

THE DESIGN OF A COMPUTER INTERFACE  
FOR A RADAR PROCESSOR

by

CARL C. ANDREASEN

B.S. Kansas State University, 1974

A MASTER'S THESIS

submitted in partial fulfillment of the  
requirements for the degree

MASTERS OF SCIENCE

Department of Electrical Engineering

KANSAS STATE UNIVERSITY

Manhattan, Kansas

1976

Approved by:

*Donald R. Hummels*

LD  
2668  
T4  
1976  
A54  
C.2  
Document

TABLE OF CONTENTS

178

	Page
TABLE OF CONTENTS . . . . .	i
LIST OF TABLES . . . . .	iii
LIST OF FIGURES . . . . .	iv
LIST OF PLATES . . . . .	v
Chapter	
1. INTRODUCTION . . . . .	1
BACKGROUND INFORMATION . . . . .	1
THE DIGITAL PROCESSOR COMPONENTS . . . . .	6
2. DESIGN OF THE COMPUTER INTERFACE . . . . .	11
THE NOVA CHARACTERISTICS . . . . .	15
Device Selection and Data Transfer . . . . .	18
The Interface Status Circuit . . . . .	23
AZIMUTH INCREMENT COUNTER . . . . .	30
THE SYNCHRONIZATION SIGNALS . . . . .	33
THE INTENSITY REGISTERS . . . . .	38
THE SELF TEST CIRCUITS . . . . .	40
3. THE CONTROL PROGRAM . . . . .	44
THE INITIALIZATION BLOCK . . . . .	44
GENERATE A FILE NAME . . . . .	49
THE DATA ACQUISITION SEQUENCE . . . . .	52
WRITING DATA TO THE DISK . . . . .	52
PREPARATION FOR THE NEXT DATA ACQUISITION SEQUENCE. . . . .	55
4. THE CONCLUSION . . . . .	61

	Page
BIBLIOGRAPHY . . . . .	64
APPENDIXES	
A. CIRCUIT DIAGRAMS . . . . .	65
B. COMPONENT PLACEMENT DIAGRAMS . . . . .	76
C. SOURCE LISTING OF RDCP . . . . .	85
ACKNOWLEDGEMENTS . . . . .	110

# **ILLEGIBLE DOCUMENT**

**THE FOLLOWING  
DOCUMENT(S) IS OF  
POOR LEGIBILITY IN  
THE ORIGINAL**

**THIS IS THE BEST  
COPY AVAILABLE**



## LIST OF TABLES

### Table

- 2.1 Correlation Between Azimuth Angle and  
Shaft Encoder Bit Position
- 3.1 Work File Parameter Table Organization
- 3.2 Work File Parameter Table Contents
- 3.3 Organization of Bytes for One Data Acquisition Sequence

## LIST OF FIGURES

### Figure

- 1.1 Weather Radar Data Gathering System Block Diagram
- 1.2 The Radar Transmitter Signal
- 1.3 Video Signal Time to Distance Relationship
- 1.4 Digital Processor Block Diagram
- 1.5 Typical Range-Azimuth Cell
- 2.1 Digital Computer Interface Block Diagram
- 2.2 Nova 1200 Architecture
- 2.3 Termination of Processor to Device Signals
- 2.4 Termination and Generation of Bidirectional Bus Lines
- 2.5 Device Code Decoding Circuit
- 2.6 Timing of an Input Instruction
- 2.7 Brute Force Data Output Gating
- 2.8 LSI Gating of Output Data
- 2.9 The Four Flipflop Control  
Network for Peripheral Interfaces
- 2.10 Busy-Done Flipflops Implimented
- 2.11 Decoding the Azimuth
- 2.12 Timing Diagram for Decoding the Azimuth
- 2.13 Generation of the Computer Interface to  
Integration and Quantization Circuit  
Synchronization Signals
- 2.14 Intensity Storage Registers
- 2.15 The Self Test Circuit Used to Simulate the Azimuth
- 2.16 The Self Test Circuit Used to Simulate the  
Integration and Quantization Circuit Output  
and Synchronization Signals
- 3.1 The High Level Flowchart of Radar Interface Control Program
- 3.2 The Initialization Block Flowchart
- 3.3 Flowchart for Determining Data File File Name
- 3.4 Timing of Program Segment to Move Data From The Interface
- 3.5 Structure of the Range Bin Intensity Words
- 3.6 Flowchart for Moving Data From Processor to Disk
- 3.7 Flowchart for Determining Next Sequence
- 4.1 Modifications to AZCCP

## LIST OF PLATES

### Plate

- I. The Computer Interface Circuit As Implemented

## Chapter 1

The Department of Electrical Engineering at Kansas State University is designing and building a system that will be used to obtain weather data from radar signals. Figure 1.1 shows a block diagram of the system. The system consists of a radar unit, a Digital Processor and the Electrical Engineering Department's Data General Nova 1200 minicomputer. The system components currently available are a surplus APS-81 radar unit, and a Nova 1200 minicomputer with the peripherals indicated in Figure 1.1. The purpose of the Digital Processor is to digitize the video signals coming from the radar unit and transfer the digital data to a mass storage device. A digital interface is required between the Nova and all of its peripherals as shown in Figure 1.1. This thesis covers a design and implementation of a Computer Interface circuit, for the Digital Processor, which controls the movement of the digitized video data to the computer's memory. Before proceeding with a description of the Computer Interface some background information on how the digitized video signal data is obtained is necessary.

### BACKGROUND INFORMATION

The principle of radar involves transmitting a pulse of energy of known magnitude towards an object, and then measuring the amount of the pulse that is reflected and the

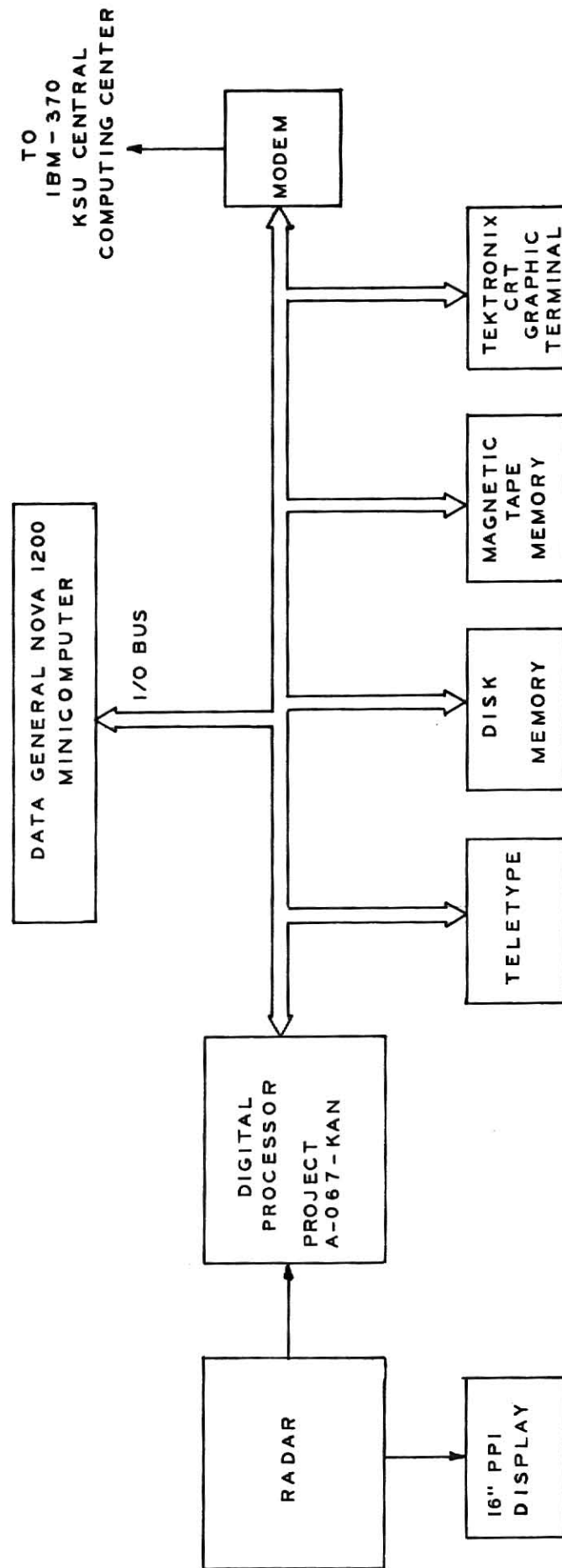


FIGURE 1.1  
WEATHER RADAR DATA GATHERING  
SYSTEM BLOCK DIAGRAM

time between transmission of the pulse and reception of a reflection, or echo. This time period is used to determine the distance between the transmitter and the object. The amount, or magnitude, of the echo is proportional to the reflectivity and size of the object. The transmitted pulse, as shown in Figure 1.2, is actually a high frequency signal which is amplified and transmitted for a short period of time.

When moisture droplets are of sufficient radius and quantity to fall as rain they will reflect the transmitted signal producing a detectable echo. In this manner an observer will "see" on the display device, a Pulse Position Indicator scope, the echoes received from falling rain. The radar antenna has a beamwidth of 5 degrees horizontally and 10 degrees vertically and is rotated clockwise 360 degrees every four seconds. With 800 high frequency pulses transmitted per second the APS-81 will receive unambiguous echos of objects that are less than 100 miles away from the antenna.

The reflected pulses are detected in the radar receiver to produce a video signal for each of the transmitted pulses. The magnitude, or voltage, at any point along the video signal is proportional to the magnitude of the echo from the associated transmitted pulse. There is a direct time to distance relationship in the video signal as shown in Figure 1.3

The antenna starts at true North and revolves 360

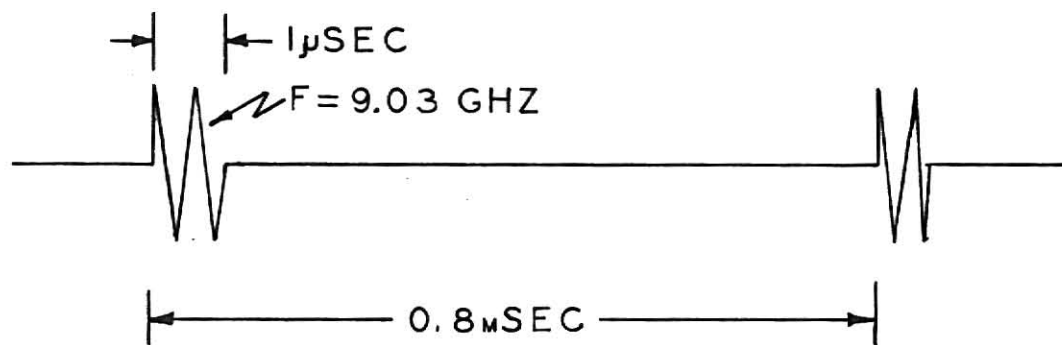


FIGURE 1.2  
THE RADAR TRANSMITTER SIGNAL

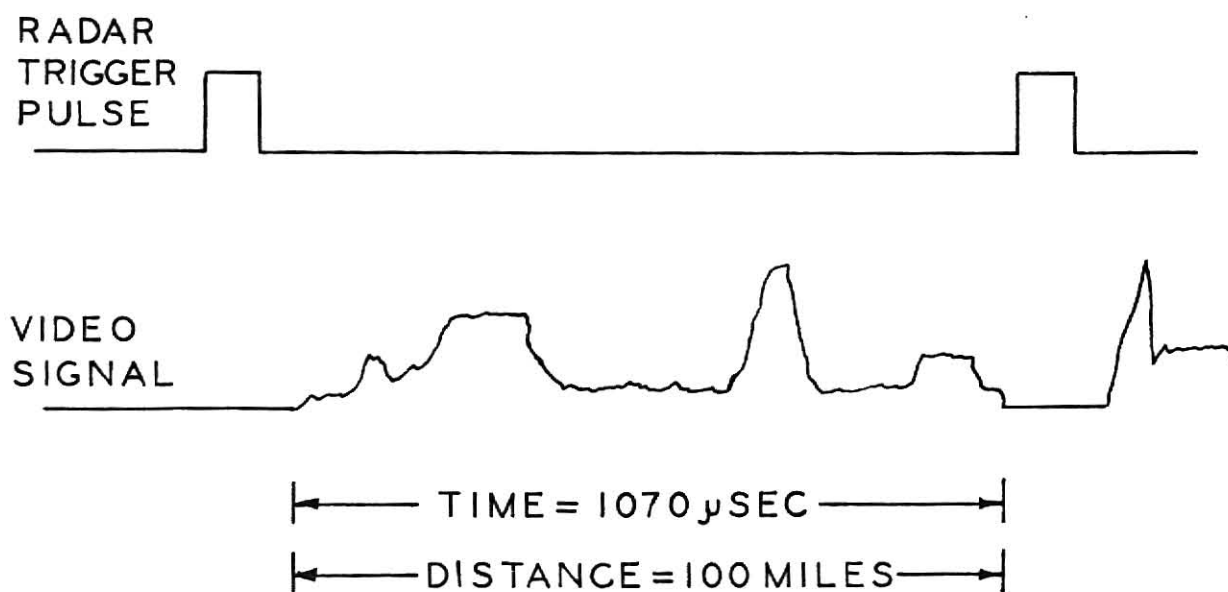


FIGURE 1.3  
VIDEO SIGNAL TIME TO  
DISTANCE RELATIONSHIP



degrees. The azimuth is defined as the angular distance from true North that the antenna is facing. The azimuth is determined by using a shaft encoder. The shaft encoder has a binary output that is equal to the angle that the encoder's shaft is with respect to a fixed point. The shaft encoder output is decoded in incremental amounts and then the video signals that occur in this incremental amount are integrated. The integral of the magnitude of the video signal over a specified interval and incremental azimuth is the echo intensity for a region, and is called a range bin integral. A range bin is defined geometrically as a sector of an annulus whose side is one-half mile and angle is the incremental azimuth value. Figure 1.4 shows the division of the radar pattern into the range bins or range azimuth cells. With the above background information a brief description of the operation of the components of the Digital Processor block can now be presented.

#### THE DIGITAL PROCESSOR COMPONENTS

The Digital Processor block is composed of three circuits as shown in Figure 1.5. The Logarithmic Video Amplifier is used to compress the large dynamic range video signal into a video signal with a ten step dynamic range. The Integration and Quantization Circuit will integrate the video signal for each of 200 separate intervals, each interval corresponding to a one-half mile segment of range.

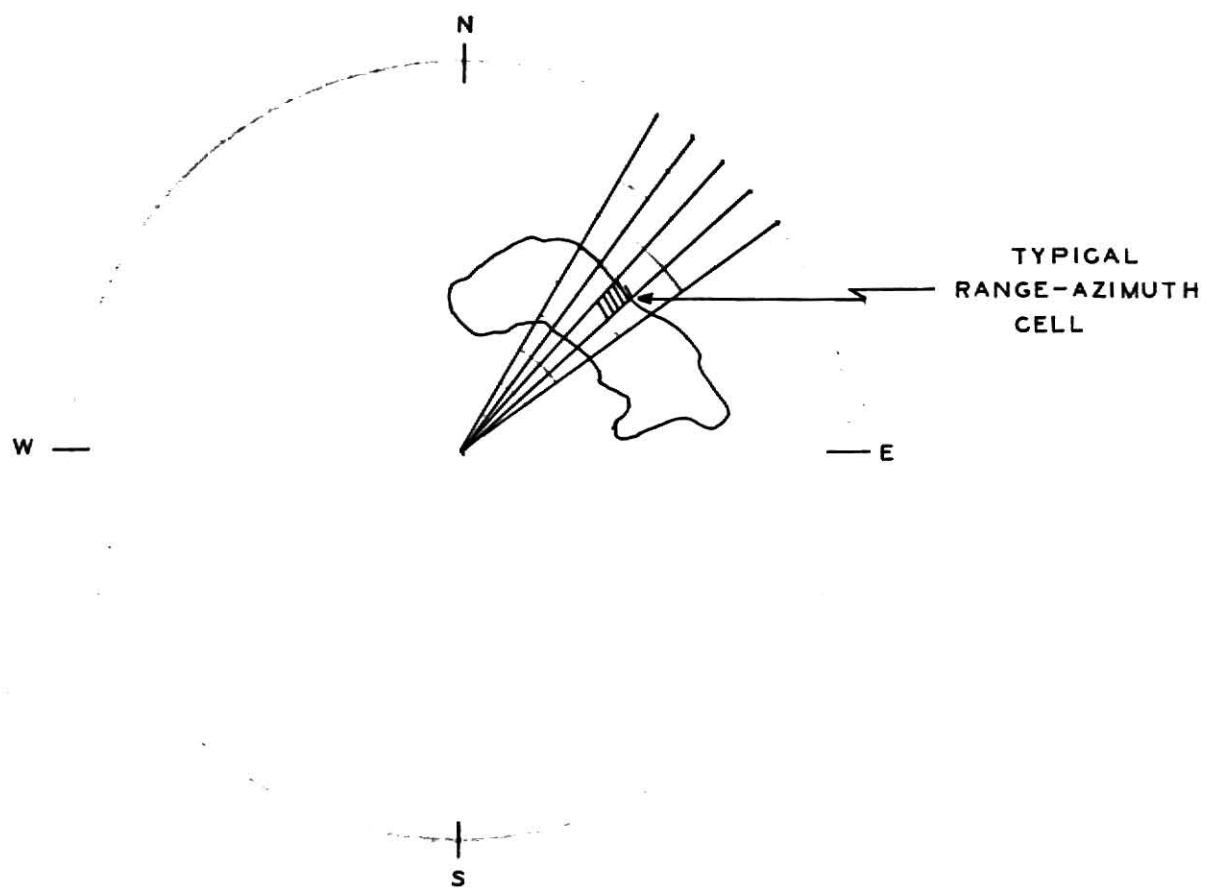


FIGURE 1.5  
TYPICAL RANGE-AZIMUTH CELL

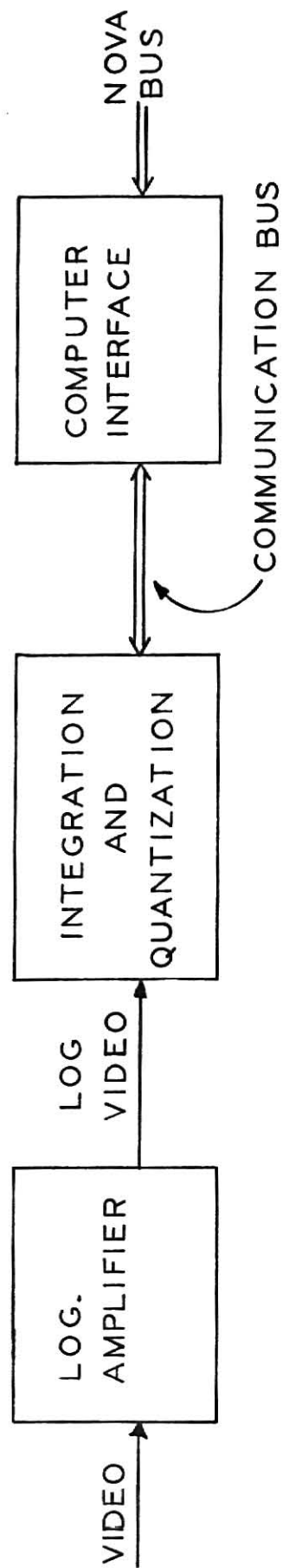


FIGURE 1.4  
DIGITAL PROCESSOR BLOCK DIAGRAM

Thus the tenth interval is that part of the video signal corresponding to 4.5 to 5.0 miles away from the radar antenna. The time of integration is the time that a point object would be within the antenna beam, during one scan. The Quantization part of the Integration and Quantization circuit converts the voltage on a capacitor, which represents the range bin integral, to a Binary Coded Decimal (BCD) value, the range bin intensity. The Integration and Quantization circuit generated the range bin intensities and now the range bin intensities are passed on to the Computer Interface circuit for eventual transfer to a mass storage device.

Each one-half mile segment of the video signal corresponds to 5.37 microseconds. After the period of integration each range bin intensity will be present on the output of the Integration and Quantization circuit for only 5.37 microseconds. This rate is faster than most mass storage devices can handle. The solution is to collect the data in a computer's memory and then move the data to a mass storage device in whatever format is needed. The sequence of events necessary to obtain and store in the computer's memory all of the range bin intensities for one revolution of the antenna is termed the "data acquisition sequence". The computer used is a Data General Nova 1200 minicomputer with 28K words of memory.

The Computer Interface circuit is needed to control the incremental azimuth over which integration of the video

signals is performed and control the movement of the range bin intensities from the Integration and Quantization Circuit to the computer's memory. This thesis covers the design and implementation of the Computer Interface circuit which accomplishes these tasks.

## Chapter 2

This chapter presents a detailed explanation of the internal operation of the Computer Interface Circuit. Figure 2.1 shows the block diagram of the Computer Interface Circuit, hereafter referred to as the interface or interface circuit. Plate 1 shows the Computer Interface circuit as implemented. The interface is necessary to synchronize the movement of the range bin intensities from the Integration and Quantization circuit to the Nova minicomputer. The interface also decodes the azimuth and provides signals which control the azimuth interval over which the video signals are integrated.

In order to describe the internal operation of the interface a description of the signals provided and required by the interface is required. The first section of this chapter presents a description of the signals required and provided by the Nova minicomputer. The next section presents a description of how the azimuth is decoded, the signals generated and their use. The next two sections provide a description of the signals used to communicate with the Integration and Quantization circuit and a description of the operation of the registers used to temporarily store the range bin intensities respectively. The final section of this chapter provides a description of the circuit which simulates the azimuth and the signals coming from the Integration and Quantization circuit.

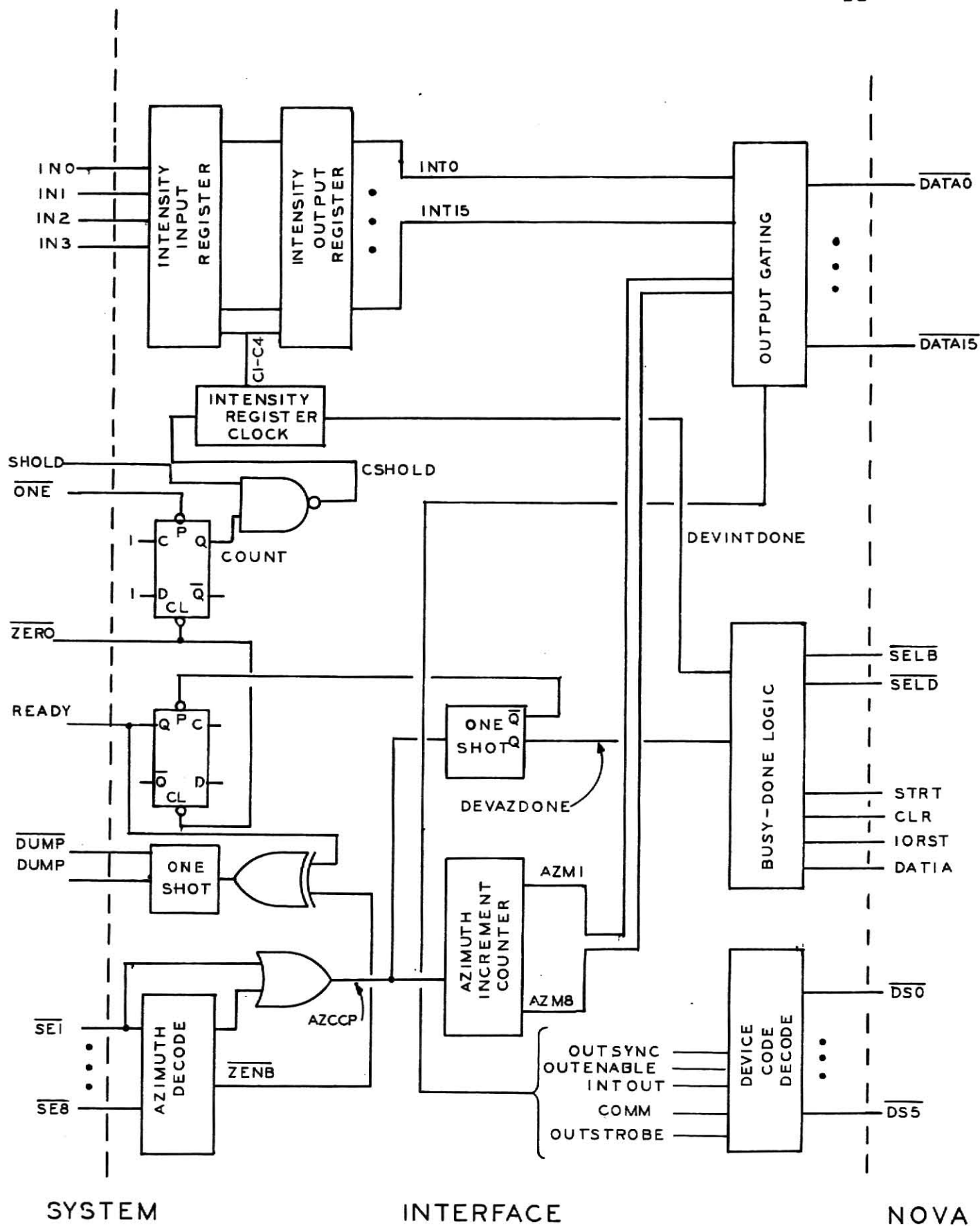
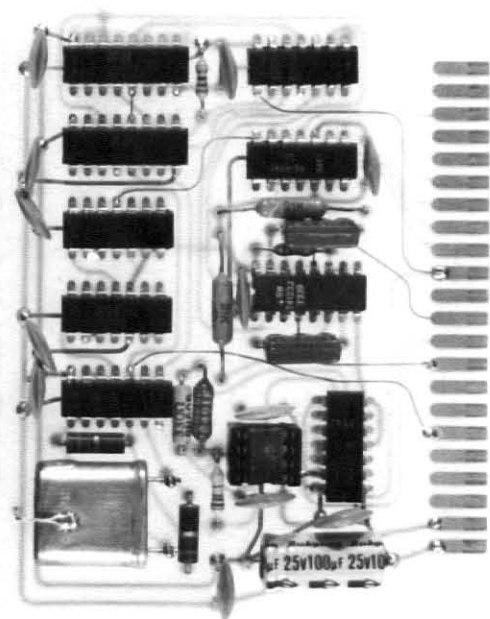
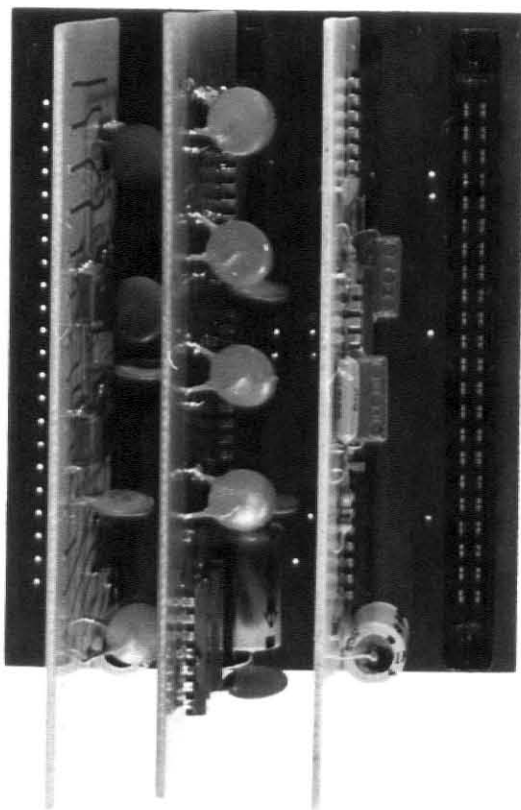
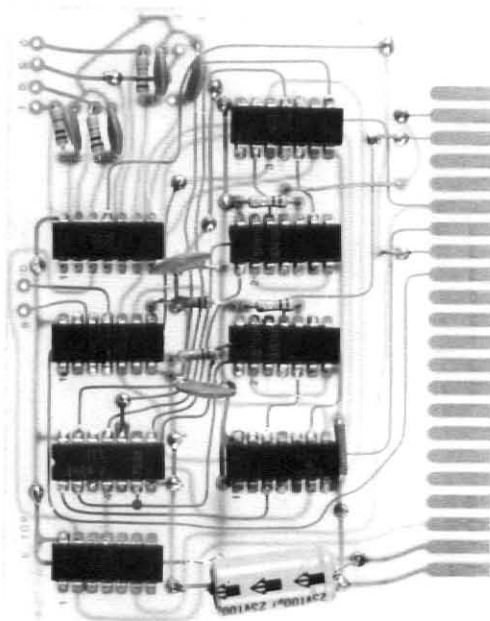
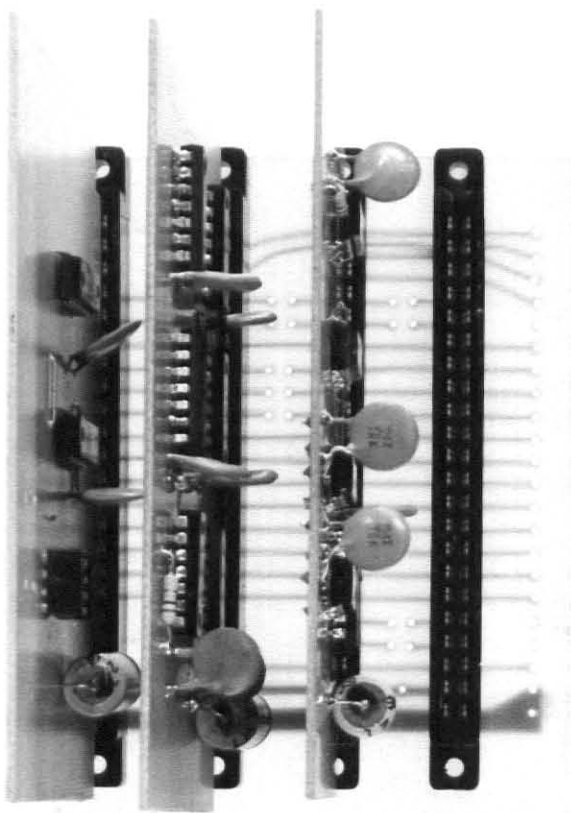


FIGURE 2.1  
COMPUTER INTERFACE BLOCK  
DIAGRAM

PLATE I

THE COMPUTER INTERFACE CIRCUIT AS IMPLEMENTED





## THE NOVA CHARACTERISTICS

One method for a computer to communicate with its peripheral devices is the use of a group of conductors, called a bus, which runs to every device. The bus carries the necessary electrical signals to allow communication between the computer and peripheral devices. The Data General Nova 1200 (hereafter is referred to as "computer") bus carries to every peripheral:

1. Six device selection lines,  $\overline{DS0}-\overline{DS5}$ , from computer to the devices.
2. Nineteen Control lines from the computer to the devices to synchronize data movement to and from the bus. For example the IORST signal that resets all peripherals.
3. Six control lines from the devices to the computer, such as the Interrupt Request line.
4. Sixteen bidirectional data lines,  $\overline{DATA0}-\overline{DATA15}$ .

Since this bus is shared by all peripherals a separate interface is provided for each peripheral device as illustrated in Figure 2.2. As shown some interfaces can be included within the computer mainframe. For additional peripheral devices the bus can be extended, as in the case of the Computer Interface circuit for the Radar Processor.

When using the external bus the control signals that originate in the computer and drive the peripheral interfaces must be terminated as shown in Figure 2.3. The

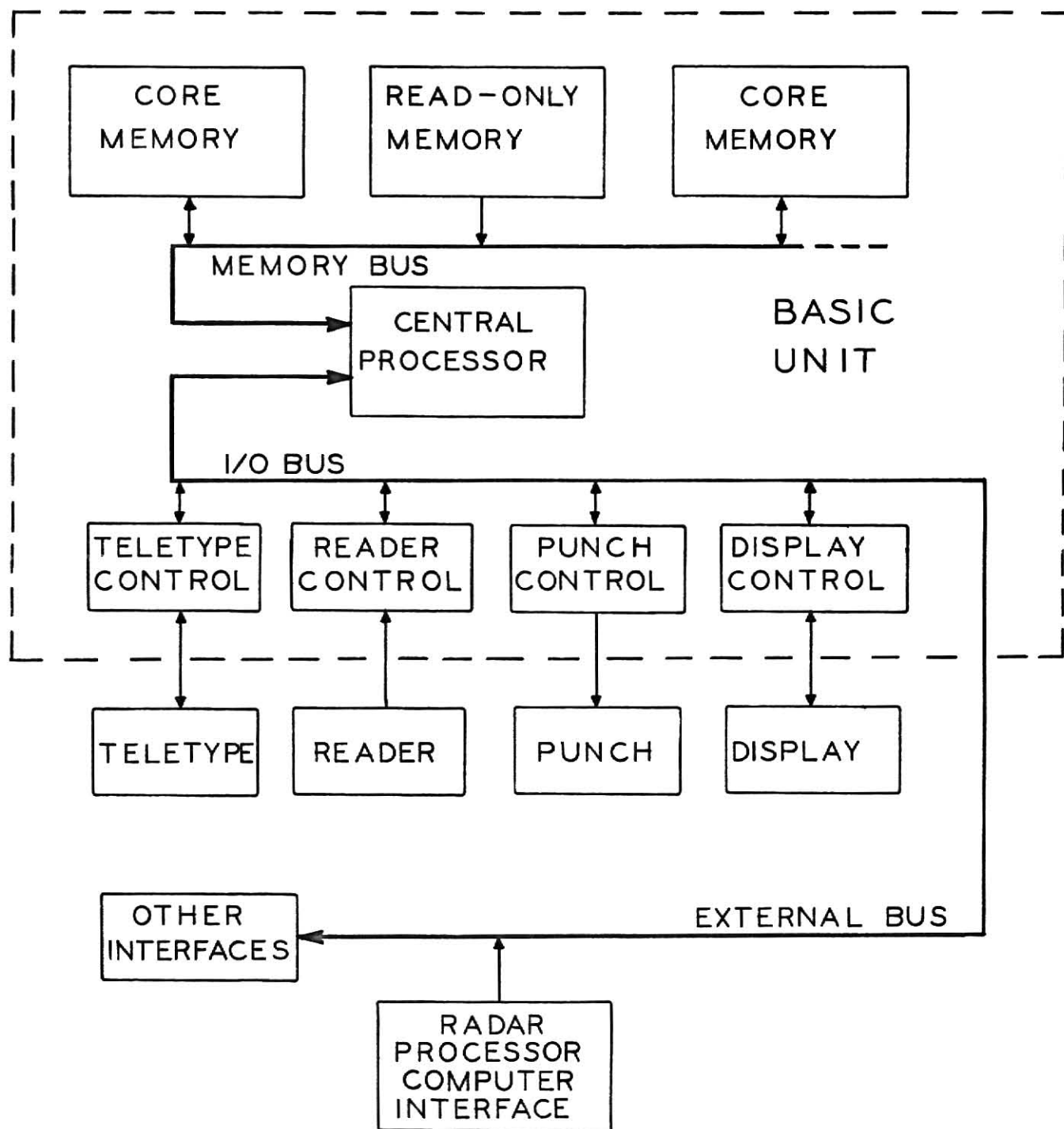


FIGURE 2.2  
NOVA 1200 ARCHITECTURE (3)

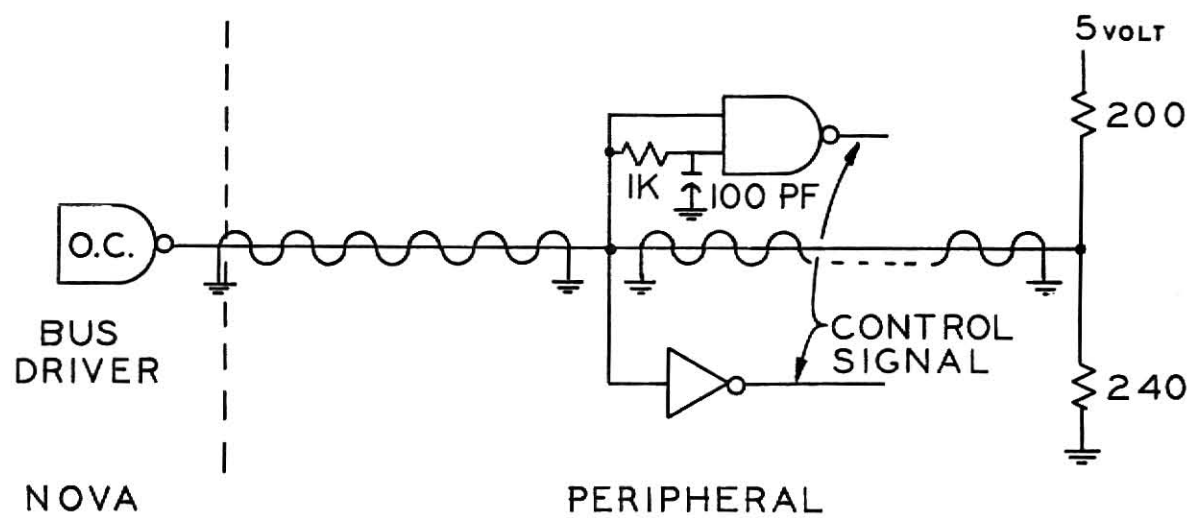


FIGURE 2.3  
TERMINATION OF PROCESSOR  
TO DEVICE SIGNALS

signal receivers are often an inverter or a nand gate with a filter on the input lines. The voltage divider at the end of each bus line gives the correct termination.

The bidirectional signals, such as the data lines, and signals which come from the peripheral interfaces to the computer, such as the  $\overline{\text{SELB}}$  line, must be terminated as shown in Figure 2.4. The bidirectional lines are "active low" in that pulling a line low is equivalent to placing a logical 1 signal on that line. Thus the output transistor of an open collector gate, such as the SN7438, is used for the line transmitter and an inverter is used for a line receiver. Note that all bus lines must be twisted pairs.

#### Device Selection and Data Transfer

The computer provides a six bit device selection code on the device selection lines  $\overline{\text{DS0}}-\overline{\text{DS5}}$ . These lines are active low and are decoded so that each interface has an enable line that is high when the interface's device code is on the device select lines. The complement of the enable line is used to synchronize the transmission or reception of data when one of the Data In (DATIA, DATIB, DATIC) or Data Out (DATOA, DATOB, DATOC) lines is active respectively. This interface transfers an azimuth value and the range bin intensity word so to distinguish between them, two device select codes are used. One device selection code (63) is used to synchronize the transfer of the output of the

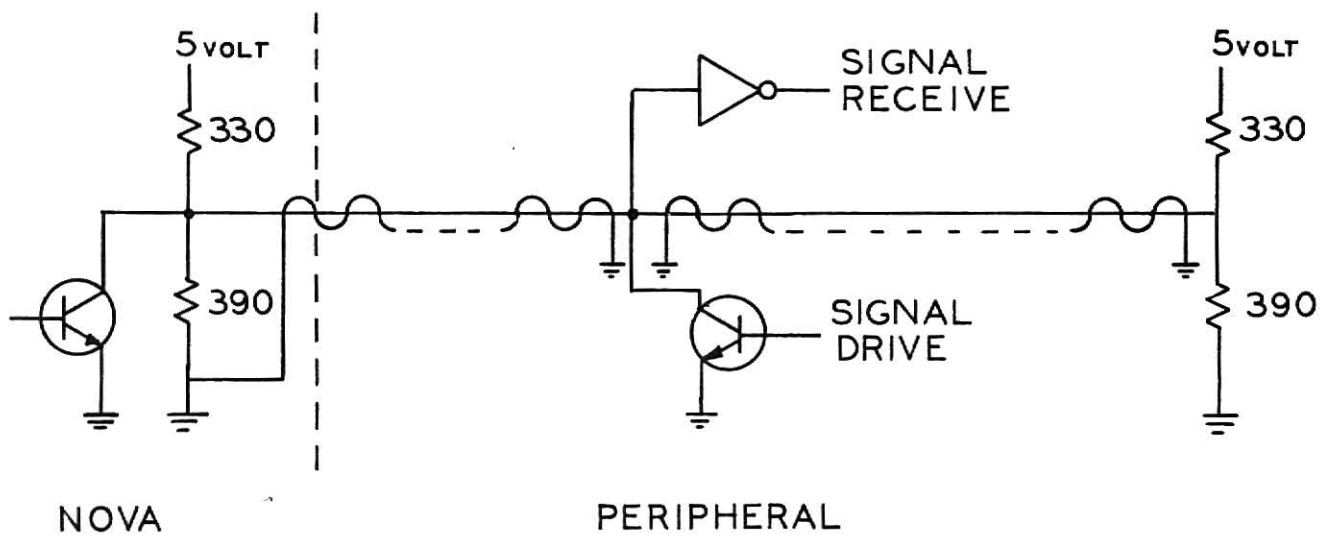


FIGURE 2.4  
TERMINATION AND GENERATION OF  
BIDIRECTIONAL BUS LINES

Azimuth Increment Counter to the data bus, and the other device selection code (67) is used to transfer the output of the Range Bin Intensity Output Register to the data bus.

Figure 2.5 shows the logic necessary to generate the enable lines DEVINT and DEVAZ required by this interface.

Figure 2.6 gives the timing diagram for a typical input instruction. This instruction is of the form:

Mnemonic(code) AC, Device Code

where the Mnemonics are:

DIA Read data from the A buffer

DIB Read data from the B buffer

DIC Read data from the C buffer

Note from the above that with each peripheral there may be up to three buffers from which data may be obtained. The optional code that can be concatenated to the mnemonic will be explained in further detail in the next section of this chapter. AC stands for one of the four accumulators (0,1,2,3), and Device Code is the two digit octal device code placed on the device selection lines.

Execution of the DIA instruction causes the device code to be placed on the device selection lines and the DATIA line to go high. When this condition is decoded by a device interface an enable line is generated which gates the data from the device interface output buffer to the data bus. Figure 2.5 shows the data output enable lines INTOUT and AZOUT that are required to gate the Azimuth Increment Counter or the Range Bin Intensity Output Register to the

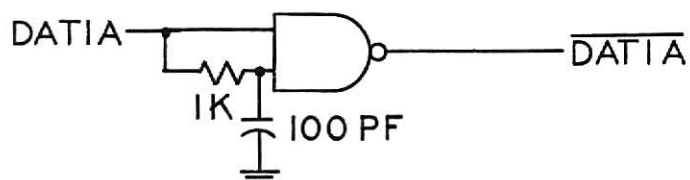
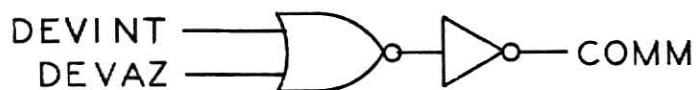
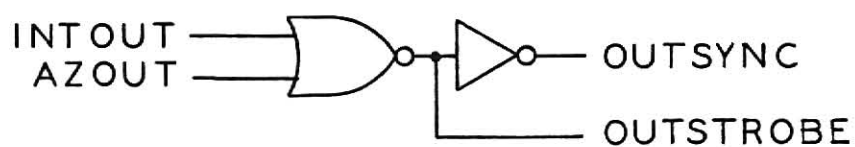
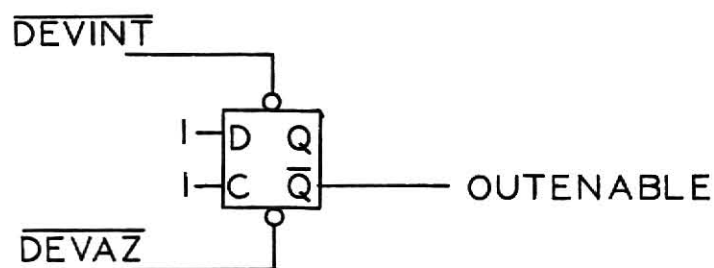
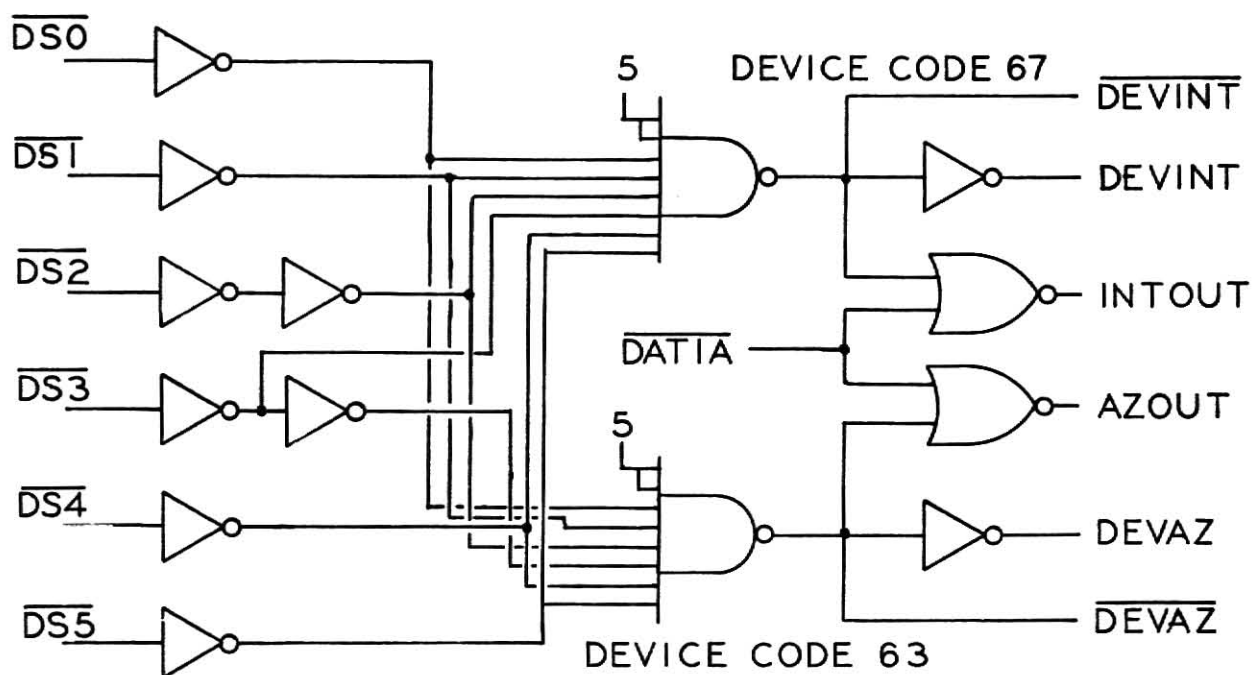
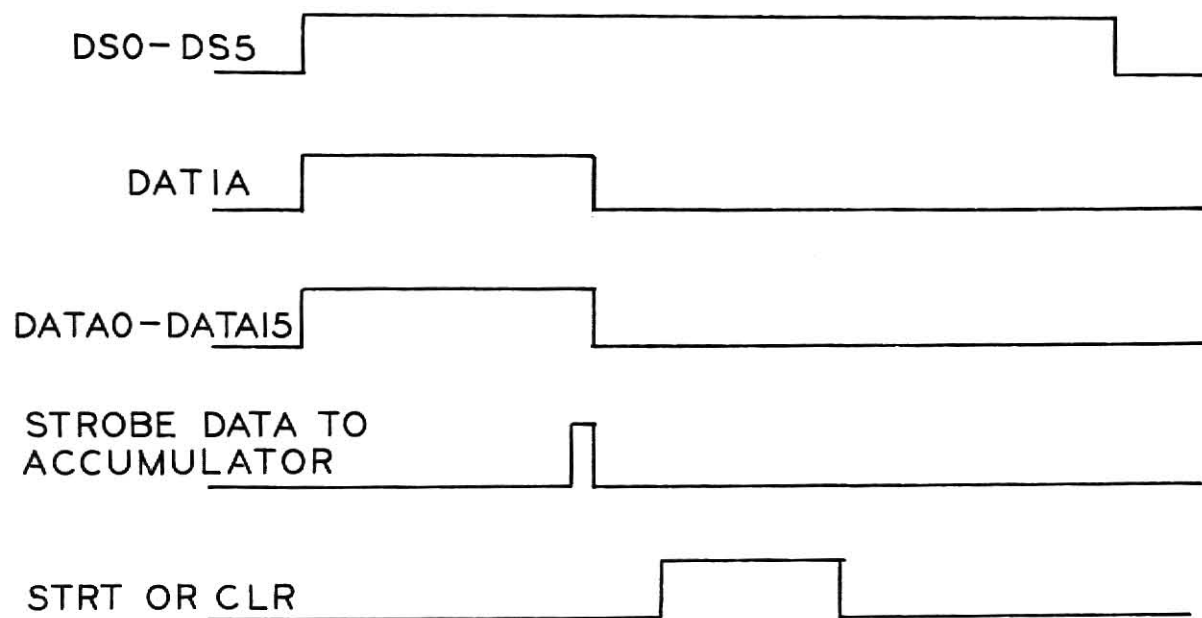


FIGURE 2.5  
DEVICE CODE DECODING CIRCUIT





**FIGURE 2.6**  
TIMING OF AN INPUT INSTRUCTION

bus.

Figure 2.7 presents a brute force method of gating the data from the Range Bin Intensity Output Register and the Azimuth Increment Counter to the data bus. A better method is shown in Figure 2.8. This method requires fewer packages and reduces the number of gate inputs that must be driven by the INTOUT signal.

### The Interface Status Circuit

Each peripheral device interface has a four flipflop control network that specifies the state of the device and requests an interrupt by that device (Busy, Done, Interrupt Disable and Interrupt Request). Figure 2.9 shows the flipflop control network suggested. When an input device has data ready to be sent to the computer, the device sets its Done flipflop. The computer will clear the Busy and clear the Done flipflops of an output device when data is placed in the buffer for that device. When the output device is ready for more data it will set its Done flipflop.

The computer can determine when a device is ready for service by two methods. The first method involves periodically checking the Selected Done ( $\overline{\text{SELD}}$ ) line which is brought low when the device Done flipflop is set and its device select code is on the device select lines. The execution of a skip-on-done instruction (SKPD!) in the program initiates this check. The second method is for the

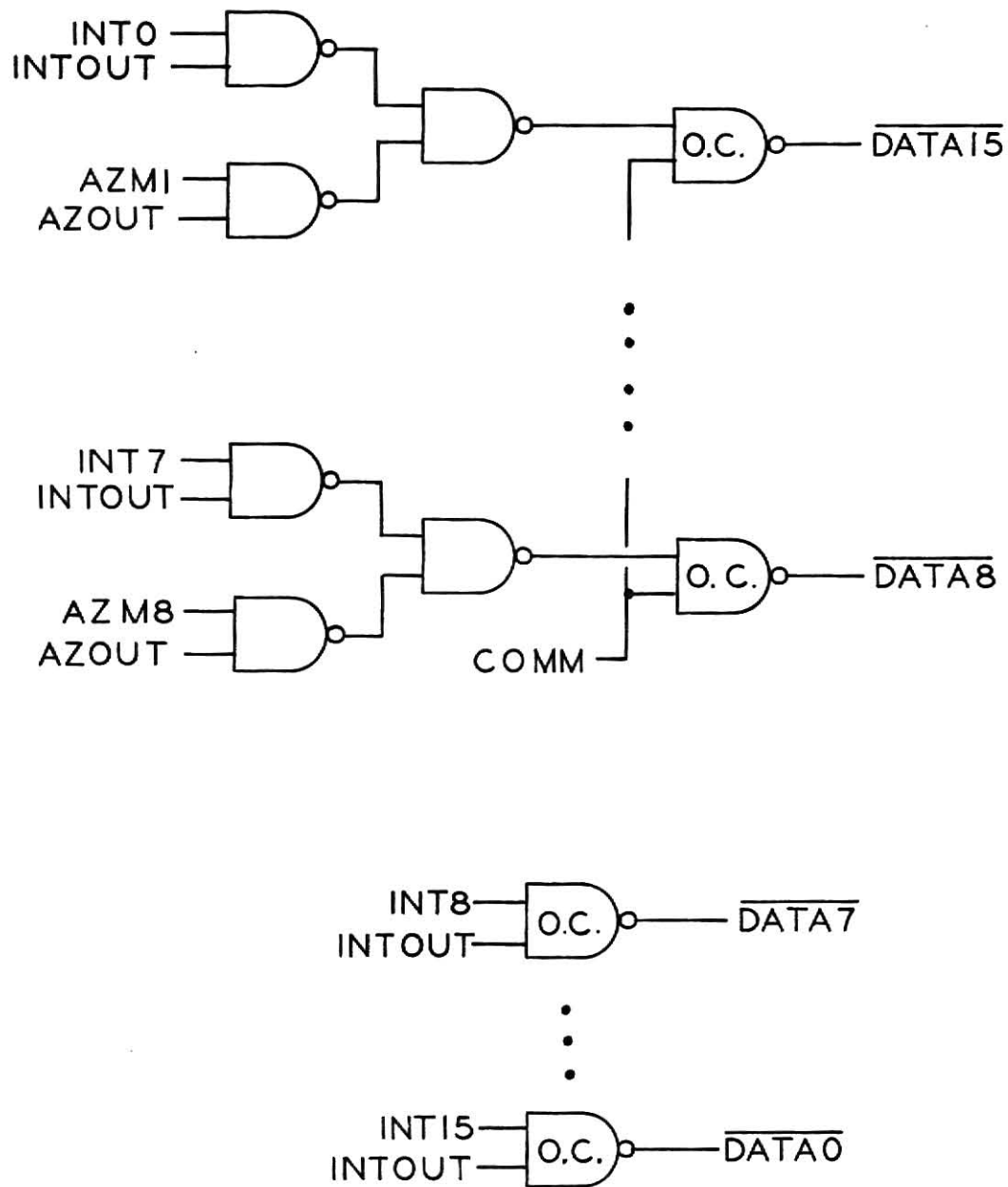


FIGURE 2.7  
BRUTE FORCE DATA OUTPUT GATING

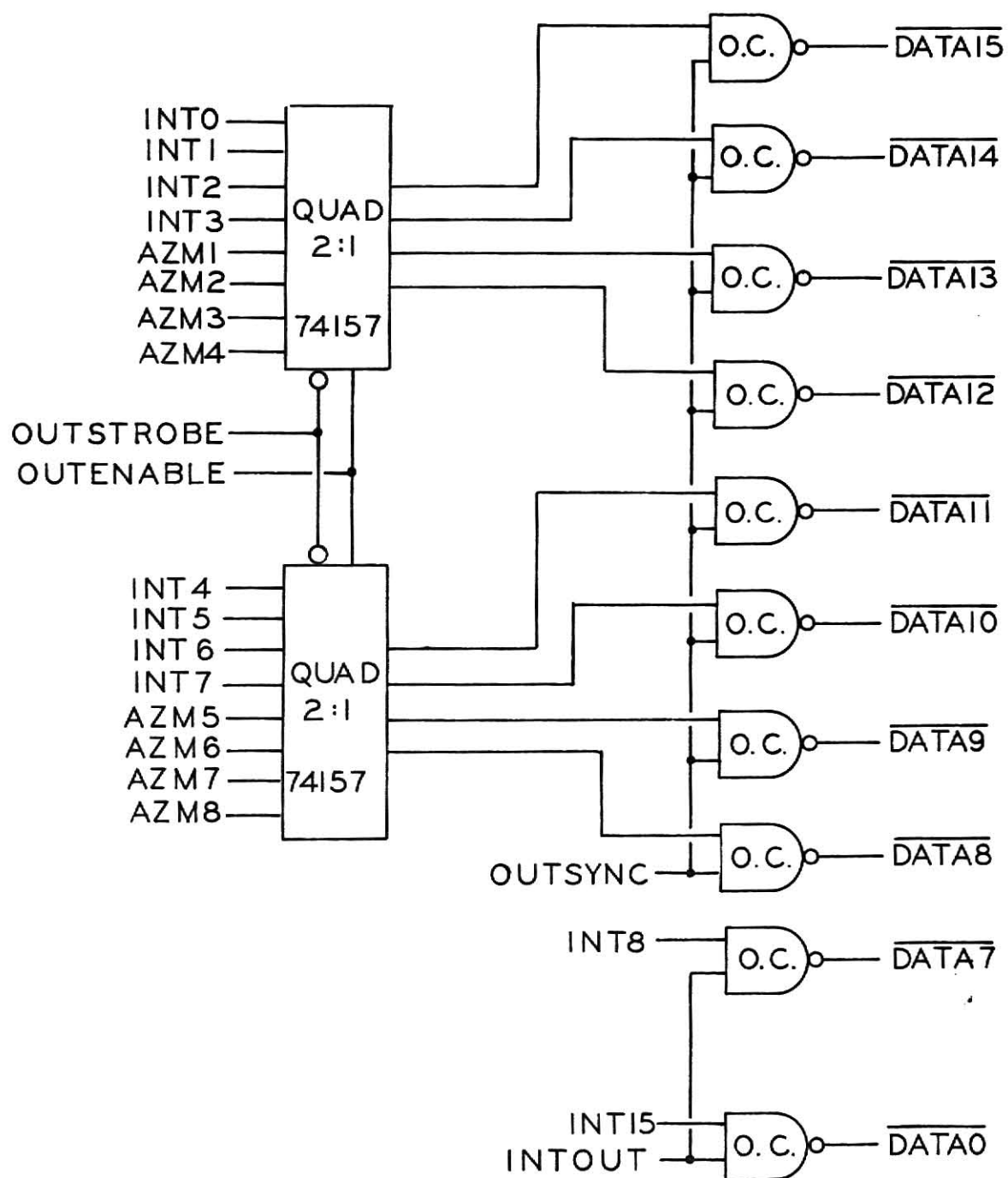


FIGURE 2.8  
LSI GATING OF OUTPUT DATA

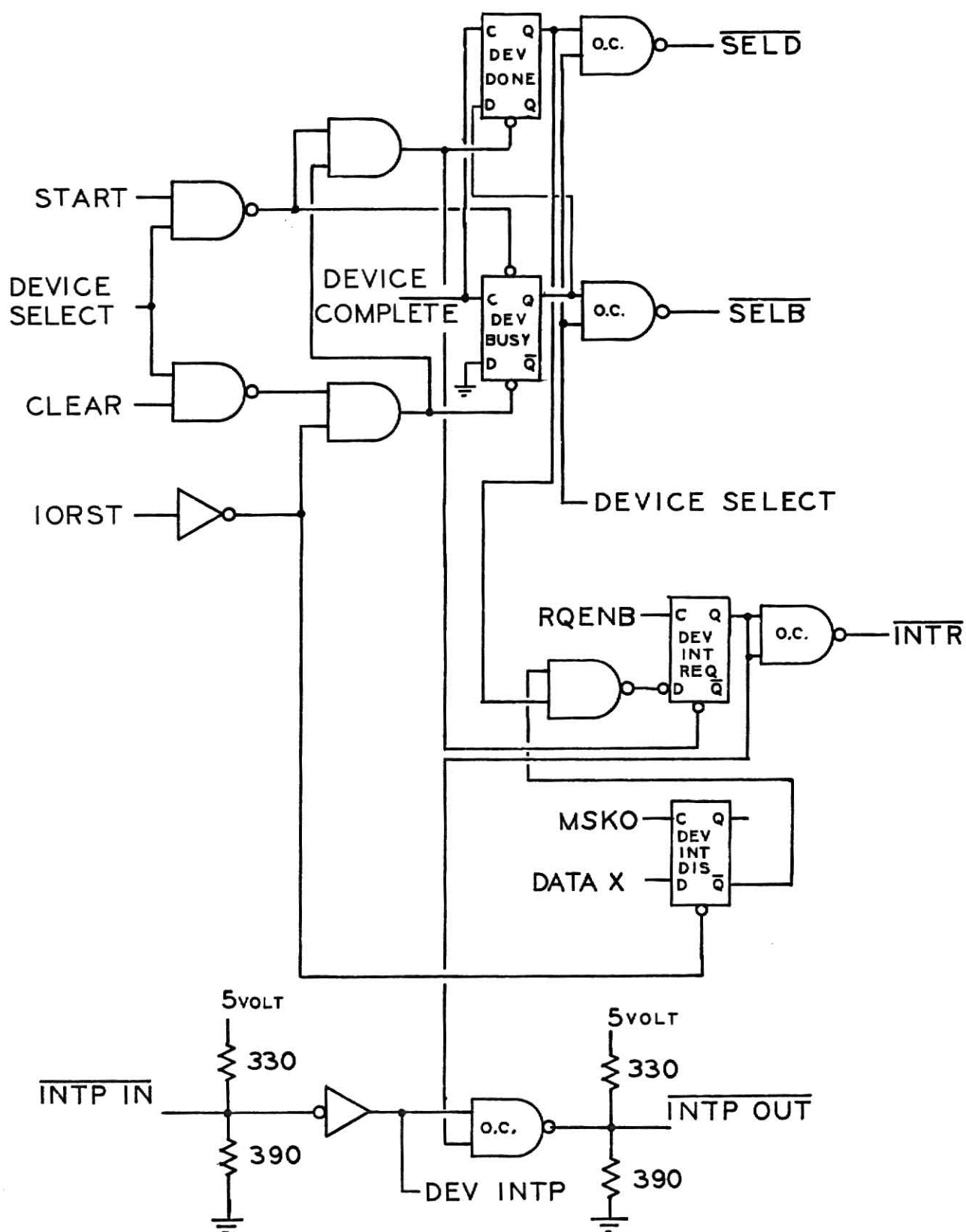


FIGURE 2.9  
THE FOUR FLIPFLOP CONTROL NETWORK  
FOR PERIPHERAL INTERFACES (3)

device interface to interrupt the computer. During each computer cycle, at a time when the computer is capable of being interrupted, the Request Enable ( $\overline{RQENB}$ ) line is pulsed. One or more devices needing service will then pull the Interrupt Request ( $\overline{INTR}$ ) line low if their associated Interrupt Disable flipflop is clear.

When an interrupt occurs the program counter (PC) is stored in location 0 and the computer begins executing instructions at the location whose address is found in location 1. Location 1 contains the address for the routine that services interrupts. The service routine must execute an Interrupt Acknowledge (INTA) instruction to determine the device code of the device closest on the bus that is interrupting. The INTA instruction pulses the  $\overline{INTP}$  IN and INTA lines and the device interrupting, which is closest to the computer places its device code on the data bus lines  $\overline{DATA10}$ - $\overline{DATA15}$ . A more detailed description of interrupt servicing can be found in (3,5).

If the interrupt feature of the Nova is used then a short interrupt handler must be written to process interrupts from the interface. The interrupt handler can be broken down into several steps. These steps and their associated execution time in microseconds are:

Finish current instruction	1.35	maximum
Store Program Counter	1.35	
Jump to service routine	1.20	
Execute:       INTA 3	2.55	

Execute: JMP @DISP,3 1.35 + 1.20

Execute program segment to read appropriate data

Execute: INTEN 3.15

The total time to execute the above, less the time to execute the program segment to read the appropriate data, is 12.15 microseconds.

The minimum time required to read the azimuth increment count is given in Figure 3.4 and is 10.80 microseconds. Thus to handle an interrupt and read the azimuth increment count would take:

$$12.15 + 10.80 = 22.95$$

microseconds. The time necessary to read the range bin intensity word is also given in Figure 3.4 and is 17.05 microseconds. Thus to handle an interrupt and read the range bin intensity word would take:

$$12.15 + 17.05 = 29.75$$

microseconds. The range bin intensity word (4 range bin intensities) must be transferred within 21.48 microseconds. Using the interrupt feature of the Nova this time requirement cannot be satisfied. Therefore the interrupt feature of the Nova cannot be used and need not be implemented in the interface. This simplifies the required implementation of the interface status circuit to that shown in Figure 2.10.

To set the Busy and Done flipflops to known states, an optional code letter can be concatenated to the input/output instruction mnemonic. The code letters are:

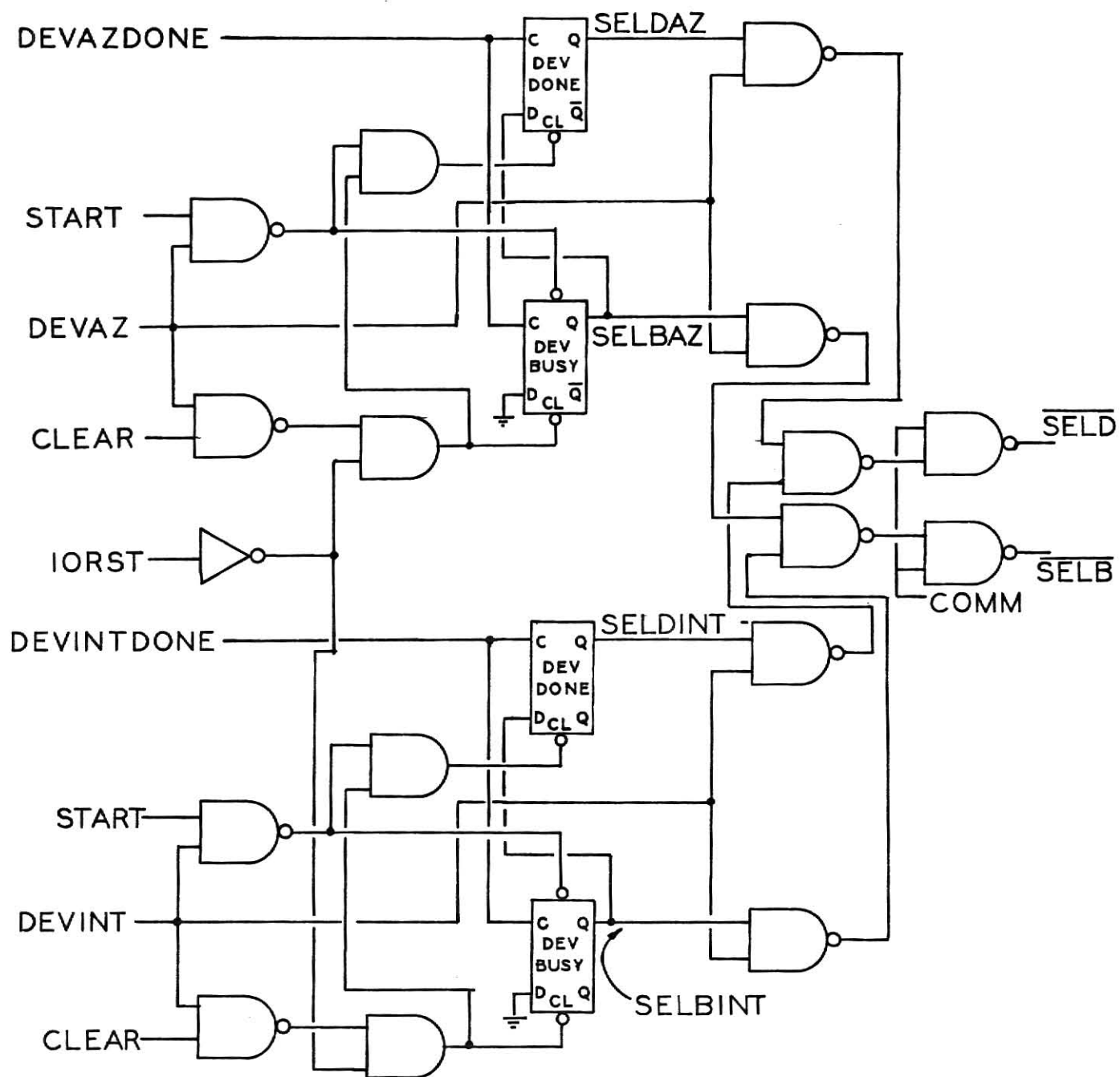


FIGURE 2.10  
BUSY-DONE FLIPFLOPS IMPLEMENTED



none	Leave Busy and Done as they are
S	Set Busy and clear Done
C	Clear both Busy and Done
P	Special pulse on IOPLS line

To distinguish which data value, azimuth increment count or range bin intensity word, is available to be transferred it is necessary to have two device codes and two Busy-Done circuit implementations or a status register to indicate which data value is available. To reduce the complexity of the circuit and increase the speed of the control program two device codes are used on this interface.

The signal DEVINTDONE is the clock pulse C4 which moves the Range Bin Intensity Input Register to the Range Bin Intensity Output Register. When this move occurs the data is ready to be transferred to the computer memory so SELDINT is set and the SELD line is pulled low when DEVINT is active.

The signal DEVAZDONE is the output of a monostable which is triggered each time the azimuth increments by 5.625 degrees. The azimuth increment count is in the Azimuth Increment Counter. A new sequence of moving the range bin intensities must start each time DEVAZDONE occurs.

#### AZIMUTH INCREMENT COUNTER

The azimuth is determined from decoding the 13 bit binary output of the shaft encoder into two pulses. The first pulse occurs when the azimuth reaches 0 degrees, and

the second pulse occurs each 5.625 degree increment. The incremental azimuth will be used to clock a counter whose output will be a binary value equal to the number of 5.625 degree increments.

The shaft encoder makes a full count in 32 revolutions of the radar antenna. The shaft encoder starts numbering with bit 1 and this convention gives the correlation between bit place and degree of revolution shown in Table 2.1. From Table 2.1 it can be seen that only SE1 to SE8 need be decoded to detect 0 degrees. When SE1-SE8 are all low, or  $\overline{\text{SE1-SE8}}$  are all high, the antenna is pointing towards North and a data acquisition sequence can now be performed. The condition of zero azimuth decoded and a flipflop is set as shown in Figure 2.11. The lowest order bit of the shaft encoder cycles from low to high to low in 5.625 degrees. The falling edge of  $\overline{\text{SE1}}$  is used to clock the Azimuth Increment Counter. This allows the azimuth counter to be clocked 2.8125 degrees before the shaft encoder shows a multiple of 5.625 degrees, thus the data is valid when the counter output is read. The Azimuth Increment Counter output will be the binary number of increments of 5.625 degrees. Thus the azimuth data word is decoded to indicate which wedge of the radar's scan the data following will represent.

In four seconds the antenna sweeps 360 degrees and the lower 8 bits of the azimuth shaft encoder counts from  $2^0$  to  $2^7$ . Thus the time per count is:

$$T = (4 \text{ sec}) / (2^7 \text{ count}) = 31.25 \text{ millisec/count}$$

Table 2.1 Correlation Between Azimuth Angle and  
Shaft Encoder Bit Position

Bit Position	Azimuth Angle
13	11520
12	5760
11	2880
10	1440
9	720
8	360
7	180
6	90
5	45
4	22.5
3	11.25
2	5.625
1	2.8125

Thus, each 5.625 degree increment will take 62.50 milliseconds.

The circuit in Figure 2.11 is the circuit for decoding the azimuth. To obtain a 1 to 0 state transition on the ZENB line all inputs to the 8-input nand gate must be high. This transition sets the Zero flipflop, which is cleared by an IORST instruction. When  $\overline{\text{ZENB}}$  is high the complement of AZ1 becomes the clock, AZCCP, to the Azimuth Increment Counter. The timing diagram for the circuit shown in Figure 2.11 is shown in Figure 2.12.

The outputs of the Azimuth Increment Counter are fed to the output gating network and placed on the data bus lines  $\overline{\text{DATA0}}-\overline{\text{DATA7}}$ . The Azimuth Increment Counter is read by executing a DIA instruction with device code 63.

### THE SYNCHRONIZATION SIGNALS

There are several signals required to synchronize transfer of the range bin intensities from the Integration and Quantization circuit to the computer interface. These are:  $\overline{\text{SHOLD}}$ ,  $\overline{\text{ZERO}}$ ,  $\overline{\text{ONE}}$ ,  $\overline{\text{READY}}$ , and  $\overline{\text{DUMP}}$ . The first three signals indicate when the intensity range bins are being transferred. The last two signals inform the Integration and Quantization circuit that the computer interface is ready to accept the range bin intensities.

Figure 2.13 shows how the signals  $\overline{\text{ZERO}}$ ,  $\overline{\text{ONE}}$ , and  $\overline{\text{SHOLD}}$  are generated on the Integration and Quantization circuit.

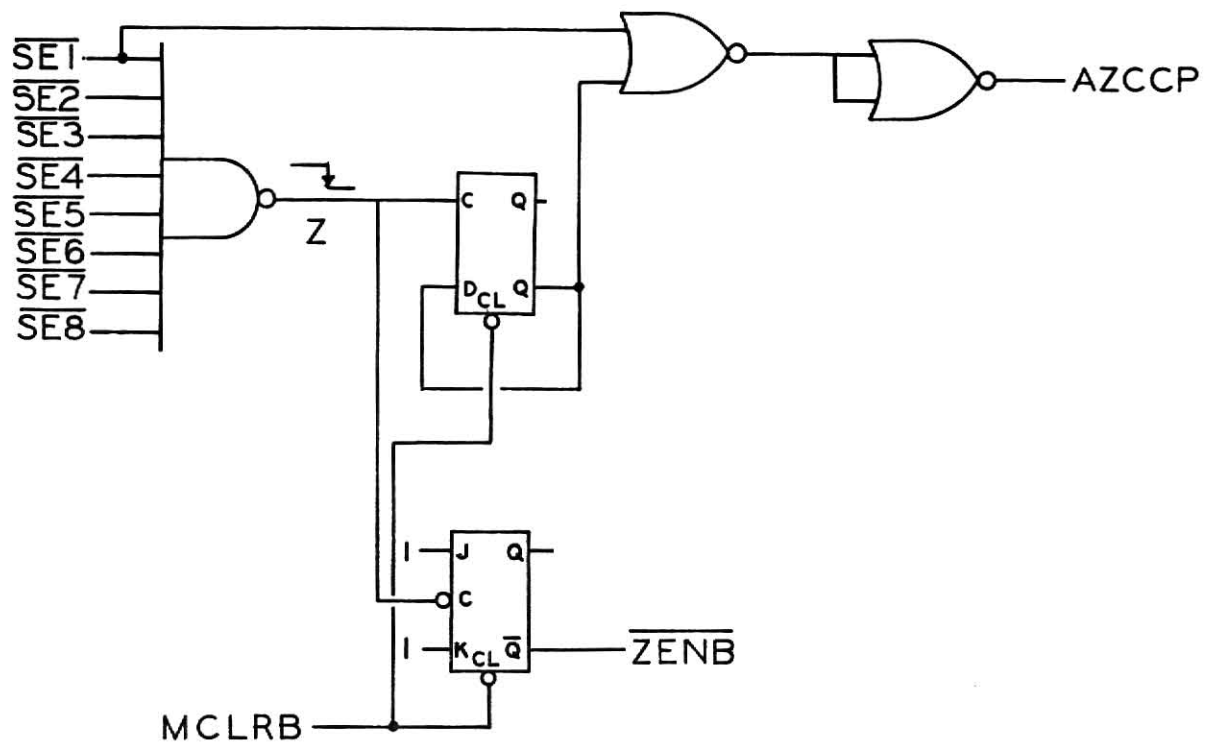


FIGURE 2.11  
DECODING THE AZIMUTH

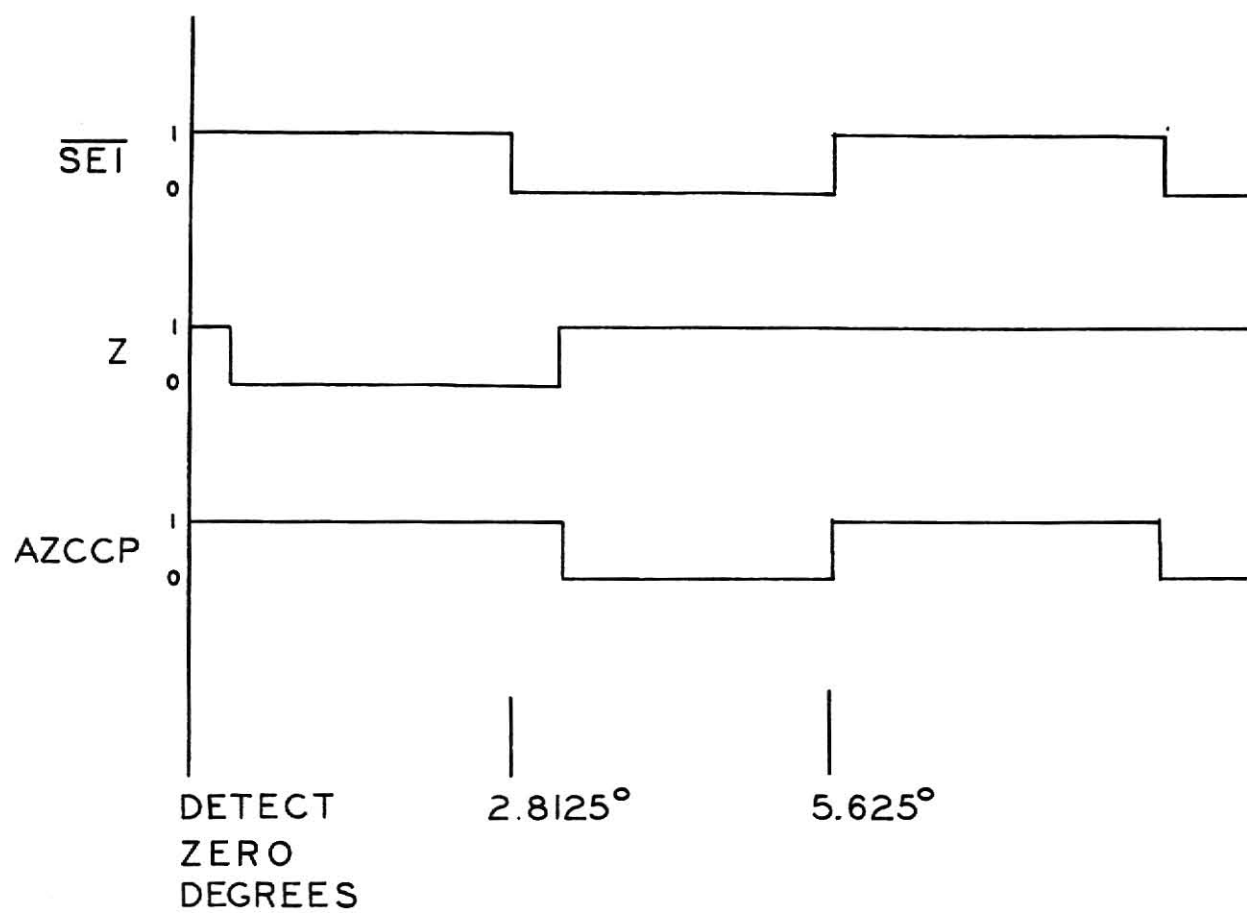
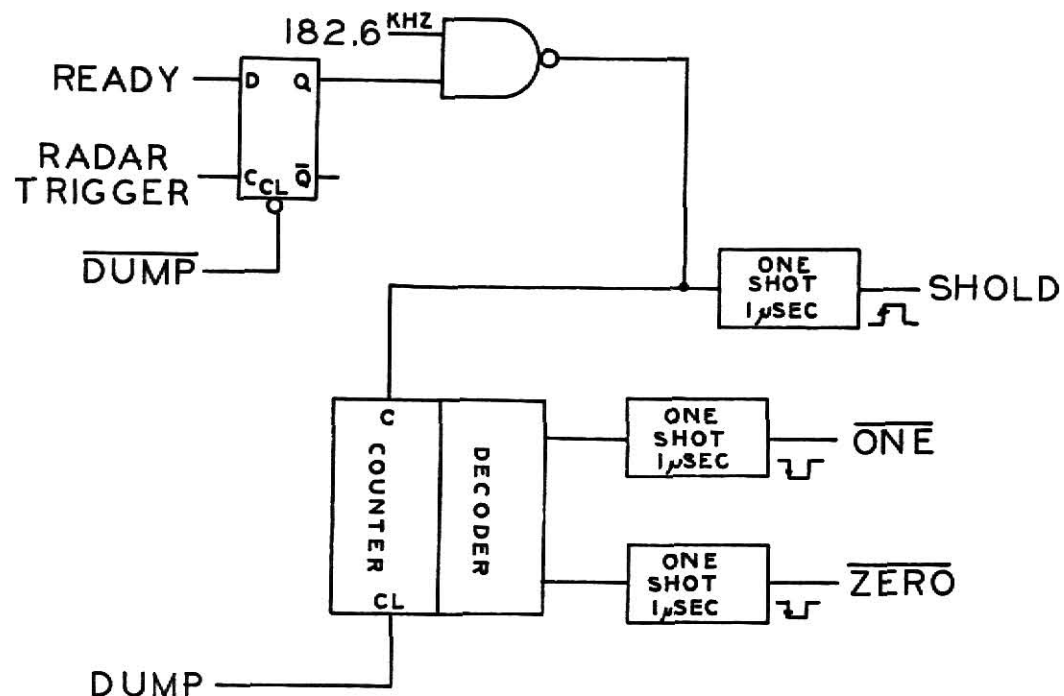
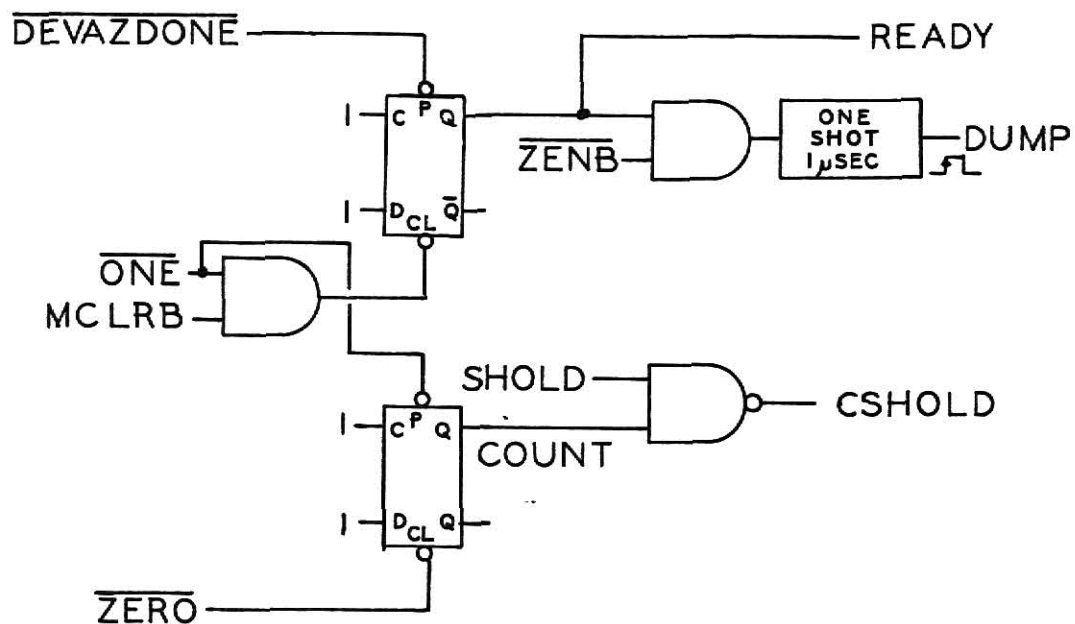


FIGURE 2.12  
TIMING DIAGRAM FOR DECODING  
THE AZIMUTH



-----  
 INTEGRATION AND QUANTIZATION  
 COMPUTER INTERFACE  
 -----



**FIGURE 2.13**  
 GENERATION OF THE COMPUTER INTERFACE  
 TO INTEGRATION AND QUANTIZATION CIRCUIT  
 SYNCHRONIZATION SIGNALS

Figure 2.11 illustrates how these are used in the computer interface and how READY and DUMP are generated.

The signal  $\overline{\text{ONE}}$  is a pulse that goes low for one microsecond when the range bin counter reaches the count of 1. The count of 1 is used because at this time the voltage on the capacitor for range bin 0 has been digitized and its value is present and stable at IN0-IN3. When the range bin count gets to one the pulse  $\overline{\text{ONE}}$  sets the count flipflop and the sample and hold signal SHOLD becomes the clock to the Range Bin Intensity Input Register Clock Counter. The counter is advanced by 1 and the data from range bin 0 is clocked into the first four bits of the Range Bin Intensity Input Register.

The signal  $\overline{\text{ZERO}}$  is a pulse that goes low for one microsecond when the range bin counter reaches a count of zero. The signal  $\overline{\text{ZERO}}$  is used to stop the sequence of movement of range bin intensities into the Range Bin Intensity Input Register by resetting the Count flipflop. When the Count flipflop is reset then SHOLD can no longer clock the Range Bin Intensity Input Register Clock Counter. The  $\overline{\text{ZERO}}$  pulse is also used to reset the Ready flipflop.

READY is synchronized with the radar trigger pulse and used as an enable line so that the  $\overline{\text{ZERO}}$  and  $\overline{\text{ONE}}$  pulse occurs only when READY is high.

The DUMP pulse goes high for one microsecond when a zero azimuth crossing is detected and after the range bin intensities for a azimuth increment have been transferred and



stored in the computer's memory. The DUMP pulse is used to reset the READY and Radar Trigger Pulse synchronization flipflops and to discharge the integration capacitors on the integration and quantization circuit.

### THE INTENSITY REGISTERS

Two registers are required in order to buffer the range bin intensities and provide enough time to move them to the computer's memory. The Range Bin Intensity Input/Output Register diagram is shown in Figure 2.14. The Range Bin Intensity Input Register is divided into four four-bit nibbles. Each nibble contains one range bin intensity. The range bin intensities come once every 5.37 microseconds, thus moving four of them into a 16 bit register giving computer 21.48 microseconds to read the data into the memory.

The Range Bin Intensity Input Clock Counter is a two bit Gray code counter whose output is decoded to provide four clock pulses as shown in Figure 2.14. The 74107 dual JK flipflop is used for the Gray code counter so that the falling edge of CSHOLD clocks the counter. When SHOLD goes high, CSHOLD goes low, advancing the counter, if READY is high. The rising edge of the clock cycles C1-C4 are used to clock the data into the D flipflops used in the Range Bin Intensity Input Register. When SHOLD goes high, so will , after a few gate delays, one of the clock cycles. The data

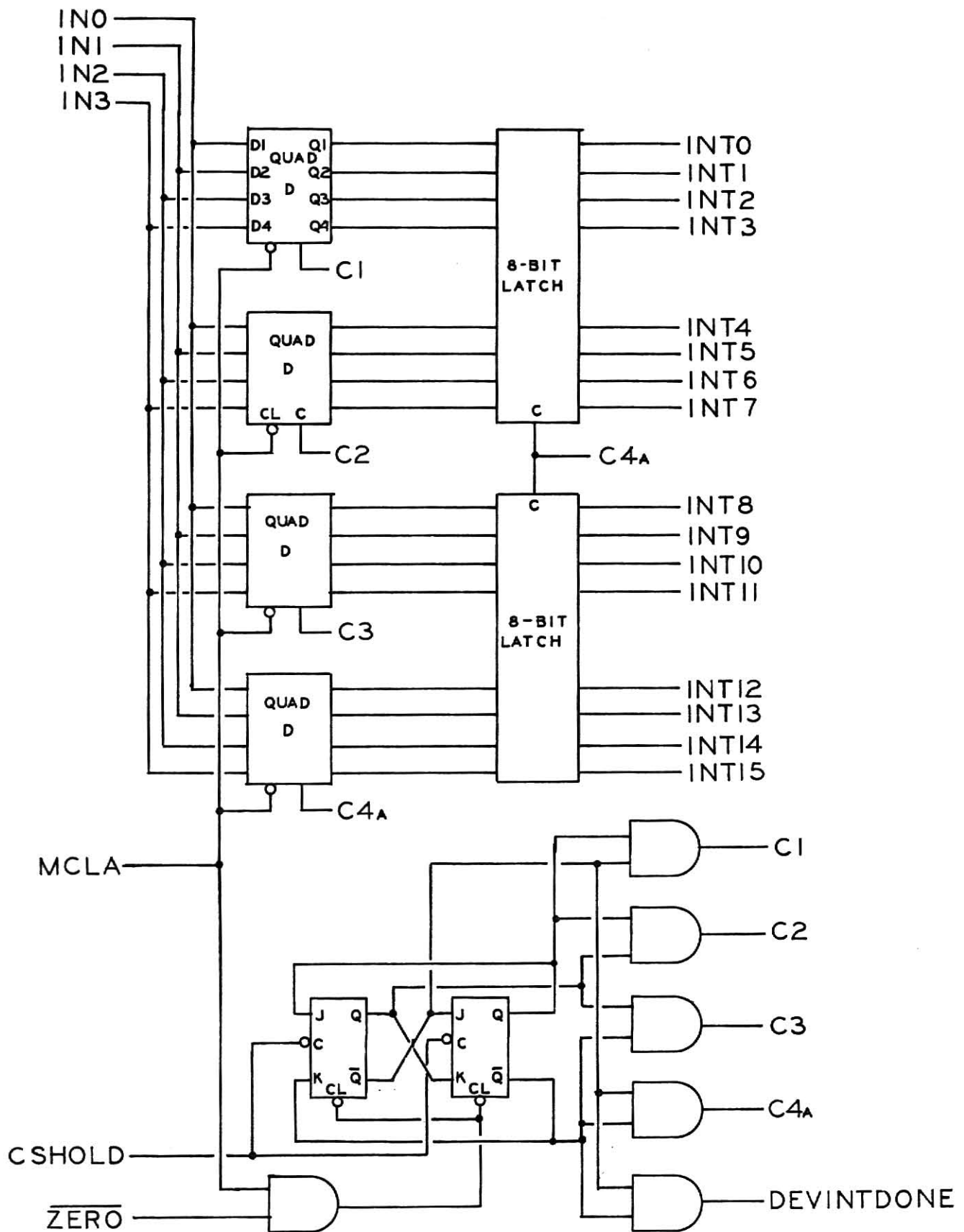


FIGURE 2.14  
INTENSITY STORAGE REGISTERS

for a range bin is stable on IN0-IN3 when SHOLD goes high and is the data for the capacitor N-1 where N=count appearing on the Range Bin Counter on the Integration and Quantization Circuit.

Clock phase C4 is used for several purposes. First C4 is used to enable the latches that compose the Range Bin Intensity Output Register. D-type flipflops could be used and a delayed C4 clock phase could clock the data from the Range Bin Intensity Input Register to the Range Bin Intensity Output Register, but use of the latches reduces package count and still maintains data transfer reliability. C4 is also used to set the Done flipflop that is associated with the intensity registers. When this Done, SELDINT, is high the data is ready to be transferred from the Range Bin Intensity Output Register. Appendix A contains the circuit diagrams for the Computer Interface as implemented. Also provided is a system timing diagram.

#### THE SELF TEST CIRCUITS

To provide a means of testing the interface and control program two self test circuits are provided. Figure 2.15 shows the circuit used to simulate the azimuth. The azimuth consists of an 8 bit binary value and is decoded to provide the proper azimuth increment. The circuit board IAQ-2 (see Appendixes A and B) contains the pseudo-azimuth circuit of Figure 2.15 and is removed when the shaft encoder provides

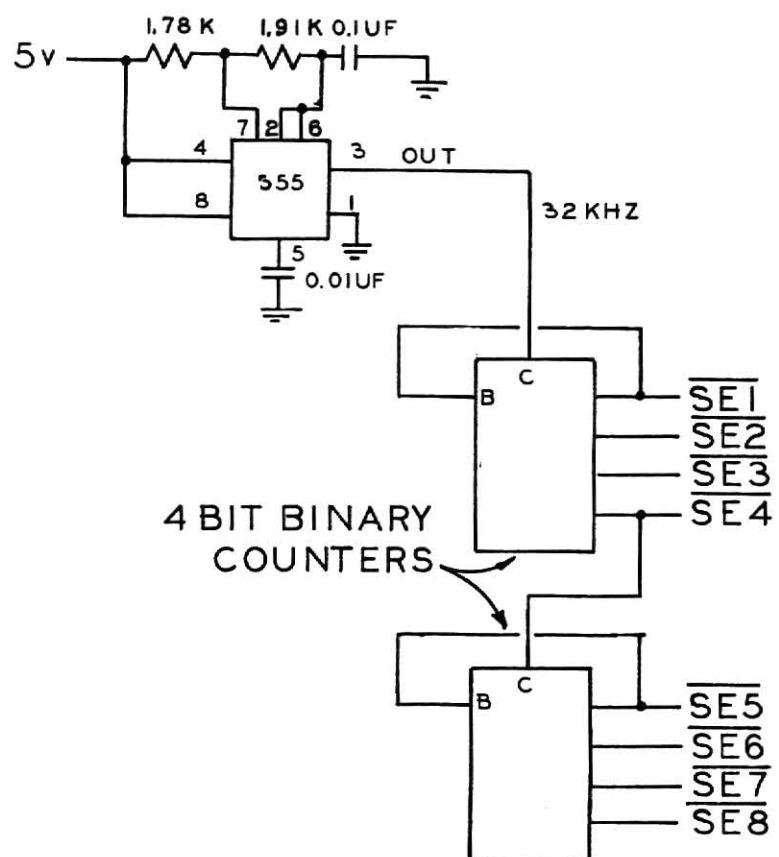


FIGURE 2.15  
THE SELF TEST CIRCUIT USED TO  
SIMULATE THE AZIMUTH

the azimuth. The voltage levels of the shaft encoder output are converted to TTL voltage levels when the shaft encoder provides the azimuth signal. The shaft encoder low order bit has a frequency of 32 Hz, and the self test circuit provides a low order bit with a frequency of 31.2 Hz.

Figure 2.16 shows the circuit that simulates the output and synchronization signals of the Integration and Quantization circuit. The synchronization signals  $\overline{\text{ONE}}$ ,  $\overline{\text{ZERO}}$ , DUMP and READY are described in the section "The Synchronization Signals". The circuit shown in Figure 2.16 is implemented on board IAQ-1. The clock frequency in the Integration and Quantization circuit that is used to select the range bins is 186.2 KHz. The self test circuit uses a 181 KHz. clock to select the range bins. This means that the self test circuit outputs a range bin once every 5.52 microseconds instead of every 5.37 microseconds as provided by the Integration and Quantization circuit.

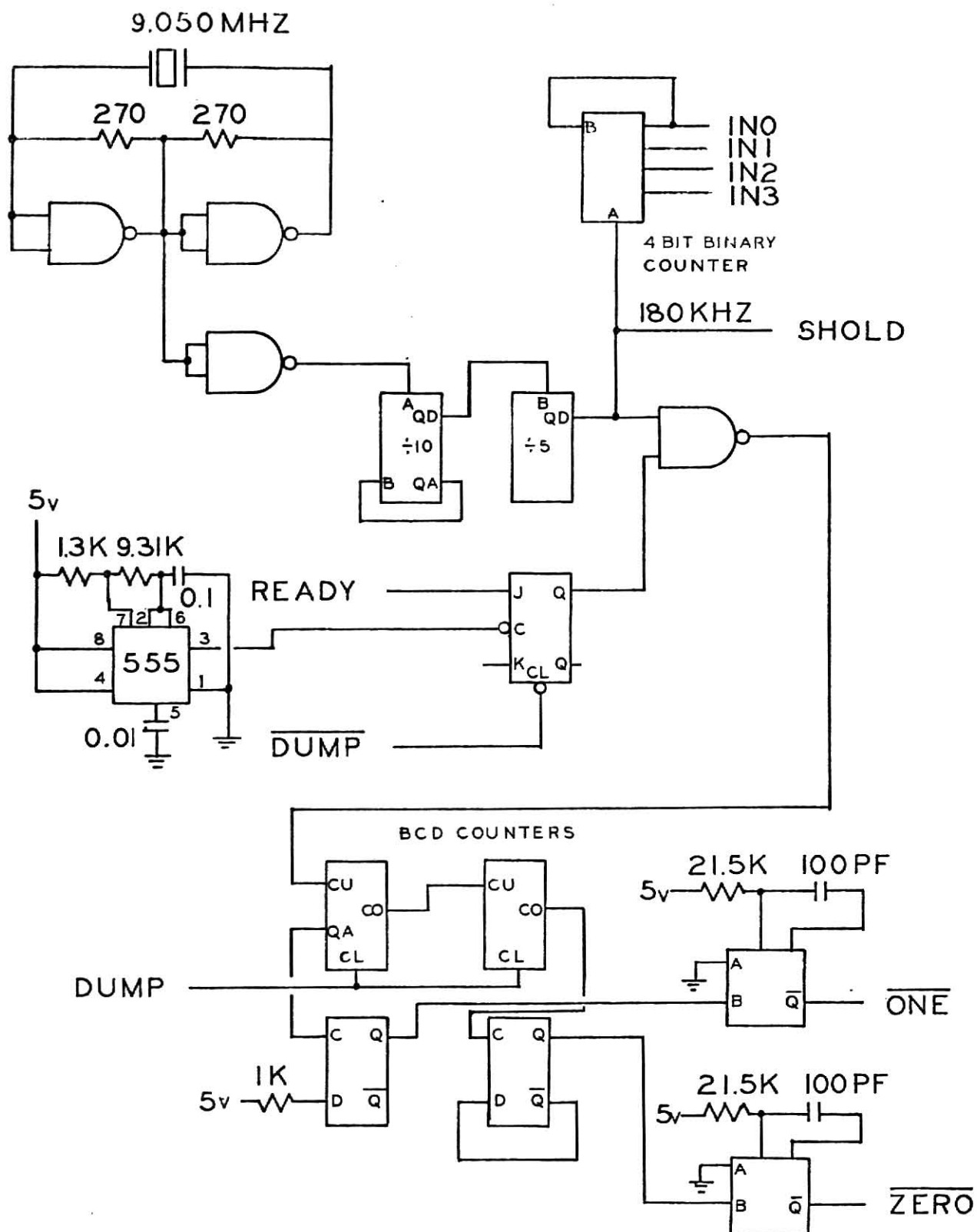


FIGURE 2.15  
GENERATION OF THE COMPUTER INTERFACE TO INTEGRATION AND  
QUANTIZATION CIRCUIT SYNCHRONIZATION SIGNALS

### Chapter 3

Since a computer is used as an intermediate step in moving the data from the interface to the disk there must be a program to control movement of the data from the interface to the computer's memory and then out to the disk. At the present time the Department of Electrical Engineering has two Caelus CD-22 disk drives and a Disk Management Operating System supporting these disk drives. Figure 3.1 is a high level flowchart of the program that moves data from the interface to the disk. With the memory size available on the Nova 1200 it is possible to store in memory all of the range bin words from one scan, and then once the scan is completed, to move the data to the disk. This program is available on the computer system's disk under the name RDCP.BI, Radar Data Collection Program (.BI stands for executable "binary" ). The function of each block of the flowchart shown in Figure 3.1 is discussed in this chapter. A source listing of RDCP is provided in Appendix C.

The term "data acquisition record" will be defined as the 64 azimuth increment counts and their 50 associated range bin intensity words.

#### THE INITIALIZATION BLOCK

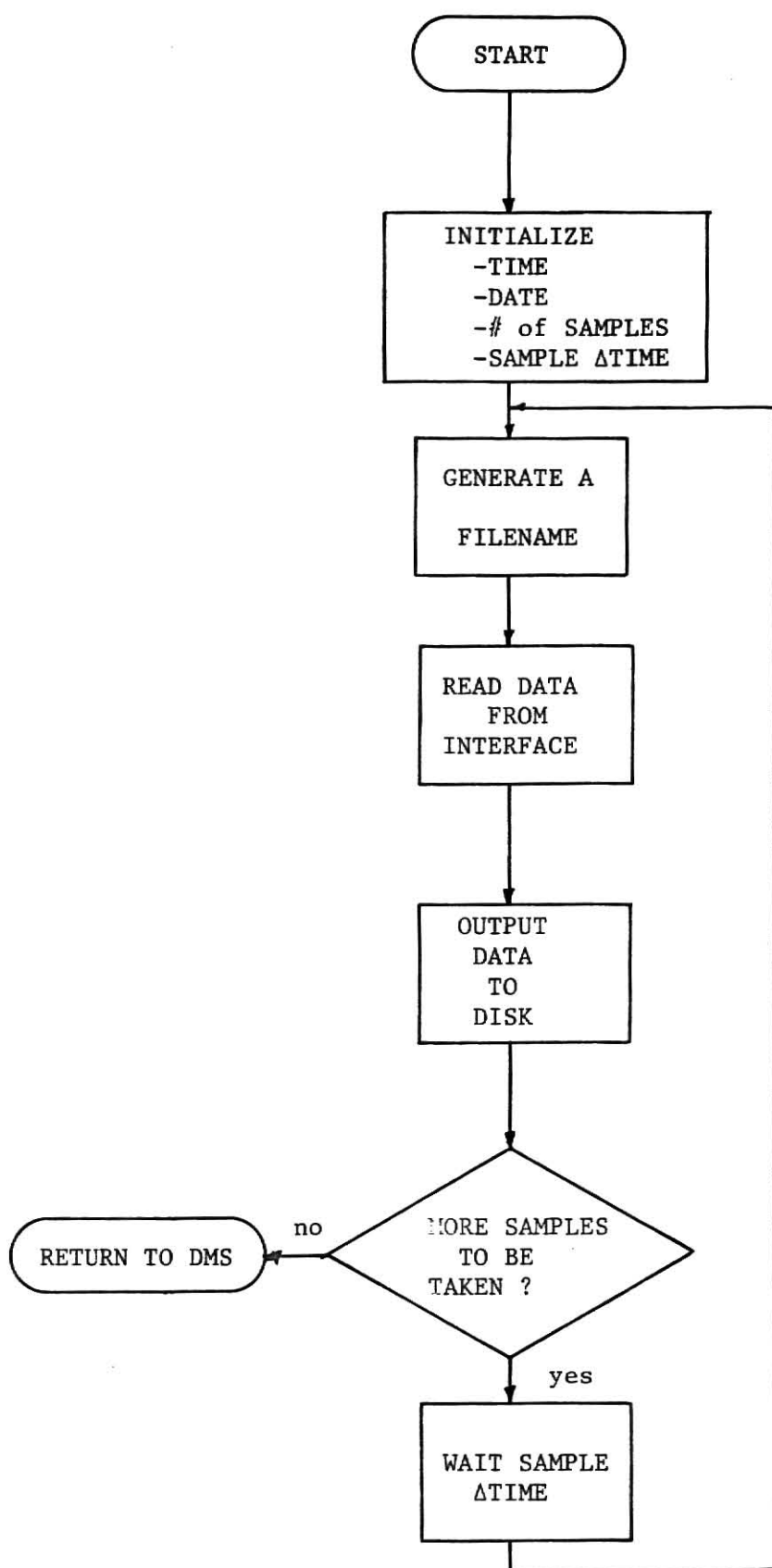


FIGURE 3.1

HIGH LEVEL FLOWCHART OF RADAR INTERFACE CONTROL PROGRAM



Prior to obtaining data from the interface the user is asked to initialize the program parameters. The user is asked the time of day and date so that each data acquisition record can be labeled for easier bookkeeping. The user can also specify how many data acquisition sequences he wishes to record and the interval between them. Figure 3.2 shows the flowchart for the initialization block of RDCP.

The program will type out the statement:

TIME OF DAY (H/M/S) :

The user must respond with the time of day in a 24 hour format. Each field (hours, minutes and seconds) must be a two digit integer value. Any invalid responses will result in the question being repeated. A typical response would be:

TIME OF DAY (H/M/S) : 09/30/25

where the "/" symbols will be provided by the program.

After obtaining the time of day the program will type out the statement:

DATE (MONTH/DAY/YEAR) :

The user must respond with the date. The year field consists of only the last two digits of the current calendar year, e.g., input 75 for 1975. The month and day responses must be two digit integer values as was required for the time fields. The response for the date February 1, 1975 would be:

DATE (MONTH/DAY/YEAR) : 02/01/75

where the "/" symbols will be provided by the program. Any incorrect response results in the repetition of the

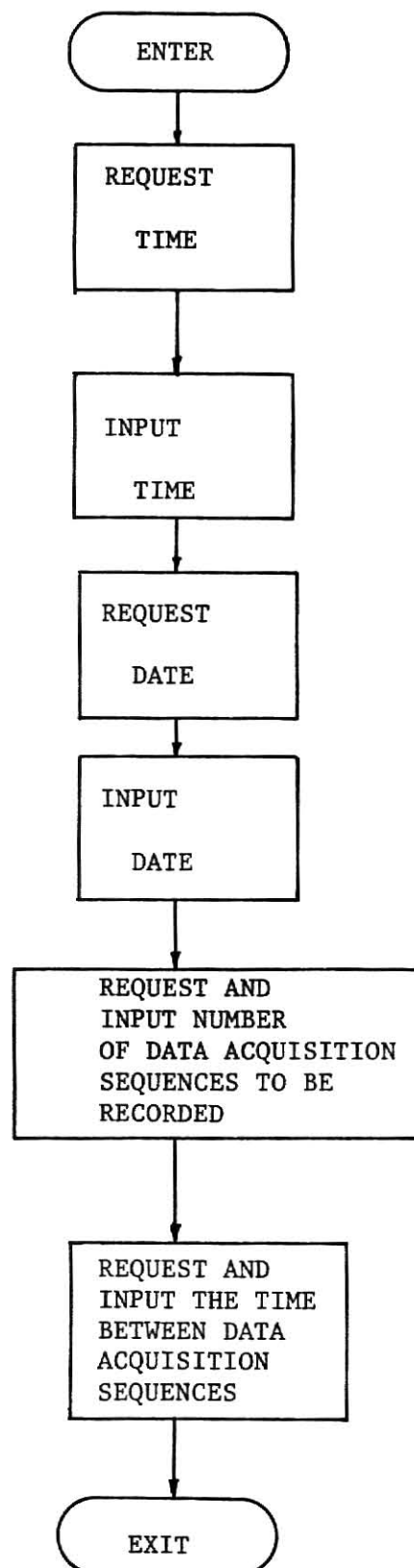


FIGURE 3.2

THE INITIALIZATION BLOCK FLOWCHART

question.

The time of day and date are stored in the first six words of the data acquisition record. The order of storage is:

Word	Contents
0	Hours
1	Minutes
2	Seconds
3	Month
4	Day
5	Year

Each word will contain the binary equivalent of the BCD integers from that field. For example the year 1975 would be entered in as 75 and appear in word 5 in binary as : 0000000001001011.

The program is initialized to obtain and store up to 100 data acquisition records. The user responds to the question:

NUMBER OF RECORDS DESIRED :

with the three digit integer number of data acquisition sequences to be recorded. Entry of a non-numeric character or a value greater than 100 is invalid and the question is repeated. The user signals the program that the entire number has been entered into the computer by pressing the RETURN key. The time interval between successive data acquisition sequences can also be specified. The time interval is set by responding to the question:

## DATA ACQUISITION SEQUENCE INTERVAL

TIME (H/M/S) :

with the desired time in a 24 hour format as required for the time of day entry.

Once the program is initialized, space must be provided on the disk for the data from the acquisition records. The next block will create a unique file name for each data acquisition record.

## GENERATE A FILE NAME

To reduce the size of the control program a specific family of file names are used and generated in this block of the control program. The data is stored on the disk and is accessed under the file name RDmn.DA where mn=00 ... 99. There is a unique file name associated with the data from each data acquisition sequence. The family of file names used will allow a maximum of 100 data acquisition sequences to be stored on the disk.

This program is written such that it must be executed under the Disk Management Operating System (DMS). Associated with each file on the disk is a table that is placed in the disk file directory. This table is called the Work File Parameter Table or WFPT. The WFPT contains the file name, the starting sector address, the location in memory where the file is placed, and the file identification. The format of the Work File Parameter Table is shown in Table 3.1.

Table 3.1 Work File Parameter Table Organization

Word	Contents
0-2	File name, up to six seven bit ASCII characters packed right to left. For an odd number of characters the last character is left justified.
3	Starting Sector Address
4	Lower Core Limit if the file is binary. Sector size if the file is source.
5	Upper Core Limit if the file is binary. Zero if the file is source or relocatable binary.
6	Start Address if the file is binary. One if the file is relocatable binary. As initialized if the file is source.
7	File ID, two ASCII characters packed right to left.

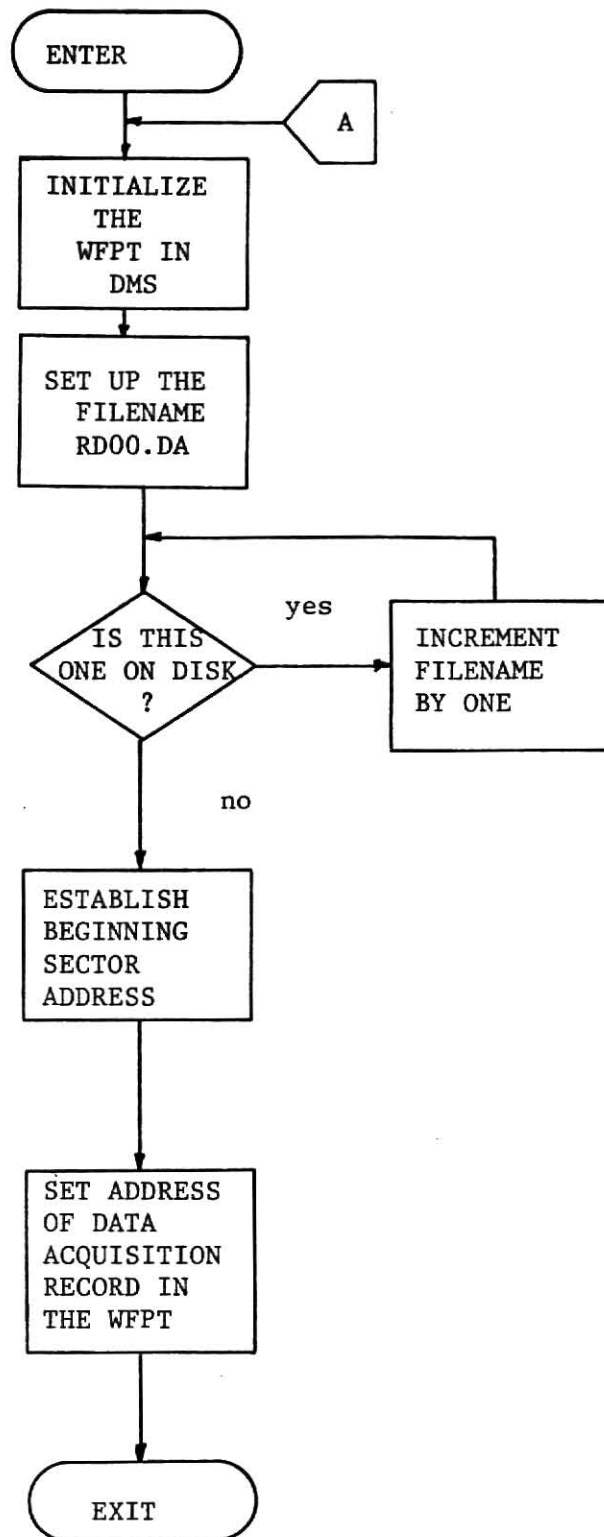


FIGURE 3.3

FLOWCHART FOR DETERMINING FILENAME

Table 3.2 relates the WFPT word to its contents. Using this format the user can read any data acquisition record one disk sector at a time to any desired buffer area. A detailed description of the use of DMS and the resident routines for reading and writing data to the disk can be found in the Caelus Memories, Inc, Disk Management System Manual (1).

### THE DATA ACQUISITION SEQUENCE

After the disk file name is created and the program parameters are initialized the data is moved from the interface to the computer's memory. Execution time in this segment is a critical factor. The interface supplies four data values in one computer word every 21.48 microseconds, thus the program must read the data from the interface, store it in the computer's memory, and be ready to read the next value in this 21.48 microsecond time frame. Figure 3.4 shows the program segment and instruction times for moving the data from the interface to the computer's memory.

Once all of the data is stored in the computer's memory it is then placed on the disk. The next block is the program segment required to accomplish this.

### WRITING DATA TO THE DISK

Figure 3.4 Instructions and Times  
For the Transfer Segment of the  
Control Program RDCP

	LDA	2,COUNT	
	COM	2,2	
START:	IORST		
	NIOS	66	
	SKPDN	66	2.55
	JMP	.-1	1.35
	DIAS	0,66	2.55
	STA	0,@TABLE	2.55 + 1.20 + 0.60

total time to read and store  
azimuth increment count : 10.80 microseconds

	NIOS	67	3.55
	SKPDN	67	2.55
	JMP	.-1	1.35
	DIAS	0,67	2.55
	STA	0,@TABLE	2.55 + 1.20 + 0.60
	INC	2,2,SZR	1.35
	JMP	.-6	1.35

total time to read and store  
range bin intensity word : 17.05 microseconds

Note: all times are given in microseconds



Table 3.2 Contents of the Work File Parameter Table

Word	Contents
0	Characters 'DR'
1	Character digits 'nm'
2	A 1 in bit 0, all other bits 0
3	Beginning sector address
4	Number of sectors used
5	A 1 in bit 0, all other bits 0
6	A 1 in bit 0, all other bits 0
7	Characters 'AD'

The WFPT was set up prior to executing the data acquisition sequence and the data is now in a memory buffer. Table 3.3 shows the organization of the buffer. The first 12 bytes contain the date and time information. Bytes 13 to 6540 contain the azimuth increment counts and range bin intensities for one 360 degree sweep of the radar antenna. Figure 3.5 shows the word structure for the intensity values.

The data is placed on the disk using the sequence of steps shown in the flowchart of Figure 3.6.

The CATAL and DSKW are resident DMS routines and their use is explained in Appendix D of (1). Once the data is on the disk it is left to the user to provide programs that will process it further.

#### PREPARATION FOR THE NEXT DATA ACQUISITION SEQUENCE

The decision block of the flowchart in Figure 3.1 determines if more data acquisition sequences are to be recorded. The Wait block provides the required interval between the data acquisition sequences. The expanded flowchart for these two blocks is shown in Figure 3.7.

During initialization the number of samples to be taken was entered by the user. The number of samples requested by the user is decremented by one each time a data acquisition sequence has been performed. After all data acquisition sequences have been performed control is returned to the

Table 3.3 ORGANIZATION OF BYTES FOR ONE DATA ACQUISITION

SEQUENCE		
AZIMUTH	AZIMUTH BYTES	INTENSITY BYTES
5.625	12 - 13	14 - 113
11.25	114 - 115	116 - 215
16.875	216 - 217	218 - 317
22.5	318 - 319	320 - 419
28.125	420 - 421	422 - 521
33.75	522 - 523	524 - 623
39.375	624 - 625	626 - 725
45	726 - 727	728 - 827
50.625	828 - 829	830 - 929
56.25	930 - 931	932 - 1031
61.875	1032 - 1033	1034 - 1133
67.5	1134 - 1135	1136 - 1235
73.125	1236 - 1237	1238 - 1337
78.75	1338 - 1339	1340 - 1439
84.375	1440 - 1441	1442 - 1541
90	1542 - 1543	1544 - 1643
95.625	1644 - 1645	1646 - 1745
101.25	1746 - 1747	1748 - 1847
106.875	1848 - 1849	1850 - 1949
112.5	1950 - 1951	1952 - 2051
118.125	2052 - 2053	2054 - 2153
123.75	2154 - 2155	2156 - 2255
129.375	2256 - 2257	2258 - 2357
135	2358 - 2359	2360 - 2459
140.625	2460 - 2461	2462 - 2561
146.25	2562 - 2563	2564 - 2663
151.875	2664 - 2665	2666 - 2765
157.5	2766 - 2767	2768 - 2867
163.125	2868 - 2869	2870 - 2969
168.75	2970 - 2971	2972 - 3071
174.375	3072 - 3073	3074 - 3173
180	3174 - 3175	3176 - 3275
185.625	3276 - 3277	3278 - 3377
191.25	3378 - 3379	3380 - 3479
196.875	3480 - 3481	3482 - 3581
202.5	3582 - 3583	3584 - 3683
208.125	3684 - 3685	3686 - 3785
213.75	3786 - 3787	3788 - 3887
219.375	3888 - 3889	3890 - 3989
225	3990 - 3991	3992 - 4091
230.625	4092 - 4093	4094 - 4193
236.25	4194 - 4195	4196 - 4295
241.875	4296 - 4297	4298 - 4397
247.5	4398 - 4399	4400 - 4499
253.125	4500 - 4501	4502 - 4601
258.75	4602 - 4603	4604 - 4703
264.375	4704 - 4705	4706 - 4805
270	4806 - 4807	4808 - 4907
275.625	4908 - 4909	4910 - 5009
281.25	5010 - 5011	5012 - 5111
286.875	5112 - 5113	5114 - 5213
292.5	5214 - 5215	5216 - 5315
298.125	5316 - 5317	5318 - 5417
303.75	5418 - 5419	5420 - 5519
309.375	5520 - 5521	5522 - 5621
315	5622 - 5623	5624 - 5723
320.625	5724 - 5725	5726 - 5825
326.25	5826 - 5827	5828 - 5927
331.875	5928 - 5929	5930 - 6029
337.5	6030 - 6031	6032 - 6131
343.125	6132 - 6133	6134 - 6233
348.75	6234 - 6235	6236 - 6335
354.375	6336 - 6337	6338 - 6437
360	6438 - 6439	6440 - 6539

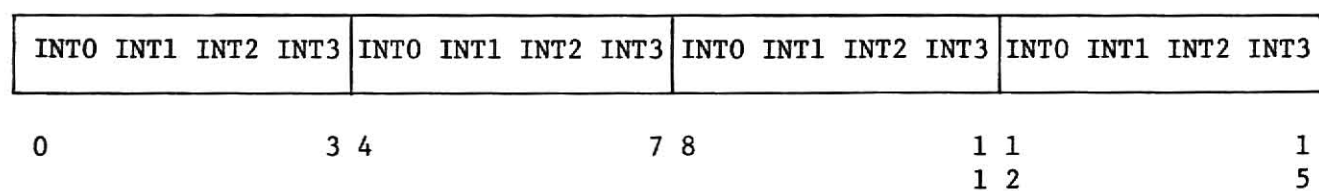


FIGURE 3.5

STRUCTURE OF THE RANGE BIN INTENSITY WORDS

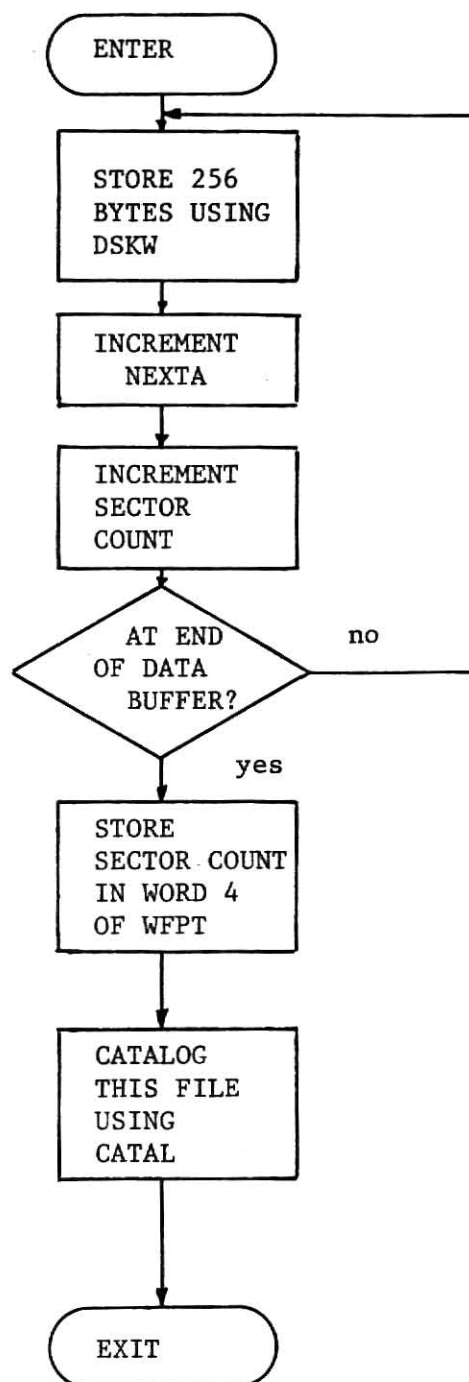


FIGURE 3.6

FLOWCHART FOR MOVING DATA FROM COMPUTER TO DISK

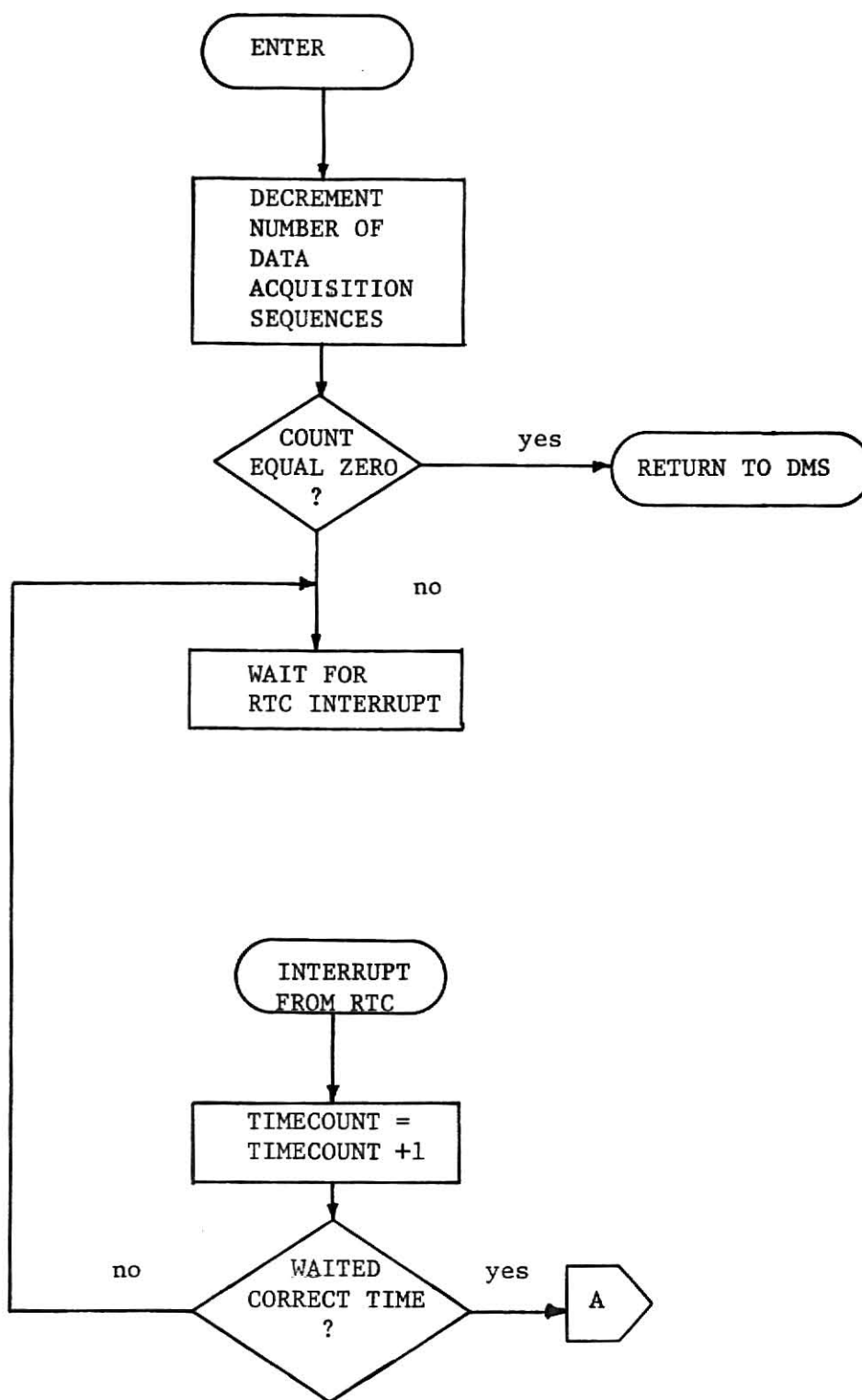


FIGURE 3.7

FLOWCHART FOR DETERMINING NEXT SEQUENCE

### Disk Management Operating System.

The Wait block is a loop that provides the time interval specified between successive data acquisition sequences. This loop counts the interrupts from the Real Time Clock (RTC) until the number of interrupts matches the MODULO 60 value of the hour, minutes, and seconds that were entered during initialization.

The time of the next data acquisition sequence is determined by adding the time to execute one data acquisition sequence and the interval time between data acquisition sequences to the time of the previous data acquisition sequence.

## Chapter 4

The data collected in one data acquisition sequence can be converted to precipitation rate and cumulative precipitation information. The cumulative precipitation information can be used in flash flood forecasting. The precipitation data collected using the system should greatly improve the quality of precipitation measurements for this area. These precipitation measurements could be input to a real time watershed modeling program. Analysis of the precipitation data over a period of time should yield better ground water predictions.

The system can also be used to record the movement of particularly large or severe storms. The computer graphics terminal can be used to display the recorded data, thus allowing analysis of such factors as the direction and velocity of movement of a storm. Work in the area of storm identification through pattern recognition is now possible with the data collected using this system.

If the beamwidth of the antenna is changed to a value less than 2.8125 degrees only minor modifications are necessary to allow collection of the range bin intensities every 2.8125 degrees. The Azimuth Interval Counter clock input, AZCCP, must be changed so that the counter will count in 2.8125 degree increments. This may be accomplished by



one of two methods.

The first method is to add a ratio gear to the shaft encoder so that the encoder will make a full count in 16 revolutions of the antenna. This modification will not require any circuit modifications as bit 1 of the shaft encoder would represent 1.40625 degree increments of the antenna. Bit 1 is multiplied by two to obtain the azimuth increment over which integration is performed.

The second method is to change the clock pulse AZCCP so that the falling and rising edges of AZCCP clock the Azimuth Increment Counter. This circuit modification is shown in Figure 4.1 and would require production of a new circuit board (Board B).

No program modifications are necessary for either of the above hardware modifications. The time between azimuth increments would still be much greater than the time to transfer all of the range bin intensities. The amount of memory used to store the range bin words is increased by a factor of 2. However the program maintains pointers to the beginning and end of the buffer containing the range bin words so all of the data would still be transferred to the disk.

Using the self test circuit, program runs indicate the control program and the interface are in working order.

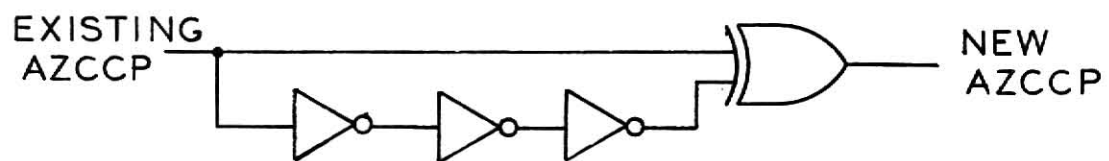


FIGURE 4.1  
MODIFICATIONS TO AZCCP

## BIBLIOGRAPHY

1. ——— Disk Management System (DMS)  
Manual, California, Caelus Memories Inc.
2. Andreasen, Carl C., A Manual for Programming  
the Data General Nova Minicomputers at the  
Assembly Language Level Under Three Operating  
Systems: Stand Alone Operating System, Disk  
Management Operating System, and Real Time Disk  
Operating System, (Unpublished Graduate problems  
paper, Kansas State University, 1975)
3. English, William, How to Use the Nova Computers,  
Massachusetts, Data General Corporation, 1971.
4. Hummels, Donald R., A Digitized Radar for  
Precipitation Measurements and Applications to Hydrology,  
Project Completion Report, Kansas State University, 1975.
5. Jesse, Richard H., The Design and Implementation  
of an Interface Between the Nova 1200 and Three  
Peripherals: An Analog-to-Digital Converter,  
A Digital-to-Analog Converter and an Incremental  
Tape Transport, (Unpublished Master's Degree Thesis,  
Kansas State University, 1974)

## APPENDIX A

### CIRCUIT DIAGRAMS

### Legend:

7404            Integrated Circuit Device Type

5 , A    Backplane Pin

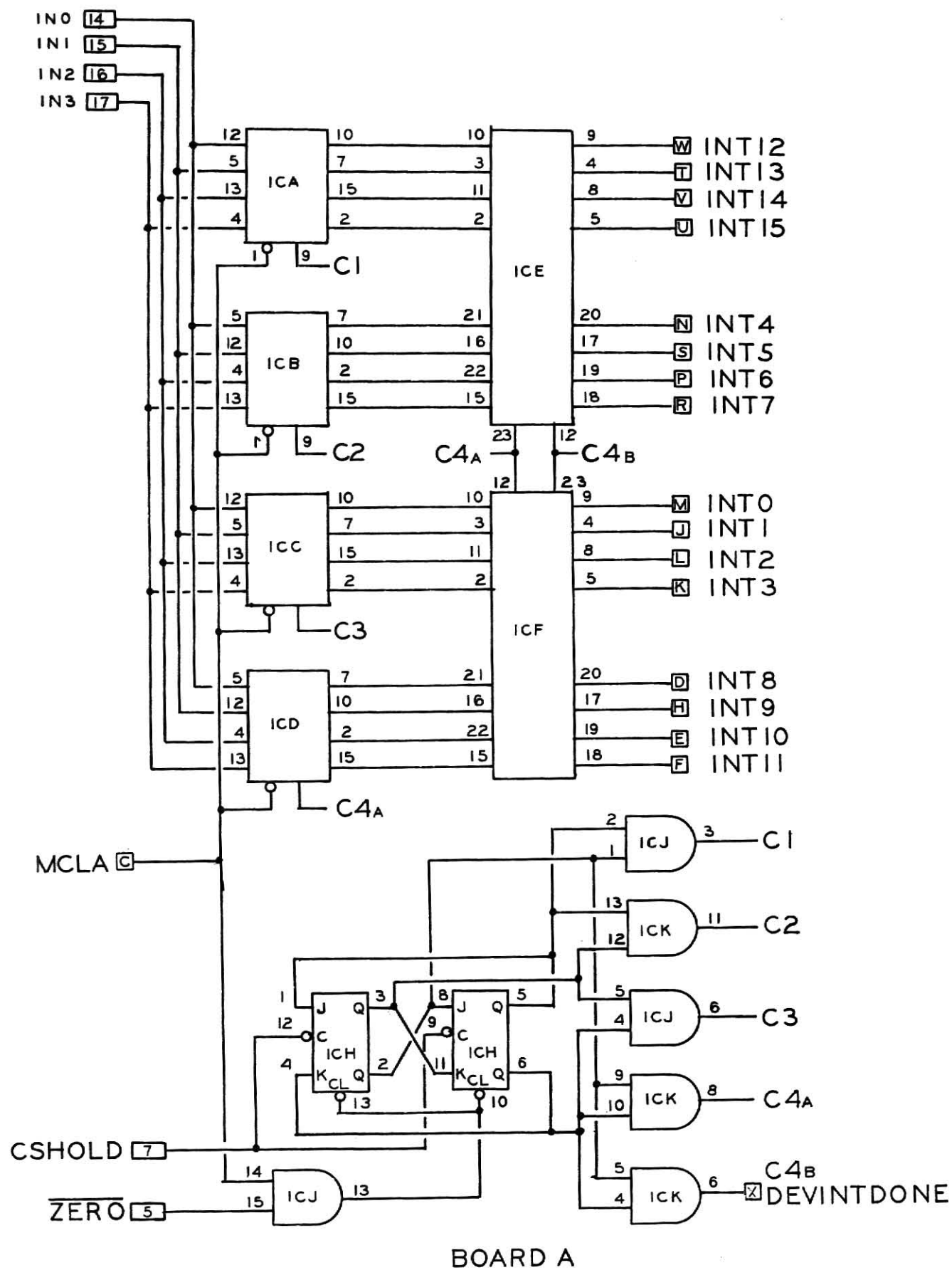
14            Computer Bus Pin

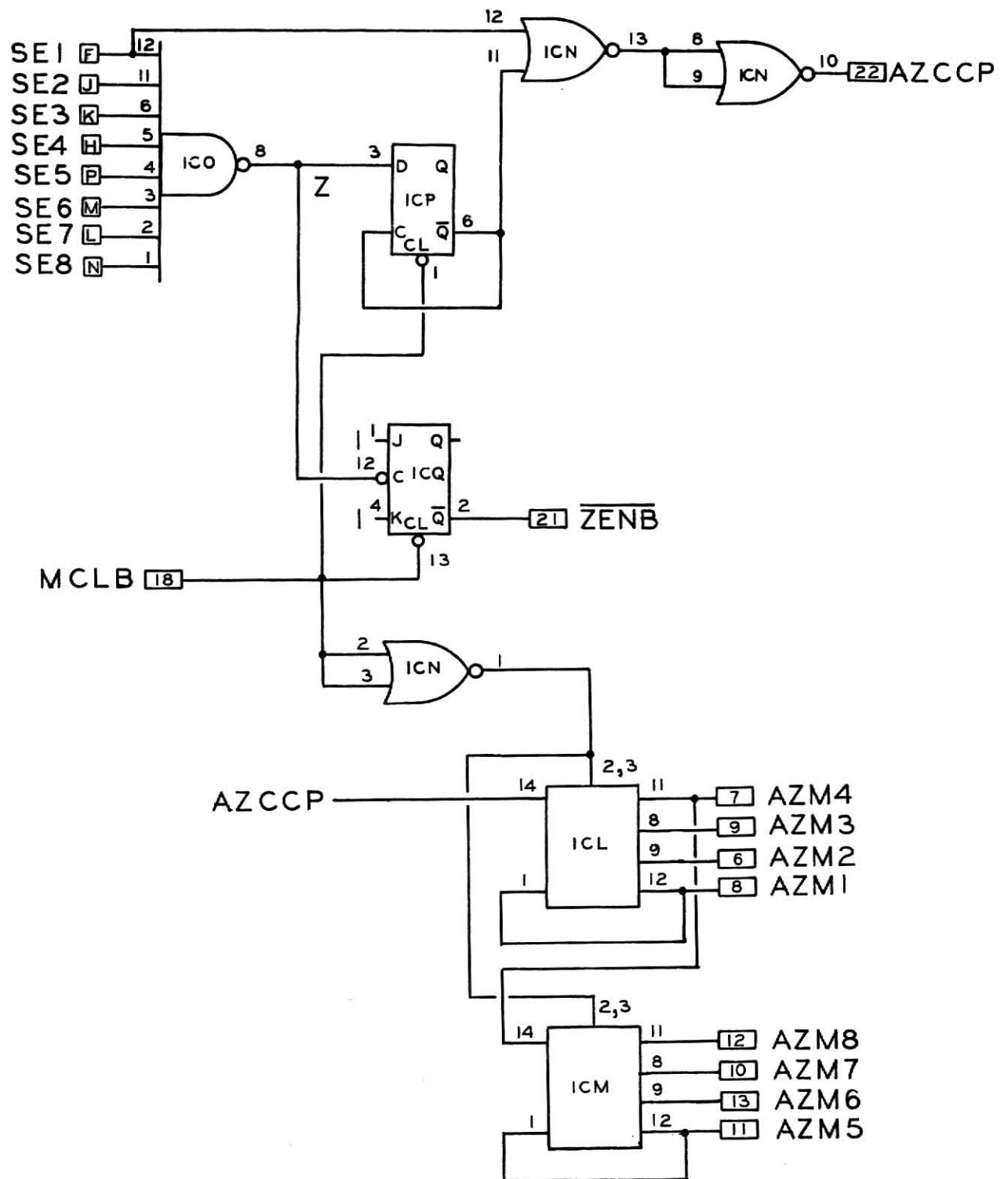
Fn            Feedthrough Number n

Jn            Jumper Number n

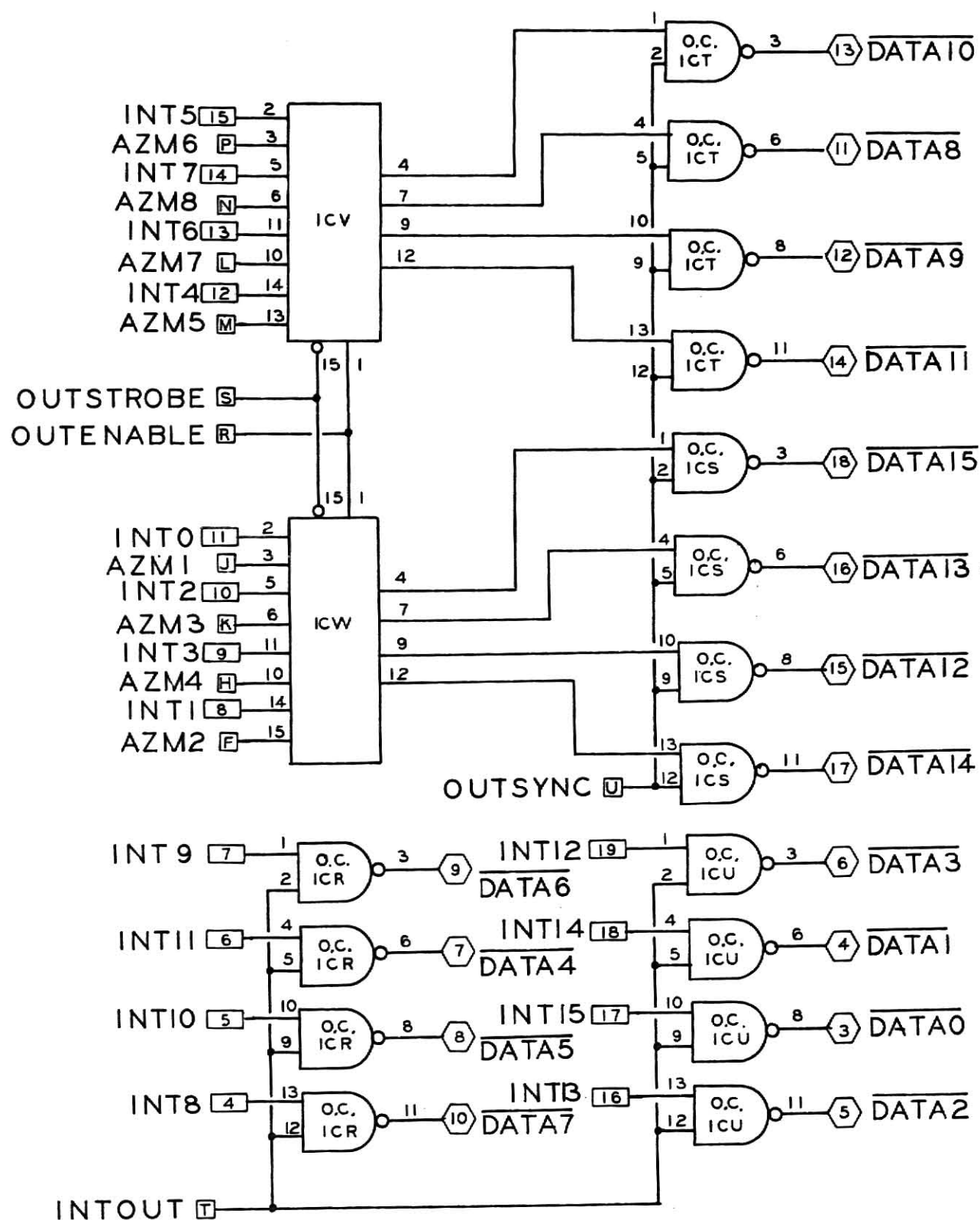
Capacitor

Resistor



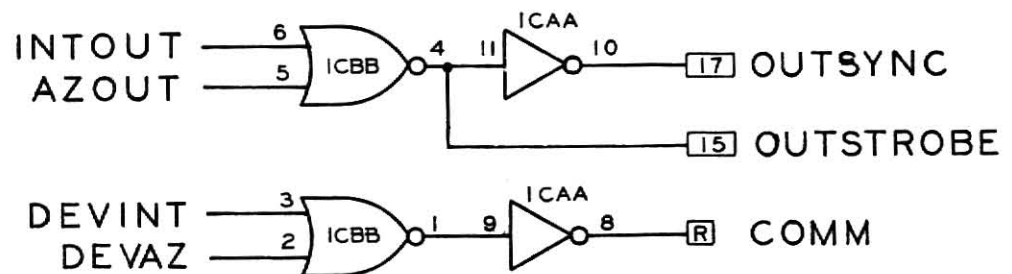
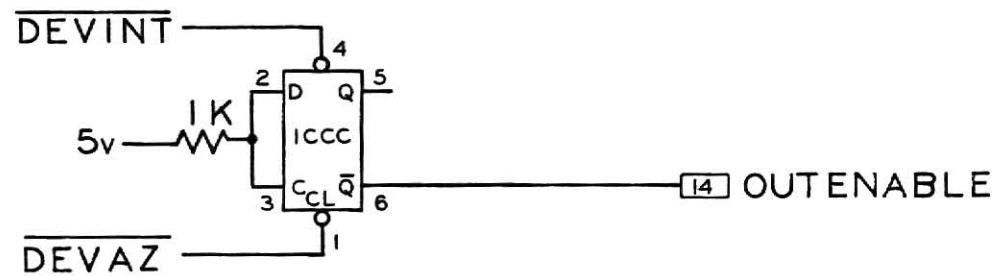
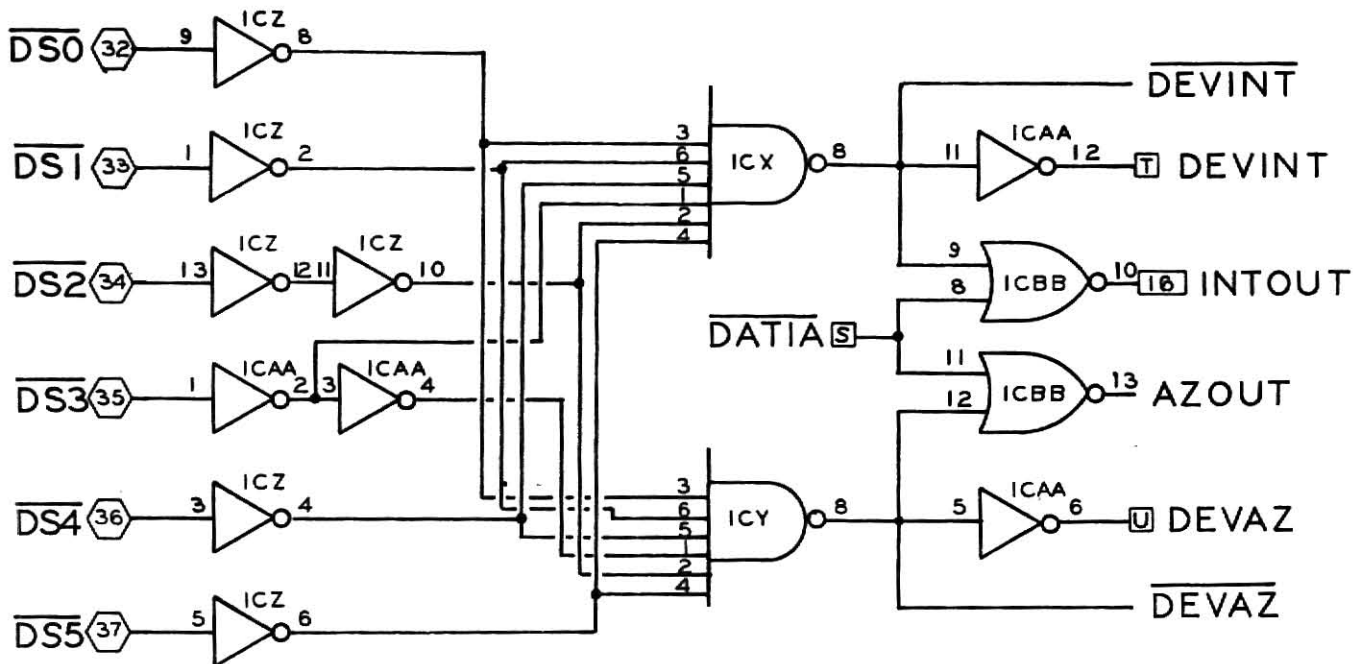


BOARD B



BOARD C

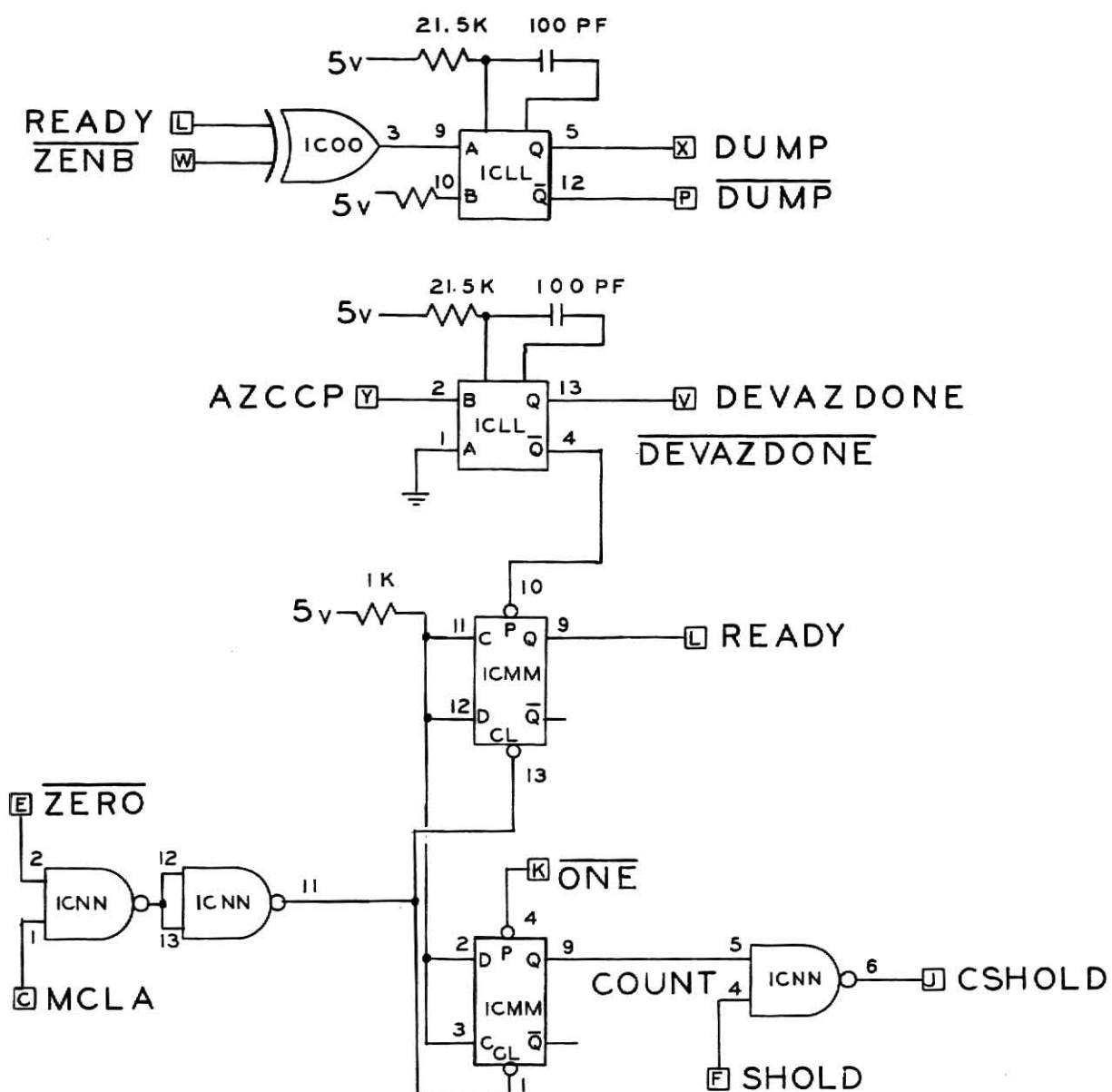




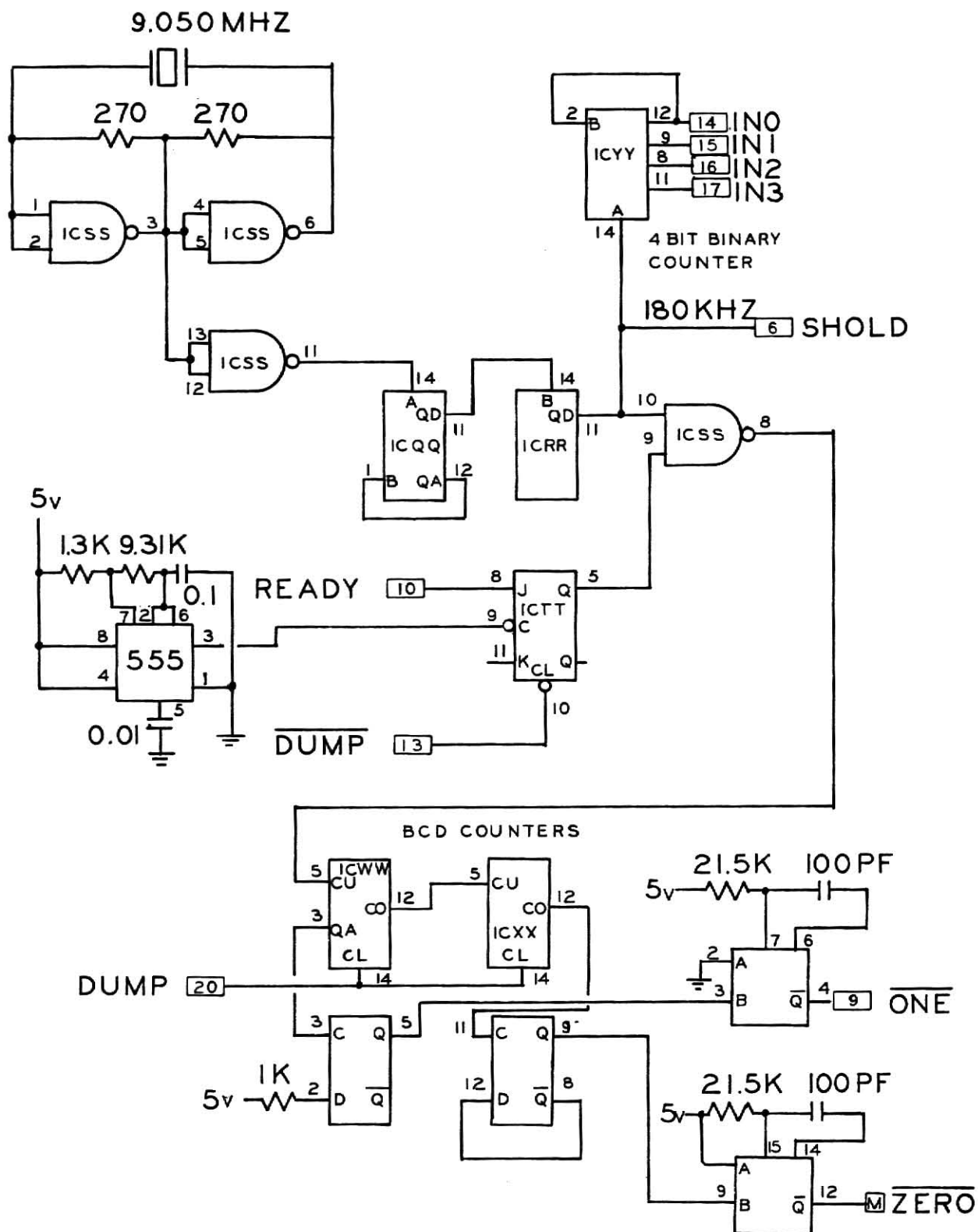
BOARD D



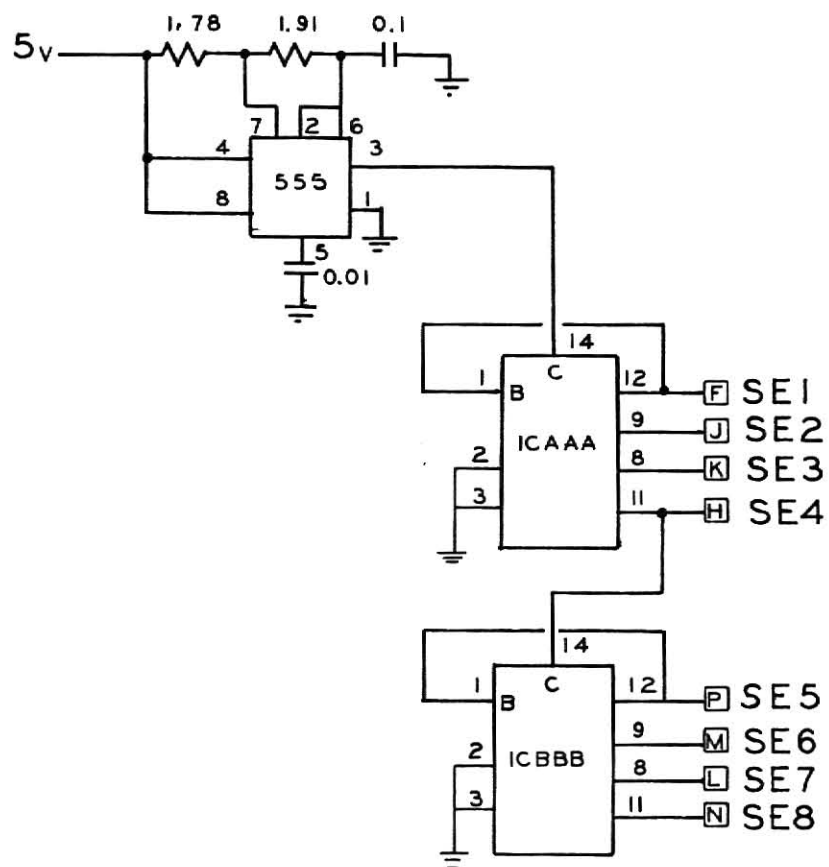
BOARD E



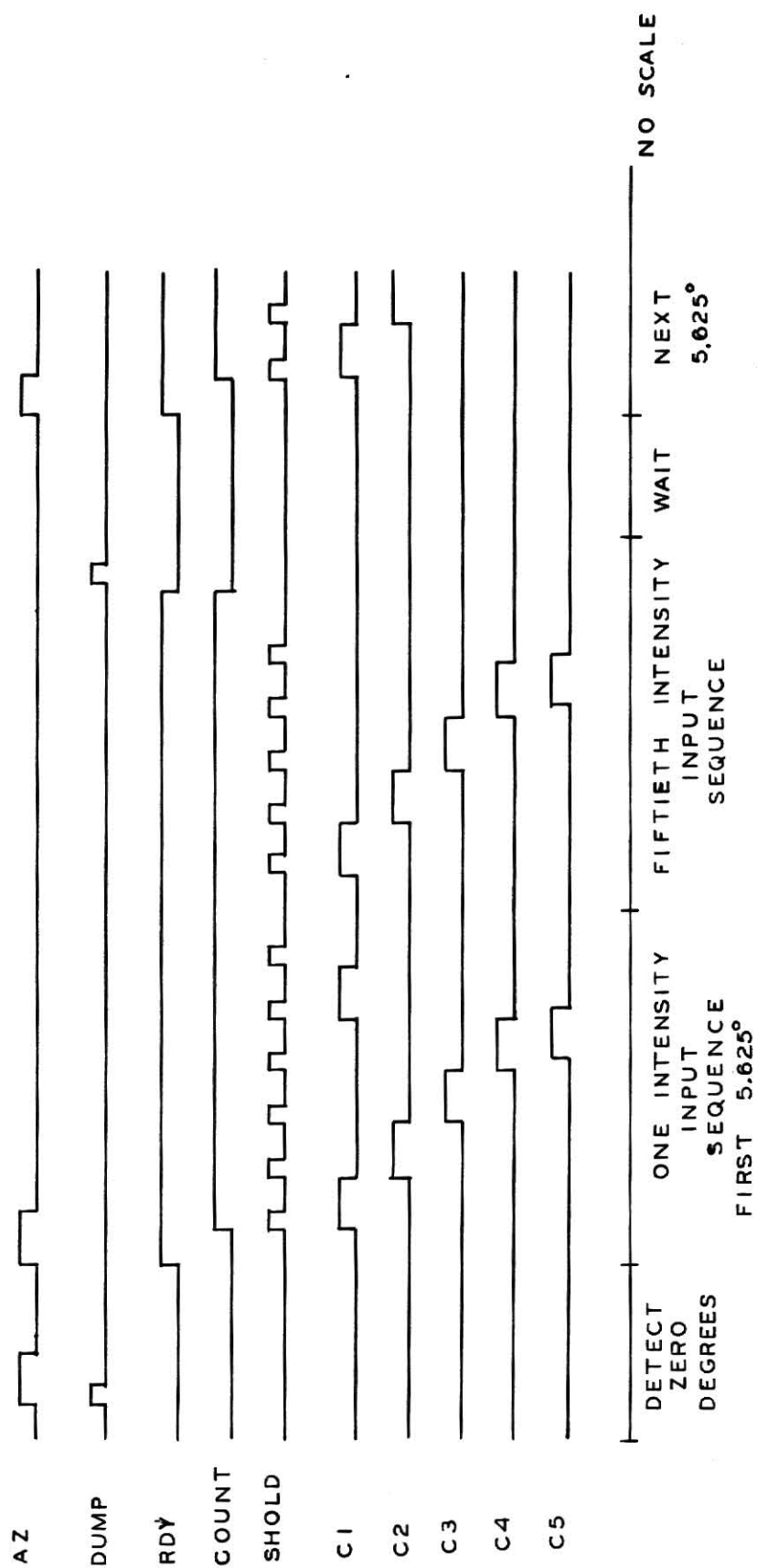
BOARD F



BOARD IAQ-1

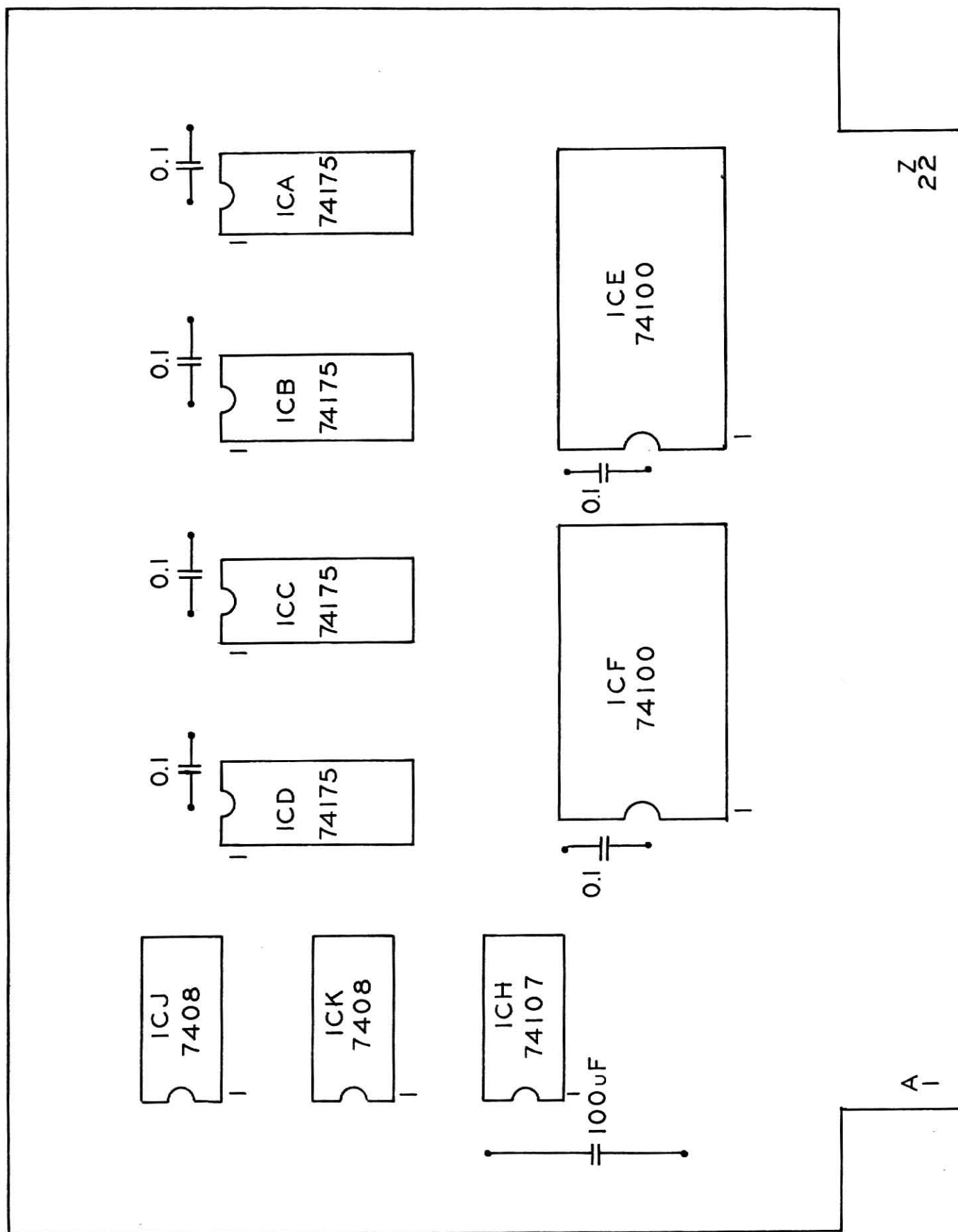


BOARD IAQ-2

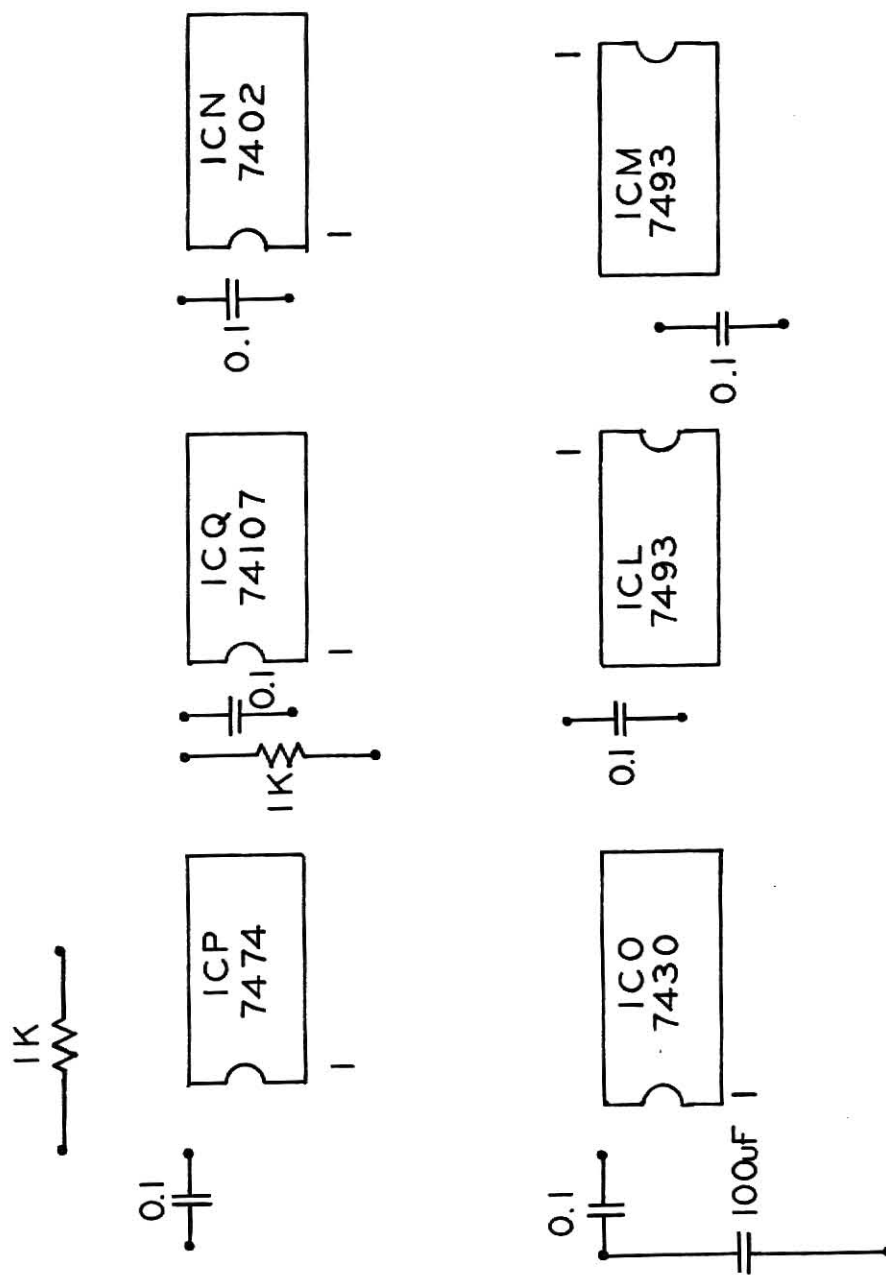


## **APPENDIX B**

### **COMPONENT PLACEMENT DIAGRAMS**

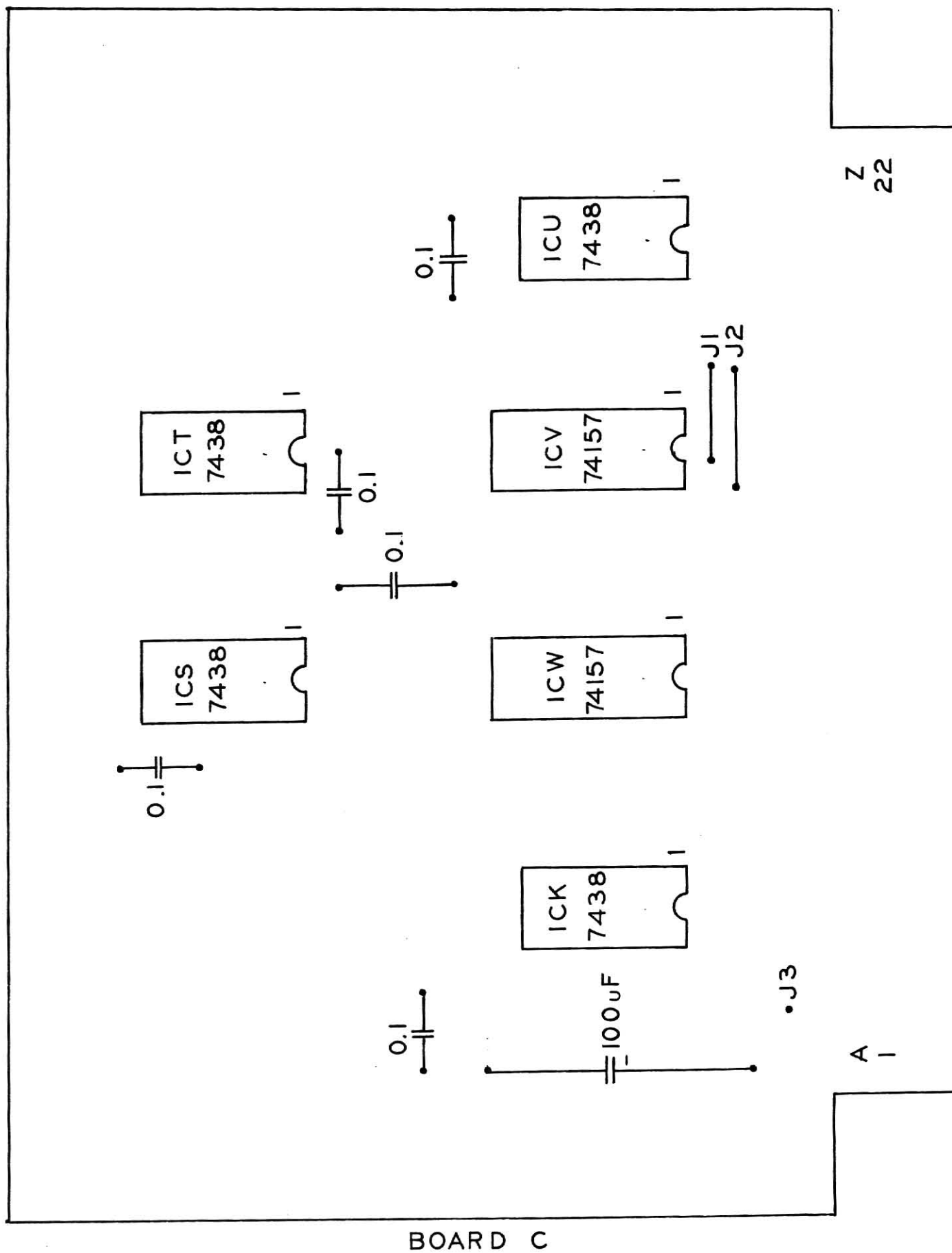


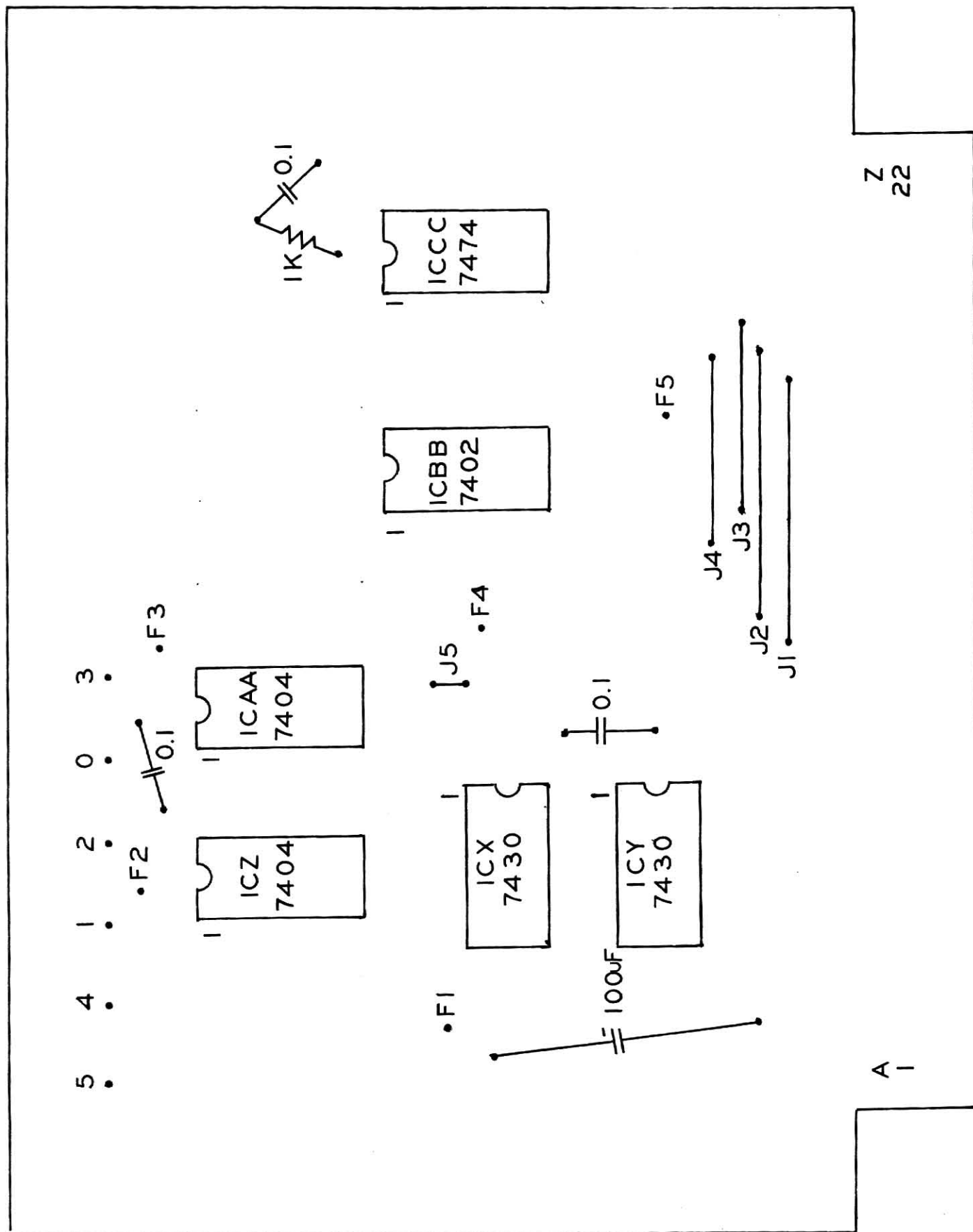


Z  
22

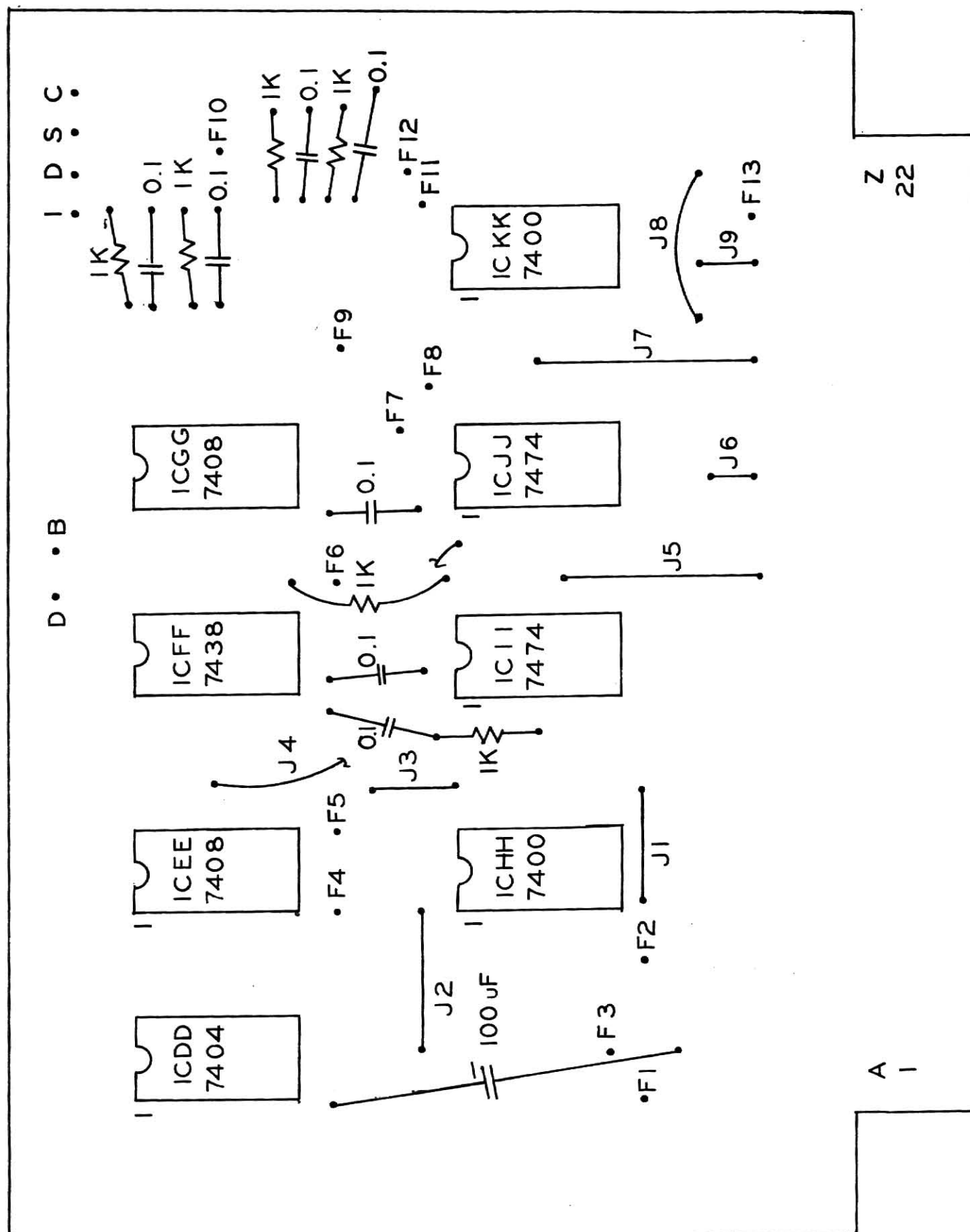
BOARD B

A  
1

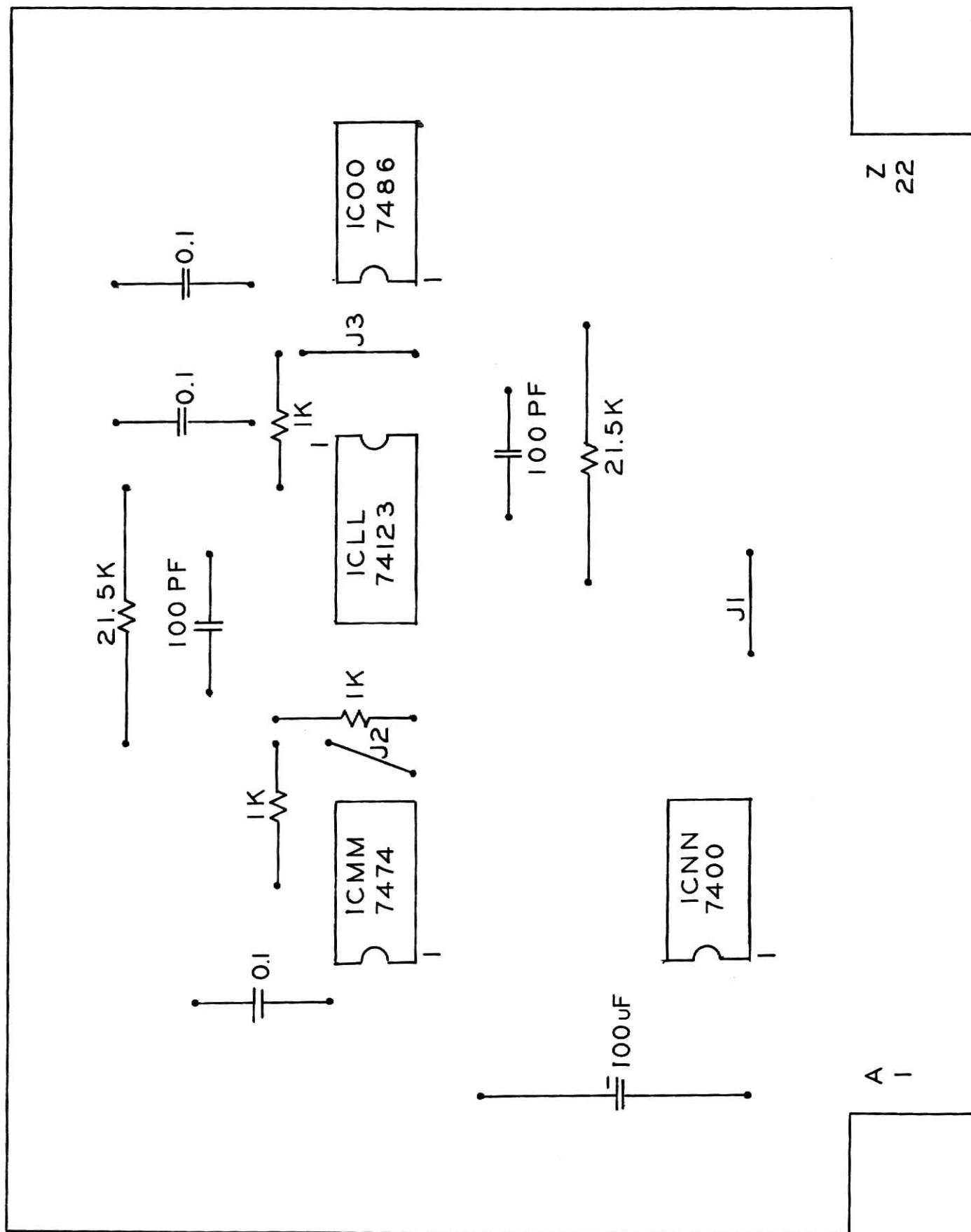




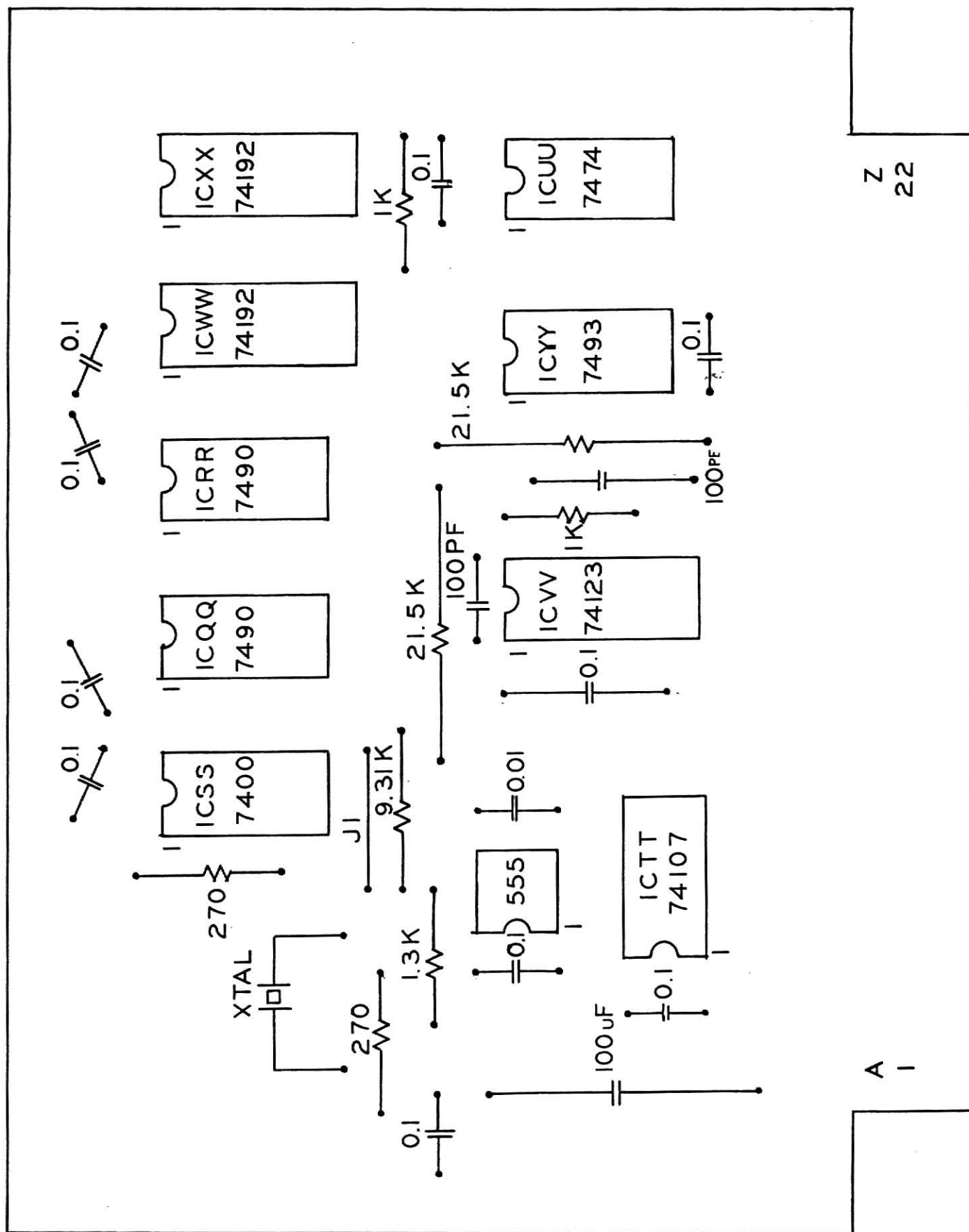
BOARD D



BOARD E

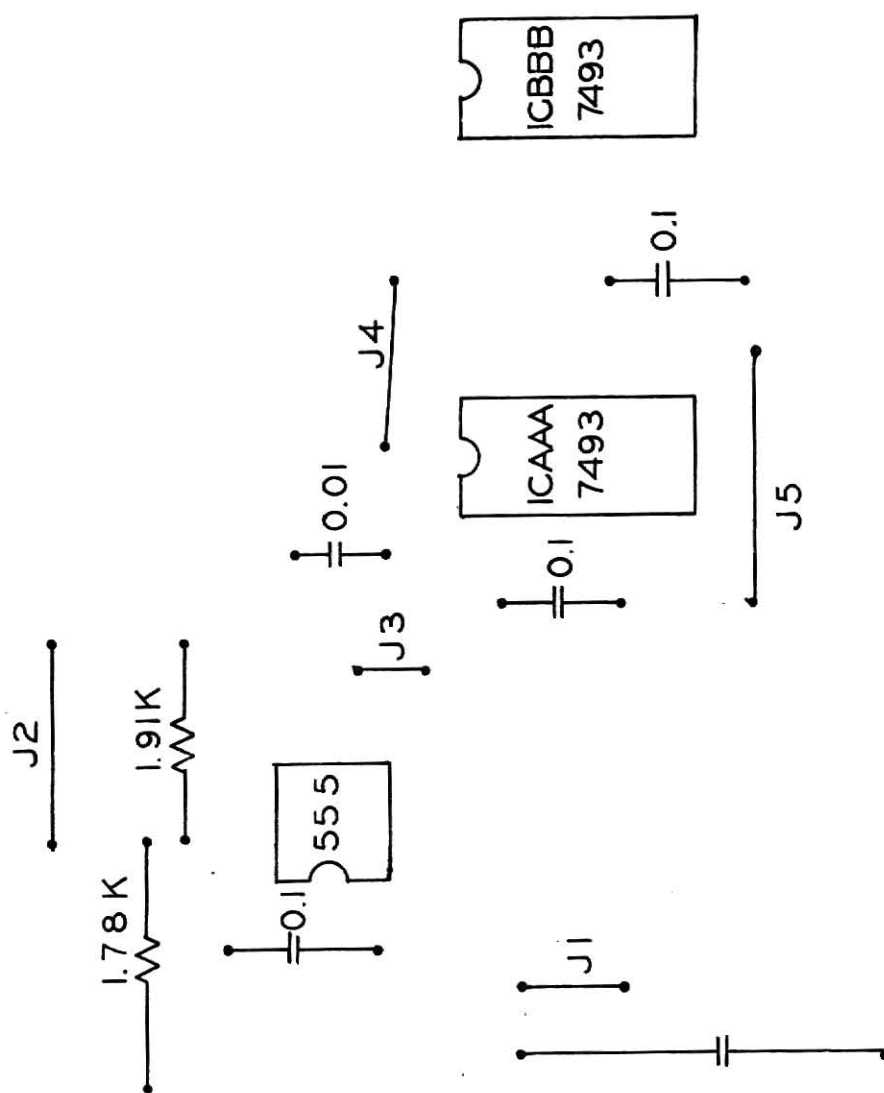


BOARD F



BOARD 1AQ-1

Z  
22A  
1

Z  
22

BOARD IAQ-2

A -

APPENDIX C  
SOURCE LISTING OF RDCP



## RDCT

;RADAR DATA COLLECTION PROGRAM  
 ;BY CARL C. ANDREASEN  
 ;JANUARY 20 1976

;RDCT IS THE PROGRAM USED TO CONTROL THE  
 ;DIGITAL COMPUTER INTERFACE OF THE DIGITAL  
 ;RADAR PROCESSOR.  
 ;REFERENCE (2) PRESENTS THE STYLE OF PROGRAMMING  
 ;USED IN THIS PROGRAM.

;CREATE POINTER TABLE ENTRY  
   .ENT DSKR,DELET,RDMS,WORKF  
   .ENT NEXTA,DSKW,CATAL,FETCH  
   .EXTD PT1,TXTOT  
   .ZREL  
 ;PROCEDURE TO BUILD ALL NECESSARY LINKAGES TO DMS.

;DISPLACEMENT VALUES FOR RESIDENT DMS REFERENCES

DIVAL:  .41     ;POINTER TO DISPLACEMENT VALUE TABLE  
         -351    ;WORKFILE PARAMETER TABLE  
         -325    ;NEXT AVAILABLE SECTOR ADDRESS IN DMS LIBRARY  
         -324    ;RETURN ADDRESS TO DMS  
         -320    ;FETCH RELOAD SECTOR ADDRESS  
         -317    ;FETCH ROUTINE  
         -340    ;DELETE ROUTINE  
         -235    ;CATALOG ROUTINE  
         -130    ;DISK READ ROUTINE  
         -126    ;DISK WRITE ROUTINE  
         -56     ;TELETYPE PRINT ROUTINE  
 C40:    40     ;END OF DIVAL TABLE FLAG &  
               ;SECTOR ADDRESS OF EXECUTIVE ROUTINE  
 ;RESIDENT DMS POINTERS

WORKF:  0       ;WORKFILE PARAMETER TABLE  
 NEXTA:  0       ;NEXT AVAILABLE SECTOR ADDRESS IN DMS LIBRARY  
 PA3EX:  0       ;RETURN ADDRESS TO DMS  
 LASTP:  0       ;FETCH RELOAD SECTOR ADDRESS  
 FETCH:  0       ;FETCH ROUTINE  
 DELET:  0       ;DELETE ROUTINE  
 CATAL:  0       ;CATALOG ROUTINE  
 DSKR:   0       ;DISK READ ROUTINE  
 DSKW:   0       ;DISK WRITE ROUTINE  
 TTOUT:  0       ;TELETYPE PRINT ROUTINE

RDMS:   57021   ;RETURN ADDRESS TO DMS  
          .NREL

;POINTER TABLE PACKING DISPLACEMENT

PTPD=WORKF-DIVAL-1

;THIS ROUTINE BUILDS ALL DMS POINTERS

START:  LDA 2,DIVAL           ;GET POINTER TO DISPLACEMENT

```

;VALUE TABLE TABLE
LDA 0,0,2      ;GET A DISPLACEMENT VALUE
MCVL# 0,0,SNC  ;IF VALUE IS POSITIVE, MUST BE END
                ;OF TABLE
JMP .+5        ;IT WAS POSITIVE, SO POINTER GENERATION
                ;IS DONE
ADD 3,0        ;WAS NOT POSITIVE, SO BUILD A POINTER
STA 0,PTPD,0   ;PACK POINTER IN DMS POINTER TABLE
INC 2,2        ;UPDATE DISPLACEMENT POINTER
JMP START+1    ;GET NEXT DISPLACEMENT

;DMS POINTER TABLE HAS BEEN BUILT, NOW SET FETCH RELOAD
;SECTOR TO 40.

STA 00,LASTP
JMP 0PT1      ;GO TO PART 1 OF THE PROGRAM

.END

```

```

; CREATE POINTER TABLE ENTRY
      .ENT PT1, TABLE, SAMNU, STMEH, STMEL
      .EXTD DATE, TIME, SAMP, DATA, STIME, RDMS, FETCH, TXTOT, CRLF
      .EXTD WORKF, DSKF, NEXTA, CATAL, HOURS, MIN, SEC
      .ZREL
PT1:   MIL1
SAMNU: 0      ; NUMBER OF DATA ACQUISITION RECORDS
STMEH: 0      ; STORAGE OF HIGH ORDER WORD OF SAMPLE TIME
STMEL: 0      ; STORAGE OF LOW ORDER WORD OF SAMPLE TIME
SECTC: 0      ; STORAGE FOR SECTOR COUNT
      .LOC 20
TABLE: 5000   ; ADDRESS OF DATA ACQUISITION RECORD
      .NREL
; INITIALIZE THE POINTER TO THE DATA ACQUISITION RECORD
; DATA BUFFER (DARDB)

MIL1:  LDA 0, MID1    ; LOAD ADDRESS OF DARDB
      STA 0, TABLE  ; STORE ADDRESS OF DARDB IN
                      ; AUTOINCREMENTING LOCATION "TABLE"
      DSZ TABLE     ; DECREMENT DARDB ADDRESS BY 1
                      ; SINCE AUTOINCREMENTING LOCATION ADDS
                      ; BEFORE STORING DATA.

; OBTAIN THE DATE AND STORE THE DATE DARDB
      JSR @DATE
; OBTAIN THE TIME AND STORE THE TIME IN DARDB
      JSR @TIME
; OBTAIN THE NUMBER OF DATA ACQUISITION RECORDS TO BE RECORDED
; THE RESULT BEING STORED IN LOW CORE LOCATION "SAMNU"
      JSR @SAMP
; OBTAIN THE TIME INTERVAL BETWEEN DATA ACQUISITION SEQUENCES.
      JSR @STIME
; ESTABLISH A FILE NAME OF THE FORM RDMN.DA
; WHERE MN = 00 ... 99
      LDA 3, WORKF    ; LOAD ADDRESS OF WFPT
      LDA 1, MID2     ; OBTAIN FIRST 2 CHARACTERS OF
                      ; FILE NAME
      STA 1, 0, 3     ; STORE IN WFPT
      LDA 1, MID3     ; LOAD SECOND 2 CHAR. OF FILE NAME
      STA 1, 1, 3     ; STORE THEM IN WFPT
MIL3:  SUBZR 0, 0      ; SET BIT 2=1
      STA 0, 2, 3     ; INITIALIZE WORDS 2-6 OF WFPT
      STA 0, 3, 3
      STA 0, 4, 3
      STA 0, 5, 3
      STA 0, 6, 3
      LDA 1, MID4     ; LOAD DISK FILE ID
      STA 1, 7, 3     ; STORE DISK FILE ID
; HAVE ESTABLISHED WFPT AT DEFAULT MODE
; NOW CHECK FOR THE NEXT RD##.DA
; WHERE # CAN BE AN INTEGER FROM 0-9
      JSR @FETCH      ; FOR THE FILE IN THE WFPT, SEE IF
                      ; IS ON THE DISK
      JMP MIL5        ; IT WAS NOT
      JMP MIL5        ; IT WAS NOT
      LDA 1, MID3     ; LOAD SECOND 2 CHARACTERS OF
                      ; FILE NAME

```

```

MOVSV      1,1      ;JUSTIFIED RIGHT-LEFT, NOW ARE
                ;LEFT-RIGHT
LDA         0,MID5   ;LOAD MASK TO KEEP RIGHT HAND SIDE
AND         0,1      ;KEEP JUST RHS
LDA         0,MID6   ;LOAD MASK FOR '9'
SUB#        0,1,SNR  ;TEST IF CHAR=MASK
JMP         MIL4     ;IS A 9 SO CHECK OTHER SIDE
LDA         1,MID3   ;LOAD SECOND 2 CHARACTERS OF FILE NAME
MOVSV      1,1      ;WAS NOT 9 SO INCREMENT
INC         1,1      ;BY 1 AND STORE THEM IN
MOVSV      1,1
LDA         3,WORKF  ;LOAD ADDRESS OF WFPT
STA         1,1,3    ;STORE SECOND 2 CHAR IN WFPT
STA         1,MID3   ;UPDATE THE SECOND 2 CHARACTERS
JMP         MIL3     ;GO CHECK IF THIS ONE EXISTS
MIL4: LDA      1,MID3   ;LOAD THE SECOND 2 CHARACTERS
LDA         0,MID5   ;LOAD MASK TO KEEP RHS
AND         0,1      ;KEEP JUST RHS
LDA         0,MID6   ;LOAD MASK OF '9'
SUB#        0,1,SR   ;TEST IF MASK=CHAR
JMP         .45      ;WAS NOT A 9
JSR         @CRLF
JSR         @TXTOT   ;WAS A 9 SO RD99.DA EXISTS
MIT1
JSR         @RDMS
LDA         1,MID3   ;LOAD THE SECOND 2 CHARACTERS
                ;OF FILE NAME
INC         1,1      ;SINCE WAS NOT 9, INCREMENT IT
LDA         0,MID7   ;LOAD MASK TO RESET CHAR '3' IN LHS
AND         0,1      ;HAVE NEXT SECOND 2 CHARACTERS
                ;OF FILE NAME
LDA         3,WORKF  ;LOAD ADDRESS OF WFPT
STA         1,1,3    ;STORE NEW SECOND 2 CHARACTERS
STA         1,MID3   ;UPDATE SECOND 2 CHARACTERS
JMP         MIL3     ;GO SEE IF THE NEW FILE NAME EXISTS
.ZREL
MID1:      5000      ;ADDRESS OF DANDB
MID2:      042122    ;CHARACTERS 'DR'
MID3:      030060    ;CHARACTERS '30'
MID4:      040504    ;CHARACTERS 'AD'
MID5:      377       ;MASK TO KEEP RHS
MID6:      071       ;CHARACTER '9'
MID7:      030377    ;CHARACTER '0' MASK FOR LHS
MID9:      0
MID10:     -60.      ;RTC COUNTER
MID15:     0         ;TEMPORARY LOCATION
MID16:     0         ;TEMPORARY LOCATION
MID12:     MIL10     ;ADDRESS OF INTERRUPT HANDLER
MID13:     177773    ;INTERRUPT MASK
MID14:     00014     ;RTC DEVICE CODE
MID17:     5         ;FUDGE FACTOR TIME CONSTANT
MID18:     -64.      ;# OF 5.625 DEGREE INCREMENTS IN 360 DEGREES
MID19:     -50.      ;#OF AZIMUTH WORDS TO BE TRANSFERED
MIA1:      MIL3
MIT1:      .TXT      &ALL FILES FROM RD00.DA TO RD99.DA EXIST&
MIT2:      .TXT      &BEGIN SAMPLING&
MIT3:      .TXT      &END SAMPLING&
.NREL

```

; READ THE DATA FROM THE INTERFACE

```
MIL5:  JSR    @CRLF
        JSR    @TXTOT
        MIT2
        JSR    @CRLF
        JSR    @DATAR
        JSR    @CRLF
        JSR    @TXTOT
        MIT3
```

; OUTPUT THE DARDB

; "TABLE" POINTS TO END OF DARDB

; "MIL1" POINTS TO BEGINNING OF DARDB

; STORE THE DATA

```
MIL6:  SUB     2,2          ;CLEAR AC2 FOR SECTOR COUNT
        STA     2,SECTC    ;CLEAR SECTOR COUNT LOCATION
        LDA     1,MID1     ;MOVE BEGINNING POINTER TO USE LOCATI
```

```
        STA     1,MID9
        LDA     3,WORKF    ;LOAD ADDRESS OF WFPT
        LDA     0,@NEXTA   ;LOAD NEXT SECTOR ADDRESS
        STA     0,3,3      ;STORE NEXT SECTOR ADDRESS IN WFPT
;                                     AS IT IS STARTING SECTOR FOR THIS FILE
```

```
        LDA     0,@NEXTA   ;LOAD NEXT SECTOR ADDRESS
        LDA     1,MID9     ;LOAD START ADDRESS OF DATA
        JSR     @DSK1      ;WRITE THE DATA
        ISZ     SECTC      ;INCREMENT THE SECTOR COUNT
        STA     0,@NEXTA   ;STORE NEXT SECTOR ADDRESS
        STA     1,MID9     ;STORE START ADDRESS OF DATA + 400
        LDA     0,TABLE    ;LOAD MAX CORE USED
        SUBZ#   1,0,SZC    ;SKIP WHEN ALL WRITTEN
        JMP     .-10
```

```
        LDA     3,WORKF    ;LOAD ADDRESS OF WFPT
        LDA     2,SECTC    ;LOAD THE SECTOR COUNT
        STA     2,4,3      ;STORE SECTOR COUNT
        SUB     2,2
        STA     2,5,3      ;CLEAR WORD 5 OF WFPT
        JSR     @CATAL     ;CATALOG THIS FILE
```

; TEST IF MORE DATA ACQUISITION SEQUENCES TO BE RECORDED

```
MIL7:  DSZ     SAMNU       ;DECREMENT # OF DATA ACQUISITION
;                                     ; SEQUENCES
```

; TO BE RECORDED.

```
        JMP     .+2        ;MORE YET TO GO
        JMP     @EDMS      ;RETURN TO OPERATING SYSTEM
```

; WAIT DESIRED INTERVAL

```
MIL8:  SUB     2,2          ;CLEAR AC2 FOR A COUNTER
        DOA     2,RTC       ;SET RTC FOR 60 HZ OPERATION
        NIOS    RTC
        LDA     1,MID12     ;LOAD AND STORE
        STA     1,1        ;ADDRESS OF INTERRUPT HANDLER
        LDA     0,STMEH
        NEG     0,0        ;FORMS 2'S COMPLEMENT OF INTERVAL TIME
        STA     0,MID15
        LDA     0,STMEL
        NEG     0,0        ;FORMS 2'S COMPLEMENT OF INTERVAL TIME
        STA     0,MID16
        LDA     1,MID13     ;LOAD INTERRUPT MASK
        DOB     1,CPU       ;EXECUTE MSKC
MIL11: LDA     0,MID10     ;LOAD -60 DEC. AS # 60 HZ
```

```

MIL9:   INTEN      ;INTERRUPTS TO COUNT
        JMP        .+1      ;ENABLE INTERRUPTS
        JMP        .-1      ;WAIT
MIL10:  INTA       2        ;INPUT DEVICE CODE OF
                           ;INTERRUPTING DEVICE
        LDA        1,MID14   ;LOAD DEVICE MASK
        NIOS       RTC      ;WAKE UP THE RTC AGAIN
        SUB#       2,1,SZR   ;TEST IF DEVICE CODE 13
        JSR        @RDMS
        INC        0,0,SZR   ;INCREMENT COUNTER
        JMP        MIL9      ;NOT ENOUGH
        LDA        2,MID16   ;LOAD LOW ORDER OF INTERVAL
                           ;TIME IN SECONDS
        LDA        1,MID15   ;LOAD HIGH ORDER OF INTERVAL
                           ;TIME IN SECONDS
        MOV        2,2,SNR   ;TEST IF LOW ORDER IS ZERO
        JMP        MIL12     ;YES, THEN PROCESS HIGH ORDER
        INC        2,2      ;NO, INCREMENT ONE SECOND
        STA        2,MID16   ;STORE NEW INTERVAL TIME COUNT
        JMP        MIL11     ;INTERRUPT FOR NEXT SECOND
MIL12:  MOV        1,1,SNR   ;TEST IF HIGH ORDER IS ZERO
        JMP        MIL13     ;YES, THEN BOTH HIGH AND
                           ;LOW ORDER INTERVAL TIME
        ;              COUNT IS ZERO, SO INTERVAL TIME
        ;              IS UP. WRITE NEXT DARD
        INC        1,1      ;NO, INCREMENT ONE SECOND
        STA        1,MID15   ;STORE NEW HIGH ORDER TIME COUNT
        JMP        MIL11     ;CONTINUE FOR NEXT SECOND
MIL13:  SUB        1,1      ;CLEAR AC!
        DOB        1,CPU     ;EXECUTE MSKO INSTRUCTION
        LDA        3,MID1   ;RESET TABLE POINTER
        STA        3,TABLE
        DSZ        TABLE   ;DECREMENT ADDRESS BY 1
        LDA        0,3,3     ;LOAD HOURS
        LDA        1,HOURS   ;LOAD INTERVAL TIME HOURS
        ADD        0,1
        STA        1,3,3     ;STORE TIMEHOURS + INTERVAL HOURS
        LDA        0,4,3     ;LOAD MINUTES
        LDA        1,MIN     ;LOAD INTERVAL TIME MINUTES
        ADD        0,1
        STA        1,4,3     ;STORE TIME MINUTES + INTERVAL MINUTES
        LDA        0,5,3     ;LOAD SECONDS
        LDA        1,SEC     ;LOAD INTERVAL SECONDS
        ADD        0,1
        LDA        0,MID17   ;LOAD INTERVAL SECONDS
        ADD        0,1
        STA        1,5,3     ;STORE TIME SECONDS + INTERVAL SECONDS
        JMP        @M1A1     ;CONTINUE FOR REST OF DESIRED
                           ;# OF RECORDS
        ;              DATA FROM DATA ACQUISITION SEQUENCE
        .END

```

```

      TXTOT
;THIS SUBROUTINE PRINTS A TEXT TO THE TELETYPE
;
;CALLING SEQUENCE:
;      JSR      @TXTOT
;      RETURN
;
;INPUTS:AC3      POINTS TO ADDRESS OF THE TEXT
;
;OUTPUTS:        NONE
;
;CREATE ENTRY TO POINTER TABLE
      .ENT      TXTOT
      .ZREL
TXTOT:  S1L1
      .NREL
;
S1L1:   STA      3,S1D0
      NIOS      TTO
      STA      2,S1D0+1
      STA      1,S1D0+2
      STA      0,S1D0+3
      LDA      1,S1D1      ;LOAD THE MASK
S1L2:   LDA      3,0,3      ;LOAD ADDRESS OF WORD OF TEXT
S1L4:   LDA      0,0,3      ;LOAD A WORD OF TEXT
      AND#      0,1,SNR     ;TEST IF BITS 0-7 ARE 0
      JMP      S1L3
      SKPBZ     TTO
      JMP      .-1
      DOAS      0,TTO
      MOVS      0,0      ;GET THE SECOND HALF
      AND#      0,1,SNR     ;TEST IF BITS 0-7 ARE 0
      JMP      S1L3
      SKPBZ     TTO
      JMP      .-1
      DOAS      0,TTO
      INC      3,3
      JMP      S1L4
S1L3:   LDA      0,S1D0+3
      LDA      1,S1D0+2
      LDA      2,S1D0+1
      ISZ      S1D0      ;INCREMENT PAST POINTER
      LDA      3,S1D0
      JMP      @S1D0
S1D0:   .BLK      4
S1D1:   00377
      .END

```

```

      CHOUT
; THIS SUBROUTINE OUTPUTS A RIGHT JUSTIFIED CHARACTER FROM ACC
;
; CALLING SEQUENCE:
;       JSR     @CHOUT
;       RETURN  HERE
;
; INPUTS: ACC0      CHARACTER IN BITS 7-15 TO GO OUT
;
; OUTPUTS:          NONE
;
; CREATE ENTRY TO POINTER TABLE
      .ENT      CHOUT
      .ZREL
CHOUT: S2L1
      .NREL
;
S2L1:  STA      3,S2D0
      STA      2,S2D0+1
      STA      1,S2D0+2
      STA      0,S2D0+3
      SKPBZ    T0      ; CHECK IF TELETYPE IS PRINTING
      JMP      -1
      DOAS     0,T0      ; OUTPUT CHARACTER
      LDA      0,S2D0+3
      LDA      1,S2D0+2
      LDA      2,S2D0+1
      LDA      3,S2D0
      JMP      @S2D0
; DATA
S2D0:  .BLK     4
      .END

```



```

      CREAD
;
; CALLING SEQUENCE:
;      JSR      @CREAD
;      RETURN HERE
;
; INPUTS: NONE
;
; OUTPUTS:      AC0      RESULT OF READ IN BITS 7-15
;
; CREATE ENTRY TO POINTER TABLE
      .ENT      CREAD
      .ZREL
CREAD:  S3L1
      .NREL
;
S3L1:  STA      3,S3D0      ;STORE ENVIRONMENT
      STA      2,S3D0+1
      STA      1,S3D0+2
      NIOS      TTI      ;START THE TELETYPE INPUT
      SKPDN     TTI      ;WAIT FOR A KEYSTROKE
      JMP      .-1
      DIAS      0,TTI      ;BRING IN A CHARACTER FROM TTI
      NI OC     TTI
      LDA      1,S3D1      ;LOAD MASK TO ELIMINATE PARITY BIT
      AND      1,0      ;MASK OFF PARITY
      LDA      2,S3D0+1
      LDA      1,S3D0+2
      JMP      @S3D0      ;AND RETURN
S3D0:  .BLK      3
S3D1:  177
      .END

```

```

      CRLF
; THIS SUBROUTINE OUTPUTS A CR AND LF
;
; CALLING SEQUENCE:
;     JSR     @CRLF
;     RETURN
;
; INPUTS: NONE
;
; OUTPUTS:      NONE
;
; SUPPORT SUBROUTINES
;     DELAY
;
; CREATE POINTER TABLE ENTRY
      .ENT     CRLF
      .EXTD    DELAY
      .ZREL
CRLF:  S3L1
      .NREL
;
S3L1:  STA     3,S3D0
      STA     2,S3D0+1
      STA     1,S3D0+2
      STA     0,S3D0+3
      LDA     1,S3D1      ; LOAD CR AND LF CHARACTERS
      SKPBZ   TTO         ; TEST IF OUTPUT BEING DONE
      JMP     .-1
      DOAS    1,TTO        ; OUTPUT CR
      JSR     @DELAY      ; GIVE A DELAY
      MOVS    1,1
      SKPBZ   TTO
      JMP     .-1
      DOAS    1,TTO        ; OUTPUT THE LF
      LDA     0,S3D0+3
      LDA     1,S3D0+2
      LDA     2,S3D0+1
      LDA     3,S3D0
      JMP     @S3D0
; DATA
S3D0:  .BLK    4
S3D1:  .TXT    '<15><12>'
      .END

```

```

      DELAY
;THIS SUBROUTINE GIVES A SHORT DELAY
;SO THAT THE CARRIAGE CAN RETURN WHEN USING THE TTY
;
;CALLING SEQUENCE
;      JSR      @DELAY
;      RETURN
;
;INPUTS:NONE
;
;OUTPUTS:      NONE
;
;CREATE POINTER TABLE ENTRY
      .ENT      DELAY
      .ZREL
DELAY:  S9L1
      .NREL
;
S9L1:   STA      3,S9D0
        STA      2,S9D0+1
        LDA      2,S9D1
        STA      2,S9D2
        DSZ      S9D2
        JMP      .-1
        LDA      2,S9D0+1
        JMP      @S9D0
S9D0:   .BLK     2
S9D1:   500
S9D2:   0
        .END

```

```

                BELL
;      OUTPUTS THE BELL CHARACTER TO THE TTO
;
;CALLING SEQUENCE
;      JSR      @BELL
;      RETURN
;
;INPUTS:NONE
;
;OUTPUTS:      NONE
;
;SUPPORT SUBROUTINES
;      DELAY
;
;CREATE POINTER TABLE ENTRY
      .ENT      BELL
      .EXTD     DELAY
      .ZREL
;
BELL:      SEL1
      .NREL
SEL1:      STA      3,SED0
           STA      2,SED0+1
           STA      1,SED0+2
           STA      0,SED0+3
           LDA      0,SED1
           DOAS     0,TTO
           JSR      @DELAY
           JSR      @DELAY
           JSR      @DELAY
           SKPBZ    TTO
           JMP      .-1
           LDA      0,SED0+3
           LDA      1,SED0+2
           LDA      2,SED0+1
           JMP      @SED0
SED0:      .BLK     4
SED1:      007
           .END

```

```

DATE
;
;THIS SUBROUTINE OBTAINS THE DATE IN THE FORMAT
;    MONTH/DAY/YEAR
;WHERE EACH FIELD IS A 2 DIGIT INTEGER VALUE
;AND IS STORED AS A BINARY VALUE.
;AN AUTOINCREMENTING LOCATION "TABLE" MUST
;EXIST. THE DATE WILL BE STORED INDIRECTLY
;THROUGH "TABLE"
;
;CALLING SEQUENCE
;    JSR    @DATE
;    RETRN
;
;SUPPORT SUBROUTINES
;    TXTOT
;    CRLF
;    GETNUM
;    CHOUT
;    ERROR
;
;CREATE POINTER TABLE ENTRY
    .ENT    DATE
    .EXTD   TXTOT,CRLF,ERROR,GETNUM,TABLE,CHOUT
    .ZREL
DATE:    SIL1
        .NREL
;
SIL1:    STA    3,SID0    ;STORE ENVIRONMENT
        STA    2,SID0+1
        STA    1,SID0+2
        STA    0,SID0+3
SIL2:    JSR    @CRLF
        JSR    @TXTOT    ;PRINT QUESTION
        SIT1
        LDA    0,SID2    ;LOAD # OF DIGITS
        LDA    1,SID4    ;LOAD MAXIMUM VALUE FOR MONTH
        JSR    @GETNUM    ;OBTAIN MONTH
        JMP    SIL3    ;ERROR,RETURN HERE
        STA    0,ETABLE    ;STORE MONTH
        LDA    0,SID7    ;LOAD CHAACTER "/"
        JSR    @CHOUT
        LDA    0,SID2    ;LOAD # OF DIGITS
        LDA    1,SID6    ;LOAD MAX FOR DAY
        JSR    @GETNUM    ;OBTAIN DAY
        JMP    SIL4    ;ERROR,RETURN HERE
        STA    0,ETABLE    ;STORE DAY
        LDA    0,SID7
        JSR    @CHOUT
        LDA    0,SID2    ;LOAD # OF DIGITS
        LDA    1,SID5    ;LOAD MAXIMUM YEAR VALUE
        JSR    @GETNUM    ;OBTAIN YEAR
        JMP    SIL5    ;ERROR,RETURN HERE
        STA    0,ETABLE    ;STORE YEAR
        LDA    0,SID0+3    ;RESTORE ENVIRONMENT
        LDA    1,SID0+2
        LDA    2,SID0+1
        LDA    3,SID0

```

```

      JMP      @SID0
SIL5:  DSZ     TABLE
SIL4:  DSZ     TABLE
SIL3:  JSR     @ERROR
      JMP      SIL2

;DATA
SID0:  .BLK    4           ;ENVIRONMENT STORAGE
SID1:  0
SID2:  2       ;MAX # OF DIGITS
SID3:  0
SID4:  12.     ;MAXIMUM MONTH VALUE
SID5:  99.     ;MAXIMUM YEAR VALUE
SID6:  31.     ;MAXIMUM DAY VALUE
SID7:  057     ;CHARACTER "/"
SIT1:  .TXT    &DATE (MONTH/DAY/YEAR) :  &
      .END

```

## GETNUM

```

;
; THIS SUBROUTINE OBTAINS A BCD NUMBER
; FROM THE TTY. VALUES ARE POSITIVE AND
; UP TO 4 BCD DIGITS IN LENGTH (MAX=9999)
; THEN THE BCD VALUE IS CONVERTED TO BINARY AND RETURNED
; IN AC0
;
; INPUTS: AC0      NUMBER OF DIGITS
;          AC1      MAXIMUM VALUE
;          MINIMUM VALUE IS 0
;
; OUTPUTS: AC0      (A BINARY NUMBER EQUAL TO BCD VALUE
;                   READ IN)
;
; CALLING SEQUENCE
;      JSR  @GETNUM
;      RETURN IF ERROR
;      NORMAL RETURN
;
; CREATE POINTER TABLE ENTRY
;      .ENT  GETNUM
;      .EXTD CREAD,CHOUT
;      .ZREL
GETNUM: SJL1
;      .NREL
;
;
SJL1:  STA  3,SJD0      ; STORE ENVIRONMENT
      STA  2,SJD0+1
      STA  1,SJD0+2
      STA  0,SJD0+3
      STA  0,SJD1      ; STORE DIGIT # COUNT
      SUBZ 2,2        ; CLEAR AC2
SJL6:  JSR  @CREAD      ; GET A CHARACTER
      JSR  @CHOUT      ; ECHO IT
      LDA  1,SJD3      ; GET MAX CHAR LIMIT
      SUBZ# 1,0,SZC    ; SKIP IF MAX>WHAT READ IN
      JMP  SJL3
      LDA  1,SJD2      ; GET MINIMUM CHAR LIMIT
      SUBZ# 1,0,SNC    ; SKIP IF MIN <OR= WHAT READ IN
      JMP  SJL3        ; INCORRECT INPUT
      LDA  1,SJD7      ; GET CHARACTER MASK
      AND  1,0        ; OBTAIN OCTAL EQUIVALENT
                        ; OF THE CHARACTER
      DSZ  SJD1        ; DECREMENT # OF DIGITS
      JMP  SJL2        ; SHIFT LEFT 4 BITS FOR NEXT INPUT
      ADD  0,2        ; WAS LAST DIGIT, SO ADD IT TO SUM
      STA  2,SJD6      ; STORE RESULT IN TEMPORARY LOCATION
      LDA  0,SJD4      ; LOAD COUNTER (# BCD DIGITS MAX)
      STA  0,SJD1      ; STORE IN COUNTER LOCATION
      SUB  2,2        ; CLEAR AC2 FOR RESULT REGISTER
SJL4:  JSR  SJL5        ; GET DIGIT
      JSR  SJL7        ; MULTIPLY SUM BY 10 AND ADD DIGIT
      ISZ  SJD1        ; DONE YET?
      JMP  SJL4        ; NO
      MOVZ 2,0        ; RESULT TO ACC0, CLEAR CARRY
      JMP  SJL8

```

```

SJL5:  LDA    1,SJD4    ;LOAD LOOP COUNTER
      STA    1,SJD5    ;STORE IT FOR USE
      LDA    1,SJD6    ;LOAD THE VALUE TO BE CONVERTED
;      TO BINARY
      SUB    0,0        ;CLEAR AC0
      MOVZL  1,1        ;SHIFT OUT A BIT
      MOVL   0,0        ;MOVE THE BIT TO AC0
      ISZ    SJD5       ;INCREMENT LOOP COUNTER AND
      JMP    -3         ;LOOP TILL A DIGIT SHIFTED TO AC0
      STA    1,SJD6    ;STORE CURRENT WORD
      JMP    0,3        ;RETURN
SJL7:  MOVZL  2,1        ;N*2
      MOVZL  1,1        ;N*4
      ADD    2,1        ;N*5
      MOVZL  1,2        ;N*12
      ADD    0,2        ;N*12+AC0
      JMP    0,3        ;RETURN
SJL2:  MOVZL  0,0        ;SHIFT LEFT 4 PLACES
      MOVZL  0,0
      MOVZL  0,0
      MOVZL  0,0
      ADD    0,2        ;FORM RESULT IN AC2
      JMP    SJL6       ;CONTINUE
SJL8:  LDA    1,SJD0+2   ;TEST RESULT AGAINST MAXIMUM
      SUBZ#   0,1,SZC    ;SKIP IF MAX<VALUE
      ISZ    SJD0       ;INCREMENT RETURN ADDRESS,
SJL3:  LDA    3,SJD0
      LDA    2,SJD0+1
      LDA    1,SJD0+2
      JMP    0,SJD0     ;RETURN
SJD0:  .BLK 4
SJD1:  0                ;LOOP COUNTER OF # OF DIGITS READ IN
SJD2:  060              ;MINIMUM CHARACTER
SJD3:  071              ;MAXIMUM CHARACTER
SJD4:  -4               ;COUNTER
SJD5:  0                ;TEMPORARY COUNTER
SJD6:  0                ;TEMPORARY LOCATION
SJD7:  177              ;CHARACTER MASK FOR ASCII TO OCTAL
      .END

```



```

      TIME
;THIS SUBROUTINE OBTAINS THE TIME IN THE FORMAT:
;    HOURS/MINUTES/SECONDS
;EACH FIELD MUST BE A TWO DIGIT INTEGER VALUE AND
;IS STORED AS A BINARY VALUE.
;AN AUTOINCRENTING LOCATION LABELED "TABLE"
;MUST EXST. THE TIME IS STORED INDIRECTLY
;THROUGH THIS LOCATION. "TABLE" MUST CONTAIN THE
;ADDRESS OF WHERE THE TIME IS TO BE STORED.
;
;INPUTS:NONE
;
;OUTPUTS:      NONE
;
;CALLING SEQUENCE
;    JSR    @TIME
;    RETURN
;
;SUPPORT SUBROUTINES
;    GETNUM
;    ERROR
;    TXTOT
;    CRLF
;
;CREATE POINTER TABLE ENTRY
      .ENT    TIME
      .EXTD   CRLF,TXTOT,GETNUM,ERROR,TABLE
      .ZREL
TIME:  SLL1
      .NREL
;
SLL1:  STA     3,SLD0      ;STORE THE ENVIRONMENT
      STA     2,SLD0+1
      STA     1,SLD0+2
      STA     0,SLD0+3
SLL2:  JSR     @CRLF
      JSR     @TXTOT      ;ASK FOR TIME INPUT
      SLT1
      LDA     0,SLD1      ;LOAD #      OF BCD DIGITS TO BE READ I
;
      LDA     1,SLD2      ;LOAD MAX OF HOURS
      JSR     @GETNUM      ;OBTAIN THE HOURS
      JMP     SLL3      ;RETURN HERE IF ERROR
      STA     0,@TABLE    ;STORE HOURS
      JSR     @TXTOT
      SLT2
      LDA     0,SLD1      ;LOAD # OF BCD DIGITS TO BE READ IN
      LDA     1,SLD3      ;LOAD MAX OF MINUTES
      JSR     @GETNUM      ;OBTAIN THE MINUTES
      JMP     SLL5      ;RETURN HERE IF ERROR
      STA     0,@TABLE    ;STORE MINUTES
      JSR     @TXTOT
      SLT2
      LDA     0,SLD1      ;LOAD # OF BCD DIGITS TO BE READ IN
      LDA     1,SLD3      ;LOAD MAX OF SECONDS
      JSR     @GETNUM      ;OBTAIN THE SECONDS

```

```

        JMP      SLL4      ;RETURN HERE IF ERROR
        STA      0, @TABLE ;STORE SECONDS
        LDA      3, SLD0   ;RESTORE ENVIRONMENT
        LDA      2, SLD0+1
        LDA      1, SLD0+2
        LDA      0, SLD0+3
        JMP      @SLD0     ;RETURN
SLL4:   DSZ      TABLE
SLL5:   DSZ      TABLE
SLL3:   JSR      @ERROR
        JMP      SLL2

; DATA
SLD0:   .BLK     4          ;ENVIRONMENT STORAGE
SLD1:   2          ;NO OF BCD CHARACTER TO BE OBTAINED
SLD2:   24.        ;MAX HOURS
SLD3:   59.        ;MAX SECONDS, MINUTES
SLT1:   .TXT      &TIME (HOURS/MINUTES/SECONDS) :&
SLT2:   .TXT      */*
        .END

```

```

      STIME
;THIS SUBROUTINE OBTAINS THE TIME BETWEEN
;DATA ACQUISITION SEQUENCES DESIRED BY THE
;USER.
;THE TIME IS A DOUBLE PRECISION
;BINARY VALUE. THE MODULO 60 OF THE HR,MIN,SEC ENTERED
;IS FOUND IN LOCATIONS STMEH,STMEL.
;
;INPUTS:NONE
;
;OUTPUTS:      NONE
;
;CALLING SEQUENCE
;      JSR      @STIME
;      RETURN
;
;
;SUPPORT SUBROUTINES
;      TXTOT
;      TIME
;      MULT,MULTA
;      CRLF
;AUTOINCREMENTING LOCATION "TABLE"
;AND LOW CORE LOCATIONS STMEH,STMEL MUST
;BE PROVIDED IN THE POINTER TABLE.
;
;CREATE POINTER TABLE ENTRY
      .ENT      STIME,HOURS,MIN,SEC
      .EXTD     CRLF,TXTOT,TIME,MULT,MULTA
      .EXTD     TABLE,STMEH,STMEL
      .ZREL
STIME:  SML1      ;SUBROUTINE TO READ TIME INTERVAL
HOURS:  0         ;STORAGE FOR INTERVAL HOURS
MIN:    0         ;STORAGE FOR INTERVAL MINUTES
SEC:    0         ;STORAGE FOR INTERVAL SECONDS
      .NREL
;
SML1:   STA      3,SMD0      ;STORE ENVIRONMENT
        STA      2,SMD0+1
        STA      1,SMD0+2
        STA      0,SMD0+3
        LDA      0,TABLE    ;CHANGE TABLE TO
        STA      0,SMD1      ;          POINT
        LDA      0,SMD2      ;          TO
        STA      0,TABLE    ;LOCATIONS IN THIS SUBROUTINE
        JSR      @CRLF
        JSR      @TXTOT      ;OUTPUT MESSAGE
        SMT1
        JSR      @TIME      ;OBTAIN THE TIME
        LDA      0,SMD3+2    ;LOAD THE SECONDS
        STA      0,SEC
        LDA      1,SMD3+1    ;LOAD THE MINUTES
        STA      1,MIN
        LDA      2,SMD5      ;LOAD MIN MULTIPLIER OF 60
        JSR      @MULTA      ;FORM SEC + (MIN*60)
        MOV      1,0        ;MOVE LOW ORDER TO AC2
        LDA      1,SMD3      ;LOAD HOURS
        STA      1,HOURS

```

```

LDA    2,SMD4    ;LOAD HOURS MULTIPLIER OF 3600
JSR    @MULTA    ;FORM (SEC + (MIN*60)) + (HRS*3600)
STA    0,STMEH   ;STORE HIGH ORDER OF RESULT
STA    1,STMEL   ;STORE LOW ORDER OF RESULT
LDA    0,SMD1    ;RESTORE ORIGINAL POINTER IN
STA    0,TABLE   ;LOCATION "TABLE"
LDA    3,SMD0    ;RESTORE THE ENVIRONMENT
LDA    2,SMD0+1
LDA    1,SMD0+2
LDA    0,SMD0+3
JMP    @SMD0     ;RETURN

;DATA
SMD0:  .BLK    4    ;ENVIRONMENT STORAGE
SMD1:  0          ;TEMPORARY STORAGE
SMD2:  SMD2       ;ADDRESS OF TIME STORAGE BLOCK - 1
SMD3:  .BLK    3    ;TIME STORAGE BLOCK
SMD4:  3600.       ;HOURS MULTIPLIER
SMD5:  60.         ;MINUTES MULTIPLIER
SMT1:  .TXT    /DATA ACQUISITION SEQUENCE INTERVAL /
      .END

```

```

MULT (AC1*AC2)
;NAME: MULTA (AC0+(AC1*AC2))
;THIS SUBROUTINE MULTIPLIES THE SINGLE PRECISION BINARY
;VALUES IN AC1 AND AC2 TOGETHER. THE RESULT IS RETURNED
;WITH HIGH ORDER IN AC0 AND LOW ORDER IN AC1.
;
;INPUT: AC1    MULTIPLIER
;        AC2    MULTIPLICAND
;
;OUTPUT AC0    HIGH ORDER OF RESULT
;        AC1    LOW ORDER OF RESULT
;
;NOTE THAT ENTRY AT SNL2 WILL
;        GIVE AC0 + (AC1*AC2).
;
;CALLING SEQUENCE
;        JSR    @MULT      ; (AC1*AC2)
;        RETURN
;    OR
;        JSR    @MULTA     ; AC0 + (AC1*AC2)
;        RETURN
;
;CREATE POINTER TABLE ENTRY
;        .ENT    MULT,MULTA
;        .ZREL
MULT:    SNL1
MULTA:   SNL2
;        .NREL
;
;
SNL1:    SUBC    0,0        ;CLEAR AC0, AND DO NOT DISTURB CARRY
SNL2:    STA     3,SND0     ;STORE RETURN ADDRESS
        LDA     3,SND1     ;LOOP COUNTER INTO AC3
SNL3:    MOVR    1,1,SNC    ;CHECK NEXT MULTIPLIER BIT
        MOVR    0,0,SKP    ;0,JUST SHIFT
        ADDZR   2,0        ;1,ADD MULTIPLICAND AND SHIFT
        INC     3,3,SZR    ;CHECK FOR 16TH TIME THROUGH
        JMP     SNL3       ;NO,CONTINUE
        MOVCR   1,1        ;YES,SHIFT LAST LOW BIT
        LDA     3,SND0     ;RESTORE AC3
        JMP     @SND0
;
;DATA
SND0:    0                ;TEMPORARY LOCATION
SND1:    -20              ;-16 IN DECIMAL, LOOP COUNTER
        .END

```

```

      ERROR
;THIS SUBROUTINE DECREASES THE ERROR COUNT
;A MAXIMUM OF N ERRORS CAN OCCUR.
;THE MAXIMUM IS STORED IN LOCATION
;"ERCNT"
;IF ERCNT=0 THEN RETURN CONTROL TO DMS
;
;INPUTS:NONE
;
;OUTPUTS:      NONE
;
;CALLING SEQUENCE
;      JSR      @ERROR
;      RETURN
;
;SUPPORT SUBROUTINES
;      TXTOT
;      CRLF
;      RDMS
;      BELL
;
;CREATE POINTER TABLE ENTRY
      .ENT      ERROR
      .EXTD     TXTOT,CRLF,RDMS,BELL
      .ZREL
ERROR: SOL1
ERCNT:  5      ;MAXIMUM NUMBER OF ENTRY ERRORS
      .NREL
;
SOL1:  STA      3,SOD0      ;STORE ENVIRONMENT
      DSZ      ERCNT
      JMP      .+5
      JSR      @CRLF
      JSR      @TXTOT
      SOT1
      JMP      @RDMS      ;OUTPUT MESSAGE
                        ;RETURN TO DMS
      JSR      @BELL      ;BLOW THE WHISTLE ONCE
      LDA      3,SOD0      ;RESTORE ENVIRONMENT
      JMP      @SOD0      ;RETURN

;DATA
SOD0:  0      ;ENVIRONMENT STORAGE
SOT1:  .TXT      /MAXIMUM NUMBER OF ENTRY ERRORS EXCEEDED/
      .END

```

```

      SAMP
;
;CALLING SEQUENCE:
;      JSR      @SAMP
;      RETURN
;
;INPUTS: NONE
;
;OUTPUTS:      NONE
;
;SUPPORT SUBROUTINES
;      CRLF
;      TXTOT
;      GETNUM
;
;CREATE POINTER TABLE ENTRY
      .ENT      SAMP
      .EXTD     CRLF, ERROR, TXTOT, GETNUM, RDMS, SAMNU
      .ZREL
SAMP:  SPL1
      .NREL
;
SPL1:  STA      3, SPD0      ;STORE ENVIRONMENT
      STA      2, SPD0+1
      STA      1, SPD0+2
      STA      0, SPD0+3
      JSR      @CRLF
      JSR      @TXTOT
      SPT1      ;ASK FOR # OF SAMPLES
      LDA      0, SPD1      ;LOAD # OF DIGITS
      LDA      1, SPD2      ;LOAD MAXIMUM VALUE
      JSR      @GETNUM      ;OBTAIN #OF SAMPLES
      JMP      SPL2      ;ERROR
      MOV      0, 1, SNR      ;TEST IF WAS ZERO
      JMP      @RDMS      ;WAS ZERO, SO QUIT
      STA      1, SAMNU      ;STORE # OF SAMPLES
      LDA      3, SPD0      ;RESTORE ENVIRONMENT
      LDA      2, SPD0+1
      LDA      1, SPD0+2
      LDA      0, SPD0+3
      JMP      @SPD0      ;RETURN
SPL2:  JSR      @ERROR      ;PROCESS ENTRY OR LIMIT ERROR
      JMP      .-20
;DATA
SPD0:  .BLK      4
SPD1:  3          ;NUMBER OF DIGITS
SPD2:  100.       ;MAXIMUM INPUT FOR # OF DATA ACQUISITION RECORDS
SPT1:  .TXT      &NUMBER OF SAMPLES DESIRED : 2
      .END

```

```

;NAME:          DATAR
;THIS SUBROUTINE READS THE DATA FROM
;THE INTERFACE
      .ENT      DATAR
      .EXTD     TABLE
      .ZREL
DATAR:  SRL1
      .NREL
SRL1:   STA      3,SRD0
      LDA      2,SRD2
      LDA      1,SRD1
      NI OC    67
      NI OC    63
      IORST
      NI OS    63
      SKPDN    63
      JMP      -1
      DIAS     0,63
      STA      0, @TABLE
      NI OS    67
      SKPDN    67
      JMP      -1
      DIAS     0,67
      STA      0, @TABLE
      INC      2,2, SZR
      JMP      -5
      LDA      2,SRD2
      INC      1,1, SZR
      JMP      -15
      JMP      @SRD0
SRD0:   0
SRD1:   -64.
SRD2:   -50.
      .END

```



## ACKNOWLEDGEMENTS

I would like to take this opportunity to thank the Kansas Water Resources Board which provided the grant under which this project was developed (4). Special thanks also go to my advisor Dr. Donald Hummels, Dr. Myron A. Calhoun, and (the late) Dr. Dale E. Kaufman for their assistance in the design phase of this project.

A special thanks is extended to my wife, Ellen, for her moral support and to Patricia Copeland and William Ryan for their assistance in the production of the printed circuit boards used in this project.

THE DESIGN OF A COMPUTER INTERFACE  
FOR A RADAR PROCESSOR

by

CARL C. ANDREASEN

B.S. Kansas State University, 1974

AN ABSTRACT OF A MASTER'S THESIS

submitted in partial fulfillment of the  
requirements for the degree

MASTERS OF SCIENCE

Department of Electrical Engineering

KANSAS STATE UNIVERSITY

Manhattan, Kansas

1976

## ABSTRACT

This thesis describes the design, implementation, and use of a special-purpose data acquisition interface which controls the transfer of binary coded decimal values representing the integral of the radar video signals in a sector of an annulus of side one-half mile and angle 5.625 degrees from the radar to the computer. This thesis is also intended to be used as an operating manual and to provide assistance for maintenance. Included is the theory of operation of the interface, a description of the interface signals that are generated from the other parts of the system and a description of how to use the associated interface control program.

A brief discussion is given on future uses of the data collected by this interface.