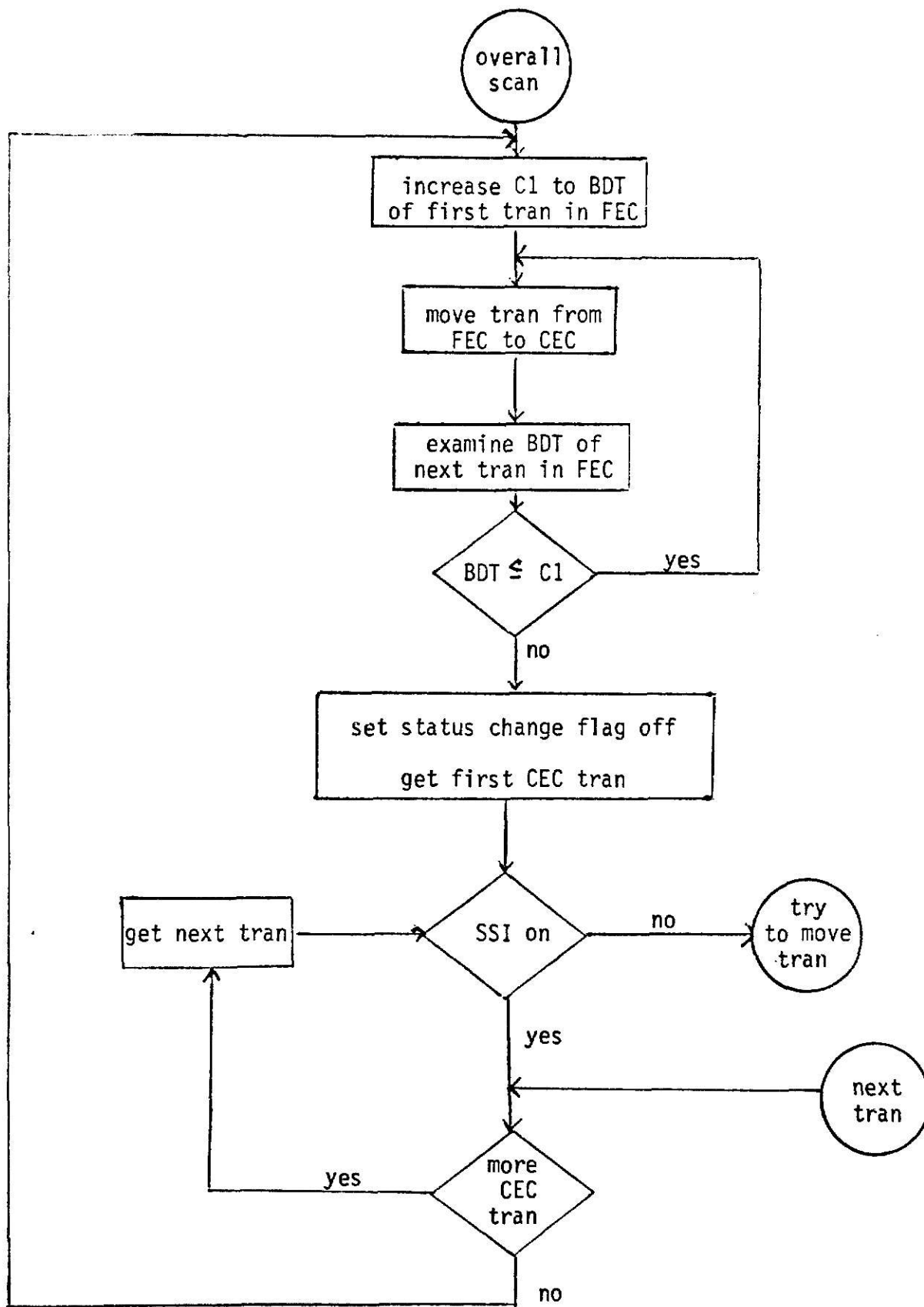
FACILITY	AVERAGE UTILIZATION	NUMBER ENTRIES	AVERAGE TIME/TRAN	SEIZING TRAN NO.
1	0.266	37	2.000	0

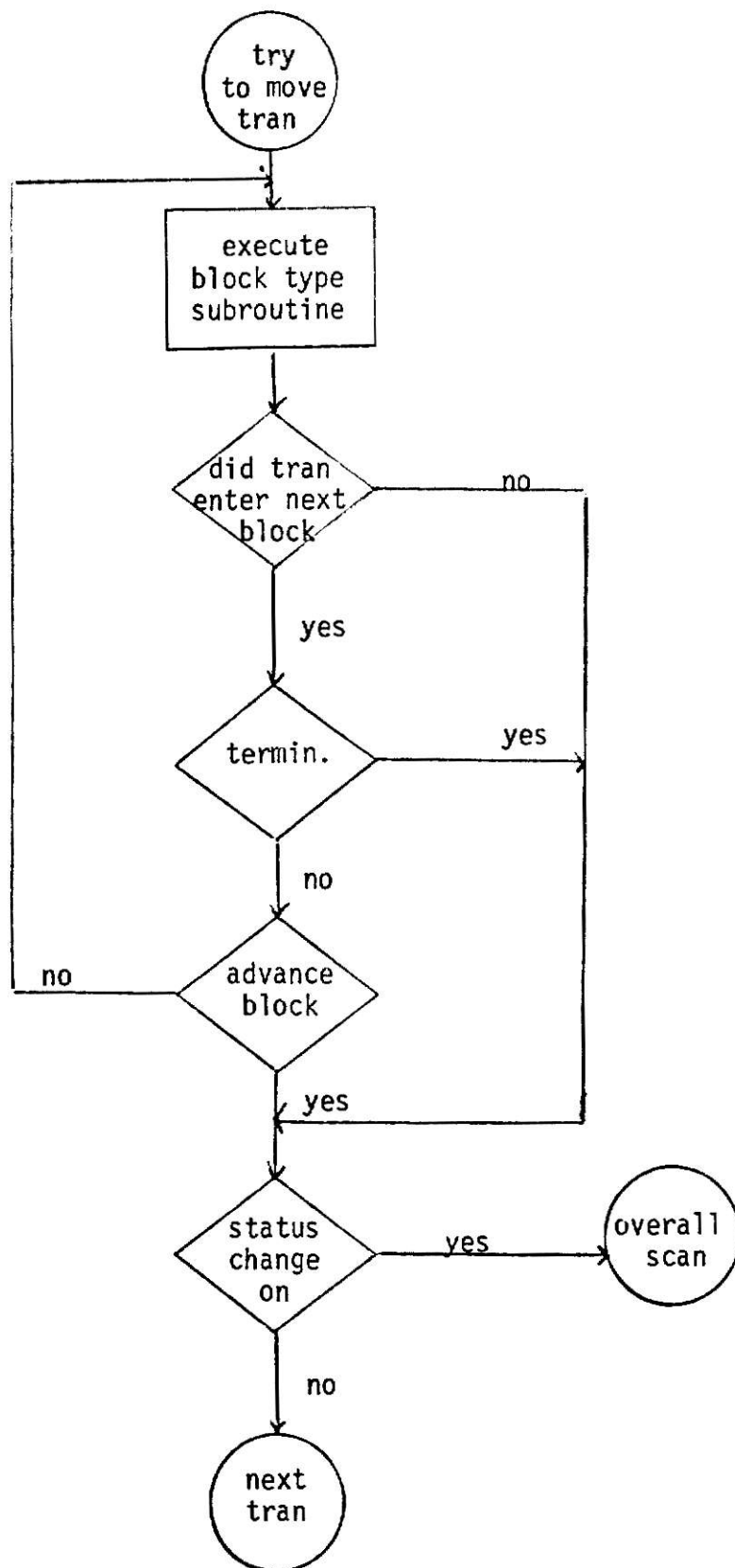
STORAGE	CAPACITY	AVERAGE CONTENTS	AVERAGE UTILIZATION	ENTRIES	AVERAGE TIME/TRAN	CURRENT CONTENTS	MAXIMUM CONTENTS
1	2	1.719	0.860	64	7.469	1	2

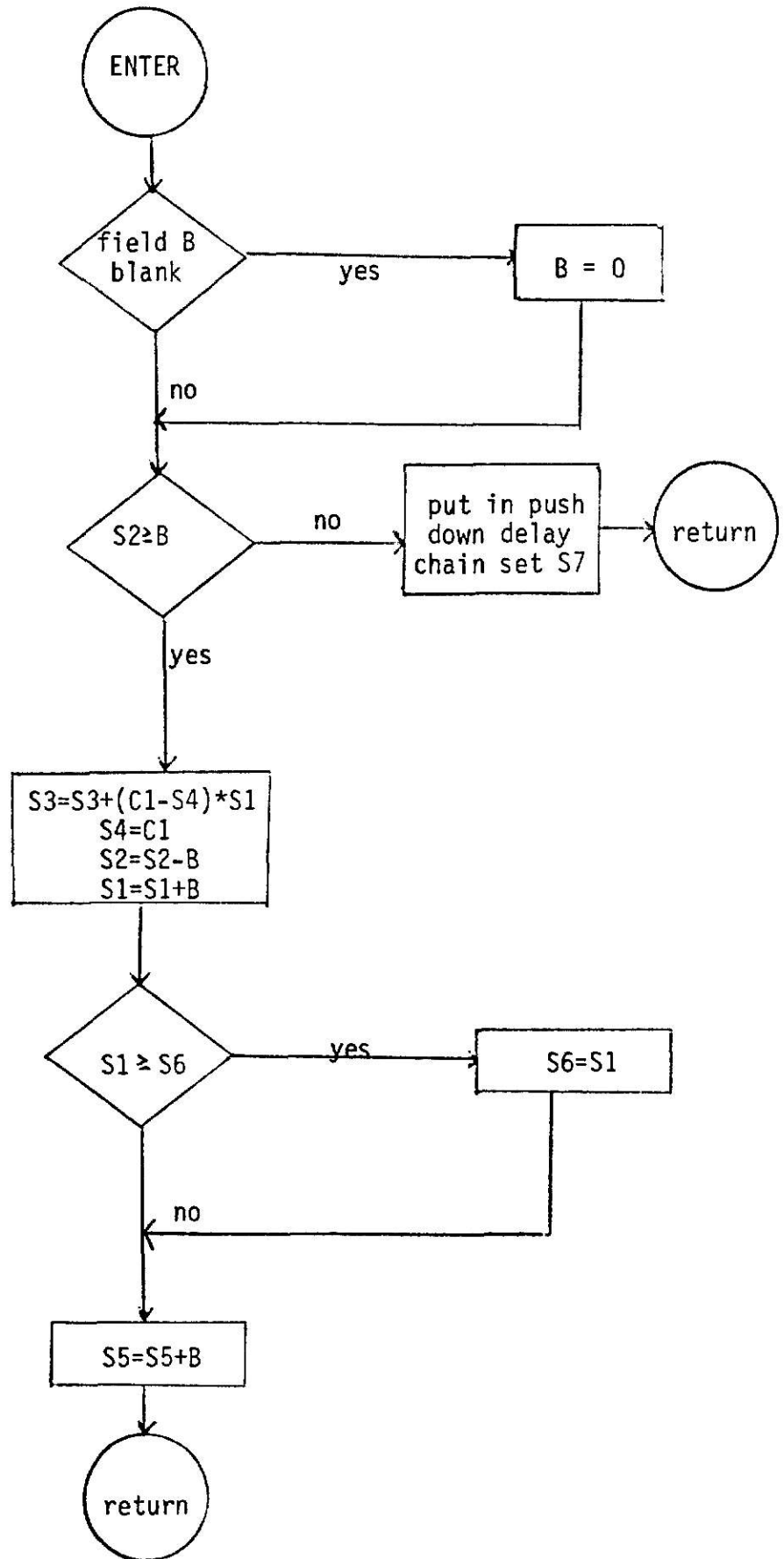
QUEUE	MAXIMUM CONTENTS	AVERAGE ENTRIES	TOTAL ENTRIES	ZERO ENTRIES	PERCENT ZEROS	AVERAGE TIME/TRAN	CURRENT CONTENTS
1	1	0.007	37	35	94.595	0.054	0
2	4	0.791	64	29	45.313	3.438	0

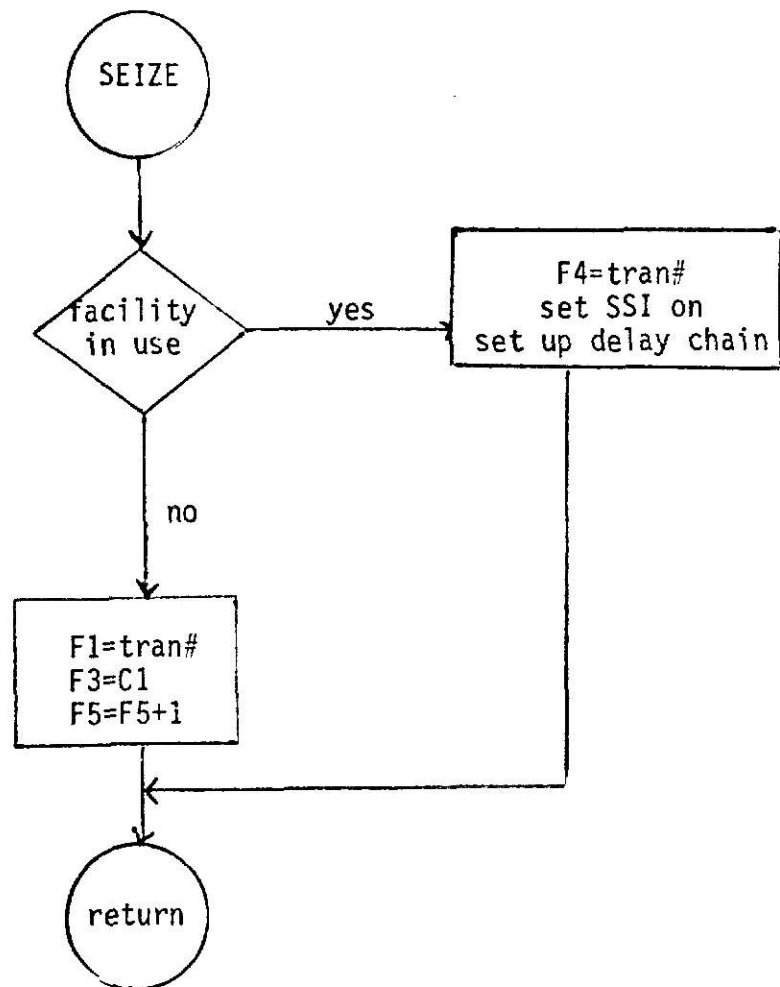
\$ AVERAGE TIME/TRAN EXCLUDES ZERO DELAY ENTRIES

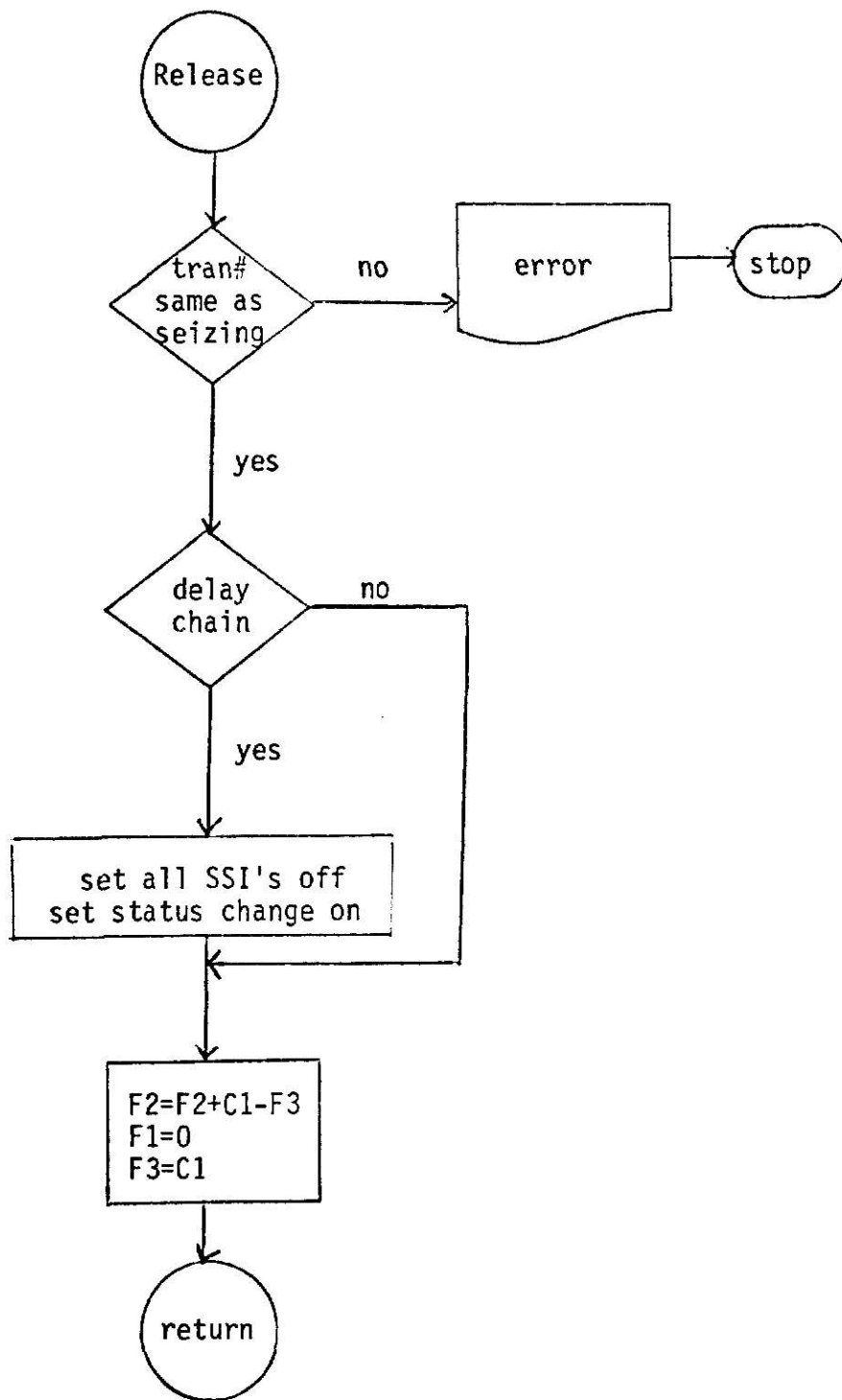
BLOCK TYPE FLOWCHARTS

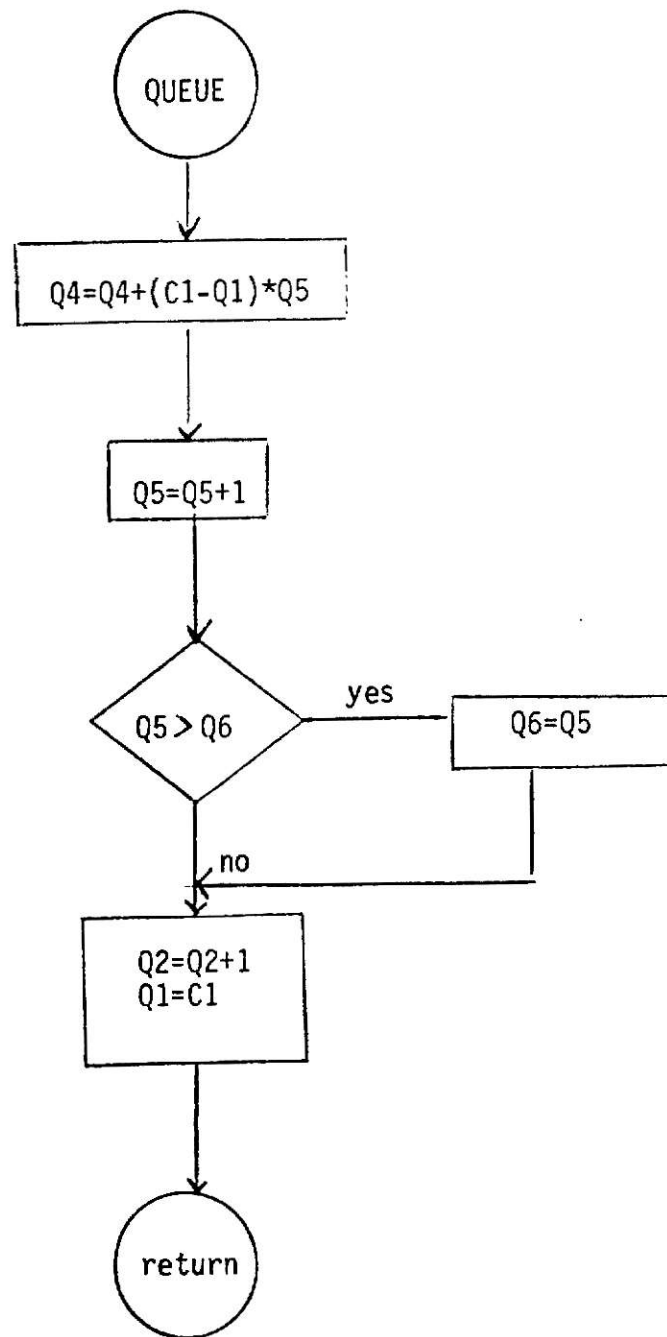


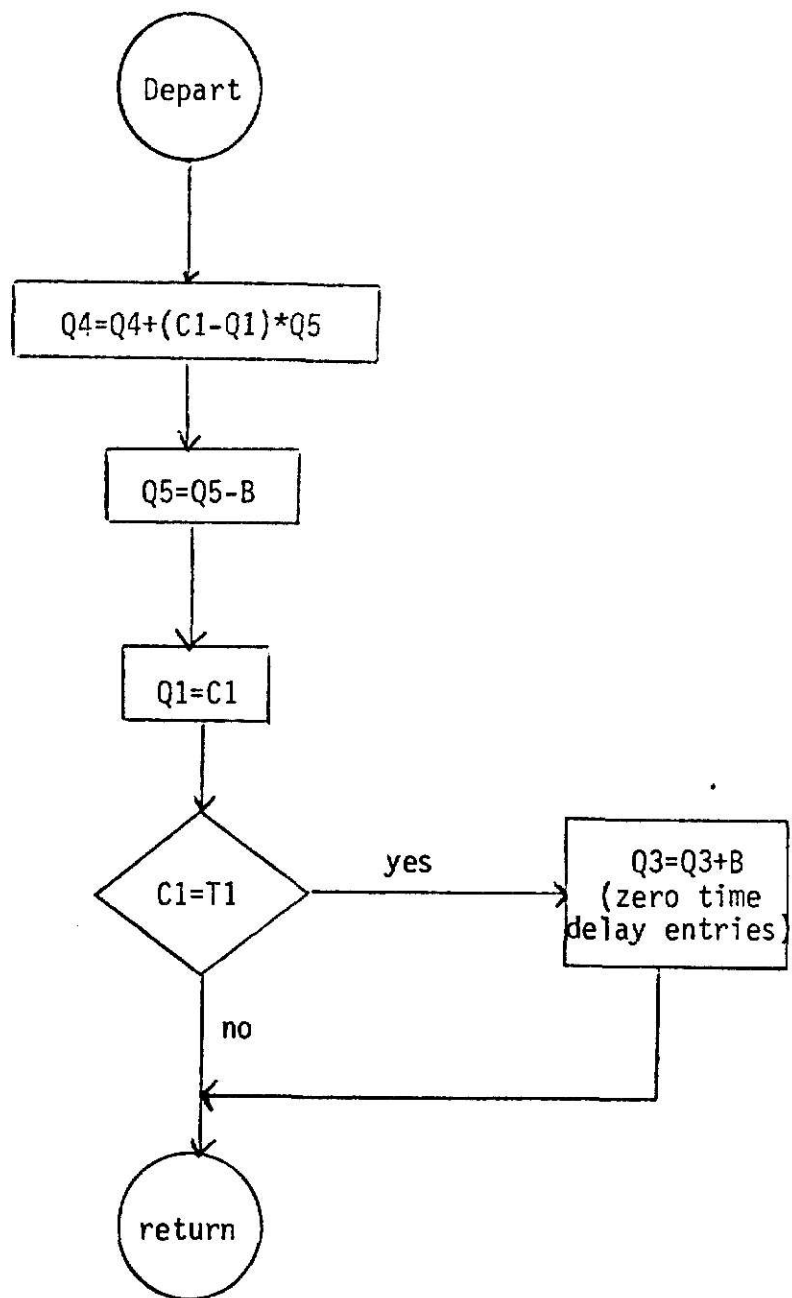


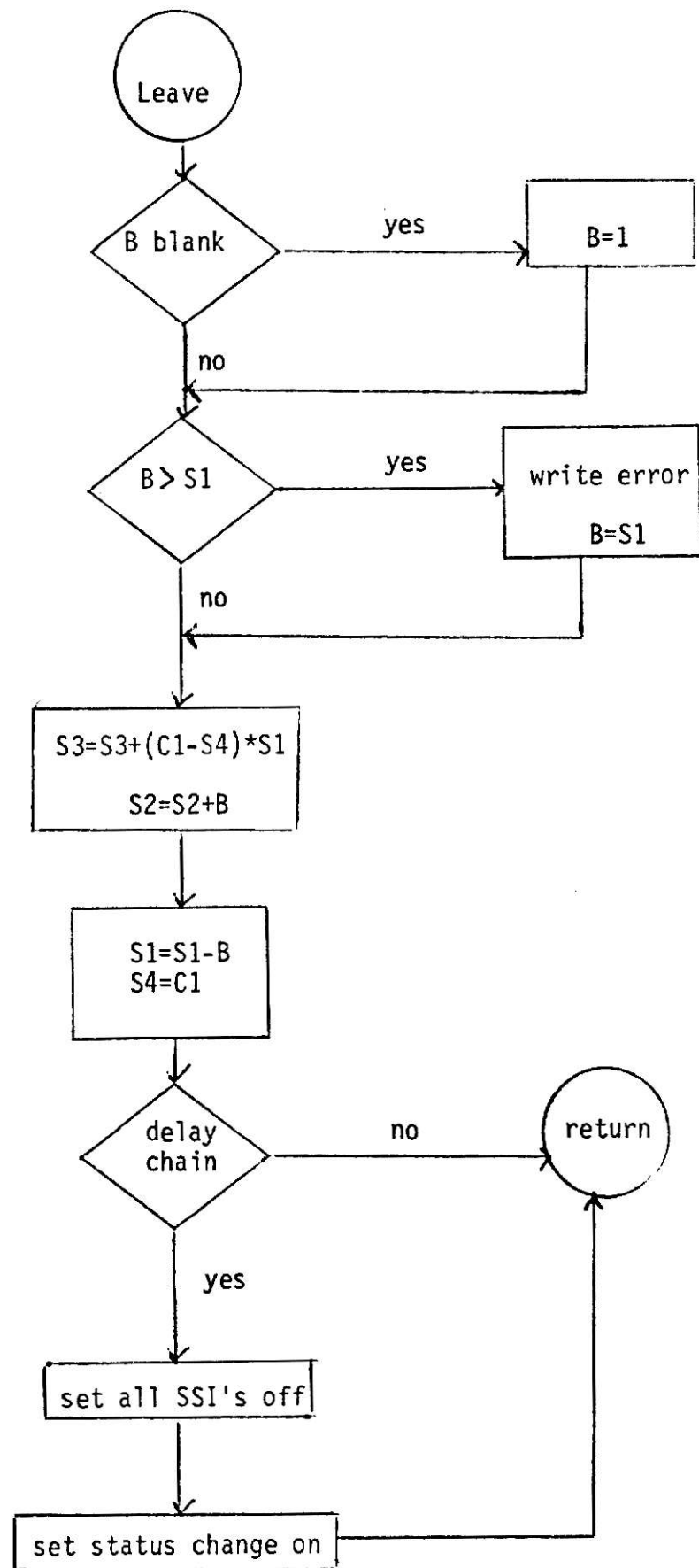


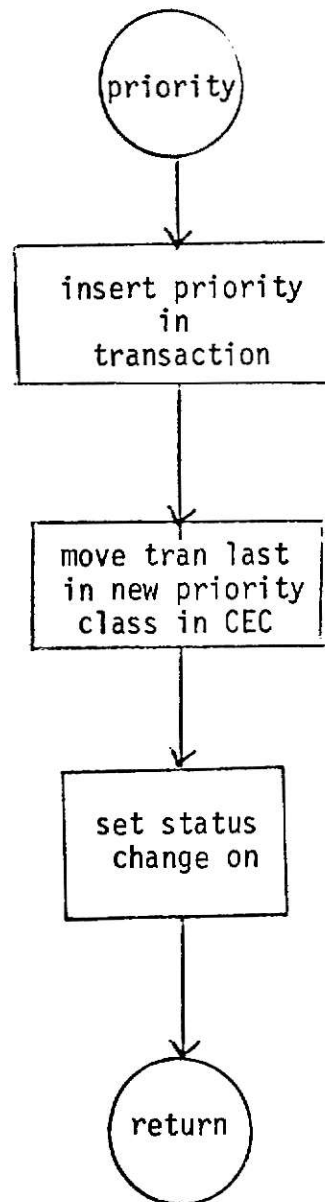


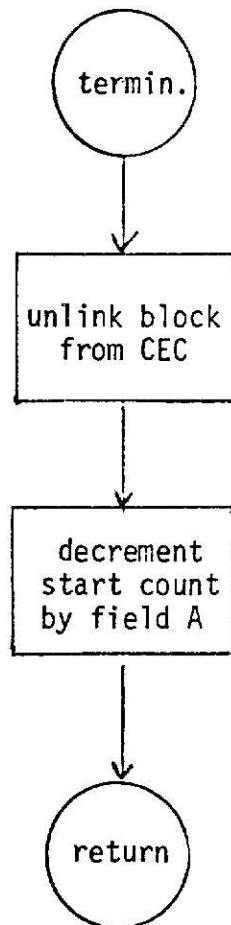


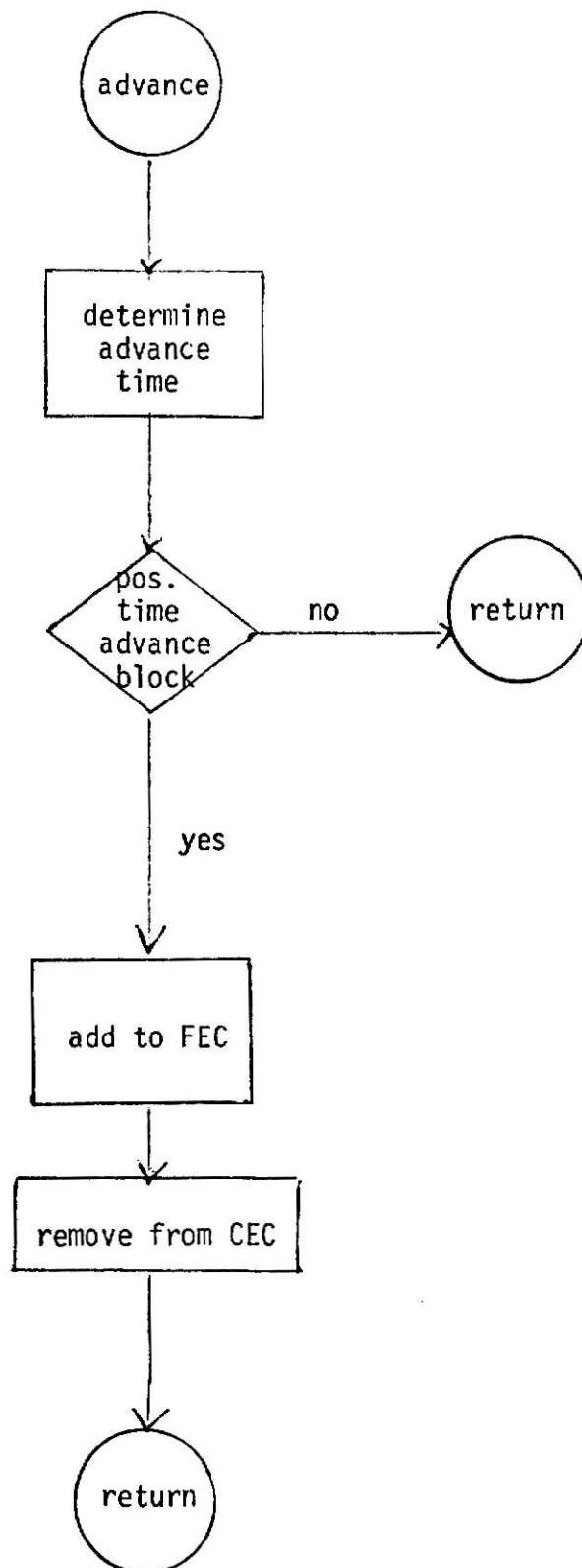


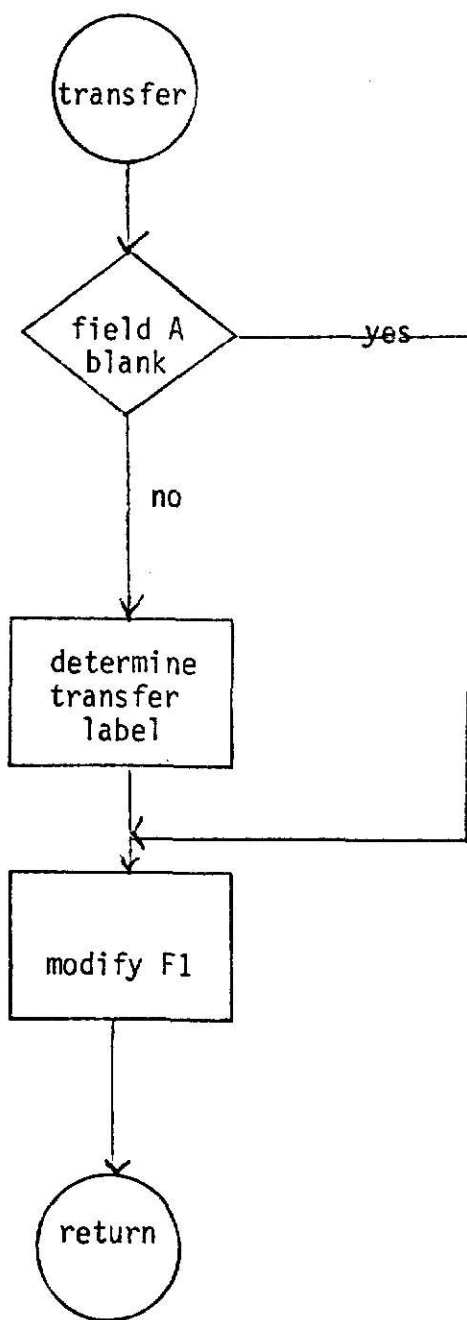


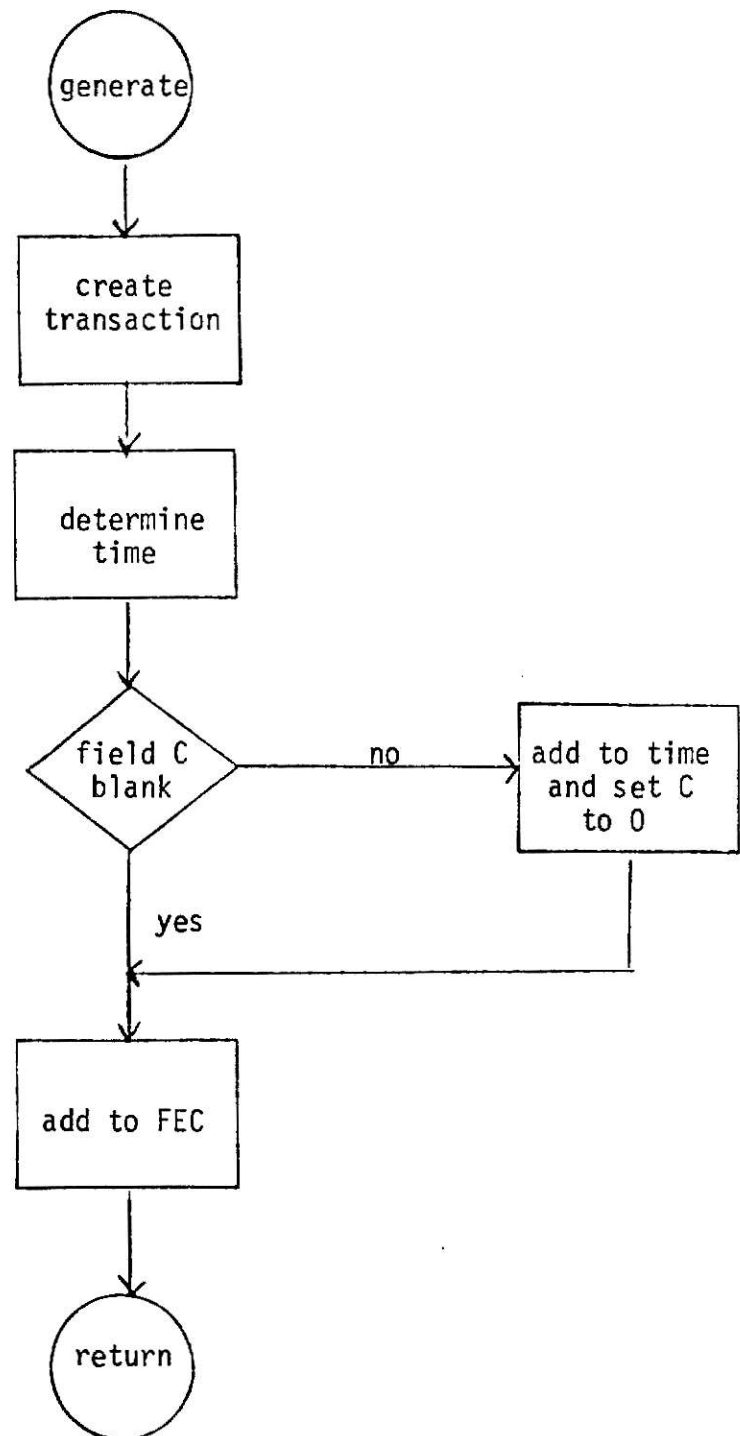












IGPSS

A General Purpose Simulation

Language for the Interdata

by

Martha Hoflich

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AN ABSTRACT OF A MASTER'S REPORT
submitted in partial fulfillment of the
requirements of the degree
MASTER OF SCIENCE

Department of Computer Science

KANSAS STATE UNIVERSITY
Manhattan, Kansas

1977

The purpose of this project is to provide a simulation language that is portable to different hardware configurations. It is designed primarily for use on the Interdata 8/32.

IGPSS is written in Fortran. It is based on GPSS and is used like GPSS with the following restrictions:

1. The input to IGPSS will be the anticipated output from the parser which remains to be written. The purpose of this project is to produce a program which will simulate a model given a set of strictly formatted input. A parser to be written as a future project will allow for more flexible input which more closely resembles the input required for GPSS.
2. The language statements available for use in IGPSS are a subset of those described in the GPSS/360 Introductory User's Guide (3).
3. Limits on number of parameters, facilities, etc. in IGPSS will be the same as the 64K version of GPSS.

The testing objective is to provide the same results using IGPSS as the same model would produce with GPSS. Documentation will consist of instructions for the use of IGPSS and any differences between it and GPSS.