ON THE ADAPTABILITY OF MULTIPASS PASCAL COMPILERS
TO VARIANTS OF (PASCAL) P-CODE MACHINE ARCHITECTURES

by

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B. S., Pittsburg State University, Pittsburg, Kansas, 1979

A MASTER'S REPORT

submitted in partial fulfillment of the
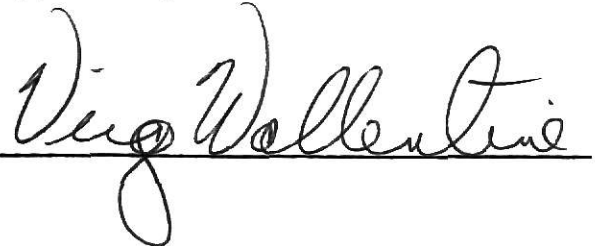
requirements for the degree

MASTER OF SCIENCE

Department of Computer Science

KANSAS STATE UNIVERSITY
Manhattan, Kansas

1981

Approved by:

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INTRODUCTION

1.1 MOTIVATION AND STATEMENT OF PROJECT

There are a number of situations where it is important that programs be portable among several machines. That is, it becomes important to be able to write a program in some high-level language and run it on a variety of different machine types (e.g., CDC 6600, IBM 370, PDP-11) without altering the source code. Such a situation occurs in a generalized network of heterogeneous computers where programs should not be unduly restricted from running on any available machine.

NADEX (Network ADaptable EXecutive) [YOUN79a, YOUN79b] is an operating system, implemented in Concurrent Pascal [BRIN77], which supports general graphs (software configurations) of communicating software nodes. The nodes of a software configuration may be sequential or concurrent programs and are not necessarily all resident on the same machine, or even the same machine type. So, for a NADEX network, there are three areas related to portability which must be addressed. First, some portion of NADEX must reside in each machine connected to the network; thus, that code must be executable on each machine type. Second, application programs should be executable on any machine in the network. Third, there may be some machines which are not capable of running compilers which produce object code for themselves, for example, because of small memory size.

One way to provide portable control and application programs is to write cross-compilers which will run on a machine and produce object code for each machine type. Source code can then be compiled for any machine type and sent via a very simple communications

protocol to the hardware on which it is to be executed. This is the approach used to provide NADEX support on the Pascal Microengine [MICR79] and described in this report.

Kansas State University currently has an Interdata 8/32 and Western Digital Pascal Microengine (among other machines) which could be included in a network running under NADEX. Concurrent and Sequential Pascal compilers run on the 8/32 and produce object code for it. This report describes an effort to modify a multipass Concurrent Pascal compiler so that it will execute on the Interdata 8/32 and produce object code for the Microengine.

The compiler to be modified generates P-code which appears to be very close to the Microengine's machine code. However, there are subtle but significant differences. A primary objective of the project was to assess the impact of these differences on a multipass compiler.

1.2 REPORT ORGANIZATION

This report consists of five chapters, apart from this introduction. Chapters two and three contain a juxtaposition of the compiler's original target machine (virtual, Concurrent Pascal) and the new one (Pascal Microengine) for which it was to be modified. The overall structure of the Concurrent Pascal compiler is discussed in chapter four. The next chapter is a description of the changes which were made to each pass of the compiler. Since pass six was so extensively changed, the modifications to it are grouped by the affected area of the pass: the objects which are generated as output, the overall structure. and the generated code sequences. Chapter six consists of an identification of the results of the project and the

work which remains. The reader is assumed to be familiar with the concepts of the Concurrent Pascal language [BRIN77].

1.3 NOMENCLATURE

Probably none of the terms used in this report will be new to the reader, but there could be some ambiguity surrounding their usage here.

In this report routine will be used to refer collectively to the Concurrent Pascal constructs PROCEDURE, FUNCTION (both ENTRY and non-ENTRY), and the initial statements (BEGIN...END. block) of processes, monitors, classes, and sequential programs. (The initial statement of a concurrent program is the same as the initial statement of a process.)

Stacks shown in the figures grow from the top of the page (high addresses) to the bottom (low addresses). The heap grows in the opposite direction. Generally, instructions can refer to three data areas in the stack-- local variables, global variables, and stack operands. Local variables are identical to Concurrent Pascal temporary variables. Global variables are the same as the permanent, or shared, variables of Concurrent Pascal. Instructions which push or pop values act on the operand stack.

Concurrent Pascal processes are synonymous with Microengine tasks-- each is the schedulable entity on the corresponding machine.

A sequential program running under a concurrent process may call certain ENTRY routines of the process. The accessible routines are named by the interface definition in the process. Those same routines are named in the sequential program's prefix. At run time, calls by the sequential program to its prefix routinesre mapped to a calls to

the corresponding process ENTRY routines. In the Concurrent Pascal virtual machine this mapping is performed by a jump table. On the Microengine, a special code segment, called the interface segment, (which contains only one routine-- the interface routine) performs the same mapping.

Actions of the two compiler systems can be controlled by specifying options. There are compiler options and driver options. Compiler options are specified at the beginning of the source code being compiled [HART76]. Driver options are specified as parameters in the CSS invocation line which the user enters at the console to start one of the compilers.

Items on the operand stack are referred to by their position relative to the top of the stack. TOS refers to the item on top of the stack, regardless of its length. For instance, the TOS item could be an integer value, occupying one word, or a real, which occupies two words. TOS-1 refers to the item which was pushed onto the stack immediately before the TOS item. Similarly, TOS-2 is the item pushed before TOS-1.

1.4 COMPILER GENEOLOGY

Figure 1 is a graphic illustration of the relationships between the various compiler versions in the geneology described below. At first the intention was to modify the code generator in HCPASCAL-- the original version of Concurrent Pascal received several years ago from the California Institute of Technology [HART76]. The object code (P-code) which it generates seemed close to the Microengine instruction set. The source code for HCPASCAL was not available, however, so an extended version of it (MCPASCAL) was used for the

FIGURE 1. Geneology of various Concurrent Pascal compilers at Kansas State University.

project. MCPASCAL [SCHM77] is the same as HCPASCAL, except that it allows use of the <u>manager</u> construct in the source language. For the purposes of this project, the portions which deal with the compilation of managers were ignored. The version of MCPASCAL which generates Microengine machine code is MEPASCAL, and is the result of the work described in this document. CPASCAL, another version of HCPASCAL, accepts the same source language as HCPASCAL but produces object code for the Interdata 8/32. CPASCAL makes more passes over the intermediate code than its parent, but also performs machine-independent and -dependent optimization.

ARCHITECTURE AND ORGANIZATION OF THE PASCAL VIRTUAL MACHINE

The MCPASCAL compiler generates code (called P-code, or virtual code) for a hypothetical machine, not for any existent hardware. P-code will execute on a real machine only if it is run interpretively, or if it is further translated to the language of a real machine. The software which implements the virtual machine has three major aspects: the virtual P-machine (that is, the machine which would be visible to a P-code programmer), the interpreter which presents that view, and the software kernel which interfaces directly with the real hardware.

2.1 VIRTUAL P-MACHINE

The virtual machine executes P-code, has a stack architecture, and uses five virtual registers, designated Q, G, B, S, and H. The Q register is the virtual program counter. G is the global base register, and can point to either of two types of data areas. When P-code in the anonymous initial process is being executed, G points to the data area containing the concurrent program's global variables. However, within system components, G points to the data area containing that component's permanent (global, shared) variables. B is the local base register, and points to the local variable area of a routine, regardless of whether or not it is a system component ENTRY routine. The stack pointer, S, always points to the top of the stack (the last word which was pushed onto it). H is the heap pointer, and points to the first byte of free space.

When a routine is called, machine state information is pushed onto the stack by the CALL instruction and the first instruction in

the called routine (for example, ENTER). This structure, called a markstack, consists of the return address in the caller's code (contents of caller's Q register), caller's global base (G), caller's local base (B), the value of S before the caller pushed the actual parameters (if any), and the source code line number of the called routine. The markstack fields are pushed on the stack in the order just given.

In general, routine calls proceed in the following manner. If the called routine is a function, the caller pushes enough space to hold the value to be returned. The caller pushes actual parameters (if any), proceeding left to right through the source code parameter list, and then the return address. The called routine builds the rest of the markstack, resets the local base register, pushes enough space onto the stack to accommodate its local variables, and resets the global base, if necessary. Figure 2 shows a stylized configuration of the data areas and registers after a function has begun execution.

The object code of a concurrent program compiled by MCPASCAL is a sequence of 16-bit integers, divided into three sections: initial process information, virtual code, and long constants. Five integers comprise the initial process information:

1) The byte address of the last word in the object code, relative to the first byte of the code file. That is, relative to the first byte of _this_ integer;

2) The number of bytes of virtual code;

3) The amount of stack space (bytes) required by the program;

4) The number of bytes required for the initial process' permanent variables;

5) The number of bytes in the long constant pool.

HIGH ADDRESSES

FIGURE 2. Stylized stack and register configuration after a Concurrent Pascal virtual P-machine function call. Calling routine and called function happen to share the same global variables.

The second section is the virtual code itself. The virtual machine instruction set (P-code) is given in [ZEPK74], along with English and psuedocode descriptions of each instruction's operation. The third section is the pool of long constants (reals, strings (arrays of CHAR), and sets). The constant pool always contains at least 16 bytes since MCPASCAL always generates the null set, even if it is never used. Figure 3 shows the object code which results from compiling a null program.

2.2 INTERPRETER

Object code generated by MCPASCAL and HCPASCAL runs interpretively. The original interpreter [ZEPK74] ran on a PDP-11/45 and consisted of 1K bytes of assembly code. Structurally, it is a jump table and a series of code pieces which carry out the actual interpretation of virtual instructions. The virtual operation codes are indices into the jump table whose entries are the addresses of the corresponding code pieces. Each code piece ends with the PDP-11 assembly instruction

$$MOV \quad @(Q)+,P$$

where P is the real machine program counter. The Q register contains the address of the next virtual operation code to be interpreted. The single move instruction uses the contents of the word to which Q points (the opcode) to locate a word in the jump table. The content of that word is then loaded into the real machine's program counter, thereby jumping from the code piece which interpreted the current virtual instruction to the code which will interpret the next one. Figure 4 shows the arrangement of virtual code, interpreter, and program counters. [ZEPK74] gives a psuedocode description of the

RELATIVE BYTE ADDRESS

|←———— 1 WORD ————→|

0	34	ADDRESS OF LAST WORD
2	10	CODE LENGTH
4	30	STACK LENGTH
6	0	PERMANENT VARIABLES LENGTH
8	16	CONSTANT POOL LENGTH
10	JUMP	
12	(2)	OPERAND FOR JUMP INSTRUCTION
14	BEGINPROCESS	
16	(1)	OPERAND FOR BEGINPROCESS INSTRUCTION
18	ENDPROCESS	
20	0	
22	0	
24	0	
26	0	
28	0	
30	0	
32	0	
34	0	

VIRTUAL CODE

CONSTANT POOL

FIGURE 3. Object code file generated by MCPASCAL from the source program: BEGIN END.

FIGURE 4. Interpreter, virtual code, and program counters before (dashed) and after (solid) execution of the MOV instruction at location 670 in the interpreter. The MOV terminates the interpretation of the GLOBADDR virtual instruction and automatically initiates the interpretation of CONSTADDR. Q points to the next word to be interpreted.

action of each code piece.

2.3 SUPPORT FOR CONCURRENT PROCESSES (KERNEL)

In the PDP-11/45 implementation of Concurrent Pascal a kernel (2.8K bytes of assembly code) handles processor multiplexing and guarantees that processes have exclusive access to monitors [BRIN75]. It also performs some hardware interface functions, such as I/O with peripherals. The interpreter has access to the kernel, via the KERNELCALL operation, in order to obtain the services which it provides. For instance, the virtual instruction INITPROC causes the creation of a new process. On encountering INITPROC, the interpreter calls on the kernel to generate a unique process identifier and process control block (PCB). The newly-created process is then placed in the ready queue for subsequent execution, and control returns to the interpreter.

Monitor access is controlled in the kernel by performing operations on a data structure, called a gate, which contains the state of the monitor (busy or free), to ensure mutual exclusion. Each monitor has its own gate, a record, the fields of which are a boolean (OPEN) and a pointer to a queue of PCBs. If a process is in the monitor (OPEN=FALSE) when another tries to enter, the PCB of the entering process is placed in the gate's queue. When the process currently in the monitor leaves, another one is selected from the queue and allowed to enter. If the queue is empty when a process exits, OPEN is made TRUE. Figure 5 depicts the logical relationship of the kernel to the interpreter and virtual code of a concurrent program.

FIGURE 5. Logical relationship of hardware, kernel. interpreter, and concurrent program.

ARCHITECTURE AND ORGANIZATION OF THE PASCAL MICROENGINE

3.1 SYSTEM OVERVIEW

The Pascal Microengine [MICR79] is a 16-bit, word addressable stack machine, manufactured by Western Digital Company. It is a desktop microcomputer built around the WD/9000 processor chip set, and is a hardware implementation of the University of California at San Diego (UCSD) virtual P-machine. The five LSI/MOS chips which constitute the processor are the Data Chip which contains the microinstruction decoder and ALU; the Control Chip, where the macroinstruction decoder, microinstruction counter, and I/O control logic are housed; and three 22 X 512-bit MICROM chips for microinstruction storage. The machine has 64K bytes of RAM, two asynchronous serial ports, one 8-bit parallel port, and controllers which give disk units direct memory access. The system at Kansas State University has two eight-inch flexible-disk drives, a CRT, and a printer attached. The hardware is supported by the Pascal Operating System which was written by the University of California at San Diego. The system supports UCSD Pascal, a variant of standard (sequential) Pascal. Assembly language is not supported.

3.2 P-MACHINE

The microengine has an architecture which is similar to that of the Concurrent Pascal virtual machine described earlier. The next three sections describe some significant aspects of that architecture.

3.2.1 Registers

The Microengine has an extensive set of logical registers. Those which correspond to the Q, G, B, S, and H registers of the virtual Concurrent Pascal machine are the IPC, BP, MP, SP, and SPLOW registers. IPC (Interpreter Program Counter) is the program counter. It contains a pointer (byte measure), relative to the start of the code segment (described below), to the instruction after the one currently under execution. BP (Base Pointer), the global base register, points to the markstack of a procedure at the outermost level of nesting in the source code. MP (current Markstack Pointer), the local base register, points to the markstack of the procedure invocation which is currently under execution. SP (Stack Pointer) points to the last word which was pushed onto the stack. SPLOW (Stack Pointer LOWer limit) points to the first free word in the heap, and also marks the last location into which the stack can grow. This is the stack's lower limit since the stack grows from high to low addresses.

In addition, there are other logical registers which have no analogs in the Concurrent Pascal virtual machine. SPUPR (Stack Pointer UPPeR limit) points to the location where stack growth begins. SEGB (SEGment Base) points to the code segment currently under execution. PRIOR (PRIORity), an 8-bit register, holds the CPU priority designation of the current task. Another 8-bit register, FLAGS, contains task state flags, but these have not yet been defined [MICR79]. If the current task gets placed in a linked data structure (ready list or semaphore wait queue, for example) the WAITQ (WAIT Queue) register contains a pointer to the next task in the structure.

The registers mentioned so far are logical registers in the

Microengine P-machine. There are three others which appear to actually exist in hardware. These are RQP (Ready Queue Pointer) which points to the first task on the ready list, CTP (Current Task Pointer) which points to the task presently running, and SDP (Segment Dictionary Pointer) which points to a vector of indirect pointers to code segments.

The machine instructions refer to all registers by number:

-3: RQP   -2: SDP   -1: CTP   0: WAITQ

1: PRIOR   1: FLAGS   2: SPLOW   3: SPUPR

4: SP   5: MP   6: BP   7: IPC

8: SEGB.

In register number 1 PRIOR is the low order byte, and FLAGS is the high order byte.

3.2.2 Object Code File Format

In disk file directory listings, object code files are identified by the extension ".CODE". Stored on disk, a code file (figure 6) consists of a sequence of contiguous 512-byte physical disk blocks. The first block in the file (block number zero) is a header which describes each of the code segments (up to sixteen) which start in block number one.

The header block (figure 6) has four areas of data in it:

1) Segment dictionary-- Sixteen entries, one per segment, regardless of the number of segments actually in the file. Each entry consists of two integer fields. The first is the block number (relative to the header which is block zero) of the first block in the corresponding segment. The second entry is the number of meaningful words in the segment. Meaningless

LOW ADDRESS

CODE FILE
HEADER BLOCK

SEGMENT

NO. 0

SEGMENT

NO. 1

SEGMENT

NO. N
(MAX. N =15)

HiGH ADDRESS

SEGMENT
DICTIONARY

| SEG 0 START BLOCK |
| SEG 0 LENGTH |
| SEG 1 START BLOCK |
| SEG 1 LENGTH |

SEG 15 START BLOCK
SEG 15 LENGTH

NAME
BLOCK

SEG 0 NAME
SEG 1 NAME

SEG 15 NAME

STATE
BLOCK

SEG 0 STATE
SEG 1 STATE

SEG 15 STATE

144 WORDS
FOR FUTURE USE

FIGURE 6. Microengine object code file general layout (left) and enlargement of header block (right). Each code segment may extend over several physical disk blocks. File may contain a variable number of segments, but each section of header contains 16 entries, some of which may be null.

words may occur at the end of a segment in order to pad to the end of a block. If the file contains fewer than sixteen segments zeros are used to indicate null entries.

2) Segment names-- Sixteen entries, as above. Each entry is a 16-byte string which is the ASCII character name of the segment, left justified, blank being the pad character. The entry is all blanks if the corresponding segment is null.

3) Segment state descriptors-- Sixteen integer entries, as above. Valid entry values are in the range zero through four. The entries encode linkage information such as the presence of external references and whether they have been resolved. For more information see MICR79, "Linker Conventions and Implementation". The entry for a null segment is zero.

4) The next 144 words are, according to [MICR79], reserved for future use. Most, but not all, of this area contains zeros. The meaning and function of the nonzero entries is unknown, and altering them seems to have no effect on the execution of the file and segment.

Each segment (figure 7) in the file may extend over any number of blocks and consists of a header word, a variable number of routines, a procedure dictionary, a final word of segment information, and possibly some padding. The header word contains the address (measured in words, relative to the start of the segment-- word zero) of the last meaningful word in the segment (the "final word of segment information" just mentioned).

Each routine (figure 7) is made up of a (possibly empty) constant pool and two words of run-time information, followed by the machine instructions themselves. The constant pool is aligned on a word

FIGURE 7. Microengine code segment layout (left) and enlargement of a typical routine (right).

boundary, and is similar to the constant pool of the Concurrent Pascal virtual machine. The first information field is the EXIT-IC, a word-aligned pointer (byte measure, relative to the start of the segment-- byte zero), the target of which is the first machine instruction of the routine's epilogue. Usually, the epilogue consists of only an RPU (Return from Procedure-- User) instruction. Presumably, either the hardware or operating system stores this value into the IPC register when a fatal run-time error occurs so that the stack gets cleaned up before returning to the operating system. The second information field, DATASIZE, is the number of words of local variable space required by the routine. The routine's code appears last, and because of prior alignments, it necessarily starts on a word boundary.

The segment's procedure dictionary (of variable length) follows the code of the last procedure and contains a pointer for every routine in the segment (see figure 7). The target of each pointer is the DATASIZE field of the corresponding routine. The pointers (word measure, relative to the start of the segment-- word zero) are arranged in what one would normally consider reverse order. That is, if there are 50 routines in the segment, the pointer for routine 50 comes first. and that for routine 1 comes last. Note that although the routine pointers are in reverse routine-number order, the routines themselves will not be in the same order since the initial statement of a routine is assigned its number before any enclosed routines are assigned theirs. The numbering scheme is further disrupted with each nesting level.

The last meaningful word in the segment contains two pieces of information. The even-address byte contains the segment's identifying

number (zero through fifteen) within the code file, and the odd-address byte contains the number of routines in the segment. If this word is not the last one in the block, the rest of the block is padded with zeros. It must be noted that the segments are the unit of execution, not the code file. During any execution, only the segments which are specifically invoked actually enter main memory to have the processor applied to them.

3.2.3 Routine Invocation and Stack Organization

Routine invocations involve a good deal of cooperation between hardware and software. If the called routine is a function, software in the calling routine pushes enough empty space onto the stack to hold the function value when it is returned. If parameters are required, the caller also pushes them onto the stack. Assume for the moment that the calling and called routines are part of the same code segment. The machine instructions which make such intrasegment calls are CPL (Call Procedure-- Local), CPG (Call Procedure-- Global), and CPI (Call Procedure-- Intermediate). During the execution of any of these instructions the hardware performs a number of operations. It uses the procedure number (fetched as an instruction operand) of the called routine as a backward index into the segment's procedure dictionary, fetches the pointer to the routine's DATASIZE field, and pushes onto the stack that number of words for use as the local variable space. It then builds a four-word markstack.

The first markstack word to be pushed on the stack contains two one-byte fields. One is the number of the code segment containing the calling routine's code (or zero, if the calling and called routine are in the same segment). The other has been reserved for future use. The

return address in the caller's code (contents of the IPC register) is stored in the next word. The third word is the dynamic link (pointer to the caller's local variables), which is a copy of the caller's MP register. Since UCSD Pascal allows nested routine definitions in the source code, the fourth word is the static link— a pointer to an enclosing routine's variables. The value used for the static link depends on the nesting level (lexical level) of the called routine, and the particular instruction used. See CPL, CPG, and CPI instructions in [MICR79]. Figure 8 depicts the general notion of static and dynamic links.

After construction of the markstack, the hardware copies the SP register into MP, making the called routine's local variables addressable. Hardware also calculates the IPC value for the first instruction of the called routine and execution continues at that point. The configuration of the stack after a function begins execution is shown in figure 9.

If the called routine is not in the same segment as the caller (that is, external) the invoking instruction is either CXL (Call eXternal Local procedure), CXG (Call eXternal Global procedure), or CXI (Call eXternal Intermediate procedure). Before the UCSD Pascal compiler generates a call to an external user routine, it generates an external call to a well-known procedure in an operating system segment which is always core-resident at run time. That routine fetches the required segment from the disk and pushes it onto the stack. Thus, the compiler and operating system work together to ensure that external procedure-call instructions will never be forced to deal with the invocation of a routine which is not in main memory. The run-time operation of external call instructions is the same as the other call

```
PROGRAM P;
"program variables"

PROCEDURE R1;
"R1 variables"

    PROCEDURE R2;
    "R2 variables"
    BEGIN "code" END;

    PROCEDURE R3;
    "R3 variables"
    BEGIN R2 END;

BEGIN R3 END;

PROCEDURE R4;
"R4 variables"
BEGIN R1 END;

BEGIN R4 END.
```


FIGURE 8. The UCSD Pascal source code on the left produces a run-time stack configuration similar to that shown on the right when procedure R2 executes. Dynamic links (solid arrows) always point to the caller. Static links (broken arrows) point to the enclosing routine activation. Since R1 and R4 are enclosed by the same routine (P) their static links point to the same markstack. Likewise, for R2 and R3 whose parent is R1. The markstacks shown here are not complete.

STACK
GROWTH

```
                        ┊                  ┊
                        ┊   CALLING        ┊
                        ┊   ROUTINE'S      ┊
                        ┊  OPERAND  STACK  ┊
                        ┊                  ┊
                   ┌────────────────────────┐
                   │    FUNCTION  VALUE      │
                   ├────────────────────────┤
                   │                        │
                   │      PARAMETERS        │
                   │                        │
                   ├────────────────────────┤
                   │                        │
                   │     FUNCTION'S         │
                   │     LOCAL  VARS        │
                   │                        │
                   ├───────────┬────────────┤
                   │ CALLER'S  │            │
                   │ SEG. NO.  │  RESERVED  │
                   ├───────────┴────────────┤
                   │   RETURN  ADDRESS      │
                   ├────────────────────────┤
                   │    DYNAMIC  LINK       │
                   ├────────────────────────┤
                   │    STATIC  LINK        │
                   └────────────────────────┘
```

MARKSTACK

MP

SP

FIGURE 9. Snapshot of a generalized Microengine function call after the function has started executing. Compare the position of the markstack relative to parameters and local variables with that in figure 2.

instructions except that the segment number operand is used to find the address of the code segment containing the called procedure before indexing into the procedure dictionary.

3.3 SUPPORT FOR CONCURRENT PROCESSES

The Microengine has several design features which can facilitate the execution of concurrent processes, although not all of them are supported by the hardware at this time.

The embodiment of a process is the Task Information Block (TIB) which is comparable to a process control block and contains the fields WAITQ, PRIOR, FLAGS, SPLOW, SPUPR, SP, MP, BP, IPC, SEGB, HANGP, XXX, SIBS, MAINTASK, and STARTMSCW. The first ten are the nonnegative logical registers described in section 3.2.1. If the task gets placed in a semaphore wait queue, HANGP can be used to hold a pointer to the semaphore data structure (described below) on which it is waiting. XXX is an unused integer field. SIBS points to an array of segment dope vectors (segment information blocks). The Western Digital documentation [MICR79] states that MAINTASK is of type BOOLEAN and that STARTMSCW is a pointer to a markstack, but does not explain their functions. The reader is left to make his own assumptions.

The target of the pointer SIBS is a vector of Segment Information Blocks (SIB). A SIB contains the fields SEGBASE, a pointer to the core-resident code segment; SEGLEN, the number of words in the segment; SEGREFS, the number of routine calls currently active in the segment; SEGADDR, the absolute physical disk address where the segment resides on secondary storage; and SEGUNIT, the number of the disk drive where the segment resides.

Figure 10 shows several tasks in, say the ready list, and how

FIGURE 10. Linkage between tasks and private code segments in the Microengine. TIBs and SIBs contain more information than shown here. Each SIB (numbered rectangles in the SIB vectors) has a SEGBASE field, although they are not all shown here.

each one has some private code segments bound to it through the SIBS field and SIB vector. The segments are "private" to each task since it would be difficult (but not impossible) for a casual programmer to execute code which traversed all the links in order to tamper with the code of some other task. In external routine call instructions, the segment number operand refers to a private segment if its value is 128 or greater.

Other segments can be easily called by any task in the system. The SIBs for these "shared" segments are not gathered together in a single data structure like a task's private segments. Instead, pointers to them reside in an array, the segment dictionary, which in turn is pointed to by the SDP register (see figure 11). In external routine call instructions, the segment number operand refers to a private segment if its value is in the range 0 through 127. However, there is an implementation restriction which limits shared segment numbers to the range 0 through 15.

Figure 12 shows how neatly the two segment schemes mesh. The three negative-numbered registers (RQP, SDP, and CTP) point to the ready queue, shared segment dictionary, and currently executing TIB respectively. When a process comes up for execution its private code segments become available through the CTP register, but it does not have easy access to the code of other processes. However, all processes have easy access to the shared code via the SDP register. This means, for example, that operating system processes can keep sensitive code in the relatively secure private segments, while making available that code which really must be public. Notice that context switching only involves resetting the CTP register and advancing the ready queue pointer since process state information is always in the

FIGURE 11. Linkage between the Segment Dictionary Pointer register and code segments which can be shared by several tasks. SIBs contain more information than shown here.

FIGURE 12. Run-time code segment configuration.

TIB.

Machine instructions WAIT and SIGNAL are provided for process synchronization. The semaphore data structure on which these act consists of a COUNT field and a pointer to the semaphore's wait queue. COUNT holds the number of outstanding SIGNALs which have been issued on the semaphore. Presumably, the two instructions behave in a straightforward manner, similar to Dijkstra's P and V operators.

At the present time the hardware does not support interrupts and hardware context switching. According to [MICR79], however, the machine is designed to provide a vectored, fixed-priority interrupt system. There will be eight two-word entries in the vector. One entry for DMA completion, six for serial and parallel port I/O completion, and one for exceptions on either of the serial ports.

3.4 P-MACHINE SEMANTICS FOR CONCURRENT PASCAL

The purpose of a translator for any high level language is to map the source code to the architecture of the target machine. In the case of HCPASCAL, the language existed first, and a virtual machine was implemented to accommodate it. The machine was built according to the needs of the language. For MEPASCAL, the source language and target machine came into existence independently of each other, so the compiler must somehow make use of the existing architectural facilities of the Microengine to carry out the semantic actions of Concurrent Pascal.

3.4.1 Code Segments

In the virtual machine there are four code spaces: kernel,
interpreter, concurrent program, and sequential programs (which exist
as permanent variables in concurrent processes). These can be
accommodated on the Microengine in the following ways.

The interpreter is trivial since its function is performed by the
compiler and hardware.

The kernel code should be independent of the concurrent program
since it is designed as a relatively stable piece of software to
provide support for any Concurrent Pascal program. Therefore, it
makes sense to put it by itself in a Microengine shared code segment.
If the kernel is always in the same segment (say, segment zero) it is
very easy to implement the kernel-call mechanism as a call to a
well-known routine within that segment. The kernel operator and
operands would have to be pushed onto the stack as parameters by
compiler-generated machine code. If it becomes necessary to modify
the kernel relatively frequently, a second approach to kernel-call
implementation might be more useful. Since the number by which the
routine is known to callers can be affected by changing the
arrangement or number of routines in the kernel segment, it would be
advantageous to isolate the kernel-call handler in a second
(well-known) kernel segment where routine number one merely examines
the kernel operator and passes the operands to the appropriate routine
in the segment. This can be done via routine calls embedded in a CASE
statement. See figure 13 for a pseudo UCSD Pascal outline of such a
segment. In this report we will assume the entire kernel is in only
one segment.

The classes, monitors, and processes of the concurrent program

```
TYPE
KERNOPTR = (INITGATE, ENTERGATE, LEAVEGATE, ENDPROCESS,
            INITPROCESS, REALTIME, DELAYGATE, CONTGATE,
            STOPJOB, WAIT, SYSERROR, IO);


SEGMENT PROCEDURE KCALLHDLR (OPTR: KERNOPTR;
                            OPND1, OPND2, OPND3, OPND4: INTEGER);
CONST
IGATERTN =                    |
EGATERTN =                    |
LGATERTN =                    |
EPROCRTN =                    |
IPROCRTN =                    |
RTIMERTN =         routine numbers in
DGATERTN =            kernel segment
CGATERTN =                    |
STJOBRTN =                    |
WAITRTN  =                    |
SYSERRTN =                    |
IORTN    =                    |

KERNSEG = 0;      "SEGMENT NUMBER OF KERNEL PROPER"

BEGIN
CASE OPTR OF
    INITGATE:    PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, IGATERTN;
    ENTERGATE:   PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, EGATERTN;
    LEAVEGATE:   PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, LGATERTN;
    ENDPROCESS:  PUSH OPERANDS; EXTERNAL CALL TO KENRSEG, EPROCRTN;
    INITPROCESS: PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, IPROCRTN;
    REALTIME:    PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, RTIMERTN;
    DELAYGATE:   PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, DGATERTN;
    CONTGATE:    PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, CGATERTN;
    STOPJOB:     PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, STJOBRTN;
    WAIT:        PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, WAITRTN;
    SYSERROR:    PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, SYSERRTN;
    IO:          PUSH OPERANDS; EXTERNAL CALL TO KERNSEG, IORTN
ENDCASE
END;
```

FIGURE 13. Pseudo UCSD Pascal outline of specialized segment to handle kernel calls. Pushing the operands and calling the external routine are separate machine operations. In high level code this would appear as a normal routine call with parameters.

can be placed in another shared segment where each initial block, routine, and entry routine of the concurrent program is one routine in the concurrent code segment. The segment must be sharable since monitors and classes can be entered by any process which has the appropriate access right. The concurrent segment will always be segment number one. Always placing the concurrent program in segment number one makes it well-known in the system and will aid the implementation of the interface mechanism (explained below).

Sequential programs and their interface with host concurrent processes present special problems. When calling a sequential program, the interpreter for the Concurrent Pascal virtual machine often makes use of a property not found on the Microengine; namely, that real memory is a single, linear address space where all addresses are alike, and that once an address has been obtained, the processor can easily be forced to jump to that location. Sequential program invocation on the virtual Pascal machine is fairly straightforward. The concurrent program must somehow load sequential code into a variable (a large array, for example), and the compiler must generate code to push the variable's address onto the stack. CALLPROG is the virtual instruction which actually starts the program. During its interpretation, the return address is saved on the stack and the virtual program counter (Q) is loaded with the address of the first sequential program instruction which is located a fixed distance from the start of the program variable (whose address is on the stack). On the Microengine, the address of the program variable (in the real memory address space) can be pushed onto the stack by any of several machine instructions. However, it cannot be placed directly into the IPC register to cause a jump to the sequential program since the

hardware expects the address in the IPC to be an offset from the beginning of a code segment. A possible solution to this problem is depicted in figure 14 and runs as follows. Assume that when the kernel creates a TIB for a process it also allocates space for two SIBs for segments private to the newly created task. The segment numbers of these can always be 128 and 129, where segment 129 is the sequential program segment. (The purpose of segment 128, the interface segment, will be described below.) As part of the program invocation code, the compiler generates instructions which follow pointers from the CTP register, through the current task's TIB to its SIB vector, and pop the address of the code variable into the SEGBASE field in SIB 129. The next instruction generated is an external call to routine number one in segment 129. The effect of all this at run time is to make the code variable look like a code segment to the hardware, and start the sequential program by means of a normal external routine invocation.

Operating system services are provided to the sequential program through an interface which names the process ENTRY routines to which the sequential program has access. The operation of the interface mechanism is best demonstrated by example. The (meaningless) concurrent program shown in figure 15 contains an interface definition (line 7) for the sequential program J. The program is actually invoked in line 27. The MCPASCAL compiler produces code which performs the following functions in order to get J started:

1) Push onto the stack the address of the first instruction of each ENTRY routine named in line 7, in reverse order of appearance in the source code. That is, push the address of PE3, then that of PE2, then that of PE1;

2) Push the parameters (the values 1 and 2, in the example);

FIGURE 14. Configuration of pointers just before calling a sequential program.

```
1    TYPE
2         PRC = PROCESS;  +1000
3         TYPE    CODE = ARRAY[1..1000] OF INTEGER;
4         VAR     CODEVAR: CODE;
5
6         PROGRAM J (A, B: INTEGER; C: CODE);
7         ENTRY PE1, PE2, PE3;
8
9         PROCEDURE ENTRY PE1 (PARAM1: INTEGER;
10                                VAR PARAM2: INTEGER);
11        BEGIN
12        PARAM2 := PARAM1 + 100
13        END;
14
15        PROCEDURE ENTRY PE2 (PARAM1: INTEGER;
16                               VAR PARAM2: INTEGER);
17        BEGIN
18        PARAM2 := PARAM1 + 200
19        END;
20
21        FUNCTION ENTRY PE3 (PARAM: INTEGER): INTEGER;
22        BEGIN
23        PE3 := PARAM + 300
24        END;
25
26        BEGIN
27        J (1, 2, CODEVAR)
28        END;
29
30
31   VAR    PRCV: PRC;
32
33
34   BEGIN
35   INIT PRCV
36   END.
```

FIGURE 15. Sample Concurrent Pascal program containing an interface definition (line 7) for the sequential program defined in line 6 and invoked in line 27.

3) Push the address of the variable containing the sequential program code;

4) Execute the CALLPROG instruction as described above.

A snapshot of the run-time stack at this point is shown in figure 16. Notice that the addresses of the process ENTRY routines (interface routines) in the stack constitute a jump table, built at run time. The sequential program calls interface routines by executing the virtual instruction CALLSYS(PREFIX_INDEX). During the interpretation of CALLSYS, the value of PREFIX_INDEX is used as an index into the jump table on the stack and the processor branches to the corresponding ENTRY routine in the concurrent process. This mechanism allows a great deal of independence between the concurrent and sequential programs. The only point on which they must agree is the order in which the interface routines are to appear.

The design of an interface mechanism for the Microengine proved to be a significant challenge. The problems stem from two Microengine architectural features: routines are known and called by their number, and the number of the called routine must be supplied to the invoking instruction as an immediate operand, not a stack operand. The first item precludes the construction of a run-time jump table, but if the sequential program is allowed to call process ENTRY routines directly, two major points of independence are lost. First, the sequential program would need to know, at compile time, the identifying numbers (not just the order) of the entry routines it is allowed to use. Second, the number and arrangement of routines in the concurrent program could be changed, independent of the sequential program, only if it could be __guaranteed__ that the routine numbers of the process entries would be unaffected-- an extremely difficult, if not

```
           |                       |
           |       OPERAND         |
           |        STACK          |
           |                       |
           |-----------------------|
           |     @  PE  3          |   INTERFACE
           |     @  PE  2          |   (JUMP TABLE)
           |     @  PE  1          |
           |-----------------------|
           |        1              |   PARAMETERS
           |        2              |
           |-----------------------|
           |   (@  CODEVAR ) /      |
           |   @ SEQ. CONSTANTS    | <------- S
           |-----------------------|
           |  //////////////////   |
           |  //////////////////   |
           |  //////////////////   |
           |  //////////////////   |
           |- - - - - - - - - - - -|
           |                       |
           |                       |
           |                       |
           |                       |
           |                       |
```

FIGURE 16. Run-time stack after execution of the virtual instruction
CALLPROG which invokes a sequential program. CALLPROG uses the
sequential program variable address which it finds on top of the stack
and replaces it with the address of the sequential program's constant
pool. The sequential program's markstack will be be built later in the
lined area.

impossible, task. If routine call instructions took their operands from the stack, then the numbers of the interface routines could be placed on the stack as parameters to the sequential program. Invocation of an interface routine would only require pushing the appropriate parameter onto the top of the stack and initiating a routine call. However, as mentioned above, the architecture prevents such a maneuver.

The required Microengine mechanism must be a run-time mapping of prefix indices to interface routine numbers in the concurrent segment. The map can be realized in an interface segment (see figure 17) built by the compiler whenever it encounters the invocation of a sequential program which requires an interface. Some details of the scheme, such as the time at which the interface segment should be loaded into main storage, and the accommodation of more than one interface in a single process have not yet been resolved. The interface segment (always segment 128) contains only one routine which itself consists largely of a CASE statement. The case labels are the prefix indices which are known to the sequential program. The code of each case is an external call to a process ENTRY routine in the concurrent segment. If the routine numbers change because of some structural change in the concurrent program, the operands to the routine calls will be adjusted accordingly when it is recompiled.

Sequential program invocation can be the same as described above for the virtual machine, except that it is no longer necessary to push the jump table onto the stack. The invocation of a process ENTRY routine will be quite different. The sequential program will push onto the stack the prefix index of the ENTRY routine as a parameter. It then will call procedure one in segment 128, the interface mapping

```
SEGMENT PROCEDURE INTERFACE (PREFIX_INDEX: INTEGER);
BEGIN

RANGECHECK (PREFIX_INDEX);
CASE PREFIX_INDEX OF
  1: EXTERNAL CALL TO CONCURRENT SEGMENT, ROUTINE W;
  2: EXTERNAL CALL TO CONCURRENT SEGMENT, ROUTINE X;
  3: EXTERNAL CALL TO CONCURRENT SEGMENT, ROUTINE Y;
  .                             .
  .                             .
  .                             .
  .                             .
  .                             .
  N: EXTERNAL CALL TO CONCURRENT SEGMENT, ROUTINE Z
ENDCASE

END;
```

FIGURE 17. Pseudo UCSD Pascal description of concurrent/sequential interface segment. PREFIX_INDEX is the index number of the prefix routine the sequential program is calling. W, X, Y, and Z are the numerical identifiers which the routine-calling instructions use to invoke the corresponding process ENTRY routine. RANGECHECK causes a run-time error if PREFIX_INDEX is not in the range 1..N.

routine. The mapping routine will use the parameter as the case selector, and call the coresponding process ENTRY routine in the concurrent segment.

Figure 18 shows the proposed arrangement of segments required for a Concurrent Pascal program to run on the Microengine. Two segments (kernel and concurrent program) are shared by all tasks, and two (interface and sequential program) are privately associated with their controlling task.

3.4.2 Data Spaces

Just as code areas must be mapped to the Microengine architecture, so must Concurrent Pascal data spaces. They will be handled in the manner described here. Figure 19 shows the general state of affairs in the stack after the initial process has started executing its initial routine. This state can be reached through the following sequence of events. At IPL time the system is powered up and a human operator presses the RESET button on the rear of the processor cabinet. This causes the hardware to read into main memory a fixed area of data from a disk unit which is well-known to it. The hardware assumes that it has just read in the first part of a bootstrap loader and proceeds to load the kernel by executing that information. Once the kernel code is in place it can begin execution and fetch the concurrent program code from a disk unit. At this point the stack consists of the kernel code segment, space for its variables, and a markstack. As concurrent object code is read into the processor, it is pushed onto the kernel's operand stack. Once the segment has been loaded, the kernel sets the segment dictionary pointer, builds a SIB, and puts the address of the segment into it. · Routines in the

FIGURE 18. Concurrent Pascal program on the Pascal Microengine.

KERNEL
CODE

KERNEL
VARIABLES

KERNEL
MARKSTACK

KERNEL
DATA
SPACE

CONCURRENT
CODE
SEGMENT(S)

OTHER INITIAL
PROCESS
VARIABLES

PARAMETERS

PERMANENT
(GLOBAL, SHARED)
VARIABLES

SYSTEM
COMPONENT
VARIABLE

INITIAL PROCESS
DATA SPACE

OTHER INITIAL
PROCESS
VARIABLES

IMAGINARY
MARKSTACK

INITIAL
PROCESS
MARKSTACK

(B)

ADDITIONAL
INITIAL
PROCESS
STACK SPACE

FIGURE 19. General layout of kernel data space and initial process stack space with a system component. While any part of the system component's code is executing, the global base register (B) points to the component record. The concurrent code segment is shown grossly out of scale for the sake of brevity.

concurrent segment can now be called. The concurrent program becomes active when the kernel calls its initial routine, routine number one.

Invocation of the initial routine causes the hardware to allocate stack space for its variables and build a markstack (see section 3.2.3). The amount of variable space to be allocated was determined by the Concurrent Pascal compiler (MEPASCAL). If any of the initial process' variables are system components, they will appear here, as shown in figure 19. The embodiment of a system component <u>variable</u> (instance of a system component type) is a record in the variable space of the initial process. The record has the same format as the top region of the stack, just after a routine call instruction has been completed. For example, suppose that a class is defined as

```
TYPE    CL = CLASS (P1, P2, P3: INTEGER);
        VAR   V1, V2: REAL;

              PROCEDURE ENTRY X;
              BEGIN
              "code"
              END;

              PROCEDURE ENTRY Y;
              BEGIN
              "code"
              END;

              BEGIN
              "initial code"
              END;
```

and that the concurrent program variables declarations include

```
        VAR    CL1, CL2, CL3, CL4: CL;  .
```

Then in the initial process variable space, along with any other concurrent program variables, there would be four records (CL1, CL2, CL3, and CL4), each containing five fields. Three integer fields at the high-address end of the record are the parameters P1, P2, and P3. Two real fields at the low-address end of each record are the permanent variables of the class, V1 and V2. Any time the class

instance CL2, for example, starts executing, the global base register (B) will be forced by compiler-generated software to point to an imaginary markstack adjacent to the high-address end of the record for CL2, as shown in figure 19. Unlike the local variable areas of procedures, functions, and **ENTRY** routines, which appear upon activation and disappear upon return, the initial process variables disappear only if the whole concurrent program returns to the kernel. Hence, the permanence of component permanent variables is realized. Before initializing a component variable, software pushes its parameters onto the stack. The actual INIT code sequence then pops them into the component variable's record. (For a process initialization, the parameter movement is handled by the kernel.) In this way access rights are "remembered" after the initial routine returns to its caller.

As figure 20 shows, process stack space is allocated contiguously in the order of process creation. The amount of space to be allocated to a process is determined at compile time. Once the space is allocated, it exists forever during execution. Within that space the stack for the process will rise and fall as code is executed and routines are activated. Since each process has its own stack operating concurrently with those of other processes, there will be several stacks and stack pointers existent in the machine at the same time. However, there is no ambiguity for the hardware regarding which stack pointer to use, nor is there any need for software to reset SP during context switches, since every process has its own SP in its TIB. When a TIB gets switched onto the processor (when the TIB becomes the target of the CTP register) the correct SP comes with it.

The amount of stack space needed by a process can be supplied to

FIGURE 20. General scheme of stack space allocation.

the kernel as a parameter to the kernel routine which creates processes. Even though the space information is _provided_ at process-creation time, it will not be _used_ until the kernel is called upon to create the _next_ process. Figure 21 shows how stack space for a new process (process 2) must be allocated in relationship to the most recently allocated space (process 1). During execution of routine call instructions the hardware apparently makes use of the SP register to determine the starting address of the routine's local variable area and, by implication, the location of the new markstack. This means that when process 2 is created, the kernel must set the SP register in process 2's TIB so that it points to the _last allocated word_ in process 1's stack space. When the new process comes up for its first execution time slice it will execute a call to its own initial routine, and the process markstack will automatically be built in the correct location. The placement of the initial routine's markstack is somewhat critical. If it spills into the previously allocated area, it could be destroyed, but placing it too far away wastes memory. The new process' global (permanent) variables have been allocated in the variable space of the anonymous initial process, so the new variable space is null. It seems possible to allocate process global variables here since they would be deallocated only if the process initial routine tried to return to its caller, but placing them in the initial process area maintains a consistency with the allocation scheme for class and monitor permanent variables.

If the method just described is to ever work successfully, a mechanism is required which will allow _one_ process to _enter_ the kernel routine which creates processes, and allow _two_ to _exit_ it _safely_. (The second process must not be allowed to cut back a nonexistent

VARIABLE
SPACE

MARKSTACK

PROCESS 1
MP

STACK SPACE
ALLOCATED TO
PROCESS 1

OPERAND
STACK

PROCESS 1
SP

ALLOCATED,
BUT NOT NOW
IN USE

PROCESS 2
SP (DURING INIT)

MARKSTACK

PROCESS 2
MP

STACK SPACE
ALLOCATED TO
PROCESS 2

OPERAND
STACK

PROCESS 2
SP (EXECUTION)

FIGURE 21. Boundary between the stack spaces allocated to two concurrent processes. When the kernel creates process 2, SP must point to the location indicated by the broken arrow. Notice that there are no variables here for the process 2 initial routine-- they are the permanent variables which have been allocated space in the data area of the anonymous initial process.

stack.) And yet, the hardware should do as much of the work as possible. This turns out to be no small task. Assume for the moment that the anonymous initial process is to start a process called PRC. The proposed mechanism calls for the compiler to generate code in the initial process which (in pseudocode) looks like:

```
CALL KERNEL TO CREATE PROCESS
SKIP NEXT INSTRUCTION
CALL PRC'S INITIAL ROUTINE
OTHER CODE,
```

and for the kernel's process-creation routine to be similar to:

```
BUILD TIB FOR PROCESS PRC
BUILD OTHER PROCESS STRUCTURES (SIB, for example)
FETCH DYNAMIC LINK (points to initial process SKIP)
ADD 2 BYTES TO DYNAMIC LINK VALUE
PUT (DYNAMIC LINK + 2) INTO PRC'S IPC
SET SP REGISTER IN PRC'S TIB
PUT PRC'S TIB IN READY LIST
RETURN TO CALLER.
```

Execution proceeds in the following manner. The anonymous process executes a normal call to the process-creation routine in the kernel. That routine builds a TIB and related data structures for PRC (see figure 18). Next, the IPC address of the SKIP instruction in the initial process is fetched. This is possible by means of the following steps:

1) Load the contents of the MP register onto the stack. This puts the memory address of the process-creation routine's markstack onto the stack.

2) Add one word to the top-of-stack value to yield the address of the dynamic link field in the markstack.

3) Push onto the stack the word pointed to by the top-of-stack word. That is, push indirectly the dynamic link.

The dynamic link is the IPC value of the SKIP instruction in the initial process. Two (the skip instruction is two bytes long) is added

to it to yield the IPC value of the CALL PRC instruction. That value (dynamic link + 2) is put into the IPC register field of PRC's TIB so that PRC will start execution there when it gets its first slice of processor time. PRC's SP register is initialized as discussed above, and the TIB is placed in the ready list. The anonymous process executes a normal return from the kernel and next executes the SKIP in the initial process. It then continues on to its other code. If the CALL were not skipped, the initial process would end up in the code of PRC, not its own. On the other hand, when PRC comes up for its first execution, it must execute the CALL so that it ends up in its own initial code. If PRC attempts to return from its initial routine, it will be prevented from entering the code of the initial process, from which it was called, by compiler-generated code which calls a kernel routine before the RETURN instruction can be executed. That routine (ENDPROCESS) simply removes PRC's TIB from the ready list.

STRUCTURE OF THE CONCURRENT PASCAL COMPILER

4.1 OVERVIEW AND SUMMARY OF PASSES ONE THROUGH FIVE

MCPASCAL is a seven-pass, recursive descent compiler written in Pascal32, a variant of sequential Pascal which generates Interdata 8/32 object code. Figure 22 shows the various parts of the compiler on Kansas State University's Interdata 8/32. MCPASCAL.CSS is an operating system command file which makes logical device assignments, allocates temporary intermediate code files and a permanent object code file, and initiates the sequential Pascal program MCPASCAL. MCPASCAL is the compiler's driving program and performs the following functions:

1) Scans the string of driver options requested by the user and saves them for future reference;

2) Invokes the passes of the compiler (MCPASSx-- Managers, Concurrent, PASS number x) in the proper sequence;

2) Invokes the program MNEM (explained below) as indicated in the driver options specified by the user;

4) Monitors the compilation and reports its progress (for example, passes completed and presence/absence of compilation errors) to the user's console.

The driver communicates with the programs it calls through variables of type ARGLIST and PROGRESULT which are defined in the prefix of each program.

The compiler passes use four disk files as shown in figure 23. The two temporary files contain the intermediate code produced by each pass. When an intermediate code file is no longer needed, it is overwritten by subsequent passes. For example, MCPASS3 writes over the output of MCPASS1 when it produces its own output.

FIGURE 22. Control structure of the MCPASCAL compiler.

FIGURE 23. Flow of intermediate code in the Mcpascal compiler system.

MNEM (MNEMonics) is a sequential program which accepts the number of a compiler pass as a parameter. It reads the intermediate code output of that pass and prints the contents on the same device (or file) as the source code listing generated by the first pass. The output is formatted, and operators appear as mnemonics instead of integers. Operands appear as integers and are enclosed by parentheses. MNEM is described more fully in Appendix A. This mechanism replaces the test mechanism built into the compiler and invoked as a compiler option. The test mechanism prints the intermediate code as integers. Operators are identified by preceding them with the letter "C".

The first pass of the compiler (MCPASS1) is the lexical scanner. In general terms, it translates the source program into a sequence of 16-bit integers which represent the program tokens, and produces a listing of the source file. It also performs three rather specific functions unrelated to lexical analysis. First, it generates the null set long constant. Second, it allocates heap space for the PASSLINK, a record data structure which remains in the heap between passes, and through which the passes communicate with each other. The third specific function is the analysis and initialization of compiler options which may appear at the beginning of the source program.

MCPASS1N (MCPASS1 No source) is the same as MCPASS1 except that it does not produce a source listing. Its existence allows the source/no-source option to be specified as a driver option when the compiler is invoked, rather than as a compiler option embedded in the source code. If "NS" is in the string of driver options, MCPASS1N is invoked as the first compiler pass instead of MCPASS1.

MCPASS2 performs recursive descent syntax analysis on the token

stream produced by the previous pass. At the end of the pass it saves the number of jump labels it created in the PASSLINK record.

Scope (name) analysis of program identifiers is performed by pass three. It enforces rules which, for example, demand that within a single block or record, identifiers have only one *meaning*, and forbid a nested component definition from referring to the parameters or variables of an enclosing component. In addition, the calculated length of the constant pool is saved in the PASSLINK, even though the constants themselves remain sprinkled through the intermediate code stream.

MCPASS4 performs the semantic processing of the declaration portions of the program. It analyzes types, assigns addresses to variables and parameters, and *assigns block labels to routines.* Semantic rules are also enforced here. For example, strings must be of even length; variables of type QUEUE must be variables of monitors only. Pass four consumes declarations, encodes their information, and distributes it wherever needed in the routine bodies. At the completion of the pass the intermediate code is merely a sequence of routine bodies. The number of routines in the program is saved in the PASSLINK.

Analysis of routine bodies *is performed by pass five and consists* of ensuring the compatability of operands with each other, and of operands with operators. For example, only an integer may be added to another integer, and the addition operator must be an _integer_ ADD, not real. The pass also generates addressing commands.

4.2 PASS SIX

The final two passes constitute a two-pass assembler. The first of these, MCPASS6, selects the final code, converts routine and jump labels to addresses, determines the stack space required for each routine and system component, and constructs the constant pool.

The main portion of the pass is a loop around a CASE statement, where the case labels are the operators of the input language. The loop consists of reading an operator (variable OP, in the MCPASS6 source code) and using it as the case selector value to perform the actions appropriate to that operator. If the operator is one which has operands in the code stream, they are read by one of the READxARG routines, where "x" is the number of operands (1-5) to be read. The global variables ARG1, ARG2, ARG3, ARG4, and ARG5 contain the operands after the read operation. The loop terminates when the EOM operator is encountered.

Some inputs translate on a simple one-for-one basis to the output language. For example, whenever PUSHCONST1 is the input operator its operand (the value to be pushed onto the stack at run time) is read into ARG1, and the output PUSHCONST2, followed by ARG1, is always emitted. Other translations are a bit more complex. For example, the intermediate code generated from PUSHVAR1 and its three operands (variable type, addressing mode, and displacement) depends on whether the variable to be pushed is of word type, whether it is a variable in a system component, and whether the displacement is positive. Based on the operand values, one or more instructions will be generated, and the displacement value may or may not be adjusted. Detailed descriptions of the translations are presented in sections 5.2.3 and 5.2.4.

In the pass six input language, the operands of jump instructions (JUMP1, FALSEJUMP1, and CASEJUMP1) are label numbers which have been generated during syntax analysis. A jump's destination is marked by a DEFLABEL1 instruction whose operand is a label number which is the same as that of the jump operand. Since jumps in the final code are in terms of displacements relative to the jump instructions themselves, the labels must be converted to displacements within the virtual code address space. As the first step in the conversion, this pass builds a table (JUMPTABLE) in the heap which will be used by the next pass. Label numbers are used as the table index to insert table entries which are object code location counter values. The global variable LOCATION serves as the location counter of the _final_ code to be produced by pass seven, and contains the address of the _next_ instruction to be generated. When a DEFLABEL1(<label>) instruction appears in the code stream, the current value of LOCATION is placed in the label-th position of JUMPTABLE. Label numbers are unique so there is no danger of overwriting previous entries. If the LINENUMBER compiler option is in effect a NEWLINE2 instruction is generated, otherwise DEFLABEL1 produces no code. When a jump instruction appears, a corresponding output instruction is generated and the location counter updated. The current value of LOCATION is then emitted, _but_ _LOCATION is not updated_ since the emitted value will be removed from the code stream by pass seven. Pass seven performs the rest of the label-to-displacement conversion.

Routine label numbers are converted to virtual addresses in a similar manner using BLOCKTABLE. CALL1 instructions correspond to jump instructions and cause the emission of the current LOCATION value as well as the output language instruction. Pass seven will use and

remove the LOCATION value placed in the code stream. ENTER1 is analogous to DEFLABEL1 and results in the insertion of the LOCATION value into the BLOCKTABLE entry indexed by the routine label. Final conversion to displacements is done by the next pass.

Pass six computes the maximum amount of run-time stack space required by each routine. This is possible since Concurrent Pascal does not allow recursive routine calls. The space requirement is the sum of the length of the routine's local variable area, the maximum size of the operand stack, the markstack size, and any additional reserved space (in the case of a process which is host to a recursive sequential program). The local variable length and size of additional reserved space are taken from the input code stream as operands to ENTER1 and are kept in the global variables VARLENGTH and STACKLENGTH, respectively. The markstack size is always five words. The maximum amount of space needed for the operand stack is determined as the pass scans the intermediate code. For each routine the variable TEMP simulates the rise and fall of the run-time operand stack. At the start of the routine (ENTER1 instruction) TEMP is set to zero. As the pass generates instructions which will push or consume stack operands at run time, TEMP is incremented or decremented by the length of the object pushed or consumed. MAXTEMP is also set to zero at the start of each routine. Each time TEMP is _increased_ its new value is compared to the current value of MAXTEMP. If TEMP is greater, its value is placed in MAXTEMP, thus recording the "high-water" mark of the operand stack up to that point in the routine. The operand stack space is also affected by routine calls. When a routine call instruction (CALL1) is encountered, the stack requirement of the called routine (obtained from STACKTABLE-- described below) is added to TEMP to simulate the

space _it_ uses during its activation. The run-time return from the called routine, and resultant release of stack space is simulated by immediately decrementing TEMP. When the compiler encounters the end of the calling routine (RETURN1 instruction) its total stack requirement is entered into a table (STACKTABLE), the index of which is the routine label, making this routine's stack requirement available to subsequent routines which might call it. Since Concurrent Pascal allows calls only to routines which have been defined earlier in the source code (and a previous pass enforces this restriction) the stack requirement values will always be in the table when needed at CALL instructions. Pass seven also uses STACKTABLE.

Long constants (compiler-generated null set, real. and string constants) appear in the input code as a CONSTANT1 operator, followed by the byte length of the constant, followed by the constant itself. The constant pool at this point is in the form of a table (CONSTTABLE), the entries of which contain one word of some constant. The global variable CONSTANTS counts the number of words of constants currently in the pool and is the index into CONSTTABLE. When a constant is found, CONSTANTS is incremented by one for every word of the constant which is read from the code stream and placed in CONSTTABLE. No output is emitted.

Logically JUMPTABLE, BLOCKTABLE, STACKTABLE, and CONSTTABLE are simple arrays of integers, but physically they appear as shown in figure 24. The LABELS, BLOCKS, and CONSTANTS values from the PASSLINK record indicate the number of entries required for each table. The tables are allocated space from the heap in pieces consisting of 100 table entries and space for a pointer to the next piece. So, for example. if the value in LABELS is 175, JUMPTABLE will consist of 200

FIGURE 24. Heap table layout. N is the number of entries actually
needed.

table entries. In general. tables have 100*⌈N/100⌉ entries allocated, where N is the number actually needed. The tables are known by their pointer (of type TABLEPTR) to the first piece, such as the JUMPTABLE variable.

Three routines manage the heap tables. ALLOCATE gets the required space from the heap, sets the pointers between table pieces, and returns to the caller a pointer to the first table piece. ENTER uses as input a pointer to a table's first piece, an index, and an integer value. It places the integer in the table named by the pointer, in the position specified by the index. The _logical_ equivalent is TABLENAME[INDEX]:=VALUE. Table entries are retrieved by the ENTR function which uses a table pointer and an index as input. It is _logically_ equivalent to ENTR:=TABLENAME[INDEX]. All four tables built by MCPASS6 are left in the heap for MCPASS7 to use. They are passed by putting the pointer to each table's first piece in the PASSLINK record. The address of the PASSLINK itself is sent to the next pass as PARAM[2].PTR. a program parameter.

The source listing of MCPASS6 contains several sections which require some comment. In the constant definitions, "VIRTUAL DATA TYPES" refers not to Pascal TYPEs but to the length attribute of an object-- objects have the same length as a byte, word, real value, or set variable. The virtual addressing constants have the following meanings:

MODE0=0; Constant-- the object is to be addressed relative to the start of the constant pool;

MODE1=1; Procedure-- the object is a parameter or local variable of a non-entry routine;

MODE2=2; Program-- the object is a parameter or variable of a

sequential program;

MODE3=3; Process entry-- the object is a parameter or local variable of an entry routine in a process type; that is, of an interface routine;

MODE4=4; Class entry-- the object is a parameter or local variable of an entry routine in a class type;

MODE5=5; Monitor entry-- the object is a parameter or local variable of an entry routine in a monitor type;

MODE6=6; Process-- the object is a parameter or permanent (global) variable of a process type;

MODE7=7; Class-- the object is a parameter or permanent (global) variable of a class type;

MODE8=8; Monitor-- the object is a parameter or permanent (global, shared) variable in a monitor type;

MODE9=9; Standard-- standard routine;

MODE10=10; Undefined-- used for error recovery;

MODE11=12; Manager-- not relevant to project;

MODE12=13; Manager entry-- not relevant to project.

Notice that the values of MODE11 and MODE12 do _not_ correspond to their names.

The section of the source code marked "COMMON TEST OUTPUT MECHANISM" contains the routines which produce an unformatted listing of the intermediate output code, perform simple pass initialization and termination functions, and handle page-buffer I/O with the intermediate code files. The "INPUT PROCEDURES" routines read from the code stream the number of instruction operands specified in the routine's name, as mentioned above. The WRITEx routines put into the output code stream the number of integers specified in the routine

name, where the first one is an operator and the others are operands. WRITEARG emits just one instruction operand (no operator). Both WRITEx and WRITEARG increment the location counter, LOCATION. Routines in "STACK PROCEDURES" are used to simulate the growth and reduction of the run-time stack. "BLOCK PROCEDURES" are called only when the compiler starts scanning a new routine or finishes the current one (that is, when ENTER1 and RETURN1 instructions, respectively, are encountered). The heap table management routines have been described above. BEGINPASS and ENDPASS take care of initialization and termination functions which are peculiar to pass six, such as allocating heap tables, starting the location counter at zero, and saving data in the PASSLINK record. In ENDPASS, PROGLENGTH is the length of the entire object file (header words, virtual code, and constants); CODELENGTH is the length of the virtual code proper; STACKLENGTH is the stack requirement for the anonymous initial process, hence the entire program; and VARLENGTH is the amount of variable space required by the initial process (that is, the length of the permanent variable space). The rest of MCPASS6 will be discussed in section 5.2.

4.3 PASS SEVEN

MCPASS7 is the code assembly phase of the compiler and is structurally much simpler than the other passes. In general terms, it performs the following functions:

1) Put the five words of header information into the object file;

2) Convert intermediate code operators to the even-number integer encoding which is intelligible to the interpreter;

3) Finish converting jump labels and routine labels to

displacements;

4) Insert stack requirement values into routine entry and process initialization instructions;

5) Remove error messages from the code stream and put them at the end of the source listing;

6) Put long constants into the object file immediately after the virtual code.

All five values for the header words are available from the PASSLINK. The translation of intermediate code operators is done on a simple one-for-one basis and operands are, for the most part, simply copied from the input file to the object file.

To convert jump and routine lables to displacements, the labels in the instructions are used as indices into JUMPTABLE and BLOCKTABLE (which were passed from MCPASS6 through the heap) to retrieve the location counter value associated with the label. The difference between that value and the current location counter value (which was emitted along with the instruction by pass six) is written out as the instruction operand.

The stack space required for each routine is in STACKTABLE. The interpreter needs the information contained there to perform run-time checks for stack overflow when routines are activated. When a routine entry or process initialization instruction is encountered, the routine's label is used as an index into STACKTABLE to fetch its stack requirement value. That value then becomes an instruction operand in the output stream.

If the source program contained any errors they were marked in previous passes by a MESSAGE(PASS, ERROR, LINE) instruction in the intermediate code. Pass seven removes MESSAGE instructions from the

code and prints the text for the error which is encoded in the ERROR operand. PASS is the number of the compiler pass which detected the error, and LINE is where it was found in the source code.

Once the virtual code proper has been generated, the constant pool is built by copying the contents of CONSTTABLE to the output file.

COMPILER MODIFICATIONS

Structurally, the MEPASCAL (MicroEngine PASCAL) compiler system of programs is directly analogous to the MCPASCAL system. The following correlation holds between the programs of the two compilers:

	MCPASCAL COMPILER	MEPASCAL COMPILER
Command File:	MCPASCAL.CSS	MEPASCAL.CSS
Driver:	MCPASCAL	MEPASCAL
Compiler Passes:	MCPASSx	MEPASSx
Pass 1, no source:	MCPASS1N	MEPASS1N
Code Mnemonics:	MNEM	MEMNEM.

Each program of MEPASCAL performs the same general function as its analog. In some cases (passes one and two, for example) the code is virtually identical.

Theoretically, passes one through five are independent of the compiler's target machine. Passes six and seven translate the machine-independent output language of pass five to the language of some particular machine. In view of this, it was originally thought that MEPASCAL could be built by rewriting just the last two passes of MCPASCAL so that they would produce Microengine machine code instead of virtual P-code. It seemed that localizing all the modifications would provide several advantages. It would eliminate the need for detailed knowledge about the workings of the other passes, design and programming errors introduced during the modification process would be confined to a small, familiar portion of the compiler, and changing just two passes in a straightforward manner seemed to offer fewer opportunities to make mistakes in the first place. The idea is simple, and it seemed workable. As it turns out, the first five passes can be slightly affected by the requirements of the target language, so they could not be used without change. However, modifications were

introduced only when it was felt that they were absolutely necessary, and so very little code was actually changed. Some changes could have been made either in pass six or one of the prior passes. As a rule, the choice was made to change pass six, even though it might have been "optimal" in some sense to do otherwise. This chapter describes the minor changes made to the first five compiler passes, and the major rewrite of pass six.

5.1 CHANGES TO PASSES ONE THROUGH FIVE

MEPASS1 (MEPASS1N) is the same as MCPASS1 (MCPASS1N) except for two changes in the PASSLINK record which is defined in the program prefix. Because of the way the Microengine XJP (case JumP) operator works, the number of words which case jumps will add to the constant pool must be exchanged between passes five and six. The field XJP_OFFSETS will hold that value. Pass five will also create and pass on a record of information concerning the concurrent/sequential program interfaces. INTERFACE will hold a pointer to that record. These two additions were also made to MCPASS2 to yield MEPASS2.

MCPASS3 incorporates the same PASSLINK changes made to the first two passes. The pass also uses a new output operator, DUPTOS2, (DUPlicate Top-Of-Stack word) when handling the input operator INIT_NAME1. The new operator is required because the Microengine code generated after INIT_NAME will, at run-time, pops the top-of-stack word, although that word will be needed later. DUPTOS2 will preserve the word for future use. See section 5.2.4.4 for more information.

The PASSLINK changes described above are included in pass four as well as the new input and output operators DUPTOS1 and DUPTOS2. DUPTOS1 appears as part of the CASE statement in the main loop of the

program. MEPASS4 also contains a modification to the BODY procedure. Before the change, the output generated in response to the BODY1 input operator varied, depending on whether the body (routine) being entered by the compiler was the initial routine of a system component. If it was **not** an initial routine, the output was the BODY2 operator and five operands, including the length of the routine's local variable area and the amount of stack space occupied by its parameters. For initial routines, the output was the same, except that the variable area and parameter length operands were always zero. Space for variables and parameters had already been allocated as a record in the data space of the anonymous initial process. For the Microengine, the two sizes need to be known in pass six for proper address displacement calculation. Since this is needed regardless of whether the routine is an initial one or not, BODY was changed so that the actual sizes are always included as BODY2 operands. More information is included in sections 5.2.4.1, 5.2.4.5, and 5.2.4.6.

Compared to the previous passes, MEPASS5 incorporates a large number of code modifications, including the new operators DUPTOS1 and DUPTOS2. In the other passes PASSLINK.INTERFACE is of type POINTER merely to reserve a fullword (32 bits) of storage, the amount of space required for **any** pointer. In pass five PASSLINK.INTERFACE is actually **used**, and so it must be declared as a pointer to some specific object type. Since it will point to an IFINFO (InterFace INFOrmation) record, it is declared to be of type IFPTR (InterFace PoinTeR). The IFINFO structure records the number of interfaces contained in the concurrent program (field INTERFACES) and the number of accessible process ENTRY routines in each one (field INTERFACESIZES). INTFACEPTR and INTFACE are instances of IFPTR and IFINFO, respectively. Figure 27 shows how

the objects and pointers of the PASSLINK relate to each other.

The instance of PASSLINK in pass five is the target of the pointer INTER_PASS_PTR. The procedure INITIALIZE handles initialization functions which are peculiar to this pass, before any intermediate code is read. In that procedure, INTER_PASS_PTR@.XJP_OFFSETS is set to zero to indicate that no case jumps (hence, no case jump offsets) have yet been found in the code. Space for the interface information record (type IFINFO) is allocated from the heap. Since no interfaces have been seen, the INTERFACES field in that record is initialized to zero. As the pass executes, it uses some space in the heap for temporary workspace. When the pass is finished, the workspace can be returned to the heap, but the interface information record must be retained for use by pass six. After allocating space for the IFINFO record, the extent of the heap is MARKed into INTER_PASS_PTR@.RESETPOINT. Temporary space is allocated beyond that point as needed, and at the end of the pass (in procedure EOM) that workspace (and only that workspace) is RELEASEd, leaving the IFINFO record intact. In preparation for pass six (procedure NEXT_PASS) the pointer to the IFINFO record (INTFACEPTR) is stored into one of the PASSLINK fields, making the record accessible to the next pass. During its execution, MEPASS5 recognizes the existence of an interface for a sequential program when it encounters a PROG_CALL1 operator with a nonzero operand. The operand is the length (in bytes) of the run-time jump table to be built before invoking the sequential program, and is a direct indicator of the number of process ENTRY routines to which the program will have access. The procedure PROG_CALL was modified to count the number of interfaces encountered, and the number of ENTRY routines made accessible in each one. That

information is saved in the IFINFO record and left in the heap for pass six. Refer to sections 5.2.4.4 and 5.2.4.8 for additional information.

The procedure CASE_LIST was modified to start the process of converting the virtual Concurrent Pascal CASE statement construct to a form suitable for the Microengine. The modification merely tallies the number of bytes by which the constant pool must be enlarged in order to accommodate the current CASE statement's list of jump offsets. The total amount by which all CASE statements enlarge the constant pool is saved in the PASSLINK record (XJP_OFFSETS) for use by pass six. More information on the Microengine CASE jump operator (XJP) and the way it is handled by MEPASCAL can be found in [MICR79, REGE79] and section 5.2.4.3.

Procedure ADDRESS contains a deletion. On the virtual Pascal machine, when a component variable is INITed, or passed as a parameter (access right) to another system component, its address is pushed onto the stack. That address is then incremented (FIELD operator) by the length of the permanent variables in the component. As a result, the top-of-stack word points to a word in the "middle" of the component variable's record, such that all the parameters (access rights) lie above it, and all the permanent variables below (see figure 25). This address will be loaded into the global base register (G) whenever a process executes the code of the component type. Since the Microengine hardware expects parameters and variables to lie above the global base (see figure 19), the FIELD operation for system components which is generated in the ADDRESS procedure must be dispensed with. Sections 5.2.4.4, 5.2.4.5, 5.2.4.6 also contain discussions related to this topic. At the time of writing it was noticed that a similar line

FIGURE 25. Relationship of global base to component variable address in the virtual Pascal machine. Compare with figure 19 which shows the analogous Microengine configuration.

of code appears at the end of the procedure SUB, and may also require deletion.

The last change to pass five involves subscript handling for strings. The virtual Pascal machine tacitly assumes that the underlying real hardware is byte addressable. That means that the address of a byte occupies the same amount of space (one word) as the address of a word. On the Microengine. not all addresses (pointers) have the same format. In fact, there are three different pointer formats [MICR79, REGE78, REGE79], although only two are of importance here; namely, word pointers (which occupy one word of memory) and byte pointers (which occupy two words). The only addresses generated by MEPASCAL which are not word pointers are pointers to elements of arrays of characters (strings). When it produces the address of an array element, pass six must know the virtual data type (BYTETYPE or not) of the pointer's target object. This information is available in MCPASS5 but not MCPASS6, so in MEPASS5 the object type (field KIND) was added as a fourth operand to the INDEX2 operator. MEPASS6 generates the correct address format based on the value of that operand. Section 5.2.4.2 contains a description of how that is accomplished.

All of the code changes made in the first five passes are shown in Appendix D.

5.2 CHANGES TO PASS SIX

Pass six contains more modifications than all of the previous passes combined. Not only were the changes numerous, but some were quite extensive. Besides changing the instructions to be generated, the format of the pass and its output structures were also changed.

5.2.1 Pass Output

As shown in figure 26, MEPASS6 takes its input from the heap and intermediate code file, and produces additional heap information and up to two intermediate code files.

The structures which pass six leaves in the heap are shown in figure 27. Space for the PASSLINK and IFINFO records was allocated in passes one and five, respectively. The fields in the PASSLINK record, which are not pointers, have the following meanings:

OPTIONS-- the set of compiler options in effect, as determined by pass one

LABELS-- the number of jump labels generated in pass two

BLOCKS-- the number of routine labels generated in pass four

CONSTANTS-- the calculated amount of constant pool space (in bytes) needed for long constants, as determined by pass three

XJP_OFFSETS-- the calculated amount of consant pool space (in bytes) needed to accommodate CASE jump offsets, as determined by pass five.

The meanings of the pointer fields is obvious from figure 27. In general terms, the IFINFO record contains information regarding the (possibly null) interface segment(s) in temporary file number four (figure 26). In particular, the INTERFACES field contains the number of interfaces found in the intermediate code, thus the number of interface segments generated. INTERFACESIZES is an array which contains the number of accessible process ENTRY routines for each interface. The interfaces are implicitly numbered in the order of their appearance in the intermediate code, and those numbers are used for the array index. In the TABLEPART record, SEGDISTANCE contains the

FIGURE 26. Input and output data for MEPASS6.
* The number of interfaces can be zero, leaving TEMP FILE 4 empty.

FIGURE 27. Structures left in the heap by MEPASS6. Although they are shown separately here, IFINFO, TABLEPART, and all of the TABLEs actually reside in the area marked "FREE HEAP SPACE". RESET points to the word just above IFINFO, since its value remains unchanged from pass five. TABLE layout is shown in detail in figure 24.

byte-measure equivalent of the header word in the _final_ concurrent code segment. See section 3.2.2 for a description of the header word. STACKLENGTH is the stack requirement for the concurrent program. It is a leftover from MCPASS6 and could have been eliminated. The meaning of the other fields is obvious from figure 27.

MEPASS6 leaves five tables in the heap: JUMPTABLE, CONSTTABLE, XJPTABLE, EXITICTABLE, and DATASIZETABLE. The first two tables are used in exactly the same way as in MCPASS6 (section 4.2). XJPTABLE contains case jump offsets. In the Pascal virtual machine these appear in the object code proper, but the Microengine requires them to be in the constant pool. Pass six removes the offsets from the code stream and places them in XJPTABLE in the same manner that constants are put into CONSTTABLE. Pass seven will combine the contents of the two tables to form a single constant pool for the concurrent segment. EXITICTABLE uses as many entries as there are routines in the concurrent segment. Routine numbers are used to index into the table, and the entries are the final-code EXIT-IC values for the routines. The meaning of the EXIT-IC field in Microengine code files is given in section 3.2.2. Routine numbers are also used as the index into DATASIZETABLE, and the entries are the DATASIZE values for the routines (see section 3.2.2).

In the conversion of MCPASS6 to MEPASS6 heap objects were added, deleted, and retained. An explanation of the reason each object was added or deleted will be helpful. The MCPASCAL compiler calculated PROGLENGTH, CODELENGTH, STACKLENGTH, and VARLENGTH in pass six so that pass seven could use those values as part of the five-word header for the object code file (see figure 3). Microengine object files (figures 6 and 7) do not require that header, so those fields could be deleted

from TABLEPART. (STACKLENGTH still remains through oversight.) In the same record, BLOCKTABLE was used to convert routine labels to the virtual addresses required by final-code CALL instructions. Microengine routine-invocation instructions use the routine labels themselves as operands, so the table was removed since the conversion became superfluous. The routine stack requirement values held in STACKTABLE are needed during MCPASS6 to help calculate the stack requirements of routines which call other routines (see section 4.2). MCPASS7 retrieves those values and inserts them into the code stream as operands to routine ENTER operators. Since Microengine object code does not require that information, STACKTABLE is not left in the heap by MEPASS6 and, as a consequence, does not appear in TABLEPART. It is used *during* pass six, as in MCPASS6, and then discarded.

The additional heap objects are INTERFACE and XJP_OFFSETS in PASSLINK, IFINFO, SEGDISTANCE in TABLEPART, XJPTABLE, EXITICTABLE, and DATASIZETABLE. Since all interface segments have the same structure (section 3.4.1) which is known to the compiler (section 5.2.4.8), the number of process ENTRY routines in an interface completely defines it. The number of interfaces and the size of each one is determined by MEPASS5 and left in the heap in the IFINFO record which is the target of the pointer INTERFACE in PASSLINK. MEPASS6 generates interface segments which are almost in final object-code form; *much* closer to that form than the concurrent segment. IFINFO is left in the heap for MEPASS7 so that when it scans the file of interface segments it will know the layout of each segment. Otherwise, there would be no way to determine, for example, the end of the constant pool and the location of the EXIT-IC field without backtracking in the file.

XJP_OFFSETS was calculated in MEPASS5 (section 5.1) and its

function is analogous to that of CONSTANTS— it tells MEPASS6 the number of table entries needed in XJPTABLE to accommodate case jump offsets which will be removed from the code. XJPTABLE is, essentially, an extension of CONSTTABLE, but case jump offsets cannot be intermixed with the long constants. The separation is required because an earlier pass calculated constant-mode displacements under the virtual-Pascal-machine assumption that case offsets would reside in the code portion of the object file. Willy-nilly inclusion of the offsets in the constant pool would invalidate those displacements in all but the most extraordinary circumstances. Pass seven will combine the contents of the two tables so that in the final code long constants will reside in the constant pool ahead of all the case offsets. This ensures that constant-mode displacements determined previously will still be valid, and yet allows pass six to easily compute the operand to the Microengine XJP (case jump) operator (see [MICR79] and section 5.2.4.3).

EXITICTABLE is passed from MEPASS6 to MEPASS7 since pass six cannot possibly know a routine's exit address before it has even started scanning its code. This situation arises because, as figure 7 indicates, the EXIT-IC field precedes the code of its routine in the object file. In MEPASCAL, pass six determines the EXIT-IC value as the routine is scanned, and pass seven places it in the proper position.

Although there is no usage-before-availability problem with object-code DATASIZE values (they are available as the VARLENGTH operand in ENTER1 instructions), emitting them directly into the code stream is undesirable. Pass seven is incapable of handling the unpredictable appearance of values which are neither operators nor operands, without extensive modification. It can, however, easily pull

the values out of the heap table and put them in place when the final code file is built.

MEPASS6 always generates a file (temporary file 2) of Microengine machine code (with a sprinkling of virtual operators) for the concurrent segment. Physically, it is a sequence of 16-bit integers, and consists of only the code for the routines which make up the concurrent program. The constant pool, procedure dictionary, and other non-code items do not appear in the file, although pass six takes into account their existence in the _final_ object code. Since the hardware requires routines (more specifically, the EXIT-IC fields of routines) to begin on word boundaries, NOP (No OPeration) instructions may appear _between_ adjacent routines for alignment purposes. Routines always end with an RPU instruction. There are three virtual (non-Microengine) instructions which can appear in the file in order to communicate information used by other parts of the compiler. MESSAGE_2 passes encoded error message data so that character error messages can be printed by pass seven. EOM_2 tells pass seven when it has reached the end of the file. NEWLIN_2 is a crutch used by MEMNEM to determine when the machine code corresponding to a new line of source code has been reached (Appendix A). MEPASS6 also inserts location counter values into the code stream (just as MCPASS6 does-- see section 4.2) as part of the mechanism for resolving jump displacements. All four of these items, which are extraneous to the final object code, will be removed in the next pass.

MEPASS6 generates an interface segment for every interface it finds in the intermediate code. All of the segments are placed end-to-end in a temporary file (file number four) as a sequence of 16-bit integers. Pass seven will concatenate the interface segments to

the concurrent code segment, pack operators and operands into words, and provide sufficient padding to ensure that each one begins on a disk block boundary, thereby constructing a single, complete object code file ready for execution on the Microengine. Section 5.2.4.8 contains an in-depth discussion of the mechanism in pass six which generates the segments.

5.2.2 Pass Structure

The structure of MEPASS6 is similar to that of MCPASS6 (section 4.2), although there are some differences. As mentioned in the preceding section, TABLEPART and PASSLINK were changed, and IFINFO was added to the package of objects to be left in the heap for pass seven. The input operators are the same except that DUPTOS1 has been added. The output operators for the Pascal virtual machine have all been replaced by Microengine operation codes. All Microengine codes are included, even though some will not be used, for the sake of completeness. Since they are defined in the CONST section of the program, the unreferenced opcodes do not add to the size of the program's object code. Three non-Microengine output operators (described in the preceding section) are also defined for handling error messages, marking the end of the concurrent code segment, and marking the start of each source line.

Several changes were made in the "COMMON TEST OUTPUT MECHANISM" portion of the pass because of the interface mechanism. Three variables (IFPAGE_OUT, IFPAGES_OUT, IFWORDS_OUT) were added to handle page-buffer output to the interface file. That output mechanism is exactly like the one used for intermediate concurrent code output, except for the file designation. The file identifiers are OUTFILE

(value=1) for the concurrent code file, and INTERFACEFILE (value=4) for the interface segment file. File 3 is the final object code file to be generated by pass seven. File initialization (INIT_PASS routine) and termination (NEXT_PASS routine) functions for all files used by pass six are performed at the same time in identical fashion.

Procedure WRITE IFL has been expanded so that output can be easily directed to either the concurrent file or interface file as necessary. When the pass starts (procedure BEGINPASS), a switch (GENNINGINTFAC) is set FALSE to indicate that an interface segment is not currently being generated (concurrent code is being generated) and that intermediate code output should be directed to the concurent file. When an interface is encountered, the switch is set TRUE (in procedure GEN_INTERFACE) to indicate the opposite state of affairs. After the interface segment has been completely generated, the switch is returned to the FALSE setting. All output to files occurs through procedure WRITE_IFL, and on every invocation it tests GENNINGINTFAC to determine the file to which the output should be directed. Duplicate sections of code handle buffer management and output to each file. The duplicate code is a violation of good programming practice, and could be easily eliminated by the use of array variables. However, at the time this mechanism was designed, the duplicate code arrangement was an easier concept to deal with, and the amount of overhead is not alarming.

The "OUTPUT PROCEDURES" section of code contains a number of modifications which are due to two causes. First, unlike the Concurrent Pascal virtual machine. Microengine operators use operands of various lengths. In fact, for some instructions, a single operand can have either of two lengths, depending on its value. The

Microengine instruction set is described in appendix B.5 of [MICR79].
Heterogeneous lengths do not affect the intermediate code produced by
MEPASS6 since it generates operators and operands which are all 16-bit
integers, but updating the final-code location counter becomes more
complicated. The new mechanism uses a new global type (TYPEOFCODE),
the values of which have the following meanings:

OPTR-- The output data item is a Microengine instruction
OPeraToR. Length is always one byte;

UB-- An Unsigned Byte operand. Length is always one byte;

SB-- A Signed Byte operand. Length is always one byte;

DB-- A "Don't care" Byte operand. Length is always one byte;

B-- A "Big" operand. Length is one or two bytes, depending on the
value;

W-- A Word operand. Length is always two bytes;

NOTME--The output data item is NOT a MicroEngine code item and
will be removed by pass seven, so length is zero.

The new procedure UPDLOC (UPDate LOCation counter) ensures that the
location counter (variable LOCATION) gets incremented correctly for
each type of intermediate code item. Every time a WRITEx or the
WRITEARG routine emits an item of output, a call is made to UPDLOC.
The numerical value of the item and its TYPEOFCODE value are provided
as parameters so that UPDLOC can determine the item's size, and
increment LOCATION accordingly. The WRITEx and WRITEARG parameter
lists were modified to accept operand TYPEOFCODE values so they can be
passed on to UPDLOC. Notice that the NOTME value provides a way to
use the standard output procedures for generating intermediate code
instructions which will not appear in the final code, without also
causing an increment of the location counter.

The second cause for modification concerns the generation of interface segments. For the sake of consistency, the interface-generation routine emits code by invoking the same output procedures as the rest of the program. Without modification of the procedures, generation of an interface would cause the concurrent segment's location counter to be incremented. This is an undesired side effect since locations within a segment must be relative to the beginning of that segment. The remedy for that side effect is to test the GENNINGINTFAC switch before calling UPDLOC. If the item just emitted was sent to the interface file the switch will be TRUE, and LOCATION will not be affected. Otherwise, it will be updated as described above.

The function DISPL (DISPLacement) is designed to return the displacement of a variable from some base address. It had to be entirely rewritten because of addressing differences between the two target machines. In the virtual Pascal machine, displacements are positive (for routine parameters) or negative (for local variables) even-byte displacements from a data-space base register (either G or B-- see figure 28). Constants are addressed relative to the starting address of the constant pool which is calculated when the concurrent program starts execution. The Microengine addresses both parameters and variables with positive word offsets from registers B or MP, and constants are addressed by word offset from the start of the segment. The virtual machine can refer to the fields of the markstack by addressing the words with offsets 0, 2, 4, 6, and 8 from, for example, the local base. The first variable word after the markstack has displacement 10, and the first parameter has displacement -2. In contrast, the Microengine hardware does not allow access to the

VIRTUAL MACHINE   MICROENGINE

FIGURE 28. Variable and parameter addressing mechanisms of the two target machines. The local stack areas are shown for a routine which has three words of parameters and four words of local variables. Addressing displacements appear to the left of each stack. In general, the Microengine markstack fields are not accessible by using the displacement mechanism.

markstack, except by the LSL (Load Static Link) instruction [MICR79]. The hardware automatically compensates for the markstack which is between the base address and the first variable, so that the first variable word has offset 1. The offset of the first parameter depends on the number of variables, since the variables separate the parameters from the base address. For example, if a routine has four words of local variables and three words of parameters, the parameter words have offsets 5, 6, and 7. The MEPASS6 version of DISPL calculates displacements in the following way. The constant block will immediately follow the segment header word in the object file, so a constant's displacement is its word displacement within the constant block plus one word. In "procedure", "program", and "entry" addressing modes (section 4.2) the displacement calculation takes into account the following factors:

1) The displacements for variables must be positive, not negative;

2) Parameters must be addressed as if they are an extension of the variables area;

3) The first two words of variables are reserved for the source line number and old global base register values, since the Microengine has no markstack fields for them specifically;

4) The Microengine uses word, not byte, displacements.

In "process", "class", and "monitor" modes, the calculation is the same, except that space for the two reserved words is of no concern since no space for them exists in the component's permanent variable record.

The displacement calculation for a "process entry" reference is no more complex than the others, but the run-time structure of the

concurrent/sequential interface, on which it is based, is somewhat complicated. Figure 29 shows the arrangement of the stack after a sequential program has invoked a process ENTRY routine and it has begun executing. That stack state was reached in the following way. The sequential program code pushed enough "blank" space onto its operand stack to hold the value which the ENTRY function will eventually return. It also pushed the parameters which the ENTRY routine requires, followed by the index of the prefix routine which the program sought to invoke. The index is itself a parameter, and it is critical that it always be pushed last onto the stack. The program then executed an external routine call instruction, causing the interface markstack to be constructed by the hardware. The interface routine pushed its local variable with offset 1 (the prefix index parameter) onto the stack for use as the selector value for its CASE statement (see figure 17). Execution of the case jump (XJP instruction) consumed the top-of-stack operand, and execution of the case code invoked the process ENTRY routine, causing the allocation of its local variable space on the stack and the construction of its markstack. As in other routines, local words one and two are reserved for the source line number and old global base. The displacement calculation for process ENTRY routine _variables_ is the same as for class and monitor ENTRY routines. _Parameters,_ however, are different since, as the reader can see in figure 29, the interface markstack sits between the local variables and parameters. Conceivably, there should never be anything _but_ the stack marker in that position, and parameter displacements could be calculated as

(BYTE_OFFSET + TWOWORDS + FOURWORDS + ONEWORD) DIV WORDLENGTH

where TWOWORDS accounts for the two reserved local words, FOURWORDS is

SEQUENTIAL
PROGRAM
OPERAND
STACK

BUILT BY
SEQUENTIAL
PROGRAM

FUNCTION VALUE

PARAMETERS

PREFIX INDEX

BUILT DURING
INTERFACE
INVOCATION

INTERFACE
MARKSTACK

BUILT DURING
ENTRY ROUTINE
INVOCATION

PROCESS
ENTRY ROUTINE
VARIABLES

PROCESS
ENTRY ROUTINE
MARKSTACK

FIGURE 29. Configuration of the run-time stack after the execution of a process ENTRY routine (a function, in this case) has begun. The interface routine has no local variables.

the size of the interface markstack, and ONEWORD accounts for the word occupied by the prefix index which is an "extra" parameter. That calculation was _not_ used, however, since it seems to depend too heavily on a favorable arrangement of the stack. Also, the machine instruction set includes a number of operators which can access the data spaces of any markstack in the chain of static links. When one of these "intermediate" instructions is used, variables are addressed relative to the markstack to which they are normally considered local. The interface routine invokes the process ENTRY routine with a CXL instruction which makes the static link field in the ENTRY routine's markstack point to the interface's markstack. By directing the "intermediate" addressing operators to traverse just one static link (operand value of 1) the interface parameters can be accessed as if they are part of the ENTRY routine's local variable space. The significance of all this is that the displacement calculation has to only take into account the single word occupied by the prefix index. Section 5.2.4.8 contains more information on the operation of the interface mechanism.

The procedures BEGINPASS and ENDPASS carry out initialization and termination functions specific to this pass of the compiler. As part of the pass setup, the sizes of all the interfaces are copied from the PASSLINK record to an array which is easier to access. Heap space is allocated for the five tables which will be left for the next pass. The current extent of the heap is MARKed, and then a temporary table (STACKTABLE) is given space. Indices into the constant and case jump offset tables are initialized to indicate empty tables. The total size of the constant pool is calculated as the sum of the constants and case jump table sizes since they will be combined in the object code

file. The location counter is initialized so that it "points" to the first byte of code proper to be generated. That byte follows the segment header word and the constant pool. ENDPASS RELEASEs the heap space which was used by STACKTABLE since that table will not be used by pass seven. Heap space is then allocated for a TABLEPART record and the table-identifying pointers are stored into it. The address of the last meaningful word in the concurrent segment can be calculated from the current location counter value and the number of routines in the segment. At the end of the pass, LOCATION is the segment-relative address of the word after the last _instruction_ in the segment. As figure 7 shows, only the procedure dictionary comes between the last routine (routine one) and the final word. There are as many entries in the procedure dictionary as there are routines in the segment, and each entry takes one word (two bytes), so the segment-relative address of the last word is

$$LOCATION + (2 * ROUTINES).$$

Four other routines were added to pass six. PICK_PUSHCONST chooses one of the three Microengine instructions which push integer constants onto the stack, given the value of the constant to be pushed. The difference between the instructions is their length, and it seemed worthwhile to make the small effort required to make optimal use of them to help reduce the size of the object code. GEN_SAVELINE_SAVEGBASE generates the prologue for routines. It generates code which, at run time, will save the source code line number and the contents of the global base register in the local variable space of the new routine. In every routine local word one is reserved for the line number, and local word two is for the old global base register. GEN_EXIT generates epilogue code which is the same for

several routine modes-- restoration of the old global base register
and the actual return to the calling routine. GEN_INTERFACE generates
interface segments, and is presented in detail in section 5.2.4.8.

Six procedures were deleted and rewritten as in-line code. Each
of them was relatively small and was called from only one point in the
program, so in the interest of comprehensibility they were moved to
the places where they more logically belong. ENTERBLOCK and EXITBLOCK
were moved into the ENTER1 and RETURN1 (respectively) cases of the
SCAN routine. The comparison procedures COMPAREWORD, COMPAREREAL, and
COMPARESET were rewritten in the COMPARE1 case of SCAN, and
COMPARESTRUCT was inserted into the COMPSTRUC1 case.

5.2.3 Straightforward Instruction Conversions

Simply stated, the major concern of this project is to change the
representation of Concurrent Pascal programs without changing their
semantics. In order to do that, each virtual instruction which can be
output by MCPASS7 was analyzed to determine what changes it causes in
the state of the virtual machine. For each one, a sequence of
Microengine instructions which would mimic those changes was chosen
for generation in place of the virtual instruction. The definitive
characterization of virtual machine state changes is the PDP-11/45
interpreter which realizes them. A source listing of the interpreter,
written in PDP-11 Assembly code, was available [ZEPK74]. The listing
is commented with the same code written in a psuedo-Pascal/assembler
"language" which provides a certain (limited) amount of insight into
the function of each virtual instruction. Since, at least
superficially, both machines have a stack architecture, some
instructions could be substituted very simply. Those virtual

instructions and their substitutes are presented here.

The instructions which operate on integer (or word-length) and real operands were among the simplest to convert since they perform virtually identical actions on both machines. Figures 30 and 31 list the MCPASS7 output instructions, their actions, and the equivalent Microengine code sequence for word and real operators. Both machines use 16-bit integers, where the low-address byte contains the low-order bits-- PDP-11 format. Real numbers are handled differently on each machine. The virtual machine uses an eight-byte format, whereas the Microengine uses the PDP-11 single precision floating point format which takes only four bytes [DIGI76].

Figure 32 is a list of the instructions for set and structure operations. Sets for the virtual machine have a fixed length of 16 bytes (128 bits). Set members are numbered from 0 to 127, and membership is represented by turning the corresponding bit "on". The Microengine hardware supports sets of variable length up to 255 bytes (4080 bits), but fixed-length sets of 128 elements are used in this project for the sake of simplicity. On the stack, Microengine sets consist of two parts. The top-of-stack word is an integer which is the number of words in the set proper, which is next on the stack. When it is on the stack, the combination of length word and membership bits is referred to as the "set". When they are not on the stack, sets consist of only the membership bits. The Microengine has no hardware instructions specifically for pushing (popping) sets and automatically appending (discarding) the length word. The instructions LDC (LoaD multiword Constant), LDM (LoaD Multiple words), and STM (STore Multiple words) can be used to push and pop the membership bits, but the length word must be handled by additional instructions. The length

VIRTUAL OPERATOR	RESULT OR ACTION	MICROENGINE INSTRUCTION(S)
ABSWORD	absolute value of TOS	ABI
ADDWORD	sum of TOS, TOS-1	ADI
ANDWORD	logical AND of TOS, TOS-1	LAND
CONVWORD	real-value equivalent of TOS	FLT
COPYWORD	store TOS indirect through TOS-1 *	STO
DECRWORD	decrement word indirect through TOS-1 *	DUP1 LDM (1) SLDCO1 SBI STO
DIVWORD	integer quotient of TOS-1/TOS	DVI
EQWORD	boolean value of (TOS-1 = TOS)	EQUI
GRWORD	boolean value of (TOS-1 > TOS)	LEQI LNOT
INCRWORD	increment word indirect through TOS-1 *	DUP1 LDM (1) SLDCO1 ADI STO

FIGURE 30. Concurrent Pascal word (integer) virtual instructions and their Microengine equivalents, where TOS is an integer which is the top-of-stack item and TOS-1 is the integer which was pushed just before TOS. Execution of an instruction causes the stack operands to be popped from the stack. The result (if any) is then pushed onto it. Instructions marked '*' do not push a result.

VIRTUAL OPERATOR	RESULT OR ACTION	MICROENGINE INSTRUCTION(S)
LSWORD	boolean value of (TOS-1 < TOS)	GEQI LNOT
MODWORD	TOS-1 MOD TOS	MODI
MULWORD	product of TOS, TOS-1	MPI
NEGWORD	2's complement of TOS	NGI
NEWORD	boolean value of (TOS-1 <> TOS)	NEQI
NGWORD	boolean value of (TOS-1 <= TOS)	LEQI
NLWORD	boolean value of (TOS-1 >= TOS)	GEQI
ORWORD	logical OR of TOS, TOS-1	LOR
PREDWORD	TOS minus 1	SLDC01 SBI
SUBWORD	TOS-1 minus TOS	SBI
SUCCWORD	TOS plus 1	SLDC01 ADI

FIGURE 30. (continued from previous page)

VIRTUAL OPERATOR	RESULT OR ACTION	MICROENGINE INSTRUCTION(S)
ABSREAL	absolute value of TOS	. ABR
ADDREAL	sum of TOS, TOS-1	ADR
COPYREAL	store TOS indirect through TOS-1 *	STM (2)
DIVREAL	quotient of TOS-1/TOS	DVR
EQREAL	boolean of (TOS-1 = TOS)	EQUREAL
GRREAL	boolean of (TOS-1 > TOS)	LEQREAL LNOT
LSREAL	boolean of (TOS-1 < TOS)	GEQREAL LNOT
MULREAL	product of TOS-1, TOS	MPR
NEGREAL	negation of TOS	NGR
NEREAL	boolean of (TOS-1 <> TOS)	EQUREAL LNOT
NGREAL	boolean of (TOS-1 <= TOS)	LEQREAL
NLREAL	boolean of (TOS-1 >= TOS)	GEQREAL
SUBREAL	TOS-1 minus TOS	SBR
TRUNCREAL	TOS, truncated and converted to integer	TNC

FIGURE 31. Concurrent Pascal virtual instructions for real values and their Microengine equivalents. TOS and TOS-1 have the same meanings as in figure 30, but are real values instead of integers. Instruction marked '*' does not push a result onto the stack.

VIRTUAL OPERATOR	RESULT OR ACTION	MICROENGINE INSTRUCTION(S)
	--- SET INSTRUCTIONS ---	
ANDSET	intersection of TOS, TOS-1	INT
BUILDSET	TOS set with a new member; New member number is given by integer TOS.	SLDC00 LDCB (127) CHK DUP1 SRS UNI
COPYSET	store TOS set indirect through TOS-1 *	FJP (0) STM (8)
EQSET	boolean value of set compare (TOS-1 = TOS)	EQUPWR
INSET	boolean value of test for inclusion of TOS-1 member number (integer) in TOS set	INN
NESET	boolean value of set comparison (TOS-1 <> TOS)	EQUPWR LNOT
NGSET	boolean value of subset test (TOS-1 <= TOS)	LEQPWR
NLSET	boolean value of superset test (TOS-1 >= TOS)	GEQPWR
ORSET	union of TOS, TOS-1 sets	UNI
SUBSET	TOS-1 set less members of TOS-1 set	DIF
	--- STRUCTURE INSTRUCTIONS ---	
COPYSTRUC	move TOS@ to TOS-1@ *	MOV (size)
EQSTRUC	boolean value of (TOS-1@ = TOS@)	EQUBYT (size)
GRSTRUC	boolean value of (TOS-1@ > TOS@)	LEQBYT (size) LNOT
LSSTRUC	boolean value of (TOS-1@ < TOS@)	GEQBYT (size) LNOT
NESTRUC	boolean value of (TOS-1@ <> TOS@)	EQUBYT (size) LNOT
NGSTRUC	boolean value of (TOS-1@ <= TOS@)	LEQBYT (size)
NLSTRUC	boolean value of (TOS-1@ >= TOS@)	GEQBYT (size)

FIGURE 32. Set and structure instructions. "TOS@" means "the object to which the top-of-stack word points". Similarly for "TOS-1@". "*" means nothing is left on the stack.

word is discarded during the execution of the ADJ (ADJust set length) instruction which forces the set to a length determined at compile time. However, in this project an FJP (False JumP) to the next instruction is used instead (see figure 32, COPYSET). The lengths of sets never really need to be adjusted since they will be fixed, and the FJP is presumed faster. Notice that structure operators use pointers as stack operands rather than the structures themselves.

The virtual machine has several instructions for process control. They are ATTRIBUTE, CONTINUE, DELAY, EMPTY, START, STOP, and WAIT [BRIN75, BRIN77, ZEPK74]. Except for EMPTY, these all interact with the kernel, either by issuing a call to a kernel routine, or by accessing a fixed location known to the kernel and interpreter. Since the kernel for the Microengine has not yet been designed, these instructions have been "commented out" of the code and must be modified at some later time. EMPTY returns a boolean value indicating whether a QUEUE variable is empty, without any kernel interaction. The code sequence SLDC00 EQUI emulates this instruction on the Microengine.

The three heap control instructions SETHEAP, NEWINIT, and NEW have not yet been changed, but rather, "commented out". SETHEAP initializes the heaptop pointer so that it points to the bottom of the heap space, thus making the heap "empty". The equivalent Microengine code is:

```
SLDC02      push heap pointer register (SPLOW) number
SWAP        reverse TOS and TOS-1 words
SPR         put TOS value in heap pointer register.
```

NEW allocates heap space for a new object and returns a pointer to it. Figure 33 shows Microengine code which performs a similar function. There are two unresolved issues here. The obvious one concerns the

```
        SLDC04              stack pointer register number (SP register)
        LPR                 push stack pointer
        LDCB (100)          keep 100 bytes between heap and stack
        SBI                 stack grows toward low addresses
        SLDC02              heap pointer register number (SPLOW register)
        LPR                 push heap pointer
        push object's length   (choose one of 3 possible instructions)
        ADI                 calculate heap pointer after allocation
        LEQI                test (SP - 100) <= (SPLOW + object length)
        FJP (ok)
        heap limit error code

ok:     SLDC02
        LPR                 push heap pointer
        STO                 store (indirect) pointer to object into
                                pointer variable

        SLDC02              SPLOW register number, used with SPR below
        SLDC02
        LPR                 push heap pointer
        push object's length
        ADI                 calculate new heaptop
        SPR                 set heap pointer register to new heaptop
```

FIGURE 33. Microengine code to emulate the NEW Pascal virtual instruction. This code performs three general functions: 1) check for collision of heap and stack, 2) store the pointer to the new object in the object's pointer variable, and 3) reset the heap pointer to the next word of free space in the heap.

nature of the "heap limit error code". Whether it should be a call on a well-known kernel routine or in-line code has not been determined. The second issue was discovered at the time of writing, and is centered on the problem of ensuring, at run time, that the stack and heap will not collide. The virtual instruction NEW has an operand which includes the stack requirement of the routine in which it appears. During execution, it compares the extent of the heap after the contemplated allocation, and the maximum extent of the stack during the routine in order to determine if a collision is possible. The MCPASCAL compiler generates the stacklength operand in pass seven by referring to the routine's entry in STACKTABLE. In MEPASCAL, however, STACKTABLE is removed by pass six, so pass seven does not have routine stack requirements available to it. The code in figure 33 assumes (rather naively) that the routine's stack will not grow by more than 100 bytes from its current position. This is a rather arbitrary "safety factor" and not nearly as desirable as the original mechanism. If possible, STACKTABLE should be passed on to MEPASS7 in order to restore that mechanism. NEWINIT is the same as NEW, except that the newly-allocated heap space is initialized to zero. Its Microengine equivalent is shown in figure 34. Besides the two issues mentioned above, an additional concern is the large amount of space consumed by the initialization code.

Sixteen other virtual instructions and their transformations will be described here. They are:

COPYTAG	EOM	FALSEJUMP	FIELD
FUNCVALUE	INITVAR	IO	JUMP
MESSAGE	NEWLINE	NOT	POINTER
POP	RANGE	REALTIME	VARIANT.

There are two Microengine instructions which are equivalent to FALSEJUMP-- FJP and FJPL. The difference is that FJP uses a signed

```
        SLDC04
        LPR             push stack pointer
        LDCB (100)      stack-heap separation
        SBI
        SLDC02
        LPR             push heap pointer
        push object's length
        ADI
        LEQI            test (SP - 100) <= (SPLOW + length)
        FJP (ok)
        heap error code

ok:     SLDC02
        LPR             push heap pointer
        STO             store pointer to new object in pointer variable

        SLDC02
        LPR             push object's base address; will be used
                            later to scan through new object
        DUP1     .      get copy of pointer to new object
        SLDC02          SPLOW register, used with SPR below
        SLDC02
        LPR             push heap pointer
        push object's length
        ADI             calculate new heaptop
        SPR             set heap pointer to new heaptop

next:   SLDC02
        LPR             push new heaptop
        GEQI            object scan pointer >= new heaptop?
        LNOT
        FJP (exit)
        SLDC02
        ADI             point to next word in object
        DUP1            copy object scan pointer for use
                            by GEQI instruction in next iteration
        DUP1            push another copy for STOring the initial value
        SLDC00          initial value to be inserted in object
        STO             set a word in the object to zero
        UJP (next)      prepare to initialize next word in object

exit:   FJP (0)         pop object scan pointer
```

FIGURE 34. Microengine code to emulate the Pascal virtual instruction NEWINIT, which allocates heap space for a new object and initializes the object to zero. Compare with figure 33. After the space has been allocated and the heap pointer updated, the heap pointer points to the word which immediately follows the new object.

byte operand whereas FJPL (False JumP Long) uses a signed word operand. While FJP takes less space, its range is also shorter. The longer instruction is used for all false jumps, even though it might be suboptimal, in order to keep code assembly simple. If both instructions were used, code addresses could not be known in pass six as they are now. JUMP also has two equivalents (UJP-- Unconditional JumP, and UJPL-- Unconditional JumP Long), and the longer instruction is always used, for the reasons just mentioned. FIELD'S equivalent is INC (INCrement field pointer), except that the offset operand must be converted to word measure. LNOT (Logical NOT) performs the same function (boolean negation) on the Microengine as the virtual machine instruction NOT.

The Microengine has no direct equivalent of POP, but FJP and NFJ (Not equal False Jump; that is, jump if equal) can be combined into a suitable substitute. POP has an operand which is the number of bytes to be discarded from the stack. An odd number of bytes is never popped since the stack consists only of whole words. Since no Microengine instruction merely discards a variable portion of the stack without side effects, the compiler generates a variable number of NFJ(0) instructions and possibly one FJP(0) instruction. NFJ compares the two top-of-stack words as integers and the jump is made if they are equal. In either outcome the two words are discarded from the stack. FJP tests the top-of-stack word, discards it, and jumps if it is FALSE. Both instructions use a signed distance operand to calculate the jump address. If the distance is zero, the next sequential instruction will be executed, regardless of the test results, and the stack operand(s) will be popped from the stack. Thus, NFJ(0) is equivalent to POP(4), and FJP(0) to POP(2). Presumably, the

number of bytes to be popped at any given time will be relatively small, so the compiler generates enough NFJ(0) instructions such that after their execution either no words or one word remains to be popped. If one word remains, an FJP(0) is generated, otherwise not.

RANGE checks the top-of-stack integer against its operands to ensure that the integer is within the range which they specify. If it is not within the range, a run-time error occurs. The Microengine operator CHK (CHecK) does the same thing (including triggering a run-time error) except that the range-specifying operands are taken from the stack, and the integer to be checked is the TOS-2 item. IO, POINTER, and REALTIME interact with the kernel, so their equivalent code has not been finalized. IO and REALTIME must call the kernel in all cases. POINTER ensures that pointer variables contain some nonzero value. A run-time error results if the value is zero, and the Microengine equivalent of the error code is undetermined at this time.

VARIANT also uses error code which is as yet undetermined, and will also require changes to one of the earlier passes in MEPASCAL. Figure 35 shows how the instruction appears in the final code. VARIANT assumes that the address of the variant record is the top-of-stack word. Its first operand is the field offset of the tag field, and the second is the set of tags (always one word long) for which access is being requested. The tag field contains the bit number of the tag which is currently "legal". During instruction execution, the bit in the request set whose number is given in the tag is examined. If that bit is not "on", a run-time error occurs, otherwise execution proceeds. Notice that VARIANT does not alter the stack, leaving the record address in place. Use of the bit number for the

```
0001 TYPE
0002 IDENTIFIER = ARRAY[1..12] OF CHAR;
0003              "0        1        2        3"
0004 ARGTAG = (NILTYP, BOOLTYP, INTTYP, IDTYP,
0005              "4        5        6        7"
0006           PTRTYP, XTYP,    YTYP,    ZTYP);
0007 RECTYPE = RECORD
0008           FIRST: INTEGER;
0009           CASE VRNTTAG: ARGTAG OF
0010              "0        1"
0011           NILTYP, BOOLTYP:          (BOOL: BOOLEAN);
0012              "2        5    6        7"
0013           INTTYP, XTYP, YTYP, ZTYP: (INT:  INTEGER);
0014              "3"
0015           IDTYP:                    (ID:    IDENTIFIER);
0016              "4"
0017           PTRTYP:                   (PTR:   REAL)
0018           END;
0019
0020 VAR   VREC: RECTYPE;
0021
0022 BEGIN
0023 VREC.VRNTTAG := INTTYP;     "TAG IS 2"
0024 VREC.INT := 100;           "USING BITS 2, 5, 6, 7"
0025
0026 VREC.VRNTTAG := YTYP;      "TAG IS 6"
0027 VREC.INT := 200;           "USING BITS 2, 5, 6, 7"
0028
0029 VREC.VRNTTAG := PTRTYP;    "TAG IS 4"
0030 VREC.PTR := 300.0          "USING BIT 4"
0031 END.
```

```
7777777777777777777777777777777777777777777777777777777777777777777
7              MCPASCAL INTERMEDIATE CODE PASS 7                   7
7777777777777777777777777777777777777777777777777777777777777777777
JUMP(2)
            BEGNPRCS(22)
LINE   23   GLOBLADD(-16)   FIELD(2)   PUSHCONS(2)   COPYWORD
LINE   24   GLOBLADD(-16)   VARIANT(2,9984)   FIELD(4)
       PUSHCONS(100)   COPYWORD
LINE   25
LINE   26   GLOBLADD(-16)   FIELD(2)   PUSHCONS(6)   COPYWORD
LINE   27   GLOBLADD(-16)   VARIANT(2,9984)   FIELD(4)
       PUSHCONS(200)   COPYWORD
LINE   28
LINE   29   GLOBLADD(-16)   FIELD(2)   PUSHCONS(4)   COPYWORD
LINE   30   GLOBLADD(-16)   VARIANT(2,2048)   FIELD(4)
LINE   31   CONSTADD(16)    PUSHREAL   COPYREAL   ENDPRCS
```

FIGURE 35. Example of the VARIANT virtual machine instruction, the source code which produces it, and the final code which surrounds it. Tag bit numbers are shown above the identifiers. Some of pass seven's output is not shown here since it is irrelevant.

current tag value is apparently desirable when running on a PDP-11
since hardware shift instructions can then be used advantageously
[DIGI76, ZEPK74]. For the Microengine, however, it will be necessary
for the tag field to hold the value of the current tag bit, not its
number. That change has not been made, and the compiler pass which
must be modified has not yet been identified. Once that change has
been made, the equivalent Microengine code would be:

```
    DUP1                save copy of record address
    IND (word displacement)   push current tag value
    LDCI (tag set)  push tags for which access is requested
    LAND                logical AND the two words together
    SLDC00              push word with all bits "off"
    EFJ (ok)            if request ok, TOS-1 will have an "on" bit
    variant error code
ok: next instruction.
```

If the NUMBER compiler option is specified, MCPASCAL generates
NEWLINE instructions at points in the object code which correspond to
the beginning of source lines. The NEWLINE instruction stores the line
number operand into a word of the local markstack. However, the
Microengine has no markstack space for the line number, and in
addition, its markstack fields are not generally accessable (section
5.2.2). The line numbering mechanism has been implemented by
specifically reserving the first word of local variable space in every
routine for the current line number. The Microengine code equivaluent
to NEWLINE consists of two instructions-- one of the "push constant"
instructions, and STL(1). The routine DISPL in MEPASS6 adjusts all
references to program variables to take into account the reserved
word. In addition to the push and store instructions, pass six also
puts the NEWLIN_2 virtual instruction into the code stream. This is a
crutch for the intermediate code mnemonics program MEMNEM. NEWLIN_2
makes it possible for MEMNEM to identify the beginning of each source
line without resorting to a lookahead technique (see Appendix A).

No Microengine code is generated when MESSAGE1 and EOM1 are encountered in pass six. They are merely inserted into the intermediate code as virtual instructions for use by pass seven.

In general terms, the virtual instruction FUNCVALUE pushes onto the stack enough space to hold the value which will be returned by a function which is about to be called. It pushes either one or four words, depending on the type of the return value. However, there is a twist. Class and monitor entry routines expect the address of the component's permanent variables to be the first "parameter" pushed on the stack. But the code generated by the compiler carries out the actions:

    1) push the address of the permanent variables record;

    2) push space to receive the function value;

    3) push parameters;

    4) call the entry routine.

If the function-value space is placed directly on the top of the stack, the permanent variables record address will not be in the proper position. The function-value space should be placed between the two top-of-stack words. FUNCVALUE performs this stack-space insertion for class and monitor function entries. There are four equivalent Microengine code sequences. As FUNCVALUE comes into MEPASS6, it has a mode operand and a type operand. If the mode is "process entry" or "procedure", one or two SLDC00 instructions are generated to push space at run time for word-type or real-type function values. If the

mode is "class entry" or "monitor entry", the code sequences are:

WORD-TYPE	REAL-TYPE
SLDC00	SLDC00
SWAP	SWAP
	SLDC00
	SWAP.

At run time, these instructions will "insert" one or two words of zero between the permanent variables record address and the next word on the stack.

Finally, INITVAR and COPYTAG were commented out of pass six since they are never produced by pass five of the concurrent compiler. They are presumably produced by the sequential compiler. Apparently, passes six and seven of the two compilers were so similar that they were combined, and are used by both translation systems.

5.2.4 "Difficult" Instruction Modifications

The virtual instruction transformations described in the preceding section were relatively easy to make since they are, for the most part, independent of the underlying hardware architecture. The transformations described here were more difficult since these instructions make use of some architectural assumptions.

5.2.4.1 Value and Address PUSH Instructions

Pass five generates four instructions for pushing values and addresses onto the operand stack. They are PUSHCONST1 (PUSH the value of a CONSTant), PUSHVAR1 (PUSH the value of a VARiable), PUSHIND1 (PUSH a value, INDirect), and PUSHADDR1 (PUSH an ADDRess). Pass six analyzes these operators and their operands in order to select the appropriate final-code instructions from nearly a dozen possibilities.

The operand for PUSHCONST1 is the 16-bit value of the constant to be pushed onto the stack. For the virtual machine there is just one instruction for pushing short (immediate) constants, but there are three for the Microengine. Based on the value of the constant, MEPASS6 selects the hardware instruction which will take the least possible amount of code space.

PUSHVAR1 has three operands: the virtual data type of the variable, the addressing mode, and its displacement from either the local or global base register of the virtual machine. If the variable is a word or less, its type will be WORDTYPE and it can be pushed directly. Otherwise, it must be pushed indirectly. The choice of the direct-push instruction to be generated depends on the variable's mode, and is made by the MEPASS6 routine PUSHVALUE, discussed below. The Microengine code sequence for an indirect push consists of two instructions; one to place the address of the variable on the stack, and another to actually use that address and replace it with that location's content. The procedure PUSHADDRESS chooses the first of these, from among several possibilities, based on the addressing mode, and PUSHINDIRECT picks the second, based on the virtual data type.

When the PUSHADDR1 and PUSHIND1 intermediate-code instructions are encountered by pass six, their operands are passed directly to the routines PUSHADDRESS and PUSHINDIRECT, respectively. Those two routines (described below) determine and generate the corresponding final code.

The procedure PUSHVALUE uses a variable's address mode to choose among three Microengine instructions which push onto the stack the value of a word in some data space. If the mode is "procedure", "class entry", or "monitor entry" then the word will be (at run time) in the

data space which is local to the routine active at that moment. The
Microengine instruction LDL (LoaD Local word) is generated for those
modes. The operand for LDL is the variable's displacement within the
local data space, and is calculated by the routine DISPL, described in
section 5.2.2. "Process", "class", and "monitor" modes indicate that
the variable is a permanent variable of a system component, and so it
will be addressable by a displacement from the global base register.
(At run time the global base register will have been set so that it
points to the _correct_ permanent-variables record in the
initial-process data space-- see sections 5.2.4.5 and 5.2.4.6.) So,
PUSHVALUE generates an LDO (LoaD glObal word) operator and the
appropriate displacement. "Program" mode variables will be in the
global data space of a sequential program, so they will be addressable
from the global base register, and an LDO instruction is generated.
Variables whose mode is "process entry" require special handling. If
the virtual displacement (the displacement as it comes from pass five)
is negative, the object being referenced is a _variable_ in the entry
routine's local data space, and an LDL instruction will serve to push
it onto the stack during execution. If the displacement is not
negative, the object is a _parameter_ (or function value) of the
routine. It cannot be addressed relative to the local base register
since it is local, not to the entry routine, but to the interface
routine which called it (see section 5.2.2 and figure 29). The object
can be addressed as an "intermediate" (neither local nor global--
lexically somewhere in between) word; that is, relative to _its own_
local base and the distance (in static links) from the current local
base to the required one. To push the variable at run time, PUSHVALUE
generates an LOD (LOaD intermediate word) instruction which has two

operands. One is the variable's displacement from its local markstack, calculated by DISPL. The other is the number of static links which must be traversed to reach the markstack to which the variable is local. The instruction used to invoke the process entry routine (CXL) makes the entry routine's static link field point to the interface routine's markstack, so only one static link needs to be traversed.

The PUSHADDRESS routine selects Microengine code to push onto the stack the address of some data object. The strategy used to make the selection is exactly the same as that used by PUSHVALUE. The instructions which can be generated are LLA (Load Local Address), LAO (Load Address, glObal), and LDA (LoaD intermediate Address). LLA is generated for "procedure", "class entry", and "monitor entry" modes. LAO is generated for "process", "class", "monitor", and "program" modes. For "process entry" mode LLA is generated if the object is a local variable. Otherwise, LDA is generated with a static-link-distance operand of one in order to access the object relative to the markstack of the interface procedure. Variables are never located in the constant pool, so PUSHVALUE never has to deal with constant-mode references. There are cases, however, where the address of a long constant must be pushed onto the stack. For constant-mode references PUSHADDRESS generates an LCA (Load Constant Address) instruction. The operand for LCA is the offset (in words) of the constant, relative to the start of the segment.

PUSHINDIRECT generates code which will push onto the stack the data item which is the target of the top-of-stack pointer. For objects of type WORDTYPE, REALTYPE, and SETTYPE that pointer occupies one stack word. It is two words for BYTETYPE objects, however. The routine generates an SINDO (Short INDex (0) and load word) for pushing

word-length objects. During its execution, the instruction adds the index (0, in this case) to the pointer and pushes the word which the new pointer indicates. For reals, an LDM (LoaD Multiple words) instruction is generated with an operand value equal to the word length of Microengine real values-- two words. An LDM is also generated to push sets indirectly. Its operand is the length of the set to be pushed. The hardware actually supports sets of varying length, but since the virtual machine only supports eight-word sets, MEPASCAL only generates fixed-length sets, and so the operand to the LDM is fixed at eight. Since the hardware expects to find the length of the set as the top-of-stack word (see section 5.2.3), the routine also generates an instruction to push the set's (fixed) length value. For byte objects an LDB (LoaD Byte) instruction is generated. At run time, however, this instruction expects the TOS-1 word to be the address of the first word of a byte array (two meaningful bytes per word), and the TOS word to be the offset (in bytes) from that address to the target byte. During execution, LDB consumes both stack words and leaves in their place a word which contains the target byte in the low-order position, and zero in the high-order position. This mechanism is quite different from that of the Concurrent Pascal virtual machine. The virtual instruction PUSHBYTE assumes that the TOS word by itself points directly to the target byte. Thus, there is an implicit assumption that the underlying hardware is byte addressable. Refer to the next section for information on how the Microengine byte pointer is built.

5.2.4.2 INDEX and BYTE Instructions

The virtual machine instructions PUSHBYTE and COPYBYTE have equivalent Microengine instructions in LDB (LoaD Byte) and STB (STore Byte). PUSHBYTE and LDB push the byte to which the top-of-stack item points. COPYBYTE and STB write the low-order byte of the top-of-stack word into the location where the TOS-1 item points. The actions of the two pairs of instructions are identical; however, the pointers required for the Microengine instructions are not the same as those for the virtual instructions. Pointers on the virtual machine are all alike-- one word long and measured in bytes. The Microengine uses three pointer formats, from one to three words long. Packed field pointers are of no concern here. Word pointers are one word long, apparently use byte measure, and point to the low-address byte (low-order bits, even address) of a word. Word pointers are generated by instructions such as LLA (Load Local Address). Byte pointers consist of two words (see figure 36). The TOS-1 word is a word pointer-- call the byte to which it points the "base". The TOS word contains the offset (byte measure) from the base to the byte which is the ultimate target [REGE78]. If MEPASCAL were to simply generate LDB and STB in place of PUSHBYTE and COPYBYTE, at run time the Microengine instructions will attempt to use the two top-of-stack words as a byte pointer. In fact, the TOS word will be a pointer in virtual machine format and the TOS-1 word will be other irrelevant data which must remain untouched.

The solution to this problem with pointer formats is in the code which generates byte pointers. Byte pointers are generated only for references to Concurrent Pascal strings (ARRAY[N..M] OF CHAR). Notice that in Concurrent Pascal and UCSD Pascal, CHAR variables (and CHAR

FIGURE 36. Microengine byte pointer to the fourth byte (offset = 3) in a string.

constants when they are on the stack) occupy only one word in which the high-order byte is zero. Concurrent Pascal arrays of characters are identical to UCSD Pascal packed arrays of characters (two characters in a word). UCSD arrays of characters have one character per word, an arrangement which is not supported by Concurrent Pascal. In the final analysis, byte pointers are generated only when a byte is referenced as an element of an array. The virtual instruction INDEX is responsible for generating the final address of an array (any array-- CHAR or not) element just before being pushed on the stack or stored, so it is the instruction which was modified to handle the Microengine byte pointer format.

INDEX assumes that the TOS-1 word is a pointer to the base of an array, and that the top-of-stack word is an integer index into the array. During its execution, the instruction compares the index value to the bounds of the array being referenced. If it falls outside, a run-time error occurs, otherwise the index (adjusted to a zero base) is multiplied by the length of an array element to yield the target element's offset (in bytes) from the base of the array. The base and offset are then added, leaving the address of the target element on top of the stack. A direct conversion to a Microengine code sequence would work, except in the case where the array elements are bytes. For byte arrays, the final arithmetic addition of base address and offset must be skipped in order to leave a Microengine byte pointer on the stack. MEPASS5 incorporates a change to the INDEX2 instruction which it emits (section 5.1). In MCPASS5, INDEX2 has three operands. In the MEPASCAL version of pass five, the instruction has a fourth operand which is the virtual data type of the array elements. When INDEX1 is encountered by MEPASS6, the type operand is used to determine if an

instruction should be generated to add the base address and element offset. In retrospect, it seems that the LENGTH operand conveys the same information as the TYPE operand-- the element is either BYTETYPE or not (element lengh is either one or not). If the element length is one byte, then it must need a byte pointer. If the length is used as the discriminant, the type operand and the change to pass five become unnecessary. The Microengine code which is equivalent to INDEX is:

```
push min.      push minimum array bound
push max.      push maximum array bound
CHK            range check
push min.      adjust index to zero base
SBI
{IXA (size)    add base address and offset,
                  yielding element address},
```

where the IXA (IndeX Array) is _not_ generated for a reference to a BYTETYPE element.

5.2.4.3 CASEJUMP

There is a radical difference in the way CASE statements are handled on each target machine. An example of a virtual machine CASE statement is shown in figure 37, and the syntax is given in [HART76]. Notice that the CASEJUMP instruction _follows_ the code for the cases. This makes for a good deal of wild branching during execution, but it also _guarantees_ that pass _six_ of the compiler will be able to resolve the displacements which are operands of the CASEJUMP operator. The displacements are the distances to the code for each case. During execution, the instruction checks the top-of-stack selector value against the legal range defined by the first two operands. If the selector value is outside the range, a run-time error is triggered. If it is in bounds, the selector is used to find the appropriate displacement operand, which is then algebraically added to the virtual

```
0001 VAR
0002 A, B, C, D, E, F, SELCTR: INTEGER;
0003
0004 BEGIN
0005 A := 100;
0006 CASE SELCTR OF
0007    3:     B := 2;
0008    8, 11: C := 3;
0009    9:     D := 4;
0010    12:    BEGIN
0011           E := 5;
0012           F := 5;
0013           END;
0014    6:     F := 6
0015 END;
0016 F := 200;
0017 END.
```

```
7777777777777777777777777777777777777777777777777777777777777777777777777777
7              MCPASCAL INTERMEDIATE CODE PASS 7                             7
7777777777777777777777777777777777777777777777777777777777777777777777777777


JUMP(2)
          BEGNPRCS(4)
LINE   5  GLOBLADD(-2)    PUSHCONS(100)   COPYWORD
LINE   6  PUSHGLBL(-14)   JUMP(138)
LINE   7
LINE   7  GLOBLADD(-4)    PUSHCONS(2)     COPYWORD    JUMP(146)
LINE   8
LINE   8  GLOBLADD(-6)    PUSHCONS(3)     COPYWORD    JUMP(124)
LINE   9
LINE   9  GLOBLADD(-8)    PUSHCONS(4)     COPYWORD    JUMP(102)
LINE  10
LINE  10
LINE  11  GLOBLADD(-10)   PUSHCONS(5)     COPYWORD
LINE  12  GLOBLADD(-12)   PUSHCONS(5)     COPYWORD
LINE  13  JUMP(58)
LINE  14
LINE  14  GLOBLADD(-12)
LINE  15  PUSHCONS(6)     COPYWORD    JUMP(32)
LINE  15
CASEJUMP(3,9,-142,18,16,-38,12,-130,-110,6,-136,-94)LINE   15
LINE  16  GLOBLADD(-12)   PUSHCONS(200)   COPYWORD
LINE  17  ENDPRCS
```

FIGURE 37. MCPASCAL CASE statement.

program counter to realize the case jump. Since there is not a fixed number of cases in a CASE statement, the operator has a variable number of operands.

The Microengine case jump operator (XJP), on the other hand, has only one immediate operand, and it points to the place in the constant pool where all the other necessary information can be found. The constant pool information occupies as many words as there are cases, plus two words. The first word is the minimum legal value for the selector, and the second word is the maximum value. The next words contain the displacements to the code for each case, relative to the byte following the XJP instruction. As mentioned in section 5.2.1, the XJP selector bounds and case offsets cannot be carelessly tossed into the constant pool. Long constants are sprinkled throughout the intermediate code stream until pass six removes and collects them in the heap table CONSTTABLE. Earlier passes have counted the number and relative positions of the long constants, and generated addressing commands based on the assumption all case jump information would be in the code stream. All case information which belongs in the constant pool is removed from the code stream and placed in XJPTABLE when MEPASS6 encounters CASEJUMP1. When MEPASS7 generates the final object code file it will emit one word which points to the end of the segment (see figure 7) and then build the constant pool by dumping the contents of CONSTTABLE and XJPTABLE, in that order. Addressing commands for constants which are not for case jumps will be valid because those constants have not changed position relative to the start of the constant pool. (The routine which calculates displacements for addressing commands-- DISPL in MEPASS6-- compensates for the final-word pointer.) The value of the XJP operand

(segment-relative offset of the selector bounds and case offsets) is the offset of the required information in XJPTABLE (variable XJPOFFSETPTR in MEPASS6) plus the size of CONSTTABLE (known from data in the PASSLINK record) plus one word (the pointer to the segment's final word). When CASEJUMP comes into pass six, the case offsets are represented by case labels. The labels are guaranteed to be resolvable into offsets since the location-counter value of each case label was inserted into JUMPTABLE as it was encountered, and all the cases appear in the intermediate code before the CASEJUMP1 operator. The labels in the CASEJUMP1 instruction are used as indices into JUMPTABLE to retrieve those location-counter values, and the differences between them and the current value of the location counter are the offsets to be inserted in XJPTABLE. The correct placement of the case offsets complicates an otherwise trivial translation of CASEJUMP to the Microengine code sequence:

```
push min.          push minimun CASE index
push max.          push maximum CASE index
CHK                check:   min. <= TOS-2 <= max
XJP (displacement) casejump using constant-pool CASE table.
```

5.2.4.4 Routine Invocation Instructions

The six virtual instructions CALL, INITCLASS, INITMON, INITPROC, CALLPROG, and CALLSYS all invoke routines, but the circumstances which surround the use of each, and the side effects produced by each one vary.

CALL is the simplest of the six. In the final code, its operand is a signed value which is the distance to be jumped in order to reach the code of the called routine. During execution, the address of the instruction after CALL is pushed onto the stack as the return address and the operand is algebraically added to the virtual program counter

so that the next instruction to be executed will be the first one in the called routine. By pushing the return address onto the stack, CALL begins the construction of the markstack (see figure 2) for the called routine. The first instruction of the called routine will finish it (see section 5.2.4.6). By the time CALL executes, other instructions have pushed onto the stack space for the function value and parameters, if required by the called routine. Only non-ENTRY, monitor ENTRY, and class ENTRY routines are invoked by CALL, so in execution sequence, CALL is always followed by either ENTER, ENTERMON, or ENTERCLASS. Along with other actions, these instructions finish building the markstack. In MEPASS6, the Microengine instruction CPL (Call Procedure Local) has been substituted for CALL, although it is not quite equivalent. The operand to CPL is the called routine's number. This happens to be the compiler-generated block label which is an operand of CALL1. During its execution, CPL not only causes a jump to the called routine, but also completely builds the markstack and sets aside local variable space for the routine.

INITCLASS is used to invoke the initial routine (BEGIN..END. block) of a class. Just before INITCLASS executes, other instuctions have pushed onto the stack the address of the "middle" of the record containing the class's permanent variables (see figure 25), and the parameters (access rights) required by the class. When the instruction executes, it pops the parameters from the stack into the component variable record, puts the return address on the stack, and jumps to the class initial code.

In order to move this instruction to the Microengine, pass five and the input syntax for pass six had to be changed. Even though INITCLASS uses the address of the class variable in copying the

parameters, it _leaves_ it on top of the stack for later use in setting the global base. For the Microengine, an STM instruction is used to pop the parameters and, as a side effect, it consumes the class variable address. The syntax change involved adding an instruction to duplicate the class address in order to leave a copy of it on top of the stack for later use. The original pass six input syntax for an _init stat_ [HART76] was

```
---> varaddr ---> FIELD(disp) ---> arg list ---|
        -------------------------------------
        |
        ---> INIT(mode, label, parm length, var length) --->  .
```

In order to preserve the class address, DUPTOS was added, giving the syntax

```
---> varaddr ---> FIELD(disp) ---> DUPTOS ---> arg list ---|
           ----------------------------------------------
           |
           ---> INIT(mode, label, parm length, var length) --->   ,
```

where DUPTOS translates to the Microengine hardware instruction DUP1. Once the access rights have been put in the class variable, the initial code for the class is invoked by a CPL instruction, so the Microengine code for INITCLASS is

```
        STM (paramsize)     pop "paramsize" words into class variable
        CPL (routine no.)   invoke initial routine of class.
```

The operation of, and equivalent code for, INITMON is the same as for INITCLASS. BEGINCLASS and BEGINMON instructions follow INITCLASS and INITMON _in execution_. They complete construction of the markstack and set the global base for the class or monitor.

INITPROCESS calls on the kernel to create a new process and get it started in its own initial routine. The kernel pops the parameters into the process variable and leaves its address on top of the stack. After the kernel call, the original process merely discards the

address. The equivalent Microengine code for the kernel call has not yet been determined, and the instruction which pops the top-of-stack word is FJP(0). The execution sequence is somewhat unclear since two processes are executing concurrently after the kernel creates the new process. During the execution of INITPROCESS, the original process enters and returns from the kernel. The new process starts out in kernel code and the first virtual instruction it executes is BEGINPROC.

CALLPROG pushes the return address onto the stack and starts the execution of a sequential program stored in a variable whose address is on top of the stack. Because of the Microengine architecture, the mechanism for starting a sequential program is entirely different, as described in section 3.4. The code for invocation of a program is given below, and assumes that the code variable address is on top of the stack. In short, the code puts the address in the sequential program's segment information block, and then calls the initial routine in that segment. Figure 14 will be helpful for following all the pointers which must be chased. The Microengine code is

```
SLDC01       push CTP register number (-1)
NGI
LPR          push pointer to TIB (CTP register content)
IND (11)     push pointer to SIB vector (11th word in TIB)
INC (5)      increment by 5 words to yield pointer to SEGBASE field
                in SIB no. 129. (SIB 129 is second record in
                SIB vector, each SIB is 5 words long, and SEGBASE
                is first field in each SIB.)
SWAP
STO          pop pointer to program variable into
                SEGBASE field of SIB 129
CXL (129, 1) invoke initial routine of sequential program.
```

The ENTRY routines of a process are the routines which provide operating system services to the sequential program running as a part of that process. In the virtual machine, the sequential program requests one of those services by executing the instruction

CALLSYS(index), where the operand is the index of the service in the sequential program prefix. CALLSYS is comparable to a conventional SVC instruction. Several events must have taken place at the time of sequential program invocation in support of CALLSYS, as described in section 3.4.1. Just before the program starts, code in the concurrent program puts on the stack a table of addresses of the process ENTRY routines. This jump table (figure 16) remains on the stack while the sequential program executes. CALLSYS uses this index operand to select an ENTRY routine address from the jump table, and then causes a jump to that address. The sequential program is, essentially, requesting system services by number. Notice that the operating system and sequential program must agree on the numbers assigned to the services. The assignments are based on the order in which the service identifiers appear in the interface definition (in the host process) and in the prefix (in the sequential program). Suppose, for example, that a process has eight ENTRY routines named OPEN, CLOSE, GET, PUT, ACCEPT. DISPLAY, MARK, and RELEASE, and that the sequential program SEQPROG will have access to the first six. That is, SEQPROG will only be able to use the first six; in fact, it will not even know of the existence of the other two. If the program and interface are defined as

```
SEQPROG (A, B: INTEGER; C: SEQL_CODE_TYPE);
ENTRY OPEN, CLOSE, GET, PUT, ACCEPT, DISPLAY;
```

then the sequential prefix must list those routines in the same order (although the identifiers can be different):

```
PROCEDURE OPENFILE (parameter list);
PROCEDURE CLOSEFILE (parameter list);
PROCEDURE GETPAGE (parameter list);
PROCEDURE PUTPAGE (parameter list);
PROCEDURE READCONSOLE (parameter list);
PROCEDURE WRITECONSOLE (parameter list);  .
```

If the two sequences fail to match (say the order of OPENFILE and CLOSEFILE is reversed), the results of a CALLSYS instruction will not be as expected (a call to OPENFILE will cause the process ENTRY routine CLOSE to be executed). As long as the interface and prefix definitions maintain the proper relationships, each program (concurrent and sequential) can be altered and recompiled independently of the other. This independence is possible because routines are known and invoked by their address, and CALLPROG indirectly takes the ENTRY routine address from the stack.

In the virtual machine the run-time jump table performs a mapping function from the prefix index number known to the sequential program to the corresponding ENTRY routine address-identifier. Since Microengine routines are known and invoked by their routine number, and the invocation instructions use only immediate operands, that mapping function must be performed by the interface segment described in sections 3.4.1 and 5.2.4.8. By the time the Microengine code equivalent to CALLSYS executes, the code segments (both shared and private to the host process) are configured as shown in figure 18. The interface segment consists of only one routine which operates as described in sections 3.4.1 and 5.2.4.8. On the Microengine, process ENTRY routines will be called by the interface routine, which is itself called by the sequential program. The interface routine requires a parameter which is the equivalent of the CALLSYS operand. The code to call the interface routine is:

```
push index      push parameter for interface routine
CXL (128, 1)    invoke interface routine in interface segment.
```

5.2.4.5 BEGIN Instructions

The virtual instructions BEGINCLASS, BEGINMON, and BEGINPROC are always the first instruction in the initial routine of classes, monitors, and processes, respectively. In _execution_ they are always preceded by the corresponding INIT instruction. The BEGIN instructions are generated by pass six when ENTER1 is encountered, based on its mode operand. (ENTER1 also translates to the ENTER instructions described in the next section.) In general terms, BEGINCLASS checks for the possibility of a heap-stack collision during the routine, finishes construction of the markstack (begun by INIT), and sets the global base register so that it points to the "middle" of the class variable (see figure 25).

Before describing the equivalent Microengine code itself, some comments must be made concerning the actions of the compiler while generating that code. When MEPASS6 encounters ENTER1, the instruction operands are copied so that their values are preserved during the compilation of the entire routine (the variables ARGx take on new values almost every time a virtual instruction is read from the input file). The "save" variables have the following meanings:

BLOCK-- the numeric label by which the routine is known and which will be used by Microengine instructions (CPG and CPL, for example) to invoke the routine;

PARAMLENGTH-- the number of bytes of parameters for the routine;

VARLENGTH-- the number of bytes of local variables required for this routine;

STACKLENGTH-- the number of bytes of extra stack space to be reserved for the routine.

In the original compiler. PARAMLENGTH and VARLENGTH for initial

routines of classes, monitors, and processes are always zero, reflecting the fact that the permanent variables are not located on the stack (rather, in the component variable record) and that the parameters (access rights) are popped off the stack into the component variable record by INIT. In MEPASCAL PARAMLENGTH and VARLENGTH always contain the corresponding length values (modification made to pass four-- see section 5.1). The length of the variables area is needed to determine the location of parameters, since variables intervene between the base (markstack) location and the parameters (see figure 19). After saving the instruction operands, AFTERBEGIN is set TRUE to indicate that the compiler is now in a routine body. In prior passes, NEWLINE instructions were generated for each source line, including lines in the declaration parts of routines. Since the intermediate code for declarations has been removed, the declaration parts consist of only NEWLINE operators. NEWLINE operators encountered outside of routine bodies (AFTERBEGIN = FALSE) are deleted from the code sent to pass seven. TEMP and MAXTEMP are used to calculate the run-time stack requirements of the routine (section 4.2). They are initialized at zero since the compiler has not yet encountered any code which will, at run time, push anything on the stack. The location counter (which is relative to the beginning of the segment) is incremented by two words to account for the space which the routine's EXIT-IC and DATASIZE fields will occupy in the final code (see figure 7).

After the compiler has completed the actions which are common to all routine modes, code is generated and other actions are taken, based on the value of the mode operand. For "class" mode the routine's DATASIZE value is inserted into its corresponding position in DATASIZETABLE. DATASIZE is the number of words this routine requires

for local variables. This is not the same as VARLENGTH. In "class" mode, VARRLENGTH measures variable space in the class record, whereas DATASIZE is the number of words which must be pushed on the stack for variables. Since the parameters (access rights) and permanent variables are in the class record, only two words of local variable space are needed on the stack. These are the words for storing the source line number and old global base value. POPLENGTH is the number of bytes of stack space which must be popped at the end of the routine in order to return the stack to its configuration before the routine invocation. It will become the operand to the RPU instruction which terminates the routine, and in general, it is the number of bytes of stack space occupied by the routine's variables and parameters. For a class routine there are two words of variables and one word of parameters (the class record address left by code equivalent to INITCLASS). Code is generated to save the source line number and the old global base address, and store the new global base address into the BP register. The address of the class record cannot be used directly as the new global base since the Microengine hardware takes into account the size of the markstack when variable and parameter references are made (see section 3.4.2, figure 9, and figure 19).

Consequently, the Microengine code sequence which is similar to BEGINCLASS includes instructions to reduce the global base address by the size of a markstack, as shown here:

```
push line no.        save souce line number in local word 1
STL (1)
SLDC06               push global base register (BP) number
LPR                  push global base value
STL (2)              save global base in local word 2
SLDC06               prepare to put new global base in BP
LDL (3)              push class record address
SLDC08               push size (bytes) of a markstack
SBI                  BP must point to bottom of imaginary markstack
SPR                  put new global base in BP register.
```

This code does _not_ include a check for heap-stack collision. Presumably, that code can be added when MEPASS6 is modified to pass STACKTABLE on to MEPASS7 as discussed in section 5.2.3. Also, there is no code here to build any part of a markstack since it is built _completely_ by the Microengine code equivalent to INITCLASS.

The code for BEGINMON is the same as BEGINCLASS, except that after the new global base has been established, a call must be made to the kernel in order to initialize the monitor's gate.

Although the compiler actions are the same, the equivalent code for BEGINPROC is much simpler than that for the corresponding class instruction. The DATASIZE value (again, two words) is entered into the DATASIZETABLE and the POPLENGTH is calculated as two words (line number and old global base). POPLENGTH does not include the "parameter" location containing the process record address because that location is in the data space of the initializing process. The equivalent code for INITPROC includes a FJP(0) instruction to pop the component address (see section 5.2.4.4). The Microengine code for BEGINPROC is

```
push line no.        save source line number in local word 1
STL(1).
```

The kernel will handle establishment of the new global base address and construction of the markstack. Notice that there is no old global base since the process did not have any prior existence.

5.2.4.6 ENTER Instructions

The virtual instructions ENTER, ENTERCLASS, ENTERMON, ENTERPROC, and ENTERPROG are always the first instructions of non-initial routines. Non-ENTRY routines on the virtual machine always begin with an ENTER instruction, and in _execution_ sequence it is always preceded by CALL. Generally, it checks for the possibility of a stack-heap collision, finishes construction of the markstack begun by CALL, and sets aside stack space for local variables. It does _not_ affect the global base. When it encounters an ENTER1 instruction with a "procedure" mode operand MEPASS6 calculates the routine's DATASIZE and POPLENGTH values and generates a simple code sequence. Unlike initial routines, local variables actually reside _on the stack_, so DATASIZE is the number of words the routine needs for explicitly declared local variables (VARLENGTH) plus two words reserved for the source line number and old global base. Although the latter is not _necessary_, it is included for the sake of consistencey-- all routines have local words one and two reserved. The global base word could be removed as long as the routines and calculations affected by that removal are also modified. The calculated values for POPLENGTH and DATASIZE would have to be reduced by one word, and the procedure DISPL would need to use a separate calculation for determining variable displacements. The number of bytes to be popped off the stack at the end of the routine is determined by the length of the routine's parameters (PARAMLENGTH), the size of the declared variables (VARLENGTH), and the two local

reserved words. The sequence of generated Microengine code is

> push line no.   put source line number in local word 1
> STL (1).

Just as for BEGINCLASS (section 5.2.4.5), there is no code here to check for a stack collision, or to explicitly build any part of the markstack.

ENTERCLASS is the first instruction of class ENTRY routines. Its function on the virtual machine is the same as BEGINCLASS (section 5.2.4.5). Consequently, the machine code which is generated in place of it is also the same. The compiler actions are different, however. The DATASIZE value for the routine is the size of the explicitly declared local variables (VARLENGTH) plus two words for the line number and old global base. Unlike the code for ENTER, the code for ENTERCLASS does use the local word reserved for the old global base. At the conclusion of the routine, its parameters, the address of the class record, its local declared, and its local reserved variables must be removed from the stack, so POPLENGTH is the sum of PARAMLENGTH, VARLENGTH, and three words.

ENTERMON is the first instruction of monitor ENTRY routines and performs the same actions as ENTERCLASS. It also makes a call on the kernel to request passage through the monitor's gate. MEPASS6 calculates DATASIZE and POPLENGTH the same way as for ENTERCLASS, then generates identical code. After the kernel call mechanism is known, code will be generated to request gate entry from the kernel.

The virtual instructions ENTERPROG and ENTERPROC were difficult to move to the Microengine because they are affected by the design and operation of the concurrent-sequential interface. ENTERPROG is the first instruction of the initial routine of a sequential program, so in execution sequence it is always preceded by CALLPROG. It completes

construction of the markstack begun by CALLPROG, checks for a possible stack collision, and allocates stack space for the initial routine's local variables. The local variables of the initial routines are the global variables of the entire sequential program, so the instruction also resets the global base register so that it points to the same place as the local base register. In addition, the instruction sets to the value 1, a word (JOB) in the kernel associated with the invoking process. This is used to indicate that the process is in sequential (user) code, not concurent (operating system) code. The Microengine code generated in place of ENTERPROG is

```
push line no.
·STL (1)        save source line number
SLDC06         push global base (BP) register number
LPR            push global base pointer
STL (2)        save old global base
SLDC06         BP register number for store
SLDC05         local base (MP) register number
LPR            push MP
SPR            store into BP.
```

The markstack was completely built by the Microengine instruction CXL which invoked the sequential program, so the code here does not do any markstack construction. The source line number (in the sequential program) and global base (pointer to the permanent variables record of the host process) are saved, as for most other routine entries. Finally, the sequential program's global base is established. No action comparable to setting the concurrent/sequential switch in the kernel is taken since the kernel has not yet been designed.

ENTERPROC is the first instruction of a process ENTRY routine, and it follows CALLSYS in execution. In the virtual machine, it checks for possible stack collision, finishes construction of the markstack, establishes the process global base, and zeroes the JOB word for the process in the kernel to indicate the execution of concurrent

(supervisor) code, not sequential (user) code. Since the concurrent/sequential interface mechanism (section 3.4) is so different, the code which emulates ENTERPROC is not quite what one would expect. On the Microengine, process ENTRY routines are not called directly from sequential programs, as they are on the virtual machine. Instead, they are invoked from the interface routine (see figure 38).

The code which MEPASS6 generates is affected by the configuration of certain pointers immediately after a process ENTRY routine is invoked (see figure 39). Before the host process invokes the sequential program the global base register points to the record containing the permanent variables and access rights of the process. Microengine code equivalent to CALLPROG starts the program, and code equivalent to ENTERPROG stores the global base in the second local variable as described above. The global base register is then adjusted to point to the program's global storage area-- the same as the local area at the time the program is started. During execution the local base will change as routines are called and return, but the global base remains fixed. The companion version of Sequential Pascal apparently does not allow nested routine definitions. The global base register is not altered when the interface routine is invoked by code emulating CALLSYS, nor is it altered by the interface routine itself. Even by the time code equivalent to ENTERPROC is about to execute, the global base still points to the global variable area in the sequential program. That code saves the source line number and global base pointer in its own local space, and re-establishes the process's global base pointer by fetching the second word in the (sequential program) global space and storing it into the global base register.

CONCURRENT
PROCESS

SEQUENTIAL
PROGRAM

ENTERPROG

CALLSYS

EXITPROG

ENTERPROCESS

ENTRY
ROUTINE

EXITPROCESS

BEGINPROCESS

INITIAL
BLOCK

CALLPROG

INTERFACE
ROUTINE

ENDPROCESS

CXL (PROCESS ENTRY)

RPU

FIGURE 38. Flow control during the life of a process which uses an ENTRY routine, on the Microengine. Wavy lines represent sequential execution of instructions; narrow lines represent sequential program invocation and return; wide lines are the path taken to invoke the ENTRY routine; and dashed lines represent the return from the ENTRY routine. All of the code is Microengine machine code. Virtual instruction names have been used here only to indicate more clearly the function of the relevant code.

SEQUENTIAL
PROGRAM

```
┌─────────────────────────┐
│        GLOBAL           │
│       VARIABLES         │
├ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ┤
│     OLD GLBL. BASE      │○───────┐
├ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ┤        │
│        LINE NO.         │        │
├─────────────────────────┤        │
│     SEQ'L  PROGRAM      │        │
│    INITIAL  ROUTINE     │        │
│       MARKSTACK         │◄───────┤
└─────────────────────────┘        │
```

GLOBAL BASE

```
┌──────────────┐
│      ○       │
└──────────────┘
```

PROCESS
RECORD

```
┌─────────────────────────┐
│                         │
│                         │
│       PARAMETERS        │
│                         │
│    PARAMETERS           │
│     FOR                 │
│     ENTRY               │
│    ROUTINE              │
├ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ┤
│        INDEX            │
├─────────────────────────┤
│      INTERFACE          │
│      MARKSTACK          │
├─────────────────────────┤
│        ENTRY            │
│       ROUTINE           │
│        LOCAL            │
│      VARIABLES          │
├ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ─ ┤
│    OLD GLBL. BASE    ○─ ┤
├─────────────────────────┤
│       LINE NO.          │
├─────────────────────────┤
│        ENTRY            │
│       ROUTINE           │
│      MARKSTACK          │
│     (STATIC  LINK)   ○  │
└─────────────────────────┘
```

PARAMETERS
(ACCESS RIGHTS)

PERMANENT
VARIABLES

IMAGINARY
MARKSTACK

FIGURE 39. Configuration of global base pointers and process ENTRY routine's static link immediately after the ENTRY routine has been invoked. Solid pointers exist before the Microengine code comparable to ENTERPROC has executed. Dashed pointers exist after that code executes.

The Microengine code to do this is:

```
push line no.
STL (1)        save source line number
SLDC06         BP register number
LPR            push global base pointer
STL (2)        save copy of global base
SLDC06         BP register number
SLD002         push process global base saved by ENTERPROG code
SPR            store into global base register.
```

The mechanism for indicating the return to supervisor code has not yet been determined. The compiler calculates the ENTRY routine's DATASIZE and POPLENGTH as the length of its declared local variables plus two words to store the global base (of the sequential program) and source line number. POPLENGTH does not include the size of the parameters since they are local to the interface markstack, not the ENTRY routine markstack (see figure 39). Also, POPLENGTH does not include an extra word to account for the component variable address "parameter" being on the stack, unlike ENTERCLASS and ENTERMON code. The component variable address never appears on the stack here since it is fetched out of the global variable space of the sequential program.

5.2.4.7 END and EXIT Instructions

Eight virtual instructions terminate routines by restoring the calling routine's environment. All of these use an RPU instruction on the Microengine. RPU (Return from Procedure--User) cleans up the stack in order to restore the stack to the condition it was in before the routine was called. It has one operand which is the number of words (not including the markstack) which must be popped in order to restore the calling routine's stacktop. The instruction automatically restores the caller's stack pointer (based on the operand) and local base pointer, and causes a jump to the return address in the caller's code. The instruction does not affect the global base pointer.

The virtual instruction EXIT terminates routines which are neither ENTRY nor initial routines. It performs the same function as RPU on the Microengine, except that it also restores the caller's global base pointer. Restoration of the global base is superfluous for non-ENTRY and non-initial routines since the global base is not changed by ENTER (section 5.2.4.6) and any system component routines called by the routine restore the caller's BP register when they return. The equivalent Microengine code is simply

RPU (POPLENGTH DIV WORDLENGTH),

where POPLENGTH is the value calculated when code for ENTER was generated at the start of the routine.

ENDCLASS and EXITCLASS terminate the initial and ENTRY routines of a class. Since the corresponding prologue instructions, BEGINCLASS and ENDCLASS, change the global base register, the caller's global base pointer must be restored before returning. The equivalent Microengine code is:

```
SLDC06              push global base register number
SLDL02              push caller's global base pointer
SPR                 restore caller's global base
RPU (poplength div wordlength)   return to calling routine.
```

ENDMON and EXITMON terminate monitor routines and are analogous to ENDCLASS and ENTERCLASS. They perform the same function, except that the monitor instructions call on the kernel to perform a gate exit before restoring the caller's global base and returning. So, the equivalent Microengine code consists of a kernel call followed by the same code which was substituted for ENDCLASS and EXITCLASS.

EXITPROC terminates process ENTRY routines. It restores the caller's (sequential program's) global base and sets the JOB switch in the kernel to indicate the execution of sequential code. The Microengine code for this is the same as for ENDCLASS plus some, as

yet undetermined, code to indicate the return to user code.

ENDPROC terminates the initial routine of processes, and is merely a return to the kernel to make the process "disappear". The Microengine code is just an RPU instruction.

EXITPROG is the last virtual instruction of a sequential program's initial statement, and in general terms, it restores the environment of the host process before the sequential program was invoked. Since it manipulates several locations in the kernel, the equivalent Microengine code has not been finalized. The code will, however, end with an RPU instruction.

Every time one of the instruction sequences discussed in this section is generated, the compiler enters the routine's STACKLENGTH into STACKTABLE. This provides a record of the routine's stack requirements for calling routines to use in calculating their own stack requirements (see section 4.2). AFTERBEGIN is set FALSE so that NEWLINE instructions for declaration source lines will not appear in the final code. If the location counter is not on a word boundary a NOP (No OPeration) is generated to force it to the required value. Alignment is required for the next routine's EXIT-IC field (see figure 7).

5.2.4.8 PUSHLABEL (Interface Generation)

The action of PUSHLABEL in the virtual machine is quite simple-- it merely pushes the address of a process ENTRY routine onto the stack. PUSHLABEL instuctions are used as part of the sequential program invocation sequence to build the jump table shown in figure 16. The jump table maps prefix routine indices to process ENTRY routine addresses during the execution of CALLSYS instructions

(sections 3.4.1 and 5.2.4.4). The concurent/sequential interface mechanism for the Microengine is entirely different.

A skeletal concurrent process and interface definition are shown in figure 40. The process can provide as many as six operating system services (P1 through P6) to sequential programs hosted by it. However, the code loaded into the variable C, and invoked as SEQLPGM, will only be offered three of these (P4, P6, and P3) for its use. The sequential program is invoked in line 16, and it is here that the run-time jump table is built to enforce access restrictions. When MEPASS6 encounters a PUSHLABEL instruction, it removes from the intermediate code the instructions which build the jump table, and generates a new code segment called the interface segment. The interface segment consists of only one routine (figure 17) and is placed in file number four (figure 26). The structure of the segment is known by MEPASS6, so the number of process ENTRY routines in the interface completely defines it. This permits the pass to generate a Microengine code segment which is extremely close to its final form. Pass seven will pack the code into machine words, pad it out to an integral multiple of the size of a disk block (512 bytes), and append it to the concurrent program's code segment.

The occurrence of a PUSHLABEL operator informs MEPASS6 that it has just found the start of an interface. The routine GEN INTERFACE reads through the incoming interface and generates the interface segment. The general machine-code form of the interface segment is shown in figure 41. Before emitting any code, the routine turns on the GENNINGINTFAC switch so that the routines which usually send generated code to the concurrent code file will start putting code into the interface segment file (see section 5.2.2). It also fetches

```
0001 TYPE PRC = PROCESS;
0002     TYPE CODE = ARRAY[1..1000] OF INTEGER;
0003     VAR PERM1, PERM2, PERM3: INTEGER;      CODEVAR: CODE;
0004
0005     PROCEDURE ENTRY P1;    BEGIN    END;
0006     PROCEDURE ENTRY P2;    BEGIN    END;
0007     PROCEDURE ENTRY P3;    BEGIN    END;
0008     PROCEDURE ENTRY P4;    BEGIN    END;
0009     PROCEDURE ENTRY P5;    BEGIN    END;
0010     PROCEDURE ENTRY P6;    BEGIN    END;
0011
0012     PROGRAM SEQLPGM(X, Y, Z: INTEGER; C: CODE);
0013     ENTRY P4, P6, P3;
0014
0015     BEGIN
0016     SEQLPGM(100, 200, 300, CODEVAR);
0017     END;
0018
0019 VAR   PRCV: PRC;
0020
0021 BEGIN
0022 INIT PRCV;
0023 END.
```

```
55555555555555555555555555555555555555555555555555555555555555555555
5              MEPASCAL INTERMEDIATE CODE PASS 5               5
55555555555555555555555555555555555555555555555555555555555555555555
   JUMP(1)   LNGCONST(16,HEX-VAL:0000,0000,0000,0000,0000,0000,
   0000,0000)

LINE    5    ENTER(3,3,0,0,0)    RETURN(3)
LINE    6    ENTER(3,4,0,0,0)    RETURN(3)
LINE    7    ENTER(3,5,0,0,0)    RETURN(3)
LINE    8    ENTER(3,6,0,0,0)    RETURN(3)
LINE    9    ENTER(3,7,0,0,0)    RETURN(3)
LINE   10    ENTER(3,8,0,0,0)    RETURN(3)

LINE   15    ENTER(6,2,0,2006,0)
LINE   16    PUSHLABL(5)    PUSHLABL(8)    PUSHLABL(6)
             PUSHCNST(100)    PUSHCNST(200)    PUSHCNST(300)
             PUSHADDR(6,-2006)    CALLPROG    POP(6)
LINE   17    RETURN(6)

LINE   21    ENTER(6,1,0,2,0)
LINE   22    PUSHADDR(6,-2)    DUP1TOS    INIT(6,2,0,2006)
LINE   23    RETURN(6)    EOM(2)
```

FIGURE 40. A concurrent process which invokes a sequential program, and the intermediate code produced by MEPASS5. The second operand of ENTER is the routine number.

```
    @ last word
mir. case index
max. case index
        |
        |
        |
RTNS words of
CASE statement
    offsets
        |
        |
        |
    exit-ic
  datasize (0)
      SLDL01                      push the prefix index parameter
      XJP (1)                     CASEJUMP--case offsets start in word 1
      NOP
      UJPL (out)                  jump if parameter out of range
      CXL (129, label 1)          call process entry routine
      UJPL (out)                  jump out of CASE statement
      CXL (129, label 2)
      UJPL (out)
            .
            .
            .
            .

            .
      CXL (129, label RTNS)
      UJPL (out)
out: NOP                          no operation, for word alignment
      RPJ (1)                     return to caller (sequential program)
    @ datasize
    1, segment id
```

FIGURE 41. General layout of an interface segment.

the size of (number of routine labels in) this interface from the array IFSEGSIZE which was loaded in procedure BEGINPASS from the PASSLINK record in the heap. The previous pass counts the number of interfaces in the concurrent program and the size of each one (see section 5.1). For the rest of this discussion the number of routines in the interface will be denoted "RTNS". The sizes (in bytes) of the various items in figure 41 are:

```
1) @ last word-- 2
2) min. case index-- 2
3) max. case index-- 2
4) offsets-- 2*RTNS
5) EXIT-IC-- 2
6) DATASIZE-- 2
7) pre-case code:
       SLDL01, XJP, NOP, and UJPL-- 7
8) code for each case:
       CXL and UJPL-- 6
9) post-case code:
       NOP and RPU-- 3
10) @ DATASIZE-- 2
11) value 1-- 1
12) segment id.-- 1.
```

The first word of a Microengine code segment points to the last word in the segment (see section 3.2.2). For an interface segment, that pointer value can be calculated as the number of words in the segment less one word. Using the list above, that is

$$(\quad (2 + 2 + 2 + 2*RTNS + 2 + 2 + 7 + 6*RTNS +$$
$$3 + 2 + 1 + 1) \text{ DIV } 2 \text{ "bytes per word") } - 1,$$

or more concisely,

$$(8*RTNS + 22) \text{ DIV } 2.$$

GEN INTERFACE puts that value in the interface file. It then starts building the constant pool which will only contain information required for the XJP instruction: the minimum and maximum case indices and the offset values for each case. The sequential program numbers its prefix routines consecutively from 1 to RTNS, so 1 and RTNS are emitted as the index limits. For each case, the distance from the

instruction after the XJP (NOP) to its code is its offset. The case offsets are generated next. The offset for any case is four bytes (NOP and UJPL after XJP) plus the size of the cases between the UJPL and itself. Thus, the first four cases have offsets 4, 10, 16, 22. If the cases are numbered from 0 to RTNS-1, a case's number (let it be KASE) can be used to calculate its offset:

$$KASE * 6 \text{ "bytes per case"} + 4.$$

The next field in the output file is the interface routine's EXIT-IC value, a segment-relative byte pointer to the epilogue code--the RPU instruction. By a calculation similar to the one above, the pointer value is

$$RTNS * 8 + 18.$$

The size of the stack space required by the routine for local variables, the routine's DATASIZE value, must be generated next. Interface routines have no local variables, so the value is always zero. The next section of the segment is the code itself. The first instruction (SLDL01--Short LoaD Local word 1) will, at run time, push onto the stack the interface (prefix) index number of the process ENTRY routine being accessed by the sequential program. The program pushed the index while executing Microengine code equivalent to CALLSYS (section 5.2.4.4). XJP (case jump) is the next instruction. It uses the top-of-stack index to select an offset from the case table which always starts in word location 1 of the segment (location of the minimum index word). That offset is added to the IPC to jump to the corresponding case code. A no-operation instruction comes next, although it might prove to be superfluous. It is generated only in imitation of the way the UCSD Pascal compiler handles case jumps. If the case index is not between the minimum and maximum values the next

instruction (UJPL) jumps around the code for all cases. Since the
cases all have six bytes of code, the jump distance is 6*RTNS. The
code for all the cases follows, but first the compiler must change the
order of the process ENTRY routine labels. In line 13 of the source
code in figure 40 the accessible ENTRY routines are given in the order
P4, P6, P3. Their corresponding routine labels are 6, 8, and 5, so the
prefix routine index numbers (which the sequential program uses to
call the routines) 1, 2, and 3 must map to concurrent routines 6, 8,
and 5, respectively. By the time the labels reach pass six, however,
their order has been reversed (figure 40, line 16 of the intermediate
code). The reverse ordering works for building the jump table for the
virtual machine, but it is the reverse of what is needed for the
Microengine. GEN_INTERFACE removes the PUSHLABEL instructions fom the
code stream and pushes the label numbers onto its own heap stack. The
labels are popped off as code is generated, so that, using the example
in figure 40, when the sequential program provides an index parameter
of 1, concurrent routine 6 (P4) is invoked. The code for each case
consists of two instructions–CXL (Call eXternal, Local routine) and
UJPL. The CXL operands are the segment number of the concurrent
segment (always 1), and the routine label of the process ENTRY routine
to be invoked when the case is executed. The jump instruction skips
around the other cases when the ENTRY routine returns. Its operand is
the number of bytes of code for the cases which follow. At compile
time, this is the size of the cases for which code has yet to be
generated. For the interface in figure 40 the case code would be

```
CXL (129, 6)
UJPL (12)
CXL (129, 8)
UJPL (6)
CXL (129, 5)
UJPL (0).
```

A NOP is generated next for word alignment. Note that this is necessary only because of the NOP which follows the case jump operator. The procedure dictionary follows, and it cnsists of just a single entry since there is only one routine in the segment. The value to be entered is the word address, relative to the start of the segment, of the routine's EXIT-IC field. From figure 41 it can be seen that the value is RTNS+3. The last word to be generated contains two one-byte fields. The high-order byte is the number of routines in the segment, and the low-order byte is the segment number for this code segment. Once they have been loaded into main memory, _all_ interface segments will have a segment number of 129 (see section 3.4.1), but the compiler numbers them consecutively after the concurrent segment number to help identify the segments within the code file. Finally, the switch GENNINGINTFAC is turned off so that output goes into the concurrent intermediate code file once again.

CONCLUSION

6.1 RESULTS AND OBSERVATIONS

Although a good deal of work remains to be done in order to produce a finished compiler, the work which has been completed so far is a major contribution toward that goal. There have been two major results of this project. First, modification of pass six is nearly complete. The structure and intended operation of the pass have been documented in this report, so completing the changes should not be difficult. Second, and more importantly, this report provides a substantial base of knowledge regarding the Concurrent Pascal virtual machine and Western Digital Pascal Microengine, and how the architecture of the former can be mapped to the latter. The software tools described in Appendix A are less significant results, but should be of great help to future workers, as they were to the author. The file transfer program will be useful-- necessary, in fact-- even after the development has been completed and the MEPASCAL compiler is used for production work.

The major observation during this project has been that small architectural differences in target stack machines can have profound consequences. The two target machines _seem_ to be similar. In some respects they are, but in others they are not; and those differences turned out to be quite significant. These ramifications impact almost every pass of a multipass compiler. Thus, porting a Concurrent Pascal compiler to a new stack machine involves a substantial amount of effort. The inability of the Microengine to call a routine whose number is on the stack forced the design of a completely new mechanism

for the concurrent/sequential program interface, and a host of changes to the compiler. The position of the markstack relative to parameters and local variables is different in each machine and that forced a change in the way pass six calculates displacements. The differences in the location of tables for CASE statements was also a source of major changes to pass six. The ENTER and BEGIN virtual instructions have no analogs on the Microengine, and finding equivalent code sequences for them was one of the hardest parts of the project, second only to designing an interface mechanism. In short, the virtual machine instruction set has been specialized for Concurrent Pascal programs, but the Microengine has a more general-purpose instruction set, and it is _not_ easy to simulate one with the other.

By month, the time spent on the project was:

October-December, 1979; 60 hours spent studying preliminary documents for the Microengine system;

February, 1980; 60 hours spent studying PDP-11 assembly code for interpreter, kernel, and low-level interaction with hardware;

March, 1980; 75 hours spent installing the Microengine, studying accompanying documents, gaining familiarity with Microengine system, and bringing up MCPASCAL on the 8/32;

April, 1980; 75 hours spent investigating Microengine architecture and working with the system;

May, 1980; 75 hours spent investigating virtual machine architecture and writing MNEM;

June, 1980; 75 hours spent studying the virtual machine architecture. integrating MNEM into MCPASCAL. Started making instruction modifications in the pass six code;

July, 1980; 75 hours spent studying the Microengine code file format, the architectures of the two target machines, and asynchronous I/O hardware on the Microengine;

August, 1980; 100 hours spent making more instruction modifications, bringing up MEPASCAL on the Interdata 8/32, writing the file transfer program (Appendix A), and designing the concurrent/sequential interface;

September, 1980; 75 hours spent making instruction modifications and writing the interface segment generator;

October, 1980; 75 hours writing and integrating MNEM into MEPASCAL, and testing the file transfer program;

November, 1980; 40 hours spent writing this report and reevaluating the code modifications;

December, 1980; 200 hours spent writing this report and reevaluating the code modifications;

January, 1981; 100 hours spent writing this report and reevaluating the code modifications.

The times shown for the months March through October are probably very conservative.

Of the 1085 hours of effort, 420 hours (38.7%) were spent in learning the stack machine semantics of both machines. If the future work (described below) to complete the porting takes two months of effort, then the learning effort is only 29.6% of the total. Discounting the learning effort, six months of effort would be required to port an operating system. We feel this is a minimal porting effort for an operating system and all of its associated Pascal utilities and applications programs.

6.2 FUTURE WORK

Although the major modifications to the compiler have already been made, work remains to produce a working compiler.

STACKTABLE should be left in the heap after pass six and used to re-implement the heap-stack collision detection mechanism built into the NEW, NEWINIT, BEGIN, and ENTER virtual instructions.

Integer and real formats will have to be changed. Since it runs on an Interdata 8/32, MEPASCAL generates integer and real constants in 8/32 format. Microengine integers have their low-order bits in the low address byte. The difference will probably only affect instructions with word operands-- LDCI (LoaD Constant Integer), UJPL (Unconditional JumP Long), and FJPL (False JumP Long)-- such that the operand bytes will need to be swapped. This could be done in pass seven. Real constants will have to be changed to the four-byte PDP-11 format. This would be best done at the point in the compiler (prior to pass six) where they are generated. Real constants are stored in the constant pool, and it would be impossible at any other point to distinguish a real constant from a string of the same length. Since Microengine reals are only half as long as those generated by the original compiler, the amount of space allocated for real variables must be adjusted in pass four.

Displacement calculations for variables will also have to be changed in pass four. The reason for the change involves the placement of variables relative to the markstack. Figure 42 shows a record variable on the virtual machine and the Microengine. Since pass six calculates variable displacements by negating the displacement it receives from pass five, the wrong "end" of multiword variables will be referenced.

FIGURE 42. Record offsets for the Concurrent Pascal virtual machine (left) and the Microengine (right). Currently the MEPASCAL compiler calculates incorrect displacement values, as shown here.

Space must be reserved in the permanent variable space of monitors for the monitor gate (semaphore) address.

A change must be made to some pass prior to the code generator to convert variant record tags from bit numbers to bit values.

The kernel call mechanism must be designed. The mechanism could be calls to kernel routines which are well-known to the compiler.

The sequential compiler must be changed to pop the parameters from the stack after a call to the interface routine. The number of parameters pushed on the stack before calling the interface routine varies with the process ENTRY routine being invoked, so the RPU instruction which terminates the interface routine cannot pop the parameters.

Some optimizations of the final code are possible. First, the compiler always generates "long" jump instructions, regardless of the size of the distance operand. A smarter compiler would use long jumps only as required, but this would also complicate address calculations considerably. Second, only LDL (LoaD Local word) and LDO (LoaD glObal word) are generated to push onto the stack words from the local and global variable spaces. Shorter, faster instructions exist for pushing the first sixteen words in each space. Finally, it seems that the local word reserved for the old global base might be done away with. The local or global base register is saved as the static link, depending on the particular instruction used to call a routine. Perhaps the static link field of the markstack could be used to store the old global base by judicious use of the various routine-calling instructions. Before returning to the caller, the global base could be restored by using the LSL (Load Static Link) instruction.

Pass seven will have to be written. The pass must do the following:

-- Generate the header block for the code file;

-- Insert non-code fields (DATASIZE, for example) into the concurrent segment;

-- Build the constant pool for the concurrent segment;

-- Pack the code into words;

-- Calculate jump displacements;

-- Pad to the next block boundary;

-- For each interface, pack the code and pad to a block boundary.

Since pass seven is, at this time, only a program stub, a message indicating compilation errors is **always** generated, even if the source code is correct.

Finally, the compiler must be tested to ensure that it generates the intended Microengine code, and that the generated code actually behaves as expected when run on the hardware.

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APPENDIX A

SOFTWARE TOOLS

INTERMEDIATE CODE MNEMONICS PROGRAMS

MNEM (MNEMonics) is a sequential Pascal32 program, adapted from Robert Young's PASST, which prints intermediate and final code mnemonics for MCPASCAL. It was written to help better understand the architecture of the Concurrent Pascal virtual machine and the operation of the MCPASCAL compiler by generating reasonably intelligible operator mnemonics instead of the integer values produced by the mechanism built into MCPASCAL.

The program is invoked by specifying pass numbers as driver options. For example, entering the command

MCPASCAL SRCPGM,PR:,5,7

will cause mnemonics to be printed for the output of passes five and seven. After each compiler pass executes, the driver program runs MNEM if the pass terminated normally and the pass was specified by the user in the list of driver options. The pass number is sent to MNEM as a program parameter so that it knows which set of mnemonic literals should be generated. Currently, MNEM will only work on the output of passes five, six, and seven, since that of the other passes was not nearly as useful to the progress of the project. If the user requests mnemonics for one of the first four passes, the driver does not call MNEM, and prints a message to that effect. This allows the compilation to proceed to its conclusion, even though the request cannot be fulfilled.

After initializing I/O buffer variables and writing a pass header, MNEM calls the routine which will actually scan the appropriate intermediate code file. Each scanning routine (PASSx) contains a structured constant table which contains the mnemonic literals to be printed for each operator and the number of operands

which the operator uses. An operator read from the input file is used as an index into the table to fetch its mnemonic character string and number of arguments, which are then passed to a generalized instruction output routine. That routine prints the character string, reads the appropriate number of operands from the input file and prints them. PASST has a mechanism for printing kind, type, mode, and context mnemonics instead of the numerical values, but it was not carried over into MNEM since it did not seem to be worth the extra effort. No great inconvenience was experienced as a result. The operator formatting routine implemented that mechanism by using four parameters (KIND_ARG, TYPE_ARG, MODE_ARG, and CONTEXT_ARG). If, for example, an operator has three operands, and the second one is the addressing mode, MODE_ARG would have the value 2, and the other parameters would have the value 0. Since the mechanism is unused in MNEM, the four parameters are always 0.

Three instructions are handled by special routines when their operators are encountered. They are NEWLINE, CASEJUMP, and LONGCONSTANT. When NEWLINE is found, the current line of output is ended and a new one started in order to format the intermediate code listing in parallel with the source listing. The other two instructions have their own routines since they have a variable number of operands. Each one prints the operator mnemonic and whatever operands the instruction always has, determines the number of variable operands, and prints them.

The output line length used in the program is 70 positions, which corresponds to no physical device length. If the output is directed to a disk file and the editor is used to peruse it, a line length of 70 ensures that all of the text for a logical line will fit on one

physical line of the CRT. If a longer line length is used, a long line of text will spill onto a second CRT line-- something this writer finds extremely annoying.

MEMNEM (MicroEngine MNEMonics) prints intermediate code mnemonics for code generated by the MEPASCAL compiler. It is analogous to MNEM in the way it is invoked and in its general structure. Currently, it works only on the output of passes five and six. Pass five mnemonics are generated just as in MNEM, but the output of pass six is handled quite differently. In overall terms, the program scans the concurrent segment, scans the interface segment, prints some of the values in the PASSLINK record, and prints the contents of the heap tables.

The scan of the concurrent segment is straightforward, except that jumps and new lines are handled a little differently from MNEM. Microengine jump operators (UJPL and FJPL) have only one operand-- the displacement to the destination. However, when pass six generates these instructions the displacement is unknown. The jump operand emitted by it is the _label_ of the destination. It then puts the current value of the location counter into the code stream for pass seven's use in calculating the displacement. MEMNEM prints jump instructions just as any other, but then prints the location value and flags it as a non-Microengine entity. Pass six also generates NEWLINE operators which are not part of the Microengine instruction set. It emits a NEWLINE virtual operator _and_ the equivalent Microengine code when it finds a NEWLINE input operator. The equivalent code consists of two instructions-- one to push the line number onto the stack, and the other to store it into the first word of the local variable space. The store into local word 1 could be used to trigger output formatting actions as in MNEM, except that the line number itself would be on the

wrong output line. Rather than use a lookahead mechanism to check every "push" for a following STL(1), the virtual operator mechanism is used. When NEWLINE is found by MEMNEM nothing is printed, but the output pointer is forced to the next line so that the "push" and STL(1) will be printed there. The non-Microengine operators MESSAGE and EOM are flagged with asterisks.

After scanning the file which contains the concurrent code, MEMNEM starts taking its input from the file of interface segments. The number of segments in the file is fetched from the PASSLINK record and then the content of each segment is printed. Printing the interface segments is straightforwad, although not smooth or elegant, since they consist of so many kinds of objects besides just straight code.

Printing some of the PASSLINK fields is another straightforward, field-by-field operation. Only non-pointer fields are printed.

The heap tables are printed last. The routine DUMP_TABLE simply reads all the entries in a table and prints them. The entries for all tables except CONSTTABLE are printed as 16-bit integer values. CONSTTABLE entries are printed as byte values.

The source code for MNEM and MEMNEM follows.

```
1  "HCPASCAL MNEM - MNEMONIC TEST OUTPUT FORMATTING"
2
3  "% CONCURRENT := TRUE"
4
5  "#########
6   # PREFIX #
7   #########"
8
9  CONST NL= '(:10:)';   FF= '(:12:)';   CR= '(:13:)';   EM= '(:25:)';
10
11 CONST PAGELENGTH = 512 DIV 2;
12 TYPE PAGE = ARRAY (.1..PAGELENGTH.) OF SHORTINTEGER;
13 CONST BYTES_PER_INTEGER = 2;
14
15 CONST LINELENGTH = 70;
16 TYPE LINE = ARRAY (.1..LINELENGTH.) OF CHAR;
17
18 CONST IDLENGTH = 12;
19 TYPE IDENTIFIER = ARRAY (.1..IDLENGTH.) OF CHAR;
20
21 TYPE FILE = 1..5;
22
23 TYPE FILEKIND = (EMPTY, SCRATCH, ASCII, SEQCODE, CONCODE);
24
25 TYPE FILEATTR = RECORD
26      KIND: FILEKIND;
27      ADDR: INTEGER;
28      PROTECTED: BOOLEAN;
29      NOTUSED: ARRAY (.1..5.) OF INTEGER
30      END;
31
32 TYPE IODEVICE =
33   (TYPEDEVICE, DISKDEVICE, TAPEDEVICE, PRINTDEVICE, CARDDEVICE);
34
35 TYPE IOOPERATION = (INPUT, OUTPUT, MOVE, CONTROL);
36
37 TYPE IOARG = (WRITEEOF, REWIND, UPSPACE, BACKSPACE);
38
39 TYPE IORESULT =
40   (COMPLETE, INTERVENTION, TRANSMISSION, FAILURE,
41    ENDFILE, ENDMEDIUM, STARTMEDIUM);
42
43 TYPE IOPARAM = RECORD
44      OPERATION: IOOPERATION;
45      STATUS: IORESULT;
46      ARG: IOARG;
47      END;
48
49 TYPE TASKKIND = (INPUTTASK, JOBTASK, OUTPUTTASK);
50
51 TYPE ARGTAG =
52   (NILTYPE, BOOLTYPE, INTTYPE, IDTYPE, PTRTYPE);
53
54 TYPE POINTER = @BOOLEAN;
55
56 TYPE ARGTYPE = RECORD
57      CASE TAG: ARGTAG OF
58        NILTYPE, BOOLTYPE: (BOOL: BOOLEAN);
59        INTTYPE: (INT: INTEGER);
```

```
60          IDTYPE: (ID: IDENTIFIER);
61          PTRTYPE: (PTR: POINTER)
62        END;
63
64   CONST MAXARG = 10;
65   TYPE ARGLIST = ARRAY (.1..MAXARG.) OF ARGTYPE;
66
67   TYPE ARGSEQ = (INP, OUT);
68
69   TYPE PROGRESULT =
70     (TERMINATED, OVERFLOW, POINTERERROR, RANGEERROR, VARIANTERROR,
71      HEAPLIMIT, STACKLIMIT, CODELIMIT, TIMELIMIT, CALLERROR);
72
73   PROCEDURE READ(VAR C: CHAR);
74   PROCEDURE WRITE(C: CHAR);
75
76   PROCEDURE OPEN(F: FILE; ID: IDENTIFIER; VAR FOUND: BOOLEAN);
77   PROCEDURE CLOSE(F: FILE);
78   PROCEDURE GET(F: FILE; P: INTEGER; VAR BLOCK: UNIV PAGE);
79   PROCEDURE PUT(F: FILE; P: INTEGER; VAR BLOCK: UNIV PAGE);
80   FUNCTION LENGTH(F: FILE): INTEGER;
81
82   PROCEDURE MARK(VAR TOP: INTEGER);
83   PROCEDURE RELEASE(TOP: INTEGER);
84
85   PROCEDURE IDENTIFY(HEADER: LINE);
86   PROCEDURE ACCEPT(VAR C: CHAR);
87   PROCEDURE DISPLAY(C: CHAR);
88
89   PROCEDURE READPAGE(VAR BLOCK: UNIV PAGE; VAR EOF: BOOLEAN);
90   PROCEDURE WRITEPAGE(BLOCK: UNIV PAGE; EOF: BOOLEAN);
91   PROCEDURE READLINE(VAR TEXT: UNIV LINE);
92   PROCEDURE WRITELINE(TEXT: UNIV LINE);
93   PROCEDURE READARG(S: ARGSEQ; VAR ARG: ARGTYPE);
94   PROCEDURE WRITEARG(S: ARGSEQ; ARG: ARGTYPE);
95
96   PROCEDURE LOOKUP(ID: IDENTIFIER; VAR ATTR: FILEATTR;
97                    VAR FOUND: BOOLEAN);
98
99   PROCEDURE IOTRANSFER
100    (DEVICE: IODEVICE; VAR PARAM: IOPARAM; VAR BLOCK: UNIV PAGE);
101
102  PROCEDURE IOMOVE(DEVICE: IODEVICE; VAR PARAM: IOPARAM);
103
104  FUNCTION TASK: TASKKIND;
105
106  PROCEDURE RUN(ID: IDENTIFIER; VAR PARAM: ARGLIST;
107               VAR LINE: INTEGER; VAR RESULT: PROGRESULT);
108
109
110  PROGRAM MNEM(VAR PARAM: ARGLIST);
111
112  "SIMPLE TYPES AND CONSTANTS"
113
114  CONST LINE_LENGTH = 70;    "OUTPUT DEV LINE LENGTH"
115
116  TYPE CHAR8 = ARRAY [1..8] OF CHAR;
117
118  CONST
119    MIN_KIND = 0;   MAX_KIND = 0;
```

```
120    MIN_TYPE = 0;    MAX_TYPE = 0;
121    MIN_MODE = 0;    MAX_MODE = 0;
122    MIN_CONTEXT = 0;    MAX_CONTEXT = 0;
123
124
125    "VARIABLES"
126
127    VAR WORDS_IN: INTEGER;        "WORD IN IFL BUFFER"
128        IN_FILE: INTEGER;        "INPUT FILE"
129        PAGES_IN: INTEGER;       "CURRENT PAGE NUMBER"
130        PAGE_IN: PAGE;           "INPUT BUFFER"
131        PASS_NO: INTEGER;        "COMPILER PASS NUMBER"
132        OUT_COL_PTR: INTEGER;    "POSN ON OUTPUT LINE FOR BUF"
133        OUT_BUF_PTR: INTEGER;    "POSN IN BUFFER"
134        OUT_BUF: LINE;           "OPERATOR OUTPUT BUFFER"
135        FIRST_COL: INTEGER;      "FIRST COL TO USE FOR NON-NEWLINE"
136        MIN_COL_SEP: INTEGER;    "MIN BLANKS BETWEEN OPS"
137        COL_SEP: INTEGER;        "SEP BETWEEN COL BOUNDARIES"
138        EOM_OP: SHORTINTEGER;    "EOM OPERATOR"
139        LINE_OP: SHORTINTEGER;   "LINE NUMBER OPERATOR"
140        LCONST_OP: SHORTINTEGER; "LCONST OPERATOR"
141        OP: SHORTINTEGER;        "CURRENT INPUT OPERATOR"
142        CODELENGTH: SHORTINTEGER;
143        CODE_READ_SO_FAR: INTEGER;
144        I: INTEGER;
145
146
147    "INPUT FILE SUPPORT ROUTINES"
148
149    PROCEDURE READ_IFL (VAR ARG: SHORTINTEGER);
150    BEGIN
151      IF WORDS_IN = PAGELENGTH THEN BEGIN
152        GET(IN_FILE,PAGES_IN,PAGE_IN);
153        PAGES_IN:=PAGES_IN+1; WORDS_IN:=0;
154      END;
155      WORDS_IN:=WORDS_IN+1;
156      ARG:=PAGE_IN[WORDS_IN];
157      CODE_READ_SO_FAR := CODE_READ_SO_FAR + BYTES_PER_INTEGER;
158    END;
159
160    "BUFFER OUTPUT ROUTINES"
161
162    PROCEDURE NEXT_COL;
163    VAR I, J: INTEGER;
164    BEGIN "PAD TO START OF NEXT OP COLUMN"
165      I:=FIRST_COL;
166      I:=((OUT_COL_PTR+MIN_COL_SEP-FIRST_COL-1) DIV COL_SEP + 1)*COL_SEP;
167      IF I<0 THEN I:=0;
168      I:=I+FIRST_COL;
169      IF I+OUT_BUF_PTR-1 > LINE_LENGTH THEN BEGIN "WON'T FIT"
170        WRITE(NL); I:=FIRST_COL; OUT_COL_PTR:=1;
171      END;
172      FOR J:=OUT_COL_PTR TO I-1 DO WRITE(' ');
173      OUT_COL_PTR:=I;
174    END;
175
176    PROCEDURE WRITE_BUF (ALIGN: BOOLEAN);
177    VAR I: INTEGER;
178    BEGIN
179      IF ALIGN THEN NEXT_COL;
```

```
180   FOR I:=1 TO OUT_BUF_PTR-1 DO WRITE(OUT_BUF[I]);
181   OUT_COL_PTR:=OUT_COL_PTR+OUT_BUF_PTR-1;
182   OUT_BUF_PTR:=1;
183   END;
184
185   PROCEDURE WRITE_CHAR (C: CHAR);
186   BEGIN
187     IF OUT_BUF_PTR>LINE_LENGTH THEN
188       WRITE_BUF(TRUE);
189     OUT_BUF[OUT_BUF_PTR]:=C;
190     OUT_BUF_PTR:=OUT_BUF_PTR+1;
191   END;
192
193   "BUFFER FORMATTING ROUTINES"
194
195   PROCEDURE WRITE_CHAR8 (TEXT: CHAR8);
196   VAR I: INTEGER;
197   BEGIN
198     FOR I := 1 TO 8 DO
199       IF (TEXT[I] <> ' ')
200         THEN WRITE_CHAR (TEXT[I])
201   END;
202
203   PROCEDURE WRITE_HEX_CHAR (VAL: SHORTINTEGER);
204   BEGIN
205     IF VAL > 9 THEN WRITE_CHAR(CHR(VAL-10+ORD('A')))
206     ELSE WRITE_CHAR(CHR(VAL+ORD('0')));
207   END;
208
209   PROCEDURE WRITE_HEX (VAL: SHORTINTEGER; N: SHORTINTEGER);
210   TYPE TAGS = (TAG1, TAG2);
211   VAR I: SHORTINTEGER;
212       X: RECORD CASE TAGS OF
213         TAG1: (INT: SHORTINTEGER);
214         TAG2: (BYTES: ARRAY [1..2] OF BYTE);
215       END;
216   BEGIN
217     X.INT:=VAL;
218     FOR I:=3-N TO 2 DO BEGIN
219       WRITE_HEX_CHAR(X.BYTES[I] DIV 16);
220       WRITE_HEX_CHAR(X.BYTES[I] MOD 16);
221     END;
222   END;
223
224   PROCEDURE WRITE_INT (VAL: SHORTINTEGER);
225   VAR A: ARRAY [1..10] OF CHAR;
226       I, J, REM: INTEGER;
227   BEGIN
228     REM:=VAL; I:=1;
229     REPEAT
230       A[I]:=CHR(ABS(REM MOD 10)+ORD('0'));
231       I:=I+1; REM:=REM DIV 10;
232     UNTIL REM = 0;
233     IF VAL<0 THEN WRITE_CHAR('-');
234     FOR J:=I-1 DOWNTO 1 DO WRITE_CHAR(A[J]);
235   END;
236
237   PROCEDURE WRITE_INTR (VAL: INTEGER; N: INTEGER);
238   "PRINTS THE VALUE OF 'VAL' IN A FIELD OF 'N' POSITIONS.
239   VALUE WILL BE RIGHT JUSTIFIED IN THE FIELD."
```

```
240  VAR A: ARRAY [1..10] OF CHAR;
241      I, J, REM: INTEGER;
242  BEGIN
243      REM:=VAL; I:=1;
244      REPEAT
245          A[I]:=CHR(ABS(REM MOD 10)+ORD('0'));
246          I:=I+1; REM:=REM DIV 10;
247      UNTIL REM = 0;
248      IF VAL < 0 THEN BEGIN
249          A[I]:='-'; I:=I+1;
250      END;
251      FOR J:=I TO N DO WRITE_CHAR(' ');
252      FOR J:=I-1 DOWNTO 1 DO WRITE_CHAR(A[J]);
253  END;
254
255
256  PROCEDURE INDEXERROR (TEXT:CHAR8; BADINDEX: SHORTINTEGER);
257  "IF THIS PROCEDURE IS INVOKED IT PROBABLY MEANS SOME OPERATOR
258   WHICH WAS ENCOUNTERED EARLIER HAS THE WRONG 'NUM_ARGS' (NO. OF
259   OPERANDS) VALUE ASSIGNED TO IT IN THE OPERATOR TABLE."
260  VAR I: INTEGER;
261  BEGIN
262      WRITE_CHAR8 ('%%%$$$$$');
263      WRITE_CHAR8 ('INDEX   ');
264      WRITE_CHAR (' ');
265      WRITE_CHAR8 ('ERROR   ');
266      WRITE_CHAR8 ('%%%$$$$$');
267      WRITE_CHAR8 (TEXT);
268      WRITE_CHAR ('=');
269      WRITE_CHAR (' ');
270      WRITE_INT (BADINDEX);
271      WRITE_CHAR (' ');
272      WRITE_CHAR (' ');
273  END;
274
275
276  "GENERAL OPERATOR FORMATTING ROUTINE"
277
278  PROCEDURE WRITE_OP (OP_NAME: CHAR8; NUM_ARGS, KIND_ARG, TYPE_ARG,
279                      MODE_ARG, CONTEXT_ARG: SHORTINTEGER);
280  VAR ARG_NO, ARG_VAL, I: SHORTINTEGER;
281  BEGIN
282      WRITE_CHAR8 (OP_NAME);
283      FOR ARG_NO := 1 TO NUM_ARGS DO
284          BEGIN
285          IF ARG_NO = 1
286              THEN WRITE_CHAR ( '(' )
287              ELSE WRITE_CHAR ( ',' );
288          READ_IFL (ARG_VAL);
289          IF
290              ARG_NO = KIND_ARG THEN
291                  BEGIN
292                  IF (ARG_VAL<MIN_KIND) OR (ARG_VAL>MAX_KIND)
293                      THEN INDEXERROR ('BAD-KIND', ARG_VAL)
294                      ELSE WRITE_CHAR8 ('NO_KIND_')
295                  END    ELSE IF
296              ARG_NO = TYPE_ARG THEN
297                  BEGIN
298                  IF (ARG_VAL<MIN TYPE) OR (ARG_VAL>MAX_TYPE)
299                      THEN INDEXERROR ('BAD-TYPE', ARG_VAL)
```

```
300         ELSE WRITE_CHAR8 ('NO_TYPE_')
301     END  ELSE IF
302  ARG_NO = MODE_ARG THEN
303     BEGIN
304     IF (ARG_VAL<MIN_MODE) OR (ARG_VAL>MAX_MODE)
305        THEN INDEXERROR ('BAD-MODE'; ARG_VAL)
306        ELSE WRITE CHAR8 ('NO_MODE_')
307     END  ELSE IF
308  ARG_NO = CONTEXT_ARG THEN
309     BEGIN
310     IF (ARG_VAL<MIN_CONTEXT) OR (ARG_VAL>MAX_CONTEXT)
311        THEN INDEXERROR ('BAD-CTXT', ARG_VAL)
312        ELSE WRITE_CHAR8 ('NO_CNTXT')
313     END  ELSE IF
314     (ARG_VAL < -32768) OR (ARG_VAL > 32767)
315        THEN WRITE_HEX (ARG_VAL, 2)
316        ELSE WRITE_INT (ARG_VAL)
317     END; "OF BEGIN"
318  IF NUM_ARGS>0 THEN WRITE_CHAR(')');
319  WRITE_BUF(TRUE);
320 END;
321
322 "SPECIAL CASE ROUTINE FOR NEW LINE OPERATOR"
323
324 PROCEDURE NEW_LINE;
325 VAR ARG: SHORTINTEGER;
326 BEGIN
327  READ_IFL(ARG); "LINE NUMBER"
328  IF OUT_COL_PTR<>1 THEN WRITE(NL);
329  WRITE(NL);
330  OUT_COL_PTR:=1;
331  WRITE_CHAR8('LINE        '); WRITE_INTR(ARG,5);
332  WRITE_BUF(FALSE);  "WRITE WITH NO COL ALIGN"
333 END;
334
335 "SPECIAL CASE ROUTINE FOR LONG CONSTANT"
336
337 PROCEDURE WRITE_LCONST;
338 VAR LEN, I, N, ARG: SHORTINTEGER;
339 BEGIN
340  WRITE_CHAR8('LNGCONST'); WRITE_CHAR('(');
341  READ_IFL(LEN);
342  WRITE_INT(LEN); WRITE_CHAR(',');
343  WRITE_CHAR8 ('HEX-VAL:');
344  NEXT_COL; N:=(LEN-1) DIV 2 + 1;
345  FOR I:=1 TO N DO BEGIN
346   READ IFL(ARG);
347   IF OUT_COL_PTR+OUT_BUF_PTR-1+9 > LINE_LENGTH THEN BEGIN
348    WRITE_BUF(FALSE); WRITE(NL); OUT_COL_PTR:=1; NEXT_COL;
349   END;
350   WRITE_HEX(ARG,2);
351   IF I < N THEN WRITE_CHAR(',') ELSE WRITE_CHAR(')');
352  END;
353  WRITE_BUF(FALSE);
354 END;
355
356 "COMMON INITIALIZATION"
357
358 PROCEDURE INITIALIZE;
```

```
360  BEGIN
361    WORDS_IN:=PAGELENGTH;
362    PAGES_IN:=1;
363    PASS_NO:=PARAM[10].INT;
364    OUT_COL_PTR:=1;
365    OUT_BUF_PTR:=1;
366    FIRST_COL:=12;
367    MIN_COL_SEP:=3;
368    COL_SEP:=12;
369  END;
370
371  "WRITE PASS LISTING HEADER"
372
373  PROCEDURE WRITE_TEXT (TEXT: LINE);
374  VAR I: INTEGER;
375  BEGIN
376    I:=1;
377    WHILE TEXT[I] <> '$' DO BEGIN
378      WRITE(TEXT[I]); I:=I+1;
379    END;
380  END;
381
382  PROCEDURE WRITE_HEADER;
383  VAR I, J: INTEGER;
384      C: CHAR;
385  BEGIN
386    C:=CHR(PASS_NO+ORD('0'));
387    WRITE (FF);
388    WRITE(NL);
389    FOR I:=1 TO LINE_LENGTH DO WRITE(C);
390    WRITE(NL);
391    WRITE(C);
392    J:=LINE_LENGTH DIV 2 - 17;
393    FOR I:=2 TO J-1 DO WRITE(' ');
394    WRITE_TEXT('MCPASCAL INTERMEDIATE CODE PASS $');
395    WRITE(C);
396    FOR I:=J+33 TO LINE_LENGTH-1 DO WRITE(' ');
397    WRITE(C); WRITE(NL);
398    FOR I:=1 TO LINE_LENGTH DO WRITE(C);
399    WRITE(NL);
400  END;
401
402  PROCEDURE WRITE_56 CASE (OP_NAME: CHAR8; PASS3: INTEGER);
403  VAR
404    I: INTEGER;
405    MIN,
406    MAX,
407    MAX_MINUS_MIN,
408    LOCATION,
409    STMT_LABEL:    SHORTINTEGER;
410
411  BEGIN
412    WRITE_CHAR8(OP_NAME);
413    WRITE_CHAR( '(' );
414    READ_IFL(MIN);
415    WRITE_INT(MIN);
416    WRITE_CHAR( ',' );
417
418    IF PASS = 5
419    THEN BEGIN
```

```
420        READ_IFL (MAX);
421        WRITE_INT (MAX);
422        MAX_MINUS_MIN := MAX - MIN
423      END
424    ELSE BEGIN  "PASS 6"
425        READ_IFL(MAX_MINUS_MIN);
426        WRITE_INT(MAX_MINUS_MIN);
427        WRITE_CHAR ( ',' );
428        READ_IFL (LOCATION);
429        WRITE_INT (LOCATION)
430      END;
431    FOR I := 0 TO MAX_MINUS_MIN DO
432      BEGIN
433        WRITE_CHAR(',');
434        READ_IFL(STMT_LABEL);
435        WRITE_INT(STMT_LABEL);
436      END;
437    WRITE_CHAR( ')' );
438  END;
439
440
441  PROCEDURE WRITE_PASS7 CASE (OP_NAME: CHAR8);
442  VAR
443      I: INTEGER;
444      MIN,
445      MAX_MINUS_MIN,
446      NUM_OF_DISTANCES,
447      ARG:               SHORTINTEGER;
448
449  BEGIN
450    WRITE_CHAR8(OP_NAME);
451    WRITE_CHAR ( '(' );
452    READ_IFL(MIN);
453    WRITE_INT(MIN);
454    WRITE_CHAR ( ',' );
455
456    READ_IFL(MAX_MINUS_MIN);
457    WRITE_INT(MAX_MINUS_MIN);
458    NUM_OF_DISTANCES := MAX_MINUS_MIN + 1;
459    FOR I := 1 TO NUM_OF_DISTANCES DO
460      BEGIN
461        WRITE_CHAR(',');
462        READ_IFL(ARG);
463        WRITE_INT(ARG);
464      END;
465    WRITE_CHAR( ')' );
466  END;
467
468
469  PROCEDURE PASS5;
470  CONST MIN_OP5 = 0;          MAX_OP5 = 48;
471  CONST
472    PASS5_TABLE = (
473      'PUSHCNST',  1,      'PUSHVAR',  3,          " 1"
474      'PUSHIND',   1,      'PUSHADDR', 2,
475      'FIELD',     1,      'INDEX',    3,
476      'POINTER',   0,      'VARIANT',  2,
477      'RANGE',     2,      'ASSIGN',   1,          " 9"
478      'ASSGNTAG',  0,      'COPY',     1,
479      'NEW',       0,      'NOT',      0,
```

```
480      'AND      ',  1,        'OR       ',  1,
481      'NEG      ',  1,        'ADD      ',  1,
482      'SUB      ',  1,        'MUL      ',  1,
483      'DIV      ',  0,        'MOD      ',  1,
484      'INVALID  ',  0,        'INVALID  ',  0,
485      'FUNCTION ',  2,        'BUILDSET ',  0,
486      'COMPARE  ',  2,        'CHKSTRUC ',  2,           "19"
487      'FUNCVAL  ',  2,        'DEFLABEL ',  1,
488      'JUMP     ',  1,        'FALSJUMP ',  1,
489      'CASEJUMP ',  0,        'INITVAR  ',  0,
490      'CALL     ',  3,        'ENTER    ',  5,           "29"
491      'RETURN   ',  1,        'POP      ',  1,
492      'NEWLINE  ',  1,        'ERROR    ',  0,
493      'LNGCONST ',  0,        'MESSAGE  ',  2,
494      'INCRMENT ',  0,        'DECRMENT ',  0,           "39"
495      'PROCDURE ',  1,        'INIT     ',  4,
496      'PUSHLABL ',  1,        'CALLPROG ',  0,           "47"
497      'EOM      ',  1);                                  "48"
498
499                ARRAY[MIN_OP5 .. MAX_OP5] OF
500            RECORD
501               OP_NAME: CHAR8;
502               NUM_ARGS: BYTE
503            END;
504
505   CONST
506      EOM5 = 48;
507      NEWLINE5 = 38;
508      CASEJUMP5 = 32;
509      LONGCONSTANT5 = 40;
510
511   BEGIN
512      IN_FILE := 2;
513      COL_SEP := 3;
514      REPEAT
515         READ_IFL (OP);
516         WITH PASS5_TABLE[OP] DO
517            IF (OP < MIN_OP5) OR (OP > MAX_OP5)
518            THEN INDEXERROR ('BAD_OP ', OP)
519            ELSE IF OP = NEWLINE5
520                 THEN NEW_LINE
521            ELSE IF OP = CASEJUMP5
522                 THEN WRITE_56_CASE (OP_NAME, PASS_NO)
523            ELSE IF OP = LONGCONSTANT5
524                 THEN WRITE LCONST
525            ELSE WRITE_OP (OP_NAME, NUM_ARGS, 0, 0, 0, 0)
526      UNTIL OP = EOM5;
527      WRITE (M.)
528   END;
529
530
531
532   PROCEDURE PASS6;
533   CONST    MIN_OP6 = 0;       MAX_OP6 = 113;
534   CONST    PASS6_TABLE = (
535
536      'CONSTADD ',  1,        'LOCALADD ',  1,        "1"
537      'GLOBLADD ',  1,        'PUSHCONS ',  1,
538      'PUSHLOCL ',  1,        'PUSHGLBL ',  1,
539      'PUSHIND  ',  0,        'PUSHBYTE ',  0,
```

```
                                    470  470
                                     •   116
                                     •   214

                                    128
                                    137

                                    149  141
                                    472  141
                                    141  470      141   470
                                    256  141
                                    141  507
                                    324
                                    141  508
                                    402  501      131
                                    141  509
                                    337
                                    278  501      502
                                    141  506
                                     74    9

                                     •   •
                                     •   •
                                     •   •
```

```
540    'PUSHREAL',  0,   'PUSHSET ',  0,    "     "
541    'FIELD   ',  1,   'INDEX   ',  1,    "   9 "
542    'POINTER ',  0,   'VARIANT ',  0,    "     "
543    'RANGE   ',  2,   'COPYBYTE',  2,    "     "
544    'COPYWORD',  0,   'COPYREAL',  0,    "     "
545    'COPYSET ',  0,   'COPYTAG ',  1,    "  19 "
546    'COPYSTRU',  1,   'NEW     ',  2,    "     "
547    'NEWINIT ',  2,   'NOT     ',  0,    "     "
548    'ANDWORD ',  0,   'ANDSET  ',  0,    "     "
549    'ORWORD  ',  0,   'ORSET   ',  0,    "     "
550    'NEGWORD ',  0,   'NEGREAL ',  0,    "  29 "
551    'ADDWORD ',  0,   'ADDREAL ',  0,    "     "
552    'SUBWORD ',  0,   'SUBREAL ',  0,    "     "
553    'SUBTRSET',  0,   'MULWORD ',  0,    "     "
554    'MULREAL ',  0,   'DIVWORD ',  0,    "     "
555    'DIVREAL ',  0,   'MODWORD ',  0,    "  39 "
556    'BUILDSET',  0,   'INSET   ',  0,    "     "
557    'LSWORD  ',  0,   'EQWORD  ',  0,    "     "
558    'GRWORD  ',  0,   'NLWORD  ',  0,    "     "
559    'NEWORD  ',  0,   'NOWORD  ',  0,    "     "
560    'LSREAL  ',  0,   'EQREAL  ',  0,    "  49 "
561    'GRREAL  ',  0,   'NLREAL  ',  0,    "     "
562    'NEREAL  ',  0,   'NOREAL  ',  0,    "     "
563    'EQSET   ',  0,   'NLSET   ',  0,    "     "
564    'NESET   ',  0,   'NGSET   ',  0,    "     "
565    'LSSTRUCT',  1,   'EQSTRUCT',  1,    "  59 "
566    'GRSTRUCT',  1,   'NLSTRUCT',  1,    "     "
567    'NESTRUCT',  1,   'NOSTRUCT',  1,    "     "
568    'FUNCVALU',  1,   'JUMP    ',  2,    "     "
569    'FALSEJMP',  2,   'CASEJUMP',  2,    "     "
570    'INITVAR ',  1,   'CALL    ',  2,    "  69 "
571    'CALLSIS ',  1,   'ENTER   ',  4,    "     "
572    'EXIT    ',  0,   'ENTRPROG',  4,    "     "
573    'EXITPROG',  0,   'BEGNCLAS',  4,    "     "
574    'ENDCLASS',  0,   'ENTRCLAS',  4,    "     "
575    'EXITCLAS',  0,   'BEGINMON',  4,    "  79 "
576    'ENDMON  ',  0,   'ENTERMON',  4,    "     "
577    'EXITMON ',  0,   'BEGNPRCS',  1,    "     "
578    'ENDPRCS ',  0,   'ENTRPRCS',  0,    "     "
579    'NEWLINE ',  1,   'POP     ',  1,    "     "
580    'DECWORD ',  0,   'INCRWORD',  1,    "  89 "
581    'INITMON ',  0,   'INITCLAS',  3,    "     "
582    'PUSHLABL',  2,   'INITPROC',  5,    "     "
583    'TRUNCREA',  0,   'CALLPROG',  0,    "     "
584    'ABSREAL ',  0,   'SUCCWORD',  0,    "     "
585    'PREDWORD',  0,   'CONVWORD',  0,    "  99 "
586    'EMPTY   ',  0,   'ATTRIBUT',  0,    "     "
587    'REALTIME',  0,   'DELAY   ',  0,    "     "
588    'CONTINUE',  0,   'I/O     ',  0,    "     "
589    'START   ',  0,   'STOP    ',  0,    " 109 "
590    'SETHEAP ',  0,   'WAIT    ',  0,    "     "
591    'MESSAGE ',  3,   'EOM     ',  0);   " 113 "

592    ARRAY [MIN_OP6..MAX_OP6] OF
593      RECORD
594        OP_NAME: CHAR8;
595        NUM_ARGS: BYTE;
596      END;
```

```
600  CONST EOM6 = 113;
601        NEWLINE6 = 80;
602        CASEJUMP6 = 67;
603
604
605  BEGIN
606    IN_FILE:=1;
607    COL_SEP:=3;
608    REPEAT
609      READ_IFL(OP);
610      WITH PASS6_TABLE[OP] DO
611        IF (OP < MIN_OP6) OR (OP > MAX_OP6)
612          THEN INDEXERROR ('BAD-OP ', OP)
613          ELSE IF OP = NEWLINE6
614            THEN NEW LINE
615            ELSE IF OP = CASEJUMP6
616              THEN WRITE_56_CASE (OP_NAME, PASS_NO)
617              ELSE WRITE_OP (OP_NAME, NUM_ARGS, 0, 0, 0, 0);
618    UNTIL OP = EOM6;
619    WRITE(M.);
620  END;
621
622
623  PROCEDURE PASS7;
624  CONST
625    PASS7_TABLE = (
626      'CONSTADD', 1,  'LOCALADD' 1,    2"
627      'GLOBLADD', 1,  'PUSHCONS', 1,    "
628      'PUSHLOCL', 1,  'PUSHGLBL', 1,    "
629      'PUSHIND ', 0,  'PUSHBYTE', 0,   10"
630      'PUSHREAL', 0,  'PUSHSET ', 0,    "
631      'FIELD   ', 1,  'INDEX   ', 3,    "
632      'POINTER ', 0,  'VARIANT ', 2,    "
633      'RANGE   ', 2,  'COPYBYTE', 0,    "
634      'COPYWORD', 0,  'COPYREAL', 0,    "
635      'COPYSET ', 0,  'COPYTAG ', 1,   20"
636      'COPYSTRU', 1,  'NEW     ', 2,    "
637      'NEWINIT ', 2,  'NOT     ', 0,    "
638      'ANDWORD ', 0,  'ANDSET  ', 0,    "
639      'ORWORD  ', 0,  'ORSET   ', 0,    "
640      'NEGWORD ', 0,  'NEGREAL ', 0,   30"
641      'ADDWORD ', 0,  'ADDREAL ', 0,    "
642      'SUBWORD ', 0,  'SUBREAL ', 0,    "
643      'SUBTRSET', 0,  'MULWORD ', 0,    "
644      'MULREAL ', 0,  'DIVWORD ', 0,    "
645      'DIVREAL ', 0,  'MODWORD ', 0,   40"
646      'BUILDSET', 0,  'INSET   ', 0,    "
647      'LSWORD  ', 0,  'EQWORD  ', 0,    "
648      'GRWORD  ', 0,  'NLWORD  ', 0,    "
649      'NEWORD  ', 0,  'NGWORD  ', 0,    "
650      'LSREAL  ', 0,  'EQREAL  ', 0,   50"
651      'GRREAL  ', 0,  'NLREAL  ', 0,    "
652      'NEREAL. ', 0,  'NGREAL  ', 0,    "
653      'EQSET   ', 0,  'NLSET   ', 0,    "
654      'NESET   ', 0,  'NGSET   ', 0,    "
655      'LSSTRUCT', 1,  'EQSTRUCT', 0,   60"
656      'GRSTRUCT', 1,  'NLSTRUCT', 1,    "
657      'NESTRUCT', 1,  'NGSTRUCT', 1,    "
658      'FUNCVALU', 1,  'JUMP    ', 1,    "
659      'FALSEJMP', 1,  'CASEJUMP', 2,    "
```

```
660          'INITVAR ',  1,  'CALL    ',  1,  " 70"
661          'CALLTS  ',  1,  'ENTER   ',  1,  "   "
662          'EXIT    ',  0,  'ENTRPROG',  =,  "   "
663          'EXITPROG',  0,  'BEGNCLAS',  =,  "   "
664          'ENDCLASS',  0,  'ENTRCLAS',  =,  " 80"
665          'EXITCLAS',  0,  'BEGINMON',  =,  "   "
666          'ENDMON  ',  0,  'ENTERMON',  =,  "   "
667          'EXITMON ',  0,  'BEGNPRCS',  -,  "   "
668          'ENDPRCS ',  0,  'ENTRPRCS',  =,  " 90"
669          'EXITPRCS',  0,  'POP     ',  0,  "   "
670          'NEWLINE ',  1,  'INCRWORD',  0,  "   "
671          'DECRWORD',  0,  'INITCLAS',  2,  "   "
672          'INITMON ',  2,  'INITPROC',  =,  "100"
673          'PUSHLABL',  1,  'CALLPROG',  0,  "   "
674          'TRUNCREA',  0,  'ABSWORD ',  0,  "   "
675          'ABSREAL ',  0,  'SUCCWORD',  0,  "   "
676          'PREDWORD',  0,  'CONVWORD',  0,  "100"
677          'EMPTY   ',  0,  'ATTRIBUT',  0,  "   "
678          'REALTIME',  0,  'DELAY   ',  0,  "   "
679          'CONTINUE',  0,  'I/O     ',  0,  "   "
680          'START   ',  0,  'STOP    ',  0,  "110"
681          'SETHEAP ',  0,  'WAIT    ',  0); "12"
682
683      ARRAY [1..112] OF
684          RECORD
685              OP_NAME: CHAR8;
686              NUM_ARGS: SHORTINTEGER;
687          END;
688
689  NEWLINE7 = 178;
690  CASEJUMP7 =136;
691
692  VAR
693      HALF_OP: SHORTINTEGER;
694      PROGLENGTH, STACKLENGTH, VARLENGTH, CONSTANTS: SHORTINTEGER;
695
696
697  BEGIN
698      IN_FILE := 3;
699      COL_SEP := 1;
700
701      READ_IFL(PROGLENGTH);
702      READ_IFL(CODELENGTH);
703      READ_IFL(STACKLENGTH);
704      READ_IFL(VARLENGTH);
705      READ_IFL(CONSTANTS);
706      CODE_READ_SO_FAR := 0;
707
708      WRITE_CHAR8 ('PROGLEN=');   WRITE_INT (PROGLENGTH);
709      WRITE_CHAR8(      (:10:)' );
710      WRITE_CHAR8 ('CODELEN=');   WRITE_INT (CODELENGTH);
711      WRITE_CHAR8(      (:10:)' );
712      WRITE_CHAR8 ('STCKLEN=');   WRITE_INT (STACKLENGTH);
713      WRITE_CHAR8(      (:10:)' );
714      WRITE_CHAR8 ('VARLEN =');   WRITE_INT (VARLENGTH);
715      WRITE_CHAR8(      (:10:)' );
716      WRITE_CHAR8 ('CONSTS =');   WRITE_INT (CONSTANTS);
717      WRITE_CHAR8 (     (:10:)' );
718
719      REPEAT
```

```
720       READ_IFL(OP);
721       HALF_OP := OP DIV 2;
722       IF (HALF_OP <= 0)  OR  (HALF_OP > 112)
723       THEN BEGIN
724           WRITE_CHAR8 ('BAD OP: ');
725           WRITE_INT (OP);
726           WRITE_CHAR8('          (:10)' );
727           FOR I := 1 TO PAGELENGTH + 10 DO WRITE_CHAR(' ');
728       END;
729       WITH PASS7_TABLE[HALF_OP] DO
730           IF OP = NEWLINE7
731           THEN NEW_LINE
732           ELSE IF OP = CASEJUMP7
733               THEN WRITE_PASS7_CASE (OP_NAME)
734               ELSE WRITE_OP (OP_NAME, NUM_ARGS, 0, 0, 0, 0)
735   UNTIL CODE_READ_SO_FAR >= CODELENGTH;
736   WRITE_CHAR8 ('      (:10)' );
737   WRITE_CHAR8 ('END-CODE');
738   WRITE_CHAR (NL);
739   WRITE_CHAR8 ('CONSTANT');
740   WRITE_CHAR8 ('S-IN-HEX');
741   WRITE_CHAR8 ('      (:10)' );
742   REPEAT
743       READ_IFL(OP);
744       IF OUT_COL_PTR + OUT_BUF_PTR - 1 + 9 > LINE_LENGTH
745       THEN BEGIN
746           WRITE_BUF(FALSE);
747           WRITE (NL);
748           OUT_COL_PTR := 1;
749           NEXT_COL;
750       END;
751       WRITE_HEX (OP, 2);
752       WRITE_CHAR (' ');
753       WRITE_CHAR (' ');
754   UNTIL CODE_READ_SO_FAR >= PROGLENGTH - 8;
755   FOR I := 1 TO PAGELENGTH +10 DO WRITE_CHAR(' ');
756   END;
757
758
759
760 BEGIN
761   INITIALIZE;
762   WRITE_HEADER;
763   CASE PASS_NO OF
764       1: "PASS1";
765       2: "PASS2";
766       3: "PASS3";
767       4: "PASS4";
768       5: PASS5;
769       6: PASS6;
770       7: PASS7;
771       8, 9: "NOT IMPLEMENTED";
772   END;
773 END.
```

CROSS REFERENCE   * IS DEF   = IS ASG

-A-

A	225*	230=	234	240*	245=	249=	252						
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ACCEPT	86*												
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-W-

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	398 399 527 619 747
WRITEARG	37
WRITEEOF	37
WRITELINE	92*
WRITEPAGE	90*
WRITE_56_CAS	402* 522 616
WRITE_BUF	176* 188 200 205 206 233 251 252 264 268 269 271 272 286 287 318 340 342
WRITE_CHAR	185* 200 205 413 416 427 433 437 451 454 461 465 727 738 752 753 755 450
WRITE_CHARB	195* 262 263 265 266 267 282 294 300 306 312 331 340 343 412
	710 711 712 713 714 715 716 717 724 726 736 737 739 740 741
WRITE_HEADER	382* 762
WRITE_HEX	209* 315 350 751
WRITE_HEX_CH	203* 219 220
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-X-

Identifier	References
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END XREF 230 IDENTIFIERS   948 TOTAL REFERENCES
129 COLLISIONS.

```
 1  "MEPASCAL MEMMEM - MNEMONIC TEST OUTPUT FORMATTING"
 2  "FOR WESTERN DIGITAL PASCAL MICROENGINE"
 3
 4  "% CONCURRENT := TRUE"
 5
 6  "#########
 7   # PREFIX #
 8   #########"
 9
10  TYPE FULLWORD = INTEGER;
11       INTEGER  = SHORTINTEGER;
12
13  CONST NL= '(:10:)';  FF= '(:12:)';  CR= '(:13:)';  EM= '(:25:)';
14
15  CONST PAGELENGTH = 512 DIV 2;
16  TYPE PAGE = ARRAY (.1..PAGELENGTH.) OF INTEGER;
17  CONST BYTES_PER_INTEGER = 2;
18
19  CONST LINELENGTH = 70;
20  TYPE LINE = ARRAY (.1..LINELENGTH.) OF CHAR;
21
22  CONST IDLENGTH = 12;
23  TYPE IDENTIFIER = ARRAY (.1..IDLENGTH.) OF CHAR;
24
25  TYPE FILE = 1..5;
26
27  TYPE FILEKIND = (EMPTY, SCRATCH, ASCII, SEQCODE, CONCODE);
28
29  TYPE FILEATTR = RECORD
30                  KIND: FILEKIND;
31                  ADDR: INTEGER;
32                  PROTECTED: BOOLEAN;
33                  NOTUSED: ARRAY (.1..5.) OF INTEGER
34                  END;
35
36  TYPE IODEVICE =
37       (TYPEDEVICE, DISKDEVICE, TAPEDEVICE, PRINTDEVICE, CARDDEVICE);
38
39  TYPE IOOPERATION = (INPUT, OUTPUT, MOVE, CONTROL);
40
41  TYPE IOARG = (WRITEEOF, REWIND, UPSPACE, BACKSPACE);
42
43  TYPE IORESULT =
44       (COMPLETE, INTERVENTION, TRANSMISSION, FAILURE,
45       ENDFILE, ENDMEDIUM, STARTMEDIUM);
46
47  TYPE IOPARAM = RECORD
48                 OPERATION: IOOPERATION;
49                 STATUS: IORESULT;
50                 ARG: IOARG
51                 END;
52
53  TYPE TASKKIND = (INPUTTASK, JOBTASK, OUTPUTTASK);
54
55  CONST MAXWORD = 100;
56
57  TYPE
58  POINTER = @INTEGER;
59
```

```
60  TABLEPTR = @TABLE;
61  TABLE = RECORD
62    NEXTPORTION: TABLEPTR;
63    CONTENTS: ARRAY[1..MAXWORD] OF INTEGER
64    END;
65
66  TABLEPART = RECORD
67    SEGDISTANCE, STACKLENGTH: INTEGER;
68    JUMPTABLE, CONSTTABLE, XJPTABLE,
69    EXITICTABLE, DATASIZETABLE: TABLEPTR
70    END;
71
72  TABLESPTR = @TABLEPART;
73
74  OPTION = 0..8;
75
76  CONST MAXINTFAC = 14;
77  TYPE
78  IFPTR = @IFINFO;
79  IFINFO = RECORD
80    INTERFACES: INTEGER;
81    INTERFACESIZES: ARRAY[1..MAXINTFAC] OF INTEGER
82    END;
83
84  PASSPTR = @PASSLINK;
85  PASSLINK = RECORD
86    OPTIONS: SET OF OPTION;
87    LABELS, BLOCKS, CONSTANTS, XJP_OFFSETS: INTEGER;
88    RESETPOINT: FULLWORD;
89    TABLES: TABLESPTR;
90    INTERFACE: IFPTR
91    END;
92
93  TYPE ARGTAG =
94    (NILTYPE, BOOLTYPE, INTTYPE, IDTYPE, PTRTYPE);
95
96  TYPE ARGTYPE = RECORD
97      CASE TAG: ARGTAG OF
98        NILTYPE, BOOLTYPE: (BOOL: BOOLEAN);
99        INTTYPE: (INT: INTEGER);
100       IDTYPE: (ID: IDENTIFIER);
101       PTRTYPE: (PTR: PASSPTR)
102      END;
103
104 CONST MAXARG = 10;
105 TYPE ARGLIST = ARRAY (.1..MAXARG.) OF ARGTYPE;
106
107 TYPE ARGSEQ = (INP, OUT);
108
109 TYPE PROGRESULT =
110   (TERMINATED, OVERFLOW, POINTERERROR, RANGEERROR, VARIANTERROR,
111   HEAPLIMIT, STACKLIMIT, CODELIMIT, TIMELIMIT, CALLERROR);
112
113 PROCEDURE READ(VAR C: CHAR);
114 PROCEDURE WRITE(C: CHAR);
115
116 PROCEDURE OPEN(F: FILE; ID: IDENTIFIER; VAR FOUND: BOOLEAN);
117 PROCEDURE CLOSE(F: FILE);
118 PROCEDURE GET(F: FILE; P: INTEGER; VAR BLOCK: UNIV PAGE);
119 PROCEDURE PUT(F: FILE; P: INTEGER; VAR BLOCK: UNIV PAGE);
```

```
120 FUNCTION LENGTH(F: FILE): INTEGER;
121
122 PROCEDURE MARK(VAR TOP: INTEGER);
123 PROCEDURE RELEASE(TOP: INTEGER);
124
125 PROCEDURE IDENTIFY(HEADER: LINE);
126 PROCEDURE ACCEPT(VAR C: CHAR);
127 PROCEDURE DISPLAY(C: CHAR);
128
129 PROCEDURE READPAGE(VAR BLOCK: UNIV PAGE; VAR EOF: BOOLEAN);
130 PROCEDURE WRITEPAGE(BLOCK: UNIV PAGE; EOF: BOOLEAN);
131 PROCEDURE READLINE(VAR TEXT: UNIV LINE);
132 PROCEDURE WRITELINE(TEXT: UNIV LINE);
133 PROCEDURE READARG(S: ARGSEQ; VAR ARG: ARGTYPE);
134 PROCEDURE WRITEARG(S: ARGSEQ; ARG: ARGTYPE);
135
136 PROCEDURE LOOKUP(ID: IDENTIFIER; VAR ATTR: FILEATTR; VAR FOUND: BOOLEAN);
137
138 PROCEDURE IOTRANSFER
139 (DEVICE: IODEVICE; VAR PARAM: IOPARAM; VAR BLOCK: UNIV PAGE);
140
141 PROCEDURE IOMOVE(DEVICE: IODEVICE; VAR PARAM: IOPARAM);
142
143 FUNCTION TASK: TASKKIND;
144
145 PROCEDURE RUN(ID: IDENTIFIER; VAR PARAM: ARGLIST;
146    VAR LINE: INTEGER; VAR RESULT: PROGRESULT);
147
148
149 PROGRAM MEMMEM(VAR PARAM: ARGLIST);
150
151 "SIMPLE TYPES AND CONSTANTS"
152
153 CONST LINE_LENGTH = 70;       "OUTPUT DEV LINE LENGTH"
154
155 TYPE CHAR8 = ARRAY [1..8] OF CHAR;
156
157 CONST
158 MIN_KIND = 0;     MAX_KIND = 0;
159 MIN_TYPE = 0;     MAX_TYPE = 0;
160 MIN_MODE = 0;     MAX_MODE = 0;
161 MIN_CONTEXT = 0;    MAX_CONTEXT = 0;
162 SFIELD = 6; LFIELD = 11; "SHORT & LONG FIELD WIDTHS"
163
164
165 "VARIABLES"
166
167 VAR WORDS_IN: INTEGER;     "WORD IN IFL BUFFER"
168 IN_FILE: INTEGER;          "INPUT FILE"
169 PAGES_IN: INTEGER;         "CURRENT PAGE NUMBER"
170 PAGE_IN: PAGE;             "INPUT BUFFER"
171 PASS_NO: INTEGER;          "COMPILER PASS NUMBER"
172 OUT_COL_PTR: INTEGER;      "POSN ON OUTPUT LINE FOR BUF"
173 OUT_BUF_PTR: INTEGER;      "POSN IN BUFFER"
174 OUT_BUF: LINE;             "OPERATOR OUTPUT BUFFER"
175 FIRST_COL: INTEGER;        "FIRST COL TO USE FOR NON-NEWLINE"
176 MIN_COL_SEP: INTEGER;      "MIN BLANKS BETWEEN OPS"
177 COL_SEP: INTEGER;          "SEP BETWEEN COL BOUNDARIES"
178 OP: INTEGER;               "CURRENT INPUT OPERATOR"
179 LINK: PASSPTR;
```

```
180
181
182      "INPUT FILE SUPPORT ROUTINES"
183
184      PROCEDURE READ_TFL (VAR ARG: INTEGER);
185      BEGIN
186        IF WORDS_IN = PAGELENGTH THEN BEGIN
187          GET(IN_FILE,PAGES_IN,PAGE_IN);
188          PAGES_IN:=PAGES_IN+1; WORDS_IN:=0;
189        END;
190        WORDS_IN:=WORDS_IN+1;
191        ARG:=PAGE_IN[WORDS_IN];
192      END;
193
194      "BUFFER OUTPUT ROUTINES"
195
196      PROCEDURE NEXT_COL;
197      VAR I, J: INTEGER;
198      BEGIN "PAD TO START OF NEXT OP COLUMN"
199        I := OUT_COL_PTR + MIN_COL_SEP;
200        IF I+OUT_BUF_PTR >= LINE_LENGTH  "WON'T FIT"
201        THEN BEGIN
202          WRITE(NL);
203          I := FIRST_COL;
204          OUT_COL_PTR := 1;
205        END;
206        FOR J:=OUT_COL_PTR TO I-1 DO WRITE(' ');
207        OUT_COL_PTR:=I;
208      END;
209
210      PROCEDURE WRITE_BUF (ALIGN: BOOLEAN);
211      VAR I: INTEGER;
212      BEGIN
213        IF ALIGN THEN NEXT_COL;
214        FOR I:=1 TO OUT_BUF_PTR-1 DO WRITE(OUT_BUF[I]);
215        OUT_COL_PTR:=OUT_COL_PTR+OUT_BUF_PTR-1;
216        OUT_BUF_PTR:=1;
217      END;
218
219      PROCEDURE WRITE_CHAR (C: CHAR);
220      BEGIN
221        IF OUT_BUF_PTR >= LINE_LENGTH THEN
222          WRITE_BUF(TRUE);
223        OUT_BUF[OUT_BUF_PTR]:=C;
224        OUT_BUF_PTR:=OUT_BUF_PTR+1;
225      END;
226
227      "BUFFER FORMATTING ROUTINES"
228
229      PROCEDURE WRITE_CHAR8 (TEXT: CHAR8);
230      VAR I: INTEGER;
231      BEGIN
232        FOR I := 1 TO 8 DO
233          IF (TEXT[I] <> ' ')
234            THEN WRITE_CHAR (TEXT[I])
235      END;
236
237      PROCEDURE WRITE_HEX_CHAR (VAL: INTEGER);
238      BEGIN
239        IF VAL > 9 THEN WRITE_CHAR(CHR(VAL-10+ORD('A')))
```

```
240    ELSE WRITE_CHAR(CHR(VAL+ORD('0')));
241  END;
242
243  PROCEDURE WRITE_HEX (VAL: INTEGER; N: INTEGER);
244  TYPE TAGS = (TAG1, TAG2);
245  VAR I: INTEGER;
246      X: RECORD CASE TAGS OF
247          TAG1: (INT: INTEGER);
248          TAG2: (BYTES: ARRAY [1..2] OF BYTE);
249      END;
250  BEGIN
251    X.INT:=VAL;
252    FOR I:=3-N TO 2 DO BEGIN
253      WRITE_HEX_CHAR(X.BYTES[I] DIV 16);
254      WRITE_HEX_CHAR(X.BYTES[I] MOD 16);
255    END;
256  END;
257
258  PROCEDURE WRITE_INT (VAL: INTEGER);
259  VAR A: ARRAY [1..10] OF CHAR;
260      I, J, REM: INTEGER;
261  BEGIN
262    REM:=VAL; I:=1;
263    REPEAT
264      A[I]:=CHR(ABS(REM MOD 10)+ORD('0'));
265      I:=I+1; REM:=REM DIV 10;
266    UNTIL REM = 0;
267    IF VAL<0 THEN WRITE_CHAR('-');
268    FOR J:=I-1 DOWNTO 1 DO WRITE_CHAR(A[J]);
269  END;
270
271  PROCEDURE WRITE_INTR (VAL: INTEGER; N: INTEGER);
272  "WRITE INTEGER, RIGHT-JUST IN FIELD OF WIDTH N"
273  VAR A: ARRAY [1..10] OF CHAR;
274      I, J, REM: INTEGER;
275  BEGIN
276    REM:=VAL; I:=1;
277    REPEAT
278      A[I]:=CHR(ABS(REM MOD 10)+ORD('0'));
279      I:=I+1; REM:=REM DIV 10;
280    UNTIL REM = 0;
281    IF VAL < 0 THEN BEGIN
282      A[I]:='-'; I:=I+1;
283    END;
284    FOR J:=I TO N DO WRITE_CHAR(' ');
285    FOR J:=I-1 DOWNTO 1 DO WRITE_CHAR(A[J]);
286  END;
287
288
289  PROCEDURE INDEXERROR (TEXT:CHAR8; BADINDEX: INTEGER);
290  VAR I: INTEGER;
291  BEGIN
292    WRITE_CHAR8 ('$$$$$$$$');
293    WRITE_CHAR8 ('INDEX   ');
294    WRITE_CHAR (' ');
295    WRITE_CHAR8 ('ERROR   ');
296    WRITE_CHAR8 ('$$$$$$$$');
297    WRITE_CHAR8 (TEXT);
298    WRITE_CHAR ('-');
299    WRITE_CHAR (' ');
```

```
300 WRITE_INT (BADINDEX);
301 WRITE_CHAR (' ');
302 WRITE_CHAR (' ');
303 END;
304
305
306 "GENERAL OPERATOR FORMATTING ROUTINE"
307
308 PROCEDURE WRITE_OP (OP_NAME: CHAR8; NUM_ARGS, KIND_ARG, TYPE_ARG,
309                     MODE_ARG, CONTEXT_ARG: INTEGER);
310 VAR ARG_NO, ARG_VAL, I: INTEGER;
311 BEGIN
312 WRITE_CHAR8 (OP_NAME);
313 FOR ARG_NO := 1 TO NUM_ARGS DO
314     BEGIN
315     IF ARG_NO = 1
316     THEN WRITE_CHAR ( '(' )
317     ELSE WRITE_CHAR ( ',' );
318     READ_IFL (ARG_VAL);
319     IF
320         ARG_NO = KIND_ARG THEN
321         BEGIN
322         IF (ARG_VAL<MIN_KIND) OR (ARG_VAL>MAX_KIND)
323         THEN INDEXERROR ('BAD-KIND', ARG_VAL)
324         ELSE WRITE_CHAR8 ('NO_KIND_')
325         END    ELSE IF
326     ARG_NO = TYPE_ARG THEN
327         BEGIN
328         IF (ARG_VAL<MIN_TYPE) OR (ARG_VAL>MAX_TYPE)
329         THEN INDEXERROR ('BAD-TYPE', ARG_VAL)
330         ELSE WRITE_CHAR8 ('NO_TYPE_')
331         END    ELSE IF
332     ARG_NO = MODE_ARG THEN
333         BEGIN
334         IF (ARG_VAL<MIN_MODE) OR (ARG_VAL>MAX_MODE)
335         THEN INDEXERROR ('BAD-MODE', ARG_VAL)
336         ELSE WRITE_CHAR8 ('NO_MODE_')
337         END    ELSE IF
338     ARG_NO = CONTEXT_ARG THEN
339         BEGIN
340         IF (ARG_VAL<MIN_CONTEXT) OR (ARG_VAL>MAX_CONTEXT)
341         THEN INDEXERROR ('BAD-CTXT', ARG_VAL)
342         ELSE WRITE_CHAR8 ('NO_CNTXT')
343         END    ELSE IF
344     (ARG_VAL < -32768) OR (ARG_VAL > 32767)
345     THEN WRITE_HEX (ARG_VAL, 2)
346     ELSE WRITE_INT (ARG_VAL)
347     END;                        "OF BEGIN"
348 IF NUM_ARGS>0 THEN WRITE_CHAR(')');
349 WRITE_BUF(TRUE);
350 END;
351
352 "SPECIAL CASE ROUTINE FOR NEW LINE OPERATOR"
353
354 PROCEDURE NEW_LINE;
355 VAR ARG: INTEGER;
356 BEGIN
357 IF PASS_NO < 6
358 THEN READ_IFL(ARG);    "LINE NUMBER"
359 IF OUT_COL_PTR<>1 THEN WRITE_CHAR(NL);
```

```
360  WRITE_CHAR(NL);   "DOUBLE SPACE BEFORE NEW LINE"
361  OUT_COL_PTR:=1;
362  IF PASS_NO < 6
363  THEN BEGIN
364      WRITE_CHAR8('LINE      ');
365      WRITE_INTR(ARG,5);
366      WRITE_BUF(FALSE)
367      END
368  ELSE WRITE_BUF(FALSE);
369  END;
370
371
372  PROCEDURE SKIP (LINES: INTEGER);
373  VAR L: INTEGER;
374  BEGIN
375  IF OUT_COL_PTR <> 1
376  THEN WRITE_CHAR (NL);
377  FOR I := 1 TO LINES DO WRITE_CHAR(NL);
378  OUT_COL_PTR := 1;
379  WRITE_BUF(FALSE)
380  END;
381
382
383  "SPECIAL CASE ROUTINE FOR LONG CONSTANT"
384
385  PROCEDURE WRITE_LCONST;
386  VAR LEN, I, N, ARG: INTEGER;
387  BEGIN
388
389  WRITE_CHAR8('LNGCONST'); WRITE_CHAR('(');
390  READ_IFL(LEN);
391  WRITE_INT(LEN); WRITE_CHAR(',');
392  WRITE_CHAR8 ('HEX-VAL:');
393  NEXT_COL; N:=(LEN-1) DIV 2 + 1;
394  FOR I:=1 TO N DO BEGIN
395      READ_IFL(ARG);
396      IF OUT_COL_PTR+OUT_BUF_PTR-1+9 > LINE_LENGTH THEN BEGIN
397          WRITE_BUF(FALSE); WRITE(NL); OUT_COL_PTR:=1; NEXT_COL;
398          END;
399      WRITE_HEX(ARG,2);
400      IF I < N THEN WRITE_CHAR(',') ELSE WRITE_CHAR(')');
401      END;
402  WRITE_BUF(FALSE)
403  END;
404
405  "COMMON INITIALIZATION"
406
407  PROCEDURE INITIALIZE;
408  BEGIN
409  WORDS_IN:=PAGELENGTH;
410  PAGES_IN:=1;
411  PASS_NO:=PARAM[10].INT;
412  OUT_COL_PTR:=1;
413  OUT_BUF_PTR:=1;
414  FIRST_COL:=12;
415  MIN_COL_SEP:=3;
416  COL_SEP:=3;
417  LINK := PARAM[2].PTR
418  END;
419  END;
```

```
420   "WRITE PASS LISTING HEADER"
421   PROCEDURE WRITE_TEXT (TEXT: LINE);
422   VAR I: INTEGER;
423   BEGIN
424       I:=1;
425       WHILE TEXT[I] <> '$' DO BEGIN
426           WRITE(TEXT[I]); I:=I+1;
427           END;
430       END;
431
432   PROCEDURE WRITE_HEADER;
433   VAR I, J: INTEGER;
434       C: CHAR;
435   BEGIN
436       C:=CHR(PASS_NO+ORD('0'));
437       WRITE (FF);
438       WRITE(NL);
439       FOR I:=1 TO LINE_LENGTH DO WRITE(C);
440       WRITE(NL);
441       WRITE(C);
442       J:=LINE_LENGTH DIV 2 - 17;
443       FOR I:=2 TO J-1 DO WRITE(' ');
444       WRITE_TEXT('HEPASCAL INTERMEDIATE CODE PASS $');
445       WRITE(C);
446       FOR I:=J+33 TO LINE_LENGTH-1 DO WRITE(' ');
447       WRITE(C); WRITE(NL);
448       FOR I:=1 TO LINE_LENGTH DO WRITE(C);
449       WRITE(NL);
450   END;
451
452   PROCEDURE WRITE_5_CASE (OP_NAME: CHAR8);
453   VAR
454       I: INTEGER;
455       MIN,
456       MAX,
457       MAX_MINUS_MIN,
458       LOCATION,
459       STMT_LABEL:     INTEGER;
460
461   BEGIN
462       WRITE_CHAR8(OP_NAME);
463       WRITE_CHAR( '(' );
464       READ_IFL(MIN);
465       WRITE_INT(MIN);
466       WRITE_CHAR( ',' );
467       READ_IFL (MAX);
468       WRITE_INT (MAX);
469       MAX_MINUS_MIN := MAX - MIN;
470
471       FOR I := 0 TO MAX_MINUS_MIN DO
472           BEGIN
473               WRITE_CHAR(',');
474               READ_IFL(STMT_LABEL);
475               WRITE_INT(STMT_LABEL);
476           END;
477       WRITE_CHAR( ')' );  WRITE_BUF(TRUE)
478   END;
479
```

```
480 PROCEDURE DUMP_TABLE (TABLE: TABLEPTR; ENTRIES:INTEGER);
481 CONST FLDSIZE = 6;
482 VAR  I, J, K: INTEGER; PORTION: TABLEPTR;
483      CONSTAB: BOOLEAN;
485 BEGIN
486 CONSTAB := TABLE = LINK@.TABLES@.CONSTTABLE;
487 PORTION := TABLE;
488 I := ENTRIES - MAXWORD;
489 WITH PORTION@ DO
490   BEGIN
491   WHILE I>0 DO
492     BEGIN
493     IF CONSTAB
494     THEN FOR J := 1 TO MAXWORD DO
495       BEGIN
496       K := CONTENTS[J];
497       WRITE_INTR(K DIV 256, FLDSIZE);
498       WRITE_BUF(TRUE);
499       WRITE_INTR(K MOD 256, FLDSIZE);
500       WRITE_BUF(TRUE);
501       END
502     ELSE FOR J := 1 TO MAXWORD DO
503       BEGIN
504       WRITE_INTR(CONTENTS[J], FLDSIZE);
505       WRITE_BUF(TRUE);
506       END;
507     I := I - MAXWORD;
508     PORTION := NEXTPORTION
509     END;
510   I := MAXWORD + I; "I IS NEGATIVE AT THIS POINT"
511   IF CONSTAB
512   THEN FOR J := 1 TO I DO
513     BEGIN
514     K := CONTENTS[J];
515     WRITE_INTR(K DIV 256, FLDSIZE);
516     WRITE_BUF(TRUE);
517     WRITE_INTR(K MOD 256, FLDSIZE);
518     WRITE_BUF(TRUE);
519     END
520   ELSE FOR J := 1 TO I DO
521     BEGIN
522     WRITE_INTR(CONTENTS[J], FLDSIZE);
523     WRITE_BUF(TRUE);
524     END;
525   END;
526 END;

529 PROCEDURE PASS5;
530 CONST MIN_OP5 = 0;      MAX_OP5 = 49;
531 CONST
532   PASS5_TABLE = (
533   'PUSHCNST', 1,        'PUSHVAR ', 3,      " 1"
534   'PUSHIND ', 1,        'PUSHADDR', 2,
535   'FIELD   ', 1,        'INDEX   ', 4,
536   'POINTER ', 0,        'VARIANT ', 2,
537   'RANGE   ', 2,        'ASSIGN  ', 1,
538   'ASSNTAG ', 0,        'COPY    ', 1,      " 9"
539   'NEW     ', 0,        'NOT     ', 0,
```

A-33

```
540         'AND     ', 1,       'OR      ', 1,
541         'NEG     ', 1,       'AND     ', 1,
542         'SUB     ', 1,       'MUL     ', 1,
543         'DIV     ', 1,       'MOD     ', 1,
544         'INVALID ', 0,       'INVALID ', 0,      "19"
545         'FUNCTION', 2,       'BUILDSET', 2,
546         'COMPARE ', 2,       'CMPSTRUC', 2,
547         'FUNCVALU', 2,       'DEFLABEL', 1,      "29"
548         'JUMP    ', 2,       'FALSJUMP', 1,
549         'CASEJUMP', 0,       'INITVAR ', 0,
550         'CALL    ', 3,       'ENTER   ', 5,
551         'RETURN  ', 1,       'POP     ', 1,
552         'NEWLINE ', 1,       'ERROR   ', 0,      "39"
553         'LONCONST', 0,       'MESSAGE ', 2,
554         'INCRMENT', 0,       'DECRMENT', 0,
555         'PROCEDURE', 1,      'INIT    ', 4,
556         'PUSHLABL', 1,       'CALLPROC', 0,      "47"
557         'EOM     ', 1,       'DUP1TOS ', 0);     "49"
558
559         ARRAY[MIN_OP5 .. MAX_OP5] OF
560         RECORD
561             OP_NAME: CHAR8;
562             NUM_ARGS: BYTE
563         END;
564
565     CONST
566         EOM5 = 48;
567         NEWLINE5 = 38;
568         CASEJUMP5 = 32;
569         LONGCONSTANT5 = 40;
570
571     BEGIN
572         IN_FILE := 2;
573         REPEAT
574             READ_IFL (OP);
575             WITH PASS5_TABLE[OP] DO
576             IF (OP < MIN_OP5) OR (OP > MAX_OP5)
577             THEN INDEXERROR ('BAD_OP ', OP)
578             ELSE IF OP = NEWLINE5
579                 THEN NEW_LINE
580                 ELSE IF OP = CASEJUMP5
581                     THEN WRITE_5_CASE (OP_NAME)
582                     ELSE IF OP = LONGCONSTANT5
583                         THEN WRITE_LCONST
584                         ELSE WRITE_OP (OP_NAME, NUM_ARGS, 0, 0, 0, 0)
585         UNTIL OP = EOM5;
586         WRITE (NL)
587     END;
588
589
590
591
592     PROCEDURE PASS6;
593     CONST   MIN_OP6 = 0;     MAX_OP6 = 255;
594     CONST   PASS56_TABLE = (
595
596         'SLDC00 ', 0,        'SLDC01 ', 0,      "001"
597         'SLDC02 ', 0,        'SLDC03 ', 0,
598         'SLDC04 ', 0,        'SLDC05 ', 0,      "005"
599         'SLDC06 ', 0,        'SLDC07 ', 0,
```

"015"   "025"   "035"   "045"   "055"   "065"   "075"   "085"   "095"   "105"   "115"   "125"

Addr	Label		Label	
600	'SLDC08'	0.	'SLDC09'	0.
601	'SLDC10'	0.	'SLDC11'	0.
602	'SLDC12'	0.	'SLDC13'	0.
603	'SLDC14'	0.	'SLDC15'	0.
604	'SLDC16'	0.	'SLDC17'	0.
605	'SLDC18'	0.	'SLDC19'	0.
606	'SLDC20'	0.	'SLDC21'	0.
607	'SLDC22'	0.	'SLDC23'	0.
608	'SLDC24'	0.	'SLDC25'	0.
609	'SLDC26'	0.	'SLDC27'	0.
610	'SLDC28'	0.	'SLDC29'	0.
611	'SLDC30'	0.	'SLDC31'	0.
612	'SLDL01'	0.	'SLDL02'	0.
613	'SLDL03'	0.	'SLDL04'	0.
614	'SLDL05'	0.	'SLDL06'	0.
615	'SLDL07'	0.	'SLDL08'	0.
616	'SLDL09'	0.	'SLDL10'	0.
617	'SLDL11'	0.	'SLDL12'	0.
618	'SLDL13'	0.	'SLDL14'	0.
619	'SLDL15'	0.	'SLDL16'	0.
620	'SLD001'	0.	'SLD002'	0.
621	'SLD003'	0.	'SLD004'	0.
622	'SLD005'	0.	'SLD006'	0.
623	'SLD007'	0.	'SLD008'	0.
624	'SLD009'	0.	'SLD010'	0.
625	'SLD011'	0.	'SLD012'	0.
626	'SLD013'	0.	'SLD014'	0.
627	'SLD015'	0.	'SLD016'	0.
628	'INVAL064'	0.	'INVAL065'	0.
629	'INVAL066'	0.	'INVAL067'	0.
630	'INVAL068'	0.	'INVAL069'	0.
631	'INVAL070'	0.	'INVAL071'	0.
632	'INVAL072'	0.	'INVAL073'	0.
633	'INVAL074'	0.	'INVAL075'	0.
634	'INVAL076'	0.	'INVAL077'	0.
635	'INVAL078'	0.	'INVAL079'	0.
636	'INVAL080'	0.	'INVAL081'	0.
637	'INVAL082'	0.	'INVAL083'	0.
638	'INVAL084'	0.	'INVAL085'	0.
639	'INVAL086'	0.	'INVAL087'	0.
640	'INVAL088'	0.	'INVAL089'	0.
641	'INVAL090'	0.	'INVAL091'	0.
642	'INVAL092'	0.	'INVAL093'	0.
643	'INVAL094'	0.	'INVAL095'	0.
644	'INVAL096'	0.	'INVAL097'	0.
645	'INVAL098'	0.	'INVAL099'	0.
646	'INVAL100'	0.	'INVAL101'	0.
647	'INVAL102'	0.	'INVAL103'	0.
648	'INVAL104'	0.	'INVAL105'	0.
649	'INVAL106'	0.	'INVAL107'	0.
650	'INVAL108'	0.	'INVAL109'	0.
651	'INVAL110'	0.	'INVAL111'	0.
652	'INVAL112'	0.	'INVAL113'	0.
653	'INVAL114'	0.	'INVAL115'	0.
654	'INVAL116'	0.	'INVAL117'	0.
655	'INVAL118'	0.	'INVAL119'	0.
656	'SIND0'	0.	'SIND1'	0.
657	'SIND2'	0.	'SIND3'	0.
658	'SIND4'	0.	'SIND5'	0.
659	'SIND6'	0.	'SIND7'	0.

```
"135"   "145"   "155"   "165"   "175"   "185"   "195"   "205"   "215"   "225"   "235"   "245"
```

```
660 'LDCR        'LDCI'
661 'LCA         'LDC'
662 'LLA         'LDO'
663 'LAO         'LDL'
664 'LDA         'LOD'
665 'UJP         'UJP.'
666 'MPI         'NVI'
667 'STM         'MODI'
668 'CPL         'CPG'
669 'CI4         'CXL'
670 'CXG         'CXI'
671 'RFU         'CPF'
672 'LDCN        'LSL'
673 'LDE         'LAE'
674 'NOP         'LPR'
675 'RPT         'RBP'
676 'LOR         'LAND'
677 'ADI         'SBI'
678 'STL         'SRO'
679 'STR         'LDB'
680 'INVAL.168'  'INVAL.169'
681 'INVAL.170'  'INVAL.171'
682 'INVAL.172'  'INVAL.173'
683 'INVAL.174'  'INVAL.175'
684 'EQUI        'NEQI'
685 'LEQI        'GEQI'
686 'LEUSW       'GEUSW'
687 'EQUPWR      'LEQPWR'
688 'GEQPWR      'EQUBYT'
689 'LRQBYT      'GEQBYT'
690 'SRS         'SWAP'
691 'TNC         'RND'
692 'ADR         'SBR'
693 'MPR         'DVR'
694 'STO         'MOV'
695 'DUP2        'ADJ'
696 'STB         'LDP'
697 'STP         'CHK'
698 'FLT         'EQUREAL'
699 'LEQREAL.    'GEQREAL'
700 'LDM         'SPR'
701 'EFJ         'NFJ'
702 'FJP         'FJPL'
703 'XJP         'IXA'
704 'IXP         'STE'
705 'TNN         'UNI'
706 'INT         'DIF'
707 'SIGNAL      'WAIT'
708 'ABI         'NGI'
709 'DUP1        'ABR'
710 'NGR         'LNOT'
711 'IND         'INC'
712 'MESAGE'     'EOM'
713 'NEWLIN'     'INVAL.235'
714 'INVAL.236'  'INVAL.237'
715 'INVAL.238'  'INVAL.239'
716 'INVAL.240'  'INVAL.241'
717 'INVAL.242'  'INVAL.243'
718 'INVAL.244'  'INVAL.245'
719 'INVAL.246'  'INVAL.247'
```

```
720  'INVAL248',  0,    'INVAL249',  0,
721  'INVAL250',  0,    'INVAL251',  0,
722  'INVAL252',  0,    'INVAL253',  0,
723  'INVAL254',  0,    'INVAL255',  0):        "255"
724
725           ARRAY [MIN_OP6..MAX_OP6] OF
726           RECORD
727              OP_NAME: CHAR8;
728              NUM_ARGS: BYTE;
729              END;
730
731
732  CONST EOM6 = 233;      UJPL = 139;      FJPL = 213;
733        NEWLINE6 = 234;
734        INTFLD = 6;
735
736  VAR  SEGS, SEGNO, SEGRTNS, ARG, I: INTEGER;
737  BEGIN
738  "SCAN CONCURRENT SEGMENT"
739     IN_FILE:=1;
740     SKIP(1); WRITE_TEXT('** CONCURRENT SEGMENT **$'); SKIP(1);
741     REPEAT
742        READ_IFL(OP);
743        WITH PASS6_TABLE[OP] DO
744        IF (OP < MIN_OP6) OR (OP > MAX_OP6)
745        THEN INDEXERROR ('BAD-OP ', OP)
746        ELSE IF OP = NEWLINE6
747           THEN NEW_LINE
748           ELSE BEGIN
749              WRITE_OP (OP_NAME, NUM_ARGS, 0, 0, 0, 0);
750              IF (OP=UJPL) OR (OP=FJPL)
751              THEN BEGIN
752                 READ_IFL(ARG);
753                 WRITE_CHAR8('#LOC= ');
754                 WRITE_INT(ARG); WRITE_CHAR('*');
755                 WRITE_BUF(TRUE)
756                 END
757              END
758     UNTIL OP = EOM6;
759
760
761
762  "SCAN INTERFACE SEGMENT"
763     IN_FILE := 4;  WORDS_IN := PAGELENGTH;   PAGES_IN := 1;
764     SKIP(2);
765     WRITE_TEXT ('** INTERFACE SEGMENT(S) **$');
766     SEGS := LINK6.INTERFACE6.INTERFACES;
767     SKIP(0);
768     IF SEGS = 0
769     THEN WRITE_CHAR8('....NONE')
770     ELSE BEGIN WRITE_CHAR8('SEGMNTS='); WRITE_INT(SEGS) END;
771     SKIP(1);
772     SEGNO := 1;
773     WHILE SEGNO <= SEGS DO
774        BEGIN
775        SEGRTNS := LINK6.INTERFACE6.INTERFACESIZES[SEGNO];
776        WRITE_TEXT('SEGMENT NO. $'); WRITE_INT(SEGNO);
777        WRITE_BUF(TRUE); NEXT_COL;
778        WRITE_TEXT('NO. OF ROUTINES: $'); WRITE_INT(SEGRTNS); SKIP(1);
779        READ_IFL(ARG); WRITE_CHAR8('SEGLEN= ');
```

```
780  WRITE_INT(ARG); WRITE_BUF(TRUE);                                              258  736  210  222
781  READ_IFL(ARG); WRITE_CHAR8('MININDX=');                                       184  736  229  222
782  WRITE_INT(ARG); WRITE_BUF(TRUE);                                              258  736  210  222
783  READ_IFL(ARG); WRITE_CHAR8('MAXINDX=');                                       184  736  229  222
784  WRITE_INT(ARG); WRITE_BUF(TRUE);                                              258  736  210
785  SKIP(0);                                                                      372
786
787  WRITE_TEXT('CASEJUMP OFFSETS:$'); WRITE_BUF(TRUE);                            423  210  222
788  FOR J := 1 TO SEGRTNS DO                                                      736  736
789     BEGIN
790     READ_IFL(ARG);                                                            184  736  210
791     WRITE_INT(ARG);  WRITE_BUF(TRUE)                                          258  736  222
792     END;
793  SKIP(0);                                                                      372
794
795  READ_IFL(ARG); WRITE_CHAR8('EXIT-IC=');                                       184  736  229  222
796  WRITE_INT(ARG); WRITE_BUF(TRUE); NEXT_COL;                                    258  736  210  222
797  READ_IFL(ARG); WRITE_CHAR8('DATASIZ=');                                       184  736  229
798  WRITE_INT(ARG);                                                               258  736
799  SKIP(0);                                                                      372
800
801  FOR I := 1 TO (4 + 2*SEGRTNS + 2) DO                                          736  736
802     BEGIN
803     READ_IFL(OP);                                                             184  178
804     WITH PASS6_TABLE[OP] DO                                                    594  178
805       IF (OP < MIN_OP6) OR (OP > MAX_OP6)                                      178  593  593
806       THEN INDEXERROR ('BAD-OP ', OP)                                         289  178
807       ELSE WRITE_OP (OP_NAME, NUM_ARGS, 0, 0, 0, 0)                           308  727  728
808     END;
809  SKIP(0);                                                                      372

810
811  READ_IFL(ARG); WRITE_CHAR8('RTN-PTR=');                                       184  736  229
812  WRITE_INT(ARG); WRITE_BUF(TRUE);                                              258  736  210  222
813  READ_IFL(ARG); WRITE_CHAR8('LO-ORDR=');                                       184  736  229
814  WRITE_INT(ARG MOD 256); WRITE_BUF(TRUE); NEXT_COL;                            258  736  210  222
815  WRITE_CHAR8('HI-ORDR='); WRITE_INT(ARG DIV 256);                             229  258  736
816
817  SKIP(2);                                                                      372
818  SEGNO := SUCC(SEGNO)                                                          736       736
819  END;
820
821
822  WITH LINK@ DO                                                                 179
823     BEGIN
824  "PRINT SOME OF THE PASSLINK FIELDS"
825  SKIP(1); WRITE_TEXT('** PASSLINK FIELDS **$'); SKIP(0);                       372  423  372
826  WRITE_CHAR8('LABELS  '); WRITE_INTR(LABELS, SFIELD);                          229  271  87   162
827  SKIP(0);                                                                      372
828  WRITE_CHAR8('BLOCKS  '); WRITE_INTR(BLOCKS, SFIELD);                          229  271  87   162
829  SKIP(0);                                                                      372
830  WRITE_TEXT('CONSTANTS$'); WRITE_INTR(CONSTANTS, SFIELD);                      423  271  87   162
831  WRITE_CHAR(' '); WRITE_CHAR('WORDS   '); SKIP(0);                             219  229  372
832  WRITE_TEXT(XJP_OFFSETS$'); WRITE_INTR(XJP_OFFSETS, SFIELD);                   423  271  87   162
833  WRITE_CHAR(' '); WRITE_CHAR8('WORDS   '); SKIP(0);                            219  229  372
834  WRITE_TEXT('SEGDISTANCE$');WRITE_INTR(TABLES@.SEGDISTANCE, SFIELD);           423  271  89   67
835  WRITE_CHAR(' '); WRITE_CHAR8('BYTES   ');                                     219  229
836
837  "PRINT CONTENTS OF THE TABLES PASSED THRU HEAP"
838  SKIP(2); WRITE_TEXT('** TABLES **$');                                         372  423
839
```

```
840
841     SKIP(1); WRITE_CHAR8('CONSTNTS'); SKIP(0);               372  229  372   87
842     IF CONSTANTS > 0                                           87   89   68
843     THEN DUMP_TABLE(TABLES@.CONSTTABLE, CONSTANTS)            481
844     ELSE WRITE_CHAR8('....NONE');                            229
845
846     SKIP(2); WRITE_CHAR8('XJPTABLE'); SKIP(0);               372  229  372   87
847     IF XJP_OFFSETS > 0                                         87   89   68
848     THEN DUMP_TABLE (TABLES@.XJPTABLE, XJP_OFFSETS)          481
849     ELSE WRITE_CHAR8('....NONE');                            229
850
851     SKIP(2);  WRITE_CHAR8('JUMPTABL'); SKIP(0);              372  229  372   87
852     IF LABELS > 0                                             .87   89   68
853     THEN DUMP_TABLE (TABLES@.JUMPTABLE, LABELS)              481
854     ELSE WRITE_CHAR8('....NONE');                            229
855
856     SKIP(2); WRITE_CHAR8('EXIT-IC '), SKIP(0);               372  229  372   87
857     IF BLOCKS > 0                                              87   89   69
858     THEN DUMP_TABLE (TABLES@.EXITICTABLE, BLOCKS)            481
859     ELSE WRITE_CHAR8('....NONE');                            229
860
861     SKIP(2); WRITE_CHAR8('DATASIZE'); SKIP(0);               372  229  372   87
862     IF BLOCKS > 0                                              87   89   69
863     THEN DUMP_TABLE (TABLES@.DATASIZETABLE, BLOCKS)          481
864     ELSE WRITE_CHAR8('....NONE');                            229
865     END;
866
867     FOR I := 1 TO PAGELENGTH + 10 DO WRITE_CHAR(' ')         736   15  219
868     END;
869
870
871
872
873
874     PROCEDURE PASS7;
875     CONST
876     MIN_OP7 = 1;      MAX_OP7 = 1;
877
878     PASS7_TABLE = ('NULLOP  ', 0):
879                    ARRAY [MIN OP7..MAX_OP7] OF              876  876  876
880                       RECORD
881                          OP_NAME: CHAR8;                         155
882                          NUM_ARGS: BYTE;                         248
883                       END;
884
885     VAR
886     I: INTEGER;                                                     5
887     BEGIN
888
889     BEGIN
890        IN_FILE := 3;                                        168
891        COL_SEP := 1;                                        177
892        FOR I := 1 TO PAGELENGTH +10 DO WRITE_CHAR(' ');     887   15  219
893     END;
894
895
896
897     BEGIN
898        INITIALIZE;                                          408
899        WRITE_HEADER;                                        432
```

```
900    CASE PASS_NO OF
901      1:"PASS1";
902      2:"PASS2";
903      3:"PASS3";
904      4:"PASS4";
905      5: PASS5;
906      6: PASS6;
907      7: PASS7;
908      8, 9: "NOT IMPLEMENTED";
909    END.
910 END.
```

171

529
592
874

CROSS REFERENCE   * IS DEF   = IS ASG

-A-

A

ABS	259*	264=	268	273*	278=	282=	285										
ACCEPT	264*	278															
ADDR	126*																
ALIGN	31*																
ARG	210*	213															
50*	133*	134*	184*	191=	355*	358	365	387*	395	399	736*	753	755	779	780	781	782
783	784	790	791	795	796	797	798	811	812	813	814	815				
ARGLIST | 105* | 105* | 145 | 149 | | | | | | | | | | | | | |
ARGSEQ | 107* | 133 | 133 | | | | | | | | | | | | | | |
ARGTAG | 93* | 97 | | | | | | | | | | | | | | | |
ARGTYPE | 96* | 105 | | | | | | | | | | | | | | | |
ARG_NO | 310* | 310= | 313= | 315 | 320 | 326 | 332 | 338 | 340 | 341 | 344 | 344 | 344 | | | | |
ARG_VAL | 310* | 318 | 322 | 323 | 328 | 328 | 329 | 334 | 335 | 340 | | | | | | | |
ASCII | 27 | | | | | | | | | | | | | | | | |
ATTR | 136* | | | | | | | | | | | | | | | | |

-B-

BACKSPACE | 41 | | | | | | | |
BADINDEX | 209* | 300 | | | | | | |
BLOCK | 118* | 119* | 129* | 130* | 139* | | | |
BLOCKS | 87* | 829 | 857 | 858 | 862 | 863 | | |
BOOL | 98 | | | | | | | |
BOOLEAN | 32 | 98 | 116 | 129 | 130 | 136 | 210 | |
BOOLTYPE | 94 | 98 | | | | | | |
BYTE | 248* | 562 | 728 | 883 | | | | |
BYTES | 248* | 253 | 254 | | | | | |
BYTES_PER_IN | 17* | | | | | | | |

-C-

C | 113* | 114* | 126* | 127* | 219* | 223 | 234= | 436= | 439 | 441 | 445 | 447 | 448 |
CALLERROR | 111 | | | | | | | | | | | | |
CARDDEVICE | 37 | | | | | | | | | | | | |
CASEJUMP5 | 568* | 580 | | | | | | | | | | | |
CHAR | 20 | 23 | 113 | 114 | 126 | 127 | 155 | 219 | 259 | 273 | | | |
CHAR8 | 155* | 229 | 289 | 308 | 452 | 561 | 727 | 882 | | | | | |
CHR | 239 | 240 | 264 | 278 | 436 | | | | | | | | |
CLOSE | 117* | | | | | | | | | | | | |
CODELIMIT | 111 | | | | | | | | | | | | |
COL_SEP | 177* | 417= | 891= | | | | | | | | | | |
COMPLETE | 44 | | | | | | | | | | | | |
CONCODE | 27 | | | | | | | | | | | | |
CONSTAB | 484* | 486= | 493 | 511 | | | | | | | | | |
CONSTANTS | 87* | 831 | 842 | 843 | | | | | | | | | |
CONSTTABLE | 68* | 486 | 843 | | | | | | | | | | |
CONTENTS | 63* | 496 | 504 | 514 | 522 | | | | | | | | |
CONTEXT_ARG | 309* | 338 | | | | | | | | | | | |
CONTROL | 39 | | | | | | | | | | | | |
CR | 13* | | | | | | | | | | | | |

-D-

DATASIZETABL | 69* | 863 | | | |
DEVICE | 139* | 141* | | | |
DISKDEVICE | 37 | | | | |
DISPLAY | 127* | | | | |
DUMP_TABLE | 481* | 843 | 848 | 853 | 858 | 863 |

-E-

EM	13*																	
EMPTY	27																	
ENDFILE	45																	
ENDMEDIUM	45																	
ENTRIES	481*	488																
EOF	129*	130*																
EOM5	566*	585																
EOM6	732*	759																
EXITICTABLE	69*	858																

-F-

F	116*	117*	118*	119*	120*		
FAILURE	44	368	379	397	402		
FALSE	366	437					
FF	13*	116	117	118	119	120	
FILE	25*	116	117	118			
FILEATTR	29*	136					
FILEKIND	27*	30					
FIRST_COL	175*	203	415=				
FJPL	78	751					
FLDSIZE	732*	497	499	504	515	517	522
FOUND	482*	136*					
FULLWORD	116*	88					
	10*	88					

-G-

GET	118*	167

-H-

HEADER	125*
HEAPLIMIT	111

-I-

I 197* 199* 200 203= 206 207 211* 214= 214 230* 232= 233 234 245* 252* 253 254 260*
262= 264 265= 265 268 274* 276= 278 279= 279 282 282= 282 284 285 290* 310* 373*
377= 387* 394* 400 424* 426= 427 428 428 428= 433* 439= 443= 446= 448= 454= 471= 483*
486= 491 507= 507 510= 510 512 520 736* 788= 801= 867= 887* 892=

ID	100	116*	136*	145=	145			
IDENTIFIER	23*	116						
IDENTIFY	125*							
IDLENGTH	22*	23	100*					
IDTYPE	94	100*						
IFINFO	78	79*						
IFPTR	78*	90						
INDEXERROR	289*	323	329	335	341	577	746	806
INITIALIZE	408*	898						
INP	107							
INPUT	39							
INPUTASK	53							
INT	99	247	251*	412				

INTEGER 10* 11* 16 31 58 63 67 80 81 87 99 118 119 120 122 123 146
167 168 169 171 172 173 175 176 177 178 184 197 211 230 237 243 243 245
247 258 260 271 271 274 289 290 309 310 355 372 373 387 424 433 454 459
481 483 736 807
766 775

INTERFACE	90*	766
INTERFACES	80*	766
INTERFACESIZ	81*	775
INTERVENTION	44	
INTFLD	734*	
INTTYPE	94	99*

```
IN_FILE        168*  187   572*  740=  763=  890=
IOARG           41*   50          141
IODEVICE        36*  139
IOMOVE         141*        141
IOOPERATION     39*   48
IOPARAM         47*  139   139
IORESULT        43*   49
IOTRANSFER     138*

-J-
J              197*  206=  260*  268=  274*  284=  285=  285
                                                   433*  442=  443   446
                                                   483*  494=  496   502=  504
               512=  514   520=  522   740=  763=  890=

JOBTASK         53
JUMPTABLE       68*  853

-K-
K              483*  496=  497   499   514=  515   517
KIND            30*  320
KIND_ARG       308*

-L-
LABELS          87*  827   852   853
LEN            387*  390   391   393
LENGTH         120*
LFIELD         162*
LINE            20*  125   131   132   146*  174   423
LINELENGTH      19*   20
LINES          372*  377
LINE_LENGTH    153*  200   221   396   439   442=  446   448
LINK           179*  418*  486   766   775   822
LOCATION       458*
LONGCONSTANT   569*  582
LOOKUP         136*

-M-
MARK           122*
MAX            456*  467   468   469
MAXARG         104*  105
MAXINTFAC       76*   81
MAXWORD         55*   63
MAX_CONTEXT    161*  340
MAX_KIND       158*  322
MAX_MINUS_MI   457*  469*  471
MAX_MODE       160*  334
MAX_OP5        530*  559   576
MAX_OP6        593*  725   745
MAX_OP7        876*  880   805
MAX_TYPE       159*  328
MEMMEM         149*
MIN            455*  464   465   469
MIN_COL_SEP    176*  199   416=
MIN_CONTEXT    161*  340
MIN_KIND       158*  322
MIN_MODE       160*  334
MIN_OP5        530*  559   576
MIN_OP6        593*  725   745
MIN_OP7        876*  880   805
MIN_TYPE       159*  328
MODE_ARG       309*  332
MOVE            39
```

-N-

N	243*	252	271*	284	387*	393=	394	400									
NEWLINE5	567*	578															
NEWLINE6	733*	747															
NEW_LINE	354*	579	579														
NEXTPORTION	62*	508															
NEXT_COL.	196*	213	397	397	777	796	814										
NIL.TYPE	94	98															
NL	13*	202	359	360	376	377	449	586									
NOTUSED	33*																
NUM_ARGS	308*	313	348	562*	728*	750	807	883*									

-O-

OP	178*	574	575	576	576	577	578	580	582	543	745	744	745	746	747	751	751
	759	803	804	805	805	806											
OPEN	116*																
OPERATION	48*																
OPTION	74*	86															
OPTIONS	86*																
OP_NAME	308*	312	452*	462	561*	581	584	727*	750	807	802*						
ORD	239	280	264	278	436												
OUT	107																
OUTPUT	39																
OUTPUTTASK	53																
OUT_BUF	174*	214	223=	215	216	221	224	396									
OUT_BUF_PTR	173*	200	214	215	207=	215	223	359	361=	375	378=	396	413=				
OUT_COL_PTR	172*	199	204=	206	207=	215	397=	413=									
OVERFLOW	110																

-P-

P	118*	119*	119	129	130	139	170		
PAGE	16*	118	16	186	410	763	867	892	
PAGELENGTH	15*	16							
PAGES_IN	169*	187	188=	188	411=	763=			
PAGE_IN	170*	187	191	145*	412	418			
PARAM	139*	141*	145*	149*					
PASS5	529*	905							
PASS5_TABLE	532*	575							
PASS6	592*	906							
PASS6_TABLE	594*	744							
PASS7	874*	804							
PASS7_TABLE	878*	907							
PASSLINK	84*	85*							
PASSPTR	84*	101	179						
PASS_NO	171*	357	362	412=	436	900			
POINTER	50*								
POINTERERROR	110								
PORTION	483*	487*	489	508=					
PRINTDEVICE	37								
PROGRESULT	109*	146							
PROTECTED	32*								
PTR	101	418							
PTRTYPE	94	101*							
PUT	119*	101*							

-R-

RANGEERROR	110
READ	113*
READARG	133*

-W-
```
WORDS_IN      167*
WRITE         114*  186   188=  190   190=  191   191   410=  763=  439   440   441   443   445   446   447   447   448   449
              202   206   214   397
              586
WRITEARG      134*
WRITEOF        41
WRITELINE     132*
WRITEPAGE     130*
WRITE_5_CASE  452*  581
WRITE_BUF     210*  222   349   366   368   379   397   402   477   498   500   505   516   518   523   756   777   780
              782   784   787   791   796   812   814
WRITE_CHAR    219*  234   239   240   267   268   280   284   285   294   298   299   301   302   316   317   348   359   360
              376   377   389   391   400   400   463   466   473   477   755   832   834   836   867   892
WRITE_CHAR8   229*  292   293   295   296   297   312   324   330   336   342   364   389   392   462   754   769   770
              779   781   783   795   797   811   813   815   827   829   832   834   836   841   844   846   849   851
              854   856   859   861   864
WRITE_HEADER  432*  899
WRITE_HEX     243*  345
WRITE_HEX_CH  237*  253   254
WRITE_INT     258*  300   386   391   465   468   475   755   770   776   778   780   782   784   791   796   798   812
              814   815
WRITE_INTR    271*  365   497   499   504   515   517   522   827   829   831   833   835
WRITE_LCONST  386*  583   750   807   765
WRITE_OP      308*  584   807
WRITE_TEXT    423*  444   741   765   776   787   826   831   833   835   839
```

-X-
```
X             246*  251   253   254
XJPTABLE       68*  848   848
XJP_OFFSETS    87*  833   847   848
```

END XREF 265 IDENTIFIERS 1174 TOTAL REFERENCES
250 COLLISIONS.

DISKETTE BLOCK DUMP PROGRAM

Some time after the arrival of the Microengine hardware it became apparent that some aspects of the machine were not clearly (if at all) described in the accompanying documents, and could be determined only by direct inspection of disk blocks. The installed system contains a disk-dumping utility (PATCH), but it is inconvenient to use. PATCH displays the contents of disk blocks as hexadecimal numbers or in a MIXED format where non-printing characters are shown in hex and the others are shown in their character representation. This was unsatisfactory since what was needed was a format which shows both the numerical and character representation of each byte. BLOCKDUMP does just that.

The program is written in UCSD Pascal and allows inspection of any block on the disk. The user is prompted for five items of information:

1) Physical disk drive unit number-- 4 and 5 are currently valid;

2) Starting disk block number-- the number of the block where the dump is to begin. The program assumes that a single-sided, double-density, eight-inch disk is mounted, so numbers between 0 and 987 are valid.

3) Number of blocks to be displayed-- must be such that block numbers greater than 987 will not be accessed. A nonpositive entry will terminate the program after the fifth piece of information has been entered.

4) Number base-- number base of the values to be displayed. The characters H (Hex) and D (Decimal) are valid.

5) Output device-- the device where the dump should be written. Valid values are the characters S (CRT Screen), P (Printer),

and Q (Quit). Entering 'Q' terminates the program.

On the printer, one block is printed per page. On the CRT, roughly one half of a block is displayed at a time. When the screen has been filled with information, typing a carriage return moves to the next screenful of information. After all the user information has been entered, there is no way to terminate the program before all the blocks entered as item 3 above have been displayed. This program was especially useful in discovering the format of Microengine code files.

The source code for BLOCKDUMP follows.

```
            (* LPRINTER: *)
     PROGRAM BLOCKDMP;
     (***************************************************************)
     (*                                                           *)
     (* PROGRAM TO DISPLAY (ON SCREEN OR PRINTER) THE CON-        *)
     (* TENTS OF DISK BLOCKS. USER TYPES IN THE UNIT NUMBER       *)
     (* OF THE DISK TO BE DUMPED, STARTING BLOCK NUMBER,          *)
     (* NUMBER OF BLOCKS TO BE DISPLAYED, AND WHERE OUTPUT        *)
     (* IS TO BE DIRECTED. CONTENTS OF EACH BLOCK ARE DIS-        *)
     (* PLAYED AS ASCII VALUES UNDER THE CHARACTER REP-           *)
     (* RESENTATIONS. BYTE ADDRESS (RELATIVE TO THE START OF      *)
     (* THE BLOCK) OF THE FIRST BYTE ON THE LINE APPEARS TO       *)
     (* THE LEFT OF THE ASCII VALUES.  .                          *)
     (*                                                           *)
     (***************************************************************)
     CONST
         CR = 13;
         LF = 10;
         BLLENGTH = 512; (* BYTES IN DISK BLOCK *)
         LINELEN  = 80; (* CHARS PER LINE *)
         MAXLINES =  22; (* LINES PER SCREEN PAGE *)
         MAXBLOCK = 987; (* MAX BLOCK # ON 8" DUAL DENS DISK *)

     TYPE
         LINE = PACKED ARRAY[1..LINELEN] OF CHAR;

     VAR
         QUIT, OK, PRTOPT, CONOPT: BOOLEAN;
         DISK, STARTBL, LENGTH, BLOCK,
             POSIT, CHORD, I, BYTE, LINES: INTEGER;
         OPT, CHARAC, CONTINUE, BASE: CHAR;
         OUTFLNAM: STRING[10];
         OUTFL: INTERACTIVE;
         BLOCKIN: PACKED ARRAY[1..BLLENGTH] OF CHAR;
         CHARLINE, NUMLINE: LINE;
         X: REAL;
```

```
PROCEDURE WRITENUM (VALUE: INTEGER; BASE: CHAR);
(* ROUTINE PUTS CHARACTER REPRESENTATION OF THE
    CONTENTS OF A BYTE INTO 'CHARLINE'.  ALWAYS
    PUTS IN THREE DIGITS.                       *)
CONST
    ORD0 = 48;
    ORDA = 65;

VAR
    TEMP: INTEGER;

BEGIN
IF BASE = 'H'
THEN BEGIN
    TEMP := VALUE DIV 256;
    IF TEMP = 0
    THEN NUMLINE[POSIT] := ' '
    ELSE IF TEMP <= 9
        THEN NUMLINE[POSIT] := CHR(TEMP + ORD0)
        ELSE NUMLINE[POSIT] := CHR(TEMP-10 + ORDA);
    POSIT := SUCC(POSIT);
    TEMP := (VALUE MOD 256) DIV 16;
    IF TEMP <= 9
    THEN NUMLINE[POSIT] := CHR(TEMP + ORD0)
    ELSE NUMLINE[POSIT] := CHR(TEMP-10 + ORDA);
    POSIT := SUCC(POSIT);
    TEMP := VALUE MOD 16;
    IF TEMP <= 9
    THEN NUMLINE[POSIT] := CHR(TEMP + ORD0)
    ELSE NUMLINE[POSIT] := CHR(TEMP-10 + ORDA);
    POSIT := SUCC(POSIT)
    END
ELSE BEGIN
    NUMLINE[POSIT] := CHR((VALUE DIV 100) + ORD0); (* 100'S *)
    POSIT := SUCC(POSIT);
    NUMLINE[POSIT] := CHR(((VALUE MOD 100) DIV 10) + ORD0); (* 10'S *)
    POSIT := SUCC(POSIT);
    NUMLINE[POSIT] := CHR((VALUE MOD 10) + ORD0); (* UNITS *)
    POSIT := SUCC(POSIT)
    END
END;
```

```
BEGIN
QUIT := FALSE;    DISK := 0;
PRTOPT := FALSE; CONOPT := FALSE;
REPEAT
    WRITELN ('ENTER UNIT # (4 OR 5)');
    READLN(DISK)
UNTIL (DISK=4) OR (DISK=5);
REPEAT
    WRITELN ('STARTING BLOCK? (0 <= BLOCK <= ', MAXBLOCK, ')');
    READLN(STARTBL)
UNTIL (STARTBL>=0) AND (STARTBL<=MAXBLOCK);
REPEAT
    WRITELN ('# OF BLOCKS TO DISPLAY?');
    READLN (LENGTH);
    IF STARTBL+LENGTH-1 >MAXBLOCK
    THEN WRITELN ('FINAL BLOCK WILL BE OFF THE DISK')
UNTIL STARTBL+LENGTH-1 <= MAXBLOCK;
IF LENGTH<=0 THEN QUIT := TRUE;
REPEAT
    WRITELN ('DECIMAL OR HEX? (D OR H)');
    READLN (BASE)
UNTIL (BASE = 'D') OR (BASE = 'H');
OK := FALSE;
WRITELN ('PRINTER OR SCREEN OUTPUT, OR QUIT?');
REPEAT
    WRITELN ('ENTER P, S, OR Q');    READLN(OPT).
    CASE OPT OF
    'P': BEGIN
        OK := TRUE;    PRTOPT := TRUE;    OUTFLNAM := 'PRINTER:'
        END;
    'S': BEGIN
        OK := TRUE;    CONOPT := TRUE;    OUTFLNAM := 'CONSOLE:';
        WRITELN; WRITELN;
        WRITELN ('CARRIAGE RETURN CONTINUES PROGRAM WHEN SCREEN FULL.');
        FOR I := 1 TO 5000 DO X := 1.1*2.2*3.3*4.4*5.5
        END;
    'Q': BEGIN
        OK := TRUE;    QUIT := TRUE
        END
    END
UNTIL OK;
```

```
IF NOT QUIT
THEN BEGIN
    RESET (OUTFL, OUTFLNAM);
    NUMLINE[4] := ' ';
    NUMLINE[5] := ' ';
    FOR POSIT := 1 TO LINELEN DO CHARLINE[POSIT] := ' ';
    FOR BLOCK := STARTBL TO (STARTBL+LENGTH-1) DO
        BEGIN
        UNITWAIT (DISK);
        UNITREAD (DISK, BLOCKIN, BLLENGTH, BLOCK);
        WRITELN (OUTFL);
        WRITELN (OUTFL, 'UNIT # ', DISK, '      BLOCK # ', BLOCK,
            '           BASE ', BASE);
        WRITELN (OUTFL);
        POSIT := 1;    (* ADDRESS STARTS IN COL 1 *)
        WRITENUM (0, BASE); (* FILL IN BYTE ADDRESS *)
        POSIT := 6;    (* VALUES START IN COL 6 *)
        LINES := 3;    (* 3 LINES DISPLAYED SO FAR *)
        FOR BYTE := 1 TO BLLENGTH DO (* FOR EVERY BYTE IN BLOCK *)
            BEGIN
            CHARAC := BLOCKIN[BYTE]; (* UNPACK BYTE & GET VALUE *)
            CHORD := ORD(CHARAC);
            IF (CHORD<=31) OR (CHORD>=126) (* IF NOT PRINTABLE *)
            THEN CHARAC := ' ';
            CHARLINE[POSIT+2] := CHARAC; (* R. JUST. CHAR OVER VALUE *)
            WRITENUM (CHORD, BASE);  (* PUT VALUE IN NUMLINE *)
            IF POSIT >= LINELEN (* IF LINE FULL THEN WRITE IT OUT *)
            THEN BEGIN          (* CENTRONICS DOESN'T DO AUTO LF-CR *)
                WRITE (OUTFL, CHARLINE);
                IF PRTOPT THEN WRITE (OUTFL, CHR(CR));
                WRITE (OUTFL, NUMLINE);
                IF PRTOPT THEN WRITE (OUTFL, CHR(CR), CHR(LF));
                POSIT := 1;
                LINES := LINES + 2; (* 2 MORE LINES PRINTED *)
                IF (CONOPT) AND (LINES >= MAXLINES)THEN
                    BEGIN (* IF SCREEN FULL, WAIT FOR CAR. RET. *)
                    LINES := 0;
                    READLN (OUTFL, CONTINUE)
                    END;
                WRITENUM (BYTE, BASE); (*FILL IN ADDR. OF NEXT BYTE*)
                POSIT := 6
                    END
            END;
```

```
        FOR I := POSIT TO LINELEN DO
            BEGIN    (* LAST LINE ISN'T FULL *)
            CHARLINE[I] := ' ';
            NUMLINE[I] := ' '
            END;
        WRITE (OUTFL, CHARLINE);
        IF PRTOPT THEN WRITE (OUTFL, CHR(CR));
        WRITE (OUTFL, NUMLINE);
        IF PRTOPT
        THEN BEGIN
            WRITE (OUTFL, CHR(CR));
            WRITE (OUTFL, CHR(CR));
            END;
        IF CONOPT THEN READLN (OUTFL, CONTINUE);
        END;
    CLOSE (OUTFL)
    END
END.
```

INTERDATA 8/32 TO MICROENGINE FILE TRANSFER PROGRAM

MEPASCAL runs on an Interdata 8/32 and never has direct contact with the Microengine. A one-way file transfer system was written to move files from the 8/32 to the Microengine. The sending program, FILEXFER, was written by Robert Young. The system uses a primitive protocol to transfer 512-byte data blocks over an asynchronous 4800 baud communication link. The UCSD Pascal program XFER is the implementation of the Microengine (receiving) side of that protocol.

The unit of transfer is a 516-byte packet in which the first 512 bytes contain the data. Byte 514 is always zero, except to indicate end-of-transmission, which is denoted by the value 7 (ASCII EOT). The packet which contains the end-of-transmission indicator is sent after the last packet of good data, so its data portion contains garbage and can be ignored. Byte 516 is a checksum on the data bytes only. The formula is the sum of the data bytes, modulus 256:

$$SUM := (SUM + BYTE) \text{ MOD } 256.$$

Bytes 513 and 515 are unused.

The Microengine side of the protocol proceeds as follows. An ENQ character (ASCII 5) is sent to indicate that the Microengine is ready to receive. The input status bit for the port is polled until a character is received, or a timeout occurs. The Microengine is expecting to receive a packet at this point. Since there could be quite a long wait between packets if the 8/32 is busy, the timeout period is quite long (a minute or more). This long period also gives the operator time to start FILEXFER on the 8/32. After the first character of the packet has arrived, the timeout period is made much shorter. The balance of the packet should be received within a few seconds of the first character. Again, the arrival of characters is

sensed by polling the status bits of the port. The number of characters received (less one) is displayed on the console, regardless of whether there is a timeout or not. So, if the entire packet is received, the number 515 is displayed. If the entie packet is received and there is no checksum error, an acknowledgement (ACK-- ASCII 6) is sent to the 8/32, otherwise a negative acknowledgement (NAK-- ASCII 21) is returned. The response consists of a single character, and is not sent until an ENQ is received, indicating that the 8/32 is ready to receive it. If the 8/32 receives a NAK, it will retransmit the packet, otherwise a new one will be sent. The Microengine sends an ENQ when it is ready for the next packet, as above.

Figure 43 shows the general format of the buffer into which the incoming data is placed. The buffer is large enough to hold 60 disk blocks of data (512 bytes each) plus four bytes. The extra four bytes are for the last four bytes of the 60th block. As characters come off the line they are placed in the input buffer. Once the checksum has been recalculated and it has been determined that the packet was received correctly, the last four bytes are unneeded, so they are overlaid by the data of the next packet. Disk accesses are quite slow on the Microengine and the large buffer size keeps the number of them to a minimum. An attempt was made to reserve a larger buffer by using negative and positive array bounds centered on zero, but the UCSD compiler crashed during the compilation. The present size has worked well, and as it is, it consumes nearly half of main storage.

Since the hardware does not yet support interrupts and does not have a timer, timeouts are calculated by decrementing a counter while polling the port's status bit.

At the present time, there is no port on the 8/32 dedicated to

FIGURE 43. Layout of the input buffer in XFER program. The last four bytes of a packet extend into the space for the data of the next block, and a data block overlays the last four bytes of the previous packet.

communication with the Microengine. This means that two plugs must be switched, and a new baud rate set on one of the ports. Unfortunately, this makes two terminals on the 8/32 unusable. Hopefully, the situation can be changed in the near future.

The file transfer system has exhibited a behavior which so far has been unexplainable. On a regular basis a small, variable number of characters never reach the Microengine. For example, in every fifth packet, from six to thirty characters are lost. This is merely an inconvenience, since the faulty packets are retransmitted and the file eventually arrives intact. Although this does slow things up, a 50K byte file can still be transferred in less than fifteen minutes.

The source code for both the sending and receiving programs follows.

```
1    "FILEXFER - ASC LINE FILE TRANSFER PROGRAM"
2    "WRITTEN BY ROBERT YOUNG"
3
4    "$INCLUDE SVC_TYPES"
5    "$INCLUDE FULL_PREFIX"
6
7    PROGRAM FILEXFER;
8    CONST EOT=7; ACK=6; NAK=21;
9    TYPE FILE_MODES = (BYTE_MODE, PAGE_MODE, ASC_MODE);
10   TYPE PAGE = ARRAY [1..512] OF BYTE;
11
12   VAR BUF: RECORD
13            DATA: PAGE;
14            EOF: SHORTINTEGER;
15            CKSUM: SHORTINTEGER;
16       END;
17   IN_MODE, OUT_MODE: FILE_MODES;
18   EM_FOUND: BOOLEAN;
19   BLOCKS: INTEGER;
20
21   PROCEDURE SVC1 (VAR PARM: SVC1 BLOCK); EXTERN;
22   PROCEDURE SVC7 (VAR PARM: SVC7 BLOCK); EXTERN;
23   PROCEDURE SVC2PAUS; EXTERN;
24
25   "DISPLAY ROUTINES"
26
27   PROCEDURE DISPLAY_TEXT (TEXT: LINE);
28   VAR I: INTEGER;
29   BEGIN
30       I:=1;
31       WHILE TEXT[I]<>'$' DO BEGIN
32           DISPLAY(TEXT[I]); I:=I+1;
33       END;
34   END;
35
36   PROCEDURE DISPLAY_HEX_CHAR (VAL: INTEGER);
37   BEGIN
38       IF VAL>9 THEN DISPLAY(CHR(VAL-10+ORD('A')))
39       ELSE DISPLAY(CHR(VAL+ORD('0')));
40   END;
41
42   PROCEDURE DISPLAY_HEX_CHARS (VAL: INTEGER);
43   BEGIN
44       DISPLAY_HEX_CHAR(VAL DIV 16);
45       DISPLAY_HEX_CHAR(VAL MOD 16);
46   END;
47
48   PROCEDURE DISPLAY_HEX (VAL: INTEGER; N: INTEGER);
49   TYPE TAGS = (TAG1, TAG2);
50   VAR I,J: INTEGER;
51       X: RECORD CASE TAGS OF
52            TAG1: (INT: INTEGER);
53            TAG2: (BYTES: ARRAY [1..4] OF BYTE);
54       END;
55   BEGIN
56       X.INT:=VAL;
57       FOR I:=5-N TO 4 DO
58           DISPLAY_HEX_CHARS(X.BYTES[I]);
59   END;
```

```
60   "INITIALIZATION"
61
62   PROCEDURE INIT_FILE (LU: INTEGER; VAR MODE: FILE_MODES);
63
64   VAR SVC7_PARM: SVC7_BLOCK;
65   BEGIN
66     WITH SVC7_PARM DO BEGIN
67       SVC7_LU:=LU; SVC7_MOD:=SVC7_FETCH_ATTR;
68       SVC7(SVC7_PARM);
69       IF SVC7_STAT<>0 THEN BEGIN
70         DISPLAY_TEXT('SVC7 ERROR $');
71         DISPLAY_HEX(SVC7_STAT,1);
72         DISPLAY_TEXT(' ON LU $');
73         DISPLAY_HEX(SVC7_LU,1);
74         DISPLAY(NL);
75         SVC7_MOD:=0; SVC7_RECLEN:=0;
76       END;
77       IF SVC7_MOD = 240 THEN MODE:=ASC_MODE
78       ELSE IF (SVC7_RECLEN=0) OR (SVC7_RECLEN>=512) THEN MODE:=PAGE_MODE
79       ELSE MODE:=BYTE_MODE;
80     END;
81   END;
82
83   "ERROR LOGGING"
84
85   PROCEDURE SVC1_ERROR_LOG (SVC1_PARM: SVC1_BLOCK);
86   BEGIN
87     WITH SVC1_PARM DO BEGIN
88       DISPLAY_TEXT('SVC1 ERROR $');
89       DISPLAY_HEX(SVC1_STAT,1);
90       DISPLAY_HEX(SVC1_DEV_STAT,1);
91       DISPLAY_TEXT(' ON LU $');
92       DISPLAY_HEX(SVC1_LU,1);
93       DISPLAY_TEXT(' FUNC $');
94       DISPLAY_HEX(SVC1_FUNC,1);
95       DISPLAY_TEXT(' XLEN $');
96       DISPLAY_HEX(SVC1_XFER_LEN,2);
97       DISPLAY_TEXT(' BLOCK $');
98       DISPLAY_HEX(BLOCKS,2);
99       DISPLAY(NL);
100    END;
101  END;
102
103  "ASC INPUT PROCESSING"
104
105  PROCEDURE READ_ASC (VAR EOF: BOOLEAN);
106  VAR SVC1_PARM: SVC1_BLOCK;
107      I, J: SHORTINTEGER;
108      REPLY: BYTE;
109      ERRORS: INTEGER;
110  BEGIN
111    ERRORS:=0;
112    WITH SVC1_PARM DO REPEAT
113      SVC1_LU:=1; SVC1_FUNC:=SVC1_READ+SVC1_IMAGE+SVC1_WAIT;
114      SVC1_BUFSTART:=ADDRESS(BUF);
115      SVC1_BUFEND:=SVC1_BUFSTART+515;
116      SVC1(SVC1_PARM);
117      IF SVC1_STAT<>0 THEN BEGIN
118        REPLY:=NAK;
119        SVC1_ERROR_LOG(SVC1_PARM);
```

```
120    END
121    ELSE REPLY:=ACK;
122    J:=0;
123    FOR I:=1 TO 512 DO
124    J:=J+(BUF.DATA[I]);
125    J:=J MOD 256;
126    IF ((J MOD 256)<>BUF.CKSUM) AND (REPLY=ACK) THEN BEGIN
127    DISPLAY_TEXT('CHECKSUM ERROR$');
128    DISPLAY_TEXT(' XPTD=$'); DISPLAY_HEX(J,2);
129    DISPLAY_TEXT(' RCVD=$'); DISPLAY_HEX(BUF.CKSUM,2);
130    DISPLAY(NL);
131    REPLY:=NAK;
132    END;
133    SVC1_FUNC:=SVC1_WRITE+SVC1_IMAGE+SVC1.WAIT;
134    SVC1_BUFSTART:=ADDRESS(REPLY);
135    SVC1_BUFEND:=SVC1_BUFSTART;
136    SVC1(SVC1_PARM);
137    IF SVC1_STAT<>0 THEN SVC1_ERROR_LOG(SVC1_PARM);
138    IF REPLY=NAK THEN ERRORS:=ERRORS+1;
139    UNTIL (ERRORS>16) OR (REPLY=ACK);
140    IF BUF.EOF = EOT THEN EOF:=TRUE ELSE EOF:=FALSE;
141    IF REPLY<>ACK THEN EOF:=TRUE; "FORCE ABORT"
142    END;
143
144    "ASC OUTPUT"
145
146    PROCEDURE WRITE_ASC (EOF: BOOLEAN);
147    VAR REPLY: BYTE;
148    SVC1_PARM: SVC1_BLOCK;
149    I: INTEGER;
150    ERRORS: INTEGER;
151    OK: BOOLEAN;
152    BEGIN
153    ERRORS:=0;
154    BUF.CKSUM:=0;
155    FOR I:=1 TO 512 DO
156    BUF.CKSUM:=BUF.CKSUM+(BUF.DATA[I]);
157    BUF.CKSUM:=BUF.CKSUM MOD 256;
158    IF EOF THEN BUF.EOF:=EOT ELSE BUF.EOF:=0;
159    ERRORS:=0;
160    WITH SVC1_PARM DO REPEAT
161    SVC1_LU:=2; SVC1_FUNC:=SVC1_WRITE+SVC1_IMAGE+SVC1_WAIT;
162    SVC1_BUFSTART:=ADDRESS(BUF);
163    SVC1_BUFEND:=SVC1_BUFSTART+515;
164    SVC1(SVC1_PARM);
165    IF SVC1_STAT<>0 THEN BEGIN
166    SVC1_ERROR_LOG(SVC1_PARM);
167    OK:=FALSE;
168    END
169    ELSE OK:=TRUE;
170    SVC1_FUNC:=SVC1_READ+SVC1_IMAGE+SVC1_WAIT;
171    SVC1_BUFSTART:=ADDRESS(REPLY);
172    SVC1_BUFEND:=SVC1_BUFSTART;
173    SVC1(SVC1_PARM);
174    IF SVC1_STAT<>0 THEN BEGIN
175    SVC1_ERROR_LOG(SVC1_PARM); OK:=FALSE;
176    END;
177    IF OK AND (REPLY <> ACK) THEN BEGIN
178    DISPLAY_TEXT('NEGATIVE ACKNOWLEDGEMENT $');
179    DISPLAY_HEX(REPLY,1); DISPLAY(NL);
```

```
180          OK:=FALSE;
181        END;
182        IF NOT OK THEN ERRORS:=ERRORS+1;
183      UNTIL (ERRORS>16) OR OK;
184    END;
185
186    PROCEDURE READ_FILE (VAR EOF: BOOLEAN);
187    VAR I: INTEGER; C: CHAR;
188    BEGIN
189      IF IN_MODE = ASC_MODE THEN READ_ASC(EOF)
190      ELSE IF IN_MODE = PAGE_MODE THEN READPAGE(BUF.DATA,EOF)
191      ELSE BEGIN
192        I:=1; EOF:=FALSE;
193        IF EM_FOUND THEN EOF:=TRUE
194        ELSE WHILE I<=512 DO BEGIN
195          READ(C);  BUF.DATA[I]:=ORD(C);
196          IF C = EM THEN BEGIN
197            I:=512; EM_FOUND:=TRUE;
198          END;
199          I:=I+1;
200        END;
201      END;
202    END;
203
204    PROCEDURE WRITE_FILE (EOF: BOOLEAN);
205    VAR I: INTEGER; C: CHAR;
206    BEGIN
207      IF OUT_MODE = ASC_MODE THEN WRITE_ASC(EOF)
208      ELSE IF OUT_MODE = PAGE_MODE THEN WRITEPAGE(BUF.DATA,EOF)
209      ELSE IF NOT EOF THEN BEGIN
210        I:=1;
211        WHILE I<=512 DO BEGIN
212          C:=CHR(BUF.DATA[I]);
213          WRITE(C);
214          IF C=EM THEN I:=512;
215          I:=I+1;
216        END;
217      END;
218    END;
219
220    PROCEDURE RUN;
221    VAR EOF: BOOLEAN;
222    BEGIN
223      EM_FOUND:=FALSE;
224      BLOCKS:=1;
225      INIT_FILE(1,IN_MODE);
226      INIT_FILE(2,OUT_MODE);
227      REPEAT
228        READ_FILE(EOF);
229        WRITE_FILE(EOF);
230        BLOCKS:=BLOCKS+1;
231      UNTIL EOF;
232    END;
233
234    BEGIN
235      RUN
236    END.
```

CROSS REFERENCE   * IS DEF   = IS ASN

-A-

ACK	8*	121	126	139	141	177
ADDRESS	114	134	162	171		
ASC_MODE	9	77	189	207		

-B-

BLOCKS	19*	98	224=	230=	230		156	157	157	158	158	162	190	195	208
BOOLEAN	18	105	146	151	186	180	154								
BMF	12*	114	124	126	129	221									
	212														
BYTE	10	53	108	147											
BYTES	53	58													
BYTE_MODE	9	79													

-C-

C	187*	195	195	196	205*	212=	213	214		157	157	158	158
CHAR	187	205	212	212=	213								
CKSUM	38	39	126	129	154=	156	157=	157					

-D-

DATA	13*	124	156	190	195	208	212	214		96	98	128	129	179
DISPLAY	32	38	39	74	99	130	179							
DISPLAY_HEX	48*	71	73	89	90	92	94		96	98	128	129	179	
DISPLAY_HEX_	36*	42*	45	58	91	95		97	127	128	129	178		
DISPLAY_TEXT	27*	70	72	88	91	93								

-E-

EM	196	214											
EM_FOUND	18*	193	197*	180	181=	186*	186*	189	190	192	192=	193=	204*	207	208
EOF	14*	105*	140	140=	223=	140=	158	158	158=	186*	189	190		
	209	221*	228	229	231									
EOT	8*	140	150*	153=	159*	182=	182	183						
ERRORS	109*	111=	130=	138	139									

-F-

FALSE	140	167	175	192	223	192								
FILEXFER	7*	17	63											
FILE_MODES	9*													

-I-

I	28*	30=	31	32	32=	32	50*	57=	58	107*	123=	124	149*	155=	156	187*	192=	194
INIT_FILE	195	197*	199*	199	205*	210=	211	212	214=	215*	213							
INT	63*	56=	225	226														
INTEGER	19	28	36	42	48	48	50	52	63	109	149	150	187	205				
IN_MODE	17*	189	190	225				189										

-J-

| J | 50* | 107* | 122= | 124 | 125= | 126 | 128 | 128 | | | | | | |

-L-

| LINE | 27 | | | | | | | | | | | | |
| LU | 63* | 67 | | | | | | | | | | | |

-M-

| MODE | 63* | 77= | 78= | 79- | | | | | | | | | |

```
-W-
WRITE              213
WRITEPAGE          208
WRITE_ASC          146*  207
WRITE_FILE         204*  229

-X-
X                  51*   56    58

END XREF   86 IDENTIFIERS   413 TOTAL REFERENCES
   33 COLLISIONS.
```

```
(* Q+ *)
(* LV1:ERRORS. TEXT *)
(* LPRINTER: *)

PROGRAM X832_WD;
(********************************************************************)
(*                                                                 *)
(* PROGRAM TO RECEIVE DATA FROM THE INTERDATA 8/32  SERIAL,        *)
(* ASYNCHRONOUS LINE TRANSMISSION.                                 *)
(*                       .                                         *)
(********************************************************************)
CONST
    BUFFSIZE   =   30724;     (* 60 DISK BLOCKS + 4 BYTES *)
    BLLEN      =     512;     (* BYTES IN DISK BLOCK *)
    BLINBUFF   =      60;     (* DISK BLOCKS IN BUFFER *)
    FLAGDIST   =     513;     (* DISTANCE FROM BEGINNING OF BLOCK TO
                                 "LASTBLOCK" FLAG *)
    CRCDIST    =     515;     (* DISTANCE TO CRC BYTE *)
                             (* 1 TIME UNIT = ABOUT .22 SECOND *)
    LONGTIME   =  (* TEST 1000 *) 200;   (* LONG TIMEOUT, ABOUT 5.5 MIN *)
    MEDTIME    =     136;     (* MEDIUM TIMEOUT, ABOUT 30 SEC *)
    SHRTTIME   =      68;     (* SHORT TIMEOUT, ABOUT 15 SEC *)
    PORT_B     =    -992;     (* PORT B (REMOTE UNIT) HAS
                                 ADDRESS -992 (FC20 HEX) *)
    PKTLEN     =     516;     (* DATA PACKET IS 516 BYTES:
                                 512 OF DATA,
                                   1 UNUSED,
                                   1 "LASTBLOCK" FLAG,
                                   1 UNUSED,
                                   1 CRC *)
```

```
(* SEE MICROENGINE USER'S MANUAL (PP 35-41) FOR DETAILS ON SERIAL
   PORT REGISTERS *)
(* BITS IN SERIAL PORT STATUS REGISTER *)
SNDEMPTY = 0;
RECFULL  = 1;
OVERRUN  = 2;
PARITY   = 3;
FRAMING  = 4;
CARRIER  = 5;
READY    = 6;
CHANGE   = 7;

   (* BITS IN SERIAL PORT CONTROL REGICTER 1 *)
RNGREADY = 0;
REQTOSND = 1;
RECENABL = 2;
PARENABL = 3;
ECHO     = 4;
STOPBITS = 5;
BREAK    = 6;
LOOPNORM = 7;

   (* BITS IN SERIAL PORT CONTROL REGISTER 2 *)
CLOCK0   = 0;
CLOCK1   = 1;
CLOCK2   = 2;
RCVRRATE = 3;
PARITSET = 4;
CHARMODE = 5;
CHARLEN6 = 6;
CHARLEN7 = 7;

TYPE
   STATBITS = SNDEMPTY..CHANGE;
   CR1BITS  = RNGREADY..LOOPNORM;
   CR2BITS  = CLOCK0  ..CHARLEN7;
   TWOCASES = 1..2;
   BYTETYPE = PACKED ARRAY [1..1] OF CHAR;

   SERPORT = RECORD    (* SERIAL PORT REGISTERS *)
               DATA:   CHAR;
               STATUS: PACKED ARRAY [STATBITS] OF BOOLEAN;
               CR2:    PACKED ARRAY [CR2BITS ] OF BOOLEAN;
               CR1:    PACKED ARRAY [CR1BITS ] OF BOOLEAN
             END;

   SERDEV = RECORD CASE TWOCASES OF
              1: (ADDRESS: INTEGER);
              2: (PORTPTR: ^SERPORT)
              END;
```

```
VAR
    BUFF: PACKED ARRAY [1..BUFFSIZE] OF CHAR;  (* INPUT BUFFER *)
    LINE: SERDEV;
    EOT, ACK, NAK, ENQ: CHAR;     (* ASCII CONTROL CHARS *)
    OK,      (* RESPONSE MEANING PACKET WAS REC'D OK *)
    NOTOK,   (* RESPONSE MEANING PACKET NOT REC'D OK *)
    FLAGBYTE: CHAR;   (* CONTAINS EOT IF THIS IS LAST PACKET
                         IN THE TRANSMISSION. ANY OTHER CODE
                         MEANS NOT LAST PACKET *)
    BLOCKS,     (* # OF BLOCKS TO BE WRITTEN TO DISK FROM INPUT BUFFER *)
    BLOKSIN,    (* # OF BLOCKS REC'D OK OFF THE LINE *)
    BLOKSOUT,   (* # OF BLOCKS SUCCESSFULLY WRITTEN TO DISK *)
    STARTBYT: INTEGER;  (* BYTE WITHIN BUFFER WHERE THE CURRENT
                           PACKET (AND BLOCK) STARTS *)
    DISKFILE: FILE;   (* NAME BY WHICH THIS PGM. KNOWS THE OUTPUT FILE *)
    OUTFLNAM: STRING[20];  (* NAME OF OUTPUT FILE WHICH WILL APPEAR IN
                              THE DISK'S VOLUME DIRECTORY *)
    ERROR,   (* TRUE= SOME SORT OF ERROR WAS DETECTED *)
    FATALERR, (* TRUE= AN ERROR WAS DETECTED & IT IS FATAL *)
    LASTPKT: BOOLEAN;  (* TRUE= JUST RECEIVED THE LAST PACKET IN THE
                          TRANSMISSION *)
       I: INTEGER;   (* TEST *)
```

```
PROCEDURE RESPOND (CCHAR: CHAR; VAR ABORT: BOOLEAN);
(* WHEN S/22 IS READY TO RECEIVE, IT SENDS AN ENQ CHAR.
   PROCEDURE WAITS FOR ENQ THEN SENDS RESPONSE MESSAGE OF A
   SINGLE CHAR *)
VAR
   TIME: INTEGER;
   EXPCTENQ: CHAR;
   ABYTE: BYTETYPE;
   ERROR: BOOLEAN;
   TICK: INTEGER;
   HAVECHAR: BOOLEAN;
BEGIN
ABORT := FALSE;
TIME := MEDTIME;
ERROR := FALSE;

REPEAT
   TICK := 1000;
   REPEAT   (* POLL LINE UNTIL CHAR ARRIVES OR TIME TO DECR. TIME PARAM *)
      HAVECHAR := LINE.PORTPTR^.STATUS [RECFULL];
      (* TEST  IF TICK <= 1000 THEN HAVECHAR := TRUE; *)
      TICK := TICK - 1
   UNTIL HAVECHAR OR (TICK <= 0);
   IF NOT HAVECHAR
   THEN TIME := TIME -1
UNTIL HAVECHAR OR (TIME <= 0);

(* GET WHATEVER IS IN RECEIVER HOLDING REGISTER *)
ABYTE[1] := LINE.PORTPTR^.DATA;
(* TEST  ABYTE[1] := ENQ; *)
WITH LINE.PORTPTR^ DO
   BEGIN
   ERROR := STATUS [OVERRUN] OR (*CHECK FOR ERRORS*)
            STATUS [PARITY] OR
            STATUS [FRAMING];
   IF STATUS[OVERRUN] THEN WRITELN ('OVERRUN');
   IF STATUS[PARITY]  THEN WRITELN('PARITY');
   IF STATUS[FRAMING] THEN WRITELN('FRAMING')
   END;
```

```
PROCEDURE READPKT (STARTBYT: INTEGER; VAR PKTERROR, ABORT: BOOLEAN);
VAR
    BYT,
    LASTBYT,
    TIME: INTEGER;
    ERROR: BOOLEAN;
    ABYTE: BYTETYPE;
    TICK: INTEGER;
    HAVECHAR: BOOLEAN;
    I: INTEGER;     (* TEST *)
    R. INTEGER;     (* TEST *)
BEGIN
PKTERROR := FALSE;
ABORT := FALSE;
BYT := STARTBYT;
LASTBYT := STARTBYT + PKTLEN;
LINE. PORTPTR^. DATA := ENQ;
TIME := LONGTIME;                          .
ERROR := FALSE;

REPEAT
    TICK := 1000;
    REPEAT   (* POLL LINE UNTIL CHAR ARRIVES OR TIME TO DECR. TIME PARAM *)
        HAVECHAR := LINE. PORTPTR^. STATUS [RECFULL];
        (* TEST  IF TICK <= 1000 THEN HAVECHAR := TRUE; *)
        TICK := TICK - 1
    UNTIL HAVECHAR OR (TICK <= 0);
    IF NOT HAVECHAR
    THEN TIME := TIME -1
UNTIL HAVECHAR OR (TIME <= 0);

(* GET WHATEVER IS IN RECEIVER HOLDING REGISTER *)
ABYTE[1] := LINE. PORTPTR^. DATA;
(* TEST  ABYTE[1] := 'A';  *)
WITH LINE. PORTPTR^ DO
    BEGIN
    ERROR := STATUS [OVERRUN] OR (*CHECK FOR ERRORS*)
             STATUS [PARITY ] OR
             STATUS [FRAMING];
    IF STATUS[OVERRUN] THEN WRITELN ('OVERRUN');
    IF STATUS[PARITY]  THEN WRITELN('PARITY');
    IF STATUS[FRAMING] THEN WRITELN('FRAMING')
    ;  IF ERROR THEN WRITELN ('ERROR 1')
    END;
BUFF[BYT] := ABYTE[1];
```

```
EXPCTENQ := ABYTE[1];
IF TIME <= 0
THEN BEGIN      (* TIMED OUT *)
   WRITELN;
   WRITELN ('PROCEDURE RESPOND:');
   WRITELN ('TIMED OUT WAITING FOR ENQ');
   ABORT := TRUE
   END;
IF EXPCTENQ <> ENQ
THEN BEGIN (* EXPECTING AN ENQ BUT ANY CHAR WILL DO FOR SYNC PURPOSES *)
   WRITELN;
   WRITELN ('PROCEDURE RESPOND:');
   WRITELN ('RECEIVED ', EXPCTENQ);
   WRITELN ('EXPECTED ENQ. CONTINUING...')
   END;
LINE.PORTPTR^.DATA := CCHAR (* PUT THE RESPONSE CHAR ON LINE *)
:   IF CCHAR=ACK THEN WRITELN ('SENDING ACK')
ELSE IF CCHAR=NAK THEN WRITELN ('SENDING NAK')
ELSE WRITELN ('SENDING NEITHER ACK NOR NAK');
END;


FUNCTION CRC_OK (STARTBYT: INTEGER): BOOLEAN;
(* RECALCULATE THE CHECKSUM & SEE IF IT MATCHES CHECKSUM IN PACKET *)
(* ONLY BYTES IN THE DATA PORTION OF THE PACKET ARE USED TO        *)
(* CALCULATE THE CHECKSUM. FORMULA IS SUM OF BYTES MOD 256.        *)
   VAR
   BYTE,
   CALC_CRC: INTEGER;   (* THE CALCULATED CHECKSUM *)
   CHARA: CHAR;
BEGIN
CALC_CRC := 0;
FOR BYTE := STARTBYT TO (STARTBYT+BLLEN-1) DO (*RECALCULATION*)
   BEGIN
   CHARA := BUFF [BYTE];
   CALC_CRC := (CALC_CRC + ORD(CHARA)) MOD 256
   END;
CHARA := BUFF [STARTBYT + CRCDIST];   (* GET CRC FROM PACKET *)
CRC_OK := CALC_CRC = ORD(CHARA)       (* COMPARE & RETURN RESULT *)
(* TEST *) ; CRC_OK := TRUE
END;
```

ı

```
PROCEDURE READPKT (STARTBYT: INTEGER; VAR PKTERROR, ABORT: BOOLEAN);
VAR
    BYT,
    LASTBYT,
    TIME: INTEGER;
    ERROR: BOOLEAN;
    ABYTE: BYTETYPE;
    TICK: INTEGER;
    HAVECHAR: BOOLEAN;
    I: INTEGER;      (* TEST *)
    R: INTEGER;      (* TEST *)
BEGIN
PKTERROR := FALSE;
ABORT := FALSE;
BYT := STARTBYT;
LASTBYT := STARTBYT + PKTLEN;
LINE. PORTPTR^. DATA := ENQ;
TIME := LONGTIME;
ERROR := FALSE;

REPEAT
    TICK := 1000;
    REPEAT   (* POLL LINE UNTIL CHAR ARRIVES OR TIME TO DECR. TIME PARAM *)
        HAVECHAR := LINE. PORTPTR^. STATUS [RECFULL];
        (* TEST  IF TICK <= 1000 THEN HAVECHAR := TRUE; *)
        TICK := TICK - 1
    UNTIL HAVECHAR OR (TICK <= 0);
    IF NOT HAVECHAR
    THEN TIME := TIME -1
UNTIL HAVECHAR OR (TIME <= 0);

(* GET WHATEVER IS IN RECEIVER HOLDING REGISTER *)
ABYTE[1] := LINE. PORTPTR^. DATA;
(* TEST  ABYTE[1] := 'A';  *)
WITH LINE. PORTPTR^ DO
    BEGIN
    ERROR := STATUS [OVERRUN] OR (*CHECK FOR ERRORS*)
             STATUS [PARITY ] OR
             STATUS [FRAMING];
    IF STATUS[OVERRUN] THEN WRITELN ('OVERRUN');
    IF STATUS[PARITY]  THEN WRITELN('PARITY');
    IF STATUS[FRAMING] THEN WRITELN('FRAMING')
    ;  IF ERROR THEN WRITELN ('ERROR 1')
    END;
BUFF[BYT] := ABYTE[1];
```

```
IF ERROR
THEN PKTERROR := TRUE;
IF TIME <= 0
THEN BEGIN
    WRITELN;
    WRITELN ('PROC READPKT:');
    WRITELN ('TIMEOUT BETWEEN PACKETS');
    ABORT := TRUE
    END;
BYT := SUCC(BYT);
R := 0;
REPEAT
    TIME := SHRTTIME;
    ERROR := FALSE;

    REPEAT
        TICK := 1000;
        REPEAT   (* POLL LINE UNTIL CHAR ARRIVES OR TIME TO DECR.  TIME PARAM *)
            HAVECHAR := LINE.PORTPTR^.STATUS [RECFULL];
            (* TEST  IF TICK <= 1000 THEN HAVECHAR := TRUE; *)
            TICK := TICK - 1
        UNTIL HAVECHAR OR (TICK <= 0);
        IF NOT HAVECHAR
        THEN TIME := TIME -1
    UNTIL HAVECHAR OR (TIME <= 0);

    (* GET WHATEVER IS IN RECEIVER HOLDING REGISTER *)
    ABYTE[1] := LINE.PORTPTR^.DATA;
    (* TEST  ABYTE[1] := 'B'; *)
    WITH LINE.PORTPTR^ DO
        BEGIN
        ERROR := STATUS [OVERRUN] OR (*CHECK FOR ERRORS*)
                 STATUS [PARITY ] OR
                 STATUS [FRAMING];
        IF STATUS[OVERRUN] THEN WRITELN ('OVERRUN');
        IF STATUS[PARITY]  THEN WRITELN('PARITY');
        IF STATUS[FRAMING] THEN WRITELN('FRAMING')
        ;  IF ERROR THEN WRITELN ('ERROR 2')
        END;

    BUFF[BYT] := ABYTE[1];
    IF ERROR
    THEN PKTERROR := TRUE;
    BYT := SUCC(BYT)
    ;  R := SUCC(R)
UNTIL (BYT >= LASTBYT) OR (TIME <= 0);

IF TIME <= 0
THEN PKTERROR := TRUE   (* TIMED OUT *)
;  (* TEST  FOR I := STARTBYT TO LASTBYT DO WRITE(BUFF[I]);   *)
WRITELN(R)
END;
```

```
PROCEDURE INITLINE;
(* INITIALIZE SERIAL PORT REGISTERS FOR PORT ASSOCIATED WITH LINE *)
BEGIN
LINE.ADDRESS := PORT_B;
WITH LINE.PORTPTR^ DO
    BEGIN
    CR1 [RNGREADY] := TRUE;     (* ENABLE CARRIER & DATA SET READY BITS *)
    CR1 [REQTOSND] := TRUE;     (* ENABLE TRANSMISSION *)
    CR1 [RECENABL] := TRUE;     (* ENABLE RECEPTION *)
    CR1 [PARENABL] := FALSE;    (* NO HDWRE PARITY GENERATION/CHECKING *)
    CR1 [ECHO    ] := FALSE;    (* DON'T ECHO RECEIVED DATA *)
    CR1 [STOPBITS] := TRUE;     (* 1 STOP BIT *)
    CR1 [BREAK   ] := FALSE;
    CR1 [LOOPNORM] := TRUE;     (* NORMAL (NOT TEST) OPERATION *)

    CR2 [CLOCK0  ] := TRUE;     (* CLOCK RATE 1 (32X) *)
    CR2 [CLOCK1  ] := FALSE;
    CR2 [CLOCK2  ] := FALSE;
    CR2 [RCVRRATE] := FALSE;    (* RECEIVER CLOCK RATE = RATE 1 *)
    CR2 [PARITSET] := FALSE;    (* EVEN PARITY, IF APPLICABLE *)
    CR2 [CHARMODE] := FALSE;    (* ASYNCHRONOUS LINE *)
    CR2 [CHARLEN6] := FALSE;    (* DATA CHARACTERS ARE 8 BITS *)
    CR2 [CHARLEN7] := FALSE;
    END
END;


PROCEDURE INIT;
BEGIN
STARTBYT := 1;
BLOCKS := 0;
BLOKSIN := 0;
BLOKSOUT :=0;

ACK := CHR(6);    OK := ACK;
NAK := CHR(21);   NOTOK := NAK;
ENQ := CHR(5);
EOT := CHR(7);

WRITELN;
WRITELN ('NAME OF OUTPUT FILE?');
READLN (OUTFLNAM);
REWRITE (DISKFILE, OUTFLNAM);
IF IORESULT <> 0
THEN BEGIN
    WRITELN;
    WRITELN ('PROCEDURE INIT:');
    WRITELN ('I/O ERROR # ', IORESULT);
    HALT
    END;
INITLINE
END;
```

```
PROCEDURE FINISH (ABORTING: BOOLEAN);
BEGIN
CLOSE (DISKFILE, LOCK);
WRITELN;
WRITELN ('NUMBER OF BLOCKS WRITTEN = ', BLOKSOUT);
IF ABORTING
THEN WRITELN ('DATA TRANSFER ABORTED')
ELSE WRITELN ('DATA TRANSFER FINISHED')
END;


BEGIN
FOR I := 1 TO BUFFSIZE DO BUFF[I] := '/';
INIT;
REPEAT
    READPKT (STARTBYT, ERROR, FATALERR);
    IF CRC_OK (STARTBYT) AND NOT (ERROR OR FATALERR)
    THEN BEGIN
       IF NOT FATALERR
       THEN BEGIN
          BLOCKS := SUCC(BLOCKS);
          FLAGBYTE := BUFF [STARTBYT+FLAGDIST];
          LASTPKT := FLAGBYTE = EOT;
          (* TEST  IF BLOCKS >= 5 THEN LASTPKT := TRUE; *)
          IF (BLOCKS >= BLINBUFF) OR LASTPKT
          THEN BEGIN  (* BUFFER FULL OR RECEIVED LAST PACKET *)
             BLOKSIN := BLOKSIN + BLOCKS;
             BLOKSOUT := BLOKSOUT + BLOCKWRITE
                           (DISKFILE, BUFF, BLOCKS, BLOKSOUT);
             (* IF IORESULT <> 0
             THEN WRITELN ('DISK WRITE ERROR. # ', IORESULT); *)
             BLOCKS := 0;
             STARTBYT := 1
             END
          ELSE STARTBYT := STARTBYT + BLLEN
          END;
       RESPOND (OK, FATALERR)
       END
    ELSE IF NOT FATALERR
         THEN RESPOND (NOTOK, FATALERR)
UNTIL LASTPKT OR FATALERR;
FINISH (FATALERR)
END.
```

APPENDIX B

SOURCE CODE FOR MEPASCAL COMPILER PASS SIX: MEPASS6

```
 1    "############################################"
 2    "# MODIFIED TO PRODUCE CODE FOR MICROENGINE. #"
 3    "############################################"
 4
 5    "CODE TO BE FINALIZED LATER IS MARKED $$$"
 6
 7    "%SRC_32 := FALSE"
 8    "%SRC_16 := NOT SRC_32"
 9
10    "PER BRINCH HANSEN
11    INFORMATION SCIENCE
12    CALIFORNIA INSTITUTE OF TECHNOLOGY
13    PASADENA, CALIFORNIA 91125
14    PDP 11/45 CONCURRENT/SEQUENTIAL PASCAL.
15    COMPILER PASS 6: CODE SELECTION
16    FOR PASCAL MICROENGINE."
17
18    TYPE FULLWORD = INTEGER;
19    TYPE INTEGER = SHORTINTEGER;
20
21    "##########
22    # PREFIX #
23    ##########"
24
25    CONST        EOL = '(:10:)';      FF = '(:12:)';      EOM = '(:25:)';
26    MAXDIGIT = 6;
27    PRINTLIMIT = 18;
28    WORDLENGTH = 2 "BYTES";
29    REALLENGTH = 4 "BYTES";
30    SETLENGTH = 16 "BYTES";
31    "ON STACK, WORD LENGTH (8) OF SET IS ALWAYS PUSHED
32    AFTER THE SET ITSELF, TO CONFORM TO MICROENGINE CONVENTION.
33    MICROE SUPPORTS VARIABLE-LENGTH SETS, BUT THEY ARE
34    NOT USED BY THIS VERSION OF CONCURRENT PASCAL."
35
36    SOURCE_WORD_LENGTH = "%IF SRC_16" 2; "%END"
37                         "%IF SRC_32" 4; "%END"
38    LISTOPTION = 0;  SUMMARYOPTION = 1;  TESTOPTION = 2;  CHECKOPTION = 3;
39    CODEOPTION = 4;  NUMBEROPTION = 5;  VARNTCHECKOPTION = 8;  "VARNT"
40    MAXWORD = 100;
41
42    TYPE FILE = 1..4;
43
44    CONST IDLENGTH = 12;
45    TYPE IDENTIFIER = ARRAY (.1..IDLENGTH.) OF CHAR;
46
47    TYPE  POINTER = @ INTEGER;
48    TABLEPTR = @TABLE;
49    TABLE = RECORD
50        NEXTPORTION: TABLEPTR;
51        CONTENTS: ARRAY (.1..MAXWORD.) OF INTEGER
52        END;
53    TABLEPART = RECORD
54        SEGDISTANCE, STACKLENGTH: INTEGER;
55        JUMPTABLE, CONSTTABLE, XJFTABLE,
56        EXITICTABLE, DATASIZETABLE: TABLEPTR
57        END;
58    TABLESPTR = @TABLEPART;
59    OPTION = LISTOPTION..VARNTCHECKOPTION;       "VARNT"
```

```
 60  CONST MAXINTFAC = 14;
 61  TYPE
 62    IFPTR = @IFINFO;
 63    IFINFO = RECORD
 64      INTERFACES: INTEGER;
 65      INTERFACESIZES: ARRAY[1..MAXINTFAC] OF INTEGER
 66      END;
 67
 68    PASSPTR = @PASSLINK;
 69    PASSLINK = RECORD
 70      OPTIONS: SET OF OPTION;
 71      LABELS, BLOCKS, CONSTANTS, XJP_OFFSETS: INTEGER;
 72      RESETPOINT: FULLWORD;
 73      TABLES: TABLESPTR;
 74      INTERFACE: IFPTR
 75      END;
 76
 77  TYPE ARGTAG =
 78    (NILTYPE, BOOLTYPE, INTTYPE, IDTYPE, PTRTYPE);
 79
 80  TYPE ARGTYPE = RECORD
 81      CASE TAG: ARGTAG OF
 82        NILTYPE, BOOLTYPE: (BOOL: BOOLEAN);
 83        INTTYPE: (INT: INTEGER);
 84        IDTYPE: (ID: IDENTIFIER);
 85        PTRTYPE: (PTR: PASSPTR)
 86      END;
 87
 88  CONST MAXARG = 10;
 89        TEXT_LENGTH = 18;
 90  TYPE ARGLIST = ARRAY (.1..MAXARG.) OF ARGTYPE;
 91       TEXT_TYPE = ARRAY (.1..TEXT_LENGTH.) OF CHAR;
 92
 93  CONST PAGELENGTH = "%IF SRC_16" 256; "%END"
 94                     "%IF SRC_32" 128; "%END"
 95  TYPE PAGE = ARRAY (.1..PAGELENGTH.) OF INTEGER;
 96
 97  PROCEDURE READ(VAR C: CHAR);
 98  PROCEDURE WRITE(C: CHAR);
 99  PROCEDURE NOTUSED1;
100  PROCEDURE NOTUSED2;
101  PROCEDURE GET(F: FILE; P: INTEGER; VAR BLOCK: UNIV PAGE);
102  PROCEDURE PUT(F: FILE; P: INTEGER; BLOCK: UNIV PAGE);
103  FUNCTION FILE_LENGTH(F:FILE): INTEGER;
104  PROCEDURE MARK(VAR TOP: FULLWORD);
105  PROCEDURE RELEASE(TOP: FULLWORD);
106
107
108  PROGRAM MAIN(VAR PARAM: ARGLIST);
109  "##########################################"
110  # PASS(VAR OK: BOOLEAN; VAR LINK: POINTER) #
111  "##########################################"
112
113  CONST
114
115
116  "INPUT OPERATORS"
117  "(ASSIGNTAG1 IS NEVER PRODUCED BY PASS 5.)"
118  "(INITVAR1 IS NEVER PRODUCED BY PASS 5.)"
119
```

```
120 PUSHCONST1 = 0;      PUSHVAR1 = 1;      PUSHIND1 = 2;       PUSHADDR1 = 3;
121 FIELD1 = 4;          INDEX1 = 5;        POINTER1 = 6;       VARIANT1 = 7;
122 RANGE1 = 8;          ASSIGN1 = 9;       ASSIGNTAG1 = 10;    COPY1 = 11;
123 NEW1 = 12;           NOT1 = 13;         AND1 = 14;          OR1 = 15;
124 NEG1 = 16;           ADD1 = 17;         SUB1 = 18;          MUL1 = 19;
125 DIV1 = 20;           MOD1 = 21;         "NOT USED"          "NOT USED"
126 FUNCTION1 = 24;      BUILDSET1 = 25;    COMPARE1 = 26;      COMPSTRUC1 = 27;
127 FUNCVALUE1 = 28;     DEFLABEL1 = 29;    JUMP1 = 30;         FALSEJUMP1 = 31;
128 CASEJUMP1 = 32;      INITVAR1 = 33;     CALL1 = 34;         ENTER1 = 35;
129 RETURN1 = 36;        POP1 = 37;         NEWLINE1 = 38;      ERROR1 = 39;
130 CONSTANT1 = 40;      MESSAGE1 = 41;     INCREMENT1 = 42;    DECREMENT1 = 43;
131 PROCEDURE1 = 44;     INIT1 = 45;        PUSHLABEL1 = 46;    CALLPROG1 = 47;
132 FOR1=48;             PUTTOS1 = 49;
133
134 "VIRTUAL DATA TYPES"
135 BYTETYPE = 0;        WORDTYPE = 1;      REALTYPE = 2;       SETTYPE = 3;
136
137 "VIRTUAL ADDRESSING MODES"
138
139 MODE0 = 0  "CONSTANT";
140 MODE1 = 1  "PROCEDURE";
141 MODE2 = 2  "PROGRAM";
142 MODE3 = 3  "PROCESS ENTRY";
143 MODE4 = 4  "CLASS ENTRY";
144 MODE5 = 5  "MONITOR ENTRY";
145 MODE6 = 6  "PROCESS";
146 MODE7 = 7  "CLASS";
147 MODE8 = 8  "MONITOR";
148 MODE9 = 9  "STANDARD";
149 MODE10=10  "UNDEFINED";
150 MODE11=12  "MANAGER";                                       "MGR"
151 MODE12=13  "MGR ENTRY";                                     "MGR"
152
153 "COMPARISON OPERATORS"
154
155 LESS = 0;            EQUAL = 1;         GREATER = 2;        NOTLESS = 3;
156 NOTEQUAL = 4;        NOTGREATER = 5;    INSET = 6;
157
158 "STANDARD FUNCTIONS"
159
160 TRUNC1 = 0;          ABS1 = 1;          SUCC1 = 2;          PRED1 = 3;
161 CONV1 = 4;           EMPTY1 = 5;        ATTRIBUTE1 = 6;     REALTIME1 = 7;
162 MIN_FUNC = 0;        MAX_FUNC = 7;
163
164 "STANDARD PROCEDURES"
165
166 DELAY1 = 0;          CONTINUE1 = 1;     IO1 = 2;            START1 = 3;
167 STOP1 = 4;           SETHEAP1 = 5;      WAIT1 = 6;
168 MIN_PROC = 0;        MAX_PROC = 6;
169
170 "OUTPUT OPERATORS (MICROENGINE OPERATION CODES)"
171
172 SLDC00_2 = 000;      SLDC01_2 = 001;    SLDC02_2 = 002;     SLDC03_2 = 003;
173 SLDC04_2 = 004;      SLDC05_2 = 005;    SLDC06_2 = 006;     SLDC07_2 = 007;
174 SLDC08_2 = 008;      SLDC09_2 = 009;    SLDC10_2 = 010;     SLDC11_2 = 011;
175 SLDC12_2 = 012;      SLDC13_2 = 013;    SLDC14_2 = 014;     SLDC15_2 = 015;
176 SLDC16_2 = 016;      SLDC17_2 = 017;    SLDC18_2 = 018;     SLDC19_2 = 019;
177 SLDC20_2 = 020;      SLDC21_2 = 021;    SLDC22_2 = 022;     SLDC23_2 = 023;
178 SLDC24_2 = 024;      SLDC25_2 = 025;    SLDC26_2 = 026;     SLDC27_2 = 027;
179 SLDC28_2 = 028;      SLDC29_2 = 029;    SLDC30_2 = 030;     SLDC31_2 = 031;
```

```
180 SLDL01_2 = 032;   SLDL02_2 = 033;   SLDL03_2 = 034;   SLDL04_2 = 035;
181 SLDL05_2 = 036;   SLDL06_2 = 037;   SLDL07_2 = 038;   SLDL08_2 = 039;
182 SLDL09_2 = 040;   SLDL10_2 = 041;   SLDL11_2 = 042;   SLDL12_2 = 043;
183 SLDL13_2 = 044;   SLDL14_2 = 045;   SLDL15_2 = 046;   SLDL16_2 = 047;
184 SLDO01_2 = 048;   SLDO02_2 = 049;   SLDO03_2 = 050;   SLDO04_2 = 051;
185 SLDO05_2 = 052;   SLDO06_2 = 053;   SLDO07_2 = 054;   SLDO08_2 = 055;
186 SLDO09_2 = 056;   SLDO10_2 = 057;   SLDO11_2 = 058;   SLDO12_2 = 059;
187 SLDO13_2 = 060;   SLDO14_2 = 061;   SLDO15_2 = 062;   SLDO16_2 = 063;
188    "NO MICROENGINE OPERATION CODES FROM 64 TO 119."
189 SIND0_2 = 120;    SIND1_2 = 121;    SIND2_2 = 122;    SIND3_2 = 123;
190 SIND4_2 = 124;    SIND5_2 = 125;    SIND6_2 = 126;    SIND7_2 = 127;
191 LDCR_2 = 128;     LDCT_2 = 129;     LCA_2 = 130;      LDC_2 = 131;
192 LLA_2 = 132;      LDO_2 = 133;      LAO_2 = 134;      LDL_2 = 135;
193 LDA_2 = 136;      LOD_2 = 137;      UJP_2 = 138;      UJPL_2 = 139;
194 MPI_2 = 140;      DVI_2 = 141;      STI_2 = 142;      MODI_2 = 143;
195 CPL_2 = 144;      CPG_2 = 145;      CPI_2 = 146;      CXL_2 = 147;
196 CXG_2 = 148;      CXI_2 = 149;      RPU_2 = 150;      CPF_2 = 151;
197 LDCN_2 = 152;     LSL_2 = 153;      LDE_2 = 154;      LAE_2 = 155;
198 NGI_2 = 156;      LPR_2 = 157;      BPT_2 = 158;      NBP_2 = 159;
199 LOR_2 = 160;      LAND_2 = 161;     ADI_2 = 162;      SBI_2 = 163;
200 STL_2 = 164;      SRO_2 = 165;      STR_2 = 166;      LDB_2 = 167;
201    "NO MICROENGINE OPERATION CODES FROM 168 TO 175."
202 EQUI_2 = 176;     NEQI_2 = 177;     LEQI_2 = 178;     GEQI_2 = 179;
203 LEUSW_2 = 180;    GEUSW_2 = 181;    EQUPWR_2 = 182;   LEQPWR_2 = 183;
204 GEQPWR_2 = 184;   EQUBYT_2 = 185;   LEQBYT_2 = 186;   GEQBYT_2 = 187;
205 SRS_2 = 188;      SWAP_2 = 189;     TNC_2 = 190;      RND_2 = 191;
206 ADR_2 = 192;      SBR_2 = 193;      MPR_2 = 194;      DVR_2 = 195;
207 STO_2 = 196;      MOV_2 = 197;      DUP2_2 = 198;     ADJ_2 = 199;
208 STB_2 = 200;      LDP_2 = 201;      STP_2 = 202;      CHK_2 = 203;
209 FLT_2 = 204;      EQREAL_2 = 205;   LEQREAL_2 = 206;  GEQREAL_2 = 207;
210 LDM_2 = 208;      SPR_2 = 209;      EFJ_2 = 210;      NFJ_2 = 211;
211 FJP_2 = 212;      FJPL_2 = 213;     XJP_2 = 214;      IXA_2 = 215;
212 IXP_2 = 216;      STE_2 = 217;      INN_2 = 218;      UNI_2 = 219;
213 INT_2 = 220;      DIF_2 = 221;      SIGNAL_2 = 222;   WAIT_2 = 223;
214 ABI_2 = 224;      NGL_2 = 225;      DUP1_2 = 226;     ABR_2 = 227;
215 MOR_2 = 228;      LNOT_2 = 229;     IND_2 = 230;      INC_2 = 231;
216    "NO MICROENGINE OPERATION CODES > 231."
217 MESSAGE_2 = 232;  EOM_2 = 233;      NEWLIN_2 = 234;
218    "OPERATORS 232, 233, 234 ARE VIRTUAL OPS TO BE REMOVED IN PASS7"
219
220 "OTHER CONSTANTS"
221
222 "MACHINE REGISTERS (NOTE: POSITIVE REGS ARE FIELDS IN ACTIVE TIB)"
223 READYQ_REG  = -3;   "POINTS TO READY QUEUE"
224 SEGDICT_REG = -2;   "POINTS TO SEGMENT DICTIONARY"
225 TASK_REG    = -1;   "POINTS TO RUNNING TASK'S TIB"
226 QLINK_REG   = 0;    "LINK FIELD FOR SEMAPHORES"
227 "NOTE: REG 1 HAS 2 FIELDS"
228 PRIORITY_REG = 1;   "TASK'S CPU PRIORITY"
229 FLAG_REG    = 1;    "STATE FLAGS"
230 LSPLIM =      2;    "LOWER STACK LIMIT (HEAPTOP)"
231 USPLIM =      3;    "UPPER STACK LIMIT"
232 STACKPTR_REG = 4;   "STACK POINTER"
233 LBASE_REG =   5;    "LOCAL BASE"
234 GBASE_REG =   6;    "GLOBAL BASE"
235 PC_REG =      7;    "PROGRAM COUNTER"
236 SEG_REG =     8;    "CURRENT SEGMENT ADDRESS"
237
238 FIXLCL_LINE = 1;    "LOCAL WORD 1 IS ALWAYS LINE #"
239 FIXLCL_OLDG = 2;    "LOCAL WORD 2 IS ALWAYS OLD GLOBAL BASE"
```

```
240 FIXGBL_OLDG = 2;    "GLOBAL WORD 2 IS OLD GLOBAL BASE"
241 MSCMSIZE = 8;       "MARK STACK TAKES 8 BYTES"
242 PDP11 = TRUE;
243 CONCURRENT = TRUE;
244 SPLITLENGTH = 4 "WORDS PER REAL";
245 HALFWORD = 1;
246 ONEWORD = 2;
247 TWOWORDS = 4;   THREEWORDS = 6;  FOURWORDS = 8;  FIVEWORDS = 10;
248 STACK_LIMIT = 32667 "GREATEST INTEGER - 100";  CODE_LIMIT = 32667;
249 THIS_PASS = 6;
250 INFILE = 2;     OUTFILE = 1;     INTERFACEFILE = 4;
251 "(FILE NUMBER 3 IS USED BY PASS7 FOR FINAL CODE OUTPUT)"
252 SIBS_OFFSET = 11;   "SIBS FIELD IS 11TH IN TIB"
253 SEQL_SEG = 129;     "AT RUN TIME, SEGMENT NO. OF SEQ'L PGM IS 129"
254 INTFSEG = 128;      "AT RUN TIME, SEGMENT NO. OF INTERFACE SEGMENT"
255 STARTPROC = 1;      "START EXECUTION OF SEGMENT WITH PROC NO. 1"
256 STACK_ERROR = 1;    CODE_ERROR = 2;
257 CALLERRORDISPL = 0; "$$$"     "BYTES OF OBJECT CODE REQUIRED TO
258                                CALL A RUN-TIME ERROR ROUTINE."
259
260 "%IF SRC_32"
261 TYPE SPLIT_INTEGER = ARRAY [1..2] OF SHORTINTEGER;
262 "%END SRC_32"
263
264 "TYPES OF OUTPUT CODE PRODUCED BY THIS PASS:
265 OPERATORS, UNSIGNED BYTE OPERANDS, SIGNED BYTE OPERANDS,
266 'DON'T CARE' BYTE OPERANDS, 'BIG' OPERANDS, WORD OPERANDS,
267 OR SOME VIRTUAL CODE (NOT MICROENGINE)."
268
269 TYPE TYPEOPCODE = (OPTR, UB, SB, DB, B, W, NOTME);
270
271 VAR
272
273 LINK: PASSPTR;
274 TABLERESET: FULLWORD;
275
276 SUMMARY,        TEST,          CHECK,
277 GENERATE,       NUMBER,        AFTERBEGIN,
278 AFTERERROR,     DONE,          VARNTCHECK:     BOOLEAN; "VARNT"
279
280 GENNINGINTFAC: BOOLEAN; "TRUE IF GENERATING AN INTERFACE SEG"
281
282 JUMPTABLE, STACKTABLE, CONSTTABLE, XJPTABLE,
283     EXITICTABLE, DATASIZETABLE: TABLEPTR;
284
285 "NUMBER OF ROUTINES IN EACH INTERFACE"
286 IFSEGSIZE: ARRAY[1..MAXINTFAC] OF INTEGER;
287
288 CONSTANTS,      "COUNTS NO. OF CONSTANTS IN CONST TABLE"
289 XJPOFFSETPTR,   "COUNTS NO. OF WORDS OF XJP OFFSETS IN CONST TABLE"
290 POOLSIZE,       "TOTAL BYTES IN CONST POOL, INCLUDING XJP OFFSETS"
291 STACKLENGTH,    "AS IT COMES FROM PASS5, STACKLENGTH IS THE AMOUNT
292                 OF EXTRA STACK SPACE REQUIRED BY A ROUTINE.
293                 IT IS INCREASED BY THE AMOUNT OF STACK SPACE
294                 THE ROUTINE NEEDS FOR VARIABLES & CALCULATIONS (IN-
295                 CLUDING STACKSPACE NEEDED BY ROUTINES CALLED BY
296                 THE CURRENT ROUTINE)."
297 VARLENGTH,      "BYTES OF LOCAL VARS NEEDED FOR CURRENT ROUTINE"
298 PARAMLENGTH,    "BYTES OF PARAMS COMING INTO CURRENT ROUTINE"
299 POPLENGTH,      "BYTES TO POP AFTER CURRENT ROUTINE"
```

```
300 TEMP,         "MAX BYTES OF CALCULATION STACK FOR CURRENT RTN"
301 MAXTEMP,      "ROUTINE (PROCEDURE) LABEL FOR CURRENT ROUTINE"
302 BLOCK,        "BYTE LOCATION COUNTER, RELATIVE TO START OF SEGMENT,
303 LOCATION,      & POINTS TO NEXT BYTE TO BE GENERATED."
304
305 IFSEGNUM,     "SEQUENTIAL NUMBER OF THE LAST INTERFACE SEG GEN'D"
306 LINE,         "SOURCE LINE NUMBER"
307 OP,           "VIRTUAL MACHINE INSTRUCTION OPERATOR, FROM PASS5"
308 ARG1, ARG2, ARG3, ARG4, ARG5   "VIRTUAL MACHINE INSTRUC. OPERANDS"
309     :INTEGER;
310
311 "##########################"
312 "COMMON TEST OUTPUT MECHANISM"
313 "##########################"
314
315 PRINTED: INTEGER;
316
317 OK: BOOLEAN;
318 "PASS1 TO 6:   OK = NOT DISK OVERFLOW
319  PASS7:       OK = NOT DISK OVERFLOW & PROGRAM CORRECT"
320
321 PAGE_IN: PAGE;   PAGES_IN, WORDS_IN: INTEGER;
322 PAGE_OUT: PAGE;  PAGES_OUT, WORDS_OUT: INTEGER;
323 IFPAGE_OUT: PAGE; IFPAGES_OUT, IFWORDS_OUT: INTEGER;
324
325 PROCEDURE PRINT_TEXT (TEXT: TEXT_TYPE);
326 VAR I: INTEGER;
327 BEGIN
328 WRITE(EOL);
329 FOR I:= 1 TO TEXT_LENGTH DO WRITE(TEXT(.I.));
330 WRITE(EOL)
331 END;
332
333 PROCEDURE FILE_LIMIT;
334 BEGIN
335 PRINT_TEXT('PASS 6: FILE_LIMIT');
336 OK:= FALSE
337 END;
338
339 PROCEDURE INIT_PASS (VAR LINK: PASSPTR);
340 BEGIN
341 LINK:= PARAM(.2.).PTR;
342 OK:= TRUE;
343 PAGES_IN:= 1; WORDS_IN:= PAGELENGTH;
344 PAGES_OUT:= 1; WORDS_OUT:= 0;
345 IFPAGES_OUT := 1; IFWORDS_OUT := 0;
346 END;
347
348 PROCEDURE NEXT_PASS (LINK: PASSPTR);
349 BEGIN
350 IF WORDS_OUT > 0
351 THEN IF PAGES_OUT > FILE_LENGTH(OUTFILE)
352      THEN FILE_LIMIT
353      ELSE PUT(OUTFILE, PAGES_OUT, PAGE_OUT);
354 IF IFWORDS_OUT > 0
355 THEN IF IFPAGES_OUT > FILE_LENGTH(INTERFACEFILE)
356      THEN FILE_LIMIT
357      ELSE PUT(INTERFACEFILE, IFPAGES_OUT, IFPAGE_OUT);
358 WITH PARAM(.1.) DO BEGIN
```

```
360      TAG:= BOOLTYPE; BOOL:=OK END;
361      WITH PARAM(.2.) DO BEGIN
362      TAG:= PTRTYPE; PTR:= LINK END;
363      WITH PARAM(.4.) DO BEGIN
364      TAG:= INTTYPE; INT:= PAGES_OUT END;
365  END;
366
367  PROCEDURE READ_IFL (VAR I: INTEGER);
368  BEGIN
369      IF WORDS_IN = PAGELENGTH THEN BEGIN
370          IF PAGES_IN > FILE_LENGTH(INFILE) THEN FILE_LIMIT
371          ELSE BEGIN
372              GET(INFILE, PAGES_IN, PAGE_IN);
373              PAGES_IN:= SUCC(PAGES_IN)
374          END;
375          WORDS_IN:= 0
376      END;
377      WORDS_IN:= SUCC(WORDS_IN);
378      I:= PAGE_IN(.WORDS_IN.)
379  END;
380
381  PROCEDURE WRITE_IFL (I: INTEGER);
382  BEGIN
383      IF GENNINGINTFAC
384      THEN BEGIN
385          IFWORDS_OUT := SUCC(IFWORDS_OUT);
386          IFPAGE_OUT[IFWORDS_OUT] := I;
387          IF IFWORDS_OUT = PAGELENGTH
388          THEN BEGIN
389              IF IFPAGES_OUT > FILE_LENGTH(INTERFACEFILE)
390              THEN FILE_LIMIT
391              ELSE BEGIN
392                  PUT (INTERFACEFILE, IFPAGES_OUT, IFPAGE_OUT);
393                  IFPAGES_OUT := SUCC(IFPAGES_OUT)
394              END;
395              IFWORDS_OUT := 0
396          END
397      END
398      ELSE BEGIN
399          WORDS_OUT:= SUCC(WORDS_OUT);
400          PAGE_OUT(.WORDS_OUT.):= I;
401          IF WORDS_OUT = PAGELENGTH
402          THEN BEGIN
403              IF PAGES_OUT > FILE_LENGTH(OUTFILE)
404              THEN FILE_LIMIT
405              ELSE BEGIN
406                  PUT(OUTFILE, PAGES_OUT, PAGE_OUT);
407                  PAGES_OUT:= SUCC(PAGES_OUT)
408              END;
409              WORDS_OUT:= 0
410          END
411      END
412  END;
413
414  PROCEDURE PRINTABS(ARG: INTEGER)
415  VAR I: ARRAY (.1..MAXDIGIT.) OF CHAR; REM, DIGIT, I: INTEGER;
416  BEGIN
417      REM:= ARG; DIGIT:= 0;
418      REPEAT
419          DIGIT:= DIGIT + 1;
```

```
420      T(.DIGIT.):= CHR(ABS(REM MOD 10) + ORD('0'));
421      REM:= REM DIV 10;
422    UNTIL (REM=0) OR (DIGIT>MAXDIGIT);
423    FOR I:= DIGIT DOWNTO 1 DO WRITE(T(.I.));
424    FOR I:= 1 TO MAXDIGIT DO WRITE(' ');
425  END;
426
427  PROCEDURE PRINTEOL;
428  BEGIN WRITE(EOL); PRINTED:= 0 END;
429
430  PROCEDURE PRINTFF;
431  VAR I:INTEGER;
432  BEGIN
433    PRINTEOL; FOR I:=1 TO 130 DO WRITE('6'); PRINTEOL
434  END;
435
436  PROCEDURE PRINTOP(OP: INTEGER);
437  BEGIN
438    IF PRINTED = PRINTLIMIT THEN PRINTEOL.;
439    WRITE('C'); PRINTABS(OP);
440    PRINTED:= PRINTED + 1;
441  END;
442
443  PROCEDURE PRINTARG(ARG: INTEGER);
444  BEGIN
445    IF PRINTED = PRINTLIMIT THEN PRINTEOL;
446    IF ARG < 0 THEN WRITE('-') ELSE WRITE(' ');
447    PRINTABS(ARG);
448    PRINTED:= PRINTED + 1;
449  END;
450  "##############"
451  "INPUT PROCEDURES"
452  "INPUT PROCEDURES"
453  "##############"
454
455  PROCEDURE READ1ARG;
456  BEGIN READ_IFL(ARG1) END;
457
458  PROCEDURE READ2ARG;
459  BEGIN READ_IFL(ARG1); READ_IFL(ARG2) END;
460
461  PROCEDURE READ3ARG;
462  BEGIN READ_IFL(ARG1); READ_IFL(ARG2); READ_IFL(ARG3) END;
463
464  PROCEDURE READ4ARG;
465  BEGIN
466    READ_IFL(ARG1); READ_IFL(ARG2);
467    READ_IFL(ARG3); READ_IFL(ARG4);
468  END;
469
470  PROCEDURE READ5ARG;
471  BEGIN
472    READ_IFL(ARG1); READ_IFL(ARG2); READ_IFL(ARG3);
473    READ_IFL(ARG4); READ_IFL(ARG5)
474  END;
475
476  "##############"
477  "OUTPUT PROCEDURES"
478  "##############"
479
```

```
480  PROCEDURE ERROR (PASS, NUMBER: INTEGER); FORWARD;
481
482  PROCEDURE UPDLOC (CODE: TYPEOFCODE; VALUE: INTEGER);
483  BEGIN  "UPDATE THE OUTPUT CODE LOCATION COUNTER."
484  IF LOCATION < CODE_LIMIT
485  THEN CASE CODE OF
486      OPTR, UB, SB, DB: LOCATION := LOCATION + HALFWORD;
487      B: IF VALUE <= 127
488         THEN LOCATION := LOCATION + HALFWORD
489         ELSE LOCATION := LOCATION + ONEWORD;
490      W: LOCATION := LOCATION + ONEWORD;
491      NOTME;
492      END
493  ELSE BEGIN
494      ERROR (THIS_PASS, CODE_ERROR);
495      LOCATION := 0;
496      END;
497  END;
498
499  "LOCATION COUNTER IS KEPT FOR THE CONCURRENT SEGMENT ONLY.
C500  IF CODE OUTPUT ROUTINES ARE CALLED WHILE GENERATING AN INTERFACE
C501  SEGMENT (GENNINGINTFAC = TRUE) THEN DON'T FOOL WITH CONCURRENT CODE
C502  LOCATION COUNTER."
503
504  PROCEDURE WRITE1(OP: INTEGER);
505  BEGIN
506  IF TEST THEN PRINTOP(OP);
507  WRITE_IFL(OP);
508  IF NOT GENNINGINTFAC THEN UPDLOC (OPTR, OP)
509  END;
510
511  PROCEDURE WRITE2(OP, ARG: INTEGER; ARGTYP: TYPEOFCODE);
512  BEGIN
513  IF TEST
514  THEN BEGIN
515      PRINTOP(OP); PRINTARG(ARG)
516      END;
517  WRITE_IFL(OP); WRITE_IFL(ARG);
518  IF NOT GENNINGINTFAC
519  THEN BEGIN
520      UPDLOC(OPTR, OP); UPDLOC(ARGTYP, ARG)
521      END;
522  END;
523
524  PROCEDURE WRITE3(OP, ARG1, ARG2: INTEGER;
525                   ARG1TYP, ARG2TYP: TYPEOFCODE);
526  BEGIN
527  IF TEST
528  THEN BEGIN
529      PRINTOP(OP);
530      PRINTARG(ARG1); PRINTARG(ARG2);
531      END;
532  WRITE_IFL(OP);
533  WRITE_IFL(ARG1); WRITE_IFL(ARG2);
534  IF NOT GENNINGINTFAC
535  THEN BEGIN
536      UPDLOC(OPTR, OP);
537      UPDLOC(ARG1TYP, ARG1); UPDLOC(ARG2TYP, ARG2)
538      END
539  END;
```

```
540  PROCEDURE WRITE4(OP, ARG1, ARG2, ARG3: INTEGER;
541                   ARG1TYP, ARG2TYP, ARG3TYP: TYPEOFCODE);
542
543  BEGIN
544    IF TEST THEN
545    BEGIN
546      PRINTOP(OP); PRINTARG(ARG1);
547      PRINTARG(ARG2); PRINTARG(ARG3);
548    END;
549    WRITE_IFL(OP); WRITE_IFL(ARG1);
550    WRITE_IFL(ARG2); WRITE_IFL(ARG3);
551    IF NOT GENNINGINTFAC
552    THEN BEGIN
553      UPDLOC(OPTR, OP); UPDLOC(ARG1TYP, ARG1);
554      UPDLOC(ARG2TYP, ARG2); UPDLOC(ARG3TYP, ARG3)
555      END;
556  END;
557
558  PROCEDURE WRITE5(OP, ARG1, ARG2, ARG3, ARG4: INTEGER;
559                   ARG1TYP, ARG2TYP, ARG3TYP, ARG4TYP: TYPEOFCODE);
560  BEGIN
561    IF TEST
562    THEN BEGIN
563      PRINTOP(OP);
564      PRINTARG(ARG1); PRINTARG(ARG2);
565      PRINTARG(ARG3); PRINTARG(ARG4);
566    END;
567    WRITE_IFL(OP);
568    WRITE_IFL(ARG1); WRITE_IFL(ARG2);
569    WRITE_IFL(ARG3); WRITE_IFL(ARG4);
570    IF NOT GENNINGINTFAC
571    THEN BEGIN
572      UPDLOC(OPTR, OP);
573      UPDLOC(ARG1TYP, ARG1); UPDLOC(ARG2TYP, ARG2);
574      UPDLOC(ARG3TYP, ARG3); UPDLOC(ARG4TYP, ARG4)
575      END
576  END;
577
578  PROCEDURE WRITEARG(ARG: INTEGER; ARGTYP: TYPEOFCODE);
579  BEGIN
580    IF TEST THEN PRINTARG(ARG);
581    WRITE_IFL(ARG);
582    IF NOT GENNINGINTFAC THEN UPDLOC(ARGTYP, ARG)
583  END;
584
585  PROCEDURE WRITELOCATION;
586  BEGIN
587    IF TEST THEN PRINTARG(LOCATION);
588    WRITE_IFL(LOCATION);
589  END;
590
591  PROCEDURE ERROR;
592  BEGIN
593    IF NOT AFTERERROR
594    THEN BEGIN
595      AFTERERROR:= TRUE;
596      WRITEARG(MESSAGE_2, NOTME);
597      WRITEARG(PASS, NOTME);
598      WRITEARG(NUMBER, NOTME);
599      WRITEARG(LINE, NOTME);
```

```
600         GENERATE:= FALSE
601             END
602 END;
603
604 "////////////////////////////"
605 "PSUEDO RUN-TIME STACK PROCEDURES"
606 "////////////////////////////"
607
608 PROCEDURE PUSHWORD;
609 BEGIN
610 IF TEMP < STACK_LIMIT THEN TEMP:= TEMP + WORDLENGTH
611     ELSE ERROR(THIS_PASS, STACK_ERROR);
612     IF TEMP > MAXTEMP THEN MAXTEMP:= TEMP;
613 END;
614
615 PROCEDURE POPWORD;
616 BEGIN TEMP:= TEMP - WORDLENGTH ;
617 END;
618
619 PROCEDURE PUSHREAL;
620 BEGIN
621 IF TEMP < STACK_LIMIT THEN TEMP:= TEMP + REALLENGTH
622     ELSE ERROR(THIS_PASS, STACK_ERROR);
623     IF TEMP > MAXTEMP THEN MAXTEMP:= TEMP;
624 END;
625
626 PROCEDURE POPREAL;
627 BEGIN TEMP:= TEMP - REALLENGTH END;
628
629 PROCEDURE PUSHSET;
630 BEGIN
631 IF TEMP < STACK_LIMIT THEN TEMP:= TEMP + SETLENGTH
632     ELSE ERROR(THIS_PASS, STACK_ERROR);
633     IF TEMP > MAXTEMP THEN MAXTEMP:= TEMP;
634 END;
635
636 PROCEDURE POPSET;
637 BEGIN TEMP:= TEMP - SETLENGTH END;
638
639 PROCEDURE PUSH(LENGTH: INTEGER);
640 BEGIN
641 IF TEMP < STACK_LIMIT - LENGTH THEN TEMP:= TEMP + LENGTH
642     ELSE ERROR(THIS_PASS, STACK_ERROR);
643     IF TEMP > MAXTEMP THEN MAXTEMP:= TEMP;
644 END;
645
646 PROCEDURE POP(LENGTH: INTEGER);
647 BEGIN TEMP:= TEMP - LENGTH END;
648
649 "////////////////////////////"
650 " SHORT CONSTANT PUSH "
651 "////////////////////////////"
652
653 PROCEDURE PICK_PUSHCONST (VALUE: INTEGER);
654 BEGIN    "PICK 1 OF 3 OPERATORS WHICH PUSH IMMEDIATE CONSTANTS"
655 IF (VALUE >= 0) AND (VALUE <= 31)
656 THEN WRITE1(VALUE)    "THE OP IS AN SLDCXX_02"
657 ELSE IF (VALUE >= 32) AND (VALUE <= 255)
658     THEN WRITE2(LDCB_2, VALUE, UB)
659     ELSE WRITE2(LDCI_2, VALUE, W)
```

```
660 END;
661 "################"
662 "VARIABLE PROCEDURES"
663 "################"
664 "################"
665
666 FUNCTION DISPL(MODE, ARG: INTEGER): INTEGER;
667 "HARTMANN'S VARIABLES HAVE NEGATIVE DISPLACEMENT, PARAMETERS HAVE
668  POSITIVE DISPLACEMENTS, BOTH IN BYTES & 2-BASED.
669  FOR MICROENGINE, BOTH VARIABLES AND PARAMETERS HAVE
670  POSITIVE WORD DISPLACEMENTS (1-BASED). VARS HAVE SMALLER DISPLS
671  SINCE THEY ARE CLOSER TO THE MSCW (PUSHED LATER) THAN PARMS."
672 BEGIN
673 CASE MODE OF
674 MODE0: "CONSTANT"
675     DISPL := (ARG + ONEWORD) DIV WORDLENGTH;
676
677 "PROCEDURE, PROGRAM, CLASS ENTRY, MONITOR ENTRY, MGR ENTRY"
678 MODE1, MODE2, MODE4, MODE5, MODE12:
679     IF ARG < 0
680     THEN DISPL := (-ARG + TWOWORDS) DIV WORDLENGTH
681     ELSE DISPL := (ARG + VARLENGTH + TWOWORDS) DIV WORDLENGTH;
682
683 MODE3: "PROCESS ENTRY"
684 "INTERFACE PROCEDURE HAS NO LOCAL VARS (NOT EVEN LINE & OLD
685  GLOBAL BASE). HOWEVER, IT DOES HAVE AN EXTRA PARM (THE NO.
686  OF THE PREFIX PROCEDURE BEING CALLED) WHICH HAS NOT BEEN
687  TAKEN INTO ACCOUNT IN DISPLACEMENT CALCULATION TO THIS POINT.
688  IF ARG < 0 THEN REFERENCE IS TO A LOCAL VARIABLE IN THE ENTRY
689  PROCEDURE. OTHERWISE, REFERENCE IS TO PARM IN THE INTERFACE
690  PROCEDURE."
691     IF ARG < 0
692     THEN DISPL := (-ARG + TWOWORDS) DIV WORDLENGTH
693     ELSE DISPL := (ARG + ONEWORD) DIV WORDLENGTH;
694
695 "PROCESS, CLASS, MONITOR, MANAGER MODES"
696 MODE6, MODE7, MODE8, MODE11:
697 "IN THESE MODES, REFERENCE IS TO 'SHARED' VARIABLES OR PARMS
698  TO THE INITIAL STATEMENT. THE RECORD CONTAINING THESE VARS
699  & PARMS IS STORED IN THE SHARED VARIABLE AREA OF THE INITIAL
700  PROCESS. SPACE FOR THE RECORD WAS ALLOCATED IN PRIOR PASSES.
701  THE GLOBAL RECORDS DO NOT CONTAIN TWO WORDS FOR LINE NO. &
702  OLD GLOBAL BASE."
703     IF ARG < 0
704     THEN DISPL := (-ARG) DIV WORDLENGTH
705     ELSE DISPL := (ARG + VARLENGTH) DIV WORDLENGTH;
706
707 MODE10: "UNDEFINED MODE"
708 END
709 END;
710
711 PROCEDURE PUSHVALUE(MODE, ARG: INTEGER);
712 BEGIN
713 CASE MODE OF
714 "PROCEDURE, CLASS ENTRY, MONITOR ENTRY, MGR ENTRY MODES"
715 MODE1, MODE4, MODE5, MODE12:          "WIR""
716     WRITE2(LDL_2, DISPL(MODE, ARG), B);
717
718 MODE2:  "PROGRAM MODE"
719
```

```
720              WRITE2(LDO_2, DISPL(MODE, ARG), B);                                   511   192   666   711   711   269
721
722        MODE3:   "PROCESS ENTRY MODE"                                               142
723            IF ARG < 0                                                              711
724            THEN WRITE2(LDN_2, DISPL(MODE, ARG), B)                                 511   192   666   711   711   269
725            ELSE WRITE3(LOD_2, 1, DISPL(MODE, ARG), DB, B);                         524   193   666   711   711   269
726
727        "PROCESS, CLASS, MONITOR, MANAGER MODES"
728        MODE6, MODE7, MODE8, MODE11:          "MGR"                                 145   146   147   150
729              WRITE2(LDO_2, DISPL(MODE, ARG), B);                                   511   192   666   711   711   269
730
731        MODE10:      "UNDEFINED MODE"                                               149
732        END;
733        PUSHWORD;                                                                   608
734      END;
735      PROCEDURE PUSHADDRESS(MODE, ARG: INTEGER);
736      BEGIN                                                                          *     *     *     S
737        CASE MODE OF                                                                736
738        MODE0:   "CONSTANT MODE"                                                    139
739              WRITE2(LCA_2, DISPL(MODE, ARG), B);                                   511   191   666   736   736   269
740
741        "PROCEDURE, CLASS ENTRY, MON. ENTRY, MGR. ENTRY"
742        MODE1, MODE4, MODE5, MODE12: "MGR"                                          140   143   144   151
743              WRITE2(LLA_2, DISPL(MODE, ARG), B);                                   511   192   666   736   736   269
744
745        MODE2:   "PROGRAM MODE"                                                     141
746              WRITE2(LAO_2, DISPL(MODE, ARG), B);                                   511   192   666   736   736   269
747
748        MODE3: "PROCESS ENTRY MODE"                                                 142
749            IF ARG < 0                                                              736
750            THEN WRITE2(LLA_2, DISPL(MODE, ARG), B)                                 511   192   666   736   736   269
751            ELSE WRITE3(LDA_2, 1, DISPL(MODE, ARG), DB, B);                         524   193   666   736   736   269
752
753        "PROCESS, CLASS, MONITOR, MANAGER MODES"
754        MODE6, MODE7, MODE8, MODE11: "MGR"                                          145   146   147   150
755              WRITE2(LAO_2, DISPL(MODE, ARG), B);                                   511   192   666   736   736   269
756
757        MODE10:      "UNDEFINED MODE"                                               149
758        END;
759        PUSHWORD;                                                                   608
760      END;
761      PROCEDURE PUSHINDIRECT(VARTYPE: INTEGER);
762      BEGIN                                                                          *     *     S
763        CASE VARTYPE OF                                                             763
764        BYTETYPE:
765          BEGIN   WRITE1(LDB_2);   POPWORD   END;                                   504   200   615
766        WORDTYPE:                                                                   135
767          WRITE1(STNDO_2);                                                          504   189
768        REALTYPE:                                                                   135
769          BEGIN WRITE2(LDN_2, REALLENGTH DIV WORDLENGTH, UB);                       511   210   29   28   269
770          POPWORD; PUSHREAL;                                                        615   619
771          END;
772        SETTYPE:                                                                    135
773          BEGIN WRITE2(LDN_2, SETLENGTH DIV WORDLENGTH, UB);                        511   210   30   28   269
774          POPWORD; PUSHSET;                                                         615   629
775          PICK_PUSHCONST(SETLENGTH DIV WORDLENGTH);                                 653   30   28
776          PUSHWORD;    "PUSH WORD LENGTH OF SET"                                    608
777          END
```

```
780  END;
781  END;
782
783
784
785  "##########"
786  "TABLE PROCEDURES"
787  "##########"
788
789  PROCEDURE ALLOCATE(VAR T: TABLEPTR; ENTRIES: INTEGER);
790  VAR PORTION: TABLEPTR; I: INTEGER;
791  BEGIN
792     NEW(T); PORTION:= T;
793     I:= ENTRIES - MAXWORD;
794     WHILE I > 0 DO
795     WITH PORTION@ DO
796     BEGIN
797        NEW(NEXTPORTION); PORTION:= NEXTPORTION;
798        I:= I - MAXWORD;
799     END;
800  END;
801
802  PROCEDURE ENTER(T: TABLEPTR; I, J: INTEGER);
803  VAR PORTION: TABLEPTR; K: INTEGER;
804  BEGIN
805     PORTION:= T; K:= I;
806     WHILE K > MAXWORD DO
807     BEGIN
808        PORTION:= PORTION@.NEXTPORTION;
809        K:= K - MAXWORD;
810     END;
811     PORTION@.CONTENTS(.K.):= J;
812  END;
813
814  FUNCTION ENTR(T: TABLEPTR; I: INTEGER): INTEGER;
815  VAR PORTION: TABLEPTR; J: INTEGER;
816  BEGIN
817     PORTION:= T; J:= I;
818     WHILE J > MAXWORD DO
819     BEGIN
820        PORTION:= PORTION@.NEXTPORTION;
821        J:= J - MAXWORD;
822     END;
823     ENTR:= PORTION@.CONTENTS(.J.);
824  END;
825
826  "##########"
827  "INITIALIZATION AND TERMINATION PROCEDURES"
828  "##########"
829
830  PROCEDURE BEGINPASS;
831  VAR I: INTEGER;
832  BEGIN
833     WITH LINK@ DO
834     BEGIN
835        SUMMARY:= SUMMARYOPTION IN OPTIONS;
836        TEST:= TESTOPTION IN OPTIONS;
837        CHECK:= CHECKOPTION IN OPTIONS;
838        VARNTCHECK:= VARNTCHECKOPTION IN OPTIONS;
839        NUMBER:= NUMBEROPTION IN OPTIONS;
```

```
840   GENERATE:= TRUE;
841   FOR I := 1 TO MAXINTFAC DO
842     IFSEGSIZE[I] := INTERFACE@.INTERFACESIZES[I];
843   ALLOCATE(JUMPTABLE, LABELS);
844   ALLOCATE(CONSTABLE, CONSTANTS DIV WORDLENGTH);
845   ALLOCATE(XJPTABLE, XJP_OFFSETS DIV WORDLENGTH);
846   ALLOCATE(EXITICTABLE, BLOCKS);
847   ALLOCATE(DATASIZETABLE, BLOCKS);
848   MARK (TABLERESET);
849   ALLOCATE(STACKTABLE, BLOCKS);
850   END;
851   CONSTANTS:= 0;
852   XJPOFFSETPTR := 0;
853   POOLSIZE := LINK@.CONSTANTS + LINK@.XJP_OFFSETS;
854   LINE := 0;
855   LOCATION := ONEWORD + POOLSIZE;
856   AFTERBEGIN := FALSE;
857   IFSEGINUM := 0;
858   GRNNINGINTFAC := FALSE;
859   "LABEL 1 (TARGET OF 1ST (DISCARDED) JUMP) IS NEVER RESOLVED"
860   ENTER(JUMPTABLE, STARTPROC, -1);
861   IF TEST THEN PRINTFF;
862   END;
863
864   PROCEDURE ENDPASS;
865   BEGIN
866     LINK@.CONSTANTS := CONSTANTS;
867     LINK@.XJP_OFFSETS := XJPOFFSETPTR;
868     WITH LINK@ DO
869       IF GENERATE THEN OPTIONS:= OPTIONS OR (.CODEOPTION.);
870     RELEASE (TABLERESET);  "REMOVE TABLES NOT NEEDED IN PASS7"
871     NEW(LINK@.TABLES)
872     WITH LINK@ DO
873       BEGIN
874       "SEGDISTANCE IS LOC. CTR. + SPACE FOR PROCEDURE DICTIONARY.
875       "LOCATION' WILL THEN POINT TO LO ORDER BYTE IN THE
876       (LAST) WORD WHICH WILL CONTAIN NO. OF PROCS. IN SEG. &
877       SEGMENT NUMBER."
878       TABLES@.SEGDISTANCE := LOCATION + 2*LINK@.BLOCKS;
879       TABLES@.JUMPTABLE:=JUMPTABLE;
880       TABLES@.CONSTTABLE:=CONSTTABLE;
881       TABLES@.XJPTABLE := XJPTABLE;
882       TABLES@.EXITICTABLE := EXITICTABLE;
883       TABLES@.DATASIZETABLE := DATASIZETABLE;
884       END
885   END;
886   "%IF SRC_32"
887   PROCEDURE SPLIT_CONST(WORD:UNIV SPLIT_INTEGER;
888               VAR HIWORD,LOWORD: INTEGER);
889   BEGIN "SPLIT A 32-BIT INTEGER INTO TWO 16-BIT INTEGERS"
890     HIWORD:=WORD[1]; LOWORD:=WORD[2];
891   END;
892   "%END"
893
894
895   PROCEDURE GEN SAVELINE SAVEBASE;
896   "PROLOGUE FOR ROUTINES"
897   BEGIN
898   IF NUMBER
899   THEN BEGIN
```

```
900      WRITEARG (HEXLIN_2, NOTHE); "CRUTCH FOR MNEMONICS PGM."
901      PICK_PUSHCONST (LINE);            PUSHWORD;
902      WRITE2 (STL_2, FIXLCL_LINE, B);   POPWORD
903      END;
904    PICK_PUSHCONST (GBASE_REG);        PUSHWORD;
905      WRITE1 (LPR_2);                   PUSHWORD;
906      WRITE2 (STL_2, FIXLCL_OLDG, B);   POP(TWOWORDS)
907      END;
908
909    PROCEDURE GEN_EXIT;
910    "EPILOGUE FOR ROUTINES"
911    BEGIN
912      PICK_PUSHCONST (GBASE_REG);       PUSHWORD;
913      WRITE1 (SLDL02_2);                PUSHWORD;
914      WRITE1 (SPR_2);                   POP(TWOWORDS);
915      WRITE2 (RPM_2, POPLENGTH DIV WORDLENGTH, B);
916      POP(POPLENGTH + MSCWSIZE)
917    END;
918
919    "#####################################"
920  C # GENERATE AN INTERFACE SEGMENT #
921    "#####################################"
922  C "#####################################"
923
924    PROCEDURE GEN_INTERFACE;
925    TYPE
926      TOSPTR = @STKENTRY;
927      STKENTRY = RECORD
928        DATA: INTEGER;
929        NXTENTRY: TOSPTR  "POINTS TO THING PUSHED JUST BEFORE"
930        END;
931    CONST
932      CONSEG = 1; "CONCURRENT SEGMENT IS SEGMENT 1"
933      HIORDER = 256; "SHIFTS SMALL INT. TO HI-ORDER BYTE IN 16-BIT WORD"
934      MININDEX = 1; "MIN INDEX FOR CASE STMT. I.E., PREFIX RTNS ARE
935                     NUMBERED STARTING AT 1"
936      DATASIZE = 0; "ROUTINE BEING GEN'D HAS NO LOCAL VAR SPACE"
937      PARMWORDS = 1; "ROUTINE BEING GEN'D NEEDS 1 PARAMETER"
938      XJPTABLE = 1; "TABLE FOR CASEJUMP INSTRUC STARTS IN WORD 1 OF SEG."
939      PROCSSEND = 1; "INTERFACE SEGS CONSIST OF ONLY 1 PROCEDURE"
940    VAR
941      OLDHEAPTOP: FULLWORD;
942      RTNS, "NO. OF ROUTINE LABELS IN THIS INTERFACE"
943      KASE, "LOOP VAR FOR RTNS & CASES IN CASE STMT"
944      IFLABEL: INTEGER; "INTERFACE ROUTINE LABEL"
945      TOS: TOSPTR; "POINTS TO CURRENT TOP OF STACK".
946      NEWTOS: TOSPTR; "POINTS TO THING TO BE PUSHED NEXT"
947    BEGIN
948      GENNINGINTFAC := TRUE; "FROM NOW TO END OF THIS PROC, SEND
                                 OUTPUT TO THE FILE OF INTERFACE SEGS"
949
950      IFSEGNUM := SUCC(IFSEGNUM);
951      RTNS := IFSEGSIZE[IFSEGNUM]; "GET NO. OF ROUTINES IN THIS INTFACE."
952
953      WRITEARG ( (RTNS*8+22) DIV 2, W); "WORDS IN SEG, LESS 1"
954
955      "GEN CONSTANT BLOCK. IT CONSISTS ONLY OF THE TABLE FOR CASE STMT."
956      WRITEARG (MININDEX, W); "MIN INDEX OF CASE STMT"
957      WRITEARG (RTNS, W); "MAX INDEX OF CASE STMT"
958      FOR KASE := 0 TO RTNS-1 DO WRITEARG (KASE*6+A, W); "OFFSETS"
959
```

```
960        "GENERATE EXIT-IC FIELD. SEGMENT-RELATIVE POINTER, IN BYTES."
961        WRITEARG (RTNS*8+16, W);
962        WRITEARG (DATASIZE, W); "GEN DATASIZE FIELD"
963
964      C "GEN CODE TO USE THE NO. OF A PREFIX ROUTINE (THE CALLING
965      C SEQUENTIAL PROGRAM MUST SUPPLY THIS AS A PARAMETER) AS THE INDEX
966      C INTO CASE STMNT AND CALL CORRESPONDING CONCURRENT PROCESS
967        ENTRY ROUTINE."
968        WRITE1 (SLDL01_2); "PUSH THE PARAMETER"
969        WRITE2 (XJP_2, XJPTABLE, D); "CASE JUMP"
970        WRITE1 (NOP_2); "GEN'D IN IMITATION OF WESTERN DIGITAL"
971        WRITE2 (UJPL_2, 6*RTNS, W); "FOR THE CASE WHERE PARM IS OUT OF RANGE"
972
973        "THE INTERFACE ROUTINE LABELS COME INTO THIS PASS IN REVERSE ORDER.
974        THEY NEED TO BE IN FORWARD ORDER. STACK THEM AND POP THEM TO
975        MAKE THE CORRECTION."
976
977        MARK (OLDHEAPTOP);
978        TOS := NIL; "THERE IS NO TOP OF STACK YET"
979        NEW (NEWTOS);
980        FOR KASE := 1 TO RTNS DO    "READ & STACK THE LABELS"
981          BEGIN
982          READ_IFL (IFLABEL);
983          WITH NEWTOS@ DO
984            BEGIN
985            DATA :=IFLABEL; "SAVE THE LABEL IN A STACK NODE"
986            NXTENTRY := TOS "TIE NODE TO EXISTING STACK (PUSH NODE)"
987            END;
988          TOS := NEWTOS; "KEEP TOS POINTER CURRENT"
989          NEW (NEWTOS);
990      C "IF THE LABEL JUST HANDLED WAS NOT THE LAST ONE IN THIS INTERFACE,
991      C THEN READ PAST THE 'PUSHLABEL' OPERATOR WHICH PRECEDES THE
992      C NEXT LABEL. IF THIS WAS THE LAST LABEL, DON'T READ THE NEXT
993      C OPERATOR, SINCE IT NEEDS TO BE HANDLED BY THE LARGE CASE
994      C STATEMENT IN THE 'SCAN' ROUTINE."
995          IF KASE <> RTNS
996          THEN READ_IFL (OP)
997          END;
998
999        FOR KASE := 1 TO RTNS DO    "GEN CODE FOR EACH CASE"
1000         BEGIN
1001     C "GEN THE CALL TO THE ENTRY ROUTINE IN THE CONCURRENT PROCESS.
1002     C MUST BE CALL EXTERNAL <LOCAL> SO THAT STATIC LINK IN PROCESS
1003     C ENTRY PROCEDURE POINTS TO MARKSTACK FOR THIS (INTERFACE)
1004     C PROCEDURE. WHEN REFERENCING PARMS, THE PROCESS ENTRY ROUTINE
1005     C WILL USE <INTERMEDIATE> LOADS AND STORES. THIS ALLOWS THE
1006     C ENTRY ROUTINE TO REFER TO PARMS IN THE INTERFACE PROCEDURE AS IF
1007       THEY HAD BEEN PASSED DIRECTLY TO THE ENTRY ROUTINE."
1008         WRITE3 (CXL_2, CONSEG, TOS@.DATA, UB, UB);
1009         WRITE2 (UJPL_2, 6*(RTNS-KASE), W); "GEN JUMP OUT OF CASE CODE"
1010         TOS := TOS@.NXTENTRY "POP THE STACK OF LABELS"
1011         END;
1012
1013       WRITE1 (NOP_2); "TO WORD-ALIGN"
1014       WRITE2 (RPU_2, DATASIZE+PARMWORDS, B);
1015       RELEASE (OLDHEAPTOP);
1016
1017       WRITEARG (RTNS+3, W); "SEG-REL WORD ADDR OF DATASIZE FIELD"
1018
1019     C "GEN A WORD IN WHICH THE HIGH ORDER BYTE CONTAINS THE NUMBER
```

```
C 1020      OF PROCEDURES IN THE SEGMENT JUST GEN'D, & THE LOW ORDER BYTE
C 1021      CONTAINS THE SEG NUMBER OF SEG JUST GEN'D. ASSUME THAT THE
C 1022      KERNEL IS SEGMENT 0, CONCURRENT PROGRAM IS SEGMENT 1, AND THAT
C 1023      INTERFACE SEGS ARE NUMBERED CONSECUTIVELY AFTER THAT. THEY WILL
C 1024      BE RENUMBERED BY THE KERNEL WHEN LOADED, BUT THESE NUMBERS
C 1025      SHOULD HELP IN TESTING SINCE MOST, DIG. SOFTWARE DOESN'T SUPPORT
C 1026      MORE THAN 16 SEGMENTS."
  1027   570 939 933 305 932 269   WRITEARG (PROCSGEND * WORDER + (IFSEGNUM + CONSEG), W);
  1028
  1029   280 337                   GENNINGINTFAC := FALSE "SEND OUTPUT TO CONCURRENT FILE"
  1030                             END;
  1031
  1032
  1033
  1034                             "########"
  1035                             "OPERATORS"
  1036                             "########"
  1037
  1038   •                         PROCEDURE SCAN;
  1039   •                         VAR ARG1LO, ARG1HI: INTEGER;
  1040       • s                   BEGIN
  1041
  1042   367 558 455               READ_IFL (OP);   READ1ARG; "DISCARD JUMP TO INITIAL BLOCK"
  1043
  1044   278 337                   DONE:=FALSE; "KEEP READING & TRANSLATING CODE UNTIL END-OF-MEDIUM."
  1045                             REPEAT
  1046   367 558                     READ_IFL(OP);
  1047   550                         CASE OP OF
  1048
  1049   120                       PUSHCONST1"(VALUE)":
  1050                             BEGIN
  1051   455 653 558 608           READ1ARG;  PICK_PUSHCONST(ARG1);  PUSHWORD;
  1052                             END;
  1053
  1054   120                       PUSHVAR1"(TYPE, MODE, DISPL)":
  1055   461                       BEGIN READ3ARG;
  1056   558 135                     IF ARG1 = WORDTYPE
  1057   711 558                       THEN PUSHVALUE(ARG2, ARG3)
  1058                             ELSE BEGIN
  1059   736 558                       PUSHADDRESS(ARG2, ARG3);
  1060   763 558                       PUSHINDIRECT(ARG1);
  1061                             END;
  1062                             END;
  1063
  1064   120                       PUSHIND1"(TYPE)":
  1065   455 763 558               BEGIN READ1ARG; PUSHINDIRECT(ARG1) END;
  1066
  1067   120                       PUSHADDR1"(MODE, DISPL)":
  1068   458 736 558 558           BEGIN READ2ARG; PUSHADDRESS(ARG1, ARG2) END;
  1069
  1070   121                       FIELD1"(DISPL)":
  1071   455                       BEGIN READ1ARG;
  1072   558                         IF ARG1 <> 0
  1073   511 215 558 28 269           THEN WRITE2(INC_2, (ARG1 DIV WORDLENGTH), B)
  1074                             END;
  1075
  1076   121                       INDEX1"(MIN, MAX, LENGTH, TYPE)":
  1077   464                       BEGIN  READ4ARG;
  1078   653 558 608                 PICK_PUSHCONST(ARG1);  PUSHWORD;
  1079   653 558 608                 PICK_PUSHCONST(ARG2);  PUSHWORD;
```

```
1080   WRITE1(CHK_2);              POP(TWOWORDS);
1081   PICK_PUSHCONST(ARG1);       PUSHWORD;
1082   WRITE1(SB1_2);              POPWORD;
1083   IF ARG4 <> BYTETYPE
1084   THEN BEGIN
1085     WRITE2(IXA_2, ARG3 DIV WORDLENGTH, B);    POPWORD
1086     END;
1087
1088
1089 POINTER1:
1090   BEGIN
1091   IF CHECK
1092   THEN BEGIN
1093     WRITE1(DUP1_2);           PUSHWORD;
1094     WRITE1(SLDC00_2);         PUSHWORD;
1095     WRITE2(EFJ_2, CALLERRORDISPL, SB);   POP(TWOWORDS);
1096     "GENERATE CALL TO POINTER ERROR ROUTINE.  $$$"
1097     END
1098   END;
1099
1100 VARIANT1"(TAGSET, DISPL)":
1101 "CHANGE PRIOR PASSES SO THAT TOS TAG IS <VALUE> OF
1102  BIT WHICH IS SET FOR TAG.   $$$"
1103   BEGIN   READ2ARG;
1104   IF (VARNTCHECK)
1105   THEN BEGIN
1106     WRITE1(DUP1_2);           PUSHWORD;
1107     WRITE2(IND_2, ARG2 DIV WORDLENGTH, B);   "POPWORD; PUSHWORD;"
1108     WRITE2(LDCI_2, ARG1, W);  PUSHWORD;
1109     WRITE1(LAND_2);           "POPWORD;"
1110     WRITE1(SLDC00_2);         "PUSHWORD;"
1111     WRITE2(EFJ_2, CALLERRORDISPL, SB);   POP(TWOWORDS);
1112     "GENERATE CALL TO VARIANT ERROR ROUTINE.   $$$"
1113     END
1114
1115   END;
1116 RANGE1"(MIN, MAX)":
1117   BEGIN READ2ARG;
1118   IF CHECK
1119   THEN BEGIN
1120     PICK_PUSHCONST(ARG1);     PUSHWORD;
1121     PICK_PUSHCONST(ARG2);     PUSHWORD;
1122     WRITE1(CHK_2);            POP(TWOWORDS);
1123     END
1124
1125   END;
1126 ASSIGN1"(TYPE)":
1127   BEGIN READ1ARG;
1128   CASE ARG1 OF
1129   BYTETYPE:
1130     BEGIN WRITE1(STB_2); POPWORD END;
1131   WORDTYPE:
1132     BEGIN WRITE1(STO_2); POPWORD END;
1133   REALTYPE:
1134     BEGIN
1135     WRITE2(STM_2, REALLENGTH DIV WORDLENGTH, UR); POPREAL,
1136     END;
1137   SETTYPE:
1138     BEGIN
1139     WRITE2(FJP_2, 0, SB); POPWORD; "POP SET LENGTH VALUE"
```

```
1140            WRITE2(STH_2, SRTLENGTH DIV WORDLENGTH, UB); POPSET;
1141            END
1142        POPWORD;
1143    END;
1144
1145
1146 C "ASSIGNTAG1 (LENGTH) :     ASSIGNTAG1 NEVER PRODUCED BY PASS 5.
1147 C  BEGIN READ1ARG;
1148 C      IF ARG1 = 0 THEN WRITE1(COPYWORD2)
1149 C      ELSE WRITE2(COPYTAG2, ARG1 DIV WORDLENGTH);
1150 C      POPWORD; POPWORD;
1151 C  END; "
1152
1153 COPY1"(LENGTH)":
1154    BEGIN READ1ARG;
1155    WRITE2(MOV_2, ARG1 DIV WORDLENGTH, B);     POP(TWOWORDS);
1156    END;
1157
1158 NEW1"(LENGTH, INITIALIZE)":     "$$$"
1159    BEGIN READ2ARG;
1160    "IF (ARG2 = 1) & CHECK
1161 C      THEN WRITE3(NEWINIT2, BLOCK, ARG1)
1162 C      ELSE WRITE3(NEW2, BLOCK, ARG1);"
1163    POPWORD;
1164    END;
1165
1166 NOT1:
1167    WRITE1(LNOT_2);
1168
1169 AND1"(TYPE)":
1170    BEGIN READ1ARG;
1171    IF ARG1 = WORDTYPE
1172    THEN BEGIN WRITE1(LAND_2); POPWORD END
1173    ELSE BEGIN WRITE1(INT_2); POPWORD; POPSET END;
1174    END;
1175
1176 OR1"(TYPE)":
1177    BEGIN READ1ARG;
1178    IF ARG1 = WORDTYPE
1179    THEN BEGIN WRITE1(LOR_2); POPWORD END
1180    ELSE BEGIN WRITE1(UNI_2); POPWORD; POPSET END;
1181    END;
1182
1183 NEG1"(TYPE)":
1184    BEGIN READ1ARG;
1185    IF ARG1 = WORDTYPE
1186    THEN WRITE1(NGI_2)
1187    ELSE WRITE1(NGR_2);
1188    END;
1189
1190 ADD1"(TYPE)":
1191    BEGIN READ1ARG;
1192    IF ARG1 = WORDTYPE
1193    THEN BEGIN WRITE1(ADI_2); POPWORD END
1194    ELSE BEGIN WRITE1(ADR_2); POPREAL END;
1195    END;
1196
1197 SUB1"(TYPE)":
1198    BEGIN READ1ARG;
1199    CASE ARG1 OF
```

Cross-reference numbers:

```
511   194   30    28    269   636
615

122
455
511   207   558   28    269   646   247
123
458

615

123
504   215

123
455
558        135   615
504        199   615   636
504        213   615
123
455
558        135   615
504        199   615   636
504        212   615
124
455
558        135
504        214
504        215
124
455
558        135   615
504        199   626
504        206
124
455
558
```

```
1200         WORDTYPE: BEGIN WRITE1(SB1_2); POPWORD END;               135  504  199  615
1201         REAL.TYPE: BEGIN WRITE1(SBR_2); POPREAL END;              135  504  206  626
1202         SETTYPE: BEGIN WRITE1(DIF_2); POPWORD; POPSET END         135  504  213  615  636
1203       END;
1204     END;
1205
1206 MUL1"(TYPE)":                                                     124
1207     BEGIN READ1ARG;                                              455
1208       IF ARG1 = WORDTYPE                                         558  135
1209       THEN BEGIN WRITE1(MPI_2); POPWORD END                      504  194  615
1210       ELSE BEGIN WRITE1(MPR_2); POPREAL END;                     504  206  626
1211     END;
1212
1213 DIV1"(TYPE)":                                                     125
1214     BEGIN READ1ARG;                                              455
1215       IF ARG1 = WORDTYPE                                         558  135
1216       THEN BEGIN WRITE1(DVL_2); POPWORD END                      504  194  615
1217       ELSE BEGIN WRITE1(DVR_2); POPREAL END;                     504  206  626
1218     END;
1219
1220 MOD1"(TYPE)":                                                     125
1221     BEGIN READ1ARG; WRITE1(MODI_2); POPWORD END;                 455  504  194  615
1222
1223 "(NOT USED)"
1224
1225 "(NOT USED)"
1226
1227 FUNCTION1"(STANDARDFUNC, TYPE)":                                  126
1228     BEGIN READ2ARG;                                              455
1229       IF (ARG1 >= MIN_FUNC) AND (ARG1 <= MAX_FUNC) THEN          558  162  162
1230       CASE ARG1 OF                                               558
1231         TRUNC1:                                                  160
1232           BEGIN WRITE1(TNC_2); POPREAL; PUSHWORD END;            160  205  626  608
1233         ABS1:
1234           IF ARG2 = WORDTYPE                                     558  135
1235           THEN WRITE1(ABI_2)                                     504  214
1236           ELSE WRITE1(ABR_2);                                    504  214
1237         SUCC1:                                                   160
1238           BEGIN
1239             WRITE1(SLDC01_2);        PUSHWORD;                   504  172  608
1240             WRITE1(ADI_2);           POPWORD                     504  199  615
1241           END;
1242         PRED1:                                                   160
1243           BEGIN
1244             WRITE1(SLDC01_2);        PUSHWORD;                   504  172  608
1245             WRITE1(SBI_2);           POPWORD;                    504  199  615
1246           END;
1247         CONV1:                                                   161
1248           BEGIN WRITE1(FLT_2); POPWORD; PUSHREAL END;            504  209  615  619
1249         EMPTY1:                                                  161
1250           BEGIN
1251             WRITE1(SLDC00_2);        PUSHWORD;                   504  172  608
1252             WRITE1(EQUI_2);          POPWORD;                    504  202  615
1253           END;
1254         ATTRIBUTE1,    "$$$"                                     161
1255             "WRITE1(ATTRIBUTE2);"
1256         REALTIME1:     "$$$"                                     161
1257             "BEGIN WRITE1(REALTIME2); PUSHWORD END"
1258       END;
1259     END;
```

```
1260  BUILDSET1:                                                                          126
1261  BEGIN
1262    WRITE1(SLDCON_2);   PUSHWORD;  "PUSH LOWER LEGAL SET BOUND."                504   172         608
1263    "PUSH UPPER LEGAL SET BOUND-- EXPECTED TO BE 127."
1264    PICK_PUSHCONST((SETLENGTH * 8 "BITS IN BYTE") - 1); PUSHWORD;               653    30   608      247
1265    WRITE1(CHK_2);      POP(TWOWORDS);                                          504   208   646
1266    WRITE1(DUP1_2);     PUSHWORD;                                               504   214   608      608
1267    WRITE1(SRS_2);      PUSHSET;  PUSHWORD;                                     504   205   629
1268    "THE SUBRANGE SET (OF 1 ITEM) JUST BUILT MAY BE LESS THAN
1269     THE FIXED SET SIZE OF 16 BYTES. CALL TO 'PUSHSET' WILL
1270     INCREASE CALCULATED SIZE OF RUNTIME STACK BY 16, REGARDLESS
1271     OF ACTUAL SET SIZE."
1272    WRITE1(UN1_2);   POPWORD;  POPSET                                           504   212   615      636
1273  END;
1274
1275  COMPARE1"(COMPARISON, TYPE)":                                                 126
1276  BEGIN READ2ARG;                                                              458
1277  CASE ARG2 OF                                                                 558
1278    WORDTYPE: BEGIN                                                            135
1279      CASE ARG1 OF                                                            558
1280        LESS:       BEGIN WRITE1(GEQ1_2); WRITE1(LNOT_2); END;             155  504  202  504  215
1281        EQUAL:            WRITE1(EQU1_2);                                    155  504  202
1282        GREATER:    BEGIN WRITE1(LEQ1_2); WRITE1(LNOT_2); END;             155  504  202  504  215
1283        NOTLESS:          WRITE1(GEQ1_2);                                    155  504  202
1284        NOTEQUAL:         WRITE1(NEQ1_2);                                    156  504  202
1285        NOTGREATER:       WRITE1(LEQ1_2)                                     156  504  202
1286      END;
1287      POPWORD;
1288    END;
1289
1290    REALTYPE: BEGIN                                                            135
1291      CASE ARG1 OF                                                            558
1292        LESS:       BEGIN WRITE1(GEQREAL_2); WRITE1(LNOT_2); END;          155  504  209  504  215
1293        EQUAL:            WRITE1(EQUREAL_2);                                 155  504  209
1294        GREATER:    BEGIN WRITE1(LEQREAL_2); WRITE1(LNOT_2); END;          155  504  209  504  215
1295        NOTLESS:          WRITE1(GEQREAL_2);                                 155  504  209
1296        NOTEQUAL: BEGIN WRITE1(EQUREAL_2); WRITE1(LNOT_2); END;            156  504  209  504  215
1297        NOTGREATER:       WRITE1(LEQREAL_2)                                  156  504  209
1298      END;
1299      POPREAL; POPREAL; PUSHWORD;
1300    END;
1301
1302    SETTYPE: BEGIN                                                             135
1303      CASE ARG1 OF                                                            558
1304        EQUAL:            WRITE1(EQUPWR_2);                                   155  504  203
1305        NOTLESS:          WRITE1(GEQPWR_2);                                   155  504  204
1306        NOTEQUAL: BEGIN WRITE1(EQUPWR_2); WRITE1(LNOT_2); END;             156  504  203  504  215
1307        NOTGREATER: WRITE1(LEQPWR_2);                                        156  504  203
1308        INSET:            WRITE1(INN_2)                                       156  504  212
1309      END;
1310      POPWORD; POPSET;                                                        615  636
1311      IF ARG1 <> INSET THEN                                                    558  156
1312        BEGIN "POPWORD;" POPSET; "PUSHWORD" END;                             636
1313    END;
1314  END;
1315
1316  COMPSTRUC1"(COMPARISON, LENGTH)":                                            126
1317  BEGIN READ2ARG;                                                             458
```

```
1320   CASE ARG1 OF
1321     LESS:        BEGIN WRITE2(GEQBYT_2, ARG2, B); WRITE1(LNOT_2); END;
1322     EQUAL:             WRITE2(EQUBYT_2, ARG2, B);
1323     GREATER:     BEGIN WRITE2(GEQBYT_2, ARG2, B); WRITE1(LNOT_2); END;
1324     NOTLESS:           WRITE2(GEQBYT_2, ARG2, B);
1325     NOTEQUAL;    BEGIN WRITE2(EQUBYT_2, ARG2, B); WRITE1(LNOT_2); END;
1326     NOTGREATER:        WRITE2(LEQBYT_2, ARG2, B)
1327   END;
1328   POPWORD;
1329 END;
1330
1331 FUNCVALUE1"(MODE, TYPE)":
1332 BEGIN READ2ARG;
1333   CASE ARG1 OF
1334
1335     MODE1, MODE3:   "PROCEDURE MODE, PROCESS ENTRY MODE"
1336        IF ARG2 = WORDTYPE
1337        THEN BEGIN WRITE1(SLDC00_2); PUSHWORD END
1338        ELSE BEGIN WRITE1(SLDC00_2); WRITE1(SLDC00_2); PUSHREAL END;
1339
1340     MODE4, MODE5, MODE12: "ENTRY MODES: CLASS, MONITOR, MGR"
1341        IF ARG2 = WORDTYPE
1342        THEN BEGIN
1343           WRITE1(SLDC00_2);   WRITE1(SWAP_2);   PUSHWORD
1344        END
1345        ELSE BEGIN
1346           WRITE1(SLDC00_2);   WRITE1(SWAP_2);
1347           WRITE1(SLDC00_2);   WRITE1(SWAP_2);   PUSHREAL
1348        END
1349   END
1350 END;
1351
1352 DEFLABEL1"(LABEL)":
1353 BEGIN READ1ARG;
1354   ENTER(JUMPTABLE, ARG1, LOCATION);
1355   IF NUMBER
1356   THEN BEGIN
1357      PICK_PUSHCONST(LINE);   PUSHWORD;
1358      WRITE2(STL_2, FIXLCL_LINE, B);   POPWORD;
1359   END;
1360 END;
1361
1362 JUMP1"(LABEL)":
1363 BEGIN READ1ARG;
1364   WRITE2(UJPL_2, ARG1, W); WRITELOCATION;
1365 END;
1366
1367 FALSEJUMP1"(LABEL)":
1368 BEGIN READ1ARG;
1369   WRITE2(FJPL_2, ARG1, W); WRITELOCATION;
1370   POPWORD;
1371 END;
1372
1373 CASEJUMP1"(MIN, MAX, LABELS)":
C 1374 "GENERATE RANGE CHECK ON CASE SELECTOR VALUE (TOS)
C 1375  AND PUT CASEJUMP OPERATOR INTO CODE STREAM.
C 1376  OPERAND TO XJP IS OFFSET (IN WORDS, FROM START
C 1377  OF SEGMENT) INTO CONSTANT POOL WHERE CASE TABLE
C 1378  WILL BE FOUND AT RUN TIME. CASE TABLE CONSISTS
C 1379  OF MIN. CASE INDEX (1 WORD), MAX. CASE INDEX
```

```
c 1380      " (1 WORD), AND JUMP DISPLACEMENTS (1 WORD EACH)."
  1381      BEGIN READ2ARG;
  1382      IF ARG2-ARG1 > 0
  1383      THEN BEGIN
  1384          PICK_PUSHICONST (ARG1);   PUSHWORD;   "PUSH MIN"
  1385          PICK_PUSHICONST (ARG2);   PUSHWORD;   "PUSH MAX"
  1386          WRITE1 (CHK_2);   POP(TWOWORDS);   "RANGECHECK"
c 1387      " 'LINK0.CONSTANTS' IS AMOUNT OF SPACE NEEDED IN CONSTANT
c 1388      POOL FOR LONG CONSTANTS ONLY (DOES NOT INCLUDE SPACE
c 1389      NEEDED FOR FOR CASE TABLES)."
  1390      " 'XJPOFFSETPTR' POINTS TO WORD LAST FILLED."
  1391      WRITE2 (XJP_2, ((ONEWORD + LINK0.CONSTANT3) DIV WORDLENGTH) +
  1392                     XJPOFFSETPTR, B);   TOPWORD;
  1393
  1394      "BUILD CASE TABLE"
  1395      XJPOFFSETPTR := SUCC(XJPOFFSETPTR);
  1396      ENTER (XJPTABLE, XJPOFFSETPTR, ARG1);   "MIN"
  1397      XJPOFFSETPTR := SUCC(XJPOFFSETPTR);
  1398      ENTER (XJPTABLE, XJPOFFSETPTR, ARG2);   "MAX"
  1399      "REMOVE CASE LABELS FROM CODE STREAM AND PUT
c 1400      CORRESPONDING OFFSETS IN CASE TABLE."
  1401      ARG3 := ARG1;
  1402      REPEAT
  1403          READ1ARG;   "READ A CASE LABEL"
  1404          ARG4 := ENTR(JUMPTABLE, ARG1) - LOCATION;
  1405          XJPOFFSETPTR := SUCC(XJPOFFSETPTR);
  1406          ENTER (XJPTABLE, XJPOFFSETPTR, ARG4);
  1407          ARG3 := ARG3 + 1;
  1408      UNTIL ARG3 > ARG2;
  1409      END;
  1410
  1411  END;
c 1412  "INITVAR1 (LENGTH) :       NEVER PRODUCED BY PASS 5.
c 1413      BEGIN READ1ARG;
c 1414      IF CHECK THEN WRITE2(INITVAR2, ARG1 DIV WORDLENGTH);
c 1415      END; "
  1416
  1417  CALL1"(MODE, LABEL, PARAMLENGTH)":       "$$$"
  1418      BEGIN READ3ARG;
  1419      IF ARG1 = MODE3       "IF PROCESS ENTRY"
  1420      THEN BEGIN
  1421          PICK_PUSHICONST(ARG2);   PUSHWORD; "PREFIX ROUTINE NO."
  1422          WRITE3(CXL_2, INTFSEG, STARTPROC, UB, UB);
  1423          ARG1:= WORDLENGTH
  1424      END
  1425      ELSE BEGIN
  1426          WRITE2(CPL_2, ARG2, UB);   PUSH(MSCWSIZE);
  1427          IF CONCURRENT
  1428          THEN ARG1:= ENTR(STACKTABLE, ARG2)
  1429          ELSE ARG1:= WORDLENGTH;
  1430      END;
  1431      PUSH(ARG1); POP(ARG1 + ARG3);
  1432      END;
  1433
  1434  ENTER1"(MODE, LABEL, PARAMLENGTH, VARLENGTH, TEMPLENGTH)":
  1435      BEGIN READ5ARG;
  1436      BLOCK := ARG2;
  1437      PARAMLENGTH := ARG3;
  1438      VARLENGTH := ARG4;
  1439      STACKLENGTH := ARG5;
```

```
1440  277  242                              AFTERBEGIN := TRUE;
1441  300                                   TEMP := 0;
1442  301                                   MAXTEMP := 0;
1443  303  303  247       247               LOCATION := LOCATION + TWOWORDS; "EXIT-IC & DATASIZE"
1444  558                                   CASE ARG1 OF
1445  140                                   MODE1: "PROCEDURE"
1446                                         BEGIN
1447  802  283  302  297  247                 ENTER (DATASIZETABLE, BLOCK, (VARLENGTH+TWOWORDS)
1448   28       247                                                          DIV WORDLENGTH);
1449  299  298  297  297                      POPLENGTH := PARAMLENGTH + VARLENGTH + TWOWORDS;
1450  653  306  608  608                      PICK_PUSHCONST(LINE);              PUSHWORD;
1451  511  200  238  269  615                 WRITE2(STL_2, FIXLCL_LINE, B); POPWORD;
1452  141                                   END;
1453                                         MODE2: "PROGRAM"
1454                                         BEGIN
1455                                         "BUMP 'JOB' "  "$$$"
1456                                         GEN_SAVELINE_SAVEGBASE;
1457  895                                   "MAKE GLOBAL BASE SAME AS PRESENT LOCAL BASE"
1458  653  234  608                           PICK_PUSHCONST(GBASE_REG);       PUSHWORD;
1459  653  233  608                           PICK_PUSHCONST(LBASE_REG);       PUSHWORD;
1460  504  198                                WRITE1 (LPR_2);                  POPWORD;
1461  504  210  646  247                      WRITE1 (SPR_2);             POP (TWOWORDS)
1462  142                                   END;
1463                                         MODE3: "PROCESS ENTRY"
1464                                         BEGIN
1465  802  283  302  297  247                 ENTER (DATASIZETABLE, BLOCK, (VARLENGTH+TWOWORDS)
1466   28  297  247                                                          DIV WORDLENGTH);
1467  299  297                                POPLENGTH := VARLENGTH + TWOWORDS;
1468  895                                     GEN_SAVELINE_SAVEGBASE;
1469                                         "RE-ESTABLISH OLD GLOBAL BASE"
1470  653  234  608                           PICK_PUSHCONST (GBASE_REG);       PUSHWORD;
1471  504  184  608                           WRITE1 (SLD002_2);               PUSHWORD;
1472  504  210  646                           WRITE1 (SPR_2);             POP (TWOWORDS);
1473                                         "SET 'JOB' := 0 "  "$$$"
1474                                         END;
1475  143                                   MODE4, "CLASS ENTRY"
1476  146                                   MODE7: "CLASS"
1477                              "  'POPLENGTH' MUST INCLUDE THE LOCATION CONTAINING
1478                                 THE ADDRESS OF THE CLASS VARIABLE.
1479                                 THE NEW GLOBAL BASE REG. MUST POINT A FEW WORDS
1480                                 BELOW FIRST GLOBAL VARIABLE. THE 'FEW WORDS'
1481                                 ARE THE MSCW. HARDWARE TAKES THE SIZE OF THE
1482                                 MSCW INTO ACCOUNT WHEN ACCESSING VARIABLES."
1483                                         BEGIN
1484  558  143                              IF ARG1 = MODE4
1485                                         THEN BEGIN "CLASS ENTRY"
1486  802  283  302  297  247                 ENTER (DATASIZETABLE, BLOCK, (VARLENGTH+TWOWORDS)
1487   28  298  246                                                          DIV WORDLENGTH);
1488  299  247                                POPLENGTH := PARAMLENGTH + ONEWORD +
1489  297                                                   VARLENGTH + TWOWORDS;
1490                                         END
1491  802  283  302  247   28               ELSE BEGIN "CLASS INITIAL STATEMENT"
1492  299  246  247                           ENTER (DATASIZETABLE, BLOCK, TWOWORDS DIV WORDLENGTH);
1493                                           POPLENGTH := ONEWORD + TWOWORDS;
1494                                         END;
1495  895                                   GEN_SAVELINE_SAVEGBASE;
1496                                         "SET NEW GLOBAL BASE"
1497  653  234  608                           PICK_PUSHCONST (GBASE_REG);       PUSHWORD;
1498  511  192  299   28  269                 WRITE2 (LDL_2, POPLENGTH DIV WORDLENGTH, B); PUSHWORD;
1499  653  241  608  608                      PICK_PUSHCONST (HSCWSIZE);        PUSHWORD;
```

```
1500        WRITE1 (SBL_2);                              POPWORD;
1501        WRITE1 (SPL_2);                              POP (TWOWORDS)
1502        END;
1503    MODE5,   "MONITOR ENTRY"
1504    MODE12:  "MANAGER ENTRY"
1505        BEGIN
1506        ENTER (DATASIZETABLE, BLOCK, (VARLENGTH + TWOWORDS)
1507                                      DIV WORDLENGTH);
1508        POPLENGTH := PARAMLENGTH + ONEWORD + VARLENGTH + TWOWORDS;
1509        GEN_SAVELINE_SAVEBASE;
1510        "NEW GLOBAL BASE"
1511        PICK_PUSHCONST (GBASE_REG);                  PUSHWORD;
1512        WRITE2 (LDL_2, POPLENGTH DIV WORDLENGTH, B); PUSHWORD;
1513        PICK_PUSHCONST (MSCWSIZE);                   PUSHWORD;
1514        WRITE1 (SBL_2);                              POPWORD;
1515        WRITE1 (SPL_2);                              POP (TWOWORDS);
1516        "KERNEL CALL TO ENTER GATE"      "$$$"
1517        END;
1518    MODE6:  "PROCESS"
1519        BEGIN
1520        ENTER (DATASIZETABLE, BLOCK, TWOWORDS DIV WORDLENGTH);
1521        POPLENGTH := TWOWORDS;
1522        PICK_PUSHCONST (LINE);                       PUSHWORD;
1523        WRITE2 (STL_2, FIXLCL_LINE, B);              POPWORD
1524        END;
1525    MODE8:   "MONITOR"
1526    MODE11:  "MANAGER"              "MGR"
1527        BEGIN
1528        ENTER (DATASIZETABLE, BLOCK, TWOWORDS DIV WORDLENGTH);
1529        POPLENGTH := ONEWORD + TWOWORDS;
1530        GEN_SAVELINE_SAVEBASE;
1531        "SET NEW GLOBAL BASE"
1532        PICK_PUSHCONST (GBASE_REG);                  PUSHWORD;
1533        WRITE2 (LDL_2, POPLENGTH DIV WORDLENGTH, B); PUSHWORD;
1534        PICK_PUSHCONST (MSCWSIZE);                   PUSHWORD;
1535        WRITE1 (SBL_2);                              POPWORD;
1536        WRITE1 (SPL_2);                              POP (TWOWORDS);
1537        "KERNEL CALL TO INITIALIZE GATE"     "$$$"
1538        END;
1539    MODE10:  "UNDEFINED"
1540        END;
1541        END;
1542
1543    RETURN1"(MODE)":
1544        BEGIN READ1ARG;
1545        "CURRENT LOCATION IS THIS ROUTINE'S 'EXIT IC' "
1546        ENTER (EXITICTABLE, BLOCK, LOCATION);
1547        CASE ARG1 OF
1548        MODE1:  BEGIN           "PROCEDURE MODE"
1549                WRITE2(RFU_2, POPLENGTH DIV WORDLENGTH, B);
1550                POP (POPLENGTH + MSCWSIZE)
1551                END;
1552        MODE2:  BEGIN           "PROGRAM MODE"     "$$$"
1553                "PUSH 'CONTINUE'"
1554 C              WRITE1(SLDC00_2);
1555 C              WRITE2(NFJ_2, "DISPLACEMENT TO CONO:", SR);
1556 C              WRITE1(SLDC00_2);
1557 C              "STORE IN 'RESULT'"
1558 C              "CONO:"
1559 C              "PUSH POINTER TO 'LINE'"
```

```
C  1560                          WRITE1(SLDLO1_2);
C  1561                          WRITE1(STO_2);
C  1562                          "PUSH 'JOB'"
C  1563                          WRITE1(SLDCO0_2);
C  1564                          WRITE2(EFJ_2, "DISPLACEMENT TO SEQL:", SB);
C  1565                          "KERNEL CALL--SYSTEM ERROR"
C  1566                          WRITE2(UJP_2, "DISPLACEMENT TO DONE:", SB);
C  1567                          "SEQL:"
C  1568                          WRITE1(SLDCO6_2);
C  1569                          WRITE1(SLDOO2_2);
C  1570                          WRITE1(SPR);
C  1571                          "CLEAR 'JOB'"
C  1572                          "DONE:"
   1573   511  196  299  28  269 WRITE2(RPU_2, POPLENGTH DIV WORDLENGTH, B)
   1574                          END;
   1575   142                 MODE3: BEGIN       "PROCESS ENTRY MODE"
   1576   909                    GEN_EXIT;       " 'JOB' := SUCC ('JOB');"     "$$$"
   1577                          END;
   1578   143  909           MODE4:  GEN_EXIT;       "CLASS ENTRY MODE"
   1579   144              MODE5,             "MONITOR ENTRY"
   1580   151              MODE12: BEGIN             "MANAGER ENTRY MODE"
   1581                          "KERNEL CALL TO LEAVE GATE"      "$$$"
   1582   909                    GEN_EXIT
   1583                          END;
   1584   145              MODE6:              "PROCESS MODE"
   1585   511  196  299  28  269 WRITE2(RPU_2, POPLENGTH DIV WORDLENGTH, B);
   1586   146  909           MODE7: GEN_EXIT;       "CLASS MODE"
   1587   147              MODE8,              "MONITOR MODE"
   1588   150              MODE11:  BEGIN             "MANAGER MODE"
   1589                          "KERNEL CALL TO LEAVE GATE"       "$$$"
   1590   909                    GEN_EXIT
   1591                          END;
   1592   149              MODE10:              "UNDEFINED MODE"
   1593                          END;
   1594   291  301  297  248    IF (STACKLENGTH + MAXTEMP + VARLENGTH) < STACK_LIMIT
   1595   291  291  301  297    THEN STACKLENGTH := STACKLENGTH + MAXTEMP + VARLENGTH
   1596   591  249  256         ELSE ERROR (THIS_PASS, STACK_ERROR);
   1597   802  282  302  291    ENTER (STACKTABLE, BLOCK, STACKLENGTH);
   1598   277  337              AFTERBEGIN := FALSE;
   1599   303                   IF (LOCATION MOD 2) = 1       "IF NOT ON WORD BOUNDARY"
   1600   504  198              THEN WRITE1 (NOP_2)       "THEN PAD TO A WORD BOUNDARY"
   1601                          END;
   1602
   1603   129              POP1"(LENGTH)":
   1604   455                BEGIN READ1ARG;
   1605   558  247           WHILE ARG1 >= TWOWORDS DO
   1606                          BEGIN
   1607   511  210  269  247    WRITE2(NFJ_2, 0, SB);    POP(TWOWORDS);
   1608   558  558  247  646    ARG1 := ARG1 - TWOWORDS;
   1609                          END;
   1610   558  246              IF ARG1 = ONEWORD
   1611   511  211  269  615    THEN BEGIN WRITE2(FJP_2, 0, SB);    POPWORD;    END;
   1612                          END;
   1613
   1614   129              NEWLINE1"(NUMBER)":
   1615                          "LOCAL WORD 1 ALWAYS RESERVED FOR LINE #."
   1616   455                BEGIN READ1ARG;
   1617   306  558           LINE := ARG1;
   1618   278  337              AFTERERROR := FALSE;
   1619   480  277              IF (NUMBER AND AFTERBEGIN)
```

```
1620       THEN BEGIN
1621         WRITEARG (NEWLIN_2, NOTHE);  "CRUTCH FOR MNEMONICS PROGRAM"
1622         PICK_PUSHCONST(ARG1);  PUSHWORD;
1623         WRITE2(STL_2, FIXICL_LINE, B);  POPWORD
1624         END;
1625       END;
1626
1627   ERROR1:
1628       GENERATE:= FALSE;
1629
1630   CONSTANT1 "(LENGTH, VALUE)":
1631   "REMOVE LONG CONSTANT FROM CODE STREAM & PUT IN CONSTANT POOL."
1632   BEGIN READ1ARG;
1633       ARG2 := ARG1;
1634       FOR ARG3 := 1 TO
1635       (ARG1+SOURCE_WORD_LENGTH-1) DIV SOURCE_WORD_LENGTH DO
1636       BEGIN
1637       CONSTANTS:= CONSTANTS + 1 "WORD";
1638       READ1ARG;
1639
1640   "$IF SRC_16"
1641       ENTER(CONSTTABLE, CONSTANTS, ARG1);
1642   "$END"
1643
1644       "$IF SRC_3"
1645       SPLIT_CONST(ARG1, ARG1HI, ARG1LO);
1646       ENTER(CONSTTABLE, CONSTANTS, ARG1HI);
1647       IF (((ARG3-1)*SOURCE_WORD_LENGTH) + WORDLENGTH) < ARG2
1648       THEN BEGIN
1649       CONSTANTS := CONSTANTS + 1;
1650       ENTER(CONSTTABLE, CONSTANTS, ARG1LO)
1651       END
1652       "$END"
1653
1654       END;
1655
1656   END;
1657   MESSAGE1"(PASS, ERROR)":
1658   BEGIN READ2ARG;
1659       ERROR(ARG1, ARG2)
1660   END;
1661
1662   INCREMENT1:
1663       BEGIN
1664       WRITE1(DUP1_2);            PUSHWORD;
1665       WRITE2(LDM1_2, 1, UB);     PUSHWORD;
1666       WRITE1(SLDCO1_2);          PUSHWORD;
1667       WRITE1(ADI_2);             POPWORD;
1668       WRITE1(STO_2);             POP(TWOWORDS);
1669       END;
1670
1671   DECREMENT1:
1672       BEGIN
1673       WRITE1(DUP1_2);            PUSHWORD;
1674       WRITE2(LDM1_2, 1, UB);     PUSHWORD;
1675       WRITE1(SLDCO1_2);          PUSHWORD;
1676       WRITE1(SBI_2);             POPWORD;
1677       WRITE1(STO_2);             POP(TWOWORDS);
1678       END;
1679
```

```
1680 PROCEDURE1"(STANDARDPROCEDURE)":        "$$$"
1681      BEGIN READ1ARG;
1682      "IF (ARG1 >= MIN_PROC) AND (ARG1 <= MAX_PROC) THEN
C 1683        CASE ARG1 OF
C 1684          DELAY1:
C 1685            BEGIN WRITE1(DELAY2); POPWORD END;
C 1686          CONTINUE1:
C 1687            BEGIN WRITE1(CONTINUE2); POPWORD END;
C 1688          IO1:
C 1689            BEGIN WRITE1(IO2); POP(THREEWORDS) END;
C 1690          START1:
C 1691            WRITE1(START2);
C 1692          STOP1:
C 1693            BEGIN WRITE1(STOP2); POP(TWOWORDS) END;
C 1694          SETHEAP1:
C 1695            BEGIN WRITE1(SETHEAP2); POPWORD END;
C 1696          WAIT1:
C 1697            WRITE1(WAIT2)
C 1698        END; "
1699      END;
1700
1701 INIT1"(MODE, LABEL, PARAM.LENGTH, VARLENGTH)":
1702      BEGIN READ4ARG;
1703      IF ARG1 = MODE6        "IF PROCESS MODE"
1704      THEN BEGIN
1705        "KERNEL CALL TO INIT PROCESS"      "$$$"
1706        WRITE2 (FJP_2, 0, SB);      POPWORD
1707        END
1708      ELSE BEGIN   "CLASS OR MONITOR"
1709        WRITE2 (STM_2, ARG3 DIV WORDLENGTH, UB);
1710        POP ((ARG3 DIV WORDLENGTH) + 1);
1711        WRITE2 (CPL_2, ARG2, UB);
1712        PUSH (MSCWSIZE);
1713        END
1714      END;
1715
1716 PUSHLABEL1"(LABEL)": GEN_INTERFACE;
1717
1718 CALLPROG1:
1719      BEGIN
1720      "ADDRESS OF SEQ'L PROGRAM VARIABLE ON TOS"
1721      "PUT IT IN SEQ'L PGM'S SIB"
1722      WRITE1(SLDC01_2); PUSHWORD; "GET @ OF CURRENTLY RUNNING TIB"
1723      WRITE1(NGI_2);
1724      WRITE1(LPR_2);
1725      WRITE2(IND_2, SIBS_OFFSET, D);   "GET ADDRESS OF SIB VECTOR"
1726      WRITE2(INC_2, 5, D);    "SEGBASE FIELD IN SIB FOR SEGMENT
1727                               129 (SEQ'L PGM) IS 5 WORDS FROM
1728                               START OF SIB VECTOR."
1729
1730      "AFTER EXECUTION OF THE 'INC (5)', @ SEQ'L PROGRAM VARIABLE
1731       IS IN THE TOS-1 WORD, AND @ OF THE SEGBASE FIELD IN SIB
1732       NO. 129 IS IN TOS WORD."
1733      WRITE1(SWAP_2);
1734      WRITE1(STO_2);  POP(TWOWORDS);    "PUT @ PGM VAR IN SIB"
1735      WRITE3(CXL_2, SEQL_SEG, STARTPROC, UB, UB); "CALL SEQ'L PGM"
1736      READ_IFL(OP); READ1ARG   "DISPOSE OF 'POP' VIRTUAL OPERATOR
1737                               AND ITS OPERAND... NOT NEEDED FOR
1738                               MICROENGINE."
1739      END;
```

```
1740  EOM1"(VARLENGTH)":
1741      BEGIN
1742          DONE:=TRUE;
1743          READ1ARG; VARLENGTH:=ARG1;
1744          ENTER (DATASIZETABLE, STARTPROC,
1745                  (VARLENGTH + TWOWORDS) DIV WORDLENGTH);
1746
1747          WRITEARG(EOM_2, NOTME);
1748      END;
1749
1750
1751  DUPPOST: BEGIN WRITE1(DUP1_2); PUSHWORD END
1752
1753
1754
1755          END
1756  UNTIL DONE
1757  END;
1758
1759  BEGIN
1760      INIT_PASS(LINK);
1761      BEGINPASS;
1762      SCAN;
1763      ENDPASS;
1764      NEXT_PASS(LINK);
1765  END.
```

CROSS REFERENCE   * IS DEF   = IS ASG

```
-A-
ABL_2          214*  1235
ABR_2          214*  1236
ABS            420
ARG1           160*  1233
ADD1           124*  1190
ADL_2          199*  1193     1240  1667
ADJ_2          207*
ADR_2          206*  1194
AFTERBEGIN     277*  856=  1440=  1598=  1619
AFTERERROR     278*  593   595=   161B-
ALLOCATE       789*  843   844    845    846    847    849
AND1           123*  1169
ARG            414*  417   443*   446    447    511*   515   517   520   578*   580   581   582   666*   675   679   680   681
               691   692   693   756    703    704    705    711*  717   720   723   724   725   729   736*  740   744   747   750
               751   752
ARG1           308*  456   459    462    466    472    524*   530   533   537   541*   546   549   553   558*   564   568   573
               1051  1056  1060   1065   1068   1072   1073   1078  1081  1108  1120   1128   1155   1171   1178   1185   1192   1199
               1208  1215  1229   1229   1230   1280   1292   1304  1312  1320  1333   1354   1364   1369   1382   1384   1396   1401
               1404  1419  1423=  1428=  1429=  1431   1431   1431  1444  1484  1547   1608=  1608   1610   1617   1622   1633   1635
               1641  1645  1659   1703
ARG1HI         1039* 1645  1646
ARG1LO         1039* 1645  1650
ARG1TYP        525*  537   542*   553    559*   573   524*  530   533   537   541*   547   550   554   558*   564   568   573
               1059  1068  1079   1107   472    1234  1278  1321  1322  1323  1324   1325   1326   1336   1341   1382   1385   1398
               1408  1121  1428   1436   1633=  1647  1659  1711
ARG2TYP        525*  537   542*   554    559*   573
ARG3           308*  462   467    472    541*   547   550   565   569
               1408  1431  1437   1639=  1647   1709  1710
ARG3TYP        542*  554   559*   57     57
ARG4           308*  467   473    558*   565
ARG4TYP        559*  574
ARG5           308*  473
ARGLIST        91*   108    1439
ARGTAG         78*   82
ARGTYP         511*  520   570*   582    578*   582
ARGTYPE        81*   91
ASSIGN1        122*  126
ASSIGNTAG1     122*
ATTRIBUTE1     161*  1254
-B-
B              269*  487   717    720    724    725    729   740   744   747   751   752   756   902   906   915   969   1014
               1073  1085  1107   1155   1321   1322  1323  1324  1325  1326  1358   1392   1451   1498   1512   1523   1533   1549
               1573  1585  1623   1725   1726
BEGINPASS      830*  1761
BLOCK          102*  103*  302=   1436=  1447    1465  1486  1492  1506  1520  1528   1546   1597
BLOCKS         72*   846   847    849    878
BOOL           83   360=
BOOLEAN        83   278   280    318
BOOL.TYPE      79   83    360
BPT_2          198*
BUILDSET1      126*  1261
BYTETYPE       135*  766   1083   1129
-C-
```

ENTR 814* 823= 1404 1428
ENTRIES 789* 793
EOL 25* 329 331
EOM 25*
EOM1 132* 1741
EOM_2 217* 1747
EQUAL 155* 1282 1294 1305 1372
EQUBYT_2 204* 1322 1325
EQUL_2 202* 1252 1282
EQUFWR_2 203* 1305 1307
EQUREAL_2 209* 1294 1297
ERROR 400* 494 591* 611 622 632 642 1596 1659
ERROR1 129* 1627
EXITICTABLE 56* 283* 846 802= 882 1546

-F-
F 102* 103* 104*
FALSE 337 600 856 858
FALSEJUMP1 127* 1367
FF 25*
FIELD1 121* 1070
FILE 42* 102 103 104
FILE_LENGTH 104* 352 356 370 389 403
FILE_LIMIT 334* 353 357 370 390 404
FIVEWORDS 247*
FIXGHL_OLDG 240*
FIXLCL_LINE 238* 902 1358 1451 1523 1623
FIXLCL_OLDG 239* 906
FJP_2 211* 1369
FJP_2 211* 1139 1611 1706
FLAG_REG 229*
FLT_2 209* 1248
FOINARD 480*
FOURWORDS 247*
FULLWORD 18* 73 105 106 274 941
FUNCTION1 126* 1227
FUNCVALUE1 127* 1331

-G-
GBASE_REG 234* 904 912 1458 1470 1497 1511 1532
GENERATE 277* 600= 840= 869 1628=
GEN_EXIT 200* 383 508 518 534 551 570 582 858= 948= 1029=
GEN_INTERFAC 924* 1716 1578 1582 1586 1590
GEN_SAVELINE 895* 1456 1468 1495 1509 1530
GEQBYT_2 204* 1321 1324
GEQL_2 202* 1281 1284
GEQPWR_2 204* 1306
GEQREAL_2 209* 1293 1296
GET 102* 372
GERUSW_2 203*
GREATER 155* 1283 1295 1323

-H-
HALFWORD 245* 406 488
HIORDER 933* 1027
HIWORD 869* 891*

-I-
I 327* 330 330= 367* 370= 381* 386 400 415* 423 423= 424= 431* 433= 790* 793= 794 798=

ID	798	802*	805	814*	817	831*	841=	842
IDENTIFIER	85	85						
IDLENGTH	45*	45						
IDTYPE	79	85*						
IFINFO	63	64*						
IFLABEL	944*	982	985					
IFPAGES_OUT	324*	346=	356	358	389	392	393=	393
IFPAGE_OUT	324*	350	386=	392				
IFPTR	63*	75						
IFSEGNUM	305*	857=	950=	950	951	1027		
IFSEGSIZE	286*	842=	951					
IFWORDS_OUT	324*	346=	355	385=	386	387	395=	
INCREMENT1	130*	1662						
INC_2	215*	1073	1726					
INDEX1	121*	1076						
IND_2	215*	1107						
INFILE	250*	370	1725					
INIT1	131*	1701	372					
INITVAR1	126*							
INIT_PASS	340*	1760						
INN_2	212*	1309						
INSET	156*	1309	1312					
INT	84	364=						
INTEGER	18	19=	47	51	54	65	66	72	84	96	102	103	104	286	309	316	322	323
	327	367	381	414	415	431	436	443	480	482	504	511	524	541	558	578	639
	646	653	666	666	711	736	763	789	790	802	803	814	814	815	831	889	928	944
	1039							
INTERFACE	75*	842						
INTERFACEFIL	250*	356	358					
INTERFACES	65*	842						
INTERFACESIZ	66*	842						
INTFSEG	254*	1422						
INTTYPE	79	84*	364					
IRT_2	213*	1173						
IO1	166*							
IXA_2	211*	1085						
IXP_2	212*							

-J-

J	802*	811	815*	817=	818	821=	821	823
JUMP1	127*	1362	843					
JUMPTABLE	55*	282*	843	860	879=	879	1354	1404

-K-

K	803*	805=	806	809=	809	811		
KASE	943*	958=	958	980=	995	999=	1009	

-L-

LABELS	72*	843
LAE_2	197*	
LAND_2	199*	1109	1172
LAO_2	192*	747	756
LBASE_REG	233*	1959	
LCA_2	191*	740	
LDA_2	193*	752	
LDB_2	200*	767	
LDCB_2	191*	658	
LDCI_2	191*	659	1108
LDCN_2	197*	

Symbol	References
LDC_2	191#
LDE_2	197#
LDL_2	192# 717 724 1498 1512 1533
LDM_2	210# 771 775 1665 1674
LDO_2	192# 720 729
LDP_2	208#
LENGTH	639# 641 646# 647
LEQBYT_2	204# 1323 1326
LEQL_2	202# 1283 1286
LEQPWR_2	203# 1308
LEQREAL_2	209# 1295 1298
LESS	155# 1281 1293 1321
LEUSH_2	203#
LINE	306# 599 854 901 1357 1450 1522 1617= 866 867 868 871 872 878 1391 1760 1764
LINK	273# 340# 342= 349# 362 833 853
LJSTOPTION	38# 59
LLA_2	192# 744 751
LNOT_2	215# 1167 1281 1283 1293 1295 1297 1307 1321 1323 1325 484 486 488 489 489= 490 490= 495= 587 588 855= 878 1354 1404 1443=
LOCATION	303# 1443 1546 1599
LOD_2	193# 725
LOR_2	199# 1179
LOWORD	889# 891=
LPR_2	198# 905 1460 1724
LSL_2	197#
LSPLIM	230#

-N-

Symbol	References
MAIN	108# 848 977
MARK	105# 91
MAXARG	89# 415 422
MAXDIGIT	26# 66 286
MAXINTFAC	61# 612 623= 633= 643= 1742= 1594
MAXTEMP	301# 51 793 798 806 809 818 821
MAXWORD	40# 1229
MAX_FUNC	162# 623 633 643 1594 1595
MAX_PROC	168#
MESSAGE1	130# 1657
MESSAGE_2	217# 596
MININDEX	934# 956
MIN_FUNC	162# 1229
MIN_PROC	168#
MOD1	125# 1220
MODE	666# 673 711# 713 717 720 724 725 729 736# 738 740 744 747 751 752 756
MODE0	139# 674 739 743 1335 1445 1548
MODE1	140# 678 716 758 1539 1592
MODE10	149# 707 731 755 1526 1588 1580
MODE11	150# 696 728 755 1504 1504
MODE12	151# 670 716 743 1340 1453 1552
MODE2	141# 678 719 746 1453 1552
MODE3	142# 683 722 749 1335 1419 1463 1575
MODE4	143# 678 716 743 1340 1475 1404 1578
MODE5	144# 678 716 743 1340 1503 1579
MODE6	145# 696 728 755 1518 1584
MODE7	146# 696 728 755 1476 1586
MODE8	147# 696 728 755 1525 1587
MODE9	148#
MODI_2	194# 1221
MOV_2	207# 1155
MPL_2	194# 1209

POINTER	1459	1470	1497	1499	1511	1513	1522	1532	1534					
POINTER1	121*	1089												
POOL.SIZE	290*	853=	855	916	1080	1095	1111	1122	1155	1266	1386	1431	1461	1472	1501	1515	1536	1550
POP	646*	906	914	1710	1734										
	1607	1668	1677												
POP1	129*	1603													
POPLENGTH	299*	915	916	1449=	1467=	1488=	1493=	1498	1508=	1512	1521=	1529=	1533	1549	1550	1573	1585
POPREAL	626*	1135	1194	1201	1210	1217	1232	1300	1300						
POPSET	636*	1140	1173	1180	1202	1273	1311	1313							
POPWORD	615*	767	772	776	902	1082	1085	1130	1132	1139	1143	1163	1172	1173	1179	1180	1193	1200
	1202	1209	1216	1221	1240	1245	1248	1252	1273	1288	1311	1328	1358	1370	1392	1451	1500	1514
	1523	1535	1611	1623	1667	1676	1706								
PORTION	790*	792=	795	797=	803*	805=	808=	808	811	815*	817=	820=	820	823	
PRED1	160*	1242													
PRINTABS	414*	439													
PRINTARG	443*	515	530	530	546	547	547	564	564	565	565	580	587		
PRINTED	316*	428=	438	440=	440	480	448=	448							
PRINTEOL	427*	433	433	438	445	445									
PRINTFF	430*	861													
PRINTLIMIT	27*	438	445												
PRINTOP	436*	506	515	529	546	563									
PRINT_TEXT	326*	336													
PRIORITY_REQ	228*														
PROCEDURE1	131*	1680													
PROCSGEND	939*	1027													
PTR	86	342	362=												
PTRTYPE	79	86*	362												
PUSH	639*	1426	1431	1712											
PUSHADDR1	120*	1067													
PUSHADDRESS	736*	1059	1068												
PUSHCONST1	120*	1049													
PUSHIND1	120*	1064													
PUSHINDIRECT	763*	1060	1065												
PUSHLABEL1	131*	1716													
PUSHREM,	619*	772	776	1338	1347										
PUSHSET	629*	776	1248												
PUSHVALUE	711*	1057	1268												
PUSHVAR1	120*	1054													
PUSHWORD	608*	733	760	778	901	904	912	913	1051	1078	1079	1081	1093	1094	1106	1108	1120	1121
	1232	1239	1244	1251	1263	1265	1267	1268	1300	1337	1343	1357	1384	1385	1421	1450	1458	1459
	1470	1471	1497	1498	1499	1511	1512	1513	1522	1532	1533	1534	1622	1664	1665	1666	1673	1674
	1675	1722	1751												
PUT	103*	354	358	392	406										
-Q-															
QLINK_REG	226*														
-R-															
RANGE1	122*	1116													
RBF_2	198*														
READ	98*														
READ1ARG	455*	459	1042	1051	1065	1071	1127	1154	1170	1177	1184	1191	1198	1207	1214	1221	1353	1363	1368
	1403	1574	1604	1616	1632	1638	1681	1736	1744						
READ2ARG	458*	1068	1103	1117	1159	1228	1277	1319	1332	1381	1658				
READ3ARG	461*	1055	1118	1702											
READ4ARG	464*	1077													
READ5ARG	470*	1435													
READYQ_REQ	223*														
READ_IFL	367*	456	459	462	462	462	466	466	467	467	467	472	472	472	473	473	982	996

Symbol							
SLDC25_2	178*						
SLDC26_2	178*						
SLDC27_2	178*						
SLDC28_2	179*						
SLDC29_2	179*						
SLDC30_2	179*						
SLDC31_2	179*						
SLDL01_2	180*	968					
SLDL02_2	180*	913					
SLDL03_2	180*						
SLDL04_2	180*						
SLDL05_2	181*						
SLDL06_2	181*						
SLDL07_2	181*						
SLDL08_2	181*						
SLDL09_2	182*						
SLDL10_2	182*						
SLDL11_2	182*						
SLDL12_2	182*						
SLDL13_2	183*						
SLDL14_2	183*						
SLDL15_2	183*						
SLDL16_2	183*						
SLDO01_2	184*						
SLDO02_2	184*	1471					
SLDO03_2	184*						
SLDO04_2	185*						
SLDO05_2	185*						
SLDO06_2	185*						
SLDO07_2	185*						
SLDO08_2	186*						
SLDO09_2	186*						
SLDO10_2	186*						
SLDO11_2	186*						
SLDO12_2	187*						
SLDO13_2	187*						
SLDO14_2	187*						
SLDO15_2	187*						
SLDO16_2	187*						
SOURCE_WORD_	36*	1635	1635	1647			
SPLIT.LENGTH	244*						
SPLIT_CONST	088*	1645					
SPLIT_INTEGE	261*	888					
SPR_2	210*	914	1461	1472	1501	1515	1536
SRO_2	200*						
SRS_2	205*	1268					
STACKLENGTH	54*	291*	1439=	1594	1595=	1595	1597
STACKPTR_REG	232*						
STACKTABLE	282*	849	1428	1597			
STACK_ERROR	256*	611	622	632	642	1596	
STACK_LIMIT	248*	610	621	631	641	1594	
START1	166*						
STARTPROC	255*	860	1422	1735	1745		
STB_2	208*	1130					
STE_2	212*						
STKENTRY	926	927*					
STL_2	200*	902	906	1358	1451	1523	1623
STL_2	194*	1135	1140	1709			
STOP1	167*						
STO_2	207*	1132	1668	1677	1734		

STP_2 208*
STR_2 200*
SUB1 124* 1197
SUCC 373 377 385 393 399 407 950 1395 1397 1405
SUCC1 160* 1237
SUMMARY 276* 835=
SUMMARYOPTIO 30* 835
SMAP_2 205* 1343 1346 1347 1733

-T-
T 415* 420= 423 789* 792 792 802* 805 814* 817
TABLE 48 49*
TABLEPART 53* 58
TABLEPTR 48* 50 56 283 789 790 802 803 814 815
TABLERESET 274* 848 870 878 879 880 801 882 883
TABLES 74* 871 878
TABLESPTR 58* 74
TAG 82* 360= 362= 364=
TASK_REG 225*
TEMP 300* 506 513 527 544 561 580 587 610 610= 612 612 616= 616 621 621= 621 623 623 627= 627 631= 631
633 633 637= 637 641 641= 643 643 647= 647 861 1441=
TEST 276* 836
TESTOPTION 38* 836
TEXT 326* 330 330
TEXT_LENGTH 90* 92
TEXT_TYPE 92* 326
THIS_PASS 249* 494
THREEWORDS 247* 611 622 632 642 1596
TNC_2 205* 1232
TOP 105* 106*
TOS 945* 978* 986 988= 945 946 948 1008 1008 1010 1010= 1010 1140 1480 1743
TOSPTR 926* 929* 303 343 595
TRUE 242 243
TRUNC1 160* 1231
TWOWORDS 247* 680 681 692 906 914 1080 1095 1111 1122 1155 1266 1386 1443 1447 1449 1461 1465
1467 1472 1486 1489 1492 1493 1501 1506 1508 1515 1520 1521 1528 1529 1536 1605 1607 1608
1668 1677 1734 1746
TYPROFCODE 269* 402 511 525 542 559 578

-U-
UB 269* 486 658 771 775 1008 1008 1135 1140 1422 1422 1426 1665 1674 1709 1711 1735
UJPL_2 193* 971 1009 1364
UJP_2 193*
UNI_2 212* 1180 1273
UPDLOC 482* 508 520 520 536 537 537 553 553 554 554
USPLIM 231*

-V-
VALUE 482* 487 653* 655 655 656 657 657 658 659 659
VARIANT1 121* 1100
VARLENGTH 297* 681 705 1430= 1447 1449 1465 1467 1489 1506 1508
VARNTCHECK 278* 838= 1104 1406 1594 1595 1744= 1746
VARNTCHECKOP 39* 59 838
VARTYPE 763* 765 572 573 573 574 582

-W-
W 269* 490 659 953 956 957 958 961 962 971 1009 1017 1027 1108 1364 1369
WAIT1 167*
WAIT_2 213*
WORD 808* 891 891

Identifier	References
WORDLENGTH	28* 610 616 675 680 681 692 693 704 705 771 775 777 844 845 915 1073 1085 1107 1135 1140 1155 1391 1423 1429 1448 1466 1487 1492 1498 1505 1507 1512 1520 1528 1533 1549 1573 1647 1709 1710 1716
WORDS_IN	322* 344= 369 375= 377= 378
WORDS_OUT	323* 345= 351 399= 399 400 401 409=
WORDTYPE	135* 768 1056 1131 1171 1234 1279 1311 1336
WRITE	99* 329 330 331 423 424 428 433 439 446 446
WRITE1	504* 656 767 769 905 913 914 968 970 1013 1080 1082 1093 1094 1106 1109 1110 1122 1130 1132 1167 1172 1173 1179 1180 1186 1187 1193 1194 1200 1201 1202 1209 1210 1216 1217 1221 1232 1235 1236 1239 1240 1244 1245 1248 1251 1252 1263 1266 1267 1268 1273 1281 1282 1283 1284 1286 1293 1294 1295 1296 1297 1298 1305 1306 1307 1307 1308 1309 1321 1323 1325 1337 1338 1343 1346 1347 1386 1460 1461 1471 1472 1500 1501 1514 1515 1535 1536 1600 1660 1664 1666 1673 1675 1676 1677 1722 1723 1724 1733 1734 1751
WRITE2	511* 658 659 717 720 724 729 740 744 747 751 756 771 775 902 906 915 969 971 1009 1014 1073 1085 1095 1107 1108 1111 1135 1139 1140 1155 1321 1322 1323 1324 1325 1326 1350 1364 1369 1391 1426 1451 1498 1512 1523 1533 1549 1573 1607 1611 1623 1665 1674 1706 1709 1711 1725 1726 1735
WRITE3	524* 725 752 1008 1017 1027
WRITE4	541*
WRITE5	550*
WRITEARG	570* 596 597 598 599 900 953 956 957 958 961 962 1017 1621 1747
WRITELOCATIO	585* 1364 1369
WRITE_IFL	381* 507 517 517 532 533 533 533 549 549 550 550 567 568 568 569 569 581 588
-X-	
XJFOFFSETPTR	289* 852= 867 1392 1395 1396 1397= 1398 1405= 1406
XJPTABLE	55* 282* 845 881= 938* 1395= 1396 1397 1398 1405
XJP_2	211* 969 1391 801 938*
XJP_OFFSETS	72* 845 853 867=

END XREF 538 IDENTIFIERS 2678 TOTAL REFERENCES
745 COLLISIONS.

APPENDIX C

SOURCE CODE FOR MEPASCAL COMPILER DRIVER: MEPASCAL

```
 1   "KANSAS STATE UNIVERSITY"
 2   "DEPARTMENT OF COMPUTER SCIENCE"
 3   "DRIVER PROGRAM FOR METASCAL COMPILER"
 4
 5   "#########
 6   #  PREFIX #
 7   #########"
 8
 9
10   TYPE FULLWORD = INTEGER;
11        INTEGER = SHORTINTEGER;
12
13   CONST NL = '(:10:)';  FF = '(:12:)';  CR = '(:13:)';  EM = '(:25:)';
14
15   CONST PAGELENGTH = 512;
16   TYPE PAGE = ARRAY (.1..PAGELENGTH.) OF CHAR;
17
18   CONST LINELENGTH = 132;
19   TYPE LINE = ARRAY (.1..LINELENGTH.) OF CHAR;
20
21   CONST IDLENGTH = 12;
22   TYPE IDENTIFIER = ARRAY (.1..IDLENGTH.) OF CHAR;
23
24   TYPE FILE = 1..2;
25
26   TYPE FILEKIND = (EMPTY, SCRATCH, ASCII, SEQCODE, CONCODE);
27
28   TYPE FILEATTR = RECORD
29                     KIND: FILEKIND;
30                     ADDR: INTEGER;
31                     PROTECTED: BOOLEAN;
32                     NOTUSED: ARRAY (.1..5.) OF INTEGER
33                   END;
34
35   TYPE IODEVICE =
36        (TYPEDEVICE, DISKDEVICE, TAPEDEVICE, PRINTDEVICE, CARDDEVICE);
37
38   TYPE IOOPERATION = (INPUT, OUTPUT, MOVE, CONTROL);
39
40   TYPE IOARG = (WRITEEOF, REWIND, UPSPACE, BACKSPACE);
41
42   TYPE IORESULT =
43        (COMPLETE, INTERVENTION, TRANSMISSION, FAILURE,
44         ENDFILE, ENDMEDIUM, STARTMEDIUM);
45
46   TYPE IOPARAM = RECORD
47                    OPERATION: IOOPERATION;
48                    STATUS: IORESULT;
49                    ARG: IOARG
50                  END;
51
52   TYPE TASKKIND = (INPUTTASK, JOBTASK, OUTPUTTASK);
53
54   TYPE ARGTAG =
55        (NILTYPE, BOOLTYPE, INTTYPE, IDTYPE, PTRTYPE);
56
57   TYPE POINTER = @INTEGER;
58
59   TYPE PASSPTR = @PASSLINK;
```

```
60
61  TYPE OPTION = (LISTOPTION,SUMMARYOPTION,TESTOPTION,CHECKOPTION,
62                CODEOPTION,NUMBEROPTION,XREFOPTION,DUMPOPTION);
63
64  TYPE PASSLINK = RECORD
65    OPTIONS: SET OF OPTION;
66    LABELS, BLOCKS, CONSTANTS, XJP_OFFSETS: INTEGER;
67    RESETPOINT: POINTER;
68    TABLES: POINTER;
69    INTERFACE: POINTER;
70    END;
71
72
73  TYPE ARGTYPE = RECORD
74    CASE TAG: ARGTAG OF
75      NILTYPE, BOOLTYPE: (BOOL: BOOLEAN);
76      INTTYPE: (INT: INTEGER);
77      IDTYPE: (ID: IDENTIFIER);
78      PTRTYPE: (PTR: PASSPTR)
79      END;
80
81  CONST MAXARG = 10;
82  TYPE ARGLIST = ARRAY (.1..MAXARG.) OF ARGTYPE;
83
84  TYPE ARGSEQ = (INP, OUT);
85
86  TYPE PROCRESULT =
87    (TERMINATED, OVERFLOW, POINTERERROR, RANGEERROR, VARIANTERROR,
88     HEAPLIMIT, STACKLIMIT, CODELIMIT, TIMELIMIT, CALLERROR);
89
90  PROCEDURE READ(VAR C: CHAR);
91  PROCEDURE WRITE(C: CHAR);
92
93  PROCEDURE OPEN(F: FILE; ID: IDENTIFIER; VAR FOUND: BOOLEAN);
94  PROCEDURE CLOSE(F: FILE);
95  PROCEDURE GET(F: FILE; P: INTEGER; VAR BLOCK: UNIV PAGE);
96  PROCEDURE PUT(F: FILE; P: INTEGER; VAR BLOCK: UNIV PAGE);
97  FUNCTION LENGTH(F: FILE): INTEGER;
98
99  PROCEDURE MARK(VAR TOP: INTEGER);
100 PROCEDURE RELEASE(TOP: INTEGER);
101
102 PROCEDURE IDENTIFY(HEADER: LINE);
103 PROCEDURE ACCEPT(VAR C: CHAR);
104 PROCEDURE DISPLAY(C: CHAR);
105
106 PROCEDURE NOTUSED;
107 PROCEDURE NOTUSED2;
108 PROCEDURE NOTUSED3;
109 PROCEDURE NOTUSED4;
110 PROCEDURE NOTUSED5;
111 PROCEDURE NOTUSED6;
112
113 PROCEDURE NOTUSED7;
114
115 PROCEDURE NOTUSED8;
116
117 PROCEDURE NOTUED9;
118
119 PROCEDURE NOTUUED10;
```

```
120
121 PROCEDURE RUN(ID: IDENTIFIER; VAR PARAM: ARGLIST;
122              VAR LINE: INTEGER; VAR RESULT: PROGRESULT);
123
124 PROCEDURE BREAKPNT(LINE: INTEGER);
125
126 PROGRAM P(PARAM: LINE);
127
128
129
130
131 "#####################################################
132 # NCPASCAL(VAR OK: BOOLEAN; SOURCE, DEST, OBJECT: IDENTIFIER) #
133 #####################################################"
134
135
136 "INSERT PREFIX HERE"
137
138 "THE PARAMETERS OF THE COMPILER PASSES
139 HAVE THE FOLLOWING MEANING:
140
141    LIST(.1.) : BOOLEAN    (COMPILATION OK)
142    LIST(.2.) : POINTER    (HEAP POINTER)
143    LIST(.3.) : INTEGER    (CODE LENGTH) "
144
145 TYPE CHARSET = SET OF CHAR;
146
147 TYPE PASSES = 0..127;
148
149 VAR
150
151 OK, PAST_EOM: BOOLEAN; SOURCE, DEST, OBJECT: ARGTYPE;
152 CODELENGTH: INTEGER;
153 WHERE: (NOWHERE, ONDISK);
154 REPORT: (MAIN, OUTP);
155 I: INTEGER;
156
157 LIST: ARGLIST;
158
159 PARAMPTR: 1..LINELENGTH;
160 TESTPASS: SET OF PASSES;
161 NOSOURCE: BOOLEAN;
162
163
164 PROCEDURE INITWRITE;
165 BEGIN
166    IDENTIFY('PASCAL.:(:10:)');
167    REPORT:= MAIN;
168 END;
169
170 PROCEDURE WRITECHAR(C: CHAR);
171 BEGIN
172    IF REPORT <> MAIN
173       THEN WRITE(C);
174    DISPLAY(C);
175 END;
176
177 PROCEDURE WRITETEXT(TEXT: LINE);
178 CONST NUL = '(:0:)';
179 VAR I: INTEGER; C: CHAR;
```

```
180  BEGIN
181    I:= 1; C:= TEXT(.1.);
182    WHILE C <> NUL DO
183    BEGIN
184      WRITECHAR(C);
185      I:= I + 1; C:= TEXT(.I.);
186    END;
187  END;
188
189  PROCEDURE WRITEINT(INT, LENGTH: INTEGER);
190  VAR NUMBER: ARRAY (.1..6.) OF CHAR;
191    DIGIT, REM, I: INTEGER;
192  BEGIN
193    DIGIT:= 0; REM:= INT;
194    REPEAT
195      DIGIT:= DIGIT + 1;
196      NUMBER(.DIGIT.):=
197        CHR(ABS(REM MOD 10) + ORD('0'));
198      REM:= REM DIV 10;
199    UNTIL REM = 0;
200    FOR I:= 1 TO LENGTH - DIGIT - 1 DO
201      WRITECHAR(' ');
202    IF INT < 0 THEN WRITECHAR('-')
203                ELSE WRITECHAR(' ');
204    FOR I:= DIGIT DOWNTO 1 DO
205      WRITECHAR(NUMBER(.I.));
206  END;
207
208  PROCEDURE WRITEID(ID: IDENTIFIER);
209  VAR I: INTEGER; C: CHAR;
210  BEGIN
211    FOR I:= 1 TO IDLENGTH DO
212    BEGIN
213      C:= ID(.I.);
214      IF C <> ' ' THEN WRITECHAR(C);
215    END;
216  END;
217
218  PROCEDURE CONVRESULT(RESULT: PROGRESULT; VAR ID: IDENTIFIER);
219  BEGIN
220    CASE RESULT OF
221      TERMINATED:    ID:= 'TERMINATED  ';
222      OVERFLOW:      ID:= 'OVERFLOW    ';
223      POINTERERROR:  ID:= 'POINTERERROR';
224      RANGEERROR:    ID:= 'RANGEERROR  ';
225      VARIANTERROR:  ID:= 'VARIANTERROR';
226      HEAPLIMIT:     ID:= 'HEAPLIMIT   ';
227      STACKLIMIT:    ID:= 'STACKLIMIT  ';
228      CODELIMIT:     ID:= 'CODELIMIT   ';
229      TIMELIMIT:     ID:= 'LOADERROR   ';
230      CALLERROR:     ID:= 'CALLERROR   ';
231    END;
232  END;
233
234  PROCEDURE WRITERESULT
235    (ID: IDENTIFIER; LINE: INTEGER; RESULT: PROGRESULT);
236  VAR ARG: IDENTIFIER;
237  BEGIN
238    WRITECHAR(NL);
239    WRITEID(ID);
```

```
240   WRITETEXT(': LINE (:0:)');
241   WRITEINT(LINE, N);
242   WRITECHAR(' ');
243   CONVRESULT(RESULT, ARG);
244   WRITEID(ARG);
245   WRITECHAR(NL);
246   OK:= (RESULT = TERMINATED);
247 END;
248
249 PROCEDURE ERROR(TEXT: LINE);
250 BEGIN
251   INITWRITE;
252   WRITETEXT(TEXT);
253   OK:= FALSE;
254 END;
255
256 FUNCTION NEXTCHAR: CHAR;
257 BEGIN
258   PARAMPTR:=PARAMPTR+1;
259   NEXTCHAR:=PARAM[PARAMPTR];
260 END;
261
262 PROCEDURE SCANPARMS;
263 VAR SKIPS,ALPHAS,NUMERICS: SET OF CHAR;
264     ID: IDENTIFIER;
265     C: CHAR;
266     N,IDPTR: INTEGER;
267 BEGIN
268   SKIPS:=[',',' '];
269   ALPHAS:=[]; FOR C:='A' TO 'Z' DO ALPHAS:=ALPHAS OR [C];
270   NUMERICS:=[]; FOR C:='0' TO '9' DO NUMERICS:=NUMERICS OR [C];
271   C:=PARAM[1]; PARAMPTR:=1;
272   WHILE C<>NL DO BEGIN
273     WHILE C IN SKIPS DO C:=NEXTCHAR;
274     IF C IN ALPHAS THEN BEGIN
275       ID:='         '; IDPTR:=1;
276       WHILE C IN (ALPHAS OR NUMERICS) DO
277         IF IDPTR<=IDLENGTH THEN BEGIN
278           ID[IDPTR]:=C; IDPTR:=IDPTR+1; C:=NEXTCHAR;
279           END ELSE C:=NEXTCHAR;
280       "IF ID[1]='H' THEN OPTIONS:=OPTIONS-[ID[2]]
281        ELSE OPTIONS:=OPTIONS OR [ID[1]];"
282       IF ID='NS'       THEN NOSOURCE:=TRUE;
283       IF ID='ALL'
284         THEN FOR N := 1 TO 7 DO TESTPASS := TESTPASS OR [N];
285       END
286     ELSE IF C IN NUMERICS THEN BEGIN
287       N:=0;
288       WHILE C IN NUMERICS DO BEGIN
289         N:=N*10+ORD(C)-ORD('0'); C:=NEXTCHAR;
290         END;
291       TESTPASS:=TESTPASS OR [N];
292       END
293     ELSE IF C<>NL THEN C:=NEXTCHAR;
294     END "WHILE";
295 END;
296
297 PROCEDURE CHECKIO;
298 VAR C:CHAR;
299 BEGIN
```

```
300     "COMPLETE SOURCE TEXT INPUT/OUTPUT:"
301     IF NOT PAST_EOM THEN REPEAT READ(C) UNTIL C=EM;
302     WRITE(EM);
303     END;
304
305   PROCEDURE INITIALIZE;
306   BEGIN
307     WRITECHAR (NL);
308     TESTPASS:=[]; NOSOURCE:=FALSE;
309     WITH LIST(.1.) DO
310     BEGIN TAG:= BOOLTYPE; BOOL:= FALSE END;
311     WITH LIST(.2.) DO
312     BEGIN TAG:=PTRTYPE; PTR:=NIL. END;
313     WITH LIST(.3.) DO
314     BEGIN TAG:= INTTYPE; INT:= 0 END;
315     SCANPARMS;
316     INITWRITE;
317     WRITETEXT (
318     'HEPASCAL COMPILER (MICROENGINE P-CODE) (:10:)(:0:)');
319   END;
320
321   PROCEDURE CALLPASS(PASSNO: INTEGER; ID: IDENTIFIER);
322   VAR LINE: INTEGER; RESULT: PROGRESULT;
323   BEGIN
324     LIST(.1.).BOOL:= FALSE;
325     RUN(ID, LIST, LINE, RESULT);
326     IF RESULT <> TERMINATED THEN
327     BEGIN
328       REPORT:= OUTP;
329       WRITERESULT(ID, LINE, RESULT);
330     END ELSE
331     BEGIN
332       OK:= LIST(.1.).BOOL;
333       CODELENGTH:= LIST(.3.).INT;
334       IF NOT OK THEN
335       ERROR('COMPILATION ERRORS(:10:)(:0:)');
336       IF OK
337       THEN BEGIN
338         WRITEID(ID); WRITETEXT(' OK(:10:)(:0:)');
339       END;
340     END;
341   END;
342
343
344   PROCEDURE RUN_HMEM (PASS: INTEGER);
345   VAR
346     L:INTEGER;
347     R: PROGRESULT;
348   BEGIN
349     LIST[10].TAG := INTTYPE;
350     LIST[10].INT := PASS;
351     IF (PASS = 5) OR
352        (PASS = 6)
353     THEN RUN ('MEMMEM     ', LIST, L, R);
354     CASE PASS OF
355     1: WRITETEXT ('NO MEMMEM 1(:10:)(:0:)');
356     2: WRITETEXT ('NO MEMMEM 2(:10:)(:0:)');
357     3: WRITETEXT ('NO MEMMEM 3(:10:)(:0:)');
358     4: WRITETEXT ('NO MEMMEM 4(:10:)(:0:)');
359     5: WRITETEXT ('MEMMEM  5 DONE(:10:)(:0:)');
```

```
360        6: WRITETEXT ('MEMBER 6  DONE(:10:)(:0:)'');
361        7: WRITETEXT ('NO MEMBER 7(:10:)(:0:)'');
362      END;
363    END;
364
365
366  BEGIN
367    INITIALIZE;
368    IF NOSOURCE THEN CALLPASS(1,'MEPASS1N      ')
369    ELSE CALLPASS(1,'MEPASS1       ');
370    PAST_EOM:= LIST(5).BOOL;
371    IF (OK AND (1 IN TESTPASS))
372    THEN RUN_MNEM (1);
373
374    IF OK THEN CALLPASS(2,'MEPASS2     ');
375    IF (OK AND (2 IN TESTPASS))
376    THEN RUN_MNEM(2);
377
378    IF OK THEN CALLPASS(3,'MEPASS3     ');
379    IF (OK AND (3 IN TESTPASS))
380    THEN RUN_MNEM(3);
381
382    IF OK THEN CALLPASS(4,'MEPASS4     ');
383    IF (OK AND (4 IN TESTPASS))
384    THEN RUN_MNEM(4);
385
386    IF OK THEN CALLPASS(5,'MEPASS5     ');
387    IF (OK AND (5 IN TESTPASS))
388    THEN RUN_MNEM(5);
389
390    IF OK THEN CALLPASS(6,'MEPASS6     ');
391    IF (OK AND (6 IN TESTPASS))
392    THEN RUN_MNEM(6);
393
394    IF OK THEN CALLPASS(7,'MEPASS7     ');
395    IF (OK AND (7 IN TESTPASS))
396    THEN RUN_MNEM(7);
397  END.
```

Cross-reference / address column:

```
360   177
361   177

367   305
368   161      321
369   321
370   151      157
371   151      160
372   344

374   151      321
375   151      160
376   344

378   151      321
379   151      160
380   344

382   151      321
383   151      160
384   344

386   151      321
387   151      160
388   344

390   151      321
391   151      160
392   344

394   151      321
395   151      160
396   344
```

(75)

CROSS REFERENCE   * IS DEF   = IS ASG

-X-
IJP_OFFSETS   66*
XREFOPTION   62

END XREF   191 IDENTIFIERS   605 TOTAL REFERENCES
25 COLLISIONS.

APPENDIX D

SOURCE CODE CHANGES TO MCPASCAL PASSES 1-5

```
45     CONST IDLENGTH = 12;
46     TYPE IDENTIFIER = ARRAY (.1..IDLENGTH.) OF CHAR;
47     TYPE LINE = ARRAY(.1..132.) OF CHAR;
48     TYPE   POINTER = @ INTEGER;
49     OPTION = LISTOPTION..VARNTCHECKOPTION;              "VARNT*"
50     PASSPTR = @PASSLINK;
51     PASSLINK =
52       RECORD
53        OPTIONS: SET OF OPTION;
                     "MICROENGINE"
54        LABELS, BLOCKS, CONSTANTS, XJP_OFFSETS: INTEGER;
55        RESETPOINT: FULLWORD;
56        TABLES: POINTER;
57        INTERFACE: POINTER    "MICROENGINE"
58       END;
```

--- PASS 1 CODE MODIFICATIONS ---

```
36     CONST IDLENGTH = 12;
37     TYPE IDENTIFIER = ARRAY (.1..IDLENGTH.) OF CHAR;
38
39     TYPE   POINTER = @ INTEGER;
40     OPTION = LISTOPTION..VARNTCHECKOPTION;              "VARNT*"
41     PASSPTR = @PASSLINK;
42     PASSLINK =
43       RECORD
44        OPTIONS: SET OF OPTION;
                     "MICROENGINE"
45        LABELS, BLOCKS, CONSTANTS, XJP_OFFSETS: INTEGER;
46        RESETPOINT: FULLWORD;
47        TABLES: POINTER;
48        INTERFACE: POINTER    "MICROENGINE"
49       END;
```

--- PASS 2 CODE MODIFICATIONS ---

```
35      CONST IDLENGTH = 12;
36      TYPE IDENTIFIER = ARRAY (.1..IDLENGTH.) OF CHAR;
37
38      TYPE   POINTER = @ INTEGER;
39      OPTION = LISTOPTION..VARNTCHECKOPTION;          "VARNT*"
40      PASSPTR = @PASSLINK;
41      PASSLINK =
42        RECORD
43          OPTIONS: SET OF OPTION;
                        "MICROENGINE"
44          LABELS, BLOCKS, CONSTANTS, XJP_OFFSETS: INTEGER;
45          RESETPOINT: FULLWORD;
46          TABLES: POINTER;
47          INTERFACE: POINTER   "MICROENGINE"
48        END;
```

```
145     SUB2=81;       INDEX2=82;      REAL2=83;       STRING2=84;
146     LCONST2=85;    MESSAGE2=86;    NEW_LINE2=87;   FWD_DEF2=88;
147     CHK_TYPE2=89;  PROCF_DEF2=90;  UNDEF2=91;      PEND2=92;
148     CASE_JUMP2=93;
149     "MANAGER:"                              "MGR*"
150     MANAGER2=94;      FROM2=95;       REFERS2=96;
151     NIL2=97;                                "MGR*"
152     VTAG_DEF2=98;    VPART_END2=99;   VVARNT_END2=100;   "VARNT*"
153     VVARIANT2=101;   "VARNT*"         DUPTOS2=102; "MICROENGINE"
154     "DUPTOS2 DUPLICATES TOP-OF-STACK WORD"
155
156     "OTHER CONSTANTS"
157
158     SPLIT_SET_LENGTH = "%IF SRC_16" 8; "%END" "WORDS/SPLIT SET"
159                        "%IF SRC_32" 4; "%END"
160     NOUN_MAX=999;
```

```
1881    FUNCTION1: FUNCTION_;
1882    GE1: BINARY(GE2);
1883    GT1: BINARY(GT2);
1884    INCLUDE1: BINARY(INCLUDE2);
1885    INIT_NAME1:BEGIN INIT_NAME; PUTO(DUPTOS2) END;"MICROENGINE"
1886    INITS_DEF1: INITS_DEF;
1887    INITS_END1: POP_LEVEL;
1888    INIT1: CALL(INIT2);
1889    INTEGER1: INDEX(XINTEGER);
1890    INTF_ID1: INTF_ID;
```

--- PASS 3 CODE MODIFICATIONS ---

```
38      CONST IDLENGTH = 12;
39      TYPE IDENTIFIER = ARRAY (.1..IDLENGTH.) OF CHAR;
40
41      TYPE    POINTER = @ INTEGER;
42      OPTION = LISTOPTION..VARNTCHECKOPTION;          "VARNT*"
43      PASSPTR = @PASSLINK;
44      PASSLINK =
45        RECORD
46          OPTIONS: SET OF OPTION;
                     "MICROENGINE"
47          LABELS, BLOCKS, CONSTANTS, XJP_OFFSETS: INTEGER;
48          RESETPOINT: FULLWORD;
49          TABLES: POINTER;
50          INTERFACE: POINTER     "MICROENGINE"
51        END;
```

```
130     CASE_JUMP1=93;
131     MANAGER1=94;      FROM1=95;        REFERS1=96;        "MGR*"
132     NIL1=97;                                   "MGR*"
133     VTAG_DEF1=98; VPART_END1=99; VVARNT_END1=100; VVARIANT1=101;
134     DUPTOS1 = 102;   "MICROENGINE"
135
136     "OUTPUT OPERATORS"
137
138     EOM2=1;          BODY2=2;        BODY_END2=3;      ADDRESS2=4;
139     RESULT2=5;       STORE2=6;       CALL_PROC2=7;     CONSTPARM2=8;
```

```
150     VCOMP2=49;       RCOMP2=50;      SUB2=51;          LCONST2=52;
151     MESSAGE2=53;     NEW_LINE2=54;   CHK_TYPE2=55;     SAVEPARM2=56;
152     CALL_GEN2=57;    NOT2=58;        UNDEF2=59;        RANGE2=60;
153     REFERS2=61;          "MGR*"
154     VVARIANT2=62;   "VARNT*"              DUPTOS2=63;"MICROENGINE"
155
156     "STANDARD SPELLING/NOUN INDICES"
157
158     XUNDEF=0;        XFALSE=1;       XTRUE=2;          XINTEGER=3;
159     XBOOLEAN=4;      XCHAR=5;        XQUEUE=6;         XABS=7;
```

--- PASS 4 CODE MODIFICATIONS ---

```
1226    PROCEDURE BODY;
1227    BEGIN
1228      WITH STACK(.T.)@ DO BEGIN
1229        VAR_SIZE:=CURRENT_DISP;
1230        IF INITIAL_ENTRY THEN BEGIN
1231          INITIAL_ENTRY:=FALSE;
1232          COMPVAR_LENGTH:=CURRENT_DISP "SAVE COMP VAR LENGTH;
1233          CURRENT_DISP:=0 "INITIAL STMNT IS VARIABLE-LESS";
1234          "REMOVED FOR MICROENGINE"
1235          "PUT5(BODY2,RMODE,RDISP,0,0,STACK_SIZE)"
1236        END;    "ELSE"  "REMOVED FOR MICROENGINE"
1237          PUT5(BODY2,RMODE,RDISP,PARM_SIZE,VAR_SIZE,STACK_SIZE)
1238      END
1239    END;
```

```
1468    CHK_TYPE1: BEGIN PUTO(CHK_TYPE2); PUT_TYPE END;
1469    CLASS1: COMP_DEF(PACKED_CLASS);
1470    CPARMLIST1: CPARM_LIST;
1471    DEF_LABEL1: IGNORE1(DEF_LABEL2);
1472    DIV1: PUTO(DIV2);
1473    DUPTOS1: PUTO(DUPTOS2);    "MICROENGINE"
1474    EMPTY_SET1: PUTO(EMPTY_SET2);.
1475    EOM1: EOM;
1476    ENUM_DEF1: ENUM_DEF;
1477    EQ1: PUTO(EQ2);
```

```
42    "************** MICROENGINE ******************"
43    CONST MAXINTFAC = 14; "14 INTERFACE SEGMENTS MAX."
44    TYPE
45       IFPTR = @IFINFO;   "INTERFACE INFORMATION POINTER"
46       IFINFO = RECORD    "INTERFACE INFORMATION RECORD"
47          INTERFACES: INTEGER; "NO. OF INTERFACE SEGMENTS"
48          INTERFACESIZES: "PROCESS ENTRY RTNS IN EACH INTERFACE"
49             ARRAY[1..MAXINTFAC] OF INTEGER;
50          END;
51    "********************************************************"
52
53    PASSPTR = @PASSLINK;
54    PASSLINK =
55      RECORD
56        OPTIONS: SET OF OPTION;
                 "MICROENGINE"
57        LABELS, BLOCKS, CONSTANTS, XJP_OFFSETS: INTEGER;
58        RESETPOINT: FULLWORD;
59        TABLES: POINTER;
60        INTERFACE: IFPTR        "MICROENGINE"
61      END;
```

```
113   VCOMP1=49;      RCOMP1=50;      SUB1=51;       LCONST1=52;
114   MESSAGE1=53;    NEW_LINE1=54;   CHK_TYPE1=55;  SAVEPARM1=56;
115   CALL_GEN1=57;   NOT1=58;        UNDEF1=59;     RANGE1=60;
116   REFERS1=61;                        "MGR*"
117   VVARIANT1=62;   "VARNT*"              DUPTOS1=63; "MICROENGINE"
118
119   "OUTPUT OPERATORS"
120
121   PUSHCONST2=0;   PUSHVAR2=1;     PUSHIND2=2;    PUSHADDR2=3;
122   FIELD2=4;       INDEX2=5;       POINTER2=6;    VARIANT2=7;
```

```
129   CASEJUMP2=32;   INITVAR2=33;    CALL2=34;      ENTER2=35;
130   RETURN2=36;     POP2=37;        NEWLINE2=38;   ERR2=39;
131   LCONST2=40;     MESSAGE2=41;    INCREMENT2=42; DECREMENT2=43;
132   PROCEDURE2=44;  INIT2=45;       PUSHLABEL2=46; CALLPROG2=47;
133   EOM2=48;            DUPTOS2=49;           "MICROENGINE"
134
135   "CONTEXT"
136
137   FUNC_RESULT=1; ENTRY_VAR=2;     VARIABLE=3;    VAR_PARM=4;
138   UNIV_VAR=5;    CONST_PARM=6;    UNIV_CONST=7;  FIELD=8;
```

--- PASS 5 CODE MODIFICATIONS ---

```
258
259      VARNT_TAGSET, VARNT_DISP: INTEGER;        "VARNT*"
260
261    INTFACEPTR: IFPTR;        "MICROENGINE"
262    INTFACE: IFINFO;        "MICROENGINE"
263
264
265    "############################"
266    "COMMON TEST OUTPUT MECHANISM"
267    "############################"



300    PROCEDURE NEXT_PASS (VAR LINK: PASSPTR);
301    BEGIN
302      LINK@.INTERFACE := INTFACEPTR;        "MICROENGINE"
303      IF WORDS_OUT > 0
304      THEN IF PAGES_OUT > FILE_LENGTH(OUTFILE)
305          THEN FILE_LIMIT
306          ELSE PUT(OUTFILE, PAGES_OUT, PAGE_OUT);
307      WITH PARAM(.1.) DO
308        BEGIN TAG:= BOOLTYPE; BOOL:=OK END;
309      WITH PARAM(.2.) DO
310        BEGIN TAG:= PTRTYPE; PTR:= LINK END;
311      WITH PARAM(.4.) DO
312        BEGIN TAG:= INTTYPE;  INT:= PAGES_OUT  END;
313    END;



461      PROCEDURE INITIALIZE;
462      BEGIN
463        DONE:=FALSE;
464        ACTIVE_VARNT:= FALSE;                "VARNT*"
465        INIT_PASS(INTER_PASS_PTR);
466        WITH INTER_PASS_PTR@ DO BEGIN
467          DEBUG:=TESTOPTION IN OPTIONS;
468          IF DEBUG THEN PRINTFF;
469          XJP_OFFSETS := 0      "MICROENGINE"
470        END;
471
472    "****************** MICROENGINE *********************"
473      NEW(INTFACEPTR); "GET SPACE FOR INTERFACE INFO RECORD"
474      MARK(INTER_PASS_PTR@.RESETPOINT); "GET @ NEW HEAPTOP"
475      INTFACEPTR@.INTERFACES := 0; "NO INTERFACES DEFINED YET"
476    "***************************************************"
477
478      ARITHMETIC:=(.INT_KIND,REAL_KIND.);
479      INDEXS:=(.INT_KIND,BOOL_KIND,CHAR_KIND,ENUM_KIND.);
```

--- PASS 5 CODE MODIFICATIONS ---

```
679     PROCEDURE ADDRESS;
680     BEGIN
681       WITH T@ DO
682         IF CLAS=VALUE THEN BEGIN
683           CASE STATE OF
684             DIRECT: BEGIN
685               IF MODE=SCONST_MODE THEN ADDR_ERROR
686               ELSE PUT2(PUSHADDR2,MODE,DISP);
687               "NEXT 'IF...THEN' DELETED FOR MICROENGINE"
688               "IF (KIND=SYSCOMP_KIND) & (CONTEXT<> FROM_VAR)"
689                 "THEN PUT1(FIELD2,LENGTH)" "OFFSET"
690             END;                          '
691             INDIRECT: PUT3(PUSHVAR2,WORD_TYP,MODE,DISP);
692             ADDR: ;
693             EXPRESSION: ADDR_ERROR
694           END;
695           STATE:=ADDR
696         END ELSE ADDR_ERROR
697     END;



901     PROCEDURE CASE_LIST;
902     VAR I,MIN,MAX:INTEGER; L:DISPLACEMENT;
903     BEGIN
904       POP "SELECTOR";
905       DEF_LABEL;
906       READ_IFL(MIN); READ_IFL(MAX); PUT2(CASEJUMP2,MIN,MAX);
907
908       "MICROENGINE"
909       "COUNT HOW MANY BYTES THIS CASE JUMP WILL ADD TO CNSTNT
910        BLOCK. MIN, MAX, OFFSET FOR EACH CASE, ARE 1 WORD EACH"
911       WITH INTER_PASS_PTR@ DO
912         XJP_OFFSETS:=XJP_OFFSETS+((MAX-MIN+1)+2)*WORDLENGTH;
913
914       FOR I:=MIN TO MAX DO BEGIN
915         READ_IFL(L); PUT_ARG(L)
916       END;
917       DEF_LABEL
918     END;
```

--- PASS 5 CODE MODIFICATIONS ---

```
968      PROCEDURE PROG_CALL;
969      VAR INTF_LENGTH:INTEGER;
970      BEGIN
971        READ_IFL(INTF_LENGTH);
972
973   "**************** MICROENGINE ******************"
974        IF INTF_LENGTH <> 0
975        THEN WITH INTFACEPTR@ DO
976          BEGIN
977          INTERFACES := SUCC(INTERFACES); "FOUND AN INTERFACE"
978          "SAVE NUMBER OF ENTRY POINTS FOR THIS INTERFACE"
979          INTERFACESIZES[INTERFACES] := INTF_LENGTH DIV
                                                  WORDLENGTH;
980          END;
981   "******************************************"
982
983        PUTO(CALLPROG2);  PUT1(POP2,INTF_LENGTH);
984        POP
985      END;




1216     PROCEDURE SUB;
1217     VAR MIN,MAX,SIZE: INTEGER;
1218     BEGIN
1219       "SUBSCRIPT" VALUE_;
1220       READ_IFL(MIN); READ_IFL(MAX); READ_IFL(SIZE);
1221
1222   "************* MICROENGINE ******************"
1223       "INDEX OPERATOR CHANGED TO ACCOMODATE MICROENGINE
1224        BYTE POINTERS. TYPE ARGUMENT (LAST ARGUMENT) WAS ADDED
1225        SO PASS 6 KNOWS TO GEN BYTE POINTERS (2 WORDS ON STACK)
1226        INSTEAD OF WORD POINTERS (1 WORD ON STACK). "
1227       PUT4(INDEX2,MIN,MAX,SIZE,T@.KIND);
1228   "******************************************"
1229
1230       PUSH; T@:=UNDEF_EXPR; "INDEX" TYPE_;
1231       IF COMPATIBLE THEN "OK";
1232       POP; POP;
1233       "ELEMENT" TYPE_;
1234       WITH T@ DO
1235         IF KIND=SYSCOMP_KIND THEN PUT1(FIELD2,LENGTH) "OFFSET"
1236     END;
```

--- PASS 5 CODE MODIFICATIONS ---

```
1277      CASE_LIST1: CASE_LIST;
1278      CHK_TYPE1: CHK_TYPE;
1279      CONSTPARM1: CONSTPARM(FALSE);
1280      DEF_LABEL1: DEF_LABEL;
1281      DIV1: DIV_MOD(DIV2);
1282      DUPTOS1: PUTO(DUPTOS2);      "MICROENGINE"
1283      EMPTY_SET1: EMPTY_SET;
1284      EOM1: EOM;
1285      EQ1: EQUALITY(EQUAL);
1286      FALSEJUMP1: FALSE_JUMP;
```

ON THE ADAPTABILITY OF MULTIPASS PASCAL COMPILERS
TO VARIANTS OF (PASCAL) P-CODE MACHINE ARCHITECTURES

by

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B. S., Pittsburg State University, Pittsburg, Kansas, 1979

———————————

AN ABSTRACT OF A MASTER'S REPORT

submitted in partial fulfillment of the

requirements for the degree

MASTER OF SCIENCE

Department of Computer Science

KANSAS STATE UNIVERSITY
Manhattan, Kansas

1981

ABSTRACT

This report is a description of an effort to modify the code-generation portion of a multipass Concurrent Pascal (CPascal) compiler so that it will produce object code for a different machine. The original compiler generated P-code (Pascal stack machine code) for a virtual machine. The modified version will generate code for the Pascal Microengine, a microcomputer whose instruction set is similar to virtual P-code. The similarity of instruction sets made it appear that the required modifications would be straightforward, and that the high-level language constructs of Concurrent Pascal (monitors, for example) would easily map onto the Microengine's architecture. However, as the project progressed the architectural differences became a significant obstacle. This report contains a description of the organization and architecture of the two machines and detail on the changes made to the original Concurrent Pascal compiler. It also contains a description of the compilation problems caused by the differences between the machines and how those difficulties were resolved.