

A STRUCTURED PROGRAMMING PREPROCESSOR
FOR A VARIETY OF BASE LANGUAGES

by

Mary L. Love

B.S., Kansas State University, 1975

A MASTER'S REPORT

submitted in partial fulfillment of the
requirements for the degree
MASTER OF SCIENCE

Department of Computer Science

KANSAS STATE UNIVERSITY
Manhattan, Kansas

1977

Approved by:


Virgil Wallentine

LD
2668
R4
1977
L68
C.2
Document

TABLE OF CONTENTS

I. Introduction	
Project Definition	1
Motivation	2
Organization of Remainder of Paper	3
II. Techniques of Design and Implementation	
Design	5
Implementation	22
III. User Information Overview	
Input.	26
Output	29
Error Conditions	31
IV. Observations	
Good and Bad Aspects of Project.	33
Overall Appraisal of Direction	34
Time, Effort, and Money Involved	35
Elements to Be Changed	35
Possible Further Developments.	36
Appendix	
I. Detailed Implementation Information.	41
II. Detailed User Information.	82
III. Job Control Language Listings.	97
IV. Example Base Language Specifications	101
V. Preprocessor Source Code Listing	105

ILLEGIBLE DOCUMENT

**THE FOLLOWING
DOCUMENT(S) IS OF
POOR LEGIBILITY IN
THE ORIGINAL**

**THIS IS THE BEST
COPY AVAILABLE**

FIGURES

Figure II.1.	Preprocessor and Language Translator Relationship . .	6
Figure II.2.	Flow of Data in Preprocessor.	18
Figure II.3.	Modularity of the Preprocessor.	20

TABLES

Table III.1.	Prototype Parameters.	27
Table III.2.	Definitions Required for Token Recognition and Code Generation.	28
Table A.I.1.	Error Exits with Additional Information Provided. . .	5A
Table A.I.2.	Scanning Routines	14A
Table A.I.3.	Production Language for Parser.	31A

ACKNOWLEDGMENT

The author wishes to express her deep thanks to her major professor, Dr. Virgil Wallentine, and the other members of her committee, Drs. Linda Shapiro and William Hankley. A sincere thank you also goes to the other students in the Department of Computer Science who listened and gave suggestions and support which contributed toward the completion of this report.

SECTION I. INTRODUCTION

PROJECT DEFINITION

This report describes the design and implementation of a preprocessor which facilitates structured programming in a variety of base languages. The preprocessor is intended to allow and encourage the use of structured programming techniques. It attempts to achieve these goals by translating structured programming control structures, not currently implemented in a particular programming language, into semantically equivalent statements which are standard in the language. The preprocessor also includes a text-replacement macro facility which allows the user to provide additional control constructs.

The user provides three types of input to the preprocessor. The first type consists of information about the base language including prototypes of several simple types of statements, operators, and other syntactic information. The second type of input consists of macro definitions. The third type of input is a program written in an extended version of a programming language. The extended version of a language includes all the usual features of that language plus the structured programming control constructs made available by the preprocessor and any macros for which the user has provided definitions.

The output of the preprocessor is a program which can be translated by the usual compiler for the language which was being extended. This program is provided both in printed form and in a temporary data set which can be used as input to the language translator. The output program is logically equivalent to the input program which was written in the extended language.

MOTIVATION

The first and foremost incentive behind this project is to make the techniques and desirable features of structured programming [Dah72,Lec74] available to programmers despite the programming language they may use. While there currently exist a number of processors designed to provide structured programming control structures for only FORTRAN [ONe74,Hig75], the author's preprocessor may be used with a variety of programming languages. The current implementation works with the PL/I, FORTRAN, and SNOBOL programming languages.

There are several reasons for using a preprocessor to accomplish this task. First, a preprocessor can be built to accomodate multiple languages. Second, the language translator currently being used may be retained because the preprocessor outputs code equivalent to the code written in the extended language but which is standard in the base language. Since the user need not acquire a new translator, programmers do not have to learn an entirely new language, just simple modifications to the current one. Furthermore, old programs do not have to be recoded in a new language. Finally, a preprocessor such as described here is easy to use in conjunction with the existing language translator because the preprocessor's output can be input directly into the existing translator.

A new preprocessor was implemented instead of using an existing general-purpose macro processor [Bro67,Bro69] because the new preprocessor is intended to encourage the use of structured programming as well as make it available. To do this, eight control structures are implemented in the preprocessor, thus relieving the programming of the task of having to define them for use with a general-purpose macro processor. This

definition of control structures would be necessary once for each language being used, whereas they are already defined for all languages in the new preprocessor.

Since the author recognizes the fact that no one person can think of all the structured programming control structures which may possibly be desirable for use at one time or another, a text-replacement macro facility [Sol74] has also been incorporated into the preprocessor. This facility allows the user to define macros which will aid in the programming function. Inclusion of the macro facility introduces the possibility that the preprocessor's goals of providing structured programming capabilities and encouraging their use may be undermined by a user who defines macros which conflict with structured programming principles. Because of this possibility a set of guidelines for macro definitions is provided in Appendix II. If the user follows these guidelines in writing macro definitions, the likelihood of violating structured programming principles in user-defined macros will be decreased.

The preprocessor is written in SNOBOL4 and run under the SPITCOL compiler [Dew71]. This language was chosen because the pattern-matching capability provided is well-suited to token recognition, a major part of this application.

ORGANIZATION OF REMAINDER OF PAPER

The remainder of this paper is organized into three sections. The first section describes the techniques of design and implementation of the preprocessor. The second section presents an overview of information required to use the preprocessor. The third section is a summary of the project. The Appendix of the paper includes details

of both the implementation and use of the preprocessor, the Job Control Language statements required to run the preprocessor, example base language specifications, and a listing of the preprocessor source code.

SECTION II. TECHNIQUES OF DESIGN AND IMPLEMENTATION

DESIGN

Design of Processor

The project was intended to effect a mapping from an extended language to the corresponding base language. The extended version of a programming language consists of all the features normally found in the language plus the structured programming control constructs made available by the processor and macros for which the user has provided definitions. To effect this mapping for a variety of languages, a generalized tool, that is, one capable of processing different languages, is required. Since a preprocessor could be designed to accept multiple languages for processing, the first design decision was to use a preprocessor.

There were other factors supporting this decision. The preprocessor could output code which is standard in the base language. As a result, the user need not acquire a new translator and programmers need to learn only simple modifications to the language they currently use. Programs already in use would not have to be recoded in a new language. Finally, since the preprocessor's output could be input directly into the language translator, it would be easy to use with the existing translator. Figure II.1 illustrates the relationship between the preprocessor and the language translator.

The decision to implement a new preprocessor rather than use an existing general-purpose macro processor stemmed from the fact that one of the goals of the project is to encourage the use of structured programming techniques. If an existing macro processor were used, the

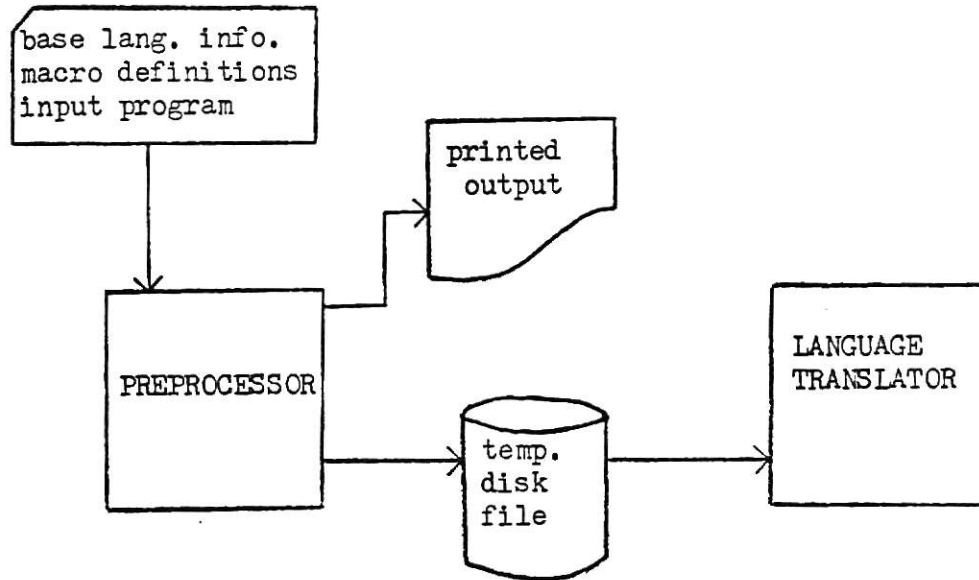


Figure II.1. Preprocessor and Language Translator Relationship.

user would have to define the extended constructs each time the macro processor was used. Furthermore, the extended constructs would have different definitions for each language. Selected structured programming control constructs could be implemented within a new preprocessor. Consequently, the need to redefine the constructs at each use and for each language would be eliminated.

Another design decision made to enable the project to be generalized was to use a table-driven structure parser. The parser was designed based on Production Language [Gri71]. Using a table-driven parser would enable the parser to recognize constructs in different languages because the tokens would be recognized by a scanner using the appropriate definitions.

Eight structured programming control structures were chosen for implementation in the preprocessor. Each of these eight was chosen

either because it receives support in the literature as being a structure necessary for structured programming [Kel,Lec74,Led75,Weg75] or because it is practical to use. The structures included are:

IF - THEN - ENDIF	IF - THEN - ELSE - ENDIF
WHILE - ENDWHILE	REPEAT UNTIL - ENDREPEAT
CASE OF - ENDCASE	DO (iterative) - ENDDO
DOGROUP - ENDDOGROUP	QUITLOOP

The decision to include a text-replacement type macro facility was a difficult one to make. On the one hand it seemed obvious that no one person could think of all the structured programming control structures which may possibly be desirable for use at one time or another. However, it also appeared that to include such a facility would defeat the purpose of the project by allowing users to define any constructs they wished, whether or not those constructs conflicted with the principles of structured programming.

After much deliberation it was decided to include the macro facility. A set of guidelines is provided along with this facility for the user to follow in defining macros. The author feels that by following these guidelines, the user can write macro definitions which have little chance of violating structured programming principles.

A final consideration in the design of the preprocessor was keeping its development and testing time and effort to a minimum. A project of this size, written in a language which does not currently support many structured programming techniques, would certainly be unmanageable if written as one large program. Therefore, to keep the coding and testing as simple as possible, the preprocessor was designed

to be modular. The final version of the preprocessor contains fifty modules in addition to the driver routine.

Design of Execution

The preprocessor may perform one or two parses in processing the input program to produce the final output. If the user has provided any macro definitions, the preprocessor will make a second parse using the output of the first parse as its input. The second parse is required because users are allowed to use the extended constructs which are built into the preprocessor within macro definitions. This second parse will expand these occurrences of the extended constructs into standard code.

Although the main processing of the input program will be the same for all languages, there are some cases for specific languages which must be handled separately. For example, in SNOBOL a logical not operator requires post processing to make the output produced by the preprocessor legal in the standard language. In many other languages the logical not may be formed in a manner similar to an equal relation, that is, with an argument followed by the operator and a second argument, but utilizing a null first argument. In SNOBOL, however, post processing is required to eliminate the comma which appears illegally when a null first argument is used in a logical not relation. That is, the form $\neg((\quad, \text{arg2}))$ must be replaced by $\neg((\text{arg2}))$. Special functions such as this are accomplished by doing post processing for each language which requires it. Currently, the post processing is written into the driver of the preprocessor. However, for generality

it could be provided as a call to an external routine. The appropriate routine would have to be linked with the preprocessor when it is loaded for execution.

Another function of the preprocessor is to insert declarations, if required, for any new variables it introduces. To allow declarations to be properly placed in a program, the user includes an indication of the appropriate location by inserting a special statement which will be replaced by declarations. If the declarations are not required, the special statement should be omitted and the prototype for a declaration statement should be blank.

Design of Extended Constructs

The code which is generated for each of the eight extended constructs is explained briefly in the following paragraphs. Along with each explanation is a simple example designed to aid the reader in visualizing the replacement strategy while eliminating the details of any particular language and the preprocessor's exact output.

IF - THEN - ENDIF: If the negation of the Boolean expression provided in the IF statement is true, a branch is taken to a unique label which will be placed on a null statement at the position of the ENDIF statement. The code following the THEN is inserted following the test of the negation.

IF bool THEN assign	<u>if not bool then go to</u> unilab ₁
	assign
ENDIF	unilab ₁ : <u>null</u>

IF - THEN - ELSE - ENDIF: If the negation of the Boolean expression provided in the IF statement is true, a branch is taken to a unique

label which will be placed on a null statement at the position of the ELSE. The code following the THEN is inserted following the test of the negation. An unconditional branch to a second unique label (unilab_{i+1}) is added at the end of this code. The second unique label will be placed on a null statement at the position of the ENDIF. The code following the ELSE is inserted after the unconditional branch to the second unique label.

IF bool THEN assign	<u>if not bool then go to</u> unilab_i
	assign
	<u>go to</u> unilab_{i+1}
ELSE branch	unilab_i : <u>null</u>
	branch
ENDIF	unilab_{i+1} : <u>null</u>

DOGROUP - ENDDOGROUP: No replacement code is generated for this construct. The code necessary to delineate the code between the DOGROUP and ENDDOGROUP statements as a separate entity is included within the replacement code for the other constructs.

DOGROUP	
assign	assign
call	call
ENDDOGROUP	

DO (iterative) - ENDDO: The four elements extracted from the DO statement are the index variable (INDEX), the beginning value (BEGIN), the ending value (END), and the increment value (INCREMENT). Three statements are inserted at the position of the DO statement. The first statement sets a unique variable (univar_i) equal to END minus INCREMENT. The second sets INDEX equal to BEGIN less INCREMENT. The third statement has a unique label (unilab_i) and assigns INDEX plus INCREMENT to INDEX.

A set of five statements is inserted at the position of the ENDDO statement. The first tests if the increment value is positive and branches to a unique label (unilab_{i+1}) if that is true. The next statement branches to unilab_i if INDEX is greater than or equal to univar_i . The third statement is an unconditional branch to another unique label (unilab_{i+2}). The fourth statement has unilab_{i+1} as the label and branches to unilab_i if INDEX is less than or equal to univar_i . The fifth statement has unilab_{i+2} as the label and assigns INDEX the value of INDEX plus INCREMENT. This rather complex set of statements at the ENDDO position is required because the increment value may be either positive or negative, and the condition for ending the loop differs according to the sign of the increment value.

DO index = begin TO end BY incr	$\text{univar}_i \leftarrow \text{end} - \text{incr}$
	$\text{index} \leftarrow \text{begin} - \text{incr}$
	$\text{unilab}_i: \text{index} \leftarrow \text{index} + \text{incr}$
assign	assign
call	call
ENDDO	<u>if</u> $\text{incr} \text{ gt } 0$ <u>then</u> <u>go to</u> unilab_{i+1}
	<u>if</u> $\text{index} \text{ ge } \text{univar}_i$ <u>then</u> <u>go to</u>
	unilab_i
	<u>go to</u> unilab_{i+2}
	$\text{unilab}_{i+1}: \text{if } \text{index} \text{ le } \text{univar}_i$
	<u>then</u> <u>go to</u> unilab_i
	$\text{unilab}_{i+2}: \text{index} \leftarrow \text{index} + \text{incr}$

WHILE - ENDWHILE: A unique label is placed on a statement which branches to a second unique label if the negation of the Boolean expression provided in the WHILE statement is true. An unconditional branch to the first unique label and a null statement with the second unique label are placed at the position of the ENDWHILE. A WHILE - ENDWHILE will not

be executed if the condition is not satisfied when the loop is encountered.

WHILE bool	unilab _i : <u>if not bool then</u>
	<u>go to</u> unilab _{i+1}
assign	assign
ENDWHILE	<u>go to</u> unilab _i
	unilab _{i+1} : <u>null</u>

REPEAT UNTIL - ENDREPEAT: A null statement with a unique label is placed at the position of the REPEAT UNTIL statement. A statement is inserted at the position of the ENDREPEAT which branches to that unique label if the negation of the Boolean expression specified in the REPEAT UNTIL statement is true. A REPEAT UNTIL - ENDREPEAT will always be executed at least once.

REPEAT UNTIL bool	unilab _i : <u>null</u>
assign	assign
ENDREPEAT	<u>if not bool then go to</u> unilab _i

CASE OF - ENDCASE: A conditional branching statement is inserted at the position of the CASE OF statement which branches to a unique label if the index of the CASE statement is not equal to a test value which is initialized to one. Two statements are inserted after each statement or statement group within the CASE which is treated as a separate entity. The first is an unconditional branch to a unique label which will be placed on the second null statement at the position of the ENDCASE. The second statement has the unique label which was the destination of the last conditional branch and branches to another unique label if the index of the CASE is not equal to the test value which has been incremented by one. These two statements are repeated until all cases for this CASE structure have been processed. At the ENDCASE a null statement with the unique label

which was the destination in the last conditional branch is placed before the null statement described above.

CASE OF index	<u>if</u> index <u>ne</u> 1 <u>then</u> <u>go to</u> unilab _{i+1}
case ₁	case ₁
	<u>go to</u> unilab _i
	unilab _{i+1} : <u>if</u> index <u>ne</u> 2 <u>then</u>
	<u>go to</u> unilab _{i+2}
case ₂	case ₂
	<u>go to</u> unilab _i
	unilab _{i+2} : <u>if</u> index <u>ne</u> 3 <u>then</u>
	<u>go to</u> unilab _{i+3}
case ₃	case ₃
ENDCASE	unilab _{i+3} : <u>null</u>
	unilab _i : <u>null</u>

QUITLOOP: A QUITLOOP is replaced by an unconditional branch to a unique label which is placed on a null statement after the inserted code for the enclosing loop.

WHILE -----	while replacement code
statement ₁	statement ₁
QUITLOOP	<u>go to</u> unilab _i
statement ₂	statement ₂
ENDWHILE	endwhile replacement code
	unilab _i : <u>null</u>

Design of Input

The following is a brief description of the input to the preprocessor. A more detailed description of the input may be found in Appendix II.

There are three types of input to the preprocessor. The first type consists of information about the base language. This information is provided by the user in a form derived from Backus-Naur Form [Gri71]. These definitions are used by the preprocessor in two ways. First, the preprocessor uses part of this information to recognize tokens in the

input program. Second, the preprocessor outputs statements standard in the base language by utilizing this information. Some of the information is used in only one manner. For example, the definition for end-of-statement is used only to recognize the end of a statement in the input program, and the statement prototypes are used only for producing output standard in the base language. On the other hand, some information is used in both ways. An example is the definition for a comment. The preprocessor uses this definition to recognize comments in the input program so they will not be analyzed as statements and also to produce comments in the output program which allow the user to locate where extended statements have been replaced.

The information required for each definition was determined and a definition form was then derived. In some cases, statement prototypes being an example, the definition simply reflects the proper form of the statement using parameter names where actual values would be placed. In other cases column numbers, lengths, and character strings are appropriate. It is also necessary to specify in some cases whether columns or character strings or both are utilized in the recognition of syntactic tokens. The comment definition is again a good example. It may contain column numbers or character strings or both. This example of a comment definition for PL/I contains an opening character string, an opening column number, and a closing character string.

EXAMPLE: Comment definition for PL/I

```
<COMMENT>::='/*' + 10 '*/'
```

The definition of unique labels contains lengths and character strings required for generation while the definition of language labels contains

character strings and column numbers needed for recognition. The following example of a unique label definition for SNOBOL defines labels of the form 'LABxxxxxxNEW' where xxxxxx is a sequence number. The example of FORTRAN language labels contains column numbers for positioning labels, a character set for the first character of a label, and a character set for the remaining characters of a label.

EXAMPLE: Unique label definition for SNOBOL

<UNIQUE LABEL>::=LAB 6 NEW

EXAMPLE: Language label definition for FORTRAN

<LANGLABEL>::=1-5 1-5 '0123456789' '0123456789'

Another piece of information relating to the base language which may be provided by the user concerns whether or not to recognize and expand extended constructs. If a base language has a construct which is recognized by keywords identical to those used by the preprocessor, the user must either specify that the extended construct keywords are to be overridden or not use the feature as provided in the base language. An override card to accomplish this task for the keywords DOGROUP and ENDDOGROUP would appear as follows:

OVERRIDE:DOGROUP,ENDDOGROUP

The second type of input consists of macro definitions. This information is optional. If the user wishes to introduce no macros, this is indicated by omission of any macro definition statements. The user provides a macro name, parameter names, and replacement text in each macro definition. The preprocessor will replace each macro call encountered with the corresponding replacement text, replacing parameter names with the actual arguments of the call.

The third type of input is the program written in the extended version of the programming language for which definitions have been provided. A special statement is included in the program to indicate the proper position for declarations of new variables.

Of some interest is the manner in which the definitions which have been provided are stored in the preprocessor. Each definition has an array associated with it. Although the specific contents of the array elements depend upon the particular definition, the concepts behind the storage of all the definitions are similar. Each array element contains a single piece of information which has been extracted from the definition. In the case of statement prototypes, the first element contains the number of statement components and each of the remaining elements contains one component, either text or a parameter name. For definitions of other syntactic elements, the definition array elements may contain information such as column numbers, character strings, or length requirements provided in the definitions.

The array for the FORTRAN language label definition above would contain the values 2, 'COL', 2, 1, 5, 1, 5, '0123456789', and '0123456789'. The first item indicates a series of column numbers is provided for the label starting position. The next two items indicate a series of column numbers is provided for the label ending position. The next four values specify the two series of columns. The two remaining character strings are the sets of valid characters for the first character of a label and the remaining characters of a label, respectively.

Design of Output

The output produced by the preprocessor is a program logically equivalent to the input program but which is standard in the base language. This output program is produced in two forms, a printed listing and a temporary disk file. The disk file can be input directly into the translator for the base language.

The output program contains at least one comment for each extended construct which is replaced by code produced by the preprocessor. There are also comments placed to indicate code inserted by the preprocessor which is not directly related to the use of a particular extended construct but of which the user should be aware. These comments indicating where new code has been produced are provided to aid the user in determining where errors have been made which relate to the error messages produced by the language translator and in correcting logic errors resulting from incorrect usage of the extended constructs.

The preprocessor's output does not go directly into the print file and the temporary disk file. There are intermediate files which hold these outputs. Figure II.2 illustrates the flow of data in the preprocessor. First the outputs are placed in intermediate files before it is determined whether a second parse will be required. If the second parse is not needed, the files are copied into the next set of intermediate files. If the second parse is required, the output of the second parse is placed in this set of intermediate files. The preprocessor then does any post processing necessary for the language being processed and places the outputs in the third set of intermediate

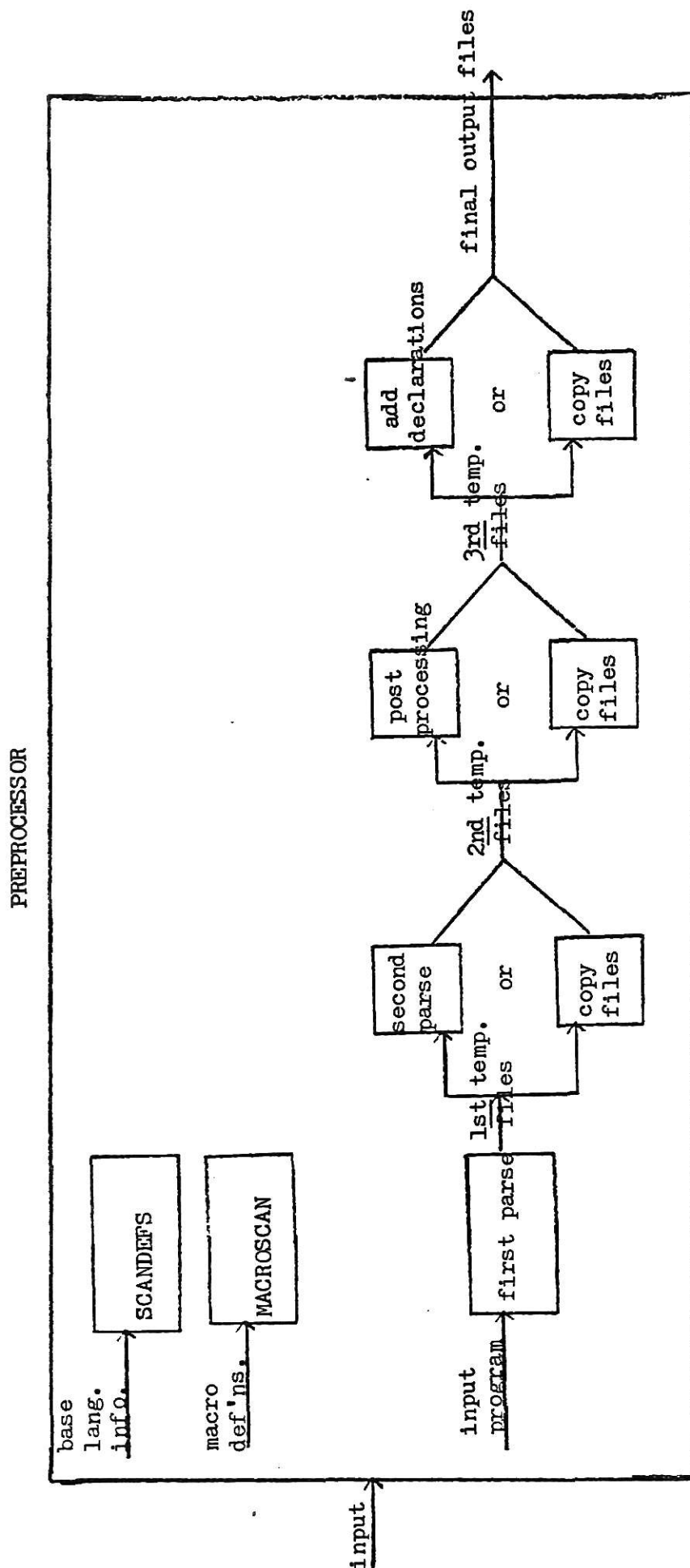


Figure II.2. Flow of Data in Preprocessor.

files. Finally, any required declarations of new variables are added and the outputs are placed in the final output files.

If an error condition occurs at any time, the preprocessor outputs an error message and prints the contents of whichever intermediate punch file contains the results of processing up to the error. This is done so the user can determine how far the preprocessor had progressed and what output had been produced thus far.

Design of Specific Modules

The following is a brief discussion of some of the modules of the implementation. Appendix I contains a detailed description of the final implementation.

Each definition provided by the user has an associated scanning routine which decodes the definition into its components and places them in the appropriate elements of the associated definition array. The scanning of all the definitions except those for macros is controlled by one routine which determines the scanning routine which should be called to process each definition.

The six statement prototypes are used to control the creation of output statements standard in the base language. Each of the prototypes has a statement creation routine with which it is associated which creates one type of statement and is controlled by the prototype.

There are three different insertion routines. One inserts a string with appropriate syntax to be a comment. A second routine inserts a string as a statement, skipping the columns specified in the ignored columns definition. The third insertion routine inserts a string which is already in the proper form for a statement.

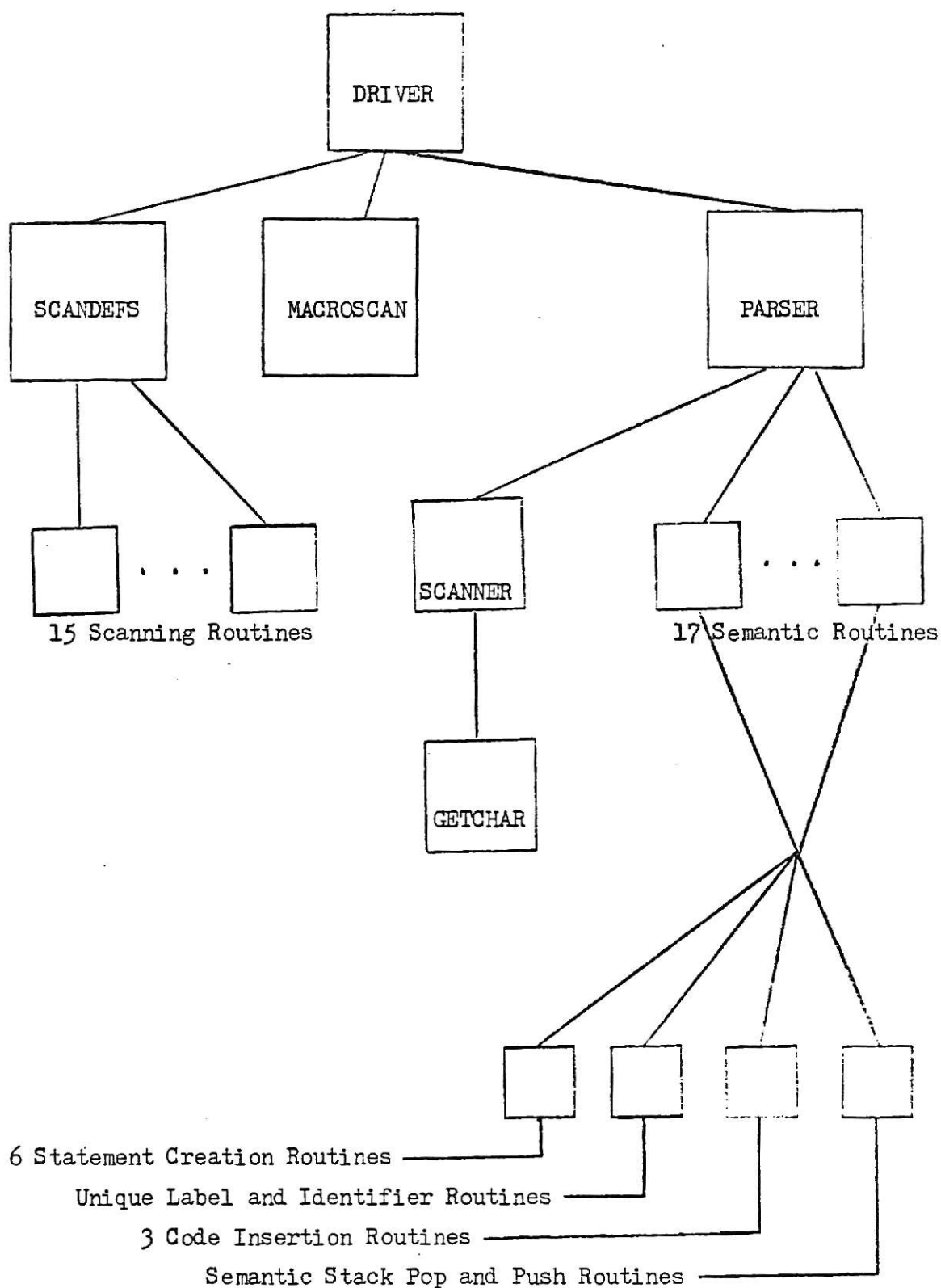


Figure II.3. Modularity of the Preprocessor.

The routine which scans macro definitions is similar in concept to the other definition scanning routines in that it decodes a definition and stores its components in an array. However, the macro scanning routine creates two arrays, one for the parameters in the calling sequence and one for the components of the replacement text. These arrays are placed in elements of a SNOBOL table and are accessed via the macro name.

The structure parser is table-driven and written in a form based on Production Language. It has one syntactic stack which is represented by a character string. It was necessary to order the tests made against this stack such that the terminating symbol of a construct, for example ENDDO, at the stack top would not be mistakenly recognized as the opening symbol, in this case DO. The parser calls a scanning routine to place new tokens at the top of the syntactic stack and seventeen semantic routines to perform code generation and replacement.

The scanner uses much of the information provided by the user about the base language. Literals and comments are recognized as such and, as a result, their contents are not scanned for extended constructs. Another token is placed on the top of the syntactic stack each time the scanner is called. The following example represents the contents of the syntactic stack following the recognition of a DOGROUP statement, an assignment statement, and a REPEAT UNTIL statement. The tokens have been examined and reduced by the parser. The DOGROUP statement is represented by the token DOGROUPSTMT. The assignment statement has been reduced to a structure list (STRLIST) token. The REPEAT UNTIL statement is represented by the REPEATSTMT token. Note that the tokens are in a character string and separated by single blanks.

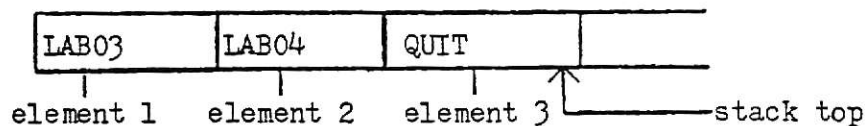
EXAMPLE: Contents of syntactic stack

DOGROUPSTMT STRLIST REPEATSTMT

There are seventeen semantic routines in the preprocessor. The basic strategy of the semantic routines is to make the extended statement into a comment and insert appropriate code which is standard in the base language. One semantic routine copies statements containing no extended constructs or macro calls directly into the output or replaces macro calls with the appropriate replacement text. Several other semantic routines, for example the one which processes the token DOGROUP, simply make extended statements into comments and produce no new code.

The semantic routines use a semantic stack represented by an array to store items to be transmitted between semantic routines. The stack may contain such items as labels, values, variables, and indications of the occurrence of the QUITLOOP construct. The following example of the semantic stack shows two labels and one indicator of the occurrence of a QUITLOOP structure.

EXAMPLE: Contents of the semantic stack.



IMPLEMENTATION

Functional Classification of Modules

There are five functional categories of modules in the preprocessor. The categories are:

- 1) definition scanners
- 2) statement, label, and variable creators
- 3) semantic routines
- 4) service routines
- 5) special purpose modules

There are three types of definition scanners. The first type of definition scanning routine consists of the six routines which decode statement prototypes. The only routine of the second type is the routine which scans macro definitions. The third type consists of the routines which decode the remainder of the definitions.

There are six statement creation routines, one label creation routine, and one variable creation routine in the creator category. Each of the statement creator routines creates one type of output statement. The label and variable creators provide unique labels and identifiers, respectively. The semantic routine category contains the seventeen semantic routines. These are the routines which generate output code sequences.

There are six service routines which perform recurring tasks. Among these are two semantic stack manipulation routines, pop.sem and push.sem, and three string insertion routines. The other routine in this category is the one which retrieves single characters from the input text.

The special purpose modules include the preprocessor driver, the structure parser, the lexical scanner, and the routine which controls definition scanning.

Examples of Output

The following three examples illustrate the preprocessor's output for specific constructs. The first example illustrates the use of a CASE construct in PL/I. The second example is a REPEAT UNTIL structure in SNOBOL. The last example uses FORTRAN to illustrate the IF-THEN-ENDIF construct.

In each of these examples, the assumption is made that the preprocessor has not created any unique labels prior to the processing of the extended construct in the example. Also assumptions are made that the specification examples given for each language in Appendix IV were used by the preprocessor.

EXAMPLE -- PL/I LANGUAGE:

Partial Input Program:

```
CASE OF I;
  B = A + 6;
  B = A * 6;
  B = A - 6;
ENDCASE;
```

Corresponding Output Program:

```

/*    CASE OF I;*/
L0003  : IF(I      =  1      )THEN GO TO L0002      ;
      B = A + 6;
      /*NEXT TWO STATEMENTS INSERTED FOR CASE PROCESSING*/
L0005  : GO TO L0001      ;
L0002  : IF(I      =  2      )THEN GO TO L0004      ;
      B = A * 6;
      /*NEXT TWO STATEMENTS INSERTED FOR CASE PROCESSING*/
L0007  : GO TO L0001      ;
L0004  : IF(I      =  3      )THEN GO TO L0006      ;
      B = A - 6;
      /*NEXT TWO STATEMENTS INSERTED FOR CASE PROCESSING*/
L0009  : GO TO L0001      ;
L0006  : IF(I      =  4      )THEN GO TO L0008      ;
      /*    ENDCASE;*/
L0008  : ;
L0001  : ;
```

EXAMPLE -- SNOBOL LANGUAGE:

Partial Input Program:

```

REPEAT UNTIL (LT(I,3))
  I = I / 6
  J = J + I
ENDREPEAT

```

Corresponding Output Program:

```

* REPEAT UNTIL (LT(I,3))
L000001NEW
  I = I / 6
  J = J + I
* ENDREPEAT
L000002NEW  ¬((LT(I,3)) ) :S(L000001NEW )

```

EXAMPLE -- FORTRAN LANGUAGE:

Partial Input Program:

```

IF (I.LE.6) THEN      X = X + I
ENDIF

```

Corresponding Output Program:

```

C IF (I.LE.6) THEN
10024 IF( .NOT. (I.LE.6) ) GO TO 10014
      X = X + I
C ENDIF
10014 CONTINUE

```

SECTION III. USER INFORMATION OVERVIEW

This overview is not intended to be used as a user's guide for the preprocessor. Appendix II contains the detailed information required for its use.

INPUT

There are three types of input to the preprocessor. The first type of input consists of information about the base language in a form derived from Backus-Naur Form. This portion of the input is described in the following paragraphs.

There are six statement prototypes provided by the user. These are patterns for the preprocessor to follow in creating output statements. The prototypes are as follows:

unconditional branch	prototype for an unconditional branching statement
conditional branch	prototype for statement which tests a condition and branches if the condition is true
declaration	prototype for a statement which identifies a variable to the compiler; if no declarations are required, this prototype is blank
null	prototype for statement which serves no function but to mark a location with a label
assignment with plus	prototype for a statement which does a single addition operation and places the result in a variable
assignment with minus	prototype for a statement which does a single subtraction operation and places the result in a variable

Prototype	Parameters
unconditional branch	label, destination
conditional branch	label, value 1, relation, value 2, destination
declaration	variable name
null	label
assignment with plus	label, variable, value 1, value 2
assignment with minus	label, variable, value 1, value 2

Table III.1 Prototype Parameters.

Each prototype has specific parameter names in it. The parameters for the prototypes are shown in Table III.1.

Each of the parameter names will be replaced by an appropriate value during generation of output code by the preprocessor. All the parameter names for each prototype must be present except in the declaration statement prototype. Since this prototype may be entirely blank under certain conditions, a test is not made for the absence of the parameter name.

The definitions required by the preprocessor for token recognition and output code generation are given in Table III.2. The definition for each of these items is represented in a special form based on the types of information and how much information is required. For example, there are four different ways in which a continuation card may be indicated. The first method is by a special continuation symbol. Another method is to start in a certain column. The third way is to

Definition	Information Provided
end of statement	recognition criteria for end of statement
continuation	recognition criteria for continuation card
language label	recognition criteria for a label in the input program
literal	recognition criteria for a literal string
comment	recognition and generation criteria for a comment
ignored columns	specification of columns whose contents are ignored in recognition and generation of statements
unique label	generation criteria for unique labels
unique variable	generation criteria for unique identifiers
greater than, greater than or equal, less than or equal, not equal, logical not	symbol strings for these Boolean operators

Table III.2. Definitions Required for Token Recognition and Code Generation.

use a special continuation symbol in a particular column. The final way is to have no indication at all, as in PL/I. Consequently, the continuation definition must convey which method is utilized and the information relevant to the method.

The user has the option of overriding the recognition of an extended construct keyword if that keyword will conflict with a

feature already present in the language. For example, if the user wishes to code the IF-THEN and IF-THEN-ELSE constructs already available in PL/I instead of using the extended constructs, it would be necessary to include the keywords IF, THEN, ELSE, and ENDIF in the list on the OVERRIDE card. This would allow the user to include the desired constructs and prevent the preprocessor from mistaking them for extended constructs.

The second type of input is macro definitions. The use of macros is optional. Therefore, this part of the input may be omitted. Each macro definition has a name, parameter names, and replacement text for the macro. When the replacement text is inserted at a macro call, the parameter names in the replacement text are replaced by the actual arguments of the macro call.

There is a set of guidelines for writing macro replacement texts given in Appendix II. The user is encouraged to follow these guidelines to reduce the possibility of writing macro definitions which conflict with the principles of structured programming.

OUTPUT

Output is produced by the preprocessor in two forms, as a printed listing and a temporary disk file. The disk file may be input directly into the translator for the base language.

Two messages will be produced in the printed output which do not appear in the disk file. These messages indicate the success or failure of the parses. If a second parse is not required, a message to that effect is printed.

Each extended statement appears as a comment immediately before the code inserted by the preprocessor at the position of the extended statement. A comment appears every place a statement is inserted by the preprocessor, even if an extended statement has not been replaced by the inserted code. An example of this situation is the insertion of statements between the cases of a CASE structure even though no extended statement appears between the cases.

Each statement label appearing in the input program is detached from the associated statement and placed on a null statement. This null statement immediately precedes the original statement. This is necessary because a label and a statement cannot be reassociated after the statement has been recognized as an extended structure or macro call and replaced by inserted code. To maintain the proper locations of labels, they are placed on the null statements and positioned as described above.

There will be redundant comments and null statements when two parses have been made in processing the program. This is a result of label processing. The preprocessor will place all labels on null statements during the first parse. On the second parse, the label on each null statement created for a label during the first parse will be detached and placed on another null statement, and a comment will be generated which notes this insertion. Consequently, each label in the input program will result in two comments and two null statements when two parses are made during processing.

Post processing may alter this situation. For instance, the post processing for the PL/I language takes the label which has been

removed from the PROCEDURE OPTIONS statements and replaces it on that statement while eliminating the null statement(s) created for the label. For FORTRAN, the statement labels are returned to FORMAT statements in post processing.

ERROR CONDITIONS

The preprocessor does not detect errors in the base language portion of the input program. This task is left to the base language translator.

The preprocessor is a rather unforgiving tool. When an error is detected, the preprocessor outputs information pertinent to the cause of the error and terminates. Depending on what type of error occurred, the contents of the syntactic stack and certain variables may or may not be output. For any type of error, an error message and the output program produced so far are printed.

The common types of errors fall into three categories. The first type of error consists of all errors in specifying information about the base language and macro definitions. Among the errors falling in this category are incorrect definition syntax, missing definitions, invalid definition contents, and a missing ADDED DECLARATIONS statement when one is required. The user must be especially careful to provide a full eighty bytes in each of the six statement prototypes. This area is the only one in which the preprocessor is at all forgiving. If the preprocessor encounters an input record not containing a definition when one is expected, the preprocessor will skip that record. Normally this condition occurs when a statement prototype is not provided as a

full eighty bytes and the preprocessor uses the beginning of the next definition as the end of the previous one. Consequently, the preprocessor will later detect a missing definition and terminate processing.

The second type of error consists of omitting a keyword in a construct -- for example, leaving out the `ENDIF` following an `IF - THEN`. The preprocessor does not recognize this error until it attempts to process the ending keyword of an enclosing construct or at the end of the input program.

The third type of error includes syntax errors in writing extended constructs. This type of error generally results in abnormal termination because some expected keywords are not recognized. A good example of this type of error is the mistake of not leaving column one of a `SNOBOL` statement blank and placing a keyword for an extended construct in that position. The keyword will be recognized as a label instead of the appropriate keyword. This type of error may be considered a subset of the second type since the preprocessor does not recognize this error as a syntax error, but as an omission error.

When the preprocessor determines that an expected keyword is missing, either by omission or syntax error, the current contents of the syntactic stack are printed. This will help the user in determining what keyword was not recognized. The output program produced so far is printed to enable the user to determine at what point in analyzing the input the error was encountered.

SECTION IV. OBSERVATIONS

GOOD AND BAD ASPECTS OF PROJECT

Naturally, there are both good and poor aspects of this project. The good points include the fact that SNOBOL has been implemented on a variety of computers produced by various manufacturers [Gri72], so the preprocessor could be ported from machine to machine without a great deal of revision. Only statements incompatible between SNOBOL and SPITBOL would have to be modified.

Another good aspect of the preprocessor is that it is generalized for use with a variety of languages. Even though the macro feature is somewhat unsophisticated, this feature contributes to the preprocessor's generality. One could also note that the preprocessor is easy to use with regard to the syntax of the input program since the original syntax of the programming language is changed only minimally.

The modular design and implementation of the preprocessor is a good point about the project. This became quite apparent during total system testing. The total system tests were completed in only eight days.

Among the poor aspects of the project is the fact that both the storage utilization and execution of the preprocessor are inefficient. A more experienced SPITBOL programmer could probably have produced a more efficient product. However, efficiency was not a major goal of this project.

The overhead involved in the use of the preprocessor is enormous at this point. To require the user to provide and the preprocessor

to decode all the base language definitions for each execution is undesirable and costly. This matter is addressed later in this section.

The fact that the user must provide the base language information for each execution is admittedly cumbersome and costly, but not overly difficult for the user. Perhaps the first development of the language specifications is somewhat difficult, but this is an initial investment only. However, the user may find it possible to use the language specifications provided as examples in Appendix IV, in which case this poor aspect is minimized.

The author would have liked to use PL/I rather than SPITBOL to code the preprocessor because SPITBOL lacks many structured programming features. However, the pattern matching feature of SPITBOL was of considerable value in writing the preprocessor.

OVERALL APPRAISAL OF DIRECTION

The author feels that the concept behind the project, that is, to create a generalized tool to provide structured programming capability, is a valid one and extremely desirable. However, it is also recognized that the current implementation will likely receive no utilization outside the academic environment. This limited use will be a result of the lack of efficiency of the final product, the cumbersome requirements for specifying base language information, and the high cost of execution. Consequently, the author feels the project was a success as far as demonstrating that the concept is viable, but the final product will receive a minimum amount of practical use.

TIME, EFFORT, AND MONEY INVOLVED

The time which the project required is difficult to provide accurately because the project was stretched out over an extended period of time. From the initial design efforts to the completion of documentation was a time span of two years. Therefore, the following figures are based on person days and not calendar days.

Design	--- 2 months
Coding	--- 1 month
Unit Tests	--- 1 month
Integrated Tests	--- 8 days
Documentation	--- 2 months

Costs for computer time utilized in the testing phases amounted to \$642.35. This figure was based on the following rates:

CPU time	1 min. = 9 units
K-byte time	1 min. = .037 unit
cards read	1 card = .0006 unit
lines printed	1 line = .0005 unit
pages printed	1 page = .007 unit
disk EXCP's	1 EXCP = .0002 unit
1 unit = \$1	

ELEMENTS TO BE CHANGED

The author hopes that one of the first changes to be implemented in the preprocessor will be to allow specifications for a language to be provided once and only once. After being provided once, they would be stored on tape or disk and retrieved when execution begins. A user would also be given the option of providing a new set of specifications to temporarily override those specifications stored in the system. This option was provided in the original design and subsequently eliminated during implementation.

Another desirable change would be to increase the efficiency of the preprocessor, both in storage utilization and execution time. A programmer knowledgeable in the intricacies of SPITBOL could accomplish this task.

A change to the keyword overriding policy would be useful. To allow the user to specify an alternate form of a keyword instead of simply eliminating it completely would alleviate a problem in using the preprocessor with the PL/I language. Currently, to use the macro facility with PL/I requires that there be no THEN keyword in the output program produced by the first parse unless it is part of an extended IF - THEN construct. This means no THEN keyword could be produced during the first parse except in the expansion of a macro call, but many of the extended constructs produce that particular keyword. It seems, therefore, that the user would be able to write very little useful code outside the macros.

POSSIBLE FURTHER DEVELOPMENTS

There are obviously possibilities for further development. One of these is to implement a more sophisticated macro facility. The macro feature should not be made elegant, because the preprocessor was not intended as a general purpose preprocessor. However, the techniques for manipulation of formal parameters and actual arguments might be upgraded to allow similar names for parameters and semantic stack manipulation.

One other possibility is to allow the user to use continuation cards in extended construct statements when continuation is indicated

by a symbol. This would involve deleting the continuation symbol when it is encountered and replacing it in all occurrences of continuation cards.

REFERENCES

- Abr75 Abrahams, Paul. "Structured Programming" Considered Harmful, SIGPLAN Notices, Vol. 10, No. 4, April 1975, pp. 13-24.
- Bau74 Bauer, F.L., and J. Eickel, eds. Compiler Construction -- An Advanced Course, Springer-Verlag, Berlin, 1974, pp. 356-365.
- Bro67 Brown, P.J. ML/I Macro Processor, Comm. ACM, Vol. 10, No. 10, October 1967, pp. 618-623.
- Bro69 Brown, P.J. A Survey of Macro Processors, Annual Review in Automatic Programming, Vol. 6, Part 2, 1969, pp. 37-88.
- Bro74 Brown, P.J. Macro Processors and Techniques for Portable Software. Wiley and Sons, New York, 1974, pp. 53-58.
- Cha71 Chandler, Gerald D. META PI -- A Language for Extensions, SIGPLAN Notices, Vol. 6, No. 12, December 1971, pp. 8-9.
- Cah72 Dahl, O.-J., E.W. Dijkstra, and C.A.R. Hoare. Structured Programming. Academic Press, London, 1972, pp. 1-82.
- Den74a Denning, Peter J. Is It Not Time to Define "Structured" Programming?, Operating Systems Review, Vol. 8, No. 1, pp. 6-7.
- Den74b Denning, Peter J. Is "Structured Programming" Any Longer the Right Term?, SIGPLAN Notices, Vol. 9, No. 11, November 1974, pp. 4-6.
- Dew71 Dewar, Robert B.K. SPITBOL, Version 2.0. Illinois Institute of Technology, February 1971.
- DiJ75 Dijkstra, E.W. Guarded Commands, Nondeterminacy, and Formal Derivation of Programs, Comm. ACM, Vol. 18, No. 8, August 1975, pp. 453-457.
- Fel66 Feldman, Jerome A. A Formal Semantics for Computer Languages and Its Application in a Compiler-Compiler, Comm. ACM, Vol. 9, No. 1, January 1966, pp. 3-9.
- Fel68 Feldman, Jerome, and David Gries. Translator Writing Systems, Comm. ACM, Vol. 11, No. 2, February 1968, pp. 77-113.
- Fer71 Fernandez, A., J.C. Heliard, and J.D. Ichbiah. Overview of the Syntax Processor SYNPROC, SIGPLAN Notices, Vol. 6, No. 12, December 1971, pp. 51-55.

- Fou74 Foulk, C.R. Yet Another Attempt to Define Structured Programming, Operating Systems Review, Vol. 8, No. 2, July 1974, pp. 4-5.
- Fri74 Friedman, Daniel P., and Stuart C. Shapiro. A Case for While-Until, SIGPLAN Notices, Vol. 9, No. 7, July 1974, pp. 7-14.
- Gri71 Gries, David. Compiler Construction for Digital Computers. Wiley and Sons, New York, 1971, pp. 11-48, 155-169.
- Gri72 Griswold, Ralph E. Macro Implementation of SNOBOL. W.H. Freeman, San Francisco, 1972, pp. 240-242.
- Hig75 Higgins, Donald S. A Structured FORTRAN Translator, SIGPLAN Notices, Vol. 10, No. 2, February 1975, pp. 42-48.
- Kel Keller, Roy F. A Modern Beginning Programming Course. Computer Science Department and Ames Laboratory, ERDA, Ames, Iowa.
- Lec74 Lecarme, Olivier. Structured Programming, Programming Teaching, and the Language Pascal, SIGPLAN Notices, Vol. 9, No. 7, July 1974, pp. 15-21.
- Led75 Ledgard, Henry, and Michael Marcotty. A Genealogy of Control Structures, Comm. ACM, Vol. 18, No. 11, November 1975, pp. 629-639.
- Luc73 Lucido, Anthony P. META-T and the Implementation of Special Application Languages, Proc. Computer Science and Statistics: 7th Annual Symposium on the Interface, 1973, pp. 412-417.
- Mat75 Mathis, Robert F. Flow Trace of a Structured Program, SIGPLAN Notices, Vol. 10, No. 4, April 1975, pp. 33-37.
- McC75 McClure, Carma L. Structured Programming in COBOL, SIGPLAN Notices, Vol. 10, No. 4, April 1975, pp. 25-33.
- Mei74 Meissner, Loren P. A Compatible "Structured" Extension to FORTRAN, SIGPLAN Notices, Vol. 9, No. 10, October 1974, pp. 29-36.
- ONe68 O'Neil, John T., Jr. META PI -- An On-Line Interactive Compiler-Compiler, Proc. Fall Joint Computer Conference, 1968, pp. 201-218.
- ONe74 O'Neill, Dennis M. SFOR -- A Precompiler for the Implementation of a FORTRAN-Based Structured Language, SIGPLAN Notices, Vol. 9, No. 12, December 1974, pp. 22-29.
- San74 Sanfield, Stuart H. The Scope of Variable Concept: The Key to Structured Programming?, SIGPLAN Notices, Vol. 9, No. 7, July 1974, pp. 22-29.

- Shn74 Shneiderman, Ben. The Chemistry of Control Structures, SIGPLAN Notices, Vol. 9, No. 12, December 1974, pp. 29-34.
- Sol74 Solntseff, N., and A. Yezerski. A Survey of Extensible Programming Languages, Annual Review in Automatic Programming, Vol. 7, Part 5, 1974, pp. 267-307.
- Vau74 Vaughn, W.C.M. Another Look at the CASE Statement, SIGPLAN Notices, Vol. 9, No. 11, November 1974, pp. 32-36.
- Weg75 Wegner, Eberhard. Control Constructs for Programming Languages, SIGPLAN Notices, Vol. 10, No. 2, February 1975, pp. 34-41.

APPENDIX I. DETAILED IMPLEMENTATION INFORMATION

This Appendix contains detailed information about the implementation of the preprocessor. It is intended as a guide to anyone wishing to make modifications to the source code. Consequently, comprehension will be severely limited unless the reader refers frequently to the source listing in Appendix V. The remainder of this Appendix describes the implementation by categories of modules as follows: driver, definition scanning routines, macro definition scanner, creation routines, service routines, semantic routines, insertion routines, parser, scanner, and character retrieval routine.

Driver

The driver of the preprocessor defines output associations, function calling sequences, and arrays. Patterns and variables are also initialized.

The driver reads the language name and checks it against those which the preprocessor can analyze. Each input card is read and appropriate processing is initiated. An OVERRIDE card is processed directly by the driver. The first definition card encountered causes the preprocessor to call SCANDEFS. The preprocessor calls the MACRCSCAN routine when the first card of a macro definition is read. The preprocessor will begin processing the input program when the first card not recognized as a definition or an override card is read. The parse number indicator is set to one, and the parser is called. When control is returned to the driver, a test is made to determine if the parse was successful.

If the parse was unsuccessful, a branch is taken to the error exit for an unsuccessful first parse. If the parse was successful, a message indicating such is printed, the output files are rewound, and new input and output associations are defined. If no macro definitions were provided, a branch is taken around the second parse and the temporary print and punch files are copied to the next intermediate files. Otherwise, variables are reset to appropriate initial values, and the parse number indicator is set to two. The parser is called again, and a test for success is made when control is returned to the driver. If an error has occurred, a branch is taken to an error exit. Whether one or two parses are made, the intermediate files are rewound and new input and output associations are defined.

If the declaration statement prototype is non-blank, declarations are added for all new variables introduced by the preprocessor. The statement containing the string 'ADDED DECLARATIONS' is found in the output files and a declaration is added at that point for each variable name in the array VARLIST. Comments are also inserted to indicate that these declarations were added by the preprocessor. The remainder of the files are then copied. If no ADDED DECLARATIONS statement is found, an error exit is taken.

The next step is to execute post-processing if required for the language being processed. The intermediate files are rewound and new input and output associations are defined. The language type is checked.

If the language is SNOBOL, each statement is scanned for the `¬(` symbols followed by blanks and `, (`. At each occurrence, the blanks and comma are deleted.

If the language is FORTRAN, all sequences of 'comment-null statement-FORMAT statement' and 'comment-comment-null statement-null statement-FORMAT statement' are found. In each case, the statement number from the first null statement is placed on the FORMAT statement. The comments and the null statements are deleted.

If the language is PL/I, the first records of the output files are scanned for either one or two comments followed by the same number of null statements and a procedure statement. The label from the first null statement is placed on the procedure statement. The comments and the null statements are deleted, and the remainders of the files are copied.

If any other language is being processed (possible only if the language type pattern, LANGPAT, is enlarged to include other languages), the files are simply copied to the final output files.

The preprocessor prints a successful completion message and terminates.

The section to print the appropriate intermediate files on an error condition tests the parse number and rewinds the appropriate intermediate punch file. An input association is defined for the rewound file, and an output association with a real printer file is defined. The contents of the temporary punch file to be displayed are read and printed. An unsuccessful termination message is also printed and execution halts.

All error exits are provided in the driver. An explanatory message is printed for each error. Other information is also provided for certain errors as shown in Table A.I.1. If the preprocessor has begun analyzing the input program, a branch is taken to the section to print the intermediate file contents.

Definition Scanning Routines

One routine, SCANDEFS, controls all definition scanning except that for macro definitions. When the first definition is encountered, the first card of that definition is passed to this controlling routine. The definition is processed as described below, and another definition is read. After all nineteen definitions have been processed, or an attempt has been made to decode nineteen definitions, control is returned to the calling routine.

To process a definition, the SCANDEFS routine determines which one of fifteen individual scanning routines to call, based on the definition name. The definition name, its enclosing angle brackets, and the ::= symbol following it are stripped away. The remainder of the definition is then passed to the routine being called. If the definition is for a statement type, the routine reads another card and passes the entire prototype to the routine. If the definition is that for a language label, the routine determines whether or not the entire definition is contained on the single card. If not, the routine reads another input card and passes the entire definition to the routine for the LANGLABEL definition. If an attempt is made to read a definition, but the card contains an invalid definition name or no definition name in angle

<u>ERROR EXIT</u>	<u>ERROR CONDITION</u>	<u>ADDITIONAL INFORMATION</u>
LANGERR	invalid language type	language type provided in input
ENDPARSEERR	first parse terminated abnormally	contents of syntactic stack
END2ERR	second parse terminated abnormally	contents of syntactic stack
GTHANDEFERR	error in definition of greather-than operation	definition provided in input
LEQUDEFERR	error in definition of less-than-or-equal oper.	definition provided in input
GEQUDEFERR	error in definition of greater-than-or-equal oper.	definition provided in input
NOTDEFERR	error in definition of not operator	definition provided in input
NEQUDEFERR	error in definition of not-equal operator	definition provided in input
MACCALLERR	error in macro calling sequence definition	macro name
PSCANERR	no value returned from scanner	input statement being processed, contents of syntactic stack
PUNIDERR	parser unable to identify contents of syntax stack	input statement being processed, contents of syntactic stack
PSTRERR	failure to parse STRUC to STRLIST reduction	input statement being processed, contents of syntactic stack
PEDOGRPERR	failure to parse ENDDOGROUP	input statement being processed, contents of syntactic stack
PEDOERR	failure to parse ENDDO	input statement being processed, contents of syntactic stack

Table A.I.1. Error Exits with Additional Information Provided.

<u>ERROR EXIT</u>	<u>ERROR CONDITION</u>	<u>ADDITIONAL INFORMATION</u>
PIFTHERR	failure to parse IF - THEN clause	input statement being processed, contents of syntactic stack
PEWHERR	failure to parse ENDWHILE	input statement being processed, contents of syntactic stack
PECASEERR	failure to parse ENDCASE	input statement being processed, contents of syntactic stack
PQUERR	failure to parse QUITLOOP	input statement being processed, contents of syntactic stack
PWHERR	failure to parse WHILE	input statement being processed, contents of syntactic stack
PREPUERR	failure to parse REPUNTIL	input statement being processed, contents of syntactic stack
PREPERR	failure to parse REPEAT	input statement being processed, contents of syntactic stack
PIFERR	failure to parse IF	input statement being processed, contents of syntactic stack
PDGRPERR	failure to parse DOGROUP	input statement being processed, contents of syntactic stack
PDOERR	failure to parse DO	input statement being processed, contents of syntactic stack
PCASEERR	failure to parse CASE	input statement being processed, contents of syntactic stack
PEREPERR	failure to parse ENDREPEAT	input statement being processed, contents of syntactic stack

Table A.I.1. Error Exits with Additional Information Provided.(cont'd).

<u>ERROR EXIT</u>	<u>ERROR CONDITION</u>	<u>ADDITIONAL INFORMATION</u>
ELSEERROR	error in semantic processing of ELSE	portion of current input statement not processed by semantic routine
CASEERROR	error in semantic processing of CASE	portion of current input statement not processed by semantic routine
DOERROR	error in semantic processing of DO	portion of current input statement not processed by semantic routine
IFERROR	error in semantic processing of IF	portion of current input statement not processed by semantic routine
REPERERROR	error in semantic processing of REPEAT	portion of current input statement not processed by semantic routine
WHERROR	error in semantic processing of WHERE	portion of current input statement not processed by semantic routine
QUITERROR	error in semantic processing of QUIT	portion of current input statement not processed by semantic routine
ECASEERROR	error in semantic processing of ENDCASE	portion of current input statement not processed by semantic routine
EIFERROR	error in semantic processing of ENDIF	portion of current input statement not processed by semantic routine
EWHILEERROR	error in semantic processing of ENDWHILE	portion of current input statement not processed by semantic routine
EREPERERROR	error in semantic processing of ENDREPEAT	portion of current input statement not processed by semantic routine
EDOERROR	error in semantic processing of ENDDO	portion of current input statement not processed by semantic routine

Table A.I.1. Error Exits with Additional Information Provided (cont'd).

<u>ERROR EXIT</u>	<u>ERROR CONDITION</u>	<u>ADDITIONAL INFORMATION</u>
EDOGRPFERROR	error in semantic processing of ENDDOGROUP	portion of current input statement not processed by semantic routine
DCLBLKERR	too few blanks in DECLARATION STATEMENT prototype to accommodate actual argument	actual argument
GOTOBLKERR	too few blanks in GO TO STATEMENT prototype to accommodate actual argument	input statement being processed
PLUSBLKERR	too few blanks in ASSIGN PLUS STATEMENT prototype to accommodate actual argument	input statement being processed
MINBLKERR	too few blanks in ASSIGN MINUS STATEMENT prototype to accommodate actual argument	input statement being processed
NULLBLKERR	too few blanks in NULL STATEMENT accommodate actual argument	input statement being processed
IFBLKERR	too few blanks in IF THEN GO TO STATEMENT prototype to accommodate actual argument	input statement being processed

Table A.I.1. Error Exits with Additional Information Provided (cont'd).

brackets, input cards are skipped until another definition is encountered. If the error occurs when the routine is attempting to read the nineteenth definition, control returns to the calling routine without reading any additional input cards.

The individual scanning routines called by SCANDEFS are summarized in Table A.I.2 and fall into four categories. The first category contains those routines which scan statement prototypes. Each of these routines has a single parameter which contains the prototype. Each prototype is broken down by the individual routines into components. Each parameter name and the blanks immediately following it comprise a component. All other parts before and between these components also constitute components. For example, the prototype for a GO TO STATEMENT contains two parameters -- a label (LAB) and a destination (DES). If the prototype for a GO TO STATEMENT were 'LAB BRANCH DES', the contents of the first component would be null. The second component would contain 'LAB '. The third component would contain 'BRANCH ', and 'DES' would be in the fourth component. Additional columns are considered to contain blanks through column eighty. The maximum number of components possible for each definition is based on the number of parameter names in the prototype. The numbers of components are as follows:

<u>DEFINITION NAME</u>	<u>MAX. NUMBER OF COMPONENTS</u>
IF THEN GO TO STATEMENT	11
GO TO STATEMENT	5
NULL STATEMENT	3
ASSIGN PLUS STATEMENT	9
ASSIGN MINUS STATEMENT	9
DECLARATION STATEMENT	3

The components for each prototype are placed in an array. The zeroeth element of each array is used to record the number of components. For the example given above, the zeroeth element would contain 4. The arrays associated with the definitions are as follows:

<u>DEFINITION NAME</u>	<u>ARRAY AND DIMENSIONS</u>
IF THEN GO TO STATEMENT	IFPAT(0:11)
GO TO STATEMENT	GOTOPAT(0:5)
NULL STATEMENT	NULLPAT(0:3)
ASSIGN PLUS STATEMENT	PLUSPAT(0:9)
ASSIGN MINUS STATEMENT	MINUSPAT(0:9)
DECLARATION STATEMENT	DCLPAT(0:3)

Each of the scanning routines which decodes a statement prototype functions in basically the same way. The parameter names for that statement type are placed in a SNOBOL pattern. The definition is searched for pairs of components, the second of which is a parameter name. The first component of the pair may be null, and the second component will include any blanks following the parameter name. If

the entire prototype has been scanned after entering the component pair into the appropriate array, the routine immediately records the number of components. If the last portion of the definition is not a parameter name, a component pair will not be found when that is all that remains to be scanned. Therefore, the remainder of the prototype is placed in the array and the number of components is then recorded. In either case, once the number of components has been placed in the array, the routine tests to see if all the necessary parameters have been specified. The exception to this procedure is in the routine which scans the DECLARATION STATEMENT prototype. Since a blank prototype may be specified for this statement type, no check is made for the appropriate number of parameter names.

Another category of definition scanning routines contains those routines which scan the definitions of unique identifiers, unique labels, literals, and continuations. These scanning routines are similar because the numbers and types of parameters are fairly stable. Both the identifier and label definitions must have a string followed by an integer. Another string may optionally follow the integer. The literal definition contains two strings, and a third string is optional. The continuation definition contains either the single word NONE or a string followed by a column number. Each of these components is placed in an element of the array associated with the definition as indicated by the following:

<u>DEFINITION NAME</u>	<u>ARRAY AND DIMENSIONS</u>
UNIQUE IDENTIFIER	VARPAT(1:3)
UNIQUE LABEL	LABELPAT(1:3)
LITERAL	LITPAT(1:3)
CONTINUATION	CONTPAT(1:2)

The third category of scanning routines contains those routines which scan the definitions for end-of-statement, ignored columns, comment, and language label. Each of these definitions has alternative forms of information, the types of which must be recorded in the array associated with the definition. The arrays and the definitions with which they are associated are:

<u>DEFINITION NAME</u>	<u>ARRAY AND DIMENSIONS</u>
END OF STATEMENT	ECSPAT(0:1)
IGNORE COLUMNS	IGNPAT(0:2)
COMMENT	COMPAT(1:6)
LANGLABEL	LANLABPAT(1:9)

The END OF STATEMENT definition may specify either a symbol string or a column number. Based on the contents of the definition, either 'STR' or 'COL' is placed in the zeroeth element of ECSPAT. The string or column number specified is placed in the next element. The IGNORE COLUMNS definition may specify either a single column number or a sequence of column numbers separated by a hyphen. If only one column number is specified, 'ONE' is placed in the zeroeth element, and the column number is placed in element one. If a sequence is specified.

'SEQ' goes into the zeroeth element, and elements one and two will contain the two column numbers given as the bounds of the sequence.

The COMMENT definition contains two sets of information. Each set may be a symbol string, a column number, or both with a '+' between them. Element one of the associated array contains an indication of what is present for the comment opener. Element four contains similar information for the comment closer. The information is specified as a string with the following meanings:

<u>STRING</u>	<u>MEANING</u>
10	symbol string only
01	column number only
11	symbol string and column number

Elements two and three contain the string and column number, if present, for the opener. Elements five and six contain the string and column number, if present, for the closer.

The language label definition contains four pieces of information. The first may be either a string, column number, or a sequence of columns represented by two column numbers separated by a hyphen. The first element of the language label array contains either 1 or 2, for one column number or a sequence of them, respectively. The fourth element contains the single column number or the first of the two column numbers specified as the sequence bounds. The fifth element contains the second column number of the sequence specification. The second piece of information may be of the same form as the first or may be a symbol string. If the column or column sequence form is specified, the

<u>ROUTINE NAME</u>	<u>'PATTERN' CREATED</u>	<u>PARAMETERS</u>
SCANNULLDEF	NULLPAT(0:3)	NULLDEF
SCANGOTODEF	GOTOPAT(0:5)	GOTODEF
SCANIFDEF	IFPAT(0:11)	IFDEF
SCANPLUSDEF	PLUSPAT(0:9)	PLUSDEF
SCANMINUSDEF	MINUSPAT(0:9)	MINUSDEF
SCANDCLDEF	DCLPAT(0:3)	DCLDEF
SCANVARDEF	VARPAT(1:3)	VARDEF
SCANLABELDEF	LABELPAT(1:3)	LABELDEF
SCANLITDEF	LITPAT(1:3)	LITDEF
SCANCONDEF	CONTPAT(1:6)	CONDEF
SCANCOMDEF	COMPAT(1:6)	COMDEF
SCANEOSDEF	EOSPAT(0:1)	EOSDEF
SCANIGNCOLDEF	IGNPAT(0:2)	IGNCOLDEF
SCANLANLABDEF	LANLABPAT(1:9)	LANLABDEF
SCANOPERDEF	GTHAN GEQU LEQU NEQU NOT	TYPEDEF, OPERDEF
SCANDEFS		CARD

Table A.I.2. Scanning Routines.

information is stored as indicated above, but in elements three, six and seven. If a symbol string is given, 'STR' is placed in element two, and the string is placed in element six. The third and fourth pieces of information in this definition are character strings. They are placed in elements eight and nine, respectively.

The fourth and final category of scanning routines contains a single routine which scans the definitions for the five relations -- greater than, less than or equal, greater than or equal, not equal, and not. Each of these definitions has a single variable associated with it in which the symbol string for the appropriate operator is placed. The definitions and their associated variables are as follows:

<u>RELATION</u>	<u>VARIABLE</u>
GREATER THAN	GTHAN
LESS THAN OR EQUAL	LEQU
GREATER THAN OR EQUAL	GEQU
NOT EQUAL	NEQU
NOT	NOT

Macro Definition Scanner

The macro definition scanner (MACROSCAN routine) is not included with the other scanning routines because it has some basic structural and technique differences from those routines. It is called directly from the driver. It is called once for each macro definition and is therefore likely to be called more than once per execution of the preprocessor. The information extracted is stored in a SNOBOL table as opposed to simply an array.

The macro scanning routine is called when the first card of a macro definition is encountered by the driver. This card contains the string '<MACRO CALL>:=' followed by the name of the macro. The macro name is immediately followed by its formal parameters separated by commas and enclosed within a single set of parantheses.

EXAMPLE:

```
<MACRO CALL>:=WHEN(X,Y)
```

MACROSCAN deletes the angle brackets and their contents before extracting the macro name and the parameters. A second card is read if the parameter list extends past the end of the first card. No blanks, imbedded or at the end of the first of two cards, are allowed. The zeroeth element of the MPARM array is set equal to the number of formal parameters provided. The first parameter is placed in element one of MPARM and succeeding parameters are placed in consecutive elements of MPARM.

The next card is read and should contain only the string '<MACRO DEFINITION>:='. The cards following are read and their contents placed in the MACREPSTR variable until the end of the definition is found. The end is signified by a card containing only the string 'END MACRO DEFINITION' beginning in column one. This string is not added to MACREPSTR.

The parameter list pattern is started with a series of 81 blanks. The formal parameters are then added to this pattern from the MPARM array. The macro scanner next breaks the MACREPSTR contents into components and stores those components in the MPAT array. The zeroeth element of MPAT is assigned the number of components in the replacement text. A component is defined as for the statement scanning routines.

The contents of the MPAT array are copied into the macro table MACROTAB indexed by the macro name in the form of a string. The parameter list for the macro, currently in the MPARM array, is copied into the element of MACROTAB indexed by the macro name in the form of a string suffixed with the character 'P'. The routine then outputs a message that this macro definition has been processed and returns to the driver.

Creation Routines

There are eight routines which create items for the preprocessor. Six of them create statements. The other two create unique identifiers and labels. Each of the statement-creating routines has a number of parameters equal to the number of parameter names which appear in the corresponding statement prototype. Each of these routines utilizes the array containing the appropriate statement components. By replacing the components containing parameter names with the corresponding arguments supplied in the call to the creating routine and including enough blanks to retain the format and length of the prototype (to eighty characters) in the output, the routines create statements which are standard in the base language.

The routines which create statements, their outputs and parameters are as follows:

<u>ROUTINE NAME</u>	<u>OUTPUT STATEMENT</u>	<u>PARAMETERS</u>
CRNULL	NULL STATEMENT	LAB
CRIF	IF THEN GO TO STATEMENT	LAB, VAL1, RL, VAL2, DES
CRGOTO	GO TO STATEMENT	LAB, DES
CRMINUS	ASSIGN MINUS STATEMENT	LAB, VAR, VAL1, VAL2
CRPLUS	ASSIGN PLUS STATEMENT	LAB, VAR, VAL1, VAL2
CRDCL	DECLARATION STATEMENT	VARNAME

The routines which create unique identifiers and labels also utilize the arrays associated with the definitions provided. VARNUM and LABELNUM are the variables used to make the identifiers and labels unique. Each routine uses the appropriate variable and then increments it by one. Enough leading zeroes are used with these variables, which are initialized to one, to create items of the lengths specified in the definitions. The routine which creates unique identifiers adds each one to the array VARLIST which is indexed using the variable NEXTVAR. This array is later used to add declarations to the program if required. Neither UNILABEL nor UNIVAR has any parameters. The routines and their outputs are as follows:

<u>ROUTINE NAME</u>	<u>OUTPUT</u>
UNIVAR	unique identifier
UNILABEL	unique label

Service Routines

Two service routines are present in the preprocessor which manipulate the semantic stack. The routines are POP.SEM and PUSH.SEM which pop items off and push items onto the semantic stack, respectively.

The semantic stack is represented by the array SEMSTACK which contains ten elements. Element 1 represents the element at the bottom of the stack. The variable TOP contains the subscript of the top element in the stack.

Each of these routines may have from two to seven parameters. The first parameter is always the number of items to be acted upon, that is, the number of items to pop from or push onto the stack. The second through seventh parameters are passed differently in the two routines. Variables are passed to the POP.SEM routine in which the popped items will be returned. These variables' names are passed as character strings, and values are assigned to the corresponding arguments using two levels of indirection. Variables passed to the PUSH.SEM routine as parameters two through seven contain values to be pushed onto the stack. The values are assigned from the corresponding arguments using only one level of indirection.

POP.SEM extracts a value from the top of the stack and decrements the stack pointer, TOP. PUSH.SEM increments TOP, the stack pointer, and assigns a value to the stack's new top element. Each routine loops until the number of values indicated by the first argument have been processed.

Semantic Routines

Most of the semantic routines call the statement creation routines. These routines are always called with a label value. This is necessary because if a symbol is normally used for terminating a label, an extraneous symbol will appear in the output if a null value is passed as the label.

Semantic Routine 0 - This semantic routine is called when the parser recognizes a label on a statement. The routine inserts a null statement with the label just recognized. If a language label is terminated by a symbol, that symbol is removed from the label before the null statement is created. This step is necessary because the terminating symbol is already provided in the null statement prototype.

Semantic Routine 01 - This routine processes items recognized as unidentified-to-end-of-statement (U-EOS) and macro calls. Any sequence of one or more items terminating in the end-of-statement condition and not containing any token recognized as part of an extended construct is reduced to the U-EOS token. This routine checks the string recognized as U-EOS to determine if it is a macro call.

If the U-EOS is a macro call, the actual arguments of the macro call are broken out and placed in an array. The macro's formal parameters and the number of parameters, as well as the macro's replacement text components are retrieved from the macro table. The parameters are placed in a SNOBOL pattern. The components of the replacement text are copied into a string, replacing the formal parameter names with actual arguments. The macro call is made into a comment and inserted into the output, followed by the expanded macro text.

If the sequence recognized as an U-EOS does not contain a macro call, the statement is copied directly from the input to the output.

Semantic Routine 1 - This routine is called by the parser when the extended construct keyword ELSE has been recognized. The routine

pops a label name from the semantic stack and creates another unique label. An unconditional branch to the new unique label and a null statement with the label popped from the semantic stack are inserted into the output, following a comment indicating that the ELSE keyword was processed at that point. The unique label created in this routine is pushed onto the semantic stack.

Semantic Routine 2 - This routine processes the statement which begins a CASE extended construct. The index of the CASE is extracted and two unique labels are created. A counter indicating the nesting level of the structure is incremented by one. Another counter which records the number of CASE structures currently started but not finished and a counter indicating which case within the current CASE structure is being processed are incremented. A single conditional branch is created which branches to the second unique label created in this routine if the value of the index of the CASE is not equal to one (the value of the case counter for this CASE structure). This statement is inserted into the output after the CASE statement in the form of a comment. The index of the CASE, the label which was the destination of the conditional branch, and the first unique label created in this routine are pushed onto the semantic stack.

Semantic Routine 3 - DO statements are processed by this semantic routine. The items of the DO -- the index, beginning value, ending value, and increment value -- are broken out of the statement. If no increment value is specified, a default value of 1 is used. A unique label and a unique identifier are created, and the counter

indicating the nesting level of the structure is incremented. Two assignments with subtraction and one assignment with addition are created and inserted in the output. The first subtraction assigns the unique identifier the value of the ending value less the increment value. The second subtraction assigns the index the beginning value less the increment value. The addition has the unique label and assigns the index the value of the index plus the increment. The DO statement is made into a comment and inserted into the output, followed by the three created statements. The increment value, the index of the DO, the unique identifier, and the unique label are pushed onto the semantic stack.

Semantic Routine 4 - A DOGROUP statement is processed by this routine. The DOGROUP statement is made into a comment and inserted in the output. The counter indicating the structure nesting level is incremented.

Semantic Routine 5 - When the parser recognizes an IF - THEN statement, this routine is called to process it. The Boolean condition is extracted from the statement and a unique label is created. A conditional branch which tests the negation of the Boolean condition and branches to the unique label on truth of the negation is inserted in the output following the IF - THEN statement in the form of a comment. The unique label is pushed onto the semantic stack.

Semantic Routine 6 - This routine processes a REPEAT UNTIL statement. The Boolean condition is extracted, a unique label is created, and the nesting level counter is incremented. A null statement with the unique label is created and inserted after a comment containing the REPEAT UNTIL. The unique label and the Boolean condition are pushed onto the semantic stack.

Semantic Routine 7 - A WHILE statement is processed by this semantic routine. The Boolean condition of the WHILE is extracted and two unique labels are created. The nesting level counter is incremented. A conditional branch with the first unique label is created which branches to the second unique label if the negation of the Boolean condition is true. The WHILE statement is converted to a comment and inserted in the output, followed by the created statement. The second and first unique labels are pushed onto the semantic stack, in that order.

Semantic Routine 8 - This semantic routine processes the QUITLOOP extended construct. A unique label is created and an unconditional branch to that label is created. The QUITLOOP statement is converted to a comment and inserted in the output, followed by the created branching statement. The unique label and the string 'QUIT' are pushed onto the semantic stack.

Semantic Routine 9 - The parser calls this routine when an ENDCASE has been recognized. Three items are popped off the semantic stack. Two null statements are created with the second and first items from the stack as labels. The number of cases for this CASE structure is reset to zero. The number of CASE structures and the nesting level are both decremented by one. The created statements are inserted following the ENDCASE in the form of a comment.

Semantic Routine 10 - This semantic routine is called by the parser to process an ENDIF. One item is popped from the semantic stack. There are two possible types of values for this item, the string 'QUIT' or a label. If the item is the character string 'QUIT', two more items are popped off the stack and the first of

those two items and the string 'QUIT' are pushed back onto the stack. If the item is not the string 'QUIT', it should be a label.

A null statement using the item from the stack is created and inserted into the output following the ENDIF in the form of a comment.

Semantic Routine 11 - This semantic routine processes the ENDWHILE statement. Two items are popped off the semantic stack and the nesting level counter is decremented by one.

If the first item popped off the semantic stack is the character string 'QUIT', the second item is the label which was the destination of the branch created for a QUITLOOP structure occurring within the WHILE - ENDWHILE construct. The variable indicating the presence of an enclosed QUITLOOP structure is set to one. Two more items are popped off the semantic stack. These items are now considered to be the first and second items popped off the stack.

Whether or not a QUITLOOP structure has occurred, an unconditional branching statement is created with the first item popped off the stack as the destination. A null statement with the second item popped off the stack as the label follows the unconditional branch. These two statements are inserted in the output following the ENDWHILE statement in the form of a comment.

If the QUITLOOP structure indicator was set, another null statement is created. This statement has the label which was the destination of the unconditional branch created for the QUITLOOP occurrence and is inserted in the output.

Semantic Routine 12 - The ENDREPEAT statement is processed by this semantic routine. Two items are popped off the semantic stack and the nesting level counter is decremented by one.

If the first item popped off the semantic stack is the character string 'QUIT', the second item is the label which was the destination of the branch created for a QUITLOOP structure occurring within the REPEAT UNTIL - ENDREPEAT construct. The QUITLOOP occurrence indicator is set to 1 and two more items are popped off the semantic stack. These items are now considered to be the first and second items popped off the stack.

Whether or not a QUITLOOP structure has occurred, the second item popped off the stack is used as the destination of a conditional branching statement which tests the negation of the first item popped off the stack. The ENDREPEAT is converted to a comment and inserted in the output, followed by the created conditional statement.

If the QUITLOOP occurrence flag was set, a null statement is created with the label used as the destination of the branching statement which replaced the QUITLOOP statement. This null statement is inserted in the output.

Semantic Routine 13 - The parser calls this semantic routine when an ENDDO has been recognized. Four items are popped from the semantic stack and the nesting level counter is decremented by one.

If the first item popped off the semantic stack is the character string 'QUIT', the second item is the label created as the branching destination during QUITLOOP processing. Two more items are popped from the semantic stack. These are added to the third and fourth

items popped previously, and the four items are considered to be the first four items popped from the stack. The QUITLOOP occurrence flag is set to one.

Whether or not a QUITLOOP structure has occurred, two unique labels are created. Five statements are created to be inserted in the output following the ENDDO statement in the form of a comment.' The first statement is a conditional branch which tests if the increment value (the fourth item popped from the stack) is positive. The destination of this statement is the first of the two new unique labels. Another conditional branching statement is the second created statement. It has the first item popped off the stack as its destination and tests if the index value (third item from the stack) is greater than or equal to the second item popped from the stack. An unconditional branch with the second unique label as its destination is created next and is followed by another conditional branch. This statement has the first unique label as its label, tests if the index value is less than or equal to the second item popped from the stack, and branches to the first item popped off the stack if the condition is true. The fifth created statement is an assignment statement with an addition operation. The statement's label is the second new unique label and it increments the index value by the increment value.

If the QUITFLAG occurrence flag was set, a null statement is created and inserted in the output. The label on this statement is the destination label of the unconditional branch created during QUITLOOP processing.

Semantic Routine 14 - This semantic routine processes the ENDDOGROUP statement. The statement is converted to a comment and inserted in the output. The nesting level counter is decremented by one.

Semantic Routine 15 - This semantic routine is called whenever the parser reduces an item or series of items to the STRLIST (structure list) token.

The routine returns immediately if there are no unfinished CASE structures or if the preprocessor is processing statements within another structure contained in a CASE structure. Otherwise, the number of cases for this CASE structure is incremented by one, and a unique label is created. Three items are popped from the semantic stack.

An unconditional branch with the first item popped from the semantic stack as the destination is created. A conditional branch with the second item from the semantic stack as the label is created next. This statement tests if the third item from the semantic stack (index of the CASE) is not equal to the counter incremented above and branches to the new unique label if the condition is true.

A comment noting that two statements are being inserted is created and inserted in the output. The two branching statements are inserted immediately following the comment.

The index of the CASE, the new unique label, and the first item popped from the stack are now pushed onto the stack, in that order.

Insertion Routines

There are three routines which insert strings into the output, each in its own way.

Insertcom - The single parameter COMSTR contains a string to be converted to a comment and inserted in the output. The routine will attempt to start the comment in column one. If this column is to be ignored, the routine determines where to start the comment. If a sequence of columns has been specified for the ignored columns in the IGNORE COLUMNS definition, the preprocessor discovers the first column which is not an ignored column. If no sequence was specified, the preprocessor uses the first column following the ignored column.

If a comment begins in a specific column, the columns preceding that column are filled with blanks. If a symbol string starts the comment in that column, the string is added, followed by the parameter. If no symbol string is needed, the parameter is added to the comment being built. The comment is truncated, if necessary, to eighty characters.

If a comment ends in a particular column, the string is truncated or padded with blanks, as necessary. If a comment ends with a certain string, the current comment string is truncated as necessary to maintain only eighty characters. If the comment will end in an ignored column, enough characters are deleted within the parameter string portion of the comment to ensure the comment ender will not appear in an ignored column. The comment closing string is added after the appropriate position has been determined. Enough blanks are added, if necessary, to make the final string eighty characters long. Finally, the comment is inserted in the output.

Insertstmt - The single parameter STR.INSRT contains the string to be inserted in the output as a statement. A counter of the positions in the output string (I) and a counter of the positions

in the parameter string (CHARKTR) are set to zero. The variable TYPE is set depending on whether a single column or a sequence of columns was specified in the IGNORE COLUMNS definition.

CHARKTR and I are both incremented by one. The final portion of the statement has been constructed if CHARKTR is equal to the number of characters in the parameter string. An entire card has been constructed if I is equal to 80. Branches are made if either of these conditions occurs.

A branch is made to check the current position in the output (value of I) against ignored columns. If only one column is ignored and the column currently indicated by I is that column, a blank is inserted. If a sequence of columns is specified and the column indicated by I is one of the ignored columns, a blank is inserted. A branch is made to increment I and continue processing this statement. In either case, if the column to which I points is a usable column, the character in the position of the parameter string indicated by CHARKTR is added to the output string being constructed and a branch is made to increment both CHARKTR and I.

When an entire card image has been constructed, it is inserted in the output. I is reset to 1 and the card image area is nullified. The preprocessor then branches to check for a column to ignore.

When the last portion of a statement has been completed, this portion is padded with blanks and inserted in the output.

Insert - The single parameter INSRT.STR contains the statement to be inserted. This string contains blanks in ignored columns. The statement is broken into sections of eighty characters and

inserted in the output one at a time. If the last section of the statement is less than eighty characters long, this portion is padded with blanks before inserting it in the output.

Parser

The structure parser has no parameters. It is called once by the driver for each parse desired. Control is returned to the point of invocation when the entire input program has been analyzed. Error conditions take special error exits out of the parser.

When the parser is called, it immediately requests a token be placed on the syntactic stack by the scanner. The parser then tests the contents of the syntactic stack, reduces those contents to appropriate tokens, and calls semantic routines and the scanner according to the strategy of the the modified Production Language in Table A.I.3.

Scanner

The scanning routine retrieves single characters from the input program and determines what syntactic tokens are to be presented to the parser. A character string is built of the input characters, and when an item is recognized, a token is placed on the syntactic stack and return is made to the parser. The scanner makes tests in the following order: comment, literal, end of statement, keyword appearing at end of statement, end of a language label, blank character, extended construct keyword followed by a blank. If the scanner is in the middle of a statement, the test for end of a language label is skipped. If the end of statement condition is not recognized, the test for a

LINE NO	STACK TOP	REDUCE TO	CALL SEMANTIC ROUTINE	SCAN	ON SUCCESS, BRANCH TO TEST ON LINE #
1	LABEL	LABEL	ROUTO	yes	1
2	ENDCASE	ENDCASE	--	yes	26
3	ENDIF	ENDIF	--	yes	27
4	ENDWHILE	ENDWHILE	--	yes	29
5	ENDREPEAT	ENDREPEAT	--	yes	30
6	ENDDO	ENDDO	--	yes	31
7	ENDDOGROUP	ENDDOGROUP	--	yes	32
8	ELSE	ELSE	ROUT1	yes	1
9	QUITLOOP	QUITLOOP	--	yes	25
10	CASE	CASE	--	yes	18
11	DO	DO	--	yes	19
12	DOGROUP	DOGROUP	--	yes	20
13	IF	IF	--	yes	21
14	REPEAT	REPEAT	--	yes	22
15	WHILE	WHILE	--	yes	24
16	STRLIST EOF	PROG	--	no	return
17	UEOS	STRUC	ROUT01	no	33
18	CASE UEOS	CASESTMT	ROUT2	yes	1
19	DO UEOS	DOSTMT	ROUT3	yes	1

Table A.I.3. Production Language for Parser.

LINE NO	STACK TOP	REDUCE TO	CALL SEMANTIC ROUTINE	SCAN	ON SUCCESS, BRANCH TO TEST ON LINE #
20	DOGROUP UEOS	DOGROUPSTMT	ROUT4	yes	1
21	IF UTHEN	IFTHEN	ROUT5	yes	1
22	REPEAT UNTIL	REPUNTIL	--	yes	23
23	REPUNTIL UEOS	REPEATSTMT	ROUT6	yes	1
24	WHILE UEOS	WHILESTMT	ROUT7	yes	1
25	QUITLOOP UEOS	STRUC	ROUT8	no	33
26	CASESTMT STRLIST ENDCASE UEOS	STRUC	ROUT9	no	33
27	IFTHCL ENDIF UEOS	STRUC	ROUT10	no	33
28	IFTHEISE ENDIF UEOS	STRUC	ROUT10	no	33
29	WHILE STMT STRLIST ENDWHILE UEOS	STRUC	ROUT11	no	33
30	REPEATSTMT STRLIST ENDREPEAT UEOS	STRUC	ROUT12	no	33
31	DOSTMT STRLIST ENDDO UEOS	STRUC	ROUT13	no	33
32	DOGROUPSTMT STRLIST ENDDOGROUP UEOS	STRUC	ROUT14	no	33
33	LABEL STRUC	STRUC	--	no	33
34	EISE STRUC	ELSE STRUC	--	no	33

Table A.I.3, Production Language for Parser (cont'd).

LINE NO	STACK TOP	REDUCE TO	CALL SEMANTIC ROUTINE	SCAN	ON SUCCESS, BRANCH TO TEST ON LINE #
35	LABEL ELSESTRUC	ELSESTRUC	--	no	33
36	IFTHCL ELSESTRUC	IFTHCLSE	--	yes	1
37	IFTHEN STRUC	IFTHCL	--	yes	1
38	STRLIST STRUC	STRLIST	ROUT15	yes	1
39	STRUC	STRLIST	ROUT15	yes	1

Table A.I.3. Production Language for Parser (cont'd).

keyword appearing at the end of a statement is skipped. If the blank character test is not satisfied, the test for an extended construct keyword followed by a blank is skipped. A minimum amount of backup is used in the scanner. A more detailed description is contained in the remainder of this section.

The scanner begins by recording whether or not an end-of-statement condition (e-o-s) was encountered in the last card scanned. If the variable EOSNEXT is equal to one, that is, the e-o-s condition should be signalled immediately, a branch is taken to return an e-o-s condition. If EOSNEXT is not one, the GETCHAR function is called to return a character and its column number. If this character is the first character of a token or the first non-blank character of a token, the starting column of the token is recorded. If this column is the first non-ignored column of a card, it is determined whether or not all blanks left over from a previous card should be discarded. If only a comment was found on the previous card, these blanks will be retained. Another condition under which the blanks will be retained is when a comment was not the only thing found on the previous card and an e-o-s condition was found in the previous call to the scanner.

Whether or not the blanks are retained, the next step is to add the new character to the current token. The label start position variable and the comment only indicator are set to zero.

The scanner then examines the token string for the beginning of a comment. There are three different forms of comment openers to look for, and the scanner uses whichever one corresponds to the definition provided by the user. When the column opening string is two characters

long, it is necessary for the preprocessor to inspect the character following an occurrence of the first character of the string. If the second character is not the second of the specified string, a backup procedure is invoked by restoring the first character thought to open the comment and its column number. If a comment opener is not recognized, the scanner branches to test for a literal. If the beginning of a comment is found, the e-o-s flag is reset to the value it contained before the call to the scanner, the comment-only indicator is set to one if nothing else has been encountered in this statement, and a search is made for the end of the comment.

The search for the end of the comment also has three possible methods and is controlled by the definition of a comment for this particular language. If the comment closer is a string of two characters, the backup procedure will again be utilized if necessary. When the end of the comment is found, the token, which consists of the text of the comment and any opening and/or closing strings, is added to the current statement string, and the token variable is set to null. A branch is then taken to the point near the beginning of the scanner where another character is retrieved, thus avoiding the manipulation and test of the e-o-s indicator. The comment is not returned to the parser.

If the test for a comment is unsuccessful, the scanner analyzes the input characters for a literal. The literal beginning may be defined as one or two characters. If a two-character opening is used, the backup procedure is used if necessary. If the beginning of a

literal is not recognized, the scanner branches to test for an e-o-s condition.

If the literal opening is found, a search is made for the end of the literal. In searching for the end, special processing is utilized for any special sequence specified in the literal definition. If the first character of the ending string, which is the first character of any special sequence, is encountered, a test is made for the special sequence if one was provided. If a special sequence is found, the scanner continues to search for the literal end. Otherwise, a test is made for the end of the literal at that point. If a literal ender is not found, the scanner continues to search for it. If the end is recognized, the literal is added to the token string and a branch is taken to the portion of the scanner which processes items not identified as extended constructs or comments.

The next section of the scanner tests for the e-o-s condition. The scanner utilizes the information provided by the user in the end of statement definition to analyze the input in this section. If the end of statement is signified by one or two symbols, the scanner tests for these characters, using the backup procedure for a two-character string if necessary. If the end of statement is signalled by a column, the scanner tests if the column number of the most recently retrieved character is that specified in the end-of-statement definition. If the end-of-statement condition is not recognized, a branch is taken to test for a language label. If what appears to be an end-of-statement condition is recognized and that condition is indicated by a column

number, the next input card is read and tested to determine if it is a continuation. If so, a branch is taken to the section of the scanner which tests for a blank character.

If the end-of-statement condition is recognized, the token string is added to the current statement. It is necessary at this point to determine if an extended construct keyword ends at the end of the statement but has not been recognized. This will be the case if e-o-s is indicated by a symbol and a keyword ends immediately preceding that symbol or if e-o-s is indicated by column number and a keyword ends immediately preceding a blank in the ending column or in the ending column itself. If a keyword is recognized in this manner, it is necessary to set the EOSNEXT variable to one so the scanner will signal e-o-s immediately upon its next invocation. The keyword found is added to the syntactic stack and control is returned to the parser. If a keyword is not found at the end of the statement, the e-o-s flag is set for this statement, the EOSNEXT flag is reset to zero, the 'unidentified to end of statement' (UEOS) token is added to the syntactic stack, the token string is set to null and control is returned to the parser.

The next portion of the scanner tests for the end of a language label if this is at the beginning of a statement. If there is anything in the current statement string followed by a blank or only blanks occur in the current statement string, the test for the end of a language label is not made, and a branch is made to the section of the scanner which tests for a blank character.

The test for the end of a language label is based on the definition provided for language labels. If the definition specifies a column

number or a series of column numbers as the determination of the end of a language label, the first test will be satisfied if the current column is the single column number or within the specified sequence. Otherwise, a branch is taken to the section of the scanner which tests for a blank character.

If the definition specifies a symbol to terminate a language label, the current character is tested against the specified symbol.

If a two-character language label separator is specified, the next character is tested and backup is used if necessary to restore the previous character as the current character. If the symbol string does not match, a branch is taken to the section of the scanner which tests for a blank character.

Once the scanner has determined that this character/column could be terminating a language label, the label is extracted from the token string. It is either the entire token or the last portion of it. The column number the label actually started in is computed. The column number computed is compared with the starting column value(s) specified in the language label definition. If this column number cannot start a language label, any necessary backup is performed and a branch is taken to the section which tests for a blank character. At this point the scanner knows a label has been found if a string is the label terminator. If the label is terminated by column numbers, the first character of the label is checked to see if it is a valid character for the start of a label. Any remaining characters are checked against the characters allowable for the remainder of the label. If either of these tests fails, any necessary backup is performed

and a branch is taken to the section of the scanner which tests for a blank character. If these two tests are passed, the next input character is tested for a label delimiter, that is, a blank or a special character which is allowed directly following a label. If this test fails, backup is performed and a branch is taken to the section which tests for a blank character.

If a label is recognized by the scanner, the label is placed in a comment indicating that it was removed from the statement and placed on a null statement inserted immediately preceding the current statement. This insertion function will be performed by the semantic routine which will be invoked when the scanner returns control to the parser. The label's character positions in the current statement are replaced by blanks, the LABEL token is placed on the syntactic stack, and return is made to the parser.

The next section of the scanner tests for a blank character. A blank encountered at this point in the scanner's execution causes a branch to the section to search for a keyword. If the character is any non-blank character, a branch is made to the BEGSCAN1 label in the scanner. The effect of this branch is to allow the scanner to read the next input character and begin its analysis once again.

The next section tests for a keyword. The first of two tests compares the token string with extended construct keywords which begin or end constructs, such as DO and ENDDO. The second test is against the extended construct keywords which appear in the middle of extended constructs, that is, the keywords THEN and UNTIL. If the token string does not satisfy either of these tests, a branch is taken to the section of the scanner which handles items not yet

recognized. If the keyword for THEN has been recognized, the unidentified-to-then (UTHEN) token is added to the syntactic stack. Otherwise, the keyword itself is added to the syntactic stack. The token string is added to the current statement, the token string is set to null, and control is returned to the parser.

The section which processes unidentified items simply branches to the BEGSCAN1 label in the scanner to read the next input character and begin its analysis again.

A section at the end of the scanner tests for an error condition when the scanner attempts to retrieve another input character at the beginning of the scanner but fails. If the token string is currently all blanks or is of zero length, a test is made to determine whether or not an e-o-s condition was found before this call to the scanner. If so, the failure to return a character indicates the end-of-file condition. The EOF token is added to the syntactic stack and control is returned to the parser. If an e-o-s condition had not been discovered, an unended statement occurs at the end of the input and an error exit is taken. If the token string contains at least one non-blank character when the failure to return a character is encountered, an error exit is taken.

Character Retrieval Routine

The GETCHAR routine retrieves a single character from the input program and returns the character and its column number to the scanner. The two parameters are CHAR and COL, which will contain the character and its column number, respectively, on a successful return. These

arguments are passed as variable names in the form of character strings and one level of indirection is used to assign the appropriate values to them.

The routine first increments the column number counter (global variable COLNUM) modulo 80. If the resultant column number is 1, the routine begins retrieving characters from the next card. This card may have already been read by the scanner in testing for a continuation card. If the card is not available, the routine tests if the end of file condition has been encountered previously, either by this routine or by the scanner in testing for a continuation card. If not, the next card is read into the program. Otherwise, the card which was read for the continuation test is copied into the variable from which the routine retrieves characters. If the end of file condition is encountered when reading a new card in this routine, the GETEND variable is set to 1 and a return with failure is made.

To retrieve a character, the routine tests if the current column is an ignored column as specified in the user's IGNORE COLUMNS definition. If this column is to be ignored, a branch is made to the beginning of the routine so the next character can be analyzed. If this column is a usable one, the character in that column and the column number are assigned to the appropriate arguments and a successful return is made to the scanner.

APPENDIX II. DETAILED USER INFORMATION

The three types of input to the preprocessor are information about the base language, macro definitions (optional), and a program written in the extended version of a programming language. The first type consists of a keyword override card and definitions given in a form similar to Backus-Naur Form. The definitions must be grouped together, and the override card may either precede or follow the definitions.

The override card contains the string 'OVERRIDE:' starting in column one. The keywords to be overridden follow the colon and are separated by commas. The first blank indicates the end of the keywords.

EXAMPLES:

```
OVERRIDE:IF,THEN,ELSE,ENDIF
```

overrides the keywords IF, THEN, ELSE, AND ENDIF

```
OVERRIDE:WHILE
```

overrides the keyword WHILE

All the definitions about the base language begin with the definition name enclosed in angle brackets and followed by ::=.

The definition name and the remainder of the contents of each definition are described below.

UNIQUE LABEL -- The information needed to create unique labels is as follows:

- 1) Uniform beginning -- a sequence of one or more characters which conforms to the rules for writing labels in the

standard language but which is shorter than the maximum allowed length. The preprocessor will use this uniform beginning to begin any labels it creates.

- 2) Variable digits -- an integer specifying how many digits should be used to make the labels unique. The preprocessor will begin counting at 1 using these digits (filling with leading zeros) and will imbed them in any labels it creates.
- 3) Remainder -- a sequence of zero or more characters which conform to the rules for ending labels. This sequence should not include any character(s) used to separate a label from the statement with which it is associated.

The uniform beginning must immediately follow the = and at least one blank must separate the components of the definition.

EXAMPLE:

(for FORTRAN) $\langle \text{UNIQUE LABEL} \rangle ::= 24\ 3$

This example will create labels such as 24001 and 24002.

UNIQUE IDENTIFIER -- The information needed to create unique variable names is as follows:

- 1) Uniform beginning -- a sequence of one or more characters which conforms to the rules for naming identifiers in the standard language but which is shorter than the maximum allowed length. The preprocessor will use this uniform beginning to begin the names of any new identifiers it creates.
- 2) Variable digits -- an integer specifying how many digits should be used to make the identifiers unique. The preprocessor will begin counting at 1 using these digits

(filling with leading zeros) and will imbed the digits in any identifier names it creates.

- 3) Remainder -- a sequence of zero or more characters which conforms to the rules for ending identifier names.

The uniform beginning must immediately follow the = and at least one blank must separate the components of the definition.

EXAMPLE:

(for PL/I) <UNIQUE IDENTIFIER> ::=LAB 3 XY

This example will create identifiers such as LAB001XY and LAB002XY.

LITERAL -- The information the user must provide about literals is as follows:

- 1) Literal opener -- a one- or two-character symbol string which begins a literal.
- 2) Literal closer -- a one- or two-character symbol string which ends a literal.
- 3) Special sequence (optional) -- a two-character symbol string which does not terminate a literal but would cause an error in determining where a literal ends if no special handling were provided because the first character of this special sequence is the same as the single character which normally terminates a literal. An example of this is the string "'" which may appear within a literal in some languages in which a single quote normally terminates a literal.

The literal opener must directly follow the = and at least one blank must separate the definition components.

EXAMPLE:

If ' is both the literal opener and literal closer, and '' is is a sequence requiring special handling (as in PL/I):

$\langle \text{LITERAL} \rangle ::= ' ' ''$

COMMENT -- The information the user must provide about comments is as follows:

- 1) Comment opener -- a one- or two-character symbol string in single quotes which begins a comment; or a column number in which a comment begins; or a one- or two-character symbol string in single quotes which begins a comment followed by a + and the column number in which the symbol string must start to indicate a comment.
- 2) Comment closer -- a one- or two-character symbol string in single quotes which ends a comment; or a column number in which a comment ends; or a one- or two-character symbol string in single quotes which ends a comment followed by a + and the column number in which the symbol string must start to terminate a comment.

The comment opener must directly follow the = and at least one blank must separate the definition components.

EXAMPLE:

If a C must appear in column 1 to indicate a comment, and a comment ends in column 72 (as in FORTRAN):

$\langle \text{COMMENT} \rangle ::= 'C' + 1 \ 72$

LANGLABEL -- The information the user must provide about labels on statements is as follows:

- 1) Label opening delimiter -- the column number of the statement in which the label must begin; or a series of column numbers within which the label must begin represented by two column numbers separated by a hyphen.
- 2) Label closing delimiter -- a one- or two-character symbol string enclosed in single quotes which separates a label from the statement with which it is associated; or a column number in which a label must end; or a series of column numbers within which a label must end represented by two column numbers separated by a hyphen.
- 3) Starting character set -- a set of characters enclosed in single quotes which may start a label,
- 4) Remaining character set -- a set of characters enclosed in single quotes which may be used in all other character positions of a label.

The opening delimiter must begin directly following the = and at least one blank must separate the components of the definition.

EXAMPLES:

(for FORTRAN)

```
<LANGLABEL> ::=1-5 1-5 '0123456789' '0123456789'
```

(for PL/I)

```
<LANGLABEL> ::=2-72 ':' 'ABCDEFGHIJKLMNOPQRSTUVWXYZ@#$_'
'ABCDEFGHIJKLMNOPQRSTUVWXYZ0123456789@#$_'
```

(for SNOBOL)

```
<LANGLABEL> ::=1 1-72 'ABCDEFGHIJKLMNOPQRSTUVWXYZ'
'ABCDEFGHIJKLMNOPQRSTUVWXYZ0123456789._'
```

GREATER THAN, GREATER THAN OR EQUAL,
 LESS THAN OR EQUAL, NOT EQUAL, NOT -- The operator for each logical
 operation should be enclosed in single quotes and directly follow
 the =.

EXAMPLES:

(for PL/I) <NOT EQUAL> ::= ' = '

(for SNOBOL) <LESS THAN> ::= 'LT'

(for FORTRAN) <NOT> ::= '.NOT.'

IGNORE COLUMNS -- The information the user must provide about columns
 to ignore in creating and scanning statements is as follows:

Column numbers -- the number of a column to be completely
 ignored; or a series of column numbers for columns to be
 ignored represented by two column numbers separated by a
 hyphen.

The definition item must immediately follow the =.

EXAMPLE:

(for SNOBOL, FORTRAN, and PL/I) <IGNORE COLUMNS> ::= 73-80

END OF STATEMENT -- The information the user must provide about the end
 of statement condition is as follows:

End of statement indicator -- a one- or two-character symbol
 string in single quotes which indicates the end of a statement;
 or a column number which terminates a statement (either a
 regular statement or on a continuation card).

The indicator must immediately follow the =.

EXAMPLES:

(for SNOBOL) <END OF STATEMENT> ::= 72

(for PL/I) <END OF STATEMENT> ::= ';' ;

CONTINUATION -- The information the user must provide about continuation cards is as follows:

Continuation symbol -- a single character followed by the column number in which that symbol must appear to indicate a continuation card; or the string NONE to indicate that statements may cross card boundaries without specifically indicating continuation.

The continuation symbol or the word NONE must immediately follow the =. If a symbol is indicated, at least one blank must separate it from the column specification.

EXAMPLES:

(for SNOBOL) <CONTINUATION> ::=+ 1

(for PL/I) <CONTINUATION> ::=NONE

GO TO STATEMENT, IF THEN GO TO STATEMENT,

ASSIGN MINUS STATEMENT, ASSIGN PLUS STATEMENT, DECLARATION STATEMENT --

In every statement definition, there are certain conventions which must be followed. At least one blank must appear between a parameter name and the items on either side of it. A parameter name may, however, begin in column 1 of the prototype, i.e., directly following the =. The parameter names must be spelled exactly as specified in the description of each definition.

The definition names of the statements, their function, and their parameters are as follows:

NULL STATEMENT -- a null statement.

Parameter names:

LAB -- label on the null statement

GO TO STATEMENT -- a branching statement.

Parameter names:

LAB -- label on the branching statement

DES -- label of the statement which is the destination
of the branch

IF THEN GO TO STATEMENT -- a conditional statement in which the
result of a true value is a branch.

Parameter names:

LAB -- label on the conditional statement

VAL1 - first value in the condition

RL -- relational operator in the condition

VAL2 - second value in the condition

DES -- label of the statement which is the destination
of the branch

ASSIGN MINUS STATEMENT -- an assignment statement in which the
expression being evaluated is a subtraction.

Parameter names:

LAB -- label on the assignment statement

VAR -- identifier to be assigned a value

VAL1 - minuend

VAL2 - subtrahend

ASSIGN PLUS STATEMENT -- an assignment statement in which the
expression being evaluated is an addition.

Parameter names:

LAB -- label on the assignment statement

VAR -- identifier to be assigned a value

VAL1 - first addend

VAL2 - second addend

DECLARATION STATEMENT -- a declaration statement for one variable which is global in scope and is of type integer.

Parameter names:

VARNAME -- identifier to be declared.

Note: If declarations of identifiers are not required, a blank prototype must be provided.

The first column following the equal sign, column 1 of the prototype, is treated as column 1 of a card image. Therefore, if a label must appear in a particular position and the remainder of the statement must appear in a particular position, the definition must appear in the same form with enough blanks included to position items correctly. Since the user also defines the length of identifier names and labels created by the preprocessor, the user must be sure to allow enough space for these to fit if the parameter name is shorter than the item which will be created to replace it. In some cases, a user's variable will be used in the place of a parameter name. Consequently, if blank spaces are critical, it may be wise to use identifiers with the same number of characters in the names. Since the definition name takes up some space and the preprocessor requires an 80-column prototype, two cards will be needed for each definition. Even if the prototype provided by the user could fit on one card with the definition name, it is necessary to provide two cards, leaving the second one blank.

All variables introduced by the preprocessor are integers and have global scope. Consequently, declarations for these new variables should be placed wherever declarations for global variables are normally

placed in a program written in the base language. To allow the preprocessor to position these declarations correctly a special input statement must be included in the extended program containing only the string ADDED DECLARATIONS. When the preprocessor is ready to add declarations, it will replace this statement with the appropriate declarations. If declarations are not required and such has been specified in the declaration definition, the 'ADDED DECLARATIONS' statement should be omitted.

EXAMPLES:

```
(for FORTRAN) <NULL STATEMENT> ::=LAB CONTINUE
(for SNOBOL) <GO TO STATEMENT> ::=LAB : (DES)
(for PL/I) <IF THEN GO TO STATEMENT> ::= LAB :
            IF VAL1      RL      VAL2      THEN GO TO DES ;
(for PL/I) <ASSIGN MINUS STATEMENT> ::= LAB :
            VAR      = VAL1      - VAL2      ;
(for SNOBOL) <ASSIGN PLUS STATEMENT> ::=LAB
            VAR      = VAL1      + VAL2
(for FORTRAN) <DECLARATION STATEMENT> ::=LAB INTEGER VARNAME
```

Macro definitions have three parts, the macro name and parameters, the replacement text, and an indication of the end of the definition. The first card of a macro definition begins in column 1 with the string ' <MACRO CALL> ::=' followed immediately by the macro name and a list of parameter names enclosed in parentheses. The parameter names are separated by commas and no imbedded blanks are allowed. The next card starts with the string ' <MACRO DEFINITION> ::=' starting in column 1. The remainder of the card is left blank. The next card is the beginning

of the replacement text for the macro. The last card of the replacement text is followed by a card containing only the string 'END MACRO DEFINITION' starting in column 1.

The parameter names appearing in the replacement text will be replaced by the actual arguments of a macro call during code expansion by the preprocessor. Because of storage limitations, there can be no more than twenty components in a macro replacement text, where a component is defined to be a parameter name and following blanks or a string of characters appearing between parameter names.

EXAMPLES:

<MACRO CALL> ::=SWAP(X,Y)

<MACRO DEFINITION> ::=

T = X

X = Y

Y = T

END MACRO DEFINITION

In the above example, SWAP is the macro name. X and Y are the formal parameter names. When the preprocessor expands the call SWAP(A,B), the following code will be inserted:

T = A

A = B

B = T

This macro definition contains the following seven components:

The diagram illustrates the seven components of the macro definition, separated by horizontal lines. The components are numbered 1 through 7:

- 1. T = X
- 2. X = Y
- 3. Y = T
- 4. (blank line)
- 5. (blank line)
- 6. (blank line)
- 7. (blank line)

If this macro were to be used in the extended version of SNOBOL, it could not be called with arguments longer than one character because the space before the equal sign would have to be retained. In FORTRAN the arguments could be one or two characters long.

EXAMPLE:

The call SWAP(A1,A2) would result in the following inserted code:

T = A1

A1= A2

A2= T

No argument may be of longer length than the corresponding parameter name and the blanks immediately following it. A call of this macro with either argument longer than two characters would cause an error during the creation of the code to be inserted.

The same macro could be defined as follows:

<MACRO CALL> ::=SWAP(X,Y)

T = X

X = Y

Y = T

END MACRO DEFINITION

Using this definition, the SWAP macro may be called with arguments of longer length than two characters.

Since the objectives of the preprocessor include encouraging the use of structured programming techniques, the following guidelines are provided for the user to use in writing macro definitions. It is felt that by following these guidelines, a user will be more

likely to write macro definitions of constructs in line with the aims of structured programming.

- 1) Sequentially execute statements within the definition wherever possible. In other words, avoid using branches among statements within the definition.
- 2) Avoid branching out of the definition wherever possible. That is, do not specify branches to labels not within the macro definition.
- 3) Use extended language statements defined in the preprocessor if possible.

Parameter names must be unique and no parameter name may be identical to the beginning of another parameter name of the same macro. This is due to the scanning technique used by the preprocessor in decoding the macro replacement text into its components. For example, the names ABCD and ABCDE may not be used as parameter names in the same macro definition. The preprocessor will fail to recognize ABCDE as a parameter name in the replacement text because it will always assume the first four characters constitutes an occurrence of the parameter name ABCD.

The syntax of the input program is nearly identical to the syntax of the base language involved. Many of the input program statements will be copied directly into the output. Extended statements should follow the same syntax as normal statements, even though the constructs are unknown in the base language. For example, when processing an extended FORTRAN program, the preprocessor will expect statement labels to be entirely numeric and appear in columns 1-5, continuation

to be indicated by a special symbol in column 6, and Boolean expressions to be expressed using the FORTRAN relational operators.

The syntax of the extended constructs statements will be similar to the syntax of other statements in the base language. There are specific requirements, however, which must be followed.

- 1) A blank must follow each extended construct keyword or the keyword must appear at the end of a statement.
- 2) The conditions specified in extended construct statements must be specified in forms appropriate for the base language because these conditions will be copied directly into the expanded text.
- 3) The conditions in the WHILE and REPEAT UNTIL statements must be enclosed in parentheses.
- 4) The following pairs of extended construct keywords must appear in the same statement: REPEAT and UNTIL, IF and THEN, CASE and OF. All other extended construct keywords may not appear in a statement with a second keyword. The ending keywords must appear in statements with no other text. The DOGROUP and QUITLOOP keywords are the only construct-opening keywords which must appear alone.

The special statement to position new declarations should be placed wherever the user wishes to have the declarations added. All variables introduced by the preprocessor are integers with global scope. The user should keep this in mind when inserting the special statement. The statement should contain only the string ADDED DECLARATIONS and any end-of-statement specification normally required by the base language translator.

EXAMPLE: (PL/I)

```
TEST:  PROC OPTIONS (MAIN);
```

```
  ADDED DECLARATIONS:
```

```
  Y = 5;
```

```
  DO I = 1 TO 3;
```

```
    Y = Y - 1;
```

```
  ENDDO;
```

```
END TEST;
```

Restrictions

The user must keep the following restrictions in mind when using the preprocessor:

- 1) Only one QUITLOOP structure may be used within another loop.
- 2) The text following the THEN keyword and its succeeding blank must appear as a normal statement on a separate card. That is, if a specific number of columns must be blank at the beginning of a statement, that same number of columns must be blank in this situation.
- 3) Array elements may not be used as actual arguments in macro calls.
- 4) No continuation cards may be used in statements containing extended constructs.

APPENDIX III. JOB CONTROL LANGUAGE LISTINGS

```

// EXFC SPITROL
/**
/**      JOB STEP TO EXECUTE PREPROCESSOR FROM SOURCE DECK
/**
//TEMPCH DD DSN=&&TEMPCH,DCB=(BLKSIZE=800,LRECL=80,RECFM=FB),
//SPACE=(TRK,(5,1)),UNIT=3330,DISP=(NEW,DELETE)
//TEMPCH2 DD DSN=&&TEMPCH2,DCB=(BLKSIZE=800,LRECL=80,RECFM=FB),
//SPACE=(TRK,(5,1)),UNIT=3330,DISP=(NEW,DELETE)
//TEMOUT DD DSN=&&TEMOUT,DCB=(LRECL=132,BLKSIZE=132,RECFM=F),
//UNIT=3330,SPACE=(TRK,(5,1)),DISP=(NEW,DELETE)
//TEMOUT2 DD DSN=&&TEMOUT2,DCB=(LRECL=132,BLKSIZE=132,RECFM=F),
//UNIT=3330,SPACE=(TRK,(5,1)),DISP=(NEW,DELETE)
//FINOUT DD SYSOUT=A,DCB=(BLKSIZE=133,LRECL=133,RECFM=FA)
//FINPUNCH DD DSN=&&FINPUNCH,DCB=(BLKSIZE=800,LRECL=80,RECFM=FB),
//DISP=(NEW,PASS),UNIT=3330,SPACE=(TRK,(5,1))
//SYSIN DD *
/**      PREPROCESSOR SOURCE AND INPUT DATA
/**

```

JCL to execute preprocessor from source deck.

```

//LKED EXEC PGM=IEWL,PARM='LET,LIST,MAP',
//**
//** JOB STEPS TO EXECUTE PREPROCESSOR FROM OBJECT DECK
//**
//** FIRST STEP LINKAGE EDITS OBJECT MODULE
//**
//** SECOND STEP EXECUTES TEMPORARY LOAD MODULE
//**
//SYSLMOD DD DSN=xxLOADMOD(PROG),DISP=(NEW,PASS),
// UNIT=3330,SPACE=(TRK,(8,5,1))
//SYSUT1 DD UNIT=(SYSDA,SEP=SYSLMOD),DCB=BLKSIZE=1024,
// SPACE=(1024,(200,20))
//SPITPROG DD DSN=SYS1.SPITBOL,DISP=SHR
//SYSLIN DD *
** PREPROCESSOR OBJECT DECK
INCLUDE SPITPROG(SPITPROG)
ENTRY DSINT
/*
//GO EXEC PGM=PROG,
// PARM='L=10K,H=100K,R=20K,C=100000,P=2000,T=120,D=10,N=58',
//STEPLIB DD DSN=xxLOADMOD,DISP=(OLD,DELETE)
//SYSPRINT DD SYSOUT=A
//SYSPUNCH DD SYSOUT=B,DCB=(BLKSIZE=800,LRECL=80,RECFM=FB,BIIFNO=1)
//TEMPCH DD DSN=xxTEMPCH,DCB=(BLKSIZE=800,LRECL=80,RECFM=FB),
// SPACE=(TRK,(5,1)),UNIT=3330,DISP=(NEW,DELETE)
//TEMPCH2 DD DSN=xxTEMPCH2,DCB=(BLKSIZE=800,LRECL=80,RECFM=FB),
// SPACE=(TRK,(5,1)),UNIT=3330,DISP=(NEW,DELETE)
//TEMOUT DD DSN=xxTEMOUT,DCB=(LRECL=132,BLKSIZE=132,RECFM=F),
// UNIT=3330,SPACE=(TRK,(5,1)),DISP=(NEW,DELETE)
//TEMOUT2 DD DSN=xxTEMOUT2,DCB=(LRECL=132,BLKSIZE=132,RECFM=F),
// UNIT=3330,SPACE=(TRK,(5,1)),DISP=(NEW,DELETE)
//FINOUT DD SYSOUT=A,DCB=(BLKSIZE=133,LRECL=133,RECFM=FA)
//FINPUNCH DD DSN=xxFINPCH,DCB=(BLKSIZE=800,LRECL=80,RECFM=FB),
// DISP=(NEW,PASS),UNIT=3330,SPACE=(TRK,(5,1))
//SYSIN DD *
** PREPROCESSOR INPUT DATA
/*

```

JCL to execute preprocessor from object deck.

```

// EXEC PGM=PREPROC,
// PARM='L=16K,H=1000K,R=20K,C=100000,P=2000,T=120,D=10,N=58'
//**
//** JOB STEP TO EXECUTE PREPROCESSOR FROM LOAD MODULE
//**
//STEPLIB DD DSN=DSPP42,PREPROC.LOADMOD,DISP=SHR
//SYSPRINT CC SYSOUT=A
//SYSPUNCH DD SYSOUT=B,DCB=(BLKSIZE=800,LRECL=80,RECFM=FB,BUFNO=1)
//TEMPCH DD DSN=SETEMPCH,UNIT=3330,DISP=(NEW,DELETE)
//SPACE=(TRK,(5,1)),DCB=(BLKSIZE=800,LRECL=80,RECFM=FB),
//TEMPCH2 DD DSN=SETEMPCH2,UNIT=3330,DISP=(NEW,DELETE)
//SPACE=(TRK,(5,1)),DCB=(LRECL=132,BLKSIZE=132,RECFM=F),
//TEMOUT DD DSN=SETEMOUT,DCB=(LRECL=132,BLKSIZE=132,RECFM=F),
//UNIT=3330,SPACE=(TRK,(5,1)),DISP=(NEW,DELETE)
//TEMOUT2 DD DSN=SETEMOUT2,DCB=(LRECL=132,BLKSIZE=132,RECFM=F),
//UNIT=3330,SPACE=(TRK,(5,1)),DISP=(NEW,DELETE)
//FINOUT DD SYSOUT=A,DCB=(BLKSIZE=133,LRECL=133,RECFM=FA)
//FINPUNCH DD DSN=SEFINPCH,DCB=(BLKSIZE=800,LRECL=80,RECFM=FB),
//DISP=(NEW,PASS),UNIT=3330,SPACE=(TRK,(5,1))
//SYSIN DD
//** PREPROCESSOR INPUT DATA
/*

```

JCL to execute preprocessor from load module.

APPENDIX IV. EXAMPLE BASE LANGUAGE SPECIFICATIONS


```

PL/I
<GREATER THAN>::='>'
<NOT>::='!'
<NOT EQUAL>::='!='
<GREATER THAN OR EQUAL>::='>='
<LESS THAN OR EQUAL>::='<='
<NULL STATEMENT>::='LAB' ;
<GO TO STATEMENT>::='LAB' : GO TO DES ;
<IF THEN GO TO STATEMENT>::='LAB' : IF (VAL1 RL VAL2 ) THEN GO TO DES
<DECLARATION STATEMENT>::='DCL VARNAME ;
<ASSIGN MINUS STATEMENT>::='LAB' : VAR =VAL1 - VAL2 ;
<ASSIGN PLUS STATEMENT>::='LAB' : VAR = VAL1 +VAL2 ;
<LITERAL>::=' ' ' '
<COMMENT>::='/* + 10 */
<UNIQUE LABEL>::='L 4
<UNIQUE IDENTIFIER>::='ID 3 ST
<END OF STATEMENT>::=';'
<CONTINUATION>::='NONE
<IGNORE COLUMNS>::='73-80
<LANG LABEL>::='2-72 ' ' ' ABCDEFGHIJKLMNOPQRSTUVWXYZ@# $ '
PQRSTUVWXYZ#_
' ABCDEFGHIJKLMNOPQRSTUVWXYZ@# $ '

```

```

SNOBOL
<NULL STATEMENT>::=LAB
<GO TO STATEMENT>::=LAB
:(DES
<DECLARATION STATEMENT>::=
  <ASSIGN MINUS STATEMENT>::=LAB
  <ASSIGN PLUS STATEMENT>::=LAB
  <IF THEN GO TO STATEMENT>::=LAB
)
<LITERAL>::='',
<COMMENT>::='*' + 1 72
<UNIQUE LABEL>::=LAB 6 NEW
<UNIQUE IDENTIFIER>::=NEWID 3
<GREATER THAN>::='GT',
<NOT>::='!',
<NOT EQUAL>::='NE',
<GREATER THAN OR EQUAL>::='GE',
<LESS THAN OR EQUAL>::='LE',
<LANGUAGE>::='1-72',
'ABCDEFGHIJKLMNQRSTUUVWXYZ0123456789._',
<CONTINUATION>::='+',
<END OF STATEMENT>::='72',
<IGNORE COLUMNS>::='73-80'

```

Sample language specifications for SNOBOL.

```

FORTRAN
<IF THEN GO TO STATEMENT>::=LAB IF(VAL1 RL VAL2 ) GO TO DES
<DECLARATION STATEMENT>::= INTEGER VARNAME
<NULL STATEMENT>::=LAB CCNTINJE
<GO TO STATEMENT>::=LAB GO TO DES
<ASSIGN MINUS STATEMENT>::=LAB VAR =VAL1 - VAL2
<ASSIGN PLUS STATEMENT>::=LAB VAR = VAL1 + VAL2
<GREATER THAN>::='.GT.'
<NOT>::='.NOT.'
<NOT EQUAL>::='.NE.'
<GREATER THAN OR EQUAL>::='.GE.'
<LESS THAN OR EQUAL>::='.LE.'
<LITERAL>::='.'
<COMMENT>::='C' + 1 72
<UNIQUE LABEL>::=10 2 4
<UNIQUE IDENTIFIER>::=1JK 2
<CONTINUATION>::=* 6
<END OF STATEMENT>::=72
<IGNORE COLUMNS>::=73-80
<LANGUAGE LABEL>::=1-5 1-5 '0123456789' '0123456789'

```

Sample language specifications for FORTRAN.

APPENDIX V. PREPROCESSOR SOURCE CODE LISTING

```

**
**
**
&STLIMIT = 150000
ASSOCIATE OUTPUT VARIABLES WITH A PRINTER FILE, A TEMPORARY
'PRINT' FILE, AND A TEMPORARY 'PUNCH' FILE
**
**
**
OUTPUT(.FINOUT,'FINOUT','')
OUTPUT(.OUTPUT,'TEMOUT;')
OUTPUT(.PUNCH,'TEMPCH;')
**
**
**
INITIALIZE KEYWORDS
**
**
**
&ANCHOR = 0
&DUMP = 1
&TRIM = 0
**
**
**
DEFINE FUNCTIONS
**
**
**
DEFINE(.SCANCOMDEF(.COMDEF),I')
DEFINE(.SCANIGNCOLDEF(.IGNCOLDEF),I')
DEFINE(.SCANLANLABDEF(.LANLABDEF),SPEC,I')
DEFINE(.SCANLITDEF(.LITDEF),I')
DEFINE(.SCANFOSDEF(.EOSDEF),I')
DEFINE(.SCANCONTDEF(.CUNTDEF),I')
DEFINE(.GETCHAR(.CHAR,COL),I')
DEFINE(.SCAN(.),I',BEGSCAN,I')
DEFINE(.INSERTC(.MICCMSTR),I')
DEFINE(.PARSE(.),I',STARTP,I')
DEFINE(.SCANDEFS(.CARD),SELECT,STMTTYPE,UNITYPE,I,REST,I')
DEFINE(.SCANOPERDEF(.TYPEDEF,OPERDEF),I')
DEFINE(.SCANNULDEF(.NULLDEF),I,J')
DEFINE(.SCANGOTDEF(.GOTDEF),I,J')
DEFINE(.SCANIFDEF(.IFDEF),I,J')
DEFINE(.SCANCLODEF(.DCLDEF),I,J')
DEFINE(.SCANMINUSDEF(.MINUSDEF),I,J')
DEFINE(.SCANPLUSDEF(.PLUSDEF),I,J')
DEFINE(.SCANABELDEF(.LABELDEF),I')
DEFINE(.SCANVARDEF(.VARDEF),I')
DEFINE(.MACROSCAN(.MACCALL,I',MACREPSTR,NAME,MACDEF,PLIST,I')
DEFINE(.CRMINUS(.LAB,VAR,VAL1,VAL2),STMTLEN,I')
DEFINE(.CRPLUS(.LAB,VAR,VAL1,VAL2),STMTLEN,I')
DEFINE(.CRNULL(.LAB),STMTLEN,I')
DEFINE(.CRGOTO(.LAB,DES),STMTLEN,I')
DEFINE(.CFIF(.LAH,VAL1,RL,VAL2,DES),STMTLEN,I')
DEFINE(.CROCL(.VARNAME),STMTLEN,I')
DEFINE(.UNILABEL(.I'),I')
DEFINE(.UNIVAR(.I'),I')
DEFINE(.POP(.SEM(NDARGS,A1,A2,A3,A4,A5,A6),I')
DEFINE(.PUSH(.SEM(NDARGS,A1,A2,A3,A4,A5,A6),I')
**
**
**
PRE000010
PRE000020
PRE000030
PRE000040
PRE000050
PRE000060
PRE000070
PRE000080
PRE000090
PRE000100
PRE000110
PRE000120
PRE000130
PRE000140
PRE000150
PRE000160
PRE000170
PRE000180
PRE000190
PRE000200
PRE000210
PRE000220
PRE000230
PRE000240
PRE000250
PRE000260
PRE000270
PRE000280
PRE000290
PRE000300
PRE000310
PRE000320
PRE000330
PRE000340
PRE000350
PRE000360
PRE000370
PRE000380
PRE000390
PRE000400
PRE000410
PRE000420
PRE000430
PRE000440
PRE000450
PRE000460
PRE000470
PRE000480

```

```

DEFINE('INSERTISMT(STR,INSRT)TYPE,OUTSTR,CHARKTR,I')
DEFINE('INSERT(INSRT,STR)')
DEFINE('ROUTO()')
DEFINE('ROUTO1()SEL MAC,STMTLEN,I,PARMLIST,PARM,MACSTR')
DEFINE('ROUT1()VAR1,VAR2')
DEFINE('ROUT2()VARO,VAR1,IND')
DEFINE('ROUT3()IND,BEGIN,END,INC,VAR1,VAR2')
DEFINE('ROUT4()')
DEFINE('ROUT5()BOOL,VAR1')
DEFINE('ROUT6()BOOL,VAR1')
DEFINE('ROUT7()BOOL,VAR1,VAR2')
DEFINE('ROUT8()VAR1')
DEFINE('ROUT9()VARO,VAR1,IND')
DEFINE('ROUT10()VAR1')
DEFINE('ROUT11()VAR1,VAR2,QUITLAB,QUITFLAG')
DEFINE('ROUT12()BOOL,VAR1,QUITLAB,QUITFLAG')
DEFINE('ROUT13()VAR1,VAR2,IND,INC,QUITLAB,QUITFLAG,VAR3,VAR4')
DEFINE('ROUT14()')
DEFINE('ROUT15()VARO,VAR1,VAR3,IND')

DECLARE ARRAYS

CONTPAT = ARRAY('1:2')
EQSPAT = ARRAY('0:1')
LITPAT = ARRAY('1:3')
LANLABPAT = ARRAY('1:9')
IGNPAT = ARRAY('0:2')
COMPAT = ARRAY('1:6')
NULLPAT = ARRAY('0:3')
GOTOPAT = ARRAY('0:5')
IFPAT = ARRAY('0:11')
MINUSPAT = ARRAY('0:9')
PLUSPAT = ARRAY('0:9')
DCLPAT = ARRAY('0:3')
LABELPAT = ARRAY('1:3')
VARPAT = ARRAY('1:3')
MACROTAB = TABLE()
SEMSTACK = ARRAY('1:10')
VARLIST = ARRAY('1:20')
MPAT = ARRAY('0:20')
MPARM = ARRAY('0:6')
CPARM = ARRAY('1:6')

INITIALIZE STRINGS AND PATTERNS

MACNAMEPAT = DUPL(' ',81)
LABDELS = ' | : : | : { ' ' . '
KEYCASE = 'CASE'

```

```

PRE000490
PRE000500
PRE000510
PRE000520
PRE000530
PRE000540
PRE000550
PRE000560
PRE000570
PRE000580
PRE000590
PRE000600
PRE000610
PRE000620
PRE000630
PRE000640
PRE000650
PRE000660
PRE000670
PRE000680
PRE000690
PRE000700
PRE000710
PRE000720
PRE000730
PRE000740
PRE000750
PRE000760
PRE000770
PRE000780
PRE000790
PRE000800
PRE000810
PRE000820
PRE000830
PRE000840
PRE000850
PRE000860
PRE000870
PRE000880
PRE000890
PRE000900
PRE000910
PRE000920
PRE000930
PRE000940
PRE000950
PRE000960

```

```

KEYDDO = 'DO'
KEYDDOGROUP = 'DOGROUP'
KEYIF = 'IF'
KEYREPEAT = 'REPEAT'
KEYWHILE = 'WHILE'
KEYELSE = 'ELSE'
KEYQUITLOOP = 'QUITLOOP'
KEYENDCASE = 'ENDCASE'
KEYENDIF = 'ENDIF'
KEYENDWHILE = 'ENDWHILE'
KEYENDREPEAT = 'ENDREPEAT'
KEYENDDO = 'ENDDO'
KEYENDDOGROUP = 'ENDDOGROUP'
KEYTHEN = 'THEN'
KEYUNTIL = 'UNTIL'
KEYWORDS = 'KEYCASE | KEYDO | KEYDOGROUP | KEYELSE |
KEYIF | KEYREPEAT | KEYWHILE | KEYELSE |
KEYQUITLOOP | KEYENDCASE | KEYENDIF |
KEYENDWHILE | KEYENDREPEAT | KEYENDDO |
KEYENDDOGROUP | KEYUNTIL'

KEYWORD2 = KEYTHEN | KEYUNTIL
NULL =
CURSIMT =
WORKING =
TOKEN =
NEXTCARD = 'FORTRAN' | 'PL/I' | 'SNOBOL'

LANGPAT = 'FORTRAN' | 'PL/I' | 'SNOBOL'

INITIALIZE SCALARS AND STRINGS TO BE USED AS SCALARS

CCLNUM = 80
VARNUM = '1'
NXTVAR = '1'
LABELNUM = '1'
TOP = 0
ENDFILE = 0
STARTCUL = 0
STINGSEP = 0
GETEND = 0
COMMENTONLY = 0
STMTFOS = 1
FOSNEXT = 0
MACROSPRES = 0
INASTRUCT = 0
ND.OF.CASES = 0
KTR0 = 0

```

```

PRE00970
PRE00980
PRE00990
PRE01000
PRE01010
PRE01020
PRE01030
PRE01040
PRE01050
PRE01060
PRE01070
PRE01080
PRE01090
PRE01100
PRE01110
PRE01120
PRE01130
PRE01140
PRE01150
PRE01160
PRE01170
PRE01180
PRE01190
PRE01200
PRE01210
PRE01220
PRE01230
PRE01240
PRE01250
PRE01260
PRE01270
PRE01280
PRE01290
PRE01300
PRE01310
PRE01320
PRE01330
PRE01340
PRE01350
PRE01360
PRE01370
PRE01380
PRE01390
PRE01400
PRE01410
PRE01420
PRE01421
PRE01422

```

```

**      DETERMINE LANGUAGE BEING PROCESSED --
**      FORTRAN, PL/I, OR SNOBOL
**
**      READ THE LANGUAGE TYPE; BRANCH IF NOT VALID
**
**      LANGUAGE = TRIM(INPUT)
**      LANGUAGE LANGPAT
**
**      ***(LANGERR)
**
**      DECODE DEFINITIONS
**
**      READ A CARD; DETERMINE IF IT IS AN OVERRIDE CARD; THE FIRST
**      DEFINITION, A MACRO CALLING FORM, OR THE FIRST CARD OF THE
**      INPUT PROGRAM
**
**      INLOOP   CARD = INPUT
**
**      DECODE AN OVERRIDE CARD; BRANCH TO READ THE NEXT CARD
**
**      CARD POS(0) 'OVERRIDE:' = 'OVERRIDE:'
**      CARD BREAK(1,1) : OVERKEY 'X' = 'X'
**      FINOUT = OVERKEY : KEYWORD OVERRIDDEN.
**      $((KEY, OVERKEY) = DUPL('X',81)
**      LASTOVERKEY CARD BREAK(1,1) : OVERKEY =
**      CAPD REM OVERKEY = DUPL('X',81)
**      $(KEY, OVERKEY) = KEYWORD OVERRIDDEN.
**      FINOUT = '
**      FINOUT = '
**      FINOUT = '
**      NOTOVERRIDE CARD POS(0) 'MACRO CALL>::='
**
**      SET THE MACROS PRESENT FLAG TO 1; CALL THE MACRO SCANNING ROUTINE;
**      BRANCH TO READ THE NEXT CARD
**
**      MACROSPRES = 1
**      MACROSCAN(TRIM(CARD))
**      CARD POS(0) '<'
**
**      NOTMACRO
**
**      CALL THE OFFINITION SCANNING ROUTINE; SET THE SEMANTIC
**      DELIMITER STRING; BRANCH TO READ ANOTHER CARD
**
**      SCANDFFS(CARD)

```

```

PRE01423
PRE01424
PRE01425
PRE01426
PRE01427
PRE01430
PRE01440
PRE01450
PRE01460
PRE01470
PRE01471
PRE01472
PRE01473
PRE01474
PRE01475
PRE01476
PRE01480
PRE01490
PRE01500
PRE01510
PRE01520
PRE01530
PRE01540
PRE01550
PRE01560
PRE01570
PRE01580
PRE01590
PRE01600
PRE01610
PRE01620
PRE01630
PRE01640
PRE01650
PRE01660
PRE01670
PRE01680
PRE01690
PRE01700
PRE01710
PRE01720
PRE01730
PRE01740
PRE01750
PRE01760
PRE01770
PRE01780
PRE01790

```



```

NOTDEFNS
**      SEMDEL = ' '
**      FOSPAT<0> 'COL'
**      SEMDEL = ' '
**      NEXTCARD = CARD
**      SUBSTR(EOSPAT<1>,1,1)
**      :S(INLOOP)
**      :S(INLOOP)
**
**      P. NUMBER = 1
**      PAKSE()
**      WORKING = POS(0) ' PROG' RPOS(0)
**      FINOUT = 'FIRST' PARSE OF PROGRAM WAS SUCCESSFUL.
**      FINOUT = ' '
**      FINOUT = ' '
**
**      REWIND THE TEMPORARY PUNCH FILE; ASSOCIATE AN INPUT VARIABLE
**      WITH THAT FILE AND OUTPUT VARIABLES WITH A SECOND TEMPORARY
**      'PRINT' FILE AND A SECOND TEMPORARY 'PUNCH' FILE
**
**      REWIND('TEMPCH')
**      INPUT('INPUT','TEMPCH')
**      OUTPUT('OUTPUT','TEMPCH2')
**      OUTPUT('PUNCH','TEMPCH2')
**
**      BRANCH IF NO SECOND PARSE IS REQUIRED (NO MACROS WERE PROVIDED)
**      EQ(MACROSPRES,0)
**      :S(NOSECPARSE)
**
**      REINITIALIZE STRINGS AND SCALARS FOR SECOND PARSE
**
**      CURSTMT =
**      WORKING =
**      TOKEN =
**      CCLNUM = 80
**      TOP = 0
**      CUSNEXT = 0
**      ENDFILE = 0
**      GETEND = 0
**      COMMENTONLY = 0
**      STIMEOS = 1
**      STARTCOL = 0
**      STRINGSEP = 0
**      INASTRUCT = 0
**      NO.OF.CASES = 0
**      KTRU = 0

```

```

PRE01800
PRE01810
PRE01820
PRE01830
PRE01840
PRE01841
PRE01842
PRE01850
PRE01851
PRE01852
PRE01860
PRE01870
PRE01880
PRE01890
PRE01900
PRE01910
PRE01920
PRE01930
PRE01940
PRE01950
PRE01960
PRE01970
PRE01980
PRE01990
PRE02000
PRE02010
PRE02020
PRE02030
PRE02040
PRE02050
PRE02060
PRE02070
PRE02080
PRE02090
PRE02100
PRE02110
PRE02120
PRE02130
PRE02140
PRE02150
PRE02160
PRE02170
PRE02180
PRE02190
PRE02200
PRE02210
PRE02220
PRE02230

```

[illegible]


```

** CKSNOBOL
**      REWIND('TEMPCH')
**      REWIND('TEMOUT')
**      INPUT('INCARD','TEMPCH')
**      INPUT('INLINE','TEMOUT')
**      OUTPUT('PUNCH','FINPUNCH')
**      OUTPUT('OUTPUT','FINOUT',' ' )
**
**      BRANCH IF THE LANGUAGE BEING PROCESSED IS NOT SNOBOL
**
**      LANGUAGE 'SNOBOL'
**
**      POST PROCESSING FOR SNOBOL:
**
**      SCAN THE FILES FOR CONDITIONALS USING ~( FOLLOWED BY
**      BLANKS AND ;( -- THESE WERE ADDED BY THE PREPROCESSOR;
**      DELETE THE INTERVENING BLANKS AND THE COMMA TO MAKE THE
**      STATEMENT A LEGAL SNOBOL STATEMENT
**
**      SNOBOLTEST CARD = INCARD
**      LINE = INLINE
**      CARD ' ~( SPAN( ' ' ) ; ; ' = ' ~( ( '
**      LINE ' ~( SPAN( ' ' ) ; ; ' = ' ~( ( '
**      PUNCH = CARD
**      OUTPUT = LINE
**
**      BRANCH IF THE LANGUAGE BEING PROCESSED IS NOT FORTRAN
**
**      NOTSNOBOL LANGUAGE 'FORTRAN'
**
**      POST PROCESSING FOR FORTRAN:
**
**      SCAN THE FILES FOR SEQUENCES OF A COMMENT, A CONTINUE, AND A
**      FORMAT STATEMENT; RETURN THE STATEMENT NUMBER FROM THE
**      CONTINUE (INDICATION IS IN THE COMMENT THAT THE STATEMENT
**      NUMBER WAS TAKEN FROM THE FORMAT AND PLACED ON THE CONTINUE)
**      TO THE FORMAT STATEMENT AND DELETE THE CONTINUE STATEMENT
**
**      FORTTEST CARD1 = INCARD
**      LINE1 = INLINE

```

```

F0RTEFST2      CARD1 POS(0) 'C' BREAK(' ') . STMT.NUMB      :F(N0F0RTECOM)
CARD2 = INCARD
LINE2 = INLINE
CARD2 POS(6) 'CONTINUE'
CARD3 = INCARD
LINE3 = INLINE
TESTCONTINUE  CARD3 POS(0) . CONTINUE .
CARD3 = INCARD
LINE3 = INLINE
THISFORM      CARD3 POS(0) . ' ARB REP 'FORMAT' =
+             DUPL(' ',5 - SIZE(STMT.NUMB)) STMT.NUMB . ' REP
+             'FORMAT' . ' ' = DUPL(' ',5 - SIZE(STMT.NUMB))
+             STMT.NUMB . '
PUNCH = CARD3
OUTPUT = LINE3
N0F0RTECOM     PUNCH = CARD1
OUTPUT = LINE1
N0F0RTECONT    PUNCH = CARD1
CARD1 = CARD2
OUTPUT = LINE1
LINE1 = LINE2
PUNCH = CARD1
PUNCH = CARD2
CARD1 = CARD3
OUTPUT = LINE1
OUTPUT = LINE2
LINE1 = LINE3
*             BRANCH IF THE LANGUAGE BEING PROCESSED IS NOT PL/I
*             :F(N0TPLI)
N0T0RTRAN LANGUAG 'PL/I'
*             ****
*             POST PROCESSING FOR PL/I:
*             SCAN THE BEGINNING OF THE FILES FOR EITHER ONE OR TWO
*             COMMENTS FOLLOWED BY THE SAME NUMBER OF NULL STATEMENTS AND
*             A PROCEDURE STATEMENT; RETURN THE LABEL WHICH WAS ORIGINALLY
*             TAKEN OFF THE PROCEDURE STATEMENT TO THE PROCEDURE STATEMENT
*             AND DELETE THE NULL STATEMENTS; COPY THE REMAINDER
*             OF THE FILES
*             CARD1 = INCARD
*             LINE1 = INLINE

```

PRE03480
 PRE03490
 PRE03500
 PRE03510
 PRE03520
 PRE03530
 PRE03540
 PRE03550
 PRE03560
 PRE03570
 PRE03580
 PRE03590
 PRE03600
 PRE03610
 PRE03620
 PRE03630
 PRE03640
 PRE03650
 PRE03660
 PRE03670
 PRE03680
 PRE03690
 PRE03700
 PRE03710
 PRE03720
 PRE03730
 PRE03740
 PRE03750
 PRE03760
 PRE03770
 PRE03780
 PRE03790
 PRE03800
 PRE03801
 PRE03802
 PRE03803
 PRE03804
 PRE03810
 PRE03820
 PRE03830
 PRE03840
 PRE03850
 PRE03860
 PRE03870
 PRE03871
 PRE03872
 PRE03880
 PRE03890

```

CARD1 POS(0) SPAN(' ') /* BREAK(' ') . STMT.NUMB
CARD2 = INCARD
LINE2 = INLINE
CARD2 POS(0) SPAN(' ') /*
      :F(THISPROC)
CARD2 = INCARD
LINE2 = INLINE
CARD2 = INCARD
LINE2 = INLINE
CARD2 = INCARD
LINE2 = INLINE
CARD2 POS(0) SPAN(' ') REP 'PROC' =
      DUPL(' ', SIZE(REP) - SIZE(STMT.NUMB))
      STMT.NUMB 'PROC'
      LINE2 POS(0) SPAN(' ') REP 'PROC' = DUPL(' ', SIZE(REP) -
      SIZE(STMT.NUMB)) STMT.NUMB 'PROC'
      PUNCH = CARD1
      PUNCH = CARD2
      OUTPUT = LINE1
      OUTPUT = LINE2
      OUTPUT = INCARD
      COPYRESTPLI OUTPUT = INLINE
      :F(COMPLETE)
      :F(COPYRESTPLI)
      THF LANGUAGE WAS NOT FORTRAN OR SNOBOL OR PL/I, SO SIMPLY
      COPY THE FILES
      NOTPLI PUNCH = INCARD
      OUTPUT = INLINE
      :F(COMPLETE)
      :F(NOTFORTRAN)
      OUTPUT A SUCCESSFUL COMPLETION MESSAGE AND TERMINATE
      COMPLETE FINOUT = 'PREPROCESSING COMPLETED SUCCESSFULLY.'
      :F(END)
      SECTION TO PRODUCE OUTPUT AFTER ERROR MESSAGES HAVE BEEN
      PRINTED FOR ERRORS DURING PROCESSING OF THE INPUT PROGRAM
      COPY OUTPUT PRODUCED UP TO THE ERROR WHICH HAS BEEN ENCOUNTERED;
      THE FILE OUTPUT DEPENDS ON THE PARSE NUMBER DURING WHICH
      THE ERROR OCCURRED

```

```

PRE03900
PRE03910
PRE03920
PRE03930
PRE03940
PRE03950
PRE03960
PRE03970
PRE03980
PRE03990
PRE04000
PRE04010
PRE04020
PRE04030
PRE04040
PRE04050
PRE04060
PRE04070
PRE04080
PRE04090
PRE04100
PRE04110
PRE04112
PRE04120
PRE04130
PRE04140
PRE04141
PRE04142
PRE04150
PRE04160
PRE04170
PRE04180
PRE04190
PRE04200
PRE04210
PRE04220
PRE04221
PRE04222
PRE04230
PRE04240
PRE04250
PRE04260
PRE04270
PRE04280
PRE04290
PRE04291
PRE04292

```



```

GEQUDEFERR      FINOUT = 'ABNORMAL TERMINATION'                : (END)
+               FINOUT = 'AN ERROR WAS ENCOUNTERED IN SCANNING THE '
+               FINOUT = 'GREATER THAN OR EQUAL OPERATOR'
+               FINOUT = 'DEFINITION = , OPERDEF
+               FINOUT = 'ABNORMAL TERMINATION'                : (END)
+               FINOUT = 'AN ERROR WAS ENCOUNTERED IN SCANNING THE '
+               FINOUT = 'NOT OPERATOR'
+               FINOUT = 'DEFINITION = , OPERDEF
+               FINOUT = 'ABNORMAL TERMINATION'                : (END)
+               FINOUT = 'AN ERROR WAS ENCOUNTERED IN SCANNING THE '
+               FINOUT = 'NOT EQUAL OPERATOR'
+               FINOUT = 'DEFINITION = , OPERDEF
+               FINOUT = 'A COLUMN SPECIFICATION DID NOT CONTAIN ONLY '
+               FINOUT = 'DIGITS OR EXTRANEOUS INFORMATION WAS FOUND IN '
+               FINOUT = 'SCANNING THE END OF STATEMENT DEFINITION'
+               FINOUT = 'ABNORMAL TERMINATION'                : (END)
+               FINOUT = 'A COLUMN SPECIFICATION DID NOT CONTAIN '
+               FINOUT = 'ONLY DIGITS IN THE IGNORE COLUMNS DEFINITION'
+               FINOUT = 'ABNORMAL TERMINATION'                : (END)
+               FINOUT = 'AN ERROR WAS ENCOUNTERED IN SCANNING THE '
+               FINOUT = 'COMMENT DEFINITION'
+               FINOUT = 'ABNORMAL TERMINATION'                : (END)
+               FINOUT = 'THE <MACRO DEFINITION> CARD FOR MACRO : NAME
+               FINOUT = 'IS MISSING OR INCORRECT'
+               FINOUT = 'ABNORMAL TERMINATION'                : (END)
+               FINOUT = 'AN ERROR WAS ENCOUNTERED IN A MACRO CALLING FORM '
+               FINOUT = 'DEFINITION'
+               FINOUT = 'MACRO NAME = , NAME
+               FINOUT = 'ABNORMAL TERMINATION'                : (END)
+               FINOUT = 'MORE THAN 6 PARAMETERS WERE SPECIFIED FOR MACRO '
+               FINOUT = 'NAME
+               FINOUT = 'ABNORMAL TERMINATION'                : (END)
+
+               EXITS FOR ERRORS WHICH OCCURRED IN THE PARSER
+
+               FINOUT = 'SCAN RETURNED TO THE PARSER WITH FAILURE '
+               FINOUT = 'WHEN A RETURNED VALUE WAS REQUIRED'
+               FINOUT = 'CURSTMT = , CURSTMT
+               FINOUT = 'WORKING = , WORKING
+               FINOUT = 'PARSER IS UNABLE TO RECOGNIZE THE SEQUENCE '
+               FINOUT = 'OF TOKENS IN THE WORKING STRING'
+               FINOUT = 'CURSTMT = , CURSTMT
+               FINOUT = 'WORKING = , WORKING
+               FINOUT = 'FAILURE TO PARSE STRUC TO STRLIST REDUCTION'
+
PRE05180
PRE05190
PRE05200
PRE05210
PRE05220
PRE05230
PRE05240
PRE05250
PRE05260
PRE05270
PRE05280
PRE05290
PRE05300
PRE05310
PRE05320
PRE05330
PRE05340
PRE05350
PRE05360
PRE05370
PRE05380
PRE05390
PRE05400
PRE05410
PRE05420
PRE05430
PRE05440
PRE05450
PRE05460
PRE05470
PRE05480
PRE05490
PRE05500
PRE05501
PRE05502
PRE05510
PRE05520
PRE05521
PRE05522
PRE05530
PRE05540
PRE05550
PRE05560
PRE05570
PRE05580
PRE05590
PRE05600
PRE05610

```

[illegible]

[illegible]

```

EDDERROR FINOUT = 'ERROR IN SEMANTIC PROCESSING OF ENDDG.'
FINOUT = 'INSTR = INSTR : (COPYTEMP OUT)'
EDOGKPEROR FINOUT = 'ERROR IN SEMANTIC PROCESSING OF ENDDOGRUP.'
FINOUT = 'INSTR = INSTR : (COPYTEMP OUT)'
MMACCALLPERR FINOUT = 'MORE THAN SIX ARGUMENTS WERE SUPPLIED IN A CALL'
+ TO THE MACRO , SELMAC : (COPYTEMP OUT)
RROUTOIERR FINOUT = 'THE STMLEN VARIABLE EQUALS 0 IN THE MACRO ,
+ DEFINITION FOR , SELMAC : (COPYTEMP OUT)
MMACARGERR FINOUT = 'A PARAMETER FOUND IN THE MACRO DEFINITION ARRAY ,
+ FOR , SELMAC , IS NOT IN THE PARAMETER ARRAY FOR
+ THAT MACRO. : (COPYTEMP OUT)
MACRLKERR FINOUT = 'NOT ENOUGH BLANKS LEFT AFTER A PARAMETER ,
+ NAME IN THE DEFINITION FOR THE MACRO , SELMAC : (COPYTEMP OUT)
TOOFULLSTK FINOUT = 'ATTEMPT TO PUSH ONTO A FULL SEMANTIC STACK.
+ : (COPYTEMP OUT)
EMPTYSTK FINOUT = 'ATTEMPT TO POP FROM AN EMPTY SEMANTIC STACK.
+ : (COPYTEMP OUT)
*****EXITIS FOR ERRORS WHICH OCCURRED IN THE CREATION ROUTINES*****
CRMINEKR FINOUT = 'AN ERROR WAS ENCOUNTERED DURING AN ATTEMPT TO ,
+ CREATE AN ASSIGN MINUS STATEMENT. : (COPYTEMP OUT)
CRCDCLERR FINOUT = 'AN ERROR WAS ENCOUNTERED DURING AN ATTEMPT TO ,
+ CREATE A DECLARATION STATEMENT. : (COPYTEMP OUT)
CRCIFERR FINOUT = 'AN ERROR WAS ENCOUNTERED DURING AN ATTEMPT TO ,
+ CREATE AN IF THEN GO TO STATEMENT. : (COPYTEMP OUT)
CRNULLERR FINOUT = 'AN ERROR WAS ENCOUNTERED DURING AN ATTEMPT TO ,
+ CREATE A NULL STATEMENT. : (COPYTEMP OUT)
CRPLERR FINOUT = 'AN ERROR WAS ENCOUNTERED DURING AN ATTEMPT TO ,
+ CREATE AN ASSIGN PLUS STATEMENT. : (COPYTEMP OUT)
CRGOTOERR FINOUT = 'AN ERROR WAS ENCOUNTERED DURING AN ATTEMPT TO ,
+ CREATE A GO TO STATEMENT. : (COPYTEMP OUT)
DCBLKERR FINOUT = 'NOT ENOUGH BLANKS LEFT AFTER A PARAMETER NAME ,
+ IN THE DECLARATION PROTOTYPE. : (COPYTEMP OUT)
GOTOBKERR FINOUT = 'VARNAME = , VARNAME : (COPYTEMP OUT)
FINOUT = 'NOT ENOUGH BLANKS LEFT AFTER A PARAMETER NAME ,
+ IN THE GO TO STATEMENT PROTOTYPE. : (COPYTEMP OUT)
PLUS9LKERR FINOUT = 'CURSTMT = , CURSTMT : (COPYTEMP OUT)
FINOUT = 'NOT ENOUGH BLANKS LEFT AFTER A PARAMETER NAME ,
+ IN THE ASSIGN PLUS PROTOTYPE. : (COPYTEMP OUT)
MINBKERR FINOUT = 'CURSTMT = , CURSTMT : (COPYTEMP OUT)
FINOUT = 'NOT ENOUGH BLANKS LEFT AFTER A PARAMETER NAME ,
+ : (COPYTEMP OUT)
PREU6500 PREU6510 PREU6520 PREU6530 PREU6540 PREU6550 PREU6560 PREU6570 PREU6580 PREU6590 PREU6600 PREU6610 PREU6620 PREU6630 PREU6640 PREU6650 PREU6660 PREU6670 PREU6680 PREU6681 PKEU6682 PREU6690 PREU6700 PREU6701 PREU6702 PREU6710 PREU6720 PREU6730 PREU6740 PREU6750 PREU6760 PREU6770 PREU6780 PREU6790 PREU6800 PREU6810 PREU6820 PREU6830 PREU6840 PREU6850 PREU6860 PREU6870 PREU6880 PREU6890 PREU6900 PREU6910 PREU6920 PREU6930

```

```

+      IN THE ASSIGN MINUS PROTOTYPE.
FINOUT = CURSMT = CURSMT : (COPYTEMPPOUT)
NULLRLKRR FINOUT = NOT ENOUGH BLANKS LEFT AFTER A PARAMETER NAME ;
+      IN THE NULL STATEMENT PROTOTYPE.
FINOUT = CURSMT = CURSMT : (COPYTEMPPOUT)
IFBLKERR FINOUT = NOT ENOUGH BLANKS LEFT AFTER A PARAMETER NAME ;
+      IN THE IF THEN GO TO STATEMENT PROTOTYPE.
FINOUT = CURSMT = CURSMT : (COPYTEMPPOUT)
TOMMANYLARELS FINOUT = ATTEMPT TO CREATE MORE UNIQUE LABELS THAN ;
+      THE DEFINITION ALLOWS.
TOMMANYVARS FINOUT = ATTEMPT TO CREATE MORE UNIQUE IDENTIFIERS THAN ;
+      THE DEFINITION ALLOWS.
BEGSCAN OLDSTMTEUS = STMTEOS
*****
ROUTINE TO SCAN INPUT FOR TOKENS RECOGNIZABLE BY THE PARSER
NO PARAMETERS
A CARD IMAGE MAY BE PRESENT IN NEXTCARD OR NOT WHEN THIS ROUTINE
IS INITIALLY CALLED
*****
IF AN EOS SHOULD BE SIGNALLED IMMEDIATELY, BRANCH TO THAT LABEL
SAVE FLAG INDICATING WHETHER OR NOT AN EOS WAS FOUND IN THE LAST
STATEMENT; SET EOS FLAG FOR THIS STATEMENT TO O
STMTEOS = O
EQ(ECSNEXT,1)
:S(FDSIMMED)
GET THE NEXT USFABLE CHARACTER AND ITS COLUMN NUMBER; IF GETCHAR
RETURNS WITH FAILURE, BRANCH
BEGSCAN GETCHAR('INCHAR','CHARCOL')
:F(ENDOFFINPUT)
IF THIS IS THE FIRST CHARACTER OF A TOKEN OR THE FIRST NON-BLANK
CHARACTER OF A TOKEN, SET THE STARTCOL VARIABLE
STARTCOL = EQ(SIZE(TOKEN),O) CHARCOL
*****
IF THIS IS THE FIRST USEABLE COLUMN ON A CARD, CHECK TO
SEE IF THE TOKEN IS CURRENTLY ALL BLANKS LEFT OVER FROM A
PREVIOUS CARD WHICH SHOULD BE DISCARDED
*****
GLOSTART NE(IGNPAT<I>,1)
:S(CKATCOLONE)
*****

```



```

      IGNPAT<0> 'ONE'
      EQ(CHARCOL,2)
      LOOKATSEQEND EQ(CHARCOL,IGNPAT<2> + 1)      :F(LEAVETOKEN)S(CKBLANKTOKEN)
      CKATCOLONE EQ(CHARCOL,1)                    :F(LEAVETOKEN)
      CKBLANKTOKEN TOKEN POS(0) SPAN(' ') RPOS(0)  :S(TESTULDEOS)
**
      IF THE TOKEN IS OF 0 LENGTH AND A COMMENT WAS NOT ALL THAT
      APPEARED ON THE LAST CARD, LEAVE THE TOKEN AND CURSTMT INTACT
      EQ(COMMENTONLY,1)
      INSERTCOM(CURSTMT)
**
      THE CURSTMT AND TOKEN ARE NULLIFIED IF AN EOS WAS FOUND BEFORE
      THIS CALL OF SCAN
**
      TESTULDEOS EQ(OLDSTMTIEOS,1)
      CURSTMT =
      TOKEN =
      :F(LEAVETOKEN)
**
      ADD THE CHARACTER TO THE TOKEN STRING
**
      LEAVETOKEN TOKEN = TOKEN INCHAR
      LABSTART = 0
      COMMENTONLY = 0
**
      LOOK FOR BEGINNING OF A COMMENT
**
      BRANCH TO APPROPRIATE SECTION DEPENDING ON OPENING
      FOR COMMENT SPECIFIED IN DEFINITION
      COMPAT<1> '10'
      :S(COMSYM)
      COLUMN NUMBER PART OR ALL OF COMMENT OPENING
      EQ(CHARCOL,COMPAT<4>)
      :F(NTCOM)
      COMMENT FOUND IF COLUMN ONLY STARTS COMMENT
      COMPAT<1> '11'
      :F(COMFOUND)S(COMCOLSYM)
      SYMBOL ONLY STARTS COMMENT
      CCMSYM COMPAT<3> POS(0) INCHAR
      :F(NTCOM)
      COMMENT FOUND IF ONLY ONE SYMBOL REQUIRED

```

```

PRE07390
PRE07400
PRE07410
PRE07420
PRE07430
PRE07440
PRE07450
PRE07460
PRE07470
PRE07480
PRE07490
PRE07500
PRE07510
PRE07520
PRE07530
PRE07540
PRE07550
PRE07560
PRE07570
PRE07580
PRE07590
PRE07600
PRE07610
PRE07620
PRE07621
PRE07622
PRE07623
PRE07624
PRE07625
PRE07626
PRE07630
PRE07640
PRE07650
PRE07660
PRE07670
PRE07680
PRE07690
PRE07700
PRE07710
PRE07720
PRE07730
PRE07740
PRE07750
PRE07760
PRE07770
PRE07780
PRE07790
PRE07800

```

```

*      EQ(SIZE(COMPAT<3>),1)                                :S(COMFOUND)
*      ERROR IF COMMENT STARTING CHARACTER STRING IS
*      LONGER THAN 2 CHARACTERS
*      EQ(SIZE(COMPAT<3>),2)                                :F(COMOPENERR)
*      SAVE THE CURRENT CHARACTER AND COLUMN
*      OLDCOL = CHARCOL
*      CLDCHAR = INCHAR
*      GET THE NEXT USEABLE CHARACTER AND ITS COLUMN
*      GETCHAR('INCHAR','CHARCOL')
*      BRANCH IF THE 2 CHARACTERS DO NOT OPEN A COMMENT;
*      OTHERWISE ADD THE NEW CHARACTER TO THE TOKEN STRING
*      COMPAT<3> OLDCHAR INCHAR                            :F(NOTCOM2CHARS)
*      TOKEN = TOKEN INCHAR                                :(COMFOUND)
*      COLUMN NUMBER AND SYMBOL START COMMENT
*      CCMCULSYM COMPAT<3> POS(0) INCHAR                    :F(NOTCOM)
*      COMMENT FOUND IF ONLY ONE SYMBOL REQUIRED
*      EQ(SIZE(COMPAT<3>),1)                                :S(COMFOUND)
*      ERROR IF COMMENT STARTING STRING LONGER THAN 2 CHARACTERS
*      FQ(SIZE(COMPAT<3>),2)                                :F(COMOPENERR)
*      SAVE THE CURRENT CHARACTER AND COLUMN
*      OLDCOL = CHARCOL
*      OLDCHAR = INCHAR
*      GET THE NEXT USEABLE CHARACTER AND ITS COLUMN
*      GETCHAR('INCHAR','CHARCOL')
*      BRANCH IF THE 2 CHARACTERS DO NOT OPEN A COMMENT;
*      OTHERWISE ADD THE NEW CHARACTER TO THE TOKEN STRING
*      COMPAT<3> OLDCHAR INCHAR                            :F(NOTCOM2CHARS)

```

```

PRE07810
PRE07820
PRE07830
PRE07840
PRE07850
PRE07860
PRE07870
PRE07880
PRE07890
PRE07900
PRE07910
PRE07920
PRE07930
PRE07940
PRE07950
PRE07960
PRE07970
PRE07980
PRE07990
PRE08000
PRE08010
PRE08020
PRE08030
PRE08040
PRE08050
PRE08060
PRE08070
PRE08080
PRE08090
PRE08100
PRE08110
PRE08120
PRE08130
PRE08140
PRE08150
PRE08160
PRE08170
PRE08180
PRE08190
PRE08200
PRE08210
PRE08220
PRE08230
PRE08240
PRE08250
PRE08260
PRE08270
PRE08280

```

```

**      TOKEN = TOKEN INCHAR
**      BACKUP BY RESTORING THE OLD CHARACTER AND OLD COLUMN NUMBER;
**      BRANCH TO SEARCH FOR SOMETHING ELSE
**      NOTCOM2CHARS INCHAR = OLDCHAR
**      CHARCOL = OLDCOL
**      COLNUM = OLDCOL
**      : (NOTCOM)
**      LOOK FOR THE END OF THE COMMENT
**      COMMENT HAS BEEN FOUND: RESET THE STMT EOS FLAG TO WHATEVER IF WAS
**      BEFORE THIS CALL OF SCAN: SET THE COMMENTONLY FLAG IF NOTHING
**      ELSE HAS BEEN ENCOUNTERED IN THIS STATEMENT; SEARCH FOR THE END
**      OF THE COMMENT
**      COMFOUND STMT EOS = OLDSTMT EOS
**      NE(SIZE(CURSTMT),0)
**      COMMENTONLY = 1
**      GET THE NEXT USEABLE CHARACTER AND ITS COLUMN NUMBER; ADD THE
**      CHARACTER TO THE TOKEN STRING
**      SKIPNOSTMT GETCHAR('INCHAR','CHARCOL')
**      TOKEN = TOKEN INCHAR
**      BRANCH TO APPROPRIATE SECTION DEPENDING ON CLOSING FOR COMMENT
**      SPECIFIED IN DEFINITION
**      COMFOUNDNOCHAR COMPAT<2> '10'
**      COLUMN NUMBER PART OR ALL OF COMMENT CLOSING
**      EQ(CHARCOL,COMPAT<6>)
**      COMMENT ENDS IF ONLY COLUMN ENDS COMMENT
**      COMPAT<2> '11'
**      COLUMN NUMBER AND SYMBOL END COMMENT
**      COMENDCOLSYM COMPAT<5> POS(0) INCHAR
**      COMMENT ENDS IF ONLY ONE SYMBOL REQUIRED

```



```

*      EQ(SIZE(COMPAT<5>),1)                :S(ENDCOM)
*      ERROR IF COMMENT CLOSING STRING IS LONGER THAN 2 CHARS.
*      FQ(SIZE(COMPAT<5>),2)                :F(COMENDERR)
*      SAVE THE CURRENT CHARACTER AND COLUMN
*      OLDCOL = CHARCOL
*      CLDCHAR = INCHAR
*
*      GET THE NEXT USEABLE CHARACTER AND ITS COLUMN NUMBER;
*      ADD THE CHARACTER TO THE TOKEN STRING
*      GETCHAR('INCHAR','CHARCOL')          :F(ENDDOFINPUT)
*      TOKEN = TOKEN INCHAR
*
*      COMMENT ENDS IF THE 2 CHARACTERS CLOSE A COMMENT;
*      OTHERWISE KEEP SEARCHING FOR THE END
*      COMPAT<5> OLDCHAR INCHAR :S(ENDCOM)F(COMFOUNDNOCHAR)
*
*      SYMBOL ONLY ENDS COMMENT, BRANCH IF THIS CHARACTER MAY
*      INDICATE THE END OF THE COMMENT
*      CUMENDSTR COMPAT<5> POS(0) INCHAR      :S(FIRSTCHAREND)
*
*      GET THE NEXT USEABLE CHARACTER AND ITS COLUMN NUMBER;
*      ADD THE CHARACTER TO THE TOKEN STRING; KEEP LOOKING
*      FOR THE CLOSING SYMBOL
*      GETCHAR('INCHAR','CHARCOL')          :F(ENDDOFINPUT)
*      TOKEN = TOKEN INCHAR :((COMENDSTR)
*
*      ONE SYMBOL HAS BEEN MATCHED; END OF COMMENT HAS BEEN FOUND
*      IF ONLY ONE CHARACTER IS NEEDED
*      FIRSTCHAREND EQ(SIZE(COMPAT<5>),1)      :S(ENDCOM)
*
*      ERROR IF COMMENT CLOSING STRING LONGER THAN 2 CHARS.
*      FQ(SIZE(COMPAT<5>),2)                :F(COMENDERR)
*
*      SAVE THE CURRENT CHARACTER AND COLUMN
*      OLDCOL = CHARCOL
*      CLDCHAR = INCHAR

```

```

PRE08710
PRE08720
PRE08730
PRE08740
PRE08750
PRE08760
PRE08770
PRE08780
PRE08790
PRE08800
PRE08810
PRE08820
PRE08830
PRE08840
PRE08850
PRE08860
PRE08870
PRE08880
PRE08890
PRE08900
PRE08910
PRE08920
PRE08930
PRE08940
PRE08950
PRE08960
PRE08970
PRE08980
PRE08990
PRE09000
PRE09010
PRE09020
PRE09030
PRE09040
PRE09050
PRE09060
PRE09070
PRE09080
PRE09090
PRE09100
PRE09110
PRE09120
PRE09130
PRE09140
PRE09150
PRE09160
PRE09170
PRE09180

```

```

* * * * *
GET THE NEXT USEABLE CHARACTER AND ITS COLUMN;
ADD THE CHARACTER TO THE TOKEN STRING
* * * * *
      GETCHAR('INCHAR','CHARCOL')
      TOKEN = TOKEN INCHAR
      :F(ENDOFINPUT)
* * * * *
END OF COMMENT FOUND IF THE 2 CHARACTERS CLOSE A COMMENT;
OTHERWISE KEEP SEARCHING
* * * * *
      COMPAT<5> OLDCHAR INCHAR
      :F(COMENDSTR)
* * * * *
END OF COMMENT FOUND
* * * * *
      CURSTMT = CURSTMT TOKEN
      TOKEN =
      : (BEGSCAN1)
* * * * *
ENDCGM
* * * * *
      NOT A COMMENT; LOCK FOR A LITERAL OPENING
      * * * * *
* * * * *
      NOTCOM
      LITPAT<1> POS(0) INCHAR
      :F(NOTLIT)
* * * * *
      LITERAL FOUND IF THE SINGLE SYMBOL STARTS A LITERAL
      :S(LITFOUND)
* * * * *
      EQ(SIZE(LITPAT<1>),1)
      ERROR IF LITERAL OPENING STRING LONGER THAN 2 CHARACTERS
      :F(LITOPENERR)
* * * * *
      SAVE THE CURRENT CHARACTER AND ITS COLUMN NUMBER
      CLDCOL = CHARCOL
      CLDCHAR = INCHAR
* * * * *
      GET THE NEXT USEABLE CHARACTER AND ITS COLUMN NUMBER
      GETCHAR('INCHAR','CHARCOL')
      :F(NOTLIT2CHARS)
* * * * *
      BRANCH IF THE 2 CHARACTERS DO NOT START A LITERAL;
      OTHERWISE ADD THE CHARACTER TO THE TOKEN STRING
      LITPAT<1> OLDCHAR INCHAR
      TOKEN = TOKEN INCHAR
      :F(NOTLIT2CHARS)
      : (LITFOUND)
* * * * *

```

```

*      BACKUP BY RESTORING OLD CHARACTER AND COLUMN NUMBER ;
*      BRANCH TO SEARCH FOR SOMETHING ELSE
*
* NOTLIT2CHARS INCHAR = OLDCHAR
*      CHARCOL = OLDCOL
*      CCLNUM = OLDCOL
*
*      : (NOTLIT)
*
*      LOOK FOR THE END OF THE LITERAL
*
*      LITERAL FOUND; GET THE NEXT USEABLE CHARACTER AND ITS
*      COLUMN NUMBER; ADD THE CHARACTER TO THE TOKEN STRING
*
*      LITFOUND  GETCHAR('INCHAR','CHARCOL')
*      TOKEN = TOKEN INCHAR
*
*      IF CHARACTER DOES NOT CLOSE A LITERAL, KEEP SEARCHING
*
*      LITPAT<2> POS(0) INCHAR
*
*      : (LITFOUND)
*
*      BRANCH TO LOOK FOR SPECIAL SEQUENCE IF THE CHARACTER MATCHES
*      THE CLOSING SINGLE CHARACTER
*
*      EQ(SIZE(LITPAT<2>),1)
*
*      ERROR IF LITERAL CLOSING STRING IS LONGER THAN 2 CHARACTERS
*
*      EQ(SIZE(LITPAT<2>),2)
*
*      SAVE THE CURRENT CHARACTER AND ITS COLUMN NUMBER
*
*      OLDCCCL = CHARCOL
*      CLDCHAR = INCHAR
*
*      GET THE NEXT USEABLE CHARACTER AND ITS COLUMN NUMBER;
*      ADD THE CHARACTER TO THE TOKEN STRING
*
*      GETCHAR('INCHAR','CHARCOL')
*      TOKEN = TOKEN INCHAR
*
*      LITERAL FOUND IF THE 2 CHARACTERS CLOSE A LITERAL;
*      OTHERWISE KEEP SEARCHING FOR END
*
*      LITPAT<2> OLDCHAR INCHAR      : F(LITFOUND)S(ENDLIT)

```

```

PRE09630
PRE09640
PRE09650
PRE09660
PRE09670
PRE09680
PRE09690
PRE09691
PRE09692
PRE09693
PRE09694
PRE09695
PRE09696
PRE09700
PRE09710
PRE09720
PRE09730
PRE09740
PRE09750
PRE09760
PRE09770
PRE09780
PRE09790
PRE09800
PRE09810
PRE09820
PRE09830
PRE09840
PRE09850
PRE09860
PRE09870
PRE09880
PRE09890
PRE09900
PRE09910
PRE09920
PRE09930
PRE09940
PRE09950
PRE09960
PRE09970
PRE09980
PRE09990
PRE10000
PRE10010
PRE10020
PRE10030
PRE10040

```

```

** ** ONE CHARACTER HAS BEEN MATCHED; LITERAL END FOUND IF NO SPECIAL
** ** SEQUENCE SPECIFIED IN DEFINITION
**
MATCHLITEND1 EQ(SIZE(LITPAT<3>),0)          *:S(ENDLIT)
**
SAVE THE CURRENT CHARACTER AND COLUMN NUMBER
**
OLDCOL = CHARCOL
OLDCHAR = INCHAR
**
GET THE NEXT USEABLE CHARACTER AND ITS COLUMN NUMBER
**
GETCHAR('INCHAR','CHARCOL')                *:F(ENDLIT1)
**
END OF LITERAL FOUND IF THIS IS NOT THE SPECIAL SEQUENCE
**
LITPAT<3> OLDCHAR INCHAR                     *:F(ENDLIT1)
**
ADD THE CHARACTER TO THE TOKEN STRING; SKIP THIS CHARACTER
AND KEEP SEARCHING FOR THE LITERAL END
**
TOKEN = TOKEN INCHAR                        *::(LITFOUND)
**
BACKUP BY RESTORING THE OLD CHARACTER AND ITS COLUMN NUMBER
**
INCHAR = OLDCHAR
CHARCOL = OLDCOL
CCOLNUM = OLDCOL
**
A LITERAL IS AN UNIDENTIFIED OBJECT FOR THE PARSER, SO BRANCH TO
THAT SECTION
**
ENDLIT                                       *::(UNIDENT)
**
NOT A LITERAL; SEARCH FOR AN END OF STATEMENT
**
ENDLIT                                     *:S(EDSCOL)
**
SYMBOL SIGNIFIES END OF STATEMENT
**
FCSPAT<1> POS(0) INCHAR                    *:F(NOTEDSSCAN)
**
END OF STATEMENT FOUND IF SINGLE SYMBOL ENDS STATEMENT
**

```

```

** **
** **      FQ(SIZE(EOSPAT<1>),1)                :S(EOSFOUND)
** **      ERROR IF END OF STATEMENT STRING IS LONGER THAN 2 CHARACTERS
** **      EQ(SIZE(EOSPAT<1>),2)                :S(EOSDEFNERR)
** **      SAVE THE CURRENT CHARACTER AND COLUMN NUMBER
** **      CLDCOL = CHARCOL
** **      OLDCHAR = INCHAR
** **
** **      GET THE NEXT USEABLE CHARACTER AND ITS COLUMN NUMBER
** **      GETCHAR(' INCHAR', 'CHARCOL')        :F(NOTEOS2CHARS)
** **      BRANCH IF THE 2 CHARACTERS DO NOT MARK END OF STATEMENT;
** **      OTHERWISE ADD THE CHARACTER TO THE TOKEN STRING
** **      EOSPAT<1> OLDCHAR INCHAR              :F(NOTEOS2CHARS)
** **      TOKEN = TOKEN INCHAR                  :S(EOSFOUND)
** **
** **      BACKUP BY RESTORING THE OLD CHARACTER AND COLUMN NUMBER;
** **      BRANCH TO SEARCH FOR SOMETHING ELSE
** **
NOTEOS2CHARS INCHAR = OLDCHAR
CHARCOL = OLDCOL
CCLNUM = OLDCCOL
** **
** **      :S(EOSFOUND)
** **
** **      END OF STATEMENT MARKED BY COLUMN
** **
** **      END OF STATEMENT FOUND IF COLUMN NUMBER IS THAT SPECIFIED
** **      IN DEFINITION; OTHERWISE GO SEARCH FOR SOMETHING ELSE
** **
EOSCOL EQ(COLNUM,EOSPAT<1>) :S(EOSCOLFOUND)F(NOTEOSSCAN)
** **
** **      END OF STATEMENT FOUND BY COLUMN NUMBER;
** **      READ IN ANOTHER CARD AND LOOK FOR THE CONTINUATION
** **      SPECIFICATION
** **
EOSCOLFOUND EQ(GETEND,1)
NEXTCARD = INPUT
ENDFILE = 1
** **
** **      :S(EOSFOUND)
** **      :S(READACARD)
** **      :S(EOSFOUND)
** **
** **      BRANCH TO SEARCH FOR SOMETHING ELSE IF IT IS FOUND THAT
** **      THE STATEMENT CONTINUES ON THE NEXT CARD
** **
READACARD SURSTR(NEXTCARD,CONTPAT<2>,SIZE(CONTPAT<1>))

```

```

PRE10490
PRE10500
PRE10510
PRE10520
PRE10530
PRE10540
PRE10550
PRE10560
PRE10570
PRE10580
PRE10590
PRE10600
PRE10610
PRE10620
PRE10630
PRE10640
PRE10650
PRE10660
PRE10670
PRE10680
PRE10690
PRE10700
PRE10710
PRE10720
PRE10730
PRE10740
PRE10750
PRE10760
PRE10770
PRE10780
PRE10790
PRE10800
PRE10810
PRE10820
PRE10830
PRE10840
PRE10850
PRE10860
PRE10870
PRE10880
PRE10890
PRE10900
PRE10910
PRE10920
PRE10930
PRE10940
PRE10950
PRE10960

```

```

+ * * * * * CONTPAT<1> :S(NOTLANLABSEP) PRE10970
+ * * * * * ADD THE TOKEN STRING TO THE CURRENT STMT STRING PRE10980
+ * * * * * PRE10990
+ * * * * * PRE11000
+ * * * * * PRE11010
+ * * * * * PRE11020
+ * * * * * PRE11030
+ * * * * * PRE11040
+ * * * * * PRE11050
+ * * * * * PRE11060
+ * * * * * PRE11070
+ * * * * * PRE11080
+ * * * * * PRE11090
+ * * * * * PRE11100
+ * * * * * PRE11110
+ * * * * * PRE11120
+ * * * * * PRE11130
+ * * * * * PRE11140
+ * * * * * PRE11150
+ * * * * * PRE11160
+ * * * * * PRE11170
+ * * * * * PRE11180
+ * * * * * PRE11190
+ * * * * * PRE11200
+ * * * * * PRE11201
+ * * * * * PRE11202
+ * * * * * PRE11203
+ * * * * * PRE11204
+ * * * * * PRE11205
+ * * * * * PRE11206
+ * * * * * PRE11210
+ * * * * * PRE11220
+ * * * * * PRE11230
+ * * * * * PRE11240
+ * * * * * PRE11250
+ * * * * * PRE11260
+ * * * * * PRE11270
+ * * * * * PRE11280
+ * * * * * PRE11290
+ * * * * * PRE11300
+ * * * * * PRE11310
+ * * * * * PRE11311
+ * * * * * PRE11312
+ * * * * * PRE11320
+ * * * * * PRE11330
+ * * * * * PRE11340
+ * * * * * PRE11350
+ * * * * * PRE11360

+ * * * * * EOSFOUND CURSTMT = CURSTMT TOKEN
+ * * * * * TEST FOR A KEYWORD ENDING DIRECTLY BEFORE THE END OF STATEMENT
+ * * * * * SYMBOL OR A KEYWORD ENDING IN THE END OF STATEMENT COLUMN
+ * * * * * OR IN THE COLUMN BEFORE A BLANK IN THE END OF STATEMENT COLUMN
+ * * * * *
+ * * * * * EOSPAT<0> 'COL'
+ * * * * * TESTSTR = SUBSTR(TOKEN,1,SIZE(TOKEN) - SIZE(EOSPAT<1>))
+ * * * * * :S(COLKEYWORD)
+ * * * * * :S(TESTKEY)
+ * * * * * COLKEYWORD TESTSTR = TOKEN
+ * * * * * TESTSTR KEYWORDS : KEYFOUND RPOS(0)
+ * * * * * TESTSTR KEYWORDS : KEYFOUND :S(KEYEOS)
+ * * * * * KEYEOS WORKING = WORKING : KEYFOUND :F(EOSIMMED)
+ * * * * * TOKEN =
+ * * * * *
+ * * * * * FOSNFTX IS SET SO AN EOS WILL BE SIGNALLED IMMEDIATELY UPON THE
+ * * * * * NEXT CALL OF SCAN; RETURN
+ * * * * *
+ * * * * * EOSNEXT = 1
+ * * * * *
+ * * * * * END OF STATEMENT HAS BEEN FOUND
+ * * * * *
+ * * * * * STMT EOS IS SET FOR THIS STATEMENT AND EOSNEXT IS SET TO 0
+ * * * * *
+ * * * * * EOSIMMED STMT EOS = 1
+ * * * * * EOSNEXT = 0
+ * * * * *
+ * * * * * ADD THE 'UNIDENTIFIED TO END OF STATEMENT' SYMBOL TO
+ * * * * * THE WORKING STRING FOR THE PARSER AND RETURN
+ * * * * *
+ * * * * * ADDUEOS WORKING = WORKING + UEOS
+ * * * * * TOKEN =
+ * * * * *
+ * * * * * NOT AN END OF STATEMENT; LOOK FOR A LANGUAGE LABEL IF AT
+ * * * * * BEGINNING OF STATEMENT; ANYTHING IN CURSTMT FOLLOWED
+ * * * * * BY A BLANK OR JUST BLANKS INDICATES A LABEL HAS BEEN PROCESSED
+ * * * * * ALREADY

```

```

NOTEOSSCAN CURSTMT SPAN(' ',) BREAK(' ',)
CURSTMT BREAK(' ',)
LANLABPAT<2> 'STR'
**
**
**
ERROR IF LANGLABEL DEFINITION DOES NOT SPECIFY EITHER
A STRING OR A COLUMN AS AN ENDING
**
**
**
LANLABPAT<2> 'COL'
**
**
**
LANGLABEL ENDED BY COLUMN
**
**
**
BRANCH IF A SEQUENCE OF COLUMN NUMBERS IS NOT SPECIFIED
AS A LANGLABEL ENDER
**
**
**
EQ(SIZE(LANLABPAT<7>),0)
**
**
**
:S(ONECOLLANLABEND)
**
**
**
MAY BE A LANGLABEL ENDER IF THE COLUMN NUMBER IS LESS THAN OR
EQUAL THE UPPER LIMIT OF THE SPECIFIED SEQUENCE
**
**
**
LE(COLUMN,LANLABPAT<7>)
**
**
**
:S(NOTLANLABSEP)
**
**
**
MAY BE A LANGLABEL ENDER IF THE COLUMN NUMBER IS GREATER THAN OR
EQUAL THE LOWER LIMIT OF THE SPECIFIED SEQUENCE; OTHERWISE IT
IS NOT A LANGLABEL ENDER
**
**
**
GE(COLUMN,LANLABPAT<6>)
**
**
**
:S(CKLANLABSEP)F(NOTLANLABSEP)
**
**
**
IF LANGLABELS END IN ONLY ONE COLUMN, AND THIS
COLUMN NUMBER IS THE RIGHT ONE, THIS MAY BE A LANGLABEL
ENDER
**
**
**
ONECOLLANLABEND EQ(COLUMN,LANLABPAT<6>)
**
**
**
:S(CKLANLABSEP)F(NOTLANLABSEP)
**
**
**
A SYMBOL STRING SEPARATES A LANGUAGE LABEL FROM A
STATEMENT; BRANCH IF THIS IS NOT THE RIGHT CHARACTER
TO BEGIN A LABEL SEPARATOR
**
**
**
STRLABSEP LANLABPAT<6> INCHAK
SUBSTR(TOKEN,1,SIZE(TOKEN) - 1) NOTANY(' ')
**
**
**
SPAN(' ',) NOTANY(' ')
**
**
**
:S(NOTLANLABSEP)
**
**
**
THIS IS A LABEL SEPARATOR IF ONLY ONE SYMBOL IS REQUIRED
**
**
**
STPINGSEP = 1
EQ(SIZE(LANLABPAT<6>),1)
**
**
**
:S(CKLANLABSEP)
**

```



```

* * * ERROR IF THE LABEL SEPARATOR STRING IS LONGER THAN 2
* * * CHARACTERS
* * *      FQ(SIZE(LANLABPAT<6>),2)
* * *      :F(LANLABDEFNERR)
* * *
* * * SAVE THE CURRENT CHARACTER AND COLUMN NUMBER
* * *      OLDCCOL = CHARCOL
* * *      OLDCHAR = INCHAR
* * *
* * * GET THE NEXT USEABLE CHARACTER AND ITS COLUMN NUMBER
* * *      GETCHAR('INCHAR','CHARCOL')
* * *      :F(ENDOFINPUT)
* * *
* * * BRANCH IF THE 2 CHARACTERS ARE NOT A LANGLABEL
* * * SEPARATOR; OTHERWISE ADD THE CHARACTER TO THE
* * * TOKEN STRING AND SET THE BACKUP FLAG TO 1 TO INDICATE THAT
* * * BACKUP MUST BE DONE LATER IF THE ITEM PRECEDING THE SEPARATOR
* * * IS NOT A LABEL
* * *
* * *      LANLABPAT<6> OLDCHAR INCHAR
* * *      TOKEN = 'TOKEN' INCHAR
* * *      BACKUP = 1
* * *      :F(NOTLANLAB2CHAR)
* * *
* * *      BACKUP BY RESTORING THE OLD CHARACTER AND COLUMN NUMBER
* * *
* * *      :F(NOTLANLAB2CHAR)
* * *      :F(NOTLANLAB2CHAR)
* * *
* * *      NOTLANLAB2CHAR INCHAR = OLDCHAR
* * *      CHARCOL = OLDCCOL
* * *      CCLNUM = OLDCCOL
* * *
* * *      :F(NOTLANLABSEP)
* * *
* * * A LANGUAGE LABEL ENDER HAS BEEN FOUND; ADD THE 'LABEL' ITEM TO
* * * THE WORKING STRING FOR THE PARSER
* * *
* * *      EXTRACT THE 'LABELNAME' FROM THE TOKEN STRING
* * *
* * * CKLANLABSEP TOKEN SPAN(' ') @LABSTART REM . LABELNAME :S(CKLABELCNT)
* * *      LABELNAME = TOKEN
* * *      LABSTART = 0
* * *
* * * COMPUTE THE COLUMN NUMBER THE LABEL ACTUALLY STARTS IN
* * *
* * * CKLABELCNT LABSTART = STARTCOL + LABSTART
* * *
* * * BRANCH IF THE LABEL MUST START WITHIN A SEQUENCE OF COLUMNS
* * *
* * *      NE(LANLABPAT<1>,1)
* * *      :S(SEQUFCOLS)
* * *
* * * IF THE LABEL STARTS IN THE CORRECT SINGLE COLUMN, BRANCH BECAUSE

```

```

PRE11830
PRE11840
PRE11850
PRE11860
PRE11870
PRE11880
PRE11890
PRE11900
PRE11910
PRE11920
PRE11930
PRE11940
PRE11950
PRE11960
PRE11970
PRE11980
PRE11990
PRE12000
PRE12010
PRE12020
PRE12030
PRE12040
PRE12050
PRE12060
PRE12070
PRE12080
PRE12090
PRE12100
PRE12110
PRE12120
PRE12130
PRE12140
PRE12150
PRE12160
PRE12170
PRE12180
PRE12190
PRE12200
PRE12210
PRE12220
PRE12230
PRE12240
PRE12250
PRE12260
PRE12270
PRE12280
PRE12290
PRE12300

```



```

** ** **
A LABEL HAS BEEN FOUND; OTHERWISE, BRANCH BECAUSE THIS IS
NOT A LABEL
** ** **
EQ(LANLABPAT<4>,LABSTART)
:S(FOUNDLABEL)F(NOLABEL)
** ** **
ERROR IF THE LANGUAGE LABEL DEFINITION SPECIFIES OTHER THAN
A SINGLE COLUMN OR A SEQUENCE OF COLUMNS (1 OR 2 IN
LANLABPAT<1>) FOR A LABEL BEGINNING
** ** **
SEQOFCOLS EQ(LANLABPAT<1>,2)
:F(LANLABDEFNERR)
** ** **
NOT A LABEL IF THE LABEL STARTS IN A COLUMN WHICH IS GREATER THAN
THE UPPER BOUND OF THE SEQUENCE OR LESS THAN THE LOWER BOUND
** ** **
GT(LABSTART,LANLABPAT<5>)
LT(LABSTART,LANLABPAT<4>)
:S(NOLABEL)
:S(NOLABEL)
** ** **
A LABEL HAS BEEN FOUND IF A STRING TERMINATED THE LABEL
** ** **
LANLABPAT<2> *STR* RPOS(0)
:S(FOUNDLABEL)
** ** **
CHECK THE FIRST CHARACTER OF THE LABEL TO SEE IF IT
IS VALID ACCORDING TO THE LABEL DEFINITION; IF NOT,
THIS IS NOT A LABEL
** ** **
SUBSTR(LABELNAME,1,1) ANY(LANLABPAT<8>) :F(NOLABEL)
** ** **
CHECK THE REMAINING CHARACTERS OF THE LABEL TO SEE IF THEY
ARE VALID ACCORDING TO THE LABEL DEFINITION; IF NOT,
THIS IS NOT A LABEL
** ** **
EQ(SIZE(LABELNAME),1)
LABELNAME LEN(1) SPAN(LANLABPAT<9>) RPOS(0)
:S(FOUNDLABEL)
:F(NOLABEL)S(FOUNDLABEL)
** ** **
THIS IS NOT A LABEL; IF NECESSARY, BACKUP SHOULD BE SET TO 0 AND
A BRANCH TAKEN TO A LABEL WHERE BACKUP WILL TAKE PLACE
** ** **
NOLABEL
RE(BACKUP,1)
BACKUP = 0
FOUNDLABEL
EC(STRINGSEP,1)
:S(NOTLANLABSEP)
:S(NOTLANLABR2CHAR)
:S(LABELCOMMENT)
** ** **
IF COLUMN WAS LABEL INDICATOR AND NO CHARACTER EXISTS BEYOND
THE COLUMN (GETCHAR RETURNS WITH FAILURE), IT IS NOT A LABEL
** ** **
CLDCCL = CHARCOL
OLDCHAR = INCHAR

```

```

**      GETCHAR(' INCHAR', 'CHARCOL')                                PRE12790
**      IF THE NEXT CHARACTER IS A LABEL DELIMITER (BLANK OR SPECIAL PRE12800
**      CHARACTER), THIS IS A LABEL                                     PRE12810
**      INCHAR LABELS                                                PRE12820
**      :S(LABELYES)                                                 PRE12830
**      BACKUP BY RESTORING THE OLD CHARACTER AND ITS COLUMN NUMBER; PRE12840
**      BRANCH BASED ON THE FINDING OF A LABEL OR NOT              PRE12850
**      BACKUPNOLAB INCHAR = OLDCHAR                                  PRE12860
**      CHARCOL = OLDCOL                                              PRE12870
**      INCHAR = OLDCHAR                                              PRE12880
**      CHARCOL = OLDCOL                                              PRE12890
**      COLNUM = OLDCOL                                              PRE12900
**      LABELYES                                                     PRE12910
**      : (NOTLANLABSEP)                                             PRE12920
**      : (LABELCOMMENT)                                             PRE12930
**      ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** ** **^ PRE12940
**      ** ** **^ PRE12950
**      ** ** **^ PRE12951
**      ** ** **^ PRE12952
**      ** ** **^ PRE12953
**      ** ** **^ PRE12954
**      ** ** **^ PRE12955
**      ** ** **^ PRE12956
**      ** ** **^ PRE12960
**      ** ** **^ PRE12970
**      ** ** **^ PRE12980
**      ** ** **^ PRE12990
**      ** ** **^ PRE13000
**      ** ** **^ PRE13010
**      ** ** **^ PRE13020
**      ** ** **^ PRE13030
**      ** ** **^ PRE13040
**      ** ** **^ PRE13050
**      ** ** **^ PRE13060
**      ** ** **^ PRE13070
**      ** ** **^ PRE13080
**      ** ** **^ PRE13090
**      ** ** **^ PRE13091
**      ** ** **^ PRE13092
**      ** ** **^ PRE13100
**      ** ** **^ PRE13110
**      ** ** **^ PRE13111
**      ** ** **^ PRE13112
**      ** ** **^ PRE13120
**      ** ** **^ PRE13130
**      ** ** **^ PRE13140
**      ** ** **^ PRE13150
**      ** ** **^ PRE13160
**      **      GETCHAR(' INCHAR', 'CHARCOL')                                PRE12790
**      IF THE NEXT CHARACTER IS A LABEL DELIMITER (BLANK OR SPECIAL PRE12800
**      CHARACTER), THIS IS A LABEL                                     PRE12810
**      INCHAR LABELS                                                PRE12820
**      :S(LABELYES)                                                 PRE12830
**      BACKUP BY RESTORING THE OLD CHARACTER AND ITS COLUMN NUMBER; PRE12840
**      BRANCH BASED ON THE FINDING OF A LABEL OR NOT              PRE12850
**      BACKUPNOLAB INCHAR = OLDCHAR                                  PRE12860
**      CHARCOL = OLDCOL                                              PRE12870
**      INCHAR = OLDCHAR                                              PRE12880
**      CHARCOL = OLDCOL                                              PRE12890
**      COLNUM = OLDCOL                                              PRE12900
**      LABELYES                                                     PRE12910
**      : (NOTLANLABSEP)                                             PRE12920
**      : (LABELCOMMENT)                                             PRE12930
**      ** ** **^ PRE12940
**      ** ** **^ PRE12950
**      ** ** **^ PRE12951
**      ** ** **^ PRE12952
**      ** ** **^ PRE12953
**      ** ** **^ PRE12954
**      ** ** **^ PRE12955
**      ** ** **^ PRE12956
**      ** ** **^ PRE12960
**      ** ** **^ PRE12970
**      ** ** **^ PRE12980
**      ** ** **^ PRE12990
**      ** ** **^ PRE13000
**      ** ** **^ PRE13010
**      ** ** **^ PRE13020
**      ** ** **^ PRE13030
**      ** ** **^ PRE13040
**      ** ** **^ PRE13050
**      ** ** **^ PRE13060
**      ** ** **^ PRE13070
**      ** ** **^ PRE13080
**      ** ** **^ PRE13090
**      ** ** **^ PRE13091
**      ** ** **^ PRE13092
**      ** ** **^ PRE13100
**      ** ** **^ PRE13110
**      ** ** **^ PRE13111
**      ** ** **^ PRE13112
**      ** ** **^ PRE13120
**      ** ** **^ PRE13130
**      ** ** **^ PRE13140
**      ** ** **^ PRE13150
**      ** ** **^ PRE13160

```

```

LOOK FOR A KEYWORD APPEARING IMMEDIATELY BEFORE THIS  

BLANK; BRANCH TO THE UNIDENTIFIED SECTION IF A  

KEYWORD IS NOT FOUND  

*****  

KEYWORDSKCH TOKEN KEYWORDS • KEYFOUND , , RPOS(O)  

+      :S(ISKEYWORD)  

      :F(UNIDENT)  

*****  

TOKEN KEYWORD2 • KEYFOUND , , RPOS(O)  

*****  

A KEYWORD HAS BEEN FOUND  

*****  

IF THE KEYWORD MATCHED IS 'THEN', ADD 'UTHEN' TO THE WORKING  

STRING FOR  

THE PARSER; BRANCH TO ADD THE TOKEN STRING TO THE CURRENT  

STATEMENT STRING  

*****  

ISKEYWORD KEYFOUND KEYTHEN  

WORKING = WORKING • UTHEN'  

:(NOTATHEN)  

:(ADDOCURSTMT)  

*****  

THE KEYWORD MATCHED IS NOT 'THEN'; ADD THE KEYWORD  

TO THE WORKING STRING FOR THE PARSER  

*****  

NOTATHEN WORKING = WORKING • • KEYFOUND  

*****  

ADD THE TOKEN STRING TO THE CURRENT STATEMENT STRING AND RETURN  

*****  

ADDOCCURSTMT CURSTMT = CURSTMT TOKEN  

TOKEN =  

:(RETURN)  

*****  

UNIDENTIFIED-TO-END-OF-STATEMENT (UEOS) HAS BEEN FOUND  

*****  

THE ITEM LOCATED IS UNIDENTIFIED; GO BACK TO THE BEGINNING OF  

THE SCANNING ROUTINE TO FIND MORE OF THE TOKEN SO  

IT MIGHT BE IDENTIFIED  

*****  

UNIDENT  

:(BEGSCANL)  

*****  

CHECK FOR ERROR AT END OF FILE  

*****

```

```
PRE13485
PRE13486
PRE13490
PRE13500
PRE13510
PRE13520
PRE13530
PRE13540
PRE13550
PRE13560
PRE13570
PRE13580
PRE13590
PRE13600
PRE13610
PRE13620
PRE13630
PRE13631
PRE13632
PRE13633
PRE13634
PRE13635
PRE13636
PRE13637
PRE13638
PPPEI13639
PRE13640
PRE13641
PRE13650
PRE13660
PRE13670
PRE13680
PRE13690
PRE13700
PRE13710
PRE13720
PRE13730
PRE13740
PRE13750
PRE13760
PRE13770
PRE13780
PRE13790
PRE13800
PRE13810
PRE13820
PRE13830
PRE13840
```

END OF FILE WAS FOUND WHEN TRYING TO GET A CHARACTER AT THE
BEGINNING OF THE SCAN ROUTINE (NOT WHEN LOOKING FOR A
SECOND CHARACTER OF A STRING); IF THE TOKEN IS ALL BLANK OR NULL,
TEST IF AN EOS HAS BEEN FOUND BEFORE THIS CALL TO SCAN

S(TESTEOS)
F(ERRREND)
F(ERRREND)

AN EOS HAD BEEN FOUND; ADD END-OF-FILE TO WORKING STRING FOR
PARSER AND RETURN WITH FAILURE

WOPKNG = WORKING + EOF.
COLNUM = COLNUM + 1

ROUTINE TO GET THE NEXT CHARACTER FROM A NON-IGNORED COLUMN;
RETURN THE CHARACTER AND ITS COLUMN NUMBER

TWO PARAMETERS -- CHAR AND COL, THE NAMES (IN STRING FORM)
OF THE VARIABLES TO HOLD THE RETURN VALUES

INCREMENT COLUMN NUMBER MODULO 80

LE(COLNUM,80)
CCOLNUM = I

READ A NEW CARD IF ONE HAS NOT BEEN SCANNED FOR CONTINUATION
OTHERWISE USE THE ONE THAT WAS SCANNED

FQ(SIZE(NEXTCARD),0)

IF THE SCANNER HAS REACHED END OF FILE PREVIOUSLY OR THIS
ROUTINE HAS REACHED END OF FILE PREVIOUSLY, RETURN WITH FAILURE

EQ(ENDOFFILE,I)
EQ(GETEND,I)

READ THE CARD AND BRANCH

CARDIMAGE = INPUT

SET THE FLAG IF JUST NOW REACHED END OF FILE AND

[illegible]


```

*      NOOPENCOL      I = '11'      : (COMENDER)
*      OPENCOL        CCMPAT<4> =    : (COMENDER)
*      CCMPAT<3> =
*
*      SPECIFICATION INDICATES ONLY A COLUMN NUMBER FOR COMMENT OPENER
*
*      I = '01'
*      COMDEF BREAK(' ' ) . COMPAT<4> SPAN(' ' ) = : F(COMERROR)
*
*      TEST FOR COLUMN SPECIFICATION CONTAINING ONLY DIGITS
*
*      CCMPAT<4> SPAN('0123456789') RPOS(0)      : F(COMERROR)
*
*      SAVE COMMENT OPENER 'DESCRIPTOR'
*
*      COMENDER      CCMPAT<1> = I
*      I = '10'
*      COMDEF " " BREAK(" ") . CCMPAT<5> " " =    : F(ENDCOL)
*      COMDEF SPAN(' ') =                          : F(NOCLOSECOL)
*
*      SPECIFICATION INDICATES A SYMBOL FOR COMMENT ENDER
*
*      COMDEF '+' SPAN(' ') REM . COMPAT<6> =      : F(NOCLOSECOL)
*
*      TEST FOR COLUMN SPECIFICATION CONTAINING ONLY DIGITS
*
*      CCMPAT<6> SPAN('0123456789') RPCS(0)      : F(COMERROR)
*
*      SPECIFICATION INDICATES A SYMBOL AND A COLUMN NUMBER FOR ENDER
*
*      I = '11'
*      NOCLOSECOL      CCMPAT<6> =    : (ENDCOMSCAN)
*      ENDCOL          CCMPAT<5> =    : (ENDCOMSCAN)
*
*      SPECIFICATION INDICATES ONLY A COLUMN NUMBER FOR COMMENT CLOSER
*
*      I = '01'
*      COMDEF REM . COMPAT<6> =          : F(COMERROR)
*
*      TEST FOR COLUMN SPECIFICATION CONTAINING ONLY DIGITS
*
*      CCMPAT<6> SPAN('0123456789') RPOS(0)      : F(COMERROR)
*
*      SAVE COMMENT ENDER 'DESCRIPTOR'
*
*      ENDCOMSCAN      CCMPAT<2> = I
*      SCANCOMDEF =
*
*      : (RETURN)

```

PRE14690
 PRE14700
 PRE14710
 PRE14720
 PRE14730
 PRE14740
 PRE14750
 PRE14760
 PRE14770
 PRE14780
 PRE14790
 PRE14800
 PRE14810
 PRE14820
 PRE14830
 PRE14840
 PRE14850
 PRE14860
 PRE14870
 PRE14880
 PRE14890
 PRE14900
 PRE14910
 PRE14920
 PRE14930
 PRE14940
 PRE14950
 PRE14960
 PRE14970
 PRE14980
 PRE14990
 PRE15000
 PRE15010
 PRE15020
 PRE15030
 PRE15040
 PRE15050
 PRE15060
 PRE15070
 PRE15080
 PRE15090
 PRE15100
 PRE15110
 PRE15120
 PRE15130
 PRE15140
 PRE15150
 PRE15160

```

SCANCONDEF CCNDEF BREAK(' ') . CONTPAT<1> SPAN(' ') = :F(NOCNT)
+
*
*
*
*
*
ROUTINE TO SCAN DEFINITION OF CONTINUATION
ONE PARAMETER -- CONDEF
CONTPAT<2> = CONDEF
CONIPAT<2> = SPAN('0123456789') RPOS(0)
CONDEF =
CONIPAT<1> = CONDEF
CCNDEF 'NONE' RPOS(0)
ENDCONTS CANCONDEF =
SCANLITDEF LITDEF BREAK(' ') . LITPAT<1> SPAN(' ') =
*
*
*
*
*
ROUTINE TO SCAN DEFINITION OF LITERAL
ONE PARAMETER -- LITDEF
LITDEF BREAK(' ') . LITPAT<2> SPAN(' ') = :F(NOSPECIAL)
LITDEF REM . LITPAT<3> = :F(ENCLITSCAN)
LITDEF REM . LITPAT<2>
LITPAT<3> =
ENDLITSCAN SCANLITDEF =
SCANLANLABDEF
*
*
*
*
*
ROUTINE TO SCAN DEFINITION OF LANGUAGE LABEL
ONE PARAMETER -- LANLABDEF
LANLABDEF BREAK(' ') . SPEC = :F(LANLABERR)
SPEC BREAK(' ') . LANLABPAT<4> '- ' REM . LANLABPAT<5> =
: F(LANLABPAT<5> =
: F(FINECOPEN)
*
*
*
*
*
SPECIFICATION INDICATES A SEQUENCE OF COLUMNS TO START LABEL
TEST FOR COLUMN SPECIFICATIONS CONTAINING ONLY DIGITS
LANLABPAT<4> SPAN('0123456789') RPOS(0) : F(LANLABERR)
LANLABPAT<5> SPAN('0123456789') RPOS(0) : F(LANLABERR)
*
*
*
*
*

```

```

PRE15170
PRE15180
PRE15190
PRE15200
PRE15210
PRE15220
PRE15230
PRE15240
PRE15250
PRE15260
PRE15270
PRE15280
PRE15290
PRE15300
PRE15310
PRE15311
PRE15312
PRE15320
PRE15330
PRE15340
PRE15350
PRE15351
PRE15352
PRE15360
PRE15370
PRE15380
PRE15390
PRE15400
PRE15410
PRE15420
PRE15421
PRE15422
PRE15430
PRE15440
PRE15450
PRE15460
PRE15461
PRE15462
PRE15470
PRE15480
PRE15490
PRE15500
PRE15510
PRE15520
PRE15530
PRE15540
PRE15550
PRE15560

```



```

* * *
* * * LANLABPAT<1> = 2                                : (CLOSELANLAB)
* * *
* * * TEST FOR COLUMN SPECIFICATION CONTAINING ONLY DIGITS
* * *
* * * ONECOLCLOSE SPEC SPAN('0123456789') RPOS(0)      : F(LANLABERR)
* * *
* * * SPECIFICATION INDICATES A COLUMN TO START LABEL
* * *
* * * LANLABPAT<1> = 1
* * * LANLABPAT<4> = SPEC
* * * LANLABPAT<5> =
* * * CLOSELANLAB LANLABDEF SPAN(' ') BREAK(' ') SPEC = : F(LANLABERR)
* * * SPEC " " BREAK(" ") . LANLABPAT<6> " " = : F(COLCLOSE)
* * *
* * * SPECIFICATION INDICATES A SYMBOL STRING TO TERMINATE LABEL
* * *
* * * LANLABPAT<2> = 'STR'
* * * LANLABPAT<7> =
* * * LANLABPAT<2> = 'COL'
* * * SPEC BREAK(' ') . LANLABPAT<6> ' ' REM . LANLABPAT<7> =
* * * : F(ONECOLCLOSE)
* * *
* * * SPECIFICATION INDICATES A SEQUENCE OF COLUMNS TO TERMINATE A LABEL
* * *
* * * TEST FOR COLUMN SPECIFICATIONS CONTAINING ONLY DIGITS
* * *
* * * LANLABPAT<6> SPAN('0123456789') RPOS(0)          : F(LANLABERR)
* * * LANLABPAT<7> SPAN('0123456789') RPOS(0)          : F(LANLABERR)
* * * LANLABPAT<3> = 2                                : (OPENCHARS)
* * *
* * * TEST FOR COLUMN SPECIFICATION CONTAINING ONLY DIGITS
* * *
* * * ONECOLCLOSE SPEC SPAN('0123456789') RPOS(0)      : F(LANLABERR)
* * *
* * * SPECIFICATION INDICATES A COLUMN TO TERMINATE LABEL
* * *
* * * LANLABPAT<3> = 1
* * * LANLABPAT<6> = SPEC
* * * LANLABPAT<7> =
* * *
* * * DETERMINE SET OF OPENING CHARACTERS
* * *
* * * OPENCHARS LANLABDEF SPAN(' ') " " BREAK(" ") . LANLABPAT<8> " " =
* * * : F(LANLABERR)
* * *
* * * DETERMINE SET OF REMAINING CHARACTERS
* * *
* * * LANLABDEF SPAN(' ') " " BREAK(" ") . LANLABPAT<9>

```

```

PRE15570
PRE15580
PRE15590
PRE15600
PRE15610
PRE15620
PRE15630
PRE15640
PRE15650
PRE15660
PRE15670
PRE15680
PRE15690
PRE15700
PRE15710
PRE15720
PRE15730
PRE15740
PRE15750
PRE15760
PRE15770
PRE15780
PRE15790
PRE15800
PRE15810
PRE15820
PRE15830
PRE15840
PRE15850
PRE15860
PRE15870
PRE15880
PRE15890
PRE15900
PRE15910
PRE15920
PRE15930
PRE15940
PRE15950
PRE15960
PRE15970
PRE15980
PRE15990
PRE16000
PRE16010
PRE16020
PRE16030
PRE16040

```



```
PRE16490
PRE16500      :F(NXT.5)
PRE16510
PRE16520      SCAN AND BRANCH
PRE16530
PRE16540      :S(EREPEATSTMT)F(PSCANERR)
PRE16550
PRE16560      TRY TO MATCH AN ENDDO
PRE16570
PRE16580      WORKING *ENDDO* RPOS(O)
PRE16590      :F(NXT.6)
PRE16600
PRE16610      SCAN AND BRANCH
PRE16620      SCAN()
PRE16630      :S(EDOSTMT)F(PSCANERR)
PRE16640
PRE16650      TRY TO MATCH AN ENDDOGRPOUP
PRE16660
PRE16670      WORKING *ENDDOGRPOUP* RPOS(O)
PRE16680      :F(NXT.7)
PRE16690
PRE16700      SCAN AND BRANCH
PRE16710      SCAN()
PRE16720      :S(EDDGRPSTMT)F(PSCANERR)
PRE16730
PRE16740      TRY TO MATCH AN ELSE
PRE16750      WORKING *ELSE* RPOS(O)
PRE16760      :F(NXT.8)
PRE16770
PRE16780      CALL SEMANTIC ROUTINE I, SCAN AND RETURN TO START
PRE16790
PRE16800      ROUTI()
PRE16810      SCAN()
PRE16820      :S(START)F(PSCANERR)
PRE16830
PRE16840      TRY TO MATCH A QUITLOOP
PRE16850      WORKING *QUITLOOP* RPOS(O)
PRE16860      :F(NXT.9)
PRE16870
PRE16880      SCAN AND BRANCH
PRE16890      SCAN()
PRE16900      :S(QUITSTMT)F(PSCANERR)
PRE16910
PRE16920      TRY TO MATCH A CASE
PRE16930      WORKING *CASE* RPOS(O)
PRE16940      :F(NXT.10)
PRE16950
PRE16960      SCAN AND BRANCH
PRE16970      SCAN()
PRE16980      :S(CASESTMT)F(PSCANERR)
PRE16990
PRE17000
```

```

* * TRY TO MATCH A DO
* *
* * NEXT.10      WORKING 'DO' RPOS(0)
* *
* * SCAN AND BRANCH
* *
* * SCAN()
* *
* * TRY TO MATCH A DOGROUP
* *
* * NEXT.11      WORKING 'DOGROUP' RPOS(0)
* *
* * SCAN AND BRANCH
* *
* * SCAN()
* *
* * TRY TO MATCH AN IF
* *
* * NEXT.12      WORKING 'IF' RPOS(0)
* *
* * SCAN AND BRANCH
* *
* * SCAN()
* *
* * TRY TO MATCH A REPEAT
* *
* * NEXT.13      WORKING 'REPEAT' RPOS(0)
* *
* * SCAN AND BRANCH
* *
* * SCAN()
* *
* * TRY TO MATCH A WHILE
* *
* * NEXT.14      WORKING 'WHILE' RPOS(0)
* *
* * SCAN AND BRANCH
* *
* * SCAN()
* *
* * TRY TO MATCH A STRUCTURE LIST AND END OF FILE; IF MATCHED,
* * RETURN TO CALLING ROUTINE
* *
* * NEXT.15      WORKING 'STRLIST EOF' = 'PROG'
* *
* * TRY TO MATCH AN UNIDENTIFIED TO END OF STATEMENT; IF NOT
* * MATCHED TAKE ERROR EXIT
* *

```

PRE16970
 PRE16980
 PRE16990
 PRE17000
 PRE17010
 PRE17020
 PRE17030
 PRE17040
 PRE17050
 PRE17060
 PRE17070
 PRE17080
 PRE17090
 PRE17100
 PRE17110
 PRE17120
 PRE17130
 PRE17140
 PRE17150
 PRE17160
 PRE17170
 PRE17180
 PRE17190
 PRE17200
 PRE17210
 PRE17220
 PRE17230
 PRE17240
 PRE17250
 PRE17260
 PRE17270
 PRE17280
 PRE17290
 PRE17300
 PRE17310
 PRE17320
 PRE17330
 PRE17340
 PRE17350
 PRE17360
 PRE17370
 PRE17380
 PRE17390
 PRE17400
 PRE17410
 PRE17420
 PRE17430
 PRE17440

```

* * * WORKING 'UEOS' RPOS(0) = 'STRUC' :F(PUNIDERR) PRE17450
* * * CALL SEMANTIC ROUTINE 01 AND BRANCH TO STRUCTURE SECTION PRE17460
* * * ROUT01() :(STRUCTURE) PRE17470
* * * CASE HAS BEEN RECOGNIZED, REDUCE TO CASESTMT; CALL SEMANTIC PRE17480
* * * ROUTINE 2, SCAN AND RETURN TO START PRE17490
* * * CASESTMT WORKING 'CASE UEOS' RPOS(0) = 'CASESTMT' :F(PCASEERR) PRE17500
* * * ROUT2() :S(START)F(RETURN) PRE17510
* * * SCAN() PRE17520
* * * DO HAS BEEN RECOGNIZED, REDUCE TO DOSTMT; CALL SEMANTIC PRE17530
* * * ROUTINE 3, SCAN AND RETURN TO START PRE17540
* * * DOSTMT WORKING 'DO UEOS' RPOS(0) = 'DOSTMT' :F(PDOPERR) PRE17550
* * * ROUT3() :S(START)F(RETURN) PRE17560
* * * SCAN() PRE17570
* * * DOGROUP HAS BEEN RECOGNIZED, REDUCE TO DOGROUPSTMT; CALL SEMANTIC PRE17580
* * * ROUTINE 4, SCAN AND RETURN TO START PRE17590
* * * DOGRPSTMT WORKING 'DOGROUP UEOS' RPOS(0) = :F(PDGRPERR) PRE17600
* * * ROUT4() :S(START)F(RETURN) PRE17610
* * * SCAN() PRE17620
* * * IF HAS BEEN RECOGNIZED, REDUCE TO IFTHEN; CALL SEMANTIC ROUTINE PRE17630
* * * 5, SCAN AND RETURN TO START PRE17640
* * * IFSTMT WORKING 'IF UTHEN' RPOS(0) = 'IFTHEN' :F(PIFERR) PRE17650
* * * ROUT5() :S(START)F(RETURN) PRE17660
* * * SCAN() PRE17670
* * * REPEAT HAS BEEN RECOGNIZED, REDUCE TO REPUNTIL; SCAN PRE17680
* * * REPSTMT WORKING 'REPEAT UNTIL' RPOS(0) = :F(PREPERR) PRE17690
* * * :F(PSCANERR) PRE17700
* * * REPUNTIL HAS BEEN RECOGNIZED, REDUCE TO REPEATSTMT; CALL SEMANTIC PRE17710
* * * ROUTINE 6, SCAN AND RETURN TO START PRE17720
* * * REPSTMT WORKING 'REPUNTIL UEOS' RPOS(0) = :F(PREPERR) PRE17730
* * * ROUT6() :F(PSCANERR) PRE17740
* * * SCAN() PRE17750
* * * :F(PREPERR) PRE17760
* * * :F(PSCANERR) PRE17770
* * * :S(START)F(PSCANERR) PRE17780
* * * :F(PREPERR) PRE17790
* * * :F(PSCANERR) PRE17800
* * * :F(PREPERR) PRE17810
* * * :F(PSCANERR) PRE17820
* * * :F(PREPERR) PRE17830
* * * :F(PSCANERR) PRE17840
* * * :F(PREPERR) PRE17850
* * * :F(PSCANERR) PRE17860
* * * :F(PREPERR) PRE17870
* * * :F(PSCANERR) PRE17880
* * * :F(PREPERR) PRE17890
* * * :F(PSCANERR) PRE17900
* * * :F(PREPERR) PRE17910
* * * :F(PSCANERR) PRE17920

```

```

*      WHILE HAS BEEN RECOGNIZED, REDUCE TO WHILESTMT; CALL SEMANTIC
*      ROUTINE 7, SCAN AND RETURN TO START
*
*WHILESTMT      WORKING 'WHILE UEOS' RPOS(0) =
+              'WHILESTMT'
+              :F(PWHERR)
+
*              :S(START)F(RETURN)
*
*      QUITLOOP HAS BEEN RECCGNIZED, REDUCE TO STRUCTURE; CALL SEMANTIC
*      ROUTINE 8 AND BRANCH TO STRUCTURE SECTION
*
*QUITSTMT      WORKING 'QUITLOOP UEOS' RPOS(0) =
+              'STRUC'
+              :F(PQUERR)
+              :(:STRUCTURE)
+
*      REDUCTION OF CASE-ENDCASE TO STRUCTURE; CALL SEMANTIC
*      ROUTINE 9 AND BRANCH TO STRUCTURE SECTION
*
*{CASESTMT      WORKING 'CASESTMT STRLIST ENDCASE UEOS'
+              RPOS(0) = 'STRUC'
+              :F(PECASEERR)
+              :(:STRUCTURE)
+
*      REDUCTION OF IF-THEN-ENDIF TO STRUCTURE; CALL SEMANTIC ROUTINE
*      10 AND BRANCH TO STRUCTURE SECTION
*
*EIFSTMT      WORKING 'IFTHCL ENDIF UEOS' RPOS(0) =
+              'STRUC'
+              :F(NXT.16)
+              :(:STRUCTURE)
+
*      REDUCTION OF IF-THEN-ELSE-ENDIF TO STRUCTURE; CALL SEMANTIC ROUTINE
*      10 AND BRANCH TO STRUCTURE SECTION
*
*NXT.16      WORKING 'IFTHELSE ENDIF UEOS' RPOS(0) =
+              'STRUC'
+              :F(PIFTHEKR)
+              :(:STRUCTURE)
+
*      REDUCTION OF WHILE-ENDWHILE TO STRUCTURE; CALL SEMANTIC
*      ROUTINE 11 AND BRANCH TO STRUCTURE SECTION
*
*FWHILESTMT      WORKING 'WHILESTMT STRLIST ENDWHILE UEOS'
+              RPOS(0) = 'STRUC'
+              :F(PEWHERR)
+              :(:STRUCTURE)
+
*      REDUCTION OF REPEAT-ENDREPEAT TO STRUCTURE; CALL SEMANTIC
*      ROUTINE 12 AND BRANCH TO STRUCTURE SECTION
*
*ERFPEATSTMT      WORKING 'REPEATSTMT STRLIST ENDREPEAT UEOS' RPOS(0) =

```

```

PRE17930
PRE17940
PRE17950
PRE17960
PRE17970
PRE17980
PRE17990
PRE18000
PRE18010
PRE18020
PRE18030
PRE18040
PRE18050
PRE18060
PRE18070
PRE18080
PRE18090
PRE18100
PRE18110
PRE18120
PRE18130
PRE18140
PRE18150
PRE18160
PRE18170
PRE18180
PRE18190
PRE18200
PRE18210
PRE18220
PRE18230
PRE18240
PRE18250
PRE18260
PRE18270
PRE18280
PRE18290
PRE18300
PRE18310
PRE18320
PRE18330
PRE18340
PRE18350
PRE18360
PRE18370
PRE18380
PRE18390
PRE18400

```

```

+      ROUT12( )      'STRUC'
*      REDUCTION OF DO-ENDDO TO STRUCTURE; CALL SEMANTIC ROUTINE
*      13 AND BRANCH TO STRUCTURE SECTION
*      RETURN TO START
+      EDOSTMT      WORKING 'DCSTMT STRLIST ENDDO UEOS'
*      ROUT13( )      RPOS(0) = 'STRUC'
*      REDUCTION OF DOGROUP-ENDDOGROUP TO STRUCTURE; CALL SEMANTIC
*      ROUTINE 14 AND BRANCH TO STRUCTURE SECTION
+      EDGRPSTMT      WORKING 'DOGROUPSMT STRLIST ENDDOGROUP UEOS'
*      ROUT14( )      RPOS(0) = 'STRUC'
*      LABEL AND A STRUCTURE REDUCE TO STRUCTURE; RETURN TO STRUCTURE
*      SECTION
+      STRUCTURE      WORKING 'LABEL STRUC' RPOS(0) =
*      'STRUC'      :S(STRUCTURE)
*      ELSE AND STRUCTURE REDUCE TO ELSESTRUCTURE; RETURN TO STRUCTURE
*      SECTION
+      WORKING 'ELSE STRUC' RPOS(0) = 'ELSESTRUC'
*      :S(STRUCTURE)
*      LABEL AND ELSE STRUCTURE REDUCE TO ELSESTRUCTURE; RETURN
*      TO STRUCTURE SECTION
+      WORKING 'LABEL ELSESTRUC' RPOS(0) =
*      'ELSESTRUC'      :S(STRUCTURE)
*      IFTHEN CLAUSE AND ELSESTRUC REDUCE TO IFTHELSE; SCAN AND
*      RETURN TO START
+      WORKING 'IFTHCL ELSESTRUC' RPOS(0) =
*      'IFTHELSE'      :S(START)F(RETURN)
*      SCAN( )
*      IF-THEN AND STRUCTURE REDUCE TO IF THEN CLAUSE; SCAN AND
*      RETURN TO START
+      WORKING 'IFTHEN STRUC' RPOS(0) =
*      'IFTHCL'      :F(NXT.18)
+      NXT.17

```

```

PRE18410
PRE18420
PRE18430
PRE18440
PRE18450
PRE18460
PRE18470
PRE18480
PRE18490
PRE18500
PRE18510
PRE18520
PRE18530
PRE18540
PRE18550
PRE18560
PRE18570
PRE18580
PRE18590
PRE18600
PRE18610
PRE18620
PRE18630
PRE18640
PRE18650
PRE18660
PRE18670
PRE18680
PRE18690
PRE18700
PRE18710
PRE18720
PRE18730
PRE18740
PRE18750
PRE18760
PRE18770
PRE18780
PRE18790
PRE18800
PRE18810
PRE18820
PRE18830
PRE18840
PRE18850
PRE18860
PRE18870
PRE18880

```



```

*      READ THE MACRO DEFINITION AND PLACE IT IN MACREPSTR
*
MACREP
  MACREPSTR = INPUT
  MACDEF POS(0) 'END MACRO DEFINITION' SPAN(' ') RPOS(0)
  +      :S(ENDMACDEF)
  +      :S(MACREP)
*
*      ADD THE MACRONAME TO THE MACRONAME PATTERN
  ENDMACDEF MACNAMEPAT = MACNAMEPAT | NAME
*
*      START THE PARAMETER PATTERN WITH A SERIES OF 81 BLANKS
  PARAM = DUPL(' ',81)
*
*      ADD THE PARAMETERS TO THE PARAMETER PATTERN
  PARMPAT
  +      PARAM = PARAM | MPARM<I>
  +      I = GT(I,1) | -1
  +      :S(PARMPAT)
*
*      BREAK THE MACRO DEFINITION INTO COMPONENTS AND PLACE THEM IN
  +      ELEMENTS 1-20 OF MPAT
  +
  +      I = -1
  +      MACDEFLOOP
  +      I = I + 2
  +      MACREPSTR ARB = COPY PARAM • PARAM =
  +      MACREPSTR POS(0) SPAN(' ') • BLKS =
  +      MPAT<I> = COPY
  +      MPAT<I + 1> = PARAM BLKS
  +      BLKS =
  +      GT(SIZE(MACREPSTR),0)
  +      MPAT<I> = MACREPSTR
  +      :S(MACDEFLOOP)F(MACDEFLEN)
  +
  +      PLACE THE NUMBER OF COMPONENTS IN MPAT<0>
  +      MPAT<0> = I
  +      MACDEFLEN MPAT<0> = I + 1
  +      :S(RETMACDEF)
  +
  +      STORE THE MPAT ARRAY IN ELEMENT 'MACRONAME' OF TABLE MACROTAB
  +
  +      STORE THE MPARM ARRAY IN ELEMENT 'MACRONAME P' OF TABLE MACROTAB
  +      AND RETURN TO CALLING ROUTINE
  +
  +      RETMACDEF
  +      MACROTAB<NAME> = COPY(MPAT)
  +      MACROTAB<NAME 'P'> = COPY(MPARM)
  +      FINOUT = NAME • MACRO DEFINITION PROCESSED. •
  +      FINOUT = • •
  +      FINOUT = • •
  +      :S(RETURNS)
  +
  +      SCANDIFS
  +      EANCHOR = 1
  +
*
PRE19330
PRE19340
PRE19350
PRE19360
PRE19370
PRE19380
PRE19390
PRE19400
PRE19410
PRE19420
PRE19430
PRE19440
PRE19450
PRE19460
PRE19470
PRE19480
PRE19490
PRE19500
PRE19510
PRE19520
PRE19530
PRE19540
PRE19550
PRE19560
PRE19570
PRE19580
PRE19590
PRE19600
PRE19610
PRE19620
PRE19630
PRE19640
PRE19650
PRE19660
PRE19670
PRE19680
PRE19690
PRE19700
PRE19710
PRE19720
PRE19730
PRE19740
PRE19750
PRE19760
PRE19770
PRE19780
PRE19790
PRE19800

```

```

**
**
** ROUTINE TO SCAN DEFINITIONS AND CALL APPROPRIATE
** INDIVIDUAL SCANNING ROUTINE
**
** ONE PARAMETER -- CARD: CARD IS THE FIRST CARD OF THE
** FIRST DEFINITION
**
**
** I = 1
** DEFNAME = 'NULL STATEMENT' | 'GO TO STATEMENT' |
** | 'IF THEN GO TO STATEMENT' | 'ASSIGN MINUS STATEMENT' |
** | 'ASSIGN PLUS STATEMENT' | 'DECLARATION STATEMENT' |
** | 'LITERAL' | 'COMMENT' | 'UNIQUE LABEL' |
** | 'UNIQUE IDENTIFIER' | 'GREATER THAN' | 'NOT' |
** | 'LESS THAN OR EQUAL' | 'GREATER THAN OR EQUAL' |
** | 'NOT EQUAL' | 'CONTINUATION' | 'END OF STATEMENT' |
** | 'IGNORE COLUMNS' | 'LANGLABEL'
**
** DETERMINE DEFINITION NAME
**
** NEXTDEF CARD '<' DEFNAME SELECT '>::=' = :F(ILLLEGALDEF)
** SELECT ARB . STMTTYPE ' STATEMENT' RPOS(0) :F(NOTASTMT)
**
** BRANCH IF A STATEMENT DEFINITION
**
** STMTTYPE 'END OF' RPOS(0) :F(NOTEOS)
**
** IF END OF STATEMENT DEFINITION, TRIM TRAILING BLANKS
** AND BRANCH
**
** CARD = TRIM(CARD)
** : (EOSSCAN)
**
** STATEMENT PROTOTYPES REQUIRE TWO INPUT CARDS, SO READ
** ANOTHER CARD AND TRIM THE TRAILING BLANKS FROM THE
** ENTIRE DEFINITION
**
** NOTEOS CARD = TRIM(CARD INPUT)
**
** BRANCH TO THE APPROPRIATE LABEL
**
** STMTTYPE 'NULL' RPOS(0) :S(NULLSCAN)
** STMTTYPE 'GO TO' RPOS(0) :S(GOTOSCAN)
** STMTTYPE 'IF THEN GO TO' RPOS(0) :S(IFSCAN)
** STMTTYPE 'ASSIGN MINUS' RPOS(0) :S(MINUSSCAN)
** STMTTYPE 'ASSIGN PLUS' RPOS(0) :S(PLUSSCAN)
**

```

```

PRE19801
PRE19802
PRE19810
PRE19820
PRE19830
PRE19840
PRE19850
PRE19860
PRE19861
PRE19862
PRE19870
PRE19880
PRE19890
PRE19900
PRE19910
PRE19920
PRE19930
PRE19940
PRE19950
PRE19960
PRE19970
PRE19980
PRE19990
PRE20000
PRE20010
PRE20020
PRE20030
PRE20040
PRE20050
PRE20060
PRE20070
PRE20080
PRE20090
PRE20100
PRE20110
PRE20120
PRE20130
PRE20140
PRE20150
PRE20160
PRE20170
PRE20180
PRE20190
PRE20200
PRE20210
PRE20220
PRE20230
PRE20240

```

```

* 1
* 2
* 3
* 4
* 5
* 6
* 7
* 8
* 9
* 10
* 11
* 12
* 13
* 14
* 15
* 16
* 17
* 18
* 19
* 20
* 21
* 22
* 23
* 24
* 25
* 26
* 27
* 28
* 29
* 30
* 31
* 32
* 33
* 34
* 35
* 36
* 37
* 38
* 39
* 40
* 41
* 42
* 43
* 44
* 45
* 46
* 47
* 48
* 49
* 50
* 51
* 52
* 53
* 54
* 55
* 56
* 57
* 58
* 59
* 60
* 61
* 62
* 63
* 64
* 65
* 66
* 67
* 68
* 69
* 70
* 71
* 72
* 73
* 74
* 75
* 76
* 77
* 78
* 79
* 80
* 81
* 82
* 83
* 84
* 85
* 86
* 87
* 88
* 89
* 90
* 91
* 92
* 93
* 94
* 95
* 96
* 97
* 98
* 99
* 100
* 101
* 102
* 103
* 104
* 105
* 106
* 107
* 108
* 109
* 110
* 111
* 112
* 113
* 114
* 115
* 116
* 117
* 118
* 119
* 120
* 121
* 122
* 123
* 124
* 125
* 126
* 127
* 128
* 129
* 130
* 131
* 132
* 133
* 134
* 135
* 136
* 137
* 138
* 139
* 140
* 141
* 142
* 143
* 144
* 145
* 146
* 147
* 148
* 149
* 150
* 151
* 152
* 153
* 154
* 155
* 156
* 157
* 158
* 159
* 160
* 161
* 162
* 163
* 164
* 165
* 166
* 167
* 168
* 169
* 170
* 171
* 172
* 173
* 174
* 175
* 176
* 177
* 178
* 179
* 180
* 181
* 182
* 183
* 184
* 185
* 186
* 187
* 188
* 189
* 190
* 191
* 192
* 193
* 194
* 195
* 196
* 197
* 198
* 199
* 200
* 201
* 202
* 203
* 204
* 205
* 206
* 207
* 208
* 209
* 210
* 211
* 212
* 213
* 214
* 215
* 216
* 217
* 218
* 219
* 220
* 221
* 222
* 223
* 224
* 225
* 226
* 227
* 228
* 229
* 230
* 231
* 232
* 233
* 234
* 235
* 236
* 237
* 238
* 239
* 240
* 241
* 242
* 243
* 244
* 245
* 246
* 247
* 248
* 249
* 250
* 251
* 252
* 253
* 254
* 255
* 256
* 257
* 258
* 259
* 260
* 261
* 262
* 263
* 264
* 265
* 266
* 267
* 268
* 269
* 270
* 271
* 272
* 273
* 274
* 275
* 276
* 277
* 278
* 279
* 280
* 281
* 282
* 283
* 284
* 285
* 286
* 287
* 288
* 289
* 290
* 291
* 292
* 293
* 294
* 295
* 296
* 297
* 298
* 299
* 300
* 301
* 302
* 303
* 304
* 305
* 306
* 307
* 308
* 309
* 310
* 311
* 312
* 313
* 314
* 315
* 316
* 317
* 318
* 319
* 320
* 321
* 322
* 323
* 324
* 325
* 326
* 327
* 328
* 329
* 330
* 331
* 332
* 333
* 334
* 335
* 336
* 337
* 338
* 339
* 340
* 341
* 342
* 343
* 344
* 345
* 346
* 347
* 348
* 349
* 350
* 351
* 352
* 353
* 354
* 355
* 356
* 357
* 358
* 359
* 360
* 361
* 362
* 363
* 364
* 365
* 366
* 367
* 368
* 369
* 370
* 371
* 372
* 373
* 374
* 375
* 376
* 377
* 378
* 379
* 380
* 381
* 382
* 383
* 384
* 385
* 386
* 387
* 388
* 389
* 390
* 391
* 392
* 393
* 394
* 395
* 396
* 397
* 398
* 399
* 400
* 401
* 402
* 403
* 404
* 405
* 406
* 407
* 408
* 409
* 410
* 411
* 412
* 413
* 414
* 415
* 416
* 417
* 418
* 419
* 420
* 421
* 422
* 423
* 424
* 425
* 426
* 427
* 428
* 429
* 430
* 431
* 432
* 433
* 434
* 435
* 436
* 437
* 438
* 439
* 440
* 441
* 442
* 443
* 444
* 445
* 446
* 447
* 448
* 449
* 450
* 451
* 452
* 453
* 454
* 455
* 456
* 457
* 458
* 459
* 460
* 461
* 462
* 463
* 464
* 465
* 466
* 467
* 468
* 469
* 470
* 471
* 472
* 473
* 474
* 475
* 476
* 477
* 478
* 479
* 480
* 481
* 482
* 483
* 484
* 485
* 486
* 487
* 488
* 489
* 490
* 491
* 492
* 493
* 494
* 495
* 496
* 497
* 498
* 499
* 500
* 501
* 502
* 503
* 504
* 505
* 506
* 507
* 508
* 509
* 510
* 511
* 512
* 513
* 514
* 515
* 516
* 517
* 518
* 519
* 520
* 521
* 522
* 523
* 524
* 525
* 526
* 527
* 528
* 529
* 530
* 531
* 532
* 533
* 534
* 535
* 536
* 537
* 538
* 539
* 540
* 541
* 542
* 543
* 544
* 545
* 546
* 547
* 548
* 549
* 550
* 551
* 552
* 553
* 554
* 555
* 556
* 557
* 558
* 559
* 560
* 561
* 562
* 563
* 564
* 565
* 566
* 567
* 568
* 569
* 570
* 571
* 572
* 573
* 574
* 575
* 576
* 577
* 578
* 579
* 580
* 581
* 582
* 583
* 584
* 585
* 586
* 587
* 588
* 589
* 590
* 591
* 592
* 593
* 594
* 595
* 596
* 597
* 598
* 599
* 600
* 601
* 602
* 603
* 604
* 605
* 606
* 607
* 608
* 609
* 610
* 611
* 612
* 613
* 614
* 615
* 616
* 617
* 618
* 619
* 620
* 621
* 622
* 623
* 624
* 625
* 626
* 627
* 628
* 629
* 630
* 631
* 632
* 633
* 634
* 635
* 636
* 637
* 638
* 639
* 640
* 641
* 642
* 643
* 644
* 645
* 646
* 647
* 648
* 649
* 650
* 651
* 652
* 653
* 654
* 655
* 656
* 657
* 658
* 659
* 660
* 661
* 662
* 663
* 664
* 665
* 666
* 667
* 668
* 669
* 670
* 671
* 672
* 673
* 674
* 675
* 676
* 677
* 678
* 679
* 680
* 681
* 682
* 683
* 684
* 685
* 686
* 687
* 688
* 689
* 690
* 691
* 692
* 693
* 694
* 695
* 696
* 697
* 698
* 699
* 700
*
```

```

* LANGLABELSCAN CARD BREAK(' ') . . . BREAK(' ') . . . REM . REST
REST SPAN(' ') =
REST " " BREAK(" ") " " =
NEEDONE CARD = TRIM(CARD INPUT)
* * CALL THE LANGUAGE LABEL DEFINITION SCANNER
* *
* * CALLSCAN SCANLANLABDEF(CARD)
* *
* * DETERMINE IF THE ENTIRE REMAINING CHARACTER SET IS CONTAINED
* * ON THE FIRST CARD: IF NOT, BRANCH TO READ IN ANOTHER CARD
* * AND TRIM THE TRAILING BLANKS: IF SO, TRIM THE TRAILING
* * BLANKS AND BRANCH TO CALL THE SCANNING ROUTINE
* *
* * LASTSET REST SPAN(' ') =
* * REST " " BREAK(" ") " "
* * CARD = TRIM(CARD)
* * SCANOPERDEF('NOT', CARD)
* * CONTINUATIONSCAN SCANCONTRDEF(CARD)
* * COMMENTSCAN SCANCOMDEF(CARD)
* * LITERALS SCAN SCANLITDEF(CARD)
* * NULLSCAN SCANNULLDEF(CARD)
* * GOTOSCAN SCANGOTODEF(CARD)
* * IFSCAN SCANIFDEF(CARD)
* * MINUSSCAN SCANMINUSDEF(CARD)
* * PLUSSCAN SCANPLUSDEF(CARD)
* * DCLSCAN SCANDCLDEF(CARD)
* * LABELSCAN SCANLABREDEF(CARD)
* * IDENTIFIERS SCAN SCANVARDEF(CARD)
* * EDSCAN SCANEDEDEF(CARD)
* * IGNSCAN SCANIGNCOLDEF(CARD)
* * NEQSCAN SCANOPERDEF('NEQU', CARD)
* * GTHAMSCAN SCANOPERDEF('GTTHAN', CARD)
* * GEQUSCAN SCANOPERDEF('GEQU', CARD)
* * LEQUSCAN SCANOPERDEF('LEQU', CARD)
* *
* * RETURN IF THIS WAS THE LAST DEFINITION
* *
* * NEXTIN EQ(I, 19)
* * &ANCHOR = 0
* *
* * INCREMENT DEFINITION COUNTER AND READ ANOTHER CARD
* *
* * INCRDEFNCTR I = I + 1
* * CARD = INPUT
* * SCANOPERDEF OPERDEF " " ARR . $TYPEDEF " " KPOS(J) =
* * :F(DEFERR)
* *
PRE20730
PRE20740
PRE20750
PRE20760
PRE20770
PRE20780
PRE20790
PRE20800
PRE20810
PRE20820
PRE20830
PRE20840
PRE20850
PRE20860
PRE20870
PRE20880
PRE20890
PRE20900
PRE20910
PRE20920
PRE20930
PRE20940
PRE20950
PRE20960
PRE20970
PRE20980
PRE20990
PRE21000
PRE21010
PRE21020
PRE21030
PRE21040
PRE21050
PRE21060
PRE21070
PRE21080
PRE21090
PRE21100
PRE21110
PRE21120
PRE21130
PRE21140
PRE21150
PRE21160
PRE21170
PRE21180
PRE21190
PRE21200

```


[illegible]

```

ROUTINE TO SCAN DEFINITION OF IF THEN GO TO STATEMENT
ONE PARAMETER -- IDEF
**
** PARM = 'LAB' | 'VAL1' | 'RL' | 'VAL2' | 'DES'
J = -1
I = 0
IFDEFLOOP
  IDEF ARB . COPY PARM . PARAM =
  IDEF POS(0) SPAN(' ') . BLKS =
  J = J + 1
  IFPAT<I>= COPY
    ** SAVE PARAMETER NAME AND BLANKS FOLLOWING IT
    IFPAT<I+1>= PARAM BLKS
    BLKS =
    GT(SIZE(IFDEF),0)
      **
      ** IDEFEND
      ** IDEFLEN
      ** UNDEFDEF
      ** TEST FOR ALL PARAMETERS SPECIFIED
      ** EQ(J,5)
      ** SCANDCLDEF
      ** ROUTINE TO SCAN DEFINITION OF DECLARATION STATEMENT
      ONE PARAMETER -- DCLDEF
      **
      ** I = -1
      ** J = 0
      ** I DCLDEF ARB . COPY 'VARNAME' =
      ** DCLDEF POS(0) SPAN(' ') . BLKS =
      **
      ** IDEFEND

```



```

**      J = J + 1
**      DCLPAT<I> = COPY
**
**      SAVE PARAMETER NAME AND BLANKS FOLLOWING IT
**
**      DCLPAT<I + 1> = 'VARNAME' BLKS
**      BLKS =
**      GT(SIZE(DCLDEF),0)
**
**      DCLDEFEND DCLPAT<I> = DCLDEF
**      DCLPAT<O> = I
**      DCLDEF =
**      DCLDEFLEN DCLPAT<O> = I + 1
**      ENDDCLDEF SCANDCLDEF =
**      SCANMINUSDEF
**
**      ROUTINE TO SCAN DEFINITION OF ASSIGN MINUS STATEMENT
**
**      ONE PARAMETER -- MINUSDEF
**
**      PARM = 'LAB' | 'VAR' | 'VAL1' | 'VAL2'
**      I = -1
**      J = 0
**      MINUSDEFLOOP I = I + 2
**      MINUSDEF ARB COPY PARM • PARAM =
**      MINUSDEF POS(0) SPAN(•,•) • BLKS =
**      J = J + 1
**      MINUSPAT<I> = COPY
**
**      SAVE PARAMETER NAME AND BLANKS FOLLOWING IT
**
**      MINUSPAT<I + 1> = PARAM BLKS
**      BLKS =
**      GT(SIZE(MINUSDEF),0)
**
**      MINUSDEFEND MINUSPAT<I> = MINUSDEF
**      MINUSPAT<O> = I
**      MINUSDEF =
**      MINUSDEFLEN MINUSPAT<O> = I + 1
**      ENDMINUSDEF SCANMINUSDEF =
**
**      TEST FOR ALL PARAMETERS SPECIFIED
**
**      EQ(J,4)

```

```

PRE222450
PRE222460
PRE222470
PRE222480
PRE222490
PRE222500
PRE222510
PRE222520
PRE222530
PRE222540
PRE222550
PRE222560
PRE222570
PRE222580
PRE222590
PRE222600
PRE222601
PRE222602
PRE222610
PRE222620
PRE222630
PRE222640
PRE222641
PRE222642
PRE222650
PRE222660
PRE222670
PRE222680
PRE222690
PRE222700
PRE222710
PRE222720
PRE222730
PRE222740
PRE222750
PRE222760
PRE222770
PRE222780
PRE222790
PRE222800
PRE222810
PRE222820
PRE222830
PRE222840
PRE222850
PRE222860
PRE222870
PRE222880

```



```

+SCANPLUSDEF
**
ROUTINE TO SCAN DEFINITION OF ASSIGN PLUS STATEMENT
ONE PARAMETER -- PLUSDEF
**
PARAM = 'LAH' | 'VAR' | 'VAL1' | 'VAL2'
I = -1
J = 0
PLUSDEFLOOP I = I + 2
PLUSDEF ARB . COPY PARAM . PARAM =
PLUSDEF POS(0) SPAN(' ') . BLKS =
J = J + 1
PLUSPAT<I> = COPY
**
SAVE PARAMETER NAME AND BLANKS FOLLOWING IT
PLUSPAT<I + 1> = PARAM BLKS
BLKS =
GT(SIZE(PLUSDEF),0)
+SPLUSDEFEND PLUSPAT<I> = PLUSDEF
PLUSPAT<O> = I
PLUSDEF =
PLUSDEFLEN PLUSPAT<O> = I + 1
ENDPLUSDEF SCANPLUSDEF =
**
TEST FOR ALL PARAMETERS SPECIFIED
EQ(J,4)
+SCANLABFDEF
**
ROUTINE TO SCAN DEFINITION OF UNIQUE LABEL
ONE PARAMETER -- LABELDEF
**
LABELDEF BREAK(' ') . LABELPAT<I> SPAN(' ') = F(LABELERROR)

```

```

LABELDEF BREAK(' ',) * LABELPAT<2> =          :F(NULABEND)
LABELDEF SPAN(' ',) = LABELDEF
LABELPAT<3> = LABELDEF
LABELDEF =
LABELPAT<3> = LABELDEF
LABELPAT<2> = LABELDEF
LABELDEF =

NULABEND

*
* TEST FOR COLUMN SPECIFICATION CONTAINING ONLY DIGITS
*
TESTLABCOL LABELPAT<2> SPAN('0123456789') RPOS(0) :F(LABERROR)
SCANLABELDEF =
ROUTO LANLABPAT<2> 'COL'
LABELNAME = SUBSTR(LABELNAME,1,SIZE(LABELNAME) - 1)
OUTPUTLABCOM STRING = CRNULL(LABELNAME)

*
* ROUTINE FOR LABEL PROCESSING -- NO PARAMETERS
*
INSERT A NULL STATEMENT USING THE LABELNAME AND RETURN
ROUTO1  &ANCHOR = 1
INSERT(STRING)
: (RETURN)

*
* ROUTINE FOR UNIDENTIFIED-TO-END-OF-STATEMENT AND MACRO CALL
* PROCESSING -- NO PARAMETERS
*
Determine if this UEOS was a macro call
CURSTMT SPAN(' ') MACNAMEPAT : SELMAC '(',
BREAK(')') : PARMLIST ')',
CURSTMT MACNAMEPAT , SELMAC '(', BREAK(')') :F(COPYUEOS)

+
+
+
MACROCALL I = 0

*
* PROCESS A MACRO CALL
*

```

```

*****
***** BREAK OUT THE ACTUAL ARGUMENTS *****
*****
BREAKPARMS I = I + 1
      PARMLIST BREAK(' ',) CPARAM<I> ',' = :S(BREAKPARMS)
      PARMLIST REM . CPARAM<I> :F(MACCALLPERR)
      ACTARG = I
*****
***** RETRIEVE THE MACRO PATTERN AND NUMBER OF COMPONENTS INTO MPAT *****
*****
RETRIEVE THE MACRO PARAMETERS AND NUMBER OF PARAMETERS INTO MPMAM
*****
      MPAT = COPY(MACROTAB<SELMAC>)
      MPMAM = COPY(MACROTAB<SELMAC 'P'>)
      I = MPMAM<O>
*****
***** START THE PARAMETER PATTERN WITH A SERIES OF 81 BLANKS *****
*****
      PARM = DUPL(' ',81)
*****
***** ADD THE PARAMETERS TO THE PARAMETER PATTERN *****
*****
MACPARMPAT PARM = PARM | MPMAM<I>
      I = GT(I,1) I - 1
      MACSTR =
          :S(MACPARMPAT)
*****
***** COPY THE STATEMENT COMPONENTS AND REPLACE THE PARAMETERS *****
***** WITH THE ARGUMENTS *****
      I = 0
      STMTLEN = MPAT<O>
      LE(STMTLEN,O)
      LT(I,STMTLEN)
      I = I + 1
      MPAT<I> PARM . SELECTPARM SPAN(' ') RPOS(O)
          :S(ROUTOLERR)
          :F(MACDONE)
*****
***** MACLOOP *****
      J = 0
      MPAT<I> PARM . SELECTPARM RPOS(O)
          :S(COPYMACPARM)
          :F(COPYMAC)
*****
***** INCR.J *****
      J = I(T(J,ACTARG) J + 1
      SELECTPARM MPMAM<J>
      MACSTR = MACSTR CPARAM<J>
*****
***** COPYMACPARM *****
      J = 0
      J = LT(J,ACTARG) J + 1
      SELECTPARM MPMAM<J>
      LT(SIZE(MPAT<I>) - SIZE(CPARAM<J>),O)
      MACSTR = MACSTR CPARAM<J> DUPL(' ',SIZE(MPAT<I>) -
          SIZE(CPARAM<J>))
*****
***** INCR2.J *****

```



```

** **
**      INSERTCOM('ELSE KEYWORD PROCESSED.')
**
**      INSERT THE STRING
**
**      INSERT(STRING)
**
**      PUSH THE LABEL CREATED IN THIS ROUTINE ONTO THE SEMANTIC STACK
**
ROUT2
**      PUSH.SFM(1,VAR2)
**      &ANCHOR = 0
**      &ANCHOR = 1
**      : (RETURN)
**
**      ROUTINE FOR CASE STATEMENT PROCESSING
**
**      ERROR IF NOT IN CORRECT FORM FOR CASE
**
**      INSTR = CURSTMT
**      INSTR SPAN(' ') =
**
**      BREAK OUT THE INDEX OF THE CASE
**
**      INSTR 'CASE OF ' BREAK(SEMDEL) . IND = :F(CASEERROR)
**
**      CREATE 2 LABELS; CREATE A STRING OF AN IF STATEMENT; SET KTR TO 1
**
GEN2
**      VARO = UNILABEL()
**      VAR1 = UNILABEL()
**      INASSTRUCT = INASSTRUCT + 1
**      NO.OF.CASES = NO.OF.CASES + 1
**      $( 'KTR' NO.OF.CASES ) = 1
**      STRING = CRIF(UNILABEL(),IND,NEQU,$('KTR' NO.OF.CASES),VAR1)
**
**      INSERT THE CURSTMT AS A COMMENT; NULLIFY CURSTMT
**
**      INSERTCOM(CURSTMT)
**      CURSTMT =
**
**      INSERT THE CREATED STRING; PUSH THE INDEX OF THE CASE AND THE 2
**      CREATED LABELS ONTO THE SEMANTIC STACK
**
**      INSERT(STRING)
**      PUSH.SEM(3,IND,VAR1,VARO)
**      &ANCHOR = 0
**      &ANCHOR = 1
**      : (RETURN)
**
ROUT3

```

```

PRE24490
PRE24500
PRE24510
PRE24520
PRE24530
PRE24540
PRE24550
PRE24560
PRE24570
PRE24580
PRE24590
PRE24600
PRE24601
PRE24602
PRE24610
PRE24620
PRE24621
PRE24622
PRE24630
PRE24640
PRE24650
PRE24660
PRE24670
PRE24680
PRE24690
PRE24700
PRE24710
PRE24720
PRE24730
PRE24740
PRE24750
PRE24760
PRE24770
PRE24780
PRE24790
PRE24800
PRE24810
PRE24820
PRE24830
PRE24840
PRE24850
PRE24860
PRE24870
PRE24880
PRE24890
PRE24900
PRE24910
PRE24920

```

```

ROUTINE FOR DO LOOP STATEMENT PROCESSING
ALL ERRORS IF NOT IN CORRECT FORM FOR DO LOOP
    INSIR = CURSTMT
    INSIR SPAN( ' ' ) =
BREAK OUT THE INDEX, BEGINNING VALUE, ENDING VALUE, AND INCREMENT
VALUE OF THE DO LOOP
    INSIR 'DO' : BREAK( ' ' ) ; IND ' = ' = : F( DOERROR )
    INSIR 'ARB' : BEGIN ' TO ' = : F( DOERROR )
    INSIR 'ARB' : END ' BY ' = : S( INCR )
    INSIR BREAK( SEMDEL ) : END = : F( DOERROR )
    INC = 1 : ( GEN3 )
    INSIR BREAK( SEMDEL ) : INC = : F( DOERROR )
INCR
CREATE 1 LABEL AND 1 VARIABLE; CREATE A STRING OF 2 ASSIGN MINUS
STATEMENTS AND AN ASSIGN PLUS STATEMENT
    VAR1 = UNILABEL( )
    VAR2 = UNIVAR( )
    INASTRUCT = INASTRUCT + 1
    STRING = CRMINUS( UNILABEL( ), VAR2, END, INC )
            CRPLUS( VAR1, IND, IND, INC )
INSERT THE CURSTMT AS A COMMENT; NULLIFY THE CURSTMT;
INSERT THE CREATED STRING
    INSERTCCM( CURSTMT )
    CURSTMT =
    INSERT( STRING )
PUSH THE INCREMENT, INDEX, AND VARIABLE AND LABEL CREATED IN
THIS ROUTINE ONTO THE SEMANTIC STACK
    PUSH. SEM( 4, INC, IND, VAR2, VAR1 )
    &ANCHOR = 0
    &ANCHOR = 1
ROUT4

```

```

**      ROUTINE FOR DOGROUP STATEMENT PROCESSING
**      ***
**      ERROR IF NOT IN CORRECT FORM FOR DOGROUP
**      ***
**      INSIR = CURSMT
**      INSTR SPAN( ' ' ) =
**      INSTR *DOGROUP* ANY( SEMDEL ) =          :F(DOGRPEROR)
**
**      INSERT CURSMT AS A COMMENT; NULLIFY CURSMT; RETURN
**
**      GEN4
**      INSERTCOM(CURSMT)
**      INASTRUCT = INASTRUCT + 1
**      CURSMT = 0
**      EANCHOR = 1
**      : (RETURN)
**
**      ROUT5
**      ***
**      ROUTINE FOR IF THEN STATEMENT PROCESSING
**      ***
**      ERROR IF NOT IN CORRECT FORM FOR IF THEN
**      ***
**      INSIR = CURSMT
**      INSTR SPAN( ' ' ) =
**
**      BREAK OUT THE CONDITION
**
**      INSTR *IF * ARB . BOOL * THEN * =          :F(IFERROR)
**
**      CREATE A LABEL FOR BRANCHING TO ON FALSE; CREATE A STRING
**      OF AN IF STATEMENT
**
**      GEN5
**      VAR1 = UNILABEL( )
**      STRING = CRIF( UNILABEL( ), NULL, NOT, ' ( ' BOOL ' ) ', VAR1 )
**
**      INSERT CURSMT AS A COMMENT; NULLIFY CURSMT
**
**      INSERTCOM(CURSMT)
**      CURSMT =
**
**      INSERT THE STRING; PUSH THE LABEL CREATED IN THIS ROUTINE
**      ONTO THE SEMANTIC STACK; RETURN
**
**      INSERT(STRING)

```

```

PRE25350
PRE25360
PRE25361
PRE25362
PRE25370
PRE25380
PRE25390
PRE25400
PRE25410
PRE25420
PRE25430
PRE25440
PRE25450
PRE25460
PRE25470
PRE25480
PRE25490
PRE25500
PRE25501
PRE25502
PRE25510
PRE25520
PRE25521
PRE25522
PRE25530
PRE25540
PRE25550
PRE25560
PRE25570
PRE25580
PRE25590
PRE25600
PRE25610
PRE25620
PRE25630
PRE25640
PRE25650
PRE25660
PRE25670
PRE25680
PRE25690
PRE25700
PRE25710
PRE25720
PRE25730
PRE25740
PRE25750
PRE25760

```



```

ROUT6
**      PUSH.SEM(1,VAR1)
**      &ANCHOR = 0
**      &ANCHOR = 1
**      ***(RETURN)
**
**      ROUTINE FOR REPEAT STATEMENT PROCESSING
**      ***(RETURN)
**
**      ERROR IF NOT IN CORRECT FORM FOR REPEAT
**
**      INSTR = CURSTMT
**      INSTR SPAN(' ') =
**      INSTR REPEAT UNTIL (' =
**
**      BREAK OUT THE CONDITION
**
**      INSTR ARB • BOOL ')* ANY(SEMDEL) =
**
**      CREATE A LABEL FOR A NULL STATEMENT; CREATE A STRING OF A NULL
**      STATEMENT
**
**      VAR1 = UNILABEL()
**      INASTRUCT = INASTRUCT + 1
**      STRING = CRNULL(VAR1)
**
**      INSERT CURSTMT AS A COMMENT; NULLIFY CURSTMT
**
**      INSERTCOM(CURSTMT)
**      CURSTMT =
**
**      INSERT THE STRING; PUSH THE LABEL CREATED IN THIS ROUTINE AND THE
**      CONDITION ONTO THE SEMANTIC STACK; RETURN
**
**      INSERT(STRING)
**      PUSH.SEM(2,VAR1,BOOL)
**      &ANCHOR = 0
**      &ANCHOR = 1
**
**      ROUTINE FOR WHILE STATEMENT PROCESSING
**
**      ERROR IF NOT IN CORRECT FORM FOR WHILE

```

```

PRE25770
PRE25780
PRE25790
PRE25800
PRE25801
PRE25802
PRE25810
PRE25820
PRE25821
PRE25822
PRE25830
PRE25840
PRE25850
PRE25860
PRE25870
PRE25880
PRE25890
PRE25900
PRE25910
PRE25920
PRE25930
PRE25940
PRE25950
PRE25960
PRE25970
PRE25980
PRE25990
PRE26000
PRE26010
PRE26020
PRE26030
PRE26040
PRE26050
PRE26060
PRE26070
PRE26080
PRE26090
PRE26100
PRE26110
PRE26120
PRE26121
PRE26122
PRE26130
PRE26140
PRE26141
PRE26142
PRE26150
PRE26160

```



```

INSTR = CURSTMT
INSTR SPAN(' ') =
INSTR 'WHILE (' =
:F(WHERROR)

BRFAK OUT THE CONDITION

INSTR ARB • BOOL ')' ANY(SEMDEL) = :F(WHERROR)

CREATE 2 LABELS; CREATE A STRING OF AN IF STATEMENT

VAR1 = UNILABEL()
VAR2 = UNILABEL()
INASTRUCT = INASTRUCT + 1
STRING = CRIF(VAR1,NULL,NOT,'(• BOOL ')',VAR2)

INSERT CURSTMT AS A CCMMENT; NULLIFY CURSTMT

INSERTCOM(CURSTMT)
CURSTMT =

INSERT THE STRING; PUSH THE 2 LABELS CREATED IN THIS
ROUTINE ONTO THE SEMANTIC STACK; RETURN

:(RETURN)

ROUT8
ANCHOR = 1

ANCHOR = 0
PUSH.SEM(2,VAR2,VAR1)
INSERT(STRING)

ROUTINE FOR QUITLOOP STATEMENT PROCESSING

ERROR IF NOT IN CORRECT FOR FOR QUITLOOP

INSTR = CURSTMT
INSTR SPAN(' ') =
INSTR 'QUITLOOP' ANY(SEMDEL) = :F(QUITERROR)

CREATE A LABEL FOR BRANCHING OUT OF THE LOOP; CREATE A
STRING OF A GO TO STATEMENT

VAR1 = UNILABEL()
STRING = CRGOTO(UNILABEL(),VAR1)

INSERT CURSTMT AS A COMMENT; NULLIFY CURSTMT

```

```

**          INSERTCOM(CURSTMT)
**          CURSTMT =
**
**          INSERT THE STRING; PUSH THE LABEL CREATED IN THIS ROUTINE
**          AND 'QUIT' ONTO THE SEMANTIC STACK; RETURN
**
ROUT9
**          INSERT(STRING)
**          PUSH_SEM(2,VAR1,'QUIT')
**          &ANCHOR = 0
**          &ANCHOR = 1
**          : (RETURN)
**
**          ROUTINE FOR ENDCASE PROCESSING
**
**          ERROR IF NOT IN CORRECT FORM FOR ENDCASE
**
**          INSTR = CURSTMT
**          INSTR SPAN( , ) =
**          INSTR 'ENDCASE' ANY(SEMDEL) =
**          :F( ECASEERROR)
**
**          POP 2 LABELS AND THE INDEX OF THE CASE OFF THE SEMANTIC STACK
**
GEN9
**          POP_SEM(3,'VARO','VAR1','IND')
**
**          CREATE A STRING OF 2 NULL STATEMENTS; RESET KTR TO 0
**
**          STRING = CRNULL(VAR1) CRNULL(VARO)
**          &KTR, NO.OF.CASES) = 0
**          NO.OF.CASES = NO.OF.CASES - 1
**          INASTRUCT = INASTRUCT - 1
**
**          INSERT CURSTMT AS A COMMENT; NULLIFY CURSTMT
**
**          INSERTCOM(CURSTMT)
**          CURSTMT =
**
**          INSERT THE STRING; RETURN
**
**          INSERT(STRING)
**          &ANCHOR = 0
**          &ANCHOR = 1
**          : (RETURN)
**
ROUT10
**          ROUTINE FOR ENDIF PROCESSING

```

```
PRE27030
PRE27031
PRE27032
PRE27040
PRE27050
PRE27060
PRE27070
PRE27080
PRE27090
PRE27100
PRE27110
PRE27120
PRE27130
PRE27140
PRE27150
PRE27160
PRE27170
PRE27180
PRE27190
PRE27200
PRE27210
PRE27220
PRE27230
PRE27240
PRE27250
PRE27260
PRE27270
PRE27280
PRE27290
PRE27300
PRE27310
PRE27320
PRE27330
PRE27340
PRE27350
PRE27360
PRE27361
PRE27362
PRE27370
PRE27380
PRE27381
PRE27382
PRE27390
PRE27400
PRE27410
PRE27420
PRE27430
PRE27440

**      **
**      * ERROR IF NOT IN CORRECT FORM FOR ENDIF          **
**      **
INSTR = CURSTMT
INSTR SPAN(.,.) =
INSTR ,ENDIF, ANY(SEMDEL) = :F(EIFERROR)

POP AN ITEM OFF THE SEMANTIC STACK; BRANCH IF IT ISN'T 'QUIT'

GENIO                                :F(NOTAQUIT)
    POP.SEM(1,'VAR1')
    VAR1,'QUIT',RPOS(0)

THE ITEM POPPED WAS 'QUIT'. SO POP 2 MORE ITEMS OFF AND PUSH THE
'QUIT' AND FOLLOWING ITEM BACK ONTO THE STACK, LEAVING A
LABEL AVAILABLE

    POP.SEM(2,'VAR2','VAR1')
    PUSH.SEM(2,VAR2,'QUIT')

CREATE A STRING OF A NULL STATEMENT

NOTAQUIT   STRING = CRNULL(VAR1)

INSERT CURSTMT AS A COMMENT; NULLIFY CURSTMT

        INSERTCOM(CURSTMT)
CURSTMT =

INSERT THE STRING; RETURN

        INSTR(SSTRING)
        &ANCHOR = 0
        &ANCHOR = 1

ROUT11                                         :(RETURN)

ROUTINE FOR ENDWHILE PROCESSING

ERROR IF NOT IN CORRECT FORM FOR ENDWHILE

            INSTR = CURSTMT
            INSTR SPAN(.,.) =
            INSTR ,ENDWHILE, ANY(SEMDEL) = :F(EWHILEERROR)

**      **
```

```

* * *
* * *
* * * GEN11
* * *
* * * SET QUITFLAG TO 0; POP 2 ITEMS OFF THE SEMANTIC STACK;
* * *   BRANCH IF THE FIRST ITEM ISN'T 'QUIT'
* * *
* * *   QUITFLAG = 0
* * *   POP.SEM(2,'VAR1','VAR2')
* * *   INSTRUCT = INSTRUCT - 1
* * *   VAR1 'QUIT' RPOS(0)
* * *
* * *   :F(ELSE11)
* * *
* * *   THE FIRST ITEM WAS 'QUIT', SO ASSIGN THE SECOND ITEM TO QUITLAB
* * *   AND SET QUITFLAG TO 1
* * *
* * *   QUITLAB = VAR2
* * *   QUITFLAG = 1
* * *
* * *   POP 2 LABELS OFF THE SEMANTIC STACK
* * *
* * *   PCP.SEM(2,'VAR1','VAR2')
* * *
* * *   CREATE A STRING OF A GO TO STATEMENT AND A NULL STATEMENT
* * *   ELSE11   STRING = CRGOTO(UNILABEL(),VAR1) CRNULL(VAR2)
* * *
* * *   INSERT CURSTMT AS A COMMENT; NULLIFY CURSTMT
* * *
* * *   INSERTCGM(CURSTMT)
* * *   CURSTMT =
* * *
* * *   INSERT THE STRING; RETURN IF QUITFLAG IS 0
* * *
* * *   INSERT(STRING)
* * *   NE(QUITFLAG,1)
* * *   &ANCHOR = 0
* * *
* * *   :F(QUITINWHILE)
* * *   :(RETURN)
* * *
* * *   QUITFLAG IS 1, CREATE A STRING OF A NULL STATEMENT; INSERT
* * *   THE STRING; RETURN
* * *
* * *   QUITINWHILE   STRING = CRNULL(QUITLAB)
* * *   INSERT(STRING)
* * *   &ANCHOR = 0
* * *   &ANCHOR = 1
* * *
* * *   ROUT12
* * *
* * *   ROUTINE FOR ENDREPEAT PROCESSING
* * *
* * *   ERROR IF NOT IN CORRECT FORM FOR ENDREPEAT
* * *

```

```

PRE27450
PRE27460
PRE27470
PRE27480
PRE27490
PRE27500
PRE27510
PRE27520
PRE27530
PRE27540
PRE27550
PRE27560
PRE27570
PRE27580
PRE27590
PRE27600
PRE27610
PRE27620
PRE27630
PRE27640
PRE27650
PRE27660
PRE27670
PRE27680
PRE27690
PRE27700
PRE27710
PRE27720
PRE27730
PRE27740
PRE27750
PRE27760
PRE27770
PRE27780
PRE27790
PRE27800
PRE27810
PRE27820
PRE27830
PRE27840
PRE27850
PRE27851
PRE27852
PRE27860
PRE27870
PRE27871
PRE27872
PRE27880

```

```

*
*      INSIR = CURSTMT
*      INSIR SPAN(1,1) =
*      INSIR 'ENDREPEAT' ANY(SEMDEL) =          :F(EREPEXOR)
*
*      SET QUITFLAG TO 0; POP 2 ITEMS OFF THE SEMANTIC STACK; BRANCH
*      IF THE FIRST ITEM ISN'T 'QUIT'
*
*      GEN12
*          QUITFLAG = 0
*          POP.SEM(2,'BCOL','VAR1')
*          INASTRUCT = INASTRUCT - 1
*          BOOL 'QUIT' RPOS(0)
*          :F(ELSE12)
*
*      THE FIRST ITEM WAS 'QUIT', SO ASSIGN THE SECOND ITEM TO QUITLAB
*      AND SET QUITFLAG TO 1
*
*          QUITLAB = VAR1
*          QUITFLAG = 1
*
*      POP A CCNDITION AND A LABEL OFF THE SEMANTIC STACK; CREATE A
*      STRING OF AN IF STATEMENT
*
*      ELSE12
*          POP.SEM(2,'BOOL','VAR1')
*          STRING = CRIF(UNILABEL(),NULL,NOT,(' ' BOOL ' '),VAR1)
*
*      INERT CURSTMT AS A COMMENT; NULLIFY CURSTMT
*
*          INSERTCOM(CURSTMT)
*          CURSTMT =
*
*      INSERT THE STRING; RETURN IF QUITFLAG IS 0
*
*          INSERT(STRING)
*          NE(QUITFLAG,1)
*          &ANCHOR = 0
*          :F(QUITINREP)
*          : (RETURN)
*
*      QUITFLAG IS 1; CREATE A STRING OF A NULL STATEMENT; INSERT THE
*      STRING; RETURN
*
*      QUITINREP
*          STRING = CRNULL(QUITLAB)
*          INSERT(STRING)
*          &ANCHOR = 0
*          : (RETURN)
*
*      ROUT13
*          &ANCHOR = 1
*
*      ROUTINE FOR ENDDO PROCESSING

```

```

PRE27890
PRE27900
PRE27910
PRE27920
PRE27930
PRE27940
PRE27950
PRE27960
PRE27970
PRE27980
PRE27990
PRE28000
PRE28010
PRE28020
PRE28030
PRE28040
PRE28050
PRE28060
PRE28070
PRE28080
PRE28090
PRE28100
PRE28110
PRE28120
PRE28130
PRE28140
PRE28150
PRE28160
PRE28170
PRE28180
PRE28190
PRE28200
PRE28210
PRE28220
PRE28230
PRE28240
PRE28250
PRE28260
PRE28270
PRE28280
PRE28290
PRE28300
PRE28310
PRE28320
PRE28321
PRE28322
PRE28330
PRE28340

```



```

**
**
**      CREATE A LABEL; POP 3 ITEMS OFF THE SEMANTIC STACK
**
**      VAR3 = UNILABEL()
**      POP.SEM(3,VARO,'VAR1','IND')
**
**      CREATE A STRING OF A GO TO STATEMENT AND IF AN STATEMENT
**
**      STRING = CRGOTO(UNILABEL(),VARO)
**      CRIF(VAR1,IND,NEQ,{('KTR' NO.OF.CASES),VAR3})
**
**      INSERT A COMMENT ABOUT THIS ROUTINE'S RESULTS
**
**      INSERTCOM('NEXT TWO STATEMENTS INSERTED FOR CASE PROCESSING')
**
**      INSERT THE STRING; PUSH THE INDEX OF THE CASE AND THE 2
**      LABELS CREATED IN THIS ROUTINE ONTO THE SEMANTIC STACK;
**      RETURN
**
**      INSERT(S(STRING)
**      PUSH.SEM(3,IND,VAR3,VARO)
**      &ANCHOR = 0
**
**      INSERTCOM I = 1
**      : (RETURN)
**
**      ROUTINE TO INSERT A STRING AS A COMMENT
**
**      ONF PARAMETER -- CCMSTR, THE STRING TO BE MADE INTO A
**      COMMENT AND INSERTED
**
**      FIND THE FIRST AVAILABLE COLUMN FOR INSERTING THE COMMENT
**
**      IGNPAT<0> 'SEQ'
**      I = EQ(I,IGNPAT<1>) I + 1
**      COLSEQ GT(I,IGNPAT<2>)
**      LT(I,IGNPAT<1>)
**      INCR.I I = LE(I,IGNPAT<2>) I + 1
**
**      FILL THE COLUMNS PRECEDING THE COLUMN FOUND WITH BLANKS,
**      BRANCH BASED ON THE COMMENT DEFINITION
**
**      FOUNDCOL OPEN OUTCOM = DUPL(' ',I - 1)
**      COMPAT<1> '10'
**
**      A COMMENT BEGINS IN A CERTAIN COLUMN, FILL WITH BLANKS TO THAT

```



```

* *      COLUMN
* *      I = COMPAT<4> - I
* *      OUTCOM = OUTCOM DUPL(' ',I)
* *      CCMAT<1> '11'
* *      :F(ADDCOMSTR)
* *
* *      A COMMENT BEGINS WITH A CERTAIN STRING IN THAT PARTICULAR
* *      COLUMN, ADD THE STRING TO THE OUTPUT COMMENT STRING
* *
* *      ADDOPENSTR OUTCOM = OUTCOM COMPAT<3>
* *
* *      ADD THE PARAMETER STRING TO THE OUTPUT COMMENT STRING,
* *      CHECK FOR THE OUTPUT STRING BEING LONGER THAN 80 CHARACTERS
* *
* *      ADDCOMSTR OUTCOM = OUTCOM COMSTR
* *      GT(SIZE(OUTCOM),80)
* *      :F(OTRUNCATE)
* *
* *      TRUNCATE TO 80 CHARACTERS
* *
* *      OUTCOM = SUBSTR(OUTCOM,1,80)
* *      NOTRUNCATE CCMAT<2> '01'
* *      :F(ENDCOMSTR)
* *
* *      A COMMENT ENDS IN A CERTAIN COLUMN, ADD BLANKS IF THE COMMENT
* *      IS TOO SHORT, TRUNCATE IF IT IS TOO LONG, BRANCH TO
* *      THE COMMENT FINISHED LABEL
* *
* *      GT(SIZE(OUTCOM),COMMAT<6>)
* *      OUTCOM = OUTCOM DUPL(' ',80 - SIZE(OUTCOM)) :S(COLTRUNC)
* *      OUTCOM = SUBSTR(OUTCOM,1,CCMAT<6>) :S(COMFIN)
* *      ENDCOMSTR COMPAT<2> '11' :S(COMENDSTRCOL)
* *
* *      A COMMENT ENDS WITH A CERTAIN STRING, IF THE STRING WON'T FIT
* *      ONTO THE OUTPUT STRING AND STILL HAVE THE STRING BE
* *      80 CHARACTERS OR LESS, TRUNCATE THE CURRENT OUTPUT STRING
* *
* *      GT(SIZE(OUTCOM),80 - SIZE(COMMAT<5>)) :F(NOSYMTUNC)
* *      OUTCOM = SUBSTR(OUTCOM,1,80 - SIZE(COMMAT<5>))
* *
* *      ADD THE TERMINATING STRING, SET I TO THE COLUMN NUMBER THE
* *      STRING MUST START IN TO END EXACTLY IN COLUMN 80, BRANCH
* *      ACCORDING TO THE IGNORE COLUMN DEFINITION SPECIFYING
* *      A SINGLE COLUMN OR A SEQUENCE OF COLUMNS
* *
* *      NOSYMTUNC OUTCOM = OUTCOM COMPAT<5>
* *      I = 81 - SIZE(CCMAT<5>)
* *      IGNPAT<0> 'SEQ'
* *      :S(COLSEQEND)
* *
* *      DECREMENT I IF THE STRING WOULD END IN AN IGNORE COLUMN AS IT

```

```

PRE29650
PRE29660
PRE29670
PRE29680
PRE29690
PRE29700
PRE29710
PRE29720
PRE29730
PRE29740
PRE29750
PRE29760
PRE29770
PRE29780
PRE29790
PRE29800
PRE29810
PRE29820
PRE29830
PRE29840
PRE29850
PRE29860
PRE29870
PRE29880
PRE29890
PRE29900
PRE29910
PRE29920
PRE29930
PRE29940
PRE29950
PRE29960
PRE29970
PRE29980
PRE29990
PRE30000
PRE30010
PRE30020
PRE30030
PRE30040
PRE30050
PRE30060
PRE30070
PRE30080
PRE30090
PRE30100
PRE30110
PRE30120

```

```

**      IS AND BRANCH
**      I = EQ(I,IGNPAT<1>) I - SIZE(CMPAT<5>)      :S(FOUNDCCOLCLOSE)
**
**      DECREMENT I UNTIL THE STRING IS SURE TO END IN A NON-IGNORED
**      COLUMN
**
**      COLSEQEND      GT(I,IGNPAT<2>)
**      LT(I + SIZE(CMPAT<5>) - 1,IGNPAT<1>)
**      I = GE(I + SIZE(CMPAT<5>) - 1,IGNPAT<1>) I - 1      :S(FOUNDCCOLCLOSE)
**      :S(FOUNDCCOLCLOSE)
**      DECR.I
**      +
**      ADD THE COMMENT TERMINATING STRING TO THE OUTPUT STRING AND
**      ADD BLANKS AS NECESSARY TO FILL 80 COLUMNS
**
**      FOUNDCCOLCLOSE OUTPUT = SUBSTR(OUTPUT,I,I - 1) CCMPAT<5>
**      DUPL(' ',81 - I - SIZE(CMPAT<5>))      :S(COMFIN)
**      COMENDSTRCOL GT(SIZE(OUTPUT),CMPAT<6>)      :S(CULSYMTRUNC)
**      OUTPUT = OUTPUT DUPL(' ',CMPAT<6> - 1 - SIZE(OUTPUT))
**      :S(ADDCOLSYM)
**      +
**      COLSYMTRUNC OUTPUT = SUBSTR(OUTPUT,I,CMPAT<6> - 1)
**      ADDCOLSYM OUTPUT = OUTPUT CCMPAT<5>
**      DUPL(' ',80 - SIZE(OUTPUT) - SIZE(CMPAT<5>))
**      +
**      PRINT AND PUNCH THE COMMENT AND RETURN
**
**      COMFIN      OUTPUT = OUTPUT
**      PUNCH = OUTPUT
**      INSEPTCOM =
**      INSERTSTMT OUTSTR =
**      +
**      ROUTINE TO INSERT A STRING AS A STATEMENT -- IGNORED COLUMNS
**      MUST BE SKIPPED
**
**      UNF PARAMETER -- STR.INSRT, THE STRING TO BE INSERTED
**      AS A STATEMENT
**      +
**      +
**      INSERTSTMT =
**      I = 0
**      CHARKTR = 0
**
**      SET TYPE APPROPRIATELY BASED ON A SINGLE COLUMN OR A
**      SEQUENCE OF COLUMNS TO BE IGNORED

```

```

PRE30130
PRE30140
PRE30150
PRE30160
PRE30170
PRE30180
PRE30190
PRE30200
PRE30210
PRE30220
PRE30230
PRE30240
PRE30250
PRE30260
PRE30270
PRE30280
PRE30290
PRE30300
PRE30310
PRE30320
PRE30330
PRE30340
PRE30350
PRE30360
PRE30370
PRE30380
PRE30390
PRE30400
PRE30410
PRE30420
PRE30430
PRE30431
PRE30432
PRE30440
PRE30450
PRE30460
PRE30470
PRE30480
PRE30490
PRE30491
PRE30492
PRE30500
PRE30510
PRE30520
PRE30530
PRE30540
PRE30550
PRE30560

```

```

SINGLFIGN      IGNPAT<0> 'SEQ'      :F(SINGLEIGN)
                TYPE = 'SEQ'        ::(INCREMENT)
                TYPE = 'ONE'

*      CHARKTR COUNTS THE POSITION IN THE PARAMETER STRING
*
* INCREMENT     CHARKTR = CHARKTR + 1
*
*      I COUNTS THE POSITION IN THE OUTPUT STRING
*
* INCREMENT2    I = I + 1
                LF(CHARKTR,SIZE(STR.INSRT))
                LE(I,80)      :F(ENDSTMTOUT)
                                :F(ENDCARD)

*      BRANCH TO CHECK FOR IGNORING COLUMNS
*
*      CHECK FOR IGNORING THIS COLUMN WHEN ONLY ONE COLUMN IS IGNORED
* CKIGNONE      NE(I,IGNPAT<1>)      :S(GOODCOL)F(IGNOREIT)
*
*      CHECK FOR IGNORING THIS COLUMN WHEN A SEQUENCE OF COLUMNS
*      IS IGNORED
* CKIGNSEQ      GT(I,IGNPAT<2>)
                LT(I,IGNPAT<1>)      :S(GOODCOL)
                                :S(GOODCOL)

*      ADD A BLANK TO THE OUTPUT STRING IN AN IGNORED COLUMN AND
*      BRANCH TO PROCESS THE NEXT CHARACTER
*
* IGNOREIT      OUTSTR = OUTSTR . ' '      ::(INCREMENT2)
*
*      ADD THE CURRENT CHARACTER TO THE OUTPUT STRING AND BRANCH
*      TO PROCESS THE NEXT CHARACTER
*
* GOODCOL      OUTSTR = OUTSTR SUBSTR(STR.INSRT,CHARKTR,1) ::(INCREMENT)
*
*      OUTPUT CARD IS FULL, PUNCH AND PRINT IT, REINITIALIZE
*      COLUMN COUNTER AND OUTPUT STRING
*
* ENDCARD      OUTPUT = OUTSTR
                PUNCH = OUTSTR
                I = 1
                OUTSTR =
                                ::(('CKIGN' TYPE))

*      END OF PARAMETER STRING, ADD ENOUGH BLANKS TO FILL
*      A CARD, PUNCH AND PRINT IT, AND RETURN
*

```

```

PRE30570
PRE30580
PRE30590
PRE30600
PRE30610
PRE30620
PRE30630
PRE30640
PRE30650
PRE30660
PRE30670
PRE30680
PRE30690
PRE30700
PRE30710
PRE30720
PRE30730
PRE30740
PRE30750
PRE30760
PRE30770
PRE30780
PRE30790
PRE30800
PRE30810
PRE30820
PRE30830
PRE30840
PRE30850
PRE30860
PRE30870
PRE30880
PRE30890
PRE30900
PRE30910
PRE30920
PRE30930
PRE30940
PRE30950
PRE30960
PRE30970
PRE30980
PRE30990
PRE31000
PRE31010
PRE31020
PRE31030
PRE31040

```

```

* ENDSTMTOUT OUTSTR = OUTSTR DUPL(' ',80 - SIZE(OUTSTR))
  OUTPUT = OUTSTR
  PUNCH = OUTSTR
  PUNCH = INSRT.STR
  *
  *
  * ROUTINE TO INSERT A STRING WHICH IS ALREADY IN THE FORM OF
  * A STATEMENT
  *
  * ONE PARAMETER -- INSRT.STR, THE STATEMENT TO BE INSERTED
  *
  * PRINT AND PUNCH THE STATEMENT AND RETURN
  *
  * MOREOUT GT(SIZE(INSRT.STR),80)
  * OUTPUT = INSRT.STR DUPL(' ',80 - SIZE(INSRT.STR))
  *          :S(OUT80)
  *          :(:RETURN)
  *
  * OUT80 OUTPUT = SUBSTR(INSRT.STR,1,80)
  * INSRT.STR = SUBSTR(INSRT.STR,81,SIZE(INSRT.STR) - 80)
  *          :(:MOREOUT)
  *
  * CRIF CRIF =
  *
  * ROUTINE TO CREATE AN IF STATEMENT
  *
  * FIVE PARAMETERS -- LAB, VAL1, RL, VAL2, AND DES
  *
  * PARMNAME = 'LAB' | 'VAL1' | 'RL' | 'VAL2' | 'DES'
  * I = 0
  * STMTLEN = IFPAT<0>
  *
  * ERROR IF NUMBER OF STATEMENT COMPONENTS IS <= 0
  *
  * CRIFLOOP LE(STMTLEN,0)
  * LT(I,STMTLEN)
  * I = I + 1
  * IFPAT<I> PARMNAME . SELECTPAM SPAN(' ') RPOS(0)
  * IFPAT<I> PARMNAME . SELECTPAM RPOS(0)
  * CRIF = CRIF $SELECTPAM
  *          :S(CRIFERR)
  *          :F(IFRETSEQ)
  *          :S(COPYIFPAM)
  *          :F(COPYCRIF)
  *          :(:CRIFLOOP)
  *
  * ERROR IF NOT ENOUGH BLANKS LEFT AFTER PARAMETER NAME IN PROTOTYPE

```

```

COPYIFPARM LT(SIZE(IFPAT<I>) - SIZE($SELECTPARM),0) :S(IFBLKERR)
COPY PARAMETER AND ENOUGH BLANKS TO EQUAL FIELD LENGTH IN PROTOTYPE
CRIF = CRIF $SELECTPARM DUPL(' ',SIZE(IFPAT<I>) -
SIZE($SELECTPARM)) :S(CRIFLOOP)
CRIF = CRIF IFPAT<I> :S(CRIFLOOP)
ADD BLANKS TO MAKE 80 COLUMNS, ERROR IF MORE THAN 80 COLUMNS
ALREADY
IFCRDCL CRIF = CRIF DUPL(' ',80 - SIZE(CRIF)) :S(RETURN)F(CRIFERR)
CRDCL =
ROUTINE TO CREATE DECLARATION STATEMENT
ONE PARAMETER -- VARNAME
I = 0
STMTLEN = DCLPAT<0>
ERROR IF NUMBER OF STATEMENT COMPONENTS IS <= 0
LE(STMTLEN,0)
LT(I,STMTLEN)
I = I + 1
DCLPAT<I> = 'VARNAME' SPAN(' ', RPOS(0))
DCLPAT<I> = 'VARNAME' RPOS(0)
CRDCL = CRDCL VARNAME
: S(CRDCLERR)
: F(DCLRETSEQ)
CRDCLLOOP
COPYDCLPARM LT(SIZE(DCLPAT<I>) - SIZE(VARNAME),0) :S(DCLBLKERR)
COPY PARAMETER AND ENOUGH BLANKS TO EQUAL FIELD LENGTH IN PROTOTYPE
CPDCL = CRDCL VARNAME DUPL(' ',SIZE(DCLPAT<I>) -
SIZE(VARNAME)) :S(CRDCLLOOP)
CRDCL = CRDCL DCLPAT<I> :S(CRDCLLOOP)
ADD BLANKS TO MAKE 80 COLUMNS, ERROR IF MORE THAN 80 COLUMNS
ALREADY

```

```

+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
+ CRMINUS      CRDCL = CRDCL DUPL(' ',80 - SIZE(CRDCL))          :S(RETURN)F(CRDCLERR)
+ CRMINUS      CRMINUS =
ROUTINE TO CREATE AN ASSIGN MINUS STATEMENT
FOUR PARAMETERS -- LAB, VAR, VAL1, AND VAL2
PARAMNAME = 'LAB' | 'VAR' | 'VAL1' | 'VAL2'
I = 0
STMTLEN = MINUSPAT<O>
ERROR IF NUMBER OF STATEMENT COMPONENTS IS <= 0
LE(STMTLEN,O)
CRMINUSLOOP LT(I,STMTLEN)
I = I + 1
MINUSPAT<I> PARAMNAME . SELECTPAM SPAN(' ') RPOS(O)
MINUSPAT<I> PARAMNAME . SELECTPAM RPCS(O)
CRMINUS = CRMINUS $SELECTPAM
:S(CRMINERR)
:F(MINRETSEQ)
ERROR IF NOT ENOUGH BLANKS LEFT AFTER PARAMETER NAME IN PROTOTYPE
COPYMINPAM LI(SIZE(MINUSPAT<I>) - SIZE($SELECTPAM),O) :S(MINBLKERR)
COPY PARAMETER AND ENOUGH BLANKS TO EQUAL FIELD LENGTH IN PROTOTYPE
CRMINUS = CRMINUS $SELECTPAM DUPL(' ',SIZE(MINUSPAT<I>) -
SIZE($SELECTPAM))
:(CRMINUSLOOP)
:(CRMINUSLOOP)
ADD BLANKS TO MAKE 80 COLUMNS, ERROR IF MORE THAN 80 COLUMNS
ALREADY
MINRETSEQ CRMINUS = CRMINUS DUPL(' ',80 - SIZE(CRMINUS))
CRPLUS      CRPLUS =
ROUTINE TO CREATE AN ASSIGN PLUS STATEMENT
FOUR PARAMETERS -- LAB, VAR, VAL1, AND VAL2

```



```

CRNULLLOOP LT(I,STMTLEN)
I = I + 1
NULLPAT<I> 'LAB' SPAN(' ') RPOS(0)
NULLPAT<I> 'LAB' RPOS(0)
CRNULL = CRNULL LAB
**
ERROR IF NOT ENOUGH BLANKS LEFT AFTER PARAMETER NAME IN PROTOTYPE
**
COPYNULLPAM LT(SIZE(NULLPAT<I>) - SIZE(LAB),0)
**
COPY PARAMETER AND ENOUGH BLANKS TO EQUAL FIELD LENGTH IN PROTOTYPE
**
CRNULL = CRNULL LAB DUPL(' ',SIZE(NULLPAT<I>) - SIZE(LAB))
+ : (CRNULLLOO)
**
COPYCRNULL CRNULL = CRNULL NULLPAT<I>
**
ADD BLANKS TO MAKE 80 COLUMNS, ERROR IF MORE THAN 80 COLUMNS
ALREADY
**
NULLRETSEQ CRNULL = CRNULL DUPL(' ',80 - SIZE(CRNULL))
+ : S(RETURN)F(CRNULLERR)
**
CRGOTO =
**
ROUTINE TO CREATE A GO TO STATEMENT
**
TWO PARAMETERS -- LAB AND DES
**
PARAMNAME = 'LAB' | 'DES'
I = 0
STMTLEN = GOTOPAT<0>
**
ERROR IF NUMBER OF STATEMENT COMPONENTS IS <= 0
**
LE(STMTLEN,0)
LT(I,STMTLEN)
I = I + 1
GOTOPAT<I> PARAMNAME . SELECTPAM SPAN(' ') RPOS(0)
+ : S(COPYGOTOPARM)
GOTOPAT<I> PARAMNAME . SELECTPAM RPOS(0)
+ : F(COPYCRGOTO)
CRGOTO = CRGOTO $SELECTPAM
+ : (CRGUTOLOO)
**
ERROR IF NOT ENOUGH BLANKS LEFT AFTER PARAMETER NAME IN PROTOTYPE
**
COPYGOTOPAM LT(SIZE(GOTOPAT<I>) - SIZE($SELECTPAM),0) : S(GOTOBLKERR)
**

```



```

**      COPY PARAMETER AND ENOUGH BLANKS TO EQUAL FIELD LENGTH IN PROTOTYPE
**
**      CRGOTO = CRGOTO $SELECTPARM DUPL(' ', SIZE(GOTOPAT<I>)) -
**      SIZE($SELECTPARM))
**      COPYCRGOTO CRGOTO = CRGOTO GOTOPAT<I>
**
**      ADD BLANKS TO MAKE 80 COLUMNS, ERROR IF MORE THAN 80 COLUMNS
**      ALREADY.
**
**      GOTORETSEQ CRGOTO = CRGOTO DUPL(' ', 80 - SIZE(CRGOTO))
**      UNILABEL UNILABEL = LABELPAT<I>
**
**      ROUTINE TO CREATE A UNIQUE LABEL
**
**      NO PARAMETERS
**
**      COMPUTE NUMBER OF ZEROS TO IMBED AND ADD THEM TO THE LABEL'S
**      UNIFORM BEGINNING
**
**      I = LABELPAT<2> - SIZE(LABELNUM)
**      LT(I,0)
**      ADDLABZEROS UNILABEL = GT(I,0) UNILABEL '0'
**      I = I - 1
**
**      ADD THE SIGNIFICANT DIGITS AND THE LABEL ENDING
**
**      ENDLABZEROS UNILABEL = UNILABEL LABELNUM LABELPAT<3>
**
**      INCREMENT THE NUMBR USED TO MAKE THE LABEL UNIQUE
**
**      LABELNUM = LABELNUM + 1
**      UNILABEL UNILABEL = UNILABEL UNILABEL
**
**      ROUTINE TO CREATE A UNIQUE VARIABLE
**
**      NO PARAMETERS
**
**      COMPUTE NUMBER OF ZEROS TO IMBED AND ADD THEM TO THE VARIABLE'S

```

```

PRE33170
PRE33180
PRE33190
PRE33200
PRE33210
PRE33220
PRE33230
PRE33240
PRE33250
PRE33260
PRE33270
PRE33280
PRE33290
PRE33300
PRE33301
PRE33302
PRE33310
PRE33320
PRE33330
PRE33340
PRE33341
PRE33342
PRE33350
PRE33360
PRE33370
PRE33380
PRE33390
PRE33400
PRE33410
PRE33420
PRE33430
PRE33440
PRE33450
PRE33460
PRE33470
PRE33480
PRE33490
PRE33500
PRE33510
PRE33511
PRE33512
PRE33520
PRE33530
PRE33540
PRE33550
PRE33551
PRE33552
PRE33560

```

```

** UNIFORM BEGINNING
**
**      I = VARPAT<2> - SIZE(VARNUM)
**      LT(I,0)
**      ADDVARZEROS UNIVAR = GT(I,0) UNIVAR '0'
**      I = I - 1
**
**      ADD THE SIGNIFICANT DIGITS AND THE VARIABLE ENDING
**      ENDVARZEROS UNIVAR = UNIVAR VARNUM VARPAT<3>
**
**      ADD THE NEW VARIABLE TO THE LIST FOR LATER DECLARATION
**      AND INCREMENT THE INDEX INTO THAT LIST AND THE NUMBER
**      USED TO MAKE THE VARIABLE UNIQUE
**
**      VARLIST<NEXTIVAR> = UNIVAR
**      NEXTVAR = NEXTVAR + 1
**      VARNUM = VARNUM + 1
**
**      : (RETURN)
**
** PUP.SFM
**
**      ROUTINE TO PUP ONE TO SIX ITEMS FROM THE TOP OF THE
**      SEMANTIC STACK
**
**      TWO TO SEVEN PARAMETERS -- NOARGS, THE NUMBER OF ITEMS TO
**      BE POPPED; A1 - A6, THE NAMES OF THE VARIABLES (IN STRING
**      FORM) INTO WHICH THE POPPED VALUES ARE TO BE PLACED
**
**      I = 1
**
**      POP A VALUE FROM THE STACK AND ASSIGN IT TO A PARAMETER
**
**      POP      1($('A' I)) = SEMSTACK<TOP>
**
**      DECREMENT THE POINTER TO THE TOP OF THE STACK
**
**      TOP = TOP - 1
**
**      IF NOT DONE POPPING, BRANCH TO POP THE NEXT VALUE
**
**      I = LT(I,NOARGS) I + 1
**
**      RETURN WITH A NULL STRING VALUE
**
**      POP.SEM =
**
**      : (RETURN)
**
PRE33570
PRE33580
PRE33590
PRE33600
PRE33610
PRE33620
PRE33630
PRE33640
PRE33650
PRE33660
PRE33670
PRE33680
PRE33690
PRE33700
PRE33710
PRE33720
PRE33730
PRE33740
PRE33750
PRE33760
PRE33770
PRE33780
PRE33790
PRE33800
PRE33810
PRE33820
PRE33830
PRE33840
PRE33841
PRE33842
PRE33850
PRE33860
PRE33870
PRE33880
PRE33890
PRE33900
PRE33910
PRE33920
PRE33930
PRE33940
PRE33950
PRE33960
PRE33970
PRE33980
PRE33990
PRE34000
PRE34010

```

```

PUSH.SEM
*****
ROUTINE TO PUSH ONE TO SIX ITEMS ON THE TOP OF THE
SEMANTIC STACK
*****
TWO TO SEVEN PARAMETERS -- NOARGS, THE NUMBER OF ITEMS TO BE
PUSHED; A1 - A6, THE VALUES TO BE PUSHED ONTO THE STACK
*****
I = 1
*****
INCREMENT THE POINTER TO THE TOP OF THE STACK
*****
TCP = TOP + 1
*****
PUSH A VALUE ONTO THE STACK
*****
SEMSTACK<TOP> = $('A' I)
*****
IF NOT DONE PUSHING, BRANCH TO PUSH THE NEXT VALUE
*****
I = LT(I,NOARGS) I + 1
*****
RETURN WITH A NULL STRING VALUE
*****
PUSH.SEM =
*****
END

```

A STRUCTURED PROGRAMMING PREPROCESSOR
FOR A VARIETY OF BASE LANGUAGES

by

Mary L. Love

B.S., Kansas State University, 1975

AN ABSTRACT OF A MASTER'S REPORT

submitted in partial fulfillment of the
requirements for the degree
MASTER OF SCIENCE

Department of Computer Science

KANSAS STATE UNIVERSITY
Manhattan, Kansas

1977

ABSTRACT

This report describes the design and implementation of a preprocessor which is intended to allow and encourage the use of structured programming techniques in a variety of base languages. Eight structured programming control structures are implemented in the preprocessor which currently works with the PL/I, SNOBOL, and FORTRAN programming languages. The extended constructs are translated by the preprocessor into statements which are standard in the base language. The input provided by the user includes information about the base language syntax, optional macro definitions, and a program written in the extended version of the language. The extended version includes the features of the base language, user-defined macros, and the built-in structured programming control structures. The output produced by the preprocessor consists of a printed listing and a temporary disk file. The disk file may be used as input to the usual translator for the base language. The macro facility is a simple text replacement macro processor, included because it is not possible for one person to think of all the control structures which may be desirable at all times. The facility is crude and allows no semantic stack manipulation. A set of guidelines is provided for use in writing macro definitions to reduce the likelihood of writing macros which conflict with the principles of structured programming and, therefore, undermine the purpose of the preprocessor.