AN ALGEBRAIC APPROACH TO COMPUTING INVERSE LAPLACE TRANSFORMS OF RATIONAL FUNCTIONS/

by

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Chapter I

Introduction

The electrical engineer must often solve linear, constant-coefficient, differential equations, especially in such fields as circuit analysis and control theory. The usual method of solution involves use of the Laplace transform. Although this method of solution is well-known [1], computing solutions to such equations often involves many tedious calculations. The most difficult step in this method of solution is performing the inverse Laplace transform of a rational function. The object of this thesis is to describe an algorithm for solving large problems of this kind.

Solving such equations is not necessarily a numerical analysis problem. We are not trying to approximate the solution to a differential equation. We already know the general form of the solution, and are simply seeking coefficients for the solution. Thus the problem is actually algebraic, rather than analytic.

There are three distinct steps in computing the inverse Laplace transform of a rational function. They are: factoring the denominator polynomial, finding the partial fraction expansion of the rational function, and computing the inverse Laplace transform of each of these partial fractions. Except for the factoring, these are not analytic problems, and even the factoring algorithm presented here makes use of an algebraic technique to speed up the finding of multiple roots.

Special concern was devoted to the partial fraction expansion portion of the algorithm in order to reduce the number of calculations. The method presented herein is based on one developed elsewhere [4] and modified to further reduce the number of calculations. The effort to minimize the number of calculations required constitutes the main contribution of this research.

As the title declares, this thesis concerns itself with computing the inverse Laplace transform of a rational function, but not with obtaining the rational function in the first place. The algorithm consists of four parts. The first part of the algorithm accepts a rational function as input and outputs the rational function with the denominator expressed as a product of irreducible factors. The next part accepts as input a rational function with a factored denominator. It outputs the partial fraction expansion of the given rational function. The third stage accepts this expansion and computes the inverse Laplace transform. The final stage evaluates the inverse transform over a desired interval and plots a graph, if desired.

Chapter II

Some Preliminaries

We first review the definition of the Laplace transform and its usefulness in solving linear, constant-coefficient, differential equations. Given a real-valued function f(t) of a real variable, its one-sided Laplace transform, F(s), is given by [1]:

$$F(s) = \mathcal{L}\{f(t)\} = \int_0^{+\infty} f(t)e^{-st}dt$$

where s is complex and hence F(s) is, in general complex-valued. This transform has two important properties which can be used to transform a differential equation in t into an algebraic equation in s. Both of these properties are easy to derive and proofs are given elsewhere. These properties are [1]:

<u>Linearity.</u> Suppose $f_1(t)$ and $f_2(t)$ are such that their Laplace transforms, $F_1(s)$ and $F_2(s)$, exist. Also, let $c_1, c_2 \in \mathbf{R}$. Then

$$\mathcal{L}\{c_1f_1(t) + c_2f_2(t)\} = c_1F_1(s) + c_2F_2(s).$$

Forward Derivative Property. Suppose that f(t) is (n-1) times continuously differentiable and that the *n*th derivative, $f^{(n)}(t)$, is such that its Laplace transform exists. Then

$$\mathcal{L}\{f^{(n)}(t)\} = s^n F(s) - \sum_{j=1}^n s^{n-j} f^{(j-1)}(0).$$

Given an nth-order linear differential equation with constant coefficients

$$\sum_{j=0}^{n} a_j f^{(j)}(t) = b(t),$$

and the initial conditions, $f(0) = c_0, f^{(1)}(0) = c_1, \ldots, f^{(n-1)}(0) = c_{n-1}$, we can take the Laplace transform of both sides of the equation and end up with [1]

$$F(s) = \frac{B(s) + \sum_{i=1}^{n} c_{i-1} \sum_{j=i}^{n} a_j s^{j-i}}{\sum_{j=0}^{n} a_j s^j},$$
(2.1)

where $B(s) = \mathcal{L}{b(t)}$.

We see that if b(t) is such that B(s) is a rational function, then F(s) will also be a rational function. Thus, the problem is reduced to finding the inverse Laplace transform of F(s). What might be done when b(t) is such that B(s) is <u>not</u> a rational function will be discussed later.

The method employed here to find the inverse Laplace transform of (2.1) is, first, to factor the denominator; next, to expand the rational function into partial fractions; and finally, by applying the linearity property of the inverse Laplace transform, to find the inverse Laplace transform of each of the partial fractions. The algorithms presented in the next three chapters accomplish each of these steps by using elementary concepts of polynomial ring theory, and recursion where possible.

Chapter III

Factoring

The form of the inverse Laplace transform of a rational function depends on the factors in the denominator polynomial. The problem of factoring polynomials is an important area of numerical analysis. The algorithm presented here relies on elementary algebraic considerations. It should be pointed out that any other factoring algorithm could be used with the rest of the programs listed in the appendix by providing suitable interfacing software.

The algorithm presented here is based on the root-finding method of D. E. Muller [2], a brief description of which follows. Given a polynomial, p, start with three initial estimates for a root: x_{-2}, x_{-1} , and x_0 . At the *n*th stage $(0 \le n \in \mathbb{N})$, fit a quadratic equation to the points $(x_{n-2}, p(x_{n-2})), (x_{n-1}, p(x_{n-1}))$, and $(x_n, p(x_n))$. Then find the root of this quadratic equation nearest x_n . This root becomes x_{n+1} . Repeat this procedure while n is less than a given upper bound, or subsequent estimates fail to improve by a given small amount. To find other roots, deflate p by the appropriate factor and reapply Muller's method if necessary.

Suppose we are given a polynomial with real coefficients and distinct complex roots. Apply Muller's method to obtain a root. Deflate the polynomial by a factor containing the root just found (and its complex conjugate, if necessary) to obtain a polynomial of smaller degree with real coefficients and distinct complex roots. Repeat this process until the degree of the deflated polynomial is of degree 0, at which point all the roots of the original polynomial have been found.

The above process works quite well if all the roots are distinct. If Muller's method is applied to find a multiple root, it converges more more slowly than it does for a single root. That is, it runs through many more iterations before the difference between successive estimates becomes as small. So we turn to abstract algebra for a way to speed things up. By using the method described below, we can ensure that we will always be searching for distinct roots.

Consider a monic polynomial $p \in \mathbf{R}[x]$, with $r \in \mathbf{C}$ a root of p. Then there exists $q \in \mathbf{C}[x]$ such that p = (x - r)q. Now consider the first derivative of p, p' = q + (x - r)q'. Observe that p'(r) = 0 if and only if q(r) = 0. Thus r appears more than once as a root of p if and only if p'(r) = 0 [3].

Now look at $g_1 = \gcd(p, p') \in \mathbf{R}[x]$. If $\deg(g_1) = 0$ then the roots of p are all distinct. If $\deg(g_1) > 0$ then each root of g_1 is also a root of p and appears more than once as a root of p. Now define $g_n = \gcd(p, p^{(n)})$. If $\deg(g_n) > 0$ then each root of g_n is also a root of p and its multiplicity is at least n + 1.

Since the multiplicity of any root is at most $\deg(p)$ there exists a minimal $k \in \mathbb{N}$ such that $\deg(g_k) = 0$. This k can be found by repeated differentiation of p and application of the Euclidean algorithm to obtain each gcd. The first gcd obtained with degree 0 is g_k .

The only roots of g_{k-1} will be all the roots of p with multiplicity k. These roots can be found using Muller's method on g_{k-1} . These roots will be distinct in g_{k-1} but of multiplicity two in g_{k-2} . But the remaining roots of g_{k-2} will be distinct, and can be found by deflating g_{k-2} by all known roots and then applying Muller's method on the deflated polynomial. Similarly, all roots of p of multiplicity n or higher will be roots of g_{n-1} . Say a root has multiplicity $m \ge n$; this root will have multiplicity m - n + 1 in g_{n-1} . Thus, we can find all roots of p by working backward from g_k to p using Muller's method and deflation. Most importantly, Muller's method is never used to search for a root of multiplicity greater than 1.

It should be pointed out that repeated polynomial divisions occur in the Euclidean algorithm. A loss of precision is inherent in this process. If one cannot make this sacrifice of accuracy of the results in favor of increased rate of convergence, realize that any other factoring algorithm will work with the rest of the programs in

the appendix. The interfacing software would be easy to write.

Chapter IV

Partial Fraction Expansion

One objective in designing a partial fraction expansion algorithm was to minimize the amount of calculation done. Chin and Steiglitz [4] devised an algorithm capable of accomplishing the expansion in N(N-1) multiplications and $\frac{3}{2}N(N-1)$ additions, where N is the degree of the denominator of the given rational function. This algorithm has a disadvantage: it requires use of complex arithmetic. Chin and Steiglitz count complex divisions as equivalent in time to complex multiplications. While this may be true for real divisions and multiplications it certainly is not true for complex ones. Observe that,

$$\frac{a+ib}{c+id} = \frac{(ac+bd)+i(bc-ad)}{c^2+d^2}$$

has 6 real multiplications, 2 real divisions and 3 real additions. Call this 8 multiplications and 3 additions. Furthermore,

$$(a+ib)(c+id) = (ac-bd) + i(bc+ad),$$

requires 4 real multiplications and 2 real additions. Finally, note that

$$(a+ib) + (c+id) = (a+c) + i(b+d),$$

consists of 2 real additions.

Examination of Chin's and Steiglitz's algorithm reveals that the expansion actually involves $\frac{3}{2}N(N-1)$ complex additions, $\frac{1}{2}N(N-1)$ complex multiplications, and $\frac{1}{2}N(N-1)$ complex divisions. Using the above calculations, this results in 6N(N-1) real multiplications and $\frac{11}{2}N(N-1)$ additions. So if complex arithmetic can be avoided and the operation count can be held lower than this, the algorithm will be improved.

It is clear that the reason Chin and Steiglitz chose to work with complex numbers is that a polynomial in $\mathbf{R}[x] \subset \mathbf{C}[x]$, splits in C. This makes the partial fraction expansion algorithm simple to describe and analyze. But $\mathbf{R}[x]$ has the property that an irreducible element is either linear or quadratic. If we use this property we can avoid complex arithmetic and thus reduce the number of calculations at the cost of complicating the algorithm a bit.

The key to adapting the algorithm of Chin and Steiglitz is finding a nice way to generalize the following problem. Let $s, r, K \in \mathbb{C}$, for $s \neq r$ and find $A, A^* \in \mathbb{C}$, such that

$$\frac{1}{x-s} \left[\frac{K}{(x-r)^n} \right] = \frac{A^*}{(x-s)(x-r)^{n-1}} + \frac{A}{(x-r)^n}$$

This problem generalizes to: Let $s, r \in \mathbf{R}[x] \setminus \mathbf{R}$, with gcd(s, r) = 1; and $K \in \mathbf{R}[x]$, with deg(k) < deg(r). Find $A, A^* \in \mathbf{R}[x]$ such that deg(A) < deg(r), and $deg(A^*) < deg(s)$, and

$$\frac{1}{s} \left[\frac{K}{r^n} \right] = \frac{A^*}{sr^{n-1}} + \frac{A}{r^n}.$$
(4.1)

First multiply through by r^{n-1} to reduce the problem to:

$$\frac{1}{s} \left[\frac{K}{r} \right] = \frac{A^*}{s} + \frac{A}{r}.$$
(4.2)

We know that since gcd(s, r) = 1, there exists $a, b \in \mathbf{R}[x]$ such that as + br = 1and so K = K[as + br] = Kas + Kbr. Apply the division algorithm to obtain the following:

$$Ka = rq + A \quad \text{such that} \quad \deg(A) < \deg(r),$$

$$Kb = sq^* + A^* \quad \text{such that} \quad \deg(A^*) < \deg(s).$$
(4.3)

Now write

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$$K = Kas + Kbr = (rq + A)s + (sq^* + A^*)r$$
$$= (q + q^*)rs + As + A^*r.$$

By hypothesis we have $\deg(K) < \deg(r) < \deg(rs)$ and from (4.3) we get $\deg(As) < \deg(rs)$ and $\deg(A^*r) < \deg(rs)$. So $q + q^* = 0$ and $K = As + A^*r$, which yields (4.2) as required.

Now we must determine how to calculate A and A^* in (4.1) [5], and the necessary number of calculations. First consider the following adaption of the Euclidean algorithm. Let $s, r \in \mathbf{R}[x]$. The division algorithm gives:

$$\begin{split} s &= rq_1 + x_1 & \deg(x_1) < \deg(r) \\ r &= x_1q_2 + x_2 & \deg(x_2) < \deg(x_1) \\ x_1 &= x_2q_3 + x_3 & \deg(x_3) < \deg(x_2) \\ \vdots & \vdots \\ x_{n-2} &= x_{n-1}q_n + x_n & \deg(x_n) < \deg(x_{n-1}) \end{split}$$

And also define $x_0 = r$. If, furthermore:

$$a_0 = 0$$

 $a_1 = 1$
 \vdots
 $a_n = a_{n-2} - q_n a_{n-1}$
 $b_n = b_{n-2} - q_n b_{n-1}$

then $a_n s + b_n r = x_n$, for all $n \ge 0$.

Proof: $a_0s + b_0r = r = x_0$. $a_1s + b_1r = s - rq_1 = x_1$. Assume the hypothesis is true for n-2 and n-1 for some $n \in \mathbb{N}$.

$$\begin{aligned} x_n &= x_{n-2} - q_n x_{n-1} \\ &= a_{n-2}s + b_{n-2}r - q_n(a_{n-1}s + b_{n-1}r) \\ &= (a_{n-2} - q_n a_{n-1})s + (b_{n-2} - q_n b_{n-1})r \\ &= a_n s + b_n r, \end{aligned}$$

as claimed.

We can use the above procedure to develop a method for evaluating A and $-A^*$. There are four cases to consider since either s or r can be of degree 1 or 2. Let us begin with the most difficult case:

Case 1: Both s and r are quadratic

We have $s = rq_1 + x$ and $r = x_1q_2 + x_2$. x_2 is a unit so $a_2s + b_2r = -q_2s + (1 + q_1q_2)r = x_2$, and we can get the following.

$$\frac{K}{sr} = \frac{\frac{1}{x_2}K(-q_2s + (1+q_1q_2)r)}{sr}$$
$$= \frac{-\frac{1}{x_2}q_2K}{r} + \frac{\frac{1}{x_2}(1+q_1q_2)K}{s}$$

By the division algorithm, $q_2K = Qr - x_2A$ for some $Q \in \mathbf{R}[x]$. The quotient, Q, does not matter as was shown in the derivation of (4.1). The remainder is the important thing. Now we need to add up all the calculations required to evaluate A.

Table (4.1)Summary of Case 1

To obtain x_1 requires			2 adds,
to obtain q_2 and x_2 requires	2 divs	2 mult	2 adds,
to obtain $q_2 K$ requires		4 mults	2 adds,
to divide by r to get x_2A requires		2 mults	2 adds,
and to evaluate A requires	2 divs		<u>1 add.</u>
Resulting in	4 divs	8 mults	9 adds.

Now we have A. We need $-A^*$ as well. Since we know that $K = As + A^*r$, write $K = k_1x + k_0$, $A = a_1x + a_0$, $A^* = a_1^*x + a_0^*$, $s = x^2 + s_1x + s_0$, and $r = x^2 + r_1x + r_0$. We obtain

$$k_1x + k_0 = (a_1x + a_0)(x^2 + s_1x + s_0) + (a_1^*x + a_0^*)(x^2 + r_1x + r_0)$$

Equating coefficients for cubes and constants results in $-a_1^* = a_1$, and $-a_0^* = (a_0s_0 - k_0)/r_0$. So $-A^*$ can be calculated using one multiplication, one division, and one addition. All told, calculation of A and $-A^*$ requires the equivalent of 14 multiplications and 11 additions.

Case 2: s is quadratic and r is linear

We have $s = rq_1 + x_1$. Since x_1 is a unit, we stop.

$$\frac{K}{sr} = \frac{\frac{1}{x_1}K(s-q_1r)}{sr}$$
$$= \frac{\frac{K}{x_1}}{r} - \frac{\frac{1}{x_1}q_1K}{s}$$

 $\frac{K}{x_1} = A$, and both K and x_1 are units, so

Table (4.2) Summary of Case 2

to obtain x_1 requires		1 mult	2 adds,
and to obtain A requires	<u>1 div.</u>	<u></u>	
The result is	$1 \mathrm{div}$	$1 \mathrm{mult}$	2 adds.

Now we have, as before,

$$k_1x + k_0 = A(x^2 + s_1x + s_0) + (a_1^*x + a_0^*)(x + r_0).$$

Equating coefficients of squares and constants gives $-a_1^* = A$ and $-a_0^* = (As_0 - k_0)/r_0$. So $-A^*$ is computed with one multiplication, one division, and one addition. All together, calculation of A and $-A^*$ requires the equivalent of 4 multiplications and 3 additions.

Case 3: s is linear and r is quadratic

We have $s = rq_1 + x_1$, and $r = x_1q_2 + x_2$. But $q_1 = 0$, thus $x_1 = s$ and notice that x_2 is then a unit, so $-q_2s + r = x_2$ and then

$$\frac{K}{sr} = \frac{\frac{1}{x_2}K(-q_2s+r)}{sr} = \frac{-\frac{1}{x_2}q_2K}{r} + \frac{\frac{1}{x_2}K}{s}$$

This time we will compute $-A^*$, a unit, first. Observe that $K = Qs - (x_2(-A^*))$. Now we do as before and write $k_1x + k_0 = (a_1x + a_0)(x + s_0) + (a_1x + a_0)(x + s_0)$

 $A^*(x^2 + r_1x + r_0)$. Equate coefficients of squares and constants to get $a_1 = -A^*$ and $a_0 = (k_0 - A^*r_0)/s_0$. Now we summarize.

Table (4.3) Summary of Case 3

To obtain x_2 requires		1 mult	2 adds,
to obtain x_2A^* requires		1 mult	2 adds,
to obtain $-A^*$ requires	1 div		1 add,
and to obtain a_0 requires	$1 \operatorname{div}$	<u>1 mult</u>	<u>1 add.</u>
The result is	2 divs	$3 \mathrm{mults}$	5 adds.

All together calculation of A and $-A^*$ requires the equivalent of 5 multiplications and 5 additions.

Case 4: r and s are both linear

This case is, of course, the simplest. $s = q_1r + x_1$, x_1 is a unit and $q_1 = 1$. So

$$\frac{K}{sr} = \frac{\frac{1}{x_1}K(s-r)}{sr}$$
$$= \frac{\frac{K}{x_1}}{r} + \frac{-\frac{K}{x_1}}{s}.$$

Thus $A = -A^* = K/x_1$, which requires 1 multiplication and 1 addition. Now that each case has been examined, the following table summarizes the preceding information:

$\deg(s)$	$\deg(r)$	multiplications	$\operatorname{additions}$
1	1	1	1
2	1	4	3
1	2	5	5
2	2	14	11

Table (4.4) Number of operations required to evaluate (4.2)

The following is essentially Chin's and Steiglitz's algorithm in $\mathbf{R}[x]$ instead of $\mathbf{C}[x]$. Let $p \in \mathbf{N}$ be given and $d_j \in \mathbf{R}[x]$, for each $1 \leq j \leq p$ and each d_j irreducible over \mathbf{R} . Let $m_j \in \mathbf{N}$ denote the multiplicity of each d_j . Also let $Q_0, D \in \mathbf{R}[x]$ be given such that $D = \prod_{j=1}^{p} (d_j)^{m_j}$ and $\deg(Q_0) < \deg(D) = \sum_{j=1}^{p} \deg(d_j)m_j$. Thus we have a proper rational function and wish to find $K_{ij} \in \mathbf{R}[x]$ such that

$$\frac{Q_0}{D} = \frac{Q_0}{\prod_{j=1}^p (d_j)^{m_j}} = \sum_{j=1}^p \sum_{i=1}^{m_j} \frac{K_{ij}}{(d_j)^i}.$$
(4.4)

Define $m_0 = 0$ and $n = \sum_{j=0}^{p} m_j$. Now, for $1 \le l \le n$, define

$$u_{l} = 1 + \min\{x \in \mathbb{N} \cup \{0\} \mid \sum_{j=0}^{x} m_{j} < l\},\$$
$$v_{j}^{l} = \begin{cases} m_{j}, & \text{if } j < u_{l};\\ l - \sum_{i=0}^{u_{l}-1} m_{i}, & \text{if } j = u_{l}, \end{cases}$$

for $1 \leq j \leq u_l$. Also, for each l, define $f_l = d_{u_l}$, and $R_l, Q_l \in \mathbf{R}[x]$ such that $Q_{l-1} = Q_l f_l + R_l$. Again define $A_{lj}(K), A_{lj}^*(K) : \mathbf{R}[x] \mapsto \mathbf{R}[x]$ by

$$\frac{K}{f_l d_j} = \frac{A_{lj}^*(K)}{f_l} + \frac{A_{lj}(K)}{d_j}$$

Now the partial fraction expansion can be obtained as below in n steps, the lth step being:

$$\frac{Q_0}{\prod_{j=1}^p d_j^{m_j}} = \frac{1}{\prod_{j=1}^n f_j} [Q_0]
= \frac{1}{\prod_{n \ge k > l} f_k} [Q_l + \sum_{j=1}^{u_l} \sum_{i=1}^{v_j^l} \frac{K_{ij}^l}{(d_j)^i}],$$
(4.5)

where it is understood that $\prod_{n \ge k > n} f_k = 1$ and

$$K_{ij}^{l} = \begin{cases} 0, & \text{if } i > v_{j}^{l} \text{ or } j > u_{l} \text{ or } l = 0; \\ A_{lj}(K_{ij}^{l-1}), & \text{if } j < u_{l} \text{ and } i = v_{j}^{l}; \\ A_{lj}(K_{ij}^{l-1} + A_{lj}^{*}(K_{(i+1)j}^{l-1})), & \text{if } j < u_{l} \text{ and } i < v_{j}^{l}; \\ K_{(i-1)j}^{l-1}, & \text{if } j = u_{l} \text{ and } i > 1; \\ R_{l} + \sum_{0 \le k < 1} A_{lk}^{*}(K_{1k}^{l-1} + A_{lk}^{*}(K_{2k}^{l-1})), & \text{if } j = u_{l} \text{ and } i = 1. \end{cases}$$

$$(4.6)$$

Now we must count operations to compare this method with Chin's and Steiglitz's method. It turns out that the number of operations depends on the number of quadratic factors in the denominator of the rational function. Let $N = \deg(D)$ and then denote the number of quadratic factors in D as q. Thus N = n + q.

First consider the number of calculations necessary to compute $\{R_l \mid 1 \leq l \leq n\}$. The *l*th stage involves dividing Q_{l-1} by f_l . Two facts are necessary: To divide an *M*-degree polynomial by a monic linear factor requires *M* multiplications and *M* additions and to divide the same polynomial by a monic quadratic factor requires 2(M-1) multiplications and additions. Note that the largest that $\deg(Q_0)$ can be is N-1. We shall prove that in this worst case it requires no more than $\frac{1}{2}N(N-1)-q$ multiplications and additions to compute $\{R_l \mid 1 \leq l \leq n\}$.

We shall use induction on N. For N = 1 we must have n = 1, and q = 0, and of course, $\deg(Q_0) = 0$ thus $\frac{1}{2}N(N-1) - q = 0$, which reflects the fact that there is really nothing to do in this case. We will also need to examine the case where N = 2, with n = 1 and q = 1. We still get $\frac{1}{2}N(N-1) - q = 0$, which again indicates that there is really no partial fraction expansion to carry out. Now assume the result for a given N.

First assume that we add a linear factor to D and increase $\deg(Q_0)$ to (N + 1) - 1 = N. So divide Q_0 by this new linear factor to get $\deg(Q_1) = N - 1$, which will require N multiplications and additions. Now apply the induction hypothesis to Q_1 . It will require $\frac{1}{2}N(N-1) - q$ multiplications and additions to obtain $\{R_l \mid n\}$

 $2 \leq l \leq n+1$. Adding up, we get

$$N + \frac{1}{2}N(N-1) - q = \frac{1}{2}(N+1)N - q.$$

Now assume that we add a quadratic factor to D and increase deg (Q_0) to (N+2) - 1 = N + 1. Dividing Q_0 by the new quadratic factor requires 2Nmultiplications and additions. We are left with deg $(Q_1) = N - 1$. By induction, to compute $\{R_l \mid 2 \leq l \leq n\}$ requires $\frac{1}{2}N(N-1) - q$ multiplications and additions. Summing, we get

$$2N + \frac{1}{2}N(N-1) - q = \frac{1}{2}(N+2)(N+1) - (q+1),$$

as required.

We must also consider the necessary number of calculations required to compute $\{K_{ij}^l \mid 1 \leq j \leq u_l, \text{and} 1 \leq i \leq v_j^l\}$ for some $1 \leq l \leq n$. It turns out that the number of operations needed to compute this set depends on $\deg(d_{u_l})$, on m_{u_l} , and on the number of quadratic factors preceding d_{u_l}

Notice that computation of $K_{iu_l}^l$ requires no calculation for i > 1. This means that the largest operation count for the partial fraction expansion algorithm occurs when all the factors of D are distinct, that is, when $m_j = 1$ for all $1 \le j \le p = n$. Consequently, $u_l = l$, so $f_l = d_l$ and $v_j^l = m_j = 1$ for all admissable l and j. Now write (4.5) as

$$\frac{Q_0}{\prod_{k=1}^n f_k} = \frac{1}{\prod_{n\geq k>l} f_k} [Q_l + \sum_{j=1}^l \frac{K_{1j}^l}{f_j}].$$

And we can also write (4.6) as

$$K_{1}^{l}j = \begin{cases} 0, & \text{if } l = 0\\ A_{lj}(K_{1j}^{l-1}), & \text{if } 1 < j < l\\ R_{l} + \sum_{0 \le k < l} A_{lk}^{*}(K_{1k}^{l-1}), & \text{if } j = l. \end{cases}$$

Calculation of $\{K_{1j}^l \mid 1 \leq j \leq l\}$ for a given l requires that A_{lj} and $-A_{lj}^*$ be determined l-1 times along with $(l-1) \deg(f_l)$ additions. Consider the number of

calculations in computation of A_{lj} and A_{lj}^* . Table (4.4) gives us this information if we let $s = f_l$ and $r = f_j$. Notice that, given f_l , it will take the most operations if f_j is quadratic. Hence, the number of calculations in computing $\{K_{1j}^l \mid 1 \le j \le l\}$ will be largest if deg $(f_k) = 2$ for all $k \le l$.

In order to find an upper bound on the number of calculations in this algorithm, assume all factors of D are distinct and ordered such that all quadratic factors appear first. Then the entire algorithm would require

$$\frac{1}{2}N(N-1) - q + \sum_{k=1}^{q} 14(k-1) + \sum_{q+1}^{n} (5q + (k - (q+1)))$$

$$= N(N-1) - q^{2} + 3Nq - 7q$$
(4.7)

multiplications. This formula also works for q = 0 and q = N/2. The result in each case is N(N-1) and $\frac{9}{4}N(N-2)$, respectively which is easily verified. Also the algorithm will require

$$\frac{1}{2}N(N-1) - q + \sum_{k=1}^{q} (11+2)(k-1) + \sum_{q+1}^{n} ((5+2)q + (1+1)(k-(q+1)))$$

$$= \frac{3}{2}N(N-1) - \frac{7}{2}q^2 + 3Nq - \frac{11}{2}q$$
(4.8)

additions. Again, for the special cases q = 0 and q = N/2, the formula yields $\frac{3}{2}N(N-1)$ and $\frac{17}{4}N(N-2)$ respectively. In fact it is quite easy to show that for all $N \ge 2$ and $N/2 \le q \le 0$ we get the following:

$$6N(N-1) > \frac{9}{4}N(N-2) \ge N(N-1) - q^2 + 3Nq - 7q.$$

This shows that the greatest number of multiplications occurs when q = n, and is still less than the number required by Chin's and Steiglitz's algorithm. Also observe

$$\frac{11}{2}N(N-1) > \frac{17}{4}N(N-2) \ge \frac{3}{2}N(N-1) - \frac{7}{2}q^2 + 3Nq - \frac{11}{2}q.$$

Which tells the same story for additions.

These results show that this adaption of Chin's and Steiglitz's algorithm saves calculations. To be fair, however, one must realize that the output of this adaptation is not the same as that of Chin and Steiglitz. Which algorithm is better will depend on the application. Partial fractions expansions can be useful in a wide scope of problems involving integrals of rational functions [6].

Chin's and Steiglitz's output differs from the one presented here in that C[x] is the ambient polynomial ring and each denominator in the partial fractions expansion is thus linear. With the adapted algorithm, R[x] is used and the denominators can be linear or quadratic. It happens that in computing inverse Laplace transforms, either form is acceptable and it is better to have fewer operations; however, this may not always be so for other applications of partial fraction expansion.

Chapter V

Inverse Laplace Transform

The most elementary approach to finding the inverse Laplace transform of a given function is to use a table of transform pairs. Indeed, large tables of transform pairs have been prepared. In particular, a popular table [7] lists the following transform pair:

$$\mathcal{L}^{-1}\left\{\frac{1}{\left(s^{2}+a^{2}\right)^{k}}\right\} = \frac{\sqrt{\pi}}{\Gamma(k)} \left(\frac{t}{2a}\right)^{k-1/2} J_{k-1/2}(at), \qquad k \in \mathbb{N}.$$
(5.1)

The functions J_p are known as the Bessel functions of half-integral order. They have the following recursive definition [8]:

$$J_{-1/2}(t) = \sqrt{\frac{2}{\pi t}} \cos t,$$

$$J_{1/2}(t) = \sqrt{\frac{2}{\pi t}} \sin t,$$

$$J_{p+1}(t) = \frac{2p}{t} J_p(t) - J_{p-1}(t).$$

Let us simplify things somewhat by defining:

$$H_{k}(at) = \sqrt{\pi} \left(\frac{t}{2a}\right)^{k-1/2} J_{k-1/2}(at).$$

We thus obtain the recursive relationship:

$$H_0(at) = \frac{2}{t}\cos(at),$$

$$H_1(at) = \frac{1}{a}\sin(at),$$

$$H_{k+1}(at) = \frac{2k-1}{2a^2}H_k(at) - \left(\frac{t}{2a}\right)^2 H_{k-1}.$$

Hence the Laplace transform pair (5.1) becomes:

$$\mathcal{L}^{-1}\left\{\frac{1}{\left(s^{2}+a^{2}\right)^{k}}\right\}=\frac{1}{\Gamma(k)}H_{k}(at).$$

If we use a well-known property of the Laplace transform, we can derive another important Laplace transform pair. Use

$$\mathcal{L}^{-1}\left\{\frac{d}{ds}F(s)\right\} = -tf(t)$$

to get

$$\mathcal{L}^{-1}\left\{\frac{d}{ds}\left(\frac{1}{\left(s^{2}+a^{2}\right)^{k}}\right)\right\} = -\frac{t}{\Gamma(k)}H_{k}(at)$$
$$\mathcal{L}^{-1}\left\{\frac{-2ks}{\left(s^{2}+a^{2}\right)^{k+1}}\right\} = -\frac{t}{\Gamma(k)}H_{k}(at)$$
$$\mathcal{L}^{-1}\left\{\frac{s}{\left(s^{2}+a^{2}\right)^{k+1}}\right\} = \frac{t}{2\Gamma(k+1)}H_{k}(at)$$

Which is equivalent to

L

$$\mathcal{L}^{-1}\left\{\frac{s}{\left(s^{2}+a^{2}\right)^{k}}\right\} = \frac{t}{2\Gamma(k)}H_{k-1}(at).$$

This would be true for $k \ge 2$ but also holds for k = 1. Again make use of a Laplace transform property, namely $\mathcal{L}^{-1} \{F(s + \tau)\} = e^{-\tau t} f(t)$ to get one of the Laplace transform pairs useful in this problem:

$$\mathcal{L}^{-1}\left\{\frac{A(s+\tau)+B}{\left((s+\tau)^{2}+a^{2}\right)^{k}}\right\} = \frac{e^{-\tau t}}{\Gamma(k)}\left[BH_{k}(at)+\frac{At}{2}H_{k-1}(at)\right].$$
(5.2)

We also make use of another often-tabulated Laplace transform pair [7]:

$$\mathcal{L}^{-1}\left\{\frac{A}{\left(s+\tau\right)^{k}}\right\} = \frac{Ae^{-\tau t}t^{k-1}}{\Gamma(k)}.$$
(5.3)

One of these two Laplace transform pairs will apply to each term in the partial fraction expansion of a rational function. From (4.3) and we have

$$\mathcal{L}^{-1}\left\{\frac{Q_{0}}{\prod_{j=1}^{p}(d_{j})^{m_{j}}}\right\} = \mathcal{L}^{-1}\left\{\sum_{j=1}^{p}\sum_{i=1}^{m_{j}}\frac{K_{ij}}{(d_{j})^{i}}\right\}$$
$$= \sum_{j=1}^{p}\sum_{i=1}^{m_{j}}\mathcal{L}^{-1}\left\{\frac{K_{ij}}{(d_{j})^{i}}\right\}.$$
(5.4)

Now if $\deg(K_{ij}) = 1$ use (5.3), and if $\deg(K_{ij}) = 2$ use (5.2).

Using (5.2), and induction it can be shown that:

$$\mathcal{L}^{-1}\left\{\frac{A(s+\tau)+B}{\left((s+\tau)^{2}+a^{2}\right)^{k}}\right\} = e^{-\tau t} \sum_{j=0}^{k-1} t^{j} \left(\alpha_{j} \cos(at) + \beta_{j} \sin(at)\right).$$

And also note that (5.3) can be written

$$\mathcal{L}^{-1}\left\{\frac{A}{\left(s+\tau\right)^{k}}\right\} = e^{-\tau t}t^{k-1}\left(A\cos(0)\right)$$

Thus we can express (5.4) in the form:

$$\sum_{j=1}^{p} \sum_{i=1}^{m_j} \mathcal{L}^{-1} \left\{ \frac{K_{ij}}{(d_j)^i} \right\} = \sum_{j=1}^{p} e^{-\tau_j t} \sum_{i=1}^{m_j} t^i \left(\alpha_{ji} \cos(a_j t) + \beta_{ji} \sin(a_j t) \right).$$
(5.5)

This is the way the inverse Laplace transform is computed by the program in the appendix. The output is simply an array of coefficients for an expression of the form (5.5).

Chapter VI

Applications

One of the most probable applications of this algorithm will be to evaluate transient responses of control systems with a known transfer function. Figure (6.1) shows a control diagram for an automatic flight control system for a supersonic aircraft [9]. The transfer function corresponding to figure (6.1) will depend on the values of K_1 and K_2 . It is simple to compute this transfer function using the coefficient values shown in figure (6.1).

The above transfer function was inverse transformed using various values for K_1 and K_2 . The transient response functions so obtained are graphed in figures (6.2) and (6.3). Notice how the response improves until the onset of instability.

It was mentioned before that sometimes we want the inverse transform of something other than a rational function. One common example arises when a control system contains time delays. We have the following transform pair.

$$\mathcal{L}\left\{u(t-T)\right\} = \frac{e^{-sT}}{s}$$

In order to apply the algorithm, we must approximate e^{sT} by a rational function. This can be done using the Padé approximant [9]. We have

$$e^{-sT} \simeq P_n(sT) = rac{\sum_{j=0}^n (-1)^j b_j(sT)^j}{\sum_{j=0}^n b_j(sT)^j},$$

where

$$b_j = \frac{\binom{n}{j}}{\binom{2n}{j}j!}$$

Using the Padé approximant, the algorithm was used to invert $\frac{e^{-s}}{s}$. Padé approximants of order 2,3,4, and 6 were used. These approximants are shown graphically in figure (6.4).

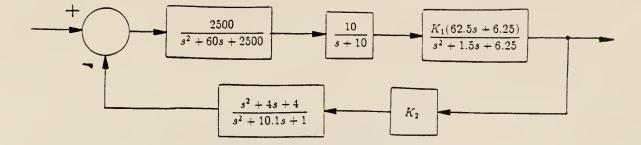


Figure (6.1) Control system diagram of SST aircraft. Adapted (9).

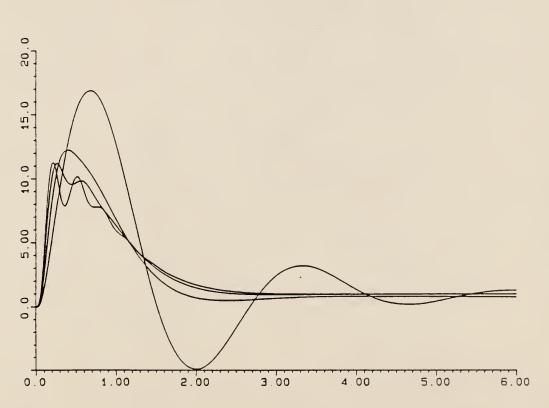


Figure (6.2) Step response of above control system for $K_1K_2 = 0, 0.2, 0.4$, and 0.6.

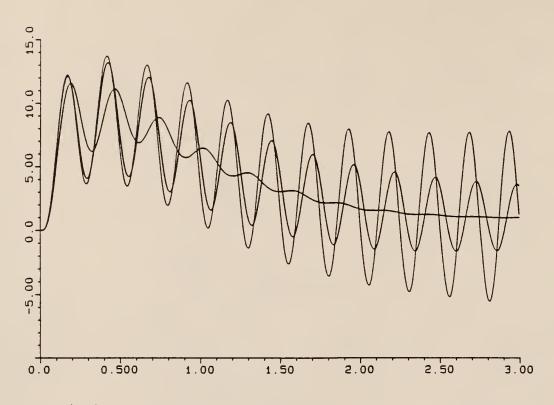


Figure (6.3) Step response of above control system for $K_1K_2 = 0.8, 1.08$, and 1.10.

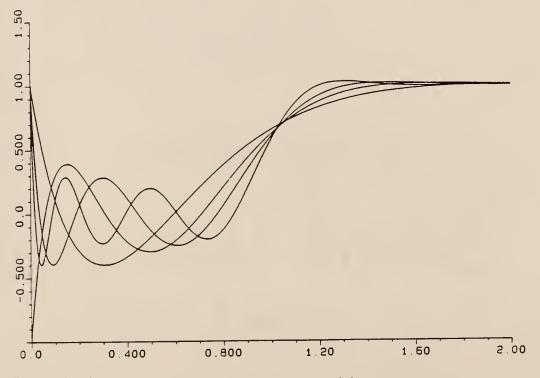


Figure (6.4) Padé approximants to unit time delay.

Chapter VII

Conclusion

This paper has presented an algorithm for computing inverse Laplace transforms of rational functions as might arise in practical electrical engineering problems. Programs written to demonstrate the algorithm follow in the appendix.

Numerical analysis aspects of the problem were not dealt with, but, except for root-finding, the problem was shown to be an algebraic one. Results from elementary abstract algebra were used to derive the methods described. Special effort was made to reduce the number of calculations in the partial fraction expansion.

Some applications were presented to show practical results. These applications also made it clear that assuming that the Laplace domain function, F(s), in (2.1) is a ratio of polynomials is not always valid. Future work on this problem should concern itself with this assumption.

Appendix I

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Appendix II

Glossary of Terms

C denotes the field of complex numbers.

C[x] denotes the ring of polynomials with coefficients in C.

deg: If $p(x) = a_0 + a_1x + \cdots + a_nx^n \neq 0$ and $a_n \neq 0$ then $\deg(p(x)) = n$.

Division algorithm: Given two polynomials $p(x), q(x) \in F[x]$, where F[x] is the ring of polynomials with coefficients in the field F and $q(x) \neq 0$, there exist two polynomials $t(x), r(x) \in F[x]$ such that f(x) = t(x)q(x) + r(x) where r(x) = 0 or deg(r(x)) < deg(q(x)). The process by which t(x) and r(x) are found is known as the division algorithm and is simply the "long-division" process everyone knows to divide one polynomial by another.

Euclidean algorithm: Given two polynomials $p(x), q(x) \in F[x]$, where F[x] is the ring of polynomials with coefficients in the field F and p(x) and q(x) are not both 0, then $gcd(p(x), q(x)) \in F[x]$ exists and there exist polynomials $m(x), n(x) \in F[x]$ such that gcd(p(x), q(x)) = m(x)p(x) + n(x)q(x). The process used to determine these special polynomials is called the Euclidean algorithm and is shown explicitly in Chapter IV.

 Γ : A sufficient definition of the function Γ for $n \in \mathbb{N} \cup \{0\}$ is

$$\Gamma(n) = \begin{cases} 0, & \text{if } n = 0; \\ (n-1)!, & \text{if } n > 0. \end{cases}$$

gcd: Let $a, b \in F[x]$. If $c \in F[x]$ satisfies:

1. c is monic.

2. c divides a and b.

3. Any other divisor of a and b divides c.

then c is called the greatest common divisor of a and b and is denoted gcd(a, b).

irreducible: Let $p \in F[x]$ be such that p = ab for some $a, b \in F[x]$ if and only if deg(a) = 0 or deg(b) = 0. Such a p is said to be irreducible over F. Note that irreducibility depends on the field, F. $x^2 + 1$ is irreducible over $\mathbf{R}[x]$ but not over $\mathbf{C}[x]$.

min is a function that operates on a well ordered set and whose value is the minimum element of that set. The well ordering property of N asserts that if $A \subset N$ then min A exists.

N denotes the ring of natural numbers.

 ${\bf R}$ denotes the field of real numbers.

 $\mathbf{R}[\mathbf{x}]$ denotes the ring of polynomials with coefficients in \mathbf{R} . u(t): Let $t \in \mathbf{R}$.

$$u(t) = \begin{cases} 0, & \text{if } t < 0; \\ 1, & \text{if } t \ge 0. \end{cases}$$

unit: $u \in F[x]$ is called a unit iff deg(u)=0.

Appendix III

FORTRAN Programs

The following programs are intended to merely demonstrate the algorithms described in the previous chapters. Specifically, they were used to evaluate the example applications mentioned in Chapter VI. No guarantee of their usefulness to any other application is implied.

These programs could certainly be made more user-friendly. There are no error handling routines, the user interaction is minimal, and file management is cumbersome. Such things are left to a better programmer. Nevertheless, the programs serve their purpose of demonstrating the algorithms.

The programs fall into four slightly overlapping categories: those associated with Chapters III through V and those for creating data files and making plots. The following is a brief categorical index of the programs. Subroutine dependence is indicated by indentions.

Factoring Program

ROOT_FIND POLY_READ FACTORER DERIV EUCLID POLDIV FIND_EM MULLER DEFVAL COMPOSE SPEC_WRITE

Partial Fraction Expansion Program

PART_FRAC SPEC_READ EXPAND TRANSFER POLDIV ALIKE TRANSFER DIFFERENT TRANSFER EUCLIDEAN TRANSFER POLDIV POLMULT POLADD PART_WRITE

Inverse Laplace Transform Program

INVERT INV2

INIT BESSEL GAMMA

Plotting Program

PLOT

READ LOTS_O_PLOTS PLOT_O_MATIC

Data Entry Routines

INPUT_RAT

SPEC_INPUT SPEC_WRITE

The programs that follow are listed in alphabetical order according to their VAX FORTRAN filenames. Each program is preceded by a header that explains the purpose of the program, and describes the variables passed to and from the program. I have tried to make each header as complete as required to be all the documentation necessary to comprehend the program it describes.

* * Department of Electrical and Computer Engineering * * Kansas State University * * * * VAX FORTRAN source filename: ALIKE.FOR * * ROUTINE: SUBROUTINE * ALIKE (I, J, X, DEGX, REM, DEGR) * * * DESCRIPTION: Refer to equation (4.6) in the main * thesis. This program computes K¹_ij * when $j=u_1$, hence the name ALIKE. * * * DOCUMENTATION * FILES: None. * * * ARGUMENTS: The following arguments are passed to * the subroutine: * * Ι (input) integer * corresponds to j in (4.6) * * J (input) integer * corresponds to i in (4.6) * * Х (input) real * is a three dimensional array, X(I,J,K) * represents to Ith coefficient of the * numerator of the (J,K)th term in the * partial fraction expansion. Namely, that term with the Jth factor of DEN * to the Kth power as denominator. * * DEGX (input) integer * is an array. DEGX(I,J) represents the * degree of the numerator of the (I, J)th * term in the partial fraction expansion. * See the description of X. * * REM (input) real * corresponds to R_1 in (4.6). * * DEGR (input) integer * is the degree of R_1 in (4.6). * * **RETURN:** Not used.

```
*
*
*
       ROUTINES
*
       CALLED:
                      PUTX
*
*
*
       AUTHOR:
                      James F. Stafford
*
*
*
       DATE CREATED:
                      8Jun87 Version 1.0
*
*
*
       REVISIONS:
                      None.
*
+
SUBROUTINE
                      ALIKE (I, J, X, DEGX, REM, DEGR)
       IMPLICIT
                      NONE
                      DEGR, I, J, K, L, DEGX(10, *)
       INTEGER
       REAL*8
                      X(0:1,10,*),REM(0:10)
       DO K=J,2,-1
          DEGX(I, K) = DEGX(I, K-1)
          DO L=0,DEGX(I,K)
              X(L, I, K) = X(L, I, K-1)
          ENDDO
       ENDDO
       CALL PUTX(I, 1, REM, DEGR, X, DEGX)
      REIURN
```

END

* Department of Electrical and Computer Engineering * * Kansas State University * * * * * VAX FORTRAN source filename: BESSEL.FOR * * ROUTINE: SUBROUTINE * BESSEL (F, A, N) * * This program computes the recursively DESCRIPTION: * defined function H_k described in * chapter V. * * **DOCUMENTATION** * FILES: None. * * * The following arguments are passed to ARGUMENTS: * the subroutine: * * F (input) real * is an array containing the coefficients * of the functions H_k. For a given j, * F(I,0,K) represents the coefficient of * the ∞ sine term of H_(j-I), with t to * the power K. F(I, l, K) represents the * coefficient of the sine term of $H_{(j-I)}$, * with t to the power K. * × Α (input) real * represents a in the recursive definition of H_k in chapter V. * * * (input) integer Ν * represents k in the recursive definition of H_k in chapter V. * * * **RETURN:** Not used. * * * ROUTINES * CALLED: None. * * * James F. Stafford AUTHOR: * * * DATE CREATED: 9Jun87 Version 1.0 *

```
*
*
       REVISIONS:
                       None.
*
*
*****
                      BESSEL(F, A, N)
       SUBROUTINE
       IMPLICIT
                       NONE
       INTEGER
                      N, J, K
                      A, F(-2:0, 0:1, -1:9)
       REAL*8
       DO J=-1, N-2
           DO K=-2,-1
               F(K,0,J) = F(K+1,0,J)
               F(K, 1, J) = F(K+1, 1, J)
           ENDDO
       ENDDO
       DO J=0,N-2
           F(0,0,J) = F(-1,0,J) * (2*N-3) /A
           F(0,1,J) = F(-1,1,J) * (2*N-3) /A
       ENDDO
       DO J=-1,N-3
           F(0,0,J+2) = F(0,0,J+2) - F(-2,0,J)
           F(0,1,J+2) = F(0,1,J+2) - F(-2,1,J)
       ENDDO
```

RETURN

* Department of Electrical and Computer Engineering * * * Kansas State University * * * VAX FORTRAN source filename:COMPOSE.FOR × * * ROUTINE: COMPOSE (X, FACTOR, NUM, DEGF, MULT) * * DESCRIPTION: This program accepts a complex-valued × root, X as input, decides whether X is * purely real or not, and updates the factor array, FACTOR, accordingly. * * * DOCUMENTATION * FILES: None. * × * ARGUMENTS: * * Х (input) complex * Is a complex-valued root of a polynomial. * * (input/output) real FACTOR * Is an array containing each factor of the * above polynomial. * * (input/output) integer NUM * Is the number of factors in FACTOR. NUM * is already incremented before calling * COMPOSE. * * (input/output) integer DEGF * Is an array specifying the degree of each * corresponding factor in FACTOR. * * MULT (input/output) integer * Is an array specifying the multiplicity of * each factor in FACTOR. * * Not used. **RETURN:** * * × ROUTINES * CALLED: None. * * * AUTHOR: James F. Stafford *

```
*
*
                     30Jun88 Version 1.0
       DATE CREATED:
*
*
*
       REVISIONS:
                      None.
*
÷
SUBROUTINE
                      COMPOSE(X, FACTOR, NUM, DEGF, MULT)
                     NONE
       IMPLICIT
       INTEGER
                     NUM, DEGF(*), MULT(*)
      REAL*8
                     FACTOR(10,0:2), SMALL
      COMPLEX*16
                     X
      LOGICAL
                    REAL
       PARAMETER
                 (SMALL=10E-4)
      REAL= FALSE.
      IF (DREAL(X).NE.0.) THEN
          IF (DABS(DIMAG(X)/DREAL(X)).LT.SMALL) THEN
              REAL=.TRUE.
              FACTOR(NUM, 1) =1
              FACTOR(NUM, 0) = -DREAL(X)
              DEGF(NUM)=1
          ENDIF
      ELSE IF (CDABS(X).LT.SMALL) THEN
          REAL=.TRUE.
          FACTOR(NUM, 1) = 1.
          FACTOR (NUM, 0) = 0.
          DEGF(NUM)=1
      ENDIF
      IF (REAL.EO..FALSE.) THEN
```

FACTOR(NUM, 2) =1
FACTOR(NUM, 1) =-2.*DREAL(X)
FACTOR(NUM, 0) =DIMAG(X)*DIMAG(X)+DREAL(X)*DREAL(X)
DEGF(NUM)=2

ENDIF MULT (NUM) =1 RETURN END

* * * * * * * * * *	Kansas State University *				
* * * *	ROUTINE:	COMPLEX*16 FUNCTION DEFVAL (POLY, DEG, FACTOR, NUM, DEGF, MULT, X)			
* * * *	DESCRIPTION:	This program evaluates a polynomial at a given complex argument, X, all known factors are divided out.			
* * *	DOCUMENTATION FILES:	None.			
* * *	ARG UMENTS:	The following arguments are passed to the function:			
* * * *	POLY	(input) real is an array containing coefficients of the polynomial to be evaluated.			
* * *	DEG	(input) integer is the degree of POLY.			
* * * *	FACIOR	(input) real is an array containg the coefficients of all known factors of POLY.			
× * *	NUM	(input) integer is the number of factors in FACTOR.			
* * *	DEGF	(input) integer is an array specifying the degree of each corresponding factor in FACTOR.			
* * *	MULT	(input) integer is an array specifying the multiplicity of each corresponding factor in FACTOR.			
* * * * *	х	(input) complex is the argument at which the polynomial is to be evaluated.			

```
Not used.
*
       RETURN:
*
*
*
       ROUTINES
*
       CALLED:
                     None.
*
*
*
       AUTHOR:
                      James F. Stafford
*
*
*
       DATE CREATED: 30Jun88 Version 1.0
*
*
*
       REVISIONS:
                     None.
*
*
COMPLEX*16 FUNCTION
                            DEFVAL (POLY, DEG, FACTOR, NUM, DEGF, MULT, X)
       IMPLICIT
                     NONE
       INTEGER
                    DEG, I, NUM, DEGF(*), MULT(*)
      REAL*8
                     POLY(0:*), FACTOR(10,0:2)
      COMPLEX*16
                    X, EVAL
      DEFVAL=POLY(DEG)
      DO I=DEG-1,0,-1
          DEFVAL=DEFVAL*X+POLY(I)
      ENDDO
      DO I=1,NUM
          IF (DEFVAL.NE.0) THEN
              DEFVAL=DEFVAL/EVAL(FACTOR, I, DEGF(I), MULT(I), X)
          ENDIF
      ENDDO
      REIURN
      END
```

COMPLEX*16 FUNCTION EVAL (FACTOR, I, DEGF, MULT, X)

INTEGER J, I, DEGF, MULT

REAL*8 FACTOR(10,0:2)

COMPLEX*16 X, VALUE

EVAL=1 VALUE=FACTOR(I,DEGF)

DO J=DEGF-1,0,-1

VALUE=VALUE*X+FACTOR(I,J)

ENDDO

DO J=1,MULT

EVAL=EVAL*VALUE

ENDDO

RETURN

*	Department of Electrical and Computer Engineering * Kansas State University *		
*	K	*	
* *******		<pre>urce filename:DERIV.FOR * ***********************************</pre>	
* * *	ROUTINE:	SUBROUTINE DERIV (POLY, DEG)	
* * *	DESCRIPTION:	This program computes the derivative of a given polynomial.	
* * *	DOCUMENTATION FILES:	None.	
* * *	ARGUMENTS:	The following arguments are passed to the routine.	
* * * * *	POLY	(input/output) real is an array containing the coefficients of the polynomial to be differentiated. On return, this array contains the coefficients of the derivative.	
* * * * *	DEG	(input/output) integer is the degree of the polynomial on input and the degree of the derivative on output.	
* * *	RETURN:	Not used.	
* * * *	ROUTINES CALLED:	None.	
* * *	AUTHOR:	James F. Stafford	
* * *	DATE CREATED:	30Jun88 Version 1.0	
* * *	REV IS IONS:	None.	

SUBROUTINE DE	ERIV(POLY,DEG)	
IMPLICIT NO	ONE	
INTEGER I,	, J, DEG	
REAL*8 PC	DLY(0:*)	
DO J=0,DEG-1		
POLY(J) = (J+1) * POLY(J+1)		

ENDDO

POLY(DEG)=0. DEG=DEG-1

RETURN

*	Department of Electrical and Computer Engineering * Kansas State University *		
*	*		
* ******	VAX FORTRAN SOU	<pre>urce filename: DIFFERENT.FOR * ***********************************</pre>	
* * *	ROUTINE:	SUBROUTINE DIFFERENT(I, B, DEGB, DEN, DEGD, MULTS, X, DEGX)	
* * * *	DESCRIPTION:	Refer to equation (4.6) in the main thesis. This program computes K^l_ij when j <u_l, different.<="" hence="" name="" th="" the=""></u_l,>	
* * *	DOCUMENTATION FILES:	None.	
* * *	ARG UMENTS:	The following arguments are passed to the subroutine:	
* * *	I	(input) integer corresponds to j in (4.6)	
^ * * * * *	В	(input) real is an array containing the coefficients of the polynomial f_1 in the notation of chapter IV.	
* * *	DEGB	(input) integer is the degree of B	
× * * * * * *	DEN	(input) real is a two-dimensional array. DEN(I,J) represents the coefficient of the Ith power of x in the Jth factor of the denominator polynomial.	
× * * * *	DEGD	(input) integer is an array. DEGD(I) represents the degree of the Ith factor in the denominator polynomial.	
* * * *	MULTS	(input) integer is an array. MULTS(I) represents the multiplicity of the Ith factor in the denominator polynomial.	

* * * * * * *		Х	(input) real is a three-dimensional array. X(I,J,K) represents the Jth coefficient of the numerator of the (J,K)th term in the partial fraction expansion. Namely, that term with the Jth factor of DEN to the Kth power as denominator.
* * * * *		DEGX	(input) integer is a two-dimensional array. DEG(I,J) represents the degree of the numerator of the (I,J)th term in the partial fraction expansion. See the description of X.
* * *		RETURN:	Not used.
* * *		ROUTINES CALLED:	EUCLID, PUTX, GETD, GETX, FOLADD
* * *		AUTHOR:	James F. Stafford
* *		DATE CREATED:	8Jun87 Version 1.0
* * *		REVISIONS:	None.
****	***		
		SUBROUTINE	DIFFERENT(I, B, DEGB, DEN, DEGD, MULTS, X, DEGX)
		IMPLICIT	NONE
	+	INTEGER	DEGD(*),MULTS(*),I,J,K,L,DEGS, DEGT,DEGX(10,*),DEGA,DEGB,DEGF
	+	REAL*8	<pre>DEN(0:2,*),X(0:1,10,*),A(0:2),B(0:2), S(0:1),T(0:1),F(0:1)</pre>
		DO J=I-1,1,-1	
		PRINT *, 'J=	',J
		CALL GETD(J	, A, DEGA, DEN, DEGD)
CALL GETX(J, MULTS(J), F, DEGF, X, DEGX)		, MULTS(J), F, DEGF, X, DEGX)	

DO K=MULTS(J)-1,1,-1

PRINT *, 'K=',K

CALL EUCLID (A, DEGA, B, DEGB, F, DEGF, S, DEGS, T, DEGT)

CALL PUTX (J, K+1, T, DEGT, X, DEGX)

CALL GETX (J, K, F, DEGF, X, DEGX)

CALL POLADD (F, DEGF, S, DEGS)

ENDDO

CALL EUCLID (A, DEGA, B, DEGB, F, DEGF, S, DEGS, T, DEGT)

CALL PUTX (J, 1, T, DEGT, X, DEGX)

CALL GETX(I, 1, F, DEGF, X, DEGX)

CALL POLADD (F, DEGF, S, DEGS)

CALL PUTX(I, 1, F, DEGF, X, DEGX)

ENDDO

RETURN

* * *	Department of Electrical and Computer Engineering * Kansas State University *				
* *******	<pre>* VAX FORTRAN source filename:EUCLID.FOR * ***********************************</pre>				
× * * *	ROUTINE:	SUBROUTINE EUCLID (POL1, DEG1, POL2, DEG2, GCD, DEGG)			
~ * * *	DESCRIPTION:	This program computes the greatest common divisor of two given polynomials.			
* * *	DOCUMENTATION FILES:	None.			
* *	ARGUMENTS:				
* * * *	POL1	(input) real is an array representing one input polynomial.			
* *	DEG1	(input) integer is the degree of POLL.			
* * *	POL2	(input) real is an array representing the other input polynomial.			
^ * * *	DEG2	(input) integer is the degree of POL2.			
* * *	GCD	(output) real is the gcd of the two input polynomials			
* * *	DEGG	(output) integer is the degree of GCD.			
* * *	RETURN:	Not used.			
* * *	ROUTINES CALLED:	POLDIV			
*	AUTHOR:	James F. Stafford			

```
*
*
*
                        30Jun88 Version 1.0
        DATE CREATED:
*
*
*
        REVISIONS:
                        None.
*
*
*****
        SUBROUTINE EUCLID (POL1, DEG1, POL2, DEG2, GCD, DEGG)
        IMPLICIT
                        NONE
        INTEGER
                        I, J, DEG1, DEG2, DEGG, DEGA, DEGB, DEGQ
       REAL*8
                        POL1(0:10), POL2(0:10), GCD(0:10),
                        A(0:10), B(0:10), Q(0:10), ZERO
    +
       LOGICAL
                        EASY, HARD
       PARAMETER
                        (Z ERO = 1.0E - 5)
       EASY=.FALSE.
       HARD=.TRUE.
       CALL GET(GCD, DEGG, POLl, DEG1)
       CALL GET(B, DEGB, POL2, DEG2)
       DO WHILE (.NOT. (DEGB. EQ. 0. AND. DABS(B(0)).LT. ZERO))
           CALL GET (A, DEGA, GCD, DEGG)
           CALL GET (GCD, DEGG, B, DEGB)
           CALL POLDIV (A, DEGA, GCD, DEGG, Q, DEGQ, B, DEGB, EASY, HARD)
           DO I=0, DEGG
               GCD(I) = GCD(I) / GCD(DEGG)
           ENDDO
       ENDDO
       RETURN
       END
       SUBROUTINE
                       GET (A, DEGA, B, DEGB)
```

IMPLICIT	NONE
INTEGER	I, J, DEGA, DEGB
REAL*8	A(0:*),B(0:*)
DEGA=DEGB	
DO I=0,DEGA	
A(I) = B(I)	
ENDDO	
RETURN	
END	

****** * * Department of Electrical and Computer Engineering * Kansas State University * * * * * EUCLIDEAN, FOR VAX FORTRAN source filename: * * ROUTINE: SUBROUTINE * EUCLID (A, DEGA, B, DEGB, F, DEGF, S, DEGS, * T, DEGT) + * DESCRIPTION: This program computes A and A^{*} in * equation (4.1) given K, s and r via * the methods discussed in chapter IV. * * * DOCUMENTATION × FILES: None. * * * The following arguments are passed to ARGUMENTS: * the subroutine: * * А (input) real is an array containing the coefficients * * of s in (4.2). * * DEGA (input) integer * is the degree of s in (4.2). * * (input) real В * is an array containing the coefficients * of r in (4.2). * * DEGB (input) integer * is the degree of r in (4.2). × * F (input) real is an array containing the coefficients * * of K in (4.2). * * (input) integer DEGF * is the degree of K in (4.2). * * S (output) real * is an array containing the coefficients * of A in (4.2). * * (output) integer DEGS * is the degree of A in (4.2).

*			
* * *		Т	(output) real is an array containing the ∞ efficients of A [*] in (4.2).
* *		DEGT	(output) integer is the degree of A^* in (4.2).
* *		REIURN:	Not used.
* * *		ROUTINES CALLED:	POLDIV, POLMULT, GET(contained in TRANSFER)
* * *		AUTHOR:	James F. Stafford
* *		DATE CREATED:	9Jun87 Version 1.0
* * *		REVISIONS:	None.
	ىلە بلە بلە بلە	و بله	****
		SUBROUTINE	EUCLID (A, DEGA, B, DEGB, F, DEGF, S, DEGS, T, DEGT)
		IMPLICIT	NONE
	+	INTEGER	I, DEGA, DEGB, DEGF, DEGT, DEGS, DEGQ, DEGR, DEG1, DEG2, ADD, MULT, DIV
	+	REAL*8	A(0:2),B(0:2),F(0:1),S(0:1),T(0:1), QUO(0:1),REM(0:2),BUFF1(0:3),BUFF2(0:3)
		LOGICAL	EASYDIV, HARD, EASYMULT, SWITCH
	1	ADD=0 DIV=0 MULT=0 EASYDIV=.TRUE. HARD=.FALSE. DEGR=1	
	:	IF (DEGA.LE.DEGE	3) THEN
		SWITCH=.TRUE	•

CALL GET (BUFF2, DEG2, A, DEGA) CALL GET (BUFF1, DEG1, B, DEGB)

ELSE

SWITCH=.FALSE.

CALL GET(BUFF1, DEG1, A, DEGA) CALL GET(BUFF2, DEG2, B, DEGB)

ENDIF

I=0

+

DO WHILE (DEGR.GT.0)

CALL POLDIV (BUFF1, DEG1, BUFF2, DEG2, QUO, DEGQ, REM, DEGR, EASYDIV, HARD)

CALL GET(BUFF1, DEG1, BUFF2, DEG2) CALL GET(BUFF2, DEG2, REM, DEGR)

EASYDIV=.FALSE. HARD=.TRUE. I=I+1

ENDDO

IF (I.LT.2) THEN

EASYMULT=.TRUE. DEGQ=0

ELSE

EASYMULT=.FALSE. BUFF2(0)=-BUFF2(0) ADD=ADD+1

ENDIF

CALL POLMULT (QUO, DEGQ, F, DEGF, BUFF1, DEG1, EASYMULT)

DO I=0,DEG1

BUFF1(I)=BUFF1(I)/BUFF2(0) DIV=DIV+1

ENDDO

END

RETURN

PRINT *, ADD, 'additions'
PRINT *, MULT, 'multiplies'
PRINT *, DIV, 'divisions'

ENDIF

- DEGT=1 S(0) =-S(0) T(1) =S(0) T(0) =(A(1) -B(0))*S(0)+F(1) ADD=ADD+3 MULT=MULT+1
- CALL POLDIV(BUFF1,DEG1,B,DEGB,QUO,DEGQ,S,DEGS, + EASYDIV,HARD)

ELSE

ENDDO

ENDIF

S(0) = S(0) + T(0)ADD=ADD+1

IF (DEGT.GT.0) THEN

S(0) =T(DEGT)*(B(DEGB-1)-A(DEGA-1)) MULT=MULT+1 ADD=ADD+1

DEGS=DEGB-1 S(DEGS)=T(DEGT) DO I=DEGS-1,0,-1

CALL POLDIV (BUFF1, DEG1, A, DEGA, QUO, DEGQ, T, DEGT, + EASYDIV, HARD)

IF (SWITCH) THEN

HARD=.FALSE.

* Department of Electrical and Computer Engineering * \star Kansas State University * * * * VAX FORTRAN source filename: EXPAND, FOR * * SUBROUTINE * ROUTINE: * EXPAND ((NUM, DEGN, DEN, DEGD, MULTS, X, DEGX, * NO_FACTS) * * * DESCRIPTION: This program performs a partial fraction * expansion on a rational function using * Chin's and Steiglitz's algorithm. * * * DOCUMENTATION * FILES: None. * * * The following arguments are passed to ARGUMENTS: * the subroutine: * * (input) real NUM * is an array containing the coefficients * of the numerator polynomial of the * rational function to be expanded. * * DEGN (input) integer * is the degree of the numerator * polynomial. * * DEN (input) real is a two-dimensional array. DEN(I, J) * * represents the coefficient of the Ith * power of x in the Jth factor of the * denominator polynomial. * * DEGD (input) integer * is an array. DEGD(I) represents the * degree of the Ith factor in the * denominator polynomial. * * MULTS (input) integer * is an array. MULTS(I) represents the * multiplicity of the Ith factor in the * denominator polynomial. * * NO_FACTS (input) integer

* *			is the number of factors in the denominator polynomial.
* * * * * * * *		Х	(output) real is a three-dimensional array. X(I,J,K) represents the Ith coefficient of the numerator of the (J,K)th term in the partial fraction expansion. Namely, that term with the Jth factor of DEN to the Kth power as denominator.
* * * * * *		DEGX	(output) integer is an array. DEGX(I,J) represents the degree of the numerator of the (I,J)th term in the partial fraction expansion. See the description of X.
* * *		RETURN:	Not used.
* * * * * *		ROUTINES CALLED:	GETD, FOLDIV, ALIKE, DIFFERENT
		AUTHOR:	James F. Stafford
* * *		DATE CREATED:	6Jun87 Version 1.0
~ * *		REV IS IONS:	None.
****	****	****	*****
	+	SUBROUTINE	EXPAND (NUM, DEGN, DEN, DEGD, MULTS, X, DEGX, NO_FACTS)
		IMPLICIT	NONE
	+ +	INTEGER	DEGN,NO_FACTS,DEGD(10),DEGB, MULTS(*),DEGR,I,DEGQ,J,K,L, DEGX(10,*)
	+	REAL*8	NUM(0:*),DEN(0:2,*),X(0:1,10,*), QUO(0:10),REM(0:10),B(0:2)
		LOGICAL	EASY, HARD
		EASY=.FALSE.	

HARD=.FALSE.

+

DO I=1, NO_FACTS

PRINT *, 'I=', I

CALL GEID (I, B, DEGB, DEN, DEGD)

DO J=1,MULTS(I)

CALL FOLDIV (NUM, DEGN, B, DEGB, QUO, DEGQ, REM, DEGR, EASY, HARD)

DEGN=DEGQ

DO K=0, DEGN

NUM(K) = QUO(K)

ENDDO

PRINT *, DEGR, 'YES'

DO K=0,DEGR

PRINT *, REM(K)

ENDDO

CALL ALIKE (I, J, X, DEGX, REM, DEGR)

CALL DIFFERENT(I, B, DEGB, DEN, DEGD, MULTS, X, DEGX)

ENDDO

ENDDO

RETURN

<pre>************************************</pre>			
* * ****	<pre>vAX FORTRAN source filename:FACTORER.FOR * **********************************</pre>		
* * * *	ROUTINE:	SUBROUTINE FACTORER (POLY, DEGP, FACTOR, NUM, DEGF, MULT)	
~ * * *	DESCRIPTION:	This program factors a given input polynomial into irreducible elements of R[x].	
* * *	DOCUMENTATION FILES:	None.	
*	ARG UMENTS:		
* * *	POLY	(input) real is an array containing the coefficients the polynomial to be factored.	
* * *	DEG P	(input) integer is the degree of POLY.	
* * * *	FACTOR	(output) real is an array containinng coefficients each factor of POLY.	
* *	NUM	(output) integer is the number of factors in FACIOR.	
* * * *	DEGF	(output) integer is an array specifying the degree of the corresponding factor in FACTOR.	
* * *	MULT	(output) integer is an array specifying the multiplicity of the corresponding factor in FACTOR.	
* * *	RETURN:	Not used.	
* * *	ROUTINES CALLED:	DERIV, EUCLID, FIND_EM	

```
*
*
                       James F. Stafford
       AUTHOR:
*
*
*
                       30Jun88 Version 1.0
       DATE CREATED:
*
*
*
       REVISIONS:
                       None.
*
SUBROUTINE
                       FACTORER (POLY, DEGP, FACTOR, NUM, DEGF, MULT)
       IMPLICIT
                       NONE
       INTEGER
                       I, J, K, DEGP, DEGF(*), MULT(*), NUM, DEGGCD(0:10),
                       DEGD, DEGG
    +
                       POLY(0:10), FACTOR(10,0:2), GCD(0:10,0:10), D(0:10)
       REAL*8
    +
                       G(0:10)
       DO I=0, DEGP
           D(I) = POLY(I)
           GCD(0, I) = POLY(I)
       ENDDO
       DEGD=DEGP
       DEGGCD(0) = DEGP
       K=0
       DO WHILE (DEGGCD(K).GT.0)
           K=K+1
           CALL DERIV (D, DEGD)
           CALL EUCLID (POLY, DEGP, D, DEGD, G, DEGG)
          DO J=0,DEGG
               GCD(K, J) = G(J)
           ENDDO
           DEGGCD(K)=DEGG
       ENDDO
       DO I=K-1,0,-1
```

```
DO J=0,DEGGCD(I)
```

```
G(J) = GCD(I, J)
```

ENDDO

```
DEGG=DEGGCD(I)
```

```
CALL FIND_EM(G, DEGG, FACTOR, NUM, DEGF, MULT)
```

ENDDO

RETURN

* *	Kansas State University *			
*	<pre>vAX FORTRAN source filename:FIND_EM.FOR * **********************************</pre>			
* * *	ROUTINE:	SUBROUTINE FIND_EM(P, DEGP, FACTOR, NUM, DEGF, MULT)		
* * *	DESCRIPTION:	A polynomial and a list of already known roots is passed to the subroutine.		
* * *	DOCUMENTATION FILES:	None.		
* * *	ARGUMENTS:			
* * *	RETURN:	Not used.		
* * *	ROUTINES CALLED:	MULLER		
* * *	AUTHOR:	James F. Stafford		
*	DATE CREATED:	5Sep86 Version 1.0		
* * *	REVISIONS:	None.		

	SUBROUTINE	FIND_EM(P, DEGP, FACTOR, NUM, DEGF, MULT)		
	IMPLICIT	NONE		
	INTEGER	J, DEGP, NUM, DEGF(*), MULT(*), SUM		
	REAL*8	P(0:*),FACTOR(10,0:2)		
	SUM=0			
	DO J=1,NUM			

MULT(J) = MULT(J)+1SUM= SUM= DEGF(J) * MULT(J)

ENDDO

DO WHILE (DEGP-SUM.GT.0)

CALL MULLER (P, DEGP, FACTOR, NUM, DEGF, MULT) SUM= SUM+ DEGF (NUM)

ENDDO

RETURN

* Department of Electrical and Computer Engineering * * * Kansas State University * * * VAX FORTRAN source filename: * GAMMA.FOR * * ROUTINE: FUNCT ION * GAMMA(K) * * This program computes the integer-valued * DESCRIPTION: * function Gamma defined in the glossary. * * DOCUMENTATION * FILES: None. * * The following argument is passed to the * ARGUMENTS: * function: * * Κ (input) integer * is any non-negative integer. * * **RETURN:** The function returns an integer according * to the definition in the glossary. * * * ROUTINES * CALLED: None. * * * James F. Stafford AUTHOR: * * * DATE CREATED: 5Sep86 Version 1.0 * * * **REV IS IONS:** None. * * INTEGER FUNCTION GAMMA(K) IMPLICIT NONE INTEGER J,K GAMMA=1

DO J=K-1,2,-1

GAMMA=GAMMA*J

ENDDO

REIURN

* × Department of Electrical and Computer Engineering * * Kansas State University * * * * VAX FORTRAN source filename: INIT.FOR * * ROUTINE: SUBROUTINE * INIT(F) * DESCRIPTION: This program initializes the recursively defined function H_k described in * * chapter V. * * × DOCUMENTATION * FILES: None. * * The following argument in passed to * ARGUMENTS: * the subroutine: * * F (output) real * is an array containing the coefficients * of the functions H_k. For a given j, * F(I,0,K) represents the ∞ efficient of the cosine term of $H_{(j-I)}$, with t to the * * power K. F(I,1,K) represents the * coefficient of the sine term of $H_{(j-I)}$ * with t to the power K. * **RETURN:** Not used. * * * ROUTINES * CALLED: None. * * * James F. Stafford AUTHOR: * * * 9Jun87 Version 1.0 DATE CREATED: * * * **REVISIONS:** None. * SUBROUTINE INIT(F)

IMPLICIT NONE I,J,K INTEGER F(-2:0,0:1,-1:9)REAL*8 DO I=-2,0 DO J=0,1 DO K=-1,9 F(I, J, K) = 0. ENDDO ENDDO ENDDO F(-1,0,-1) = 1.F(0,1,0) = 1. RETURN

.

* Department of Electrical and Computer Engineering * * * Kansas State University * * * * VAX FORTRAN source filename: INPUT_RAT.FOR * ROUTINE: PROGRAM * * DESCRIPTION: This program allows one to establish a * data file compatible with the inverse * transform package programs containing the necessary data to describe a rational * * function. * * * DOCUMENTATION * FILES: None. * * * ARGUMENTS: Not used. * * * **RETURN:** Not used. * * * ROUTINES * None. CALLED: * * * AUTHOR: James F. Stafford * * * 27May87 Version 1.0 DATE CREATED: * * * **REVISIONS:** None. * IMPLICIT NONE INTEGER I, N_DEG, D_DEG, NO_ROOTS, MLTPLCTS(10) REAL*8 NUM(0:15), DEN(0:15) COMPLEX*16 ROOTS(10)CHARACTER*15 FILENAME, YESNO

PRINT *, 'Are numerator roots known? (Y/N) '

READ (*,200) YESNO

IF (YESNO. EQ. 'Y') THEN

CALL INPUT_FACT (NO_ROOTS, ROOTS, MLTPLCTS)

CALL RECONSTRUCT (NO_ROOTS, ROOTS, MLTPLCTS, NUM, N_DEG)

ELSE

+

CALL INPUT_NONFACT (NUM, N_DEG)

ENDIF

PRINT *, 'Are denominator roots known? (Y/N) '

READ (*,200) YESNO

IF (YESNO. EQ. 'Y') THEN

CALL INPUT_FACT (NO_ROOTS, ROOTS, MLTPLCTS)

CALL RECONSTRUCT (NO_ROOTS, ROOTS, MLTPLCTS, DEN, D_DEG)

ELSE

+

CALL INPUT_NONFACT (DEN, D_DEG)

ENDIF

```
PRINT *, 'Enter filename.'
```

READ (*,200) FILENAME

200 FORMAT (A15)

OPEN (UNIT=1, FILE=FILENAME, STATUS='NEW')

WRITE (1,*) N_DEG

DO I=0, N_DEG

WRITE (1,*) NUM(I)

ENDDO

WRITE (1,*) D_DEG

DO I=0,D_DEG

WRITE (1,*) DEN(I)

ENDDO

CLOSE (UNIT=1, STATUS='KEEP')

END

SUBROUTINE INPUT_FACT (NO_ROOTS, ROOTS, MLTPLCTS)

IMPLICIT NONE

INTEGER NO_ROOTS, MLTPLCTS(*), I

COMPLEX*16 ROOTS(*)

PRINT *, 'Input number of roots'

READ (*,*) NO_ROOTS

100 FORMAT (F8.5,F8.5)

DO I=1, NO_ROOTS

PRINT *,'Input root number ',I READ (*,100) RCOTS(I) PRINT *,'Input corresponding multiplicity' READ (*,*) MLTPLCTS(I)

PRINT *, ROOTS(I), MLTPLCTS(I)

ENDDO

RETURN

END

+

SUBROUTINE	RECONSTRUCT (NO_ROOTS, ROOTS, MLTPLCTS, RESULT, ORDER)
IMPLICIT	NONE

- INTEGER NO_ROOTS, MLTPLCTS(*), I, J, ORDER, MULT
- COMPLEX*16 ROOTS(*), BUFFER(0:50)

REAL*8

RESULT(0:*)

BUFFER(0) = DCMPLX(1.0,0.0)

DO I=1,50

BUFFER(I) = DCMPLX(0.,0.)

ORDER=ORDER+MLTPLCTS(I)

DO J=ORDER, 1,-1

BUFFER(0) = -BUFFER(0) * ROOTS(1)

INPUT_NONFACT (POLY, DEG)

BUFFER(J) = BUFFER(J-1) - BUFFER(J) * ROOTS(I)

MULT=MLTPLCTS(I)

ENDDO

ENDDO

DO I=0, ORDER

ENDDO

ENDDO

RETURN

SUBROUTINE

IMPLICIT

INTEGER

REAL*8

END

DO WHILE (MULT. NE. 0)

MULT=MULT-1

RESULT(I) = REAL(BUFFER(I))

NONE

DEG,I

POLY(0:*)

ENDDO

ORDER=0

DO I=1, NO_ROOTS

PRINT *,'Input degree'
READ (*,*) DEG
DO I=0,DEG
PRINT *,'Input coeff. of power ',I
READ (*,*) POLY(I)
ENDDO

RETURN

******	*****	***************************************
* * *	Department of H	Electrical and Computer Engineering * Ansas State University *
		arce filename: INV2.FOR *
* * *	ROUTINE:	SUBROUTINE INV2 (I, J, RESP, OMEGA, A, B)
* * * * * * * * * * * * *	DESCRIPTION:	This program computes the inverse Laplace transform of a rational function such that equation (5.2) applies. Note that since the algorithm described in chapter V is recursive, that for a given k, all of the previous transforms for $k>j>=1$ must be already computed. The function H_k is computed each time the subroutine is called, using H_(k-1) and H_(k-2) which are held in an array intrinsic to this routine, namely, F. On K=1, F is initialized to hold the coefficients of H_0 and H1. The inverse transform coefficients are accumulated in an array called RESP.
* * *	DOCUMENTATION FILES:	None.
* * *	ARGUMENTS:	The following arguments are passed to the subroutine:
* * * * *	I	(input) integer is an index variable specifying which factor of the denominator polynomial of the original rational function is currently being inverse transformed.
* * *	J	(input) integer corresponds to k in (5.2)
* * * * * *	RESP	(input/output) real is an array to accumulate the computed time-domain response. RESP(j,0,i) represents alpha_ji in equation (5.5) and RESP(j,1,i) respresents beta_ji in equation (5.5).
*	OMEGA	(input) real

```
*
                        corresponds to a in (5.2).
*
*
                        (input) real
            А
*
                        corresponds to A in (5.2).
*
*
            В
                        (input) real
*
                        corresponds to B in (5.2).
*
*
        RETURN:
                        Not used.
*
*
*
        ROUTINES
*
        CALLED:
                        INIT, BESSEL, GAMMA
*
*
*
                        James F. Stafford
       AUTHOR:
*
*
*
                        5Sep86 Version 1.0
       DATE CREATED:
*
*
*
       REVISIONS:
                        None.
*
*
SUBROUTINE
                        INV2(I, J, RESP, OMEGA, A, B)
       IMPLICIT
                       NONE
       INTEGER
                        I, J, K, L, M, GAMMA
       REAL*8
                       A, B, OMEGA, F(-2:0,0:1,-1:9), RESP(10,0:1,0:9),
    +
                        ADJ
       IF (J.EO.1) THEN
           CALL INIT(F)
       ELSE
           CALL BESSEL (F, OMEGA, J)
       ENDIF
       ADJ=GAMMA(J)*(2*OMEGA)**(J-1)
       DO K=0,J-1
           RESP(I, 0, K) = RESP(I, 0, K) + (F(-1, 0, K-1) * A + F(0, 0, K) * B)
    +
           /ADJ
```

RESP(I,1,K) = RESP(I,1,K) + (F(-1,1,K-1)*A+F(0,1,K)*B) + /ADJ ENDDO

RETURN

***** Department of Electrical and Computer Engineering * * * * Kansas State University * * VAX FORTRAN source filename: × INVERT. FOR * * * ROUTINE: PROGRAM * INVERT * DESCRIPTION: This program computes the inverse * Laplace transform of a rational * function that has been expanded into * partial fractions. The user is prompted for a filename under which the * output of the partial fraction expander * * program has been stored. Again, the * user is prompted for another filename * under which to store the parameters of * the inverse transform function. * * DOCUMENTATION * FILES: None. * * ARGUMENTS: None. * * * **RETURN:** Not used. * * * ROUTINES * CALLED: INV2 * * * AUTHOR: James F. Stafford * * * 10Jun87 Version 1.0 DATE CREATED: * * * REVISIONS: None. * PROGRAM INVERT NONE IMPLICIT INTEGER I, J, K, NO_TERMS, ORDER, MULT(10), GAMMA

REAL*8 TAU(10),OMEGA(10),A,B,F(-2:0,0:1,-1:10), + RESP(10,0:1,0:9)

CHARACTER*15 FILENAME PRINT *,'Enter filename.' READ (*,200) FILENAME

200 FORMAT (A15)

OPEN (UNIT=1, FILE=FILENAME, STATUS='OLD')

READ (1,*) NO_TERMS

DO I=1, NO_TERMS

READ (1,*) ORDER PRINT *,ORDER READ (1,*) MULT(I) PRINT *,MULT(I)

IF (ORDER. EQ.1) THEN

READ (1,*) TAU(I) PRINT *,TAU(I)

ELSE

```
READ (1,*) TAU(I)
PRINT *,TAU(I)
READ (1,*) OMEGA(I)
PRINT *,OMEGA(I)
```

ENDIF

DO J=1,MULT(I)

IF (ORDER. EQ.1) THEN

READ (1,*) A PRINT *,A RESP(I,0,J-1) =A/GAMMA(J)

ELSE

READ (1,*) A PRINT *,'A =',A READ (1,*) B PRINT *, 'B =', B

CALL INV2(I, J, RESP, OMEGA(I), A, B)

ENDIF

ENDDO

ENDDO

CLOSE (UNIT=1, STATUS='KEEP')

PRINT *, 'Enter filename.'

READ (*,200) FILENAME

OPEN (UNIT=1, FILE=FILENAME, STATUS='NEW')

WRITE (1,*) NO_TERMS

DO I=1, NO_TERMS

```
WRITE(*,300) 'exp(',-TAU(I),'t)*'
WRITE(1,*) TAU(I),OMEGA(I)
WRITE(1,*) MULT(I)
```

DO J=1, MULT(I)

```
WRITE(*,301) '(',RESP(I,0,J-1),'t^',J-1,
'COS',OMEGA(I),'t + '
WRITE(1,*) RESP(I,0,J-1)
WRITE(*,301) ' ',RESP(I,1,J-1),'t^',J-1,
'SIN',OMEGA(I),'t) '
WRITE(1,*) RESP(I,1,J-1)
```

ENDDO

ENDDO

+

+

CLOSE (UNIT=1, STATUS='KEEP')

300 FORMAT (A5,El2.4E3,A3)
301 FORMAT (A2,El2.4E3,A2,I2,A3,El2.4E3,A4)

* Department of Electrical and Computer Engineering * * * Kansas State University * * * VAX FORTRAN source filename: LOTS O PLOTS.FOR * * * ROUTINE: subroutine * FIRST PLOT (DEVICE, NUM POINTS, X DATA, * Y_DATA, X_AXIS_TITLE, X_AXIS_UNITS, * Y_AXIS_TITLE, Y_AXIS_UNITS, PLOT_TITLE, * INFO) * * DESCRIPTION: Makes a plot using X_DATA as abscissa and Y_DATA as ordinate. The axes are * labelled with titles and units. The * plot is also titled. * * * DOCUMENTATION * FILES: None. * * * ARGUMENTS: * * (input) integer DEVICE * is the device type to display the plot * * 7475 for plotter * 4014 for terminal (Selanar only) * * (input) integer NUM POINTS * is the number of data points to * be plotted * * X DATA (input) real * is the array of abscissa values for the * data to be plotted * * Y DATA (input) real * is the array of ordinate values for the * data to be plotted * * X AXIS TITLE (input) character*(*) * is the title to be placed on the x_axis * * X AXIS UNITS (input) character*(*) * is the name to be given to the units

*			associated with the x-axis
^ * *		Y_AXIS_TITL	E (input) character*(*) is the title to be placed on the y-axis
* * *		Y_AXIS_UNIT	S (input) character*(*) is the name to be given to the units associated with the y-axis
* *		PLOT_TITLE	(input) character*(*) is the title to be placed on the plot
*		REIURN:	
* * * *		INFO	(output) real(6) is the information necessary to make subsequent plots on the same axes.
* *		ROUTINES	
* * *		CALLED:	P System of Generalized Plot Routines
* *		AUTHOR:	James F. Stafford
* * *		DATE CREATED:	24May86 Version 1.0
* * *		REVISIONS:	None.
*	****	*****	***************************************
	+ +	X	ST_PLOT (DEVICE, NUM_POINTS, X_DATA, Y_DATA, _AXIS_TITLE, X_AXIS_UNITS, Y_AXIS_TITLE, AXIS_UNITS, PLOT_TITLE, INFO)
		IMPLICIT NONE	
	+		NUM_POINTS, FORLAB, FORTIC, NEGFLG, FORM, LENSTR, UPDOWN
	+ +	FIRSTX	(*),Y_DATA(*),FACIOR,VEL,X,Y,LENGTH, ,DELTAX,ANGLE,CLEN,FIRDEL(4), ,DIVLNY,WIDTH,HEIGHT,INFO(6)
	+		<pre>X_AXIS_TITLE,Y_AXIS_TITLE,X_AXIS_UNITS, Y_AXIS_UNITS,PLOT_TITLE</pre>

CHARACTER*(1) BLANK, SIZE

*INTIALIZE PLOT DEVICE

FACTOR=1.0 BLANK='' SIZ E='A'

CALL PINIT (DEVICE, BLANK, FACTOR, SIZE)

*SET PEN VELOCITY

VEL=10.0

CALL PSTVEL(VEL)

*ESTABLISH ORIGIN

X=4.5 Y=4.5

CALL PORIG(X, Y)

*SET OFFSETS FOR AXIS ROUTINES (RELATIVE TO ORIGIN)

X=0.0 Y=0.0

*DRAW Y-AXIS AND LABEL

+

LENGTH=12.0

CALL PSCALE (Y_DATA, NUM_POINTS, LENGTH, FIRSTX, DELTAX, DIVLNY)

FIRDEL(3) =FIRSTX FIRDEL(4) =DELTAX FORLAB=110 FORTIC=1001 ANGLE=90.0

CALL PAXIS(X,Y,Y_AXIS_TITLE,Y_AXIS_UNITS,FORLAB, + FORTIC,LENGTH, ANGLE,FIRSTX,DELTAX,DIVLNY)

*DRAW X-AXIS AND LABEL

LENGTH=18

CALL PSCALE (X_DATA, NUM_POINTS, LENGTH, FIRSTX,

+ DELTAX, DIVLNX)

FIRDEL(1)=FIRSTX FIRDEL(2)=DELTAX FORLAB=211 FORTIC=2001 ANGLE=0.0

CALL PAXIS(X,Y,X_AXIS_TITLE,X_AXIS_UNITS,FORLAB, + FORTIC, LENGTH, ANGLE, FIRSTX, DELTAX, DIVLNX)

*DRAW CURVE

SONTL=0

CALL PLINE(X_DATA, Y_DATA, NUM_POINTS, FIRDEL, SONTL, + BLANK, DIVLNX, DIVLNY)

*TITLE THE PLOT

UPDOWN=0 X=9.0 Y=13.0 INFO(1) =FIRDEL(1) INFO(2) =FIRDEL(2) INFO(3) =FIRDEL(3) INFO(4) =FIRDEL(4) INFO(5) =DIVLNX INFO(6) =DIVLNY

CALL PPLOT (X, Y, UPDOWN)

CALL PTXTLN(PLOT_TITLE, LENSIR)

WIDTH=-LENSIR/2 HEIGHT=0.0

CALL PCHRPL(WIDTH, HEIGHT)

CALL PTEXT (PLOT_TITLE)

* CALL POLOSP

REIURN

*	Department of Electrical and Computer Engineering * Kansas State University *	
* *	VAX FORTRAN SOU	rce filename: MULLER.FOR *
******	*****	***************************************
* * *	ROUTINE:	SUBROUTINE MULLER (POLY, DEG, X)
* * * *	DESCRIPTION:	This program uses Muller's method (page 262, Numerical Recipes) to find a root of a polynomial.
* * *	DOCUMENTATION FILES:	None.
* * *	ARG UMENTS:	The following arguments are passed to the subroutine:
* * * *	POLY	(input) real is an array containing the coefficients of the polynomial of interest.
*	DEG	(input) integer is the degree of the above polynomial.
* * *	FACTOR	(input/output) is an array containing known roots of the polynomial represented by POLY.
* * *	NUM	(input/output) is the number of factors in FACTOR.
× * * *	DEGF	(input/output) is an array containing the degree of each corresponding factor in FACTOR.
* * *	MULT	(input/output) is an array containing the multiplicity of each corresponding factor in FACTOR
* * *	RETURN:	Not used.
* * *	ROUTINES CALLED:	DEFVAL, COMPOSE

```
*
*
*
                         James F. Stafford
        AUTHOR:
*
*
*
        DATE CREATED:
                         28 May 87 Version 1.0
*
*
*
                         30Jun88 Added deflation and factor table
        REVISIONS:
*
                         updating.
*
+
*****
                         MULLER (POLY, DEG, FACTOR, NUM, DEGF, MULT)
        SUBROUTINE
        IMPLICIT
                        NONE
        INTEGER
                         DEG, NUM, DEGF(*), MULT(*), I, NO_ITERATIONS, MAX
        REAL*8
                        POLY(0:*), ZERO, FACTOR(10,0:2)
        COMPLEX*16
                        X(-2:1), Q, A, B, C, D, P(-2:0), DEFVAL
        PARAMETER
                        (ZERO=1.0E-12)
       PARAMETER
                         (MAX=200)
        NO ITERATIONS=0
       X(-2) = DCMPLX(1.,1.)
       X(-1) = DCMPLX(1.,0.)
       X(0) = DCMPLX(1, -1)
       DO WHILE ((CDABS(X(0) - X(-1))) \cdot GT \cdot CDABS(X(0)) * Z ERO)
    +
                 . AND. (CDABS(X(0) - X(-2))). GT. CDABS(X(0)) * Z ERO)
    +
                 . AND. (NO_ ITERATIONS. LT. MAX))
            NO ITERATIONS=NO ITERATIONS+1
            B=DCMPLX(0.,0.)
           D=DCMPLX(0.,0.)
           DO WHILE ((D. EQ. DCMPLX(0.,0.)).AND.(B. EQ. DCMPLX(0.,0.)))
                DO I = -2, 0
                    P(I) = DEFVAL (POLY, DEG, FACTOR, NUM, DEGF, MULT, X(I))
                ENDDO
                O = (X(0) - X(-1)) / (X(-1) - X(-2))
```

```
A=Q*P(0) -Q*(1+Q)*P(-1)+Q*Q*P(-2) 
B=(2*Q+1)*P(0) -((1+Q)**2)*P(-1)+Q*Q*P(-2) 
C=(1+Q)*P(0) 
D=SQRT(B*B-4*A*C)
```

IF ((D.EQ.DCMPLX(0.,0.)).AND.(B.EQ.DCMPLX(0.,0.))) THEN

X(-1) = (X(0) + X(-1))/2.X(-2) = (X(0) + X(-2))/2.

ENDIF

ENDDO

IF (CDABS(B+D).GT.CDABS(B-D)) THEN

X(1) = X(0) - (X(0) - X(-1)) * 2*C/(B+D)

ELSE

X(1) = X(0) - (X(0) - X(-1)) + 2 C/(B-D)

ENDIF

DO I=-2,0

X(I) = X(I+1)

ENDDO

ENDDO

```
PRINT *, 'I MADE IT HERE', NO_ITERATIONS
```

NUM=NUM+1

CALL COMPOSE(X(1), FACTOR, NUM, DEGF, MULT)

RETURN

* Department of Electrical and Computer Engineering * * * Kansas State University * * * VAX FORTRAN source filename: PART FRAC. FOR * * * ROUTINE: PROGRAM * TEST * * DESCRIPTION: This program computes the partial * fraction expansion of a rational * function using the method described in * the thesis. The user is prompted * for a filename under which the factored * form of the rational function has been stored. The user is prompted again * for a filename under which to store the * * partial fraction expansion. * * * DOCUMENTATION * FILES: None. * * * ARGUMENTS: None. * * * Not used. **RETURN:** * * * ROUTINES * CALLED: SPEC_READ, PART_WRITE, EXPAND * * * AUTHOR: James F. Stafford * * * DATE CREATED: 10Jun87 Version 1.0 * * * **REVISIONS:** None. * ÷ PROGRAM TEST IMPLICIT NONE

INTEGER I, J, K, DEGN, DEGD(10), NO_FACTS, MULTS(10), + DEGX(10,5) REAL*8 NUM(0:15), DEN(0:2,10), X(0:1,10,5), + FACT(0:2)

LOGICAL EASY, HARD

CALL SPEC_READ (NUM, DEGN, DEN, DEGD, MULTS, NO_FACTS)

CALL EXPAND (NUM, DEGN, DEN, DEGD, MULTS, X, DEGX, NO_FACTS)

DO I=1,NO_FACTS

DO J=1,MULTS(I)

PRINT *, I, J

```
DO K=0, DEGX(I, J)
```

PRINT *, X(K, I, J)

ENDDO

ENDDO

ENDDO

CALL PART_WRITE (NO_FACTS, MULTS, X, DEGX, DEN, DEGD)

****** * * Department of Electrical and Computer Engineering * * Kansas State University * * VAX FORTRAN source filename: PART_WRITE.FOR * * * ROUTINE: SUBROUTINE * PART_WRITE (NO_FACTS, MULTS, X, DEGX, DEN, * DEGD) * * DESCRIPTION: This program writes the partial fraction * expansion into a file. The user is * prompted for a filename. * * DOCUMENTATION * FILES: None. * * * ARGUMENTS: The following arguments are passed to * the subroutine: * NO_FACTS (input) integer * is the number of factors in the * denominator polynomial. * * MULTS (input) integer * is an array containing the multiplicities of each factor in * * the denominator polynomial. * * Х (input) real * is a three-dimensional array. X(I, J, K)* represents the Ith coefficient of the * numerator of the (J,K)th term in the * partial fraction expansion. Namely, * that term with the Jth factor of DEN * to the Kth power as denominator. * * DEGX (input) integer * is an array. DEGX(I,J) represents the * degree of the numerator of the (I, J)th * term in the partial fraction expansion. * See the description of X. * * (input) real DEN * is a two-dimensional array. DEN(I,J) * represents the coefficient of the Ith * power of x in the Jth factor of the * denominator polynomial.

*		
* * * *	DEGD	(input) integer is an array. DEGD(I) represents the degree of the Ith factor of the denominator polynomial.
*	RETURN:	Not used.
* * *	ROUTINES CALLED:	None.
* * *	AUTHOR:	James F. Stafford
* * *	DATE CREATED:	7Jun87 Version 1.0
* * *	REVISIONS:	None.
******	*****	*******
	SUBROUTINE	PART_WRITE(NO_FACIS, MULTS, X, DEGX, DEN, DEGD)
	IMPLICIT	NONE
+	INTEGER	I, J, K, DEGD(10), NO_FACTS, MULTS(10), DEGX(10,5)
	REAL*8	DEN(0:2,10),X(0:1,10,5),ALPHA,BETA,A,B
	CHARACTER*15	FILENAME
	PRINT *, 'Enter :	filename.'
	READ (*,200) FILENAME	
200	FORMAT (A15)	
	OPEN (UNIT=1, F	ILE=FILENAME, STATUS='NEW')
	WRITE (1,*) NO_FACTS	
	DO I=1, NO_FACTS	
	WRITE (1,*)	DER =', DEGD(I)

IF (DEGD(I).EQ.1) THEN

WRITE (1,*) DEN(0,I)
PRINT *,'ALPHA =',DEN(0,I)

ELSE

```
ALPHA=DEN(1,I)/2
BETA=DSQRT(DEN(0,I)-ALPHA**2)
WRITE (1,*) ALPHA
PRINT *,'ALPHA =',ALPHA
WRITE (1,*) BETA
PRINT *,'BETA =',BETA
```

ENDIF

```
DO J=1, MULTS(I)
```

IF (DEGD(I).EQ.1) THEN

WRITE (1,*) X(0,I,J) PRINT *,'A =',X(0,I,J)

ELSE

```
A=X(1,I,J)
B=(X(0,I,J)-A*ALPHA)/BETA
WRITE (1,*) A
PRINT *,'A =',A
WRITE (1,*) B
PRINT *,'B =',B
```

ENDIF

ENDDO

ENDDO

```
CLOSE (UNIT=1, STATUS='KEEP')
```

RETURN

* Department of Electrical and Computer Engineering * * * Kansas State University * * * VAX FORTRAN source filename: * PLOT.FOR * * ROUTINE: PROGRAM * PLOT * * DESCRIPTION: This program makes plots of time * domain response functions computed * by the inverse transform program. * * **DOCUMENTATION** * FILES: None. * * * ARGUMENTS: None * * * **RETURN:** Not used. * * * ROUTINES * CALLED: FIRST_PLOT, READ, PLOT, PCLOSP (contained in the P System * * of Generalized Plotting Routines) * * James F. Stafford AUTHOR: * * * 24May87 Version 1.0 DATE CREATED: * * * **REVISIONS:** None. * * PROGRAM TEST NONE IMPLICIT INTEGER DEVICE, NUM_POINTS, I, NUM_FILES REAL $X_DATA(1000), Y_DATA(1000),$ + INFO(6), TONE, TTWO X_TITLE, Y_TITLE, X_UNITS, Y_UNITS, CHARACTER*(15) TITLE, FILES(5) +

PRINT *, 'INPUT <7475> FOR PLOTTER OR <4014> FOR TERMINAL' READ (*, *) DEVICE NUM_POINTS=1000 X TITLE='TIME' Y TITLE='VALUE' X_UNITS='INTERVALS' Y_UNITS='UNITS' TITLE='TEST PLOT'

PRINT *, 'Input number of files to plot' READ (*, *) NUM_FILES

DO I=1, NUM_FILES

PRINT *, 'Input name of file number ', I READ (*,200) FILES(I)

ENDDO

200 FORMAT (A15)

> PRINT *, 'Input initial time' READ (*,*) TONE

PRINT *, 'Input final time' READ (*, *) TIWO

DO I=1,NUM_POINTS

 $X_DATA(I) = TONE+ (TTWO-TONE) *$ + (FLOATJ(I-1)/FLOATJ(NUM POINTS))

ENDDO

CALL READ (NUM_POINTS, X_DATA, Y_DATA, FILES(1))

CALL FIRST PLOT (DEVICE, NUM POINTS, X_DATA, Y_DATA, + X_TITLE, X_UNITS, Y_TITLE, Y_UNITS, TITLE, INFO)

DO I=2, NUM FILES

CALL READ (NUM_POINTS, X_DATA, Y_DATA, FILES(I)) CALL PLOT (DEVICE, NUM_POINTS, X_DATA, Y_DATA, INFO)

ENDDO

CALL POLOSP

Department of Electrical and Computer Engineering * * * Kansas State University * * * * * VAX FORTRAN source filename: PLOT_O_MATIC.FOR * * ROUTINE: subroutine * PLOT (DEVICE, NUM_POINTS, X_DATA, * Y_DATA, INFO) * * * Makes a plot using X_DATA as abscissa DESCRIPTION: * and Y_DATA as ordinate. The axes are * assumed to be already drawn in * accordance with INFO. * * * DOCUMENTATION FILES: None. * * * ARGUMENTS: * * DEVICE (input) integer * is the device type to display the plot * * 7475 for plotter * 4014 for terminal (Selanar only) * * NUM_POINTS (input) integer * is the number of data points to * be plotted * * X_DATA (input) real * is the array of abscissa values for the * data to be plotted * * Y_DATA (input) real * is the array of ordinate values for the * data to be plotted * * * **RETURN:** * * (output) real(6) INFO * is the information necessary to make * subsequent plots on the same axes. *

* * ROUTINES * CALLED: P System of Generalized Plot Routines * * * James F. Stafford AUTHOR: * * * DATE CREATED: 24May86 Version 1.0 * * * **REVISIONS:** None. * *

SUBROUTINE PLOT (DEVICE, NUM_POINTS, X_DATA, Y_DATA, + INFO)

IMPLICIT NONE

INTEGER DEVICE, NUM_FOINTS, FORLAB, FORTIC, NEGFLG, FORM, + SCNTL, LENSTR, UPDOWN

REAL X_DATA(*),Y_DATA(*),FACTOR,VEL,X,Y,LENGTH, + FIRSTX,DELTAX,ANGLE,CLEN,FIRDEL(4), + DIVLNX,DIVLNY,WIDTH,HEIGHT,INFO(6)

CHARACTER*(1) BLANK, SIZE

*INTIALIZE PLOT DEVICE

FACTOR=1.0 BLANK='' SIZ E='A'

* CALL PINIT (DEVICE, BLANK, FACTOR, SIZE)

*SET PEN VELOCITY

VEL=10.0

CALL PSTVEL(VEL)

*ESTABLISH ORIGIN

X=4.5 Y=4.5 CALL PORIG(X,Y)

*SET OFFSETS FOR AXIS ROUTINES (RELATIVE TO ORIGIN)

X=0.0 Y=0.0

*ESTABLISH INFORMATION FOR PLOTTING SUBROUTINE

FIRDEL(1)=INFO(1) FIRDEL(2)=INFO(2) FIRDEL(3)=INFO(3) FIRDEL(4)=INFO(4) DIVLNX=INFO(5) DIVLNY=INFO(6)

*DRAW CURVE

+

SCNTL=0

CALL PLINE (X_DATA, Y_DATA, NUM_POINTS, FIRDEL, SONTL, BLANK, DIVLNX, DIVLNY)

RETURN

*****	****	****
* * *		Electrical and Computer Engineering * ansas State University *
	VAX FORTRAN SOU	rce filename: POLADD.FOR ************************************
* * * *	ROUTINE:	SUBROUTINE POLADD (A, DEGA, B, DEGB)
^ * * * *	DESCRIPTION:	This program subtracts on polynomial from another, replacing the first addend with the difference.
* * *	DOCUMENTATION FILES:	None.
* * *	ARGUMENTS:	The following arguments are passed to the subroutine.
* * * * *	A	(input/output) real is an array containing the coefficients for a polynomial on input and the coefficients for the difference on return.
* * * * *	DEGA	(input/output) integer is the degree of the above polynomial on input and the degree of the difference on return.
* * *	В	(input) real is an array containing the coefficients of the polynomial to be subtracted from A.
* * *	DEGB	(input) integer is the degree of the above polynomial.
* * *	RETURN:	Not used.
* * *	ROUTINES CALLED:	None.

```
*
       AUTHOR:
                      James F. Stafford
*
*
*
       DATE CREATED:
                      6Jun87 Version 1.0
×
*
*
       REV IS IONS:
                      None.
*
*
SUBROUTINE
                      POLADD (A, DEGA, B, DEGB)
       IMPLICIT
                      NONE
       INTEGER
                      I, DEGA, DEGB, ADDS
                      A(0:*),B(0:*)
       REAL*8
       ADDS=0
       IF (DEGA.GE.DEGB) THEN
          DO I=0,DEGB
              A(I) = A(I) - B(I)
              ADDS=ADDS+1
          ENDDO
       ELSE
          DO I=0, DEGA
              A(I) = A(I) - B(I)
              ADDS=ADDS+1
          ENDDO
          DO I=DEGA+1, DEGB
              A(I) = -B(I)
              ADDS=ADDS+1
          ENDDO
       ENDI F
      DEGA=JMAX0(DEGA, DEGB)
      PRINT *, ADDS, 'additions'
```

RETURN

* * Department of Electrical and Computer Engineering * * Kansas State University * * * VAX FORTRAN source filename: POLDIV.FOR * * * ROUTINE: SUBROUTINE * POLDIV (NUM, N_DEG, DEN, D_DEG, QUO, DEGQ, * REM, DEGR, EASY, HARD) * * * DESCRIPTION: This program performs the division * algorithm on two input polynomials * to obtain a quotient and remainder. * * * DOCUMENTATION * FILES: None. * * * ARGUMENTS: The following arguments are passed to * the subroutine. * * (input) real NUM * is an array containing the coefficients * of the numerator polynomial. * * N DEG (input) integer * is the degree of the numerator polynomial. * * DEN (input) real * is an array containing the coefficients * of the denominator polynomial. * * D_DEG (input) integer * is the degree of the denominator × polynomial. * * QUO (output) real * is an array containing the coefficients * of the quotient polynomial. * * DEGO (output) integer * is the degree of the quotient polynomial. * * REM (output) real * is an array containing the coefficients * of the remainder polynomial. *

* * *	DEGR	(output) integer is the degree of the remainder polynomial.
* * * *	EASY	(input) logical should be set to .TRUE. if both the numerator and denominator are monic polynomials. This will save calculations.
* * * * * *	HARD	(input) logical must be set to .TRUE. if the denominator is not a monic polynomial. Otherwise leave it false to save calculations.
* * *	RETURN:	Not used.
* * * *	ROUTINES CALLED:	None.
* *	AUTHOR:	James F. Stafford
* * *	DATE CREATED:	6Jun87 Version 1.0
* * *	REV ISIONS:	27Jul87 Added calculation-saving.
******	************	****
+	SUBROUTINE	POLDIV (NUM, N_DEG, DEN, D_DEG, QUO, DEGQ, REM, DEGR, EASY, HARD)
	IMPLICIT	NONE
	INTEGER	J, K, N_DEG, D_DEG, DEGQ, DEGR, MULT, ADD, DIV
	REAL*8	NUM(0:*),DEN(0:*),QUO(0:*),REM(0:*),ZERO
	LOGICAL	EASY, HARD
	PARAMETER	(ZERO=1.0E-5)
	DEGQ=N_DEG-D_DEX MULT=0 DIV=0 ADD=0	3

DO J=0, N_DEG

 $\operatorname{REM}(J) = \operatorname{NUM}(J)$

ENDDO

IF (DEGQ.GE.0) THEN

DO K=DEGQ,0,-1

IF (HARD) THEN

QUO(K) = REM(D_DEG+K)/DEN(D_DEG) DIV=DIV+1

ELSE

 $QUO(K) = REM(D_DEG+K)$

ENDIF

DO $J=D_DEG+K-1, K, -1$

IF (EASY) THEN

REM(J) = REM(J) - DEN(J-K)ADD=ADD+1

ELSE

```
REM(J) =REM(J) -QUO(K)*DEN(J-K)
ADD=ADD+1
MULT=MULT+1
```

ENDIF

ENDDO

EASY=.FALSE.

ENDDO

IF $(D_DEG.EQ.0)$ REM(0) = 0.

 $DEGR=JMAX0(0, D_DEG-1)$

DO WHILE ((DABS(REM(DEGR)).LT.ZERO).AND.DEGR.GT.0)

DEGR=DEGR-1

ENDDO

ELSE

$$DEGR=N_DEG$$

 $DEGQ=0$
 $OUO(0)=0$.

ENDIF

- PRINT *, MULT, 'multiplies' PRINT *, ADD, 'additions' PRINT *, DIV, 'divisions' *
- *
- *

RETURN

* Department of Electrical and Computer Engineering * * * Kansas State University ÷ * * * VAX FORTRAN source filename: POLMULT.FOR * * ROUTINE: SUBROUTINE * POLMULT (POL1, DEG1, POL2, DEG2, PROD, DEGP, * EASY) ÷ * * DESCRIPTION: This program multiplies two polynomials ÷ and returns their product. * * * DOCUMENTATION * FILES: None. * * * ARGUMENTS: The following arguments are passed to * the subroutine. * * POL1, POL2 (input) real * are arrays containing the coefficients * of the polynomials to be multiplied. * * DEG1,DEG2 (input) integer * are the degrees of the above * polynomials. * * PROD (output) real * is an array containing the coefficients * of the product polynomial. * * DEGP (output) integer * is the degree of the product * polynomial. * * EASY (input) logical should be set to .TRUE. only if POLL * is monic in order to save calculations. * * * * Not used. **RETURN:** * * * ROUTINES * CALLED: None. *

* * *	AUTHOR:	James F. Stafford
* * *	DATE CREATED:	7Jun87 Version 1.0
* * *	REVISIONS:	20Jul87 Added calculation-saving.
*	**************************************	<pre>************************************</pre>
	IMPLICIT	NONE
	INTEGER	J, K, DEG1, DEG2, DEGP, ADD, MULT
	REAL*8	POL1(0:*),POL2(0:*),PROD(0:*)
	LOGICAL	EASY
	ADD=0 MULT=0 DEGP=DEG1+DEG2	
	DO J=DEG2,0,-1	
	IF (EASY) T	HEN
	PROD (DEG1+J) =POL2 (J) ELSE	
	PROD(DE MULT=MU	Gl+J)=POL1(DEG1)*POL2(J) LT+1
	ENDIF	
	ENDDO	
	DO J=DEG1-1,0,-	1
	DO K=DEG2,1	,-1

```
PROD(J+K)=PROD(J+K)+POL1(J)*POL2(K)
ADD=ADD+1
MULT=MULT+1
```

ENDDO

PROD(J)=POL1(J)*POL2(0) MULT=MULT+1

ENDDO

PRINT *, ADD, 'additions' PRINT *, MULT, 'multiplies'

RETURN

* Department of Electrical and Computer Engineering * * * Kansas State University * * * VAX FORTRAN source filename: POLY_READ.FOR * * * ROUTINE: SUBROUTINE * POLY_READ (POLY, DEG) * * DESCRIPTION: * This program reads a polynomial from * a data file. The user must enter the data file name from the keyboard. * * * * DOCUMENTATION * FILES: None. * * * The following arguments are passed to ARGUMENTS: * the subroutine. * * (output) real POLY * is an array to receive the coefficients * of the polynomial. * * DEG (output) integer * is the degree of the polynomial. * * * **RETURN:** Not used. * * * ROUTINES * CALLED: None. * * * James F. Stafford AUTHOR: * * * 6Jun87 Version 1.0 DATE CREATED: * * * **REVISIONS:** None. * * SUBROUTINE POLY_READ (POLY, DEG)

IMPLICIT NONE INTEGER DEG,I REAL*8 POLY(0:10) READ (1,*) DEG READ (1,*) (POLY(I), I=0,DEG) RETURN

•

*	-	Electrical and Computer Engineering *
*	Ka	ansas State University *
	VAX FORTRAN sou	
* * * *	ROUTINE:	SUBROUTINE READ(NUM_POINTS,X_DATA,Y_DATA,FILENAME)
* * * * * *	DESCRIPTION:	This program evaluates a response function computed by the inverse transform program. The particular function parameters are stored in a data file under FILENAME.
* * *	DOCUMENTATION FILES:	None.
× * *	ARGUMENTS:	The following arguments are passed to the subroutine:
* * *	NUM_POINTS	(input) integer is the number of points in X_DATA at which response values are desired.
* * * *	X_DATA	(input) real is an array containing the values of time that the response function is to be evaluated at.
* * *	Y_DATA	(output) real is an array to accumulate the computed function values.
* * * *	FILENAME	(input) character is the filename of the response function to be evaluated.
* * *	RETURN:	Not used.
* * * *	ROUTINES CALLED:	None.
*	AUTHOR:	James F. Stafford

```
*
*
*
                       24May86 Version 1.0
       DATE CREATED:
*
*
*
       REVISIONS:
                       None.
*
READ (NUM_POINTS, X_DATA, Y_DATA, FILENAME)
       SUBROUTINE
       IMPLICIT
                       NONE
       INTEGER
                       NUM_POINTS, I, J, K, NO_TERMS, MULTS
                       X_DATA(1000), Y_DATA(1000), EXPONENTIAL,
       REAL
    +
                       TONE, TIWO, POWEROFT
       REAL*8
                       TAU, OMEGA, RESP(0:1)
       CHARACTER*(15) FILENAME
       DO I=1, NUM_POINTS
           Y_DATA(I)=0.
       ENDDO
       OPEN (UNIT=1, FILE=FILENAME, STATUS='OLD')
       READ (1,*) NO_TERMS
       DO I=1, NO_TERMS
           READ(1,*) TAU, OMEGA
           READ(1,*) MULTS
           POWEROFT=1
          DO J=1,MULTS
              READ(1, *) RESP(0)
              READ(1, \star) RESP(1)
              DO K=1, NUM_POINTS
                  IF (J-1.GT.0) POWEROFT=X_DATA(K)**(J-1)
                  Y_DATA(K) = Y_DATA(K) +
                  SNGL(DEXP(-TAU*X_DATA(K))*
    +
```

+	$(RESP(0) * DCOS(OMEGA*X_DATA(K)) +$
+	$RESP(1) *DSIN(OMEGA*X_DATA(K))))*$
+	POWEROFT

ENDDO

ENDDO

ENDDO

CLOSE (UNIT=1)

RETURN

******** × Department of Electrical and Computer Engineering * * * Kansas State University * * * VAX FORTRAN source filename: ÷ ROOT_FIND.FOR * * ROUTINE: PROGRAM * TEST * * DESCRIPTION: This program factors the denominator of * a rational function with real coefficients * into irreducible polynomials in R[x]. * The user is prompted for a filename * under which a rational function has * been stored. The user is again prompted * for a filename to store the result * under. * * DOCUMENTATION * FILES: None. * * * ARGUMENTS: None. * * * **RETURN:** Not used. * * * ROUTINES * CALLED: POLY_READ, FACTORER, SPEC_WRITE * * * AUTHOR: James F. Stafford * * * DATE CREATED: 8Jun87 Version 1.0 * * * **REVISIONS:** None. * * PROGRAM TEST IMPLICIT NONE INTEGER DEGN, DEGD, DEGF(10), NUM_FACTS, MULT(10), I, J REAL*8 NUM(0:10), DENOM(0:10), FACTOR(10,0:2)

CHARACTER*15 FILENAME

PRINT *, 'Input data file name'

READ (*,200) FILENAME

200 FORMAT (A15)

OPEN (UNIT=1, FILE=FILENAME, STATUS='OLD')

CALL POLY_READ (NUM, DEGN)

CALL POLY_READ (DENOM, DEGD)

CLOSE (UNIT=1)

CALL FACTORER (DENOM, DEGD, FACTOR, NUM_FACTS, DEGF, MULT)

DO I=1, NUM_FACTS

PRINT *, 'FACTOR NUMBER ', I

DO J=0, DEGF(I)

PRINT *, FACTOR(I, J)

ENDDO

PRINT *, 'MULTIPLICITY', MULT(I)

ENDDO

CALL SPEC_WRITE (NUM, DEGN, NUM_FACTS, FACTOR, DEGF, MULT) END

* *		Electrical and Computer Engineering * * ansas State University *
* *		urce filename: SIMPLE PLOT.FOR *

* R * *	OUTINE:	subroutine SIMPLE_PLOT(DEVICE, NUM_POINTS, X_DATA, Y_DATA, X_AXIS_TITLE, X_AXIS_UNITS, Y_AXIS_TITLE, Y_AXIS_UNITS, PLOT_TITLE,
* * *		PLOT_TYPE)
* * * * *	DESCRIPTION:	Makes a plot using X_DATA as abscissa and Y_DATA as ordinate. The axes are labelled with titles and units. The plot is also titled. One of four plot types can be selected.
* * *	DOCUMENTATION FILES:	None.
*	ARGUMENTS:	
~ * *	DEVICE	(input) integer is the device type to display the plot
~ * *		<pre>7475 for plotter 4014 for terminal (Selanar only)</pre>
* * *	NUM_ PO INTS	(input) integer is the number of data points to be plotted
* * *	X_DATA	(input) real is the array of abscissa values for the data to be plotted
* * *	Y_DATA	(input) real is the array of ordinate values for the data to be plotted
* * *	X_AXIS_TITL	E (input) character*(*) is the title to be placed on the x_axis
*	X_AXIS_UNIT	S (input) character*(*)

*	is the name to be given to the units associated with the x-axis
*	associated with the x-axis
* Y_AXIS_TITLE * *	(input) character*(*) is the title to be placed on the y-axis
	(input) character*(*) is the name to be given to the units associated with the y-axis
	(input) character*(*) is the title to be placed on the plot
* PLOT_TYPE * * *	(input) character*(*) is a character string which specifies the type of plot to be generated. The following are valid:
* * * *	'LINEAR' for linear-linear 'LOG_LINEAR' for log-linear 'LINEAR_LOG' for linear-log 'LOG_LOG' for log-log
* RETURN: 1	Not used.
* ROUTINES * CALLED: 1	P System of Generalized Plot Routines
* AUTHOR: 0	James F. Stafford
*	5Sep86 Version 1.0
*	None.
* *************************************	*************
+ X_A	LE_PLOT (DEV ICE, NUM_POINTS, X_DATA, Y_DATA, XIS_TITLE, X_AXIS_UNITS, Y_AXIS_TITLE, XIS_UNITS, PLOT_TITLE, PLOT_TYPE)
IMPLICIT NONE	
INTEGER DEVICE, N	JUM_ PO INTS, FORLAB, FORTIC, NEGFLG, FORM,

+ SONTL, LENSIR, UPDOWN

REAL X_DATA(*),Y_DATA(*),FACTOR,VEL,X,Y,LENGTH, + FIRSTX,DELTAX,DIVLEN,ANGLE,CLEN,FIRDEL(4), + DIVLNX,DIVLNY,WIDTH,HEIGHT

CHARACTER*(*) X_AXIS_TITLE, Y_AXIS_TITLE, X_AXIS_UNITS, + Y_AXIS_UNITS, PLOT_TITLE, PLOT_TYPE

CHARACTER*(1) BLANK, SIZE

*INTIALIZE PLOT DEVICE

FACTOR=1.0 BLANK='' SIZ E='A'

CALL PINIT (DEVICE, BLANK, FACTOR, SIZE)

*SET PEN VELOCITY

VEL=10.0

CALL PSIVEL (VEL)

*ESTABLISH ORIGIN

X=4.5 Y=4.5

CALL PORIG(X, Y)

*SET OFFSETS FOR AXIS ROUTINES (RELATIVE TO ORIGIN)

X=0.0 Y=0.0

*DRAW Y-AXIS AND LABEL

IF (PLOT_TYPE.EQ.'LINEAR'.OR.PLOT_TYPE.EQ.'LINEAR_LOG') + THEN

LENGTH=12.0

CALL PSCALE (Y_DATA, NUM_POINTS, LENGTH, FIRSTX, + DELTAX, DIVLEN)

FIRDEL(3)=FIRSTX FIRDEL(4)=DELTAX DIVLNY=DIVLEN FORLAB=110 FORTIC=1001 ANGLE=90.0

CALL PAXIS(X,Y,Y_AXIS_TITLE,Y_AXIS_UNITS,FORLAB, + FORTIC, LENGTH, ANGLE, FIRSTX, DELTAX, DIVLEN)

ELSE

LENGTH=12.0

CALL PLOGSC (Y_DATA, NUM_POINTS, LENGTH, FIRSTX, CLEN, + NEGFLG)

FIRDEL(3)=FIRSTX FIRDEL(4)=CLEN FORM=-1010 ANGLE=90.0

CALL PLGAXS(X,Y,Y_AXIS_TITLE,Y_AXIS_UNITS,FORM, + LENGTH, ANGLE, FIRSTX, CLEN)

ENDIF

*DRAW X-AXIS AND LABEL

IF (PLOT_TYPE.EQ.'LINEAR'.OR.PLOT_TYPE.EQ.'LOG_LINEAR') + THEN

LENGTH=18

CALL PSCALE(X_DATA, NUM_POINTS, LENGTH, FIRSTX, + DELTAX, DIVLEN)

> FIRDEL(1) =FIRSTX FIRDEL(2) =DELTAX DIVLNX=DIVLEN FORLAB=211 FORTIC=2001 ANGLE=0.0

CALL PAXIS(X,Y,X_AXIS_TITLE,X_AXIS_UNITS,FORLAB, + FORTIC, LENGTH, ANGLE, FIRSTX, DELTAX, DIVLEN)

ELSE

LENGTH=18.0

CALL PLOGSC (X_DATA, NUM_POINTS, LENGTH, FIRSTX, CLEN, + NEGFLG)

FIRDEL(1)=FIRSTX FIRDEL(2)=CLEN FORM=+2011 ANGLE=0.0

CALL PLGAXS(X,Y,X_AXIS_TITLE,X_AXIS_UNITS,FORM, + LENGTH, ANGLE, FIRSTX, CLEN)

ENDI F

*DRAW CURVE

IF (PLOT_TYPE.EQ.'LINEAR') THEN

SONTL=0

CALL PLINE(X_DATA, Y_DATA, NUM_POINTS, FIRDEL, SONTL, + BLANK, DIVLNX, DIVLNY)

ELSE IF (PLOT_TYPE. EQ. 'LOG_LINEAR') THEN

SONTL=0

CALL PLGLIN (X_DATA, Y_DATA, NUM_FOINTS, FIRDEL, SONIL, + BLANK, DIVLEN)

ELSE IF (PLOT_TYPE. EQ. 'LINEAR_LOG') THEN

SONTL=0

CALL PLNLOG (X_DATA, Y_DATA, NUM_POINTS, FIRDEL, SONTL, BLANK, DIVLEN)

ELSE IF (PLOT_TYPE. EQ. 'LOG_LOG') THEN

SONTL=0

CALL PLGLOG (X_DATA, Y_DATA, NUM_POINTS, FIRDEL, SONTL, + BLANK)

ENDIF

*TITLE THE PLOT

+

UPDOWN=0 X=9.0 Y=13.0

CALL PPLOT (X, Y, UPDOWN)

CALL PTXTLN(PLOT_TITLE, LENSIR)

WIDTH=-LENSTR/2 HEIGHT=0.0

CALL PCHRPL (WIDTH, HEIGHT)

CALL PTEXT (PLOT_TITLE)

CALL PCLOSP

RETURN END

* Department of Electrical and Computer Engineering * * * Kansas State University * * VAX FORTRAN source filename: * SPEC_INPUT.FOR * * PROGRAM ROUTINE: * SPEC_INPUT * * This program allows the user to enter DESCRIPTION: * the specifications for a rational * function already in factored form. * * DOCUMENTATION * FILES: None. * * * ARGUMENTS: None. * * * Not used. **RETURN:** * * * ROUTINES * CALLED: None. * * * AUTHOR: James F. Stafford * * * DATE CREATED: 15Jun87 Version 1.0 * * * **REVISIONS:** None. * * IMPLICIT NONE INTEGER $I, J, N_DEG, D_DEG(10), NO_ROOTS, MULTS(10),$ + DEG, NO_FACTS REAL NUM(0:15), DEN(0:2,10), X(0:1,10,5), + FACT(0:2), POLY(0:20)PRINT *, 'Enter numerator information.' CALL INPUT_NONFACT (NUM, N_DEG)

PRINT *, 'Enter denominator information.'

PRINT *, 'Input number of relatively prime irreducible factors.'
READ (*,*) NO_FACTS

DO I=1, NO_FACTS

PRINT *,'Enter information on factor number',I

CALL INPUT_NONFACT (FACT, D_DEG(I))

DO $J=0, D_DEG(I)$

DEN(J, I) = FACT(J)

ENDDO

PRINT *, 'Enter number of times this factor appears.'

READ (*,*) MULTS(I)

ENDDO

```
CALL SPEC_WRITE (NUM, N_DEG, NO_FACTS, DEN, D_DEG, MULTS)
```

END

SUBROUTINE INPUT_NONFACT (POLY, DEG)

IMPLICIT NONE

INTEGER DEG, I

REAL POLY(0:*)

PRINT *, 'Input degree'

READ (*,*) DEG

DO I=0,DEG

```
PRINT *,'Input coeff. of power ',I
READ (*,*) POLY(I)
```

ENDDO

RETURN

*	Department of H	Electrical and Computer Engineering *
*	Ka	ansas State University *
		<pre>prce filename: SPEC_READ.FOR * ***********************************</pre>
* * * *	ROUTINE:	SUBROUTINE SPEC_READ(NUM, DEGN, DEN, DEGD, MULTS, NO_FACTS)
* * * * * * * *	DESCRIPTION:	This program reads data for a rational function out of a file created by the program for factoring the denominator. There is also a program called SPEC_WRITE that will create a data file in the proper format.
* * * *	DOCUMENTATION FILES:	None.
* * *	ARGUMENTS:	The following arguments are passed to the subroutine:
* * * *	NUM	(output) real is an array containing the coefficients of the numerator polynomial.
~ * * *	DEGN	(output) integer is the degree of the numerator polynomial.
* * * *	DEN	(output) real is a two-dimensional array. DEN(I,J) represents the coefficient of the Ith power of x in the Jth factor of the denominator polynomial.
* * * * * *	DEGD	(output) integer is an array. DEGD(I) represents the degree of the Ith factor in the denominator polynomial.
* * *	MULTS	(output) integer is an array. MULTS(I) represents the multiplicity of the Ith factor in the

```
*
                        denominator polynomial.
*
*
            NO_FACTS
                        (output) integer
*
                        is the number of factors in the
*
                        denominator polynomial.
*
*
*
        RETURN:
                       Not used.
*
*
*
        ROUTINES
*
        CALLED:
                       None.
*
*
*
                       James F. Stafford
        AUTHOR:
*
*
*
        DATE CREATED:
                       6Jun87 Version 1.0
*
*
*
        REVISIONS:
                       None.
*
*
SPEC_READ (NUM, DEGN, DEN, DEGD, MULTS, NO_FACTS)
        SUBROUTINE
        IMPLICIT
                       NONE
        INTEGER
                       I, J, DEGN, DEGD(10), NO_FACTS, MULTS(10)
                       NUM(0:15), DEN(0:2,10), X(0:1,10,5),
       REAL*8
    +
                       FACT(0:2)
       CHARACTER*15
                       FILENAME
       PRINT *, 'Input data file name'
       READ (*,200) FILENAME
200
       FORMAT (A15)
       OPEN (UNIT=1, FILE=FILENAME, STATUS='OLD')
       READ (1,*) DEGN
       DO I=0, DEGN
           READ (1, *) NUM(I)
       ENDDO
```

```
READ (1,*) NO_FACTS
 DO I=1,NO_FACTS
    READ (1, *) DEGD(I)
    DO J=0, DEGD(I)
        READ (1, *) DEN(J, I)
    ENDDO
    READ (1,*) MULTS(I)
ENDDO
CLOSE (UNIT=1)
PRINT *, DEGN
DO I=0, DEGN
    PRINT *, NUM(I)
ENDDO
PRINT *, NO_FACTS
DO I=1, NO_FACTS
   PRINT *, DEGD(I)
   DO J=0, DEGD(I)
        PRINT *, DEN(J, I)
    ENDDO
    PRINT *, MULTS(I)
ENDDO
```

RETURN

END

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* Department of Electrical and Computer Engineering * * Kansas State University * * * * VAX FORTRAN source filename: SPEC WRITE.FOR * * * ROUTINE: SUBROUTINE * SPEC_WRITE (NUM, N_DEG, NO_FACTS, FACTOR, * D_DEG, MULTS) * * DESCRIPTION: This program writes out the factored * form of a rational function to a file specified by the user. * * * * DOCUMENTATION * FILES: None. * * * ARGUMENTS: The following arguments are passed to * the subroutine: * * (input) real NUM * is an array containing the coefficients * of the numerator polynomial. * * N_DEG (input) integer * is the degree of NUM * * NO FACTS (input) integer * is the number of factors in the * denominator polynomial. * (input) real * FACTOR * is an array containing the coefficients of each factor in * * the denominator polynomial. * * D_DEG (input) integer * is an array specifying the * degree of each factor in * the denominator polynomial. * * MULTS (input) integer * is an array specifying the * multiplicity of each factor in * the denominator polynomial. * \star **RETURN:** Not used.

```
*
*
*
       ROUTINES
*
       CALLED:
                     None.
*
*
*
                     James F. Stafford
       AUTHOR:
*
*
*
       DATE CREATED: 30Jun88 Version 1.0
*
*
*
       REVISIONS: None.
*
SUBROUTINE
                    SPEC_WRITE (NUM, N_DEG, NO_FACTS, FACTOR, D_DEG, MULTS
       IMPLICIT
                    NONE
       INTEGER
                     I, J, N_DEG, D_DEG(*), NO_FACTS, MULTS(*)
       REAL*8
                     NUM(0:*), FACTOR(10, 0:2)
       CHARACTER*15 FILENAME
       PRINT *, 'Enter filename.'
       READ (*,200) FILENAME
200 FORMAT (A15)
      OPEN (UNIT=1, FILE=FILENAME, STATUS='NEW')
      WRITE (1,*) N_DEG
      DO I=0, N_DEG
          WRITE (1,*) NUM(I)
      ENDDO
      WRITE (1,*) NO_FACTS
      DO I=1,NO_FACTS
          WRITE (1,*) D_DEG(I)
          DO J=0, D_DEG(I)
```

WRITE (1,*) FACTOR(I,J)

٠

ENDDO

WRITE (1,*) MULTS(I)

ENDDO

CLOSE (UNIT=1, STATUS='KEEP')

RETURN

* Department of Electrical and Computer Engineering * * * Kansas State University * * * VAX FORTRAN source filename: TRANSFER. FOR * * * ROUTINE: There are actually three programs * in this file: * SUBROUTINE * GEID (J, A, DEGA, DEN, DEGD) * * SUBROUTINE * GETX(J, K, F, DEGF, X, DEGX)* * SUBROUTINE * PUTX(J, K, F, DEGF, X, DEGX)* * * DESCRIPTION: These programs are a substitute for * a more sophisticated data structuring method. They copy the coefficients * for a polynomial embedded in a * higher dimensional array into a one-* dimensional array, or vice versa. * * * DOCUMENTATION * FILES: None. * * * ARG UMENTS: The following arguments are passed to * GETD: * * J (input) integer * is a number representing which factor * of the denominator is sought. * * DEN (input) real * is a two-dimensional array. DEN(I, J) * represents the coefficient of the Ith * power of x in the Jth factor of the * denominator polynomial. * * DEGD (input) integer * is an array. DEGD(I) represents the * degree of the Ith factor in the * denominator polynomial. * * Α (output) real

- 25ª

* * *		is an array to receive the coefficients of the factor polynomial.
* * *	DEGA	(output) integer is the degree of the factor polynomial.
* * *		The following arguments are passed to GETX:
* * *	J,K	(input) integer are the coordinates of the term in the partial fraction expansion that is sought. See the description of X.
< * * * * * * * *	Х	(input) real is a three-dimensional array. X(I,J,K) represents the Ith coefficient of the numerator of the (J,K)th term in the partial fraction expansion. Namely, that term with the Jth factor of DEN to the Kth power as denominator.
* * * *	DEGX	(input) integer is an array. DEGX(I,J) represents the degree of the numerator of the (I,J)th term in the partial fraction expansion. See the description of X.
* * *	F	(output) real is an array to receive the coefficients of the desired polynomial.
* * *	DEGF	(output) integer is the degree of the polynomial, F.
* * *		The following arguments are passed to PUTX:
* * * *	J,K	(input) integer are the coordinates of the term in the partial fraction expansion that is to be updated. See the description of X.
* * * * * * * *	х	(input) real is a three-dimensional array. X(I,J,K) represents the Kth coefficient of the numerator of the (I,J)th term in the partial fraction expansion. Namely, that term with the Ith factor of DEN to the Jth power as denominator.

* * * * * * * *	DEGX	(input) integer is an array. DEGX(I,J) represents the degree of the numerator of the (I,J)th term in the partial fraction expansion. See the description of X.
* * * * * *	F	(input) real is an array to containing the coefficients of the polynomial to be embedded in the partial fraction matrix.
~ * *	DEGF	(input) integer is the degree of the polynomial, F.
* *	REIURN:	Not used.
* * *	ROUTINES CALLED:	None.
* * *	AUTHOR:	James F. Stafford
* * *	DATE CREATED:	6Jun87 Version 1.0
* * *	REVISIONS:	None.
* *****	****	******
	SUBROUTINE	GETD (J, A, DEGA, DEN, DEGD)
	IMPLICIT	NONE
	INTEGER	J, K, DEGA, DEGD(*)
	REAL*8	A(0:*), DEN(0:2,*)
	DEGA=DEGD (J)	
	DO K=0,DEGA	
	A(K)=DEN(K, S	נ)
	ENDDO	
	RETURN	

```
END
```

(TTDDCT/TTD)	
SUBROUTINE	GETX(J,K,F,DEGF,X,DEGX)
IMPLICIT	NONE
INTEGER	J, K, L, DEGF, DEGX(10,*)
REAL*8	F(0:*),X(0:1,10,*)
DEGF=DEGX(J,K)	
DO L=0,DEGF	
F(L) = X(L, J, I)	K)
ENDDO	
RETURN	
END	
SUBROUTINE	PUTX (J, K, F, DEGF, X, DEGX)
IMPLICIT	NONE
INTEGER	J,K,L,DEGF,DEGX(10,*)
REAL*8	F(0:*),X(0:1,10,*)
DEGX(J,K)=DEGF	
DO L=0,DEGF	
X(L, J, K) = F(I)	L)
ENDDO	
RETURN	
END	
SUBROUTINE	GET(A, DEGA, B, DEGB)
IMPLICIT	NONE
INTEGER	I, DEGA, DEGB
REAL*8	A(0:*),B(0:*)

DEGA=DEGB DO I=0,DEGA A(I)=B(I) ENDDO RETURN

AN ALGEBRAIC APPROACH TO COMPUTING INVERSE LAPLACE TRANSFORMS OF RATIONAL FUNCTIONS

by

JAMES FLOYD STAFFORD

B.S. Kansas State University, 1986

AN ABSTRACT OF A MASTER'S THESIS submitted in partial fulfillment of the requirements for the degree

MASTER OF SCIENCE

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Abstract

The usual method for solving linear, constant-coefficient. differential equations involves use of the Laplace transform. The most difficult step in this method of solution is computing the inverse Laplace transform of a rational function. The object of this thesis is to describe an algorithm for solving large systems of this kind. The thesis demonstrates that the problem of solving such systems can be treated completely algebraically once the denominator of the rational function is factored. It is shown that the number of operations required to compute a suitable partial fraction expansion of a rational function can be reduced by factoring the denominator into irreducible linear and quadratic factors in $\mathbf{R}[\mathbf{x}]$. Applications to control theory are discussed. The algorithms are derived with mathematical rigor. Working FORTRAN programs implementing the derived algorithms are given in an appendix. Electrical engineers solving practical problems in circuit analysis and control theory might find these algorithms useful. ~