

OPTIMIZATION OF INDUSTRIAL SYSTEMS WITH
THE SEPARABLE PROGRAMMING AND
THE GENERALIZED REDUCED GRADIENT METHODS

by

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CHAPTER ONE
INTRODUCTION

The constrained mathematical optimization problem has had an important role in the analysis of complex systems. Use of linear programming is common in the management decision process. Its significant success leads to the development of more sophisticated nonlinear models that attempt to describe in more detail the intricacies of the system.

In the past two decades the barrier to nonlinear programming research was the absence of computer coded algorithms. Today, modern technology has blessed (some may prefer to say burdened) us with the necessary computer hardware and software for coding and executing in reasonable time the procedures needed to attack the nonlinear programming optimization problem. With reasonably simple computer languages such as FORTRAN, ALGOL, and PL-1 the researcher, or systems analyst has the ability to express his specialized algorithm in computer code. Ideally then, most of the time of the problem solving process can be devoted to the problem analysis and not to its computational aspects.

In the case of linear programming the application of the computer for computation is a simple process when compared to the modeling process. The systems analyst is allowed to devote practically all of his effort into the problem definition and analysis. More time is available for redefining the assumptions and measuring the resultant reaction of the system. For this reason linear models are often developed in spite of nonlinearities.

In contrast, the development of nonlinear programming has been hampered by theoretical difficulties. The problems of nonconvex sets and local optima are encountered and prevalent. Consequently, the coding of existing algorithms has been limited to special cases for which the behavior of the functions are predictable. Indeed, the algorithm superior to all algorithms is probably nonexistent.

The point is that in a study involving constrained nonlinear optimization

the modeling process is often secondary to the computational aspects of the problem. Efficient algorithms are either not available in computer code or are unknown. A solution may then be achieved by some obscure, problem oriented, possibly inefficient technique that is known only to a few experts in the field. This makes it difficult for others to verify or extend work. It is this difficulty that obscures academic progress, hinders project evaluation, and may cause relevant and important studies to be ignored. For this reason it is important to realize existing and efficient nonlinear programming codes of algorithms that are well based theoretically. When a particular known code can be applied to a problem it should be used in lieu of attempting to develop a new technique.

It is the purpose of this thesis to study in depth two existing and apparently efficient computer coded nonlinear programming algorithms. The rational is that this is another small step in the collection and presentation of information that is vital to the individual involved in constrained nonlinear optimization.

The separable programming method is a special purpose procedure in nonlinear optimization. It is an extension of the simplex method of linear programming. This is important because of the overwhelming success of coded simplex procedures. The separable programming method is, for this reason, very reliable.

One linear programming code that contains the separable programming method as an option is the Mathematical Programming System/360 (MPS/360) [37]. The usage of the code along with the theoretical development of separable programming is discussed in Chapter Two. The method requires much arithmetic in the formulation of the problem in the format acceptable by MPS/360. This necessitates the development of supplementary programs to generate the separable

programming data. For this reason a FORTRAN program is presented that aids in the usage of MPS/360. Several examples are worked including a very interesting extension. The application of separable programming to the geometric programming dual problem with N degrees of difficulty is discussed. This may be important to the nonlinear programming layman who has occasional experiences with nonlinear problems.

One important aspect of a simplex related nonlinear method is the potential for sensitivity analysis. Of course, MPS/360 contains any procedure desired by the user for such analysis and they can be applied to the separable programming option. It is this aspect coupled with the reliability of the separable programming method that warrants its presentation in Chapter Two.

The generalized reduced gradient (GRG) method is an elaborate extension of the hill-climbing gradient techniques. It is, therefore, limited to applications containing continuous, differentiable functions. The method is designed to optimize a variety of nonlinear programming problems and is, as suspected, very complicated.

The method has been coded in FORTRAN in a program named GREG[30]. The code is available but is still considered experimental in spite of its excellent performance. Because both the technique and code are new, the remaining chapters (three thru six) of this thesis are devoted to discussions on the CRG theory and on the application of the GREG code. The theory is presented in Chapter Three along with an illustrative numerical example. This is followed in Chapter Four by introducing the usage of the GREG code. Applications are discussed in Chapter Five in the form of numerical comparisons with previously worked problems from the literature. The finally is the application of the code to solve a multistage water quality management problem. In all cases the

superiority of the GRG method is apparent even with the minor shortcomings of the GREG code which are accordingly pointed out. The GRG method promises to remain an excellent nonlinear programming tool.

CHAPTER TWO
SEPARABLE PROGRAMMING

2.1 INTRODUCTION

Separable programming is a special class of nonlinear programming that is adaptable to linear programming. The problems are constructed of separable functions which have the form

$$\phi(\bar{x}) = \sum_{i=1}^m h_i(x_i) \quad (1)$$

The separable programming problem can be defined as finding a set of x_i , $i=1, 2, \dots, m$ which maximizes (or minimizes)

$$c(\bar{x}) = \sum_{i=1}^m f_i(x_i) \quad (2)$$

subject to the constraints

$$\sum_{i=1}^m g_{ki}(x_i) \leq b_k, \quad k=1, \dots, p \quad (3)$$

and

$$x_i \geq 0$$

By approximating a nonlinear function of one variable by a piecewise linear function, the problem becomes a restricted linear programming problem, and can be solved by a slightly revised simplex method.

This chapter reviews the separable programming theory used in the Mathematical Programming System /360 (MPS/360) [37]. It shall be apparent that the arithmetic involved in setting up a separable programming problem for the revised simplex method becomes unbearably cumbersome. For this reason a FORTRAN program is presented that does all of the necessary calculations and produces input data for MPS/360. The program is logically simple and is exemplified.

2.2 AN APPROXIMATION OF SEPARABLE FUNCTIONS

A continuous nonlinear function of a single variable, x_i , can be approximated by a piecewise linear function over a specified interval domain. This is done by partitioning this interval domain into n_i disjoint, but contiguous, intervals. The $(n_i + 1)$ points of the partitions are represented by the set

$$S = \{ x_i^0, x_i^1, x_i^2, \dots, x_i^{n_i} \}$$

There are two methods of representing the piecewise linear approximation of a continuous nonlinear function of one variable. The method employed in this chapter is known as the "delta method." Both methods are developed in G. Hadley's Nonlinear and Dynamic Programming [31]. The "delta method" uses the differences of adjacent points of the set, S , and the differences of the functional values at the adjacent points in developing the approximating equation of a function, $f_i(x_i)$. The differences are represented by

$$\begin{aligned} \Delta x_i^j &= x_i^j - x_i^{j-1}, \\ &\quad i = 1, 2, \dots, m \\ &\quad j = 1, 2, \dots, n_i \end{aligned} \tag{4}$$

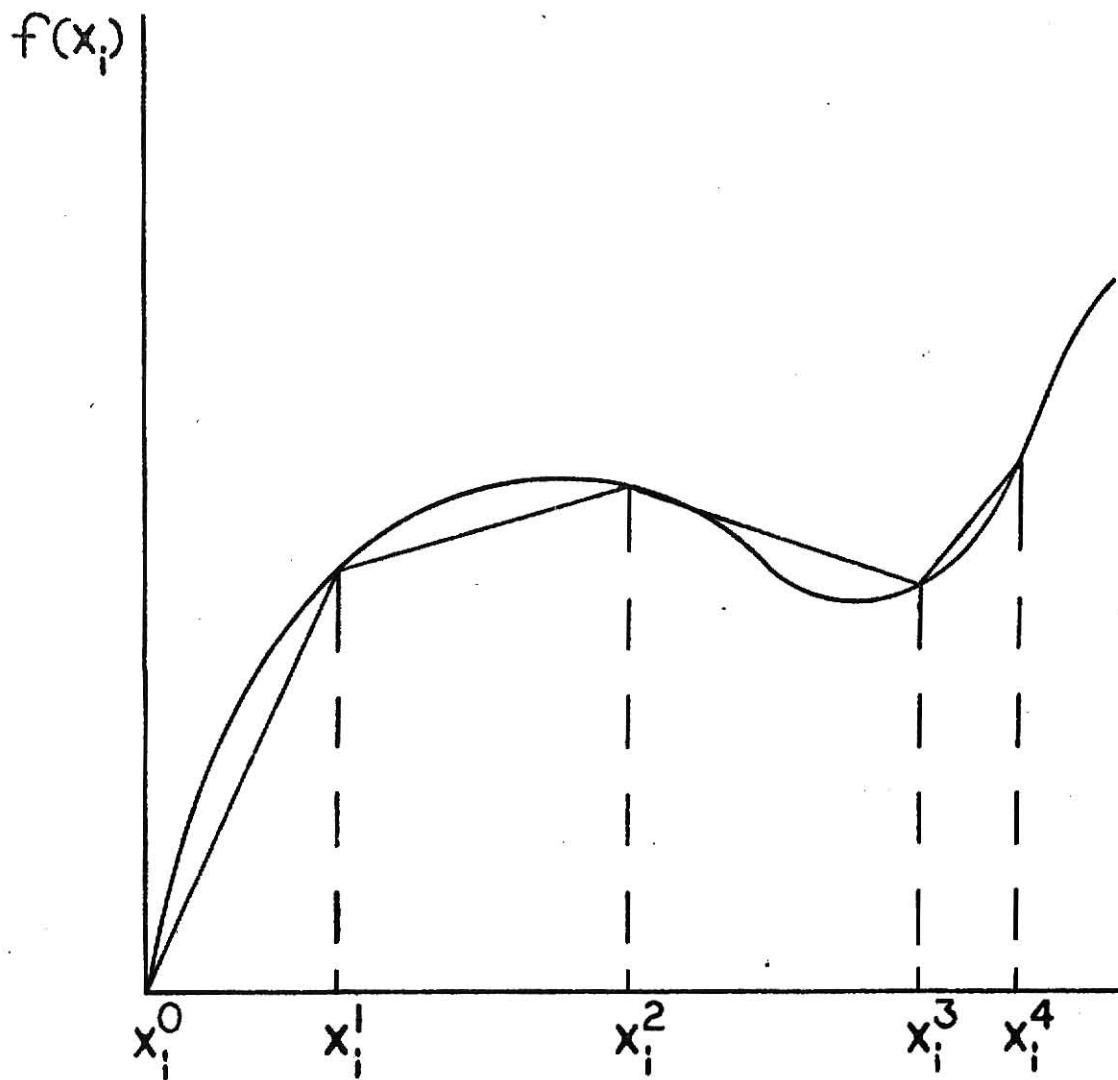
$$\Delta f_i^j = f_i(x_i^j) - f_i(x_i^{j-1}),$$

where the subscript refers to a function and/or variable such as x_i , $f_i(x_i)$, and $g_{ki}(x_i)$, and the superscript refers to a partitioning of a variable. That is, $f_i(x_i^j)$ is the value of $f_i(x_i)$ at $x_i = x_i^j$. The differences for adjacent points and the corresponding functional values for a function with $n_i = 4$ are shown in Fig. 2.1.

To represent the variable x_i and the approximation of $f_i(x_i)$, a set of variables, D_i^j , $j=1, 2, \dots, n_i$, is created that follows what is known as the "restricted-basis-entry rule." The rule is satisfied for any one of the following conditions.

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$$\Delta x_i^1 = x_i^1 - x_i^0$$

$$\Delta f_i^1 = f_i(x_i^1) - f_i(x_i^0)$$

$$\Delta x_i^2 = x_i^2 - x_i^1$$

$$\Delta f_i^2 = f_i(x_i^2) - f_i(x_i^1)$$

$$\Delta x_i^3 = x_i^3 - x_i^2$$

$$\Delta f_i^3 = f_i(x_i^3) - f_i(x_i^2)$$

$$\Delta x_i^4 = x_i^4 - x_i^3$$

$$\Delta f_i^4 = f_i(x_i^4) - f_i(x_i^3)$$

Fig. 2.1. Linear approximation of $f_i(x_i)$.

$$(i) 0 \leq D_i^1 \leq 1 \text{ iff}$$

$$D_i^j = 0, j = 2, 3, \dots, n_i$$

$$(ii) 0 \leq D_i^j \leq 1 \text{ iff}$$

$$D_i^\ell = 1, \ell = 1, 2, \dots, j-1$$

and

$$D_i^k = 0, k = j+1, \dots, n_i$$

$$(iii) 0 \leq D_i^{n_i} \text{ iff}$$

$$D_i^j = 1, j = 1, 2, \dots, n_i - 1$$

where n_i is the number of partitioning intervals for a variable x_i . D_i^j represents a special variable created for the partition j of variable x_i .

Intuitively, for any $0 \leq D_i^j \leq 1$ all previous D_i^ℓ variables ($\ell = 1, \dots, j-1$) must have a value of one and all following values ($\ell = j+1, \dots, n_i$) must be zero.

Example 1

Let $n_i = 4$. Cases a, b, c, and d satisfy the "restricted-basis-entry rule."

	D_i^1	D_i^2	D_i^3	D_i^4
a	1/2	0	0	0
b	1	1/4	0	0
c	1	1	7/8	0
d	1	1	1	2

Note that in case d variable D_i^4 does not have to be less than or equal to one because it is the last variable in the set.

Using the "restricted-basis-entry rule" the variable x_i in Fig. 2.1 is represented by

$$x_i = \Delta x_i^1 D_i^1 + \Delta x_i^2 D_i^2 + \Delta x_i^3 D_i^3 + \Delta x_i^4 D_i^4 \quad (5)$$

and the function $f_i(x_i)$ is approximated by

$$f_i \approx \Delta f_i^1 D_i^1 + \Delta f_i^2 D_i^2 + \Delta f_i^3 D_i^3 + \Delta f_i^4 D_i^4 \quad (6)$$

It is apparent that x_i is represented exactly and that $f_i(x_i)$ is approximated by what is essentially linear interpolation.

Example 2

Approximate $f_i(x_i) = (x_i)^3$ using three partitions over $(0,2)$. Let $S = \{0, 1/2, 1, 2\}$.

$$\begin{array}{ll} x_i^0 = 0 & f_i(x_i^0) = 0 \\ x_i^1 = 1/2 & f_i(x_i^1) = 1/8 \\ x_i^2 = 1 & f_i(x_i^2) = 1 \\ x_i^3 = 2 & f_i(x_i^3) = 8 \end{array}$$

The differences can now be calculated.

j	1	2	3
Δx_i^j	$1/2$	$1/2$	1
Δf_i^j	$1/8$	$7/8$	7

Using the differences, the exact equation for x_i and the approximating equation for $f_i(x_i)$ can be written as

$$x_i = 1/2 D_i^1 + 1/2 D_i^2 + 1 \cdot D_i^3$$

and

$$f_i \approx 1/8 D_i^1 + 7/8 D_i^2 + 7 \cdot D_i^3$$

Some sample calculations comparing the approximate and the exact representations of $f_i(x_i)$ are given in Table 2.1. Notice that when x_i is outside the interval domain ($x_i = 4$) the approximation becomes very poor.

The development so far has implicitly assumed $x_i^0 = 0$ and $f_i(x_i^0) = 0$.

This is not generally true. The general expressions for x_i and $f_i(x_i)$ given that $x_i^0 \neq 0$ and $f_i(x_i^0) \neq 0$ with n_i partitions are

$$x_i = x_i^0 + \sum_{j=1}^{n_i} \Delta x_i^j D_i^j \quad (7)$$

and

$$f_i \approx f_i(x_i^0) + \sum_{j=1}^{n_i} \Delta f_i^j D_i^j \quad (8)$$

This is extremely important when working with logarithmic and exponential functions.

Example 3

Approximate $f_i(x_i) = e^{\frac{x_i}{2}}$ with two uniform partitions over $(0, 2)$ (see Fig. 2.2). Notice that when $x_i^0 = 0$ the value of $e^{\frac{x_i}{2}} = 1.0$. The piecewise linear approximation function must have the general form of equation (8).

The function values at the partition boundaries are

$$\begin{array}{ll} x_i^0 = 0 & f_i(x_i^0) = 1 \\ x_i^1 = 1 & f_i(x_i^1) = 2.7183 \\ x_i^2 = 2 & f_i(x_i^2) = 7.3891 \end{array}$$

The differences are now calculated.

j	1	2
Δx_i^j	1.0	1.0
Δf_i^j	1.7183	4.6708

(given x_i)	D_i^1	D_i^2	D_i^3	f_i (approximate)	f_i (exact)
1	1	1	0	1	1
$\frac{3}{4}$	1	$\frac{1}{2}$	0	$\frac{9}{16} \approx .563$	$\frac{27}{64} \approx .422$
$\frac{3}{2}$	1	1	$\frac{1}{2}$	$4 \frac{1}{2}$	$3 \frac{3}{8}$
4	1	1	3	22	64

Table 2.1. Sample calculations comparing some exact and approximate values.

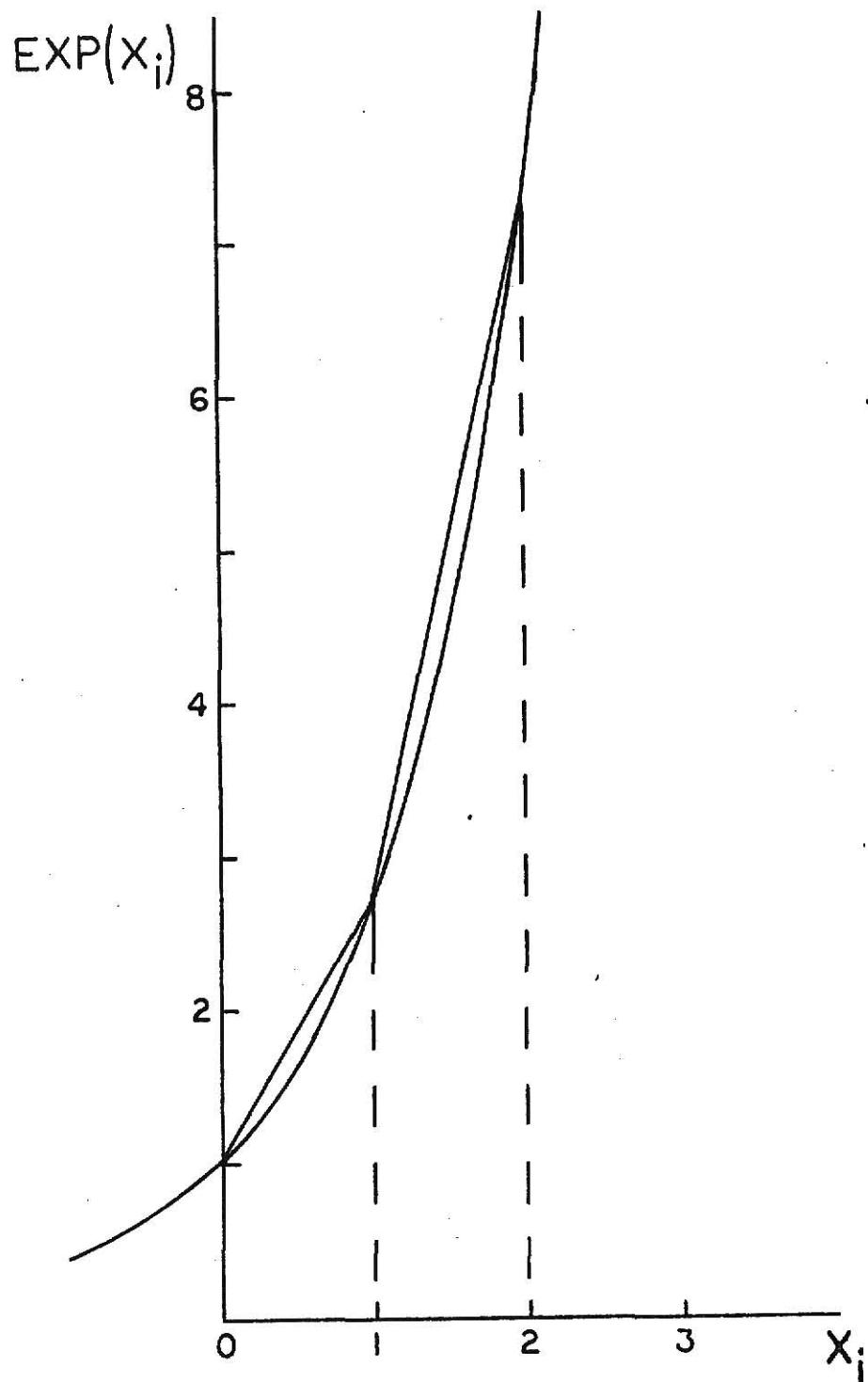


Fig. 2.2. $\text{Exp}(x_i)$ with two uniform partitions over $(0, 2)$.

The piecewise linear functions are

$$x_i = 1.0 D_i^1 + 1.0 D_i^2$$

and

$$\bar{x}_i \approx 1.0 + 1.7183 D_i^1 + 4.6708 D_i^2$$

Now consider a constraint of the form

$$\bar{x}_i \leq b_k$$

The above approximation is applied in the following manner.

$$1.0 + 1.7183 D_i^1 + 4.6708 D_i^2 \leq b_k$$

or

$$1.7183 D_i^1 + 4.6708 D_i^2 \leq b_k - 1.0$$

The constant term is moved to the right hand side of the inequality without altering its net effect. The constraint is now a linear function subject to the "restricted-basis-entry rule."

2.3 THE SEPARABLE PROGRAMMING PROBLEM

The separable programming problem as defined before is
maximize (or minimize)

$$c(\bar{x}) = \sum_{i=1}^m f_i(x_i) \quad (2)$$

subject to the constraints

$$\sum_{i=1}^m g_{ki}(x_i) \leq b_k, \quad k = 1, 2, \dots, p \quad (3)$$

$$x_i \geq 0$$

The number of variables equals m and the number of constraints, p .

If n_i is the number of partitions for the variable x_i then the variable x_i and the approximations for the functions $f_i(x_i)$ and $g_{ki}(x_i)$ can be written as

$$x_i = x_i^0 + \sum_{j=1}^{n_i} \Delta x_i^j D_i^j \quad (9)$$

$$f_i(x_i) \approx f_i(x_i^0) + \sum_{j=1}^{n_i} \Delta f_i^j D_i^j \quad (10)$$

$$g_{ki}(x_i) \approx g_{ki}(x_i^0) + \sum_{j=1}^{n_i} \Delta g_{ki}^j D_i^j \quad (11)$$

$$k = 1, 2, \dots, p, \text{ and } i = 1, 2, \dots, m$$

The approximate form of the problem becomes maximize (or minimize)

$$c = \sum_{i=1}^m \sum_{j=1}^{n_i} \Delta f_i^j D_i^j + \sum_{i=1}^m f_i(x_i^0) \quad (12)$$

subject to the constraints

$$\sum_{i=1}^m \sum_{j=1}^{n_i} \Delta g_{ki}^j D_i^j \leq b_k - \sum_{i=1}^m g_{ki}(x_i^0), \quad k = 1, 2, \dots, p \quad (13)$$

$$x_i - \sum_{j=1}^{n_i} \Delta x_i^j D_i^j = x_i^0, \quad i = 1, 2, \dots, m \quad (14)$$

$$0 \leq D_i^j \leq 1, \quad j = 1, 2, \dots, n_i \quad (15)$$

$$x_i \geq 0 \quad (16)$$

If $f_i(x_i^0) \neq 0$ for any i value, equation (12) has a constant term,

$\sum_{i=1}^m f_i(x_i^0)$, which may be dropped without effect on the optimum solution.

If any $g_{ki}(x_i^0) \neq 0$ then the bound, b_k , must be adjusted. This is important because the problem changes if they are dropped. And finally, if $x_i^0 \neq 0$ the grid equations, equation (14), will not be equal to zero. The number of original constraints equals p , the number of grid equations equals m , and the number of constraints due to the piecewise linear approximation functions, $D_i^j \leq 1$, equals $\sum_{i=1}^m n_i$. The problem now has $p + m + \sum_{i=1}^m n_i$ constraints.

The separable programming problem represented by equations (12) through (16) is now a linear programming problem that is subject to the "restricted-basis-entry rule." Any basis transformation cannot violate the "restricted-basis-entry rule", hence, the origin of this title. Put simply, a variable D_i^j cannot enter the basis unless all of the previous D_i^j values are equal to one and all of the following D_i^j values are zero. This rule insures that the approximated functions are represented correctly.

Example 4.

A simple separable programming problem shall be solved. The problem is maximize

$$c(\bar{x}) = (2x_1 - (x_1)^2) + (x_2 - (x_2)^2)$$

subject to the constraint

$$(x_1)^2 + 2(x_2)^2 \leq 1$$

$$x_1 \geq 0, \quad x_2 \geq 0$$

The functions of this problem are separable.

$$f_1(x_1) = 2x_1 - (x_1)^2, \quad f_2(x_2) = x_2 - (x_2)^2$$

$$g_{11}(x_1) = (x_1)^2, \quad g_{12}(x_2) = 2(x_2)^2$$

The variable domains are chosen as (0, 1) and (0, 1/2) for x_1 and x_2 , respectively. The number of partitions for x_1 is $n_1 = 4$ and for x_2 is $n_2 = 2$. The partitions are uniform for simplicity. The function values at the interval boundaries and the calculated differences are given in

Table 2.2. The approximating equations are determined as

Function values

j	x_1^j	f_1^j	g_{11}^j	x_2^j	f_2^j	g_{12}^j
0	0	0	0	0	0	0
1	$\frac{1}{4}$	$\frac{7}{16}$	$\frac{1}{16}$	$\frac{1}{4}$	$\frac{3}{16}$	$\frac{1}{8}$
2	$\frac{1}{2}$	$\frac{3}{4}$	$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{2}$
3	$\frac{3}{4}$	$\frac{15}{16}$	$\frac{9}{16}$			
4	1	1	1			

Function differences

j	Δx_1^j	Δf_1^j	Δg_{11}^j	Δx_2^j	Δf_2^j	Δg_{12}^j
1	$\frac{1}{4}$	$\frac{7}{16}$	$\frac{1}{16}$	$\frac{1}{4}$	$\frac{3}{16}$	$\frac{1}{8}$
2	$\frac{1}{4}$	$\frac{5}{16}$	$\frac{3}{16}$	$\frac{1}{4}$	$\frac{1}{16}$	$\frac{3}{8}$
3	$\frac{1}{4}$	$\frac{3}{16}$	$\frac{5}{16}$			
4	$\frac{1}{4}$	$\frac{1}{16}$	$\frac{7}{16}$			

Table 2.2. The function values and differences for example 4.

$$f_1(x_1) \approx \frac{7}{16} D_1^1 + \frac{5}{16} D_1^2 + \frac{3}{16} D_1^3 + \frac{1}{16} D_1^4$$

$$f_2(x_2) \approx \frac{3}{16} D_2^1 + \frac{1}{16} D_2^2$$

$$g_{11}(x_1) \approx \frac{1}{16} D_1^1 + \frac{3}{16} D_1^2 + \frac{5}{16} D_1^3 + \frac{7}{16} D_1^4$$

$$g_{12}(x_2) \approx \frac{1}{8} D_2^1 + \frac{3}{8} D_2^2$$

The grid equations are

$$x_1 = \frac{1}{4} D_1^1 + \frac{1}{4} D_1^2 + \frac{1}{4} D_1^3 + \frac{1}{4} D_1^4$$

$$x_2 = \frac{1}{4} D_2^1 + \frac{1}{4} D_2^2$$

It must be remembered that all $D_i^j \leq 1$. This series of constraints must be incorporated in the final L.P. problem. The final problem set-up is maximize

$$c = \frac{7}{16} D_1^1 + \frac{5}{16} D_1^2 + \frac{3}{16} D_1^3 + \frac{1}{16} D_1^4 + \frac{3}{16} D_2^1 + \frac{1}{16} D_2^2$$

subject to the constraints

$$\frac{1}{16} D_1^1 + \frac{3}{16} D_1^2 + \frac{5}{16} D_1^3 + \frac{7}{16} D_1^4 + \frac{1}{8} D_2^1 + \frac{3}{8} D_2^2 \leq 1$$

$$x_1 - \frac{1}{4} D_1^1 - \frac{1}{4} D_1^2 - \frac{1}{4} D_1^3 - \frac{1}{4} D_1^4 = 0$$

$$x_2 - \frac{1}{4} D_2^1 - \frac{1}{4} D_2^2 = 0$$

$$0 \leq D_i^j \leq 1, \quad i = 1, 2, \quad j = 1, \dots, n_i$$

$$x_i \geq 0$$

This problem is easily solved by incorporating the "restricted-basis-entry rule" with any standard simplex tableau. The problem may be solved in five iterations. The starting tableau is given for reference in Table 2.3. In the starting tableau notice that slack variable S_1 is for the original constraint. Slack variables S_i , $i = 2, \dots, 7$, are for the separable variable conditions, $D_i^j \leq 1$. Also notice that it is imperative that the first iteration involves D_1^1 or D_2^1 . If the first iteration involves D_1^1 then the second iteration can involve D_2^1 , or D_1^2 if $D_1^1 = 1$. The iteration process is given in Table 2.4. The final optimum solution is

$$x_1 = \frac{3}{4} = .75000$$

$$x_2 = \frac{11}{24} = .45833$$

$$c(\bar{x}) = \frac{113}{96} = 1.17708$$

It is apparent that the total number of arithmetic calculations could become astronomical for multidimensional problems involving many constraints. An upper bound of the number of differences calculated is

$$\sum_{i=1}^m n_i p + \sum_{i=1}^m n_i$$

where

m = number of variables used in separable functions

n_i = number of partitions for variable i

p = number of nonlinear constraints plus 1 (for a nonlinear objective function)

c_i	b_k	x_1	x_2	D_1^1	D_1^2	D_1^3	D_2^1	D_2^2	S_1	S_2	S_3	S_4	S_5	S_6	S_7
0	S_1	1	0	0	$\frac{1}{16}$	$\frac{5}{16}$	$\frac{3}{16}$	$\frac{1}{16}$	$\frac{3}{16}$	$\frac{1}{16}$	0	0	0	0	0
0	x_1	0	1	0	$-\frac{1}{4}$	$-\frac{1}{4}$	$-\frac{1}{4}$	$-\frac{1}{4}$	0	0	0	0	0	0	0
0	x_2	0	0	1	0	0	0	0	0	$-\frac{1}{4}$	$-\frac{1}{4}$	0	0	0	0
0	S_2	1			1				1						
0	S_3	1				1				1					
0	S_4	1					1				1				
0	S_5	1						1				1			
0	S_6	1							1				1		
0	S_7	1								1				1	
$c = 0$															
0 0 - $\frac{7}{16}$ - $\frac{5}{16}$ - $\frac{3}{16}$ - $\frac{1}{16}$ - $\frac{3}{16}$ - $\frac{1}{16}$ 0 0 0 0 0 0 0 0															

Table 2.3. Starting simplex tableau. (All blank spaces are zero.)

Iteration Number	Vector into Basis	Vector out of Basis
1	D_1^1	s_2
2	D_1^2	s_3
3	D_2^1	s_6
4	D_1^3	s_4
5	D_2^2	s_1

Table 2.4. Iteration process for the revised simplex algorithm. Notice that no basis transformation violates the "restricted-basis-entry rule."

For the reliability problem [45] worked later involving five variables, three constraints, a nonlinear objective function, and $n_i = 11$ for all i , the total number of differences calculated for the given solution were 275. The problem of arithmetic is what justifies a FORTRAN routine to aid in separable programming efforts.

2.4 A REVIEW OF IBM MPS/360 VERSION 2, LINEAR AND SEPARABLE PROGRAMMING

MPS/360 is a set of procedures designed to solve problems involving linear and separable programming. The user controls the procedures by means of a control language program. This section is to briefly review the necessities with emphasis on problem formulation and data input.

Figure 2.3 shows an MPS/360 control language program which formulates the problem. The statements shall be briefly discussed. The "PROGRAM" and "INITIALZ" statements initiate and initialize the system. They are required control statements. The "MOVE (XDATA, 'name')" statement defines the name of the input data. Closely related is the "MOVE (XPBNAME, 'name')" statement. "CONVERT" is the input procedure. It must be supplied with the name of the data and the problem name by the two previous "MOVE" statements . The optional procedure "BCDOUT" is a data echo check procedure. The next two "MOVE" statements define the objective function (XOBJ) and the right hand side (XRHS). They are required for linear and separable programming problems. "TITLE" is an optional statement supplying page headings. The next procedure is "SETUP." This procedure transforms the problem into the form that can be optimized by an optimization procedure. In this case the optimization procedure is "PRIMAL." In "SETUP" a "BOUND" vector is defined as "SEPBBOUND." The problem is maximization since the parameter "MAX" is specified (default is minimization). "SOLUTION" is the final output from "PRIMAL" and the program is ended with

CONTROL PROGRAM CCMPILER - MPS/360 V2-M8

```
(C01      PRCGRAM
(C02      INITIALZ
(C64      MOVE(XDATA,'DATA-SET')
(C65      MOVE(XPBNAME,'EXAMPLE')
(C66      CCNVERT
(C67      BCDOUT
(C68      MOVE(XCBJ,'OBJECT')
(C69      MOVE(XRHS,'LIMITS')
(C70      TITLE('EXAMPLE')
(C71      SETUP('BOUNE','SEPPBOUND','MAX')
(C72      PRIMAL
(C73      SOLUTION
(C74      EXIT
(C75      PEND
```

Fig. 2.3. Listing of an MPS/360 Control Language Program.

"EXIT" and "PEND." In review, the control language program shown in Fig. 2.3 defines a problem named EXAMPLE with input data named DATA-SET. The objective function is OBJECT and the right hand side is LIMITS. This problem, titled EXAMPLE, will have some of its variables bounded by a vector called SEPBOUND and then it will be maximized.

The input data are delivered, for the purposes of this thesis, to MPS/360 in five sections. They are NAME, ROWS, COLUMNS, RHS, and BOUNDS. MPS/360 distinguishes each section by its name. They must be supplied in the given order and each section is defined by the above names, left-justified, starting in card column one. Each section follows a general card format.

Figure 2.4 shows the format for the general MPS/360 input data card. It is divided into six fields with each field being designated for a special purpose. Field 1 is the code field containing information relating to the data on the card. If the data card happens to define a row then the code will indicate N, L, G, or E for "objective function", "less than or equal to", "greater than or equal to", or "equal to", respectively. Fields 2, 3, and 5 are name fields. They can contain an alphabetic name that is 1-8 characters in length and is left-justified. Fields 4 and 6 are numeric fields with a FORTRAN format equivalent of F12.X.

The next two examples demonstrate the input data format. The 'MARKER' cards necessary for distinguishing the special separable variables are described in the second example.

EXAMPLE 5

This example describes how the data is formulated for the simple linear problem of maximize $y = 3X_1 + X_2$

	Field 1	Field 2	Field 3	Field 4	Field 5	Field 6
Columns	2-3	5-12	15-22	25-36	40-47	50-61
Type of Contents	code	name	name	value	name	value

Fig. 2.4. General card format for MPS/360 Linear and Separable Programming [37].

subject to the constraints

$$x_1 + x_2 \leq 2$$

$$0 \leq x_1 \text{ and } x_2 \leq 3/2$$

The input data are shown in Fig. 2.5 and consists of five parts: NAME, ROWS, COLUMNS, RHS, and BOUNDS. The name of the data is defined in the NAME, section. The ROWS section defines the rows of the L.P. problem to MPS/360. A row is coded as N, L, G, or E (objective, \leq , \geq , or $=$) and given a name. In this case OBJECT is the objective function and ROW1 is the constraint. The COLUMNS section defines the elements of the L.P. matrix by variable (column). The RHS section defines the right hand side of the constraints. In this case the RHS vector is named LIMITS. The last section is BOUNDS. The variables are given upper bounds by UP and the BOUNDS vector is named BOUND. The input data is concluded by ENDDATA.

Example 6.

Example 4 is worked using MPS/360. The control language program is shown in Fig. 2.3. Figure 2.6 shows the MPS/360 data echo check. Separable data are denoted by 'MARKER' cards. A 'MARKER', 'SEPORG' card occurs prior to each set of separable variables as shown. Each set is given a name. They are SET1 and SET2 in this example. At the end of all of the separable data occurs a 'MARKER', 'SEPEND' card. This card tells MPS/360 that no more separable variables follow. The 'MARKER' card name is contained in field 2, the 'MARKER' symbol is contained in field 3 and the 'SEPORG' or 'SEPEND' symbol is contained in field 5. All linear data occur before or after the separable data. The 'MARKER' cards are the only differences when using MPS/360 for separable programming. The final solution is shown in Fig. 2.7.

	<u>Columns</u>											
12345	1	1	2	2	3	3	4	4	5	5	6	
	0	5	0	5	0	5	0	5	0	5	0	

NAME DATA-SET
N OBJECT
L ROW1
COLUMNS
 X1 OBJECT 3.0 ROW1 1.0
 X2 OBJECT 1.0 ROW1 1.0
RHS
 LIMITS ROW1 2.0
BOUNDS
 UP BOUND X1 1.5
 UP BOUND X1 1.5
ENDATA

Fig. 2.5. Sample MPS/360 input data. Data are assumed zero unless specified otherwise.

EXECUTCR. MPS/360 V2-M8

NAME	DATA-SET			
ROWS				
N OBJECT				
L ROW1				
E GRID1				
E GRID2				
COLUMNS				
X1	GRID1	1.00000		
X2	GRID2	1.00000		
SET1	'MARKER'		'SEPERG'	
D11	OBJECT	.43750	RCW1	.06250
D11	GRID1	-	.25000	
D12	OBJECT	.31250	ROW1	.18750
D12	GRID1	-	.25000	
D13	OBJECT	.18750	RCW1	.31250
D13	GRID1	-	.25000	
D14	OBJECT	.06250	ROW1	.43750
D14	GRID1	-	.25000	
SET2	'MARKER'		'SEPERG'	
D21	OBJECT	.18750	RCW1	.12500
D21	GRID2	-	.25000	
D22	OBJECT	.06250	RCW1	.37500
D22	GRID2	-	.25000	
ENCSET	'MARKER'		'SEPFND'	
RHS				
LIMITS	ROW1	1.00000		
BOUNDS				
UP SEPBOUND	D11	1.00000		
UP SEPBOUND	D12	1.00000		
UP SEPBOUND	D13	1.00000		
UP SEPBOUND	D14	1.00000		
UP SEPBOUND	D21	1.00000		
UP SEPBOUND	D22	1.00000		
ENDATA				

Fig. 2.6. Input echo check from "BCDOUT".

Fig. 2.7. MPS/360 solution for Example 4.

EXAMPLE

SOLUTION (OPTIMAL)

TIME = 0.30 MINS. ITERATION NUMBER = 5

```
...NAME... ACTIVITY... DEFINED AS
FUNCTIONAL 1.177C8 OBJECT LIMITS
RESTRAINTS SEPBUND
ACUNDS...
```

EXAMPLE

SECTION 1 - ROWS

NUMBER	ROW	AT	ACTIVITY...	SLACK ACTIVITY	LOWER LIMIT.	UPPER LIMIT.	DUAL ACTIVITY
1	OBJECT	BS	1.177C8		NONE	NONE	1.C0000
2	RCH1	LL	1.C0000		NONE	1.00000	-16667-
A	GRID1	EC					.
A	GRIC2	EC					.
							.
							.

EXAMPLE

SECTION 2 - COLUMNS

NUMBER	COLUMN	AT	ACTIVITY...	INPUT COST..	LOWER LIMIT.	UPPER LIMIT.	REDUCED COST.
5	X1	BS	•75CCC	*	*	*	*
6	X2	BS	•45833	*	*	*	.
7	C11	LL	1.CCCCC	•43750	*	*	427C8
8	C12	LL	1.CCCCC	•31250	*	*	•28125
9	C13	LL	1.CCCCC	•18750	*	*	•13542
1C	C14	LL	1.CCCCC	•C625C	*	*	•01042-
11	C21	LL	1.CCCCC	•18750	*	*	•16667-
12	C22	BS	•83333	•06250	*	*	.
							1.00000

2.5 A FORTRAN SUBROUTINE FOR GENERATING SEPARABLE PROGRAMMING DATA TO BE USED WITH MPS/360

One difficulty in using the separable programming subroutine of MPS/360 is calculating the values of Δx_i^j , Δf_i^j , and Δg_{ki}^j of equations (9), (10), and (11) and punching these values in the input data format. This section presents a FORTRAN program which performs the necessary calculations and punches the data. The general flow chart is shown in Fig. 2.8.

There are three main loops within the flow chart logic. The first loop controls the variable, x_i , $i = 1, \dots, m$, and the computational process on it. Each variable is admitted to the computational process and is partitioned according to four input parameters. The first is x_i^0 , the starting point. The next three parameters are used to partition a predetermined feasible region uniformly. A partition is automatically created between x_i^0 and a specified lower bound, unless the distance between them is zero. From the lower bound to an upper bound the user specifies the number of uniform partitions. So each variable requires x_i^0 , a lower bound, an upper bound, and the number of uniform partitions. The second loop controls the partitioning process of the variable, x_i , and is contained within the first loop. It is at this point that a special separable variable is created, one for each partition. After the partition is created the program punches the data according to the MPS/360 format. The third loop calculates the function differences, Δf_i^j or Δg_{ki}^j , which are the coefficients for the special separable variables, D_i^j . Within this loop a subroutine, named FUNC, defines the separable functions of the problem. For the main program listing, refer to Appendix 1.

The general flow chart for FUNC is shown in Fig. 2.9. The subroutine

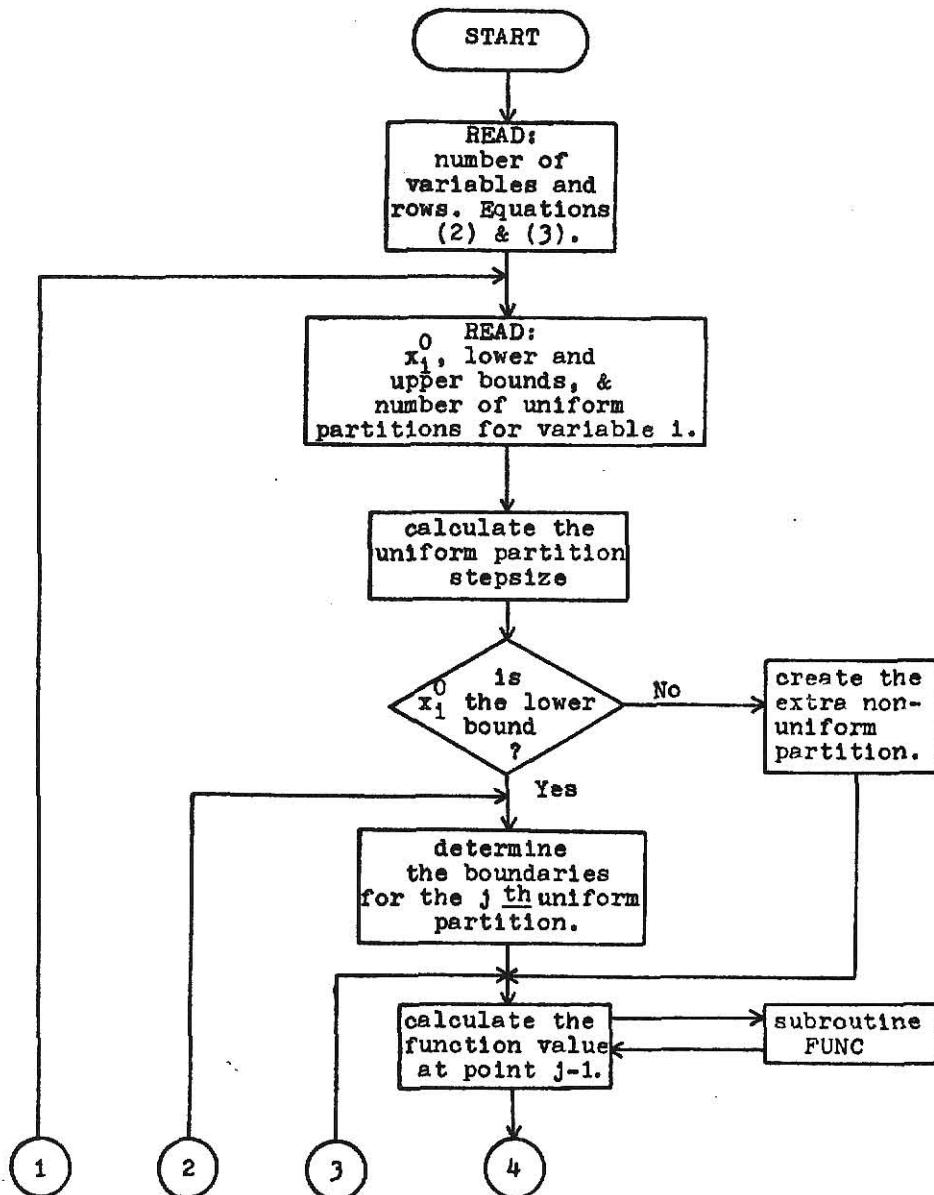


Fig. 2.8. The flow chart of the program that generates the separable programming data.

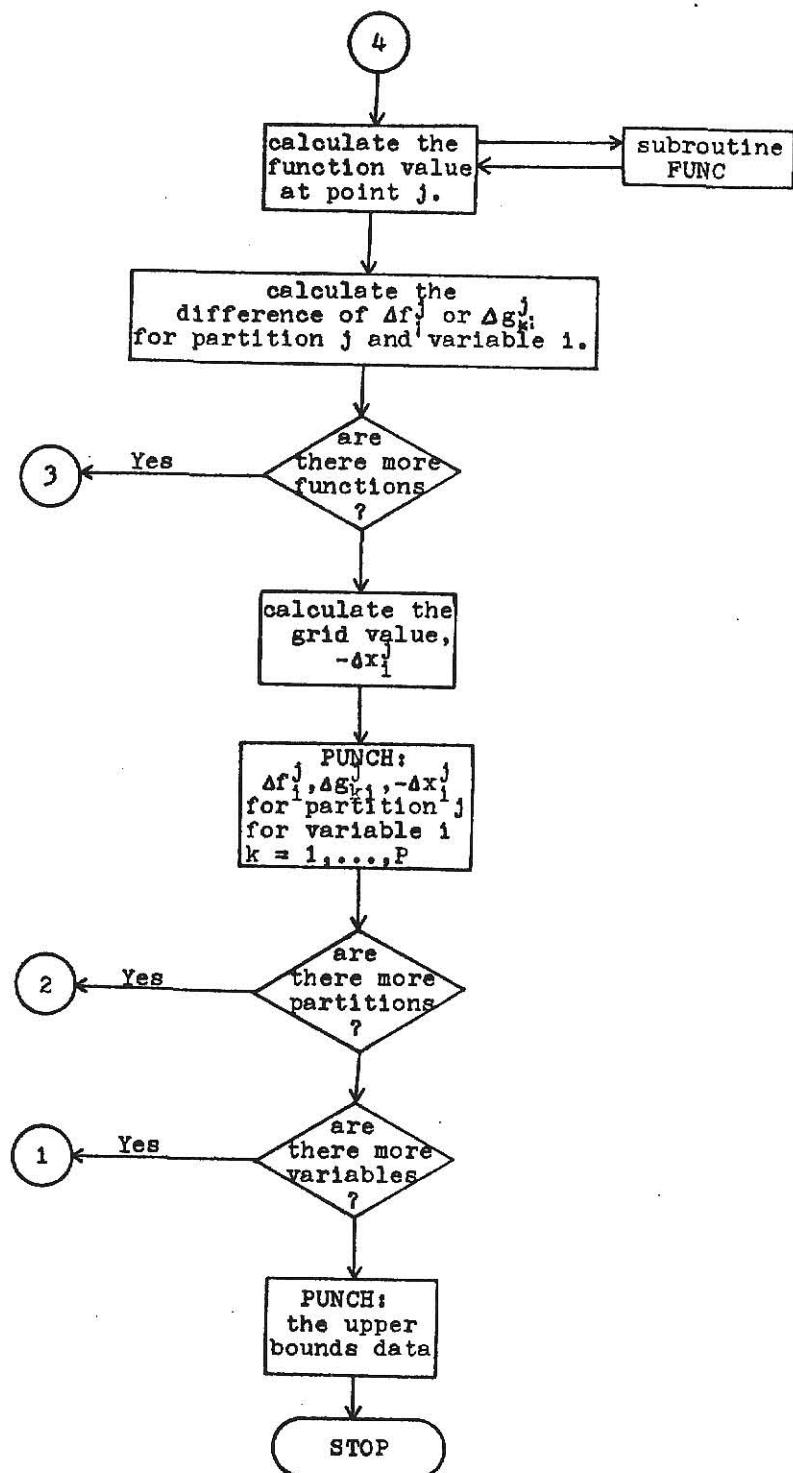


Fig. 2.8. (continued)

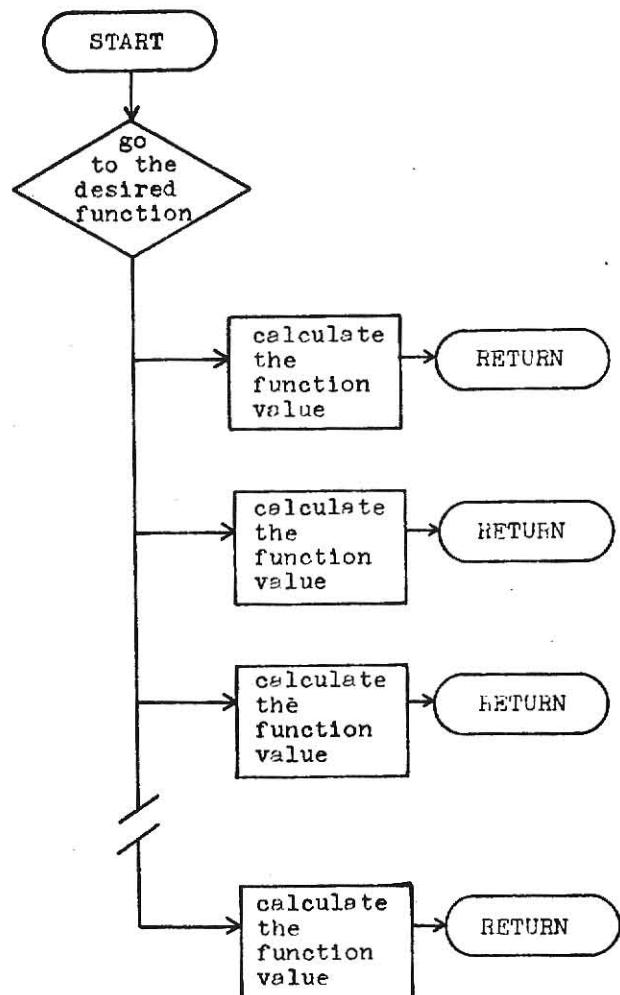


Fig. 2.9. Flow chart for the subroutine FUNC.

contains the separable functions in the form of a matrix, $F(I, J)$. The first row of the matrix, $I = 1$, is for the objective function (or possibly the first constraint if the objective function is linear). The remaining rows are for the constraint functions. The matrix columns represent functions for the variable, x_1, \dots, x_m . Consider Example 4. Figure 2.10 shows the separable function breakdown with row 1 as the objective function and row 2 as the constraint. When FUNC is called by the main program it is supplied with the desired row and column number. FUNC returns the computed value for that function. This may be done in several ways. A practical procedure is to use a "locator function" and a computed GO TO statement to locate the correct function. This is demonstrated in Fig. 2.11. The arithmetic statement $(J - 1) * NROWS + I$, where I is the row, J is the column, and $NROWS$ is the total number of rows in the matrix, will locate $F(I, J)$ when the matrix is ordered linearly by columns. The main program supplies the desired row and column numbers and the value of the dummy variable, x . FUNC returns the desired function value. Figure 2.12 exemplifies the construction of FUNC for Example 4. Note that it is not necessary to treat $F(I, J)$ as a matrix, but it aids the programmer.

The next item is the input data. The main program must be supplied two parameter values for initialization purposes and four values for each problem variable. The first two are the dimensions of the separable function matrix, $F(I, J)$. They are $NFUN$, the number of variables, and $NROWS$, the number of matrix rows. They are read only once on the first data card. A set of the next four values is needed for each variable. They are $START$, $INTER$, $ALOW$, and $UPPER$. $START$ is the initial starting point, x_i^0 . Its value is usually zero unless the user is working with a function not defined at

	column 1	column 2
row 1	$F(1, 1) = 2x_1 - (x_1)^2$	$F(1, 2) = x_2 - (x_2)^2$
row 2	$F(2, 1) = (x_1)^2$	$F(2, 2) = 2(x_2)^2$

Fig. 2.10. The matrix of separable functions, $F(I, J)$, for example 4.

```
SUBROUTINE FUNC(I,J,X,F,NFUN,NROWS)
DIMENSION F(50,50)
MM=(J-1)*NROWS+I
GO TO(1,2,...),MM
1 F(1,1)=
  RETURN
2 F(2,1)=
  RETURN
3 F(3,1)=
  RETURN
.
.
.
END
```

Fig. 2.11. A proper construction of FUNC.

```
SUBROUTINE FUNC(I,J,X,F,NFUN,NROWS)
DIMENSION F(50,50)
GO TO(1,2,3,4),MM
1 F(1,1)=2.0*X-X**2
   RETURN
2 F(2,1)=X**2
   RETURN
3 F(1,2)=X-X**2
   RETURN
4 F(2,2)=2.0*X**2
   RETURN
END
```

Fig. 2.12. FUNC for Example 4.

zero. The other three values are concerned with partitioning only the region where the optimum solution is expected to be located. Between ALOW and UPPER, the lower and upper bounds, there will occur INTER uniform partitions. The total number of partitions will be either INTER or INTER + 1. If ALOW is not equal to START an extra partition is created between them. The formats for the input variables are given in Table 2.5 and in the comment section of the main program. It is not necessary but recommended that START \leq ALOW < UPPER. Figure 2.13 shows some sample input data for Example 4.

The output of the program is given in three parts. The first and easiest to identify is the data echo check. All of the input data are echo checked for user convenience. The second part is the generated separable programming data. It is punched in the correct MPS/360 format. Lastly is the bounds section of the data. This punched data is necessary to satisfy the condition that the special separable variables, D_i^j , must be less than or equal to one. (The upper bound of the last special separable variable of each set may be removed.) Figure 2.14 shows the echo check and a printout of the data cards punched for Example 4. Notice that all of the 'MARKER' cards are punched and that the bounds section is separated from the rest of the deck by a line of asterisks for easy separation. ROW101 is the objective function and ROW102 is the constraint. The matrix is always related to the data as follows.

<u>Matrix</u>	<u>Data</u>
row 1	ROW101
row 2	ROW102
row 3	ROW103
.	.
.	.
.	.
row N	ROW100+N

N = # matrix rows = NROWS

FORTRAN SYMBOLS	EXPLANATION	FORMAT
NFUN	Number of separable variables	(2I5)
NROWS	Number of matrix (F(I,J)) rows	
START	Initial starting point (X_1^0)	
INTER	Number of uniform partitions between ALOW and UPPER	(F10.0,I10,2F10.0)
ALOW	Lower bound of the uniform region	
UPPER	Upper bound of the uniform region	

Table 2.5. Input data formats for the FORTRAN program.

	<u>Columns</u>				
	1	2	3	4	5
card 1.	12345	0	0	0	0
	2	2			
	↑	↑			
	(NFUN)	(NROWS)			
card 2.	0.0	4	0.0	1.0	
card 3.	0.0	2	0.0	0.5	
	↑	↑	↑	↑	
	(START)	(INTER)	(ALOW)	(UPPER)	

Fig. 2.13. Sample input data to generate separable data for Example 4.

SET101	'MARKER'		'SEPORG'	
P1001	ROW101	0.43750	ROW102	0.06250
P1001	GRID101	-0.25000		
P1002	ROW101	0.31250	ROW102	0.18750
P1002	GRID101	-0.25000		
P1003	ROW101	0.18750	ROW102	0.31250
P1003	GRID101	-0.25000		
P1004	ROW101	0.06250	ROW102	0.43750
P1004	GRID101	-0.25000		
SET102	'MARKER'		'SEPORG'	
P1005	ROW101	0.18750	ROW102	0.12500
P1005	GRID102	-0.25000		
P1006	ROW101	0.06250	ROW102	0.37500
P1006	GRID102	-0.25000		
ENDSET	'MARKER'		'SEPEND'	

UP SEPBOUND P1001	1.0			
UP SEPBOUND P1002	1.0			
UP SEPBOUND P1003	1.0			
UP SEPBOUND P1004	1.0			
UP SEPBOND P1005	1.0			
UP SEPBOUND P1006	1.0			

ECHO CHECK FOR NFLN & NROWS
 NFUN = 2
 NROWS= 2

ECHO CHECK FCR VARIABLE 1
 START= 0.0
 INTER= 4
 ALLOW = 0.0
 UPPER= 1.00

ECHO CHECK FCR VARIABLE 2
 START= 0.0
 INTER= 2
 ALLOW = 0.0
 UPPER= 0.50

Fig. 2.14. FORTRAN output for Example 4. The Echo Check is printed and the separable data is punched.

The grid equations are GRID101 and GRID102. The grid equation rows in general are

variable 1	GRID101
variable 2	GRID102
.	.
.	.
.	.
variable m	GRID100+m

m = # variables = NFUN

The special variables (D_i^j) are P1001, P1002, etc. For each variable there will be either INTER or INTER+1 special variables depending on the values of START and ALOW.

With the program set-up and ready to go the user might want to know approximately how many data cards will be punched.

An upper bound may be established.

$$\text{the number of cards} \leq \sum_{i=1}^m n_i \left(\frac{P}{2}\right)^* + \sum_{i=1}^m n_i + (m + 2)$$

P = number of matrix rows (NROWS)

n_i = number of partitions for variable i (either INTER
or INTER+1)

m = number of variables used in separable functions (NFUN)

* - rounded up to the nearest integer.

There are minor limitations on this program. They are

NFUN \leq 899, NROWS \leq 899

and

$$\sum_{i=1}^m n_i \leq 8999$$

This will allow a problem with 400 variables, averaging 20 partitions per variable, and 400 constraints to be worked without altering the program (except for increasing the dimension of F(50, 50) and Row (50), and the possibility of exceeding core). The output could consist of 1,616,402 cards. This enormous output could be overcome. The MPS/360 data could be stored sequentially on a direct access device. The MPS/360 control language program could then receive the data directly without reading enormous amounts of cards. The given FORTRAN routine can easily be altered to do this type of work, but the subject will not be pursued further this thesis.

Example 7

Example 4 is worked using the proposed subroutine. Figure 2.15 shows the MPS/360 control language program. The results are shown in Fig. 2.16. Remember how to determine which row is the objective function and which grid goes with a given variable.

2.6 SOLUTION TO A RELIABILITY MODEL

The following problem is from an article titled "Systems Reliability Subject to Multiple Nonlinear Constraints" [45], and will be discussed in Chapter 5. The problem is maximize

$$\begin{aligned} S = \ln(1 - .2^{x_1}) + \ln(1 - .15^{x_2}) + \ln(1 - .10^{x_3}) + \ln(1 - .35^{x_4}) \\ + \ln(1 - .25^{x_5}) \end{aligned}$$

subject to the following constraints.

$$x_1^2 + 2x_2^2 + 3x_3^2 + 4x_4^2 + 2x_5^2 \leq 110$$

CONTROL PROGRAM COMPILE - MPS/360 V2-M8

(C01	PROGRAM
(C02	INITIALZ
(C064	MOVE(XDATA,'DATA-SET')
(C065	MOVE(XPRNAME,'EXAMPLE')
(C066	CONVERT
(C067	BCCOUT
(C068	MOVE(XOBJ,'ROW101')
(C069	MOVE(XRHS,'LIMITS')
(C070	TITLE('EXAMPLE')
(C071	SETUP('BOUND','SEPBOUND','MAX')
(C072	PRIMAL
(C073	SOLUTION
(C074	EXIT
(C075	PEND

Fig. 2.15. MPS/360 Control Language Program for Example 4.
Notice that the objective function is defined
as ROW101.

ILLEGIBLE DOCUMENT

**THE FOLLOWING
DOCUMENT(S) IS OF
POOR LEGIBILITY IN
THE ORIGINAL**

**THIS IS THE BEST
COPY AVAILABLE**

Fig. 2.16. MPS/360 solution for example 4 using the data generated by the FORTRAN program.

EXAMPLE
SOLUTION (OPTIMAL)

TIME = 0.28 MINS. ITERATION NUMBER = 6

```
      ••• NAME •••   ••• ACTIVITY •••   DEFINE AS
      FUNCTIONAL          1.17708   RCW1C1
      RESTRAINTS          1.00000   LIMITS
      ROUNDS •••          1.00000   SEPCUND
```

EXAMPLE

SECTION 1 - RCWS

NUMBER	•••RCW••	AT	•••ACTIVITY•••	SLACK ACTIVITY	••LOWER LIMIT•	••UPPER LIMIT•	••DUAL ACTIVITY•
1	RCW101	PS	1.17708	1.177CR-	NONE	NONE	1.00000
2	RCW1C2	LL	1.CCCCC0		NONE		1.16667-
A	GRIC101	EC					
A	GRIC1C2	EC					

EXAMPLE

SECTION 2 - COLUMNS

NUMBER	•COLUMN•	AT	••ACTIVITY•••	••INPUT COST••	••LOWER LIMIT•	••UPPER LIMIT•	••REDUCED COST•
5	X1	ES	•75000	•75000	•	•	•
6	X2	ES	•45833	•45833	•	•	•
7	P1CC1	LL	1.CCCCC0	1.CCCCC0	•43750	•43750	•42703
8	P1CC2	LL	1.CCCCC0	1.CCCCC0	•31250	•31250	•28125
9	P1CC3	LL	1.CCCCC0	1.CCCCC0	•18750	•18750	•13542
10	P1CC4	LL	1.CCCCC0	1.CCCCC0	•C6250	•C6250	•01042-
11	P1CC5	LL	1.CCCCC0	1.CCCCC0	•18750	•18750	•16667
12	P1CC6	ES	•83333	•83333	•C6250	•C6250	•

$$7(x_1 + e^{x_1/4}) + 7(x_2 + e^{x_2/4}) + 5(x_3 + e^{x_3/4}) + 9(x_4 + e^{x_4/4}) \\ + 4(x_5 + e^{x_5/4}) \leq 175$$

$$7(x_1 e^{x_1/4} + 8x_2 e^{x_2/4} + 8x_3 e^{x_3/4} + 6x_4 e^{x_4/4} + 9x_5 e^{x_5/4}) \leq 200$$

Notice that START (x_i^0) cannot equal zero for any x_i since the objective function is a logarithmic sum. This means that the bounds 110, 175, and 200 must be adjusted and the objective function value will be off by a constant (see equations (12) thru (14)). Let START = 1.0 for each variable. With this starting value the adjusted bounds are 98, 101.992, and 151.218, respectively. Figure 2.17 shows the FORTRAN input data echo check. The subroutine FUNC is given in Appendix 2.

The MPS/360 control language program used to solve this problem is given in Appendix 2. This program is somewhat different from the programming used to solve Example 4. The procedure is

1. solve the problem with the special separable variables set at their lower bounds,
2. solve the problem with the special separable variables set at their upper bounds (this is the MPS/360 default procedure),
3. solve the problem by the dual algorithm with the special variables set at their upper bounds.

The procedure is recommended when using MPS/360 [37] to determine the existance of a local optimum solution, if it exists. Separable programming, at its best, will guarantee only a local optimum. One reason is that, unlike linear inequality constraints, nonlinear inequality constraints do not necessarily form a convex set. A second reason is that a nonlinear function is not

```
ECHO CHECK FOR NFUN & NROWS
NFUN =      5
NROWS=      4

ECHO CHECK FOR VARIABLE      1
START=      1.00
INTER=      10
ALOW =      2.00
UPPER=      3.00

ECHO CHECK FOR VARIABLE      2
START=      1.00
INTER=      10
ALOW =      2.00
UPPER=      3.00

ECHO CHECK FOR VARIABLE      3
START=      1.00
INTER=      10
ALOW =      1.50
UPPER=      2.50

ECHO CHECK FOR VARIABLE      4
START=      1.00
INTER=      10
ALOW =      3.00
UPPER=      4.00

ECHO CHECK FOR VARIABLE      5
START=      1.00
INTER=      10
ALOW =      2.50
UPPER=      3.50
```

Fig. 2.17 The FORTRAN data echo check for the reliability problem.

necessarily concave or convex. The only way to guarantee a stationary point is a global maximum is for a function to be concave, or if it is a global minimum the function to be convex. Since the linear approximation function of a separable nonlinear function will reflect its particular concave and convex properties, separable programming will, at its best, produce a local optimum solution. Notice that the problem is not encountered in Example 4 because the constraint forms a convex set and the objective function is a concave function producing a global maximum. The solution of the reliability problem is determined to yield a global maximum because each of the three procedures yields the same result. The control language program data echo check and the solution outputs are shown in Appendix 2.

The solution to the problem is

$$\begin{aligned}x_1 &= 2.70000 \\x_2 &= 2.32929 \\x_3 &= 2.10000 \\x_4 &= 3.50000 \\x_5 &= 2.80000\end{aligned}$$

This is compared to the given solution [45].

$$\begin{aligned}x_1 &= 2.6000 \\x_2 &= 2.2816 \\x_3 &= 2.0075 \\x_4 &= 2.6882 \\x_5 &= 3.3981\end{aligned}$$

It must be remembered that separable programming is an approximate method depending on the fineness of the grid equations for accuracy. The uniform grid for this solution is only .10. The effects of grid size on problem accuracy is dependent on the properties of the approximated functions.

2.7 THE APPLICATION OF THE SEPARABLE PROGRAMMING ALGORITHM TO THE GEOMETRIC PROGRAMMING PROBLEM

Geometric programming is a relatively new optimization technique. The method was originally developed by Duffin, Peterson, and Zener [15]. They defined the "posynomial" optimization problem which was later extended by Passy [40], and Wilde and Beightler [46] to include negative coefficients. The generalized problem is defined as

minimize

$$y_0 = \sum_{t=1}^{T_0} \sigma_{0t} c_{0t} \prod_{n=1}^N x_n^{a_{0tn}}, \quad \sigma_{0t} \equiv \pm 1, \quad c_{0t} > 0$$

subject to the constraints

$$y_m = \sum_{t=1}^{T_m} \sigma_{mt} c_{mt} \prod_{n=1}^N x_n^{a_{mtn}} \leq \sigma_m \equiv \pm 1$$

$$\sigma_{mt} \equiv \pm 1, \quad m = 1, \dots, M$$

$$c_{mt} > 0, \quad x_n > 0$$

This problem is termed the primal problem. It is greatly simplified by working with the associated dual problem, then solving for the primal variables in terms of the dual variables. The dual program is [46]

maximize

$$V(\bar{w}) = \sigma \left(\prod_{m=1}^M \prod_{t=1}^{T_m} \left(\frac{c_{mt} w_{m0}}{w_{mt}} \right)^{\sigma_{mt}} w_{mt}^{\sigma} \right)^\sigma$$

subject to the linear constraints

$$\sum_{t=1}^{T_0} \sigma_{0t} w_{0t} = \sigma \equiv \pm 1 \quad (\text{normality condition})$$

$$\sum_{m=0}^M \sum_{t=1}^{T_m} \sigma_{mt} a_{mtn} w_{mt} = 0, \quad n = 1, \dots, N \quad (\text{orthogonality condition})$$

$$w_{m0} \equiv \sigma_m \sum_{t=1}^{T_m} \sigma_{mt} w_{mt} \geq 0, \quad m = 1, \dots, M$$

$$w_{mt} \geq 0$$

The following conditions are assumed.

$$w_{00} \equiv 1$$

$$\lim_{w_{mt} \rightarrow 0} \left(\frac{c_{mt} w_{m0}}{w_{mt}} \right)^{\sigma_{mt} w_{mt}} = 1$$

The relationship between the primal and the dual variables are

$$c_{0t} \prod_{n=1}^N x_n^{a_{0tn}} = w_{0t} \sigma V(\bar{w}), \quad t = 1, \dots, T_0$$

and

$$c_{mt} \prod_{n=1}^N x_n^{a_{mtn}} = \frac{w_{mt}}{w_{m0}}, \quad t = 1, \dots, T_m, \quad m = 1, \dots, M$$

Note that these equations are linear in $\ln(x_n)$ and can be easily solved.

For the case when

$$\sigma = 1$$

$$\sigma_m = 1, \quad m = 1, \dots, M$$

$$\sigma_{mt} = 1, \quad t = 1, \dots, T_m$$

a strict inequality relationship exists between y_0 and $V(\bar{w})$ as

$$y_0(\bar{x}) \geq V(\bar{w})$$

At the optimum solution, \bar{x}^* ,

$$y_0(\bar{x}^*) = V(\bar{w}^*)$$

Under this condition the solution is a global optimum. This is not true of the generalized version which yields only a local optimum solution [46].

The beauty of geometric programming occurs when the total number of terms is one greater than the number of variables. In this case the solution is determined by solving the linear constraints without reference to the objective function.

Let

$$T = \sum_{m=0}^M T_m$$

The total degrees of difficulty is defined as

$$D = T - N - 1$$

When D is greater than zero the dual problem is not so easily optimized. The remaining portion of this section describes how to apply separable programming to optimize the dual geometric program with D degrees of difficulty.

Since the constraints are linear there is no problem in applying a separable programming algorithm to them. The only nonlinear portion is the objective function which is not obviously separable. The dual objective function can be made separable by a simple linear transformation.

$$V(\bar{w}) = \sigma \left(\prod_{m=0}^M \prod_{t=1}^{T_m} \left(\frac{c_{mt}}{w_{mt}} \right)^{\sigma} w_{mt} \right) \sigma$$

$$= \sigma \left(\prod_{m=0}^M \prod_{t=0}^{T_m} \left(\frac{c_{mt}}{w_{mt}} \right)^{\sigma} w_{mt} \right) \sum_{t=1}^{T_m} \sigma_{mt} w_{mt}$$

Since $\sigma_m = \pm 1$,

$$\sigma_m w_{m0} = \frac{w_{m0}}{\sigma_m} = \sum_{t=1}^{T_m} \sigma_{mt} w_{mt}$$

$$w_{00} \equiv 1$$

The final result is

$$V(\bar{w}) = \sigma \left(\prod_{m=0}^M \prod_{t=1}^{T_m} \left(\frac{c_{mt}}{w_{mt}} \right)^{\sigma} w_{mt} \right) \sigma w_{m0}$$

The linear transformation is

$$w_{m0} = \sigma_m \sum_{t=0}^{T_m} \sigma_{mt} w_{mt}$$

If $\sigma = 1$ then $V(\bar{w})$ is made separable by taking its natural log. If $\sigma = -1$ then the objective function can be maximized by minimizing (the positive function)

$$- V(\bar{w})$$

This function can also be made separable by taking its natural log.

Example 8

Minimize

$$y_0 = 4x_1 + 10x_2 + 4x_3 + 2(x_1^2 + x_2^2)^{\frac{1}{2}}$$

subject to the constraint

$$x_1 x_2 x_3 \geq 100, x_i \geq 0$$

The problem may be formulated as a geometric programming problem by allowing

$$x_4 = (x_1^2 + x_2^2)^{\frac{1}{2}}$$

It now has the posynomial form of Duffin, Peterson, and Zener [15].

Minimize

$$y_0 = 4x_1 + 10x_2 + 4x_3 + 2x_4$$

subject to the constraints

$$y_1 = x_1^2 x_4^{-2} + x_2^2 x_4^{-2} \leq 1$$

$$y_2 = \frac{100}{x_1 x_2 x_3} \leq 1$$

The inequality constraint

$$x_1^2 + x_2^2 \leq x_4^2$$

is not unreasonable since the objective function is minimized only when $x_1^2 + x_2^2 = x_4^2$. The problem has 7-4-1 = 2 degrees of difficulty. The dual problem becomes maximize the product function

$$V = \left(\frac{4}{w_{01}}\right)^{w_{01}} \left(\frac{10}{w_{02}}\right)^{w_{02}} \left(\frac{4}{w_{03}}\right)^{w_{03}} \left(\frac{2}{w_{04}}\right)^{w_{04}} \left(\frac{1}{w_{11}}\right)^{w_{11}}$$

$$\left(\frac{1}{w_{12}}\right)^{w_{12}} \left(\frac{100}{w_{21}}\right)^{w_{21}} (w_{11} + w_{12})^{(w_{11} + w_{12})} (w_{21})^{w_{21}}$$

subject to the constraints

$$\begin{array}{rcl} w_{01} + w_{02} + w_{03} + w_{04} & & = 1 \\ w_{01} & + 2w_{11} & - w_{21} = 0 \\ w_{02} & + 2w_{12} & - w_{21} = 0 \\ w_{03} & & - w_{21} = 0 \\ w_{04} - 2w_{11} - 2w_{12} & & = 0 \end{array}$$

To put the objective function in the separable form let

$$A_1 = w_{11} + w_{12}$$

$$A_2 = w_{21}$$

The objective function to be maximized becomes

$$\begin{aligned} \ln V &= w_{01} \ln\left(\frac{4}{w_{01}}\right) + w_{02} \ln\left(\frac{10}{w_{02}}\right) + w_{03} \ln\left(\frac{4}{w_{03}}\right) + w_{04} \ln\left(\frac{2}{w_{04}}\right) \\ &\quad + w_{11} \ln\left(\frac{1}{w_{11}}\right) + w_{12} \ln\left(\frac{1}{w_{12}}\right) + w_{21} \ln\left(\frac{100}{w_{21}}\right) + A_1 \ln(A_1) + A_2 \ln(A_2) \end{aligned}$$

In addition to the previous linear constraints the constraints

$$A_1 - w_{11} - w_{12} = 0$$

$$A_2 - w_{12} = 0$$

are imposed on the problem.

A FORTRAN program was developed to generate data to solve dual geometric programming problems using MPS/360 and is given in Appendix 3. A solution was determined with a final uniform grid size of 0.0005. Values for the dual variables were determined as

$$w_{01} = 0.23117$$

$$w_{02} = 0.30500$$

$$w_{03} = 0.33333$$

$$w_{04} = 0.13050$$

$$w_{11} = 0.05108$$

$$w_{12} = 0.01417$$

$$w_{21} = 0.33333$$

The dual objective function optimal value was given as

$$V = \text{Exp}(4.47719) = 87.98708$$

Solving for the primal variables involved the solution of the following simultaneous equations (quite trivial for this simple example)

$$4X_1 = (0.23117)(87.98708)$$

$$10X_2 = (0.30500)(87.98708)$$

$$4X_3 = (0.33333)(87.98708)$$

$$2X_4 = (0.13050)(87.98708)$$

yielding,

$$\begin{aligned}x_1 &= 5.085 \\x_2 &= 2.684 \\x_3 &= 7.332 \\x_4 &= 5.741 \approx (x_1^2 + x_2^2)^{\frac{1}{2}}\end{aligned}$$

Substituting these values into the primal objective function yielded

$$y_0 = 87.990$$

Since this is a posynomial problem

$$87.98708 < v(\bar{w}^*) = y_0(\bar{x}^*) < 87.990$$

The given solution is apparently very close to the true optimum.

2.8 CONCLUDING REMARKS

There are several important topics that should be mentioned. One already briefly mentioned in section 2.6 on the reliability problem is the subject of local versus global optimum solutions. Another is the question of convergence by employing finer grids for better approximations. Finally there is the question of creating separability. These topics have not been developed because they are problem dependent. That is, the methods employed in solving or alleviating the above problems cannot be generalized to include all cases.

There are two reasons why separable programming will yield a local optimum solution. One is that a nonlinear separable inequality constraint does not necessarily constitute a convex set. Likewise, the piecewise linear approximation function does not produce a convex set. It is basic to linear programming theory that a set of linear constraints necessarily form a convex set. How, then, can a piecewise linear approximation function form a non-convex set? The "restricted basis entry rule" is the reason. This rule

makes the piecewise linear approximation function nonlinear in nature. This function is not linear but psuedo-linear. A similar condition occurs with the objective function, the second reason why a local optimum solution can occur. The piecewise linear approximation function will in general be concave over one region and convex over another. The result is the possibility of a local optimum solution. When the inequality constraints form a convex set and the objective function is concave for a maximum and convex for a minimum, the solution will always be global. In general though, the separable programming algorithm produces, at its best, a local optimum solutuion. For an example involving a local optimum solution refer to the MPS/360 Linear and Separable Programming User's Manual [37], chapter 5.

Grid size is another important consideration in separable programming. Two things must be considered when determining grid size. The first is computational time and the second is computational precision. In theory, a sequence of piecewise linear approximation functions exists that converge uniformly to the separable function. In reality, we are limited in time and precision and thus, can settle only for an approximation. A good procedure is to use approximately 20 partitions per variable and solve the problem. Use this solution to partition a region around the approximate solution and solve the problem again. If the solution changes sufficiently try this again, but remember, as the grid becomes finer the coefficients of the special separable variables become smaller and the problem will suffer from lack of machine precision. Because of this, the proposed FORTRAN routine will allow no more than 50 partitions per variable. Also, more partitions mean more iterations for the revised simplex method. Computational time can become excessive.

A final consideration when using separable programming is that it can be extended to other nonlinear problems by creating separability. The reliability problem of section 2.6 and the geometric programming extension are the result of a transformation. The objective functions were both originally a product function and were made separable by taking their logarithm. Transformation techniques are discussed in the literature [31, 39] and can be important since many nonlinear programming techniques are not as simple as separable programming.

In closing, separable programming is a powerful nonlinear programming technique with the same sensitivity analysis potential as linear programming. As with any nonlinear technique it will yield at least a local optimum solution, if it exists. With the proposed FORTRAN routine and MPS/360, separable programming can be performed with little difficulty. Its main disadvantage is precision, but this is of little consequence since most engineering problems need only a good approximation. Separable programming should be considered as a solution technique when a problem involves separable functions or functions that can be made separable by a transformation.

CHAPTER THREE

INTRODUCTION TO THE GENERALIZED REDUCED GRADIENT METHOD

3.1 INTRODUCTION

Attempts to solve nonlinear programming problems have resulted in many algorithms that work perfectly for special cases. Examples are separable programming, quadratic programming, geometric programming, and the Wolfe reduced gradient method[47,48]. In an attempt to develope a procedure that will handle the general nonlinear programming problem, the generalized reduced gradient method (GRG) has been proposed by Abadie and Carpentier [3, 4, 5]. The algorithm is generalized from the Wolfe reduced gradient method[47, 48].

The Wolfe reduced gradient method solves problems with a nonlinear objective function and linear constraints. It classes the variables as independent and dependent, and substitutes into the objective function the expressions obtained from the linear constraints in independent variables for the dependent variables. This essentially reduces the original problem to an unconstrained problem of reduced dimension. From this point, a variety of optimization techniques may be used. Applying the same concept to a nonlinear set of constraints complicates matters, but it is possible by using numerical methods.

The GRG method has been studied extensively and coded in FORTRAN by Abadie and Guigou [1, 6, 7, 29, 30]. Three generations of programs have been developed. The first was an experimental code called GRG 66 which was followed by the second code, GRG 69. Both procedures compared favorably to other nonlinear programming codes by ranking highly in the Colville study [7, 10]. GRG 69 was an exceptional standout among the other participants. An improved code, GREG, is the result of the previous work and promises to remain among the highly regarded nonlinear programming procedures.

3.2 DEFINITION OF NOTATION

Notation used in the generalized reduced gradient (GRG) algorithm is explicitly defined in this section. Symbols representing functions, vectors, matrices, gradients, and Jacobians are necessary to describe the algorithm.

An N-dimensional vector, \bar{x} , is represented as

$$\bar{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ \vdots \\ x_N \end{pmatrix}$$

The transpose of a vector, \bar{x} , is \bar{x}^T . The objective function is represented by

$$f_0(\bar{x})$$

and a vector function of M constraints is given as

$$\bar{f}(\bar{x}) = 0$$

where

$$\bar{f}(\bar{x}) = \begin{pmatrix} f_1(\bar{x}) \\ f_2(\bar{x}) \\ \vdots \\ \vdots \\ f_M(\bar{x}) \end{pmatrix}$$

Notice that a bar above a symbol indicates the presence of a vector. The gradient of a function, $f_i(\bar{x})$, is given as a row vector

$$\frac{\partial f_i}{\partial \bar{x}} = \left(\frac{\partial f_i}{\partial x_1}, \frac{\partial f_i}{\partial x_2}, \dots, \frac{\partial f_i}{\partial x_N} \right), \quad i = 0, 1, \dots, M$$

The "gradient of a vector function", is represented by

$$\frac{\partial \bar{f}}{\partial \bar{x}} = \begin{pmatrix} \frac{\partial f_1}{\partial \bar{x}} \\ \frac{\partial f_2}{\partial \bar{x}} \\ \vdots \\ \frac{\partial f_M}{\partial \bar{x}} \end{pmatrix} = \begin{pmatrix} \frac{\partial f_1}{\partial x_1} & \frac{\partial f_1}{\partial x_2} & \dots & \frac{\partial f_1}{\partial x_N} \\ \frac{\partial f_2}{\partial x_1} & \frac{\partial f_2}{\partial x_2} & \dots & \frac{\partial f_2}{\partial x_N} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial f_M}{\partial x_1} & \frac{\partial f_M}{\partial x_2} & \dots & \frac{\partial f_M}{\partial x_N} \end{pmatrix}$$

This operation results in an $M \times N$ matrix.

Another symbol used is related to the real number inequality relations.

The notation

$$\underline{a} \leq \bar{x} \leq \bar{b}$$

implies

$$a_i \leq x_i \leq b_i, \quad i = 1, \dots, N$$

This relation forms the boundary conditions for the problem variables.

3.3 STATEMENT OF THE PROBLEM

The general nonlinear programming problem may be defined in the following form. Maximize

$$f_0(\bar{x})$$

subject to the constraints

$$\bar{f}(\bar{x}) = 0$$

$$\underline{a} \leq \bar{x} \leq \bar{b}$$

Note that an inequality constraint

$$G_i(\bar{x}) \leq 0$$

can be redefined as

$$f_i(\bar{x}) = G_i(\bar{x}) + s_i = 0, \quad s_i \geq 0$$

The addition of a slack variable, s_i , is analogous to the linear programming technique.

Let the set formed by the constraint functions be described as

$$V = \{\bar{X} \mid f_i(\bar{X}) = 0, i = 1, \dots, M\}$$

and the set formed by the boundary conditions as

$$P = \{\bar{X} \mid a_j \leq X_j \leq b_j, j = 1, \dots, N\}$$

The solution set of the problem becomes

$$S = V \cap P$$

the intersection of V and P . For a solution to exist, the set S must not be empty.

3.4 THE PRINCIPLE OF THE REDUCED GRADIENT

The generalized reduced gradient algorithm is based on a very basic optimization procedure which transforms a constrained optimization problem into one that is unconstrained. This is done by dividing the solution vector components into two groups, independent and dependent. The dependent variables, denoted by the dummy vector \bar{y} , are solved for in terms of the independent variables, represented by \bar{x} , via the constraint functions.

In terms of dependent and independent variables the constraint may be written as

$$\bar{f}(\bar{X}) = \bar{f}(\bar{x}, \bar{y}) = 0$$

The strategy of representing \bar{y} in terms of \bar{x} attempts to find a set of functions such that

$$\bar{y} = \phi(\bar{x})$$

The number of components of \bar{y} is M , and the number of components of \bar{x} is

(N-M). The vector function, $\bar{\phi}(\bar{x})$, is M-dimensional. Values for \bar{y} are then substituted into the objective function such that

$$f_0(\bar{x}) = f_0(\bar{x}, \bar{y}) = f_0\left(\bar{x}, \bar{\phi}(\bar{x})\right) = F(\bar{x})$$

The problem is simplified to maximize

$$F(\bar{x})$$

subject to the constraint

$$\bar{a} \leq \bar{x} \leq \bar{b}$$

Notice that the dimension of the solution vector \bar{x} of F , is only (N-M). The problem has been "reduced."

Example 1

Maximize

$$f_0(\bar{x}) = (x_1 - x_1^2) + (x_2 - x_2^2)$$

subject to the constraints

$$x_1^2 + x_2^2 \leq 1$$

$$0 \leq x_1 \leq 5$$

$$0 \leq x_2 \leq 3$$

The constraint, $f_1(\bar{x})$, may be defined as an equality by adding a slack variable,

x_3 .

$$f_1(\bar{x}) = x_1^2 + x_2^2 + x_3 - 1 = 0$$

$$0 \leq x_1 \leq 5$$

$$0 \leq x_2 \leq 3$$

$$0 \leq x_3 \leq \infty$$

Now, choose x_1 as the dependent variable with the assumption that the boundary condition, $0 \leq x_1 \leq 5$ will not be violated at the optimum solution. Thus,

$$y_1 = x_1 = \sqrt{1 - x_2^2 - x_3}$$

Substituting this into the objective function the problem becomes

maximize

$$F(\bar{x}) = F(x_2, x_3) = \sqrt{1 - x_2^2 - x_3} - (1 - x_2^2 - x_3) + (x_2 - x_2^2)$$

subject to the conditions

$$0 \leq x_2 \leq 3$$

$$0 \leq x_3 \leq \infty$$

The gradient of this problem and what is termed the "reduced gradient" of the original problem is

$$\begin{aligned}\frac{\partial F}{\partial \bar{x}} &= \left\{ \frac{\partial F}{\partial x_2}, \frac{\partial F}{\partial x_3} \right\} \\ &= \left\{ \frac{-x_2}{\sqrt{1 - x_2^2 - x_3}}, \frac{\frac{1}{2}}{\sqrt{1 - x_2^2 - x_3}} + 1 \right\}\end{aligned}$$

The term "reduced gradient" is used because the gradient of the original objective function, $f_0(\bar{x})$, has N components while the gradient of $F(\bar{x})$ has a "reduced" dimension of (N-M).

The conditions that determine an optimum solution, \bar{x}^* , are (for all j)

$$\frac{\partial F}{\partial x_j^*} = 0 \quad \text{if} \quad a_j < x_j^* < b_j$$

$$\frac{\partial F}{\partial x_j^*} \leq 0 \quad \text{if} \quad x_j^* = a_j$$

$$\frac{\partial F}{\partial x_j^*} \geq 0 \quad \text{if} \quad x_j^* = b_j$$

The optimal solution may be found trivially for this problem by solving the equations

$$\frac{-x_2}{\sqrt{1 - x_2^2 - x_3}} + 1 = 0$$

and

$$\frac{-\frac{1}{2}}{\sqrt{1 - x_2^2 - x_3}} + 1 = 0$$

This yields

$$x_2 = \frac{1}{2}$$

and

$$x_3 = \frac{1}{2}$$

By substituting these values into the constraint function it is determined that $x_1 = \frac{1}{2}$, satisfying the condition $0 \leq x_1 \leq 5$.

The direction of movement in the GRG algorithm is always determined by the independent variables and the resulting reduced gradient. The objective function may be written as

$$f_0(\bar{x}) = f_0(\bar{x}, \bar{y}) = F(\bar{x})$$

The reduced gradient is

$$\frac{\partial F}{\partial \bar{x}} = \frac{\partial f_0}{\partial \bar{x}} + \frac{\partial f_0}{\partial \bar{y}} \frac{\partial \bar{y}}{\partial \bar{x}}$$

$$\frac{\partial \bar{F}}{\partial \bar{x}} = \begin{pmatrix} \frac{\partial f_0}{\partial x_1} & \dots & \frac{\partial f_0}{\partial x_j} & \dots & \frac{\partial f_0}{\partial x_{N-M}} \end{pmatrix} + \begin{pmatrix} \frac{\partial f_0}{\partial y_1} & \dots & \frac{\partial f_0}{\partial y_i} & \dots & \frac{\partial f_0}{\partial y_M} \end{pmatrix} \begin{pmatrix} \frac{\partial y_1}{\partial x_1} & \dots & \frac{\partial y_1}{\partial x_{N-M}} \\ \vdots & \ddots & \vdots \\ \frac{\partial y_M}{\partial x_1} & \dots & \frac{\partial y_M}{\partial x_{N-M}} \end{pmatrix}$$

The matrix, $\frac{\partial \bar{y}}{\partial \bar{x}}$, is difficult to determine explicitly. This can be calculated indirectly from the constraints.

$$\bar{f}(\bar{x}) = \bar{f}(\bar{x}, \bar{y}) = 0$$

$$\frac{\partial \bar{f}}{\partial \bar{x}} + \frac{\partial \bar{f}}{\partial \bar{y}} \frac{\partial \bar{y}}{\partial \bar{x}} = 0$$

$$\frac{\partial \bar{f}}{\partial \bar{x}} = - \begin{pmatrix} \frac{\partial \bar{f}}{\partial \bar{y}} \end{pmatrix}^{-1} \begin{pmatrix} \frac{\partial \bar{f}}{\partial \bar{x}} \end{pmatrix}$$

$$= \begin{pmatrix} \frac{\partial f_1}{\partial y_1} & \dots & \frac{\partial f_1}{\partial y_M} \\ \vdots & \ddots & \vdots \\ \frac{\partial f_M}{\partial y_1} & \dots & \frac{\partial f_M}{\partial y_M} \end{pmatrix}^{-1} \begin{pmatrix} \frac{\partial f_1}{\partial x_1} & \dots & \frac{\partial f_1}{\partial x_{(N-M)}} \\ \vdots & \ddots & \vdots \\ \frac{\partial f_M}{\partial x_1} & \dots & \frac{\partial f_M}{\partial x_{(N-M)}} \end{pmatrix}$$

The reduced gradient in the final form becomes,

$$g^T = \frac{\partial \bar{F}}{\partial \bar{x}} = \frac{\partial f_0}{\partial \bar{x}} - \frac{\partial f_0}{\partial \bar{y}} \begin{pmatrix} \frac{\partial \bar{f}}{\partial \bar{y}} \end{pmatrix}^{-1} \begin{pmatrix} \frac{\partial \bar{f}}{\partial \bar{x}} \end{pmatrix} \quad (1)$$

Example 2

The reduced gradient for example 1 will be computed directly from the reduced gradient formula. Let

$$x_1 = x_2$$

$$x_2 = x_3$$

$$y_1 = x_1$$

$$\frac{\partial \bar{f}_0}{\partial \bar{x}} = \begin{pmatrix} \frac{\partial f_0}{\partial x_1}, & \frac{\partial f_0}{\partial x_2} \end{pmatrix} = \begin{pmatrix} \frac{\partial f_0}{\partial x_2}, & \frac{\partial f_0}{\partial x_3} \end{pmatrix} = [1 - 2x_2, 0]$$

$$\frac{\partial \bar{f}_0}{\partial \bar{y}} = \begin{pmatrix} \frac{\partial f_0}{\partial y_1} \end{pmatrix} = \begin{pmatrix} \frac{\partial f_0}{\partial x_1} \end{pmatrix} = [1 - 2x_1]$$

$$\frac{\partial \bar{f}}{\partial \bar{y}} = \begin{pmatrix} \frac{\partial f_1}{\partial y_1} \end{pmatrix} = \begin{pmatrix} \frac{\partial f_1}{\partial x_1} \end{pmatrix} = [2x_1]$$

$$\frac{\partial \bar{f}}{\partial \bar{x}} = \begin{pmatrix} \frac{\partial f_1}{\partial x_1}, & \frac{\partial f_1}{\partial x_2} \end{pmatrix} = \begin{pmatrix} \frac{\partial f_1}{\partial x_2}, & \frac{\partial f_1}{\partial x_3} \end{pmatrix} = [2x_2, 1]$$

The reduced gradient is

$$\begin{aligned} \frac{\partial \bar{F}}{\partial \bar{x}} &= [1 - 2x_2, 0] - [1 - 2x_1] [2x_1]^{-1} [2x_2, 1] \\ &= \begin{pmatrix} 1 - \frac{x_2}{x_1}, & 1 - \frac{1}{x_1} \end{pmatrix} \end{aligned}$$

To put it in terms of Example 1 with $x_1 = \sqrt{1 - x_2^2 - x_3^2}$, the reduced gradient becomes

$$\frac{\partial \bar{F}}{\partial \bar{x}} = \left(\frac{-\bar{x}_2}{\sqrt{1 - \bar{x}_2^2 - \bar{x}_3}} + 1, \frac{\frac{1}{2}}{\sqrt{1 - \bar{x}_2^2 - \bar{x}_3}} + 1 \right)$$

There are two problems in the development of the GRG algorithm. The first is indicated in example 1. If the boundary condition for the dependent variable, $y_1 = x_1$, would have been, say, $1 \leq x_1 \leq 5$, the method would have yielded an optimal solution outside of the boundary, P. In the development of the GRG algorithm this is one condition that must be corrected. The second problem is the inversion of the square matrix, or Jacobian, $\frac{\partial \bar{f}}{\partial \bar{y}}$. This matrix must be non-singular.

As for the non-singularity property it is interesting to note the Implicit Function Theorem [8, 13, 36] which is relevant to the reduced gradient.

The Implicit Function Theorem

Suppose there exists a M-dimensional vector function

$$\bar{f}(\bar{x}) = 0$$

with continuous first order partial derivatives. The vector $\bar{x}^T = [x_1, x_2, \dots, x_N]$ ($N > M$) is partitioned into $\bar{x} = [\bar{x}, \bar{y}]$, where

$$\bar{x}^T = [x_1, \dots, x_{N-M}]$$

and

$$\bar{y}^T = [y_1, \dots, y_M]$$

Also assume that the determinant of the Jacobian, $\frac{\partial \bar{f}}{\partial \bar{y}}$, is not zero.

Then there exists one, and only one, vector function of M-dimension such that

$$\bar{y} = \bar{\phi}(\bar{x})$$

and each component has continuous first order partial derivatives.

This theorem does not insure a functional relation between \bar{y} and \bar{x} if the Jacobian is singular. It also states the existance, but not the form, of the function, $\bar{f}(\bar{x})$. How the values are obtained is not mentioned. In reality, they may not be solved for analytically as in Example 1, but numerical methods do exist making it possible to develop a generalized procedure.

It is worthwhile to note the relation of the Lagrange multipliers to the reduced gradient. The Lagrangian function is

$$L(\bar{x}, \bar{u}) = f_0(\bar{x}) - \bar{u}^T \bar{f}(\bar{x})$$

The Kuhn-Tucker condition for optimality at \bar{x} is

$$\frac{\partial L}{\partial x_j} = \begin{cases} \leq 0 & \text{if } x_j = a_j \\ \geq 0 & \text{if } x_j = b_j \\ = 0 & \text{if } a_j < x_j < b_j \end{cases}$$

Therefore

$$\frac{\partial L}{\partial \bar{x}} = \frac{\partial f_0}{\partial \bar{x}} - \bar{u}^T \frac{\partial \bar{f}}{\partial \bar{x}} = \bar{v}^T \quad (2)$$

$$\frac{\partial L}{\partial y} = \frac{\partial f_0}{\partial y} - \bar{u}^T \frac{\partial \bar{f}}{\partial y} = 0 \quad (3)$$

From equation (3) the vector of Lagrange multipliers is given as

$$\bar{u}^T = \frac{\partial f_0}{\partial y} \left(\frac{\partial \bar{f}}{\partial y} \right)^{-1} \quad (4)$$

Substituting equation (4) into equation (2) yields

$$\frac{\partial L}{\partial \bar{x}} = \frac{\partial f_0}{\partial \bar{x}} - \frac{\partial f_0}{\partial \bar{y}} \left(\frac{\partial \bar{f}}{\partial \bar{y}} \right)^{-1} \left(\frac{\partial \bar{f}}{\partial \bar{x}} \right) = \bar{v}^T$$

which is the reduced gradient given by equation (1).

3.5 THE GENERALIZED REDUCED GRADIENT ALGORITHM

The GRG algorithm depends on partitioning the solution vector, \bar{x} , into \bar{x} and \bar{y} (N-M and M-dimensioanl vectors, respectively). The partitioned vector, \bar{x} , is referred to as the set of independent variables. The vector, \bar{y} , is denoted as containing the dependent variables, which are often called basic variables in analogies to linear programming theory. The behavior of the dependent variables is of extreme importance to the GRG algorithm. The non-degeneracy assumption of the algorithm is that for a given iteration there exists a set of dependent variables that are contained within the boundary conditions, P, and the Jacobian, $\frac{\partial \bar{f}}{\partial \bar{y}}$, is non-singular.

Using the above information, the basic GRG algorithm can be stated in five steps [2].

Step 1

The direction of movement is calculated at the starting point \bar{x}^0 and is denoted as \bar{h}^0 . It is calculated in three sub-steps.

Compute the reduced gradient at the point $\bar{x}^0 = [\bar{x}^0, \bar{y}^0]$.

$$\bar{g}^{0T} = \frac{\partial F}{\partial \bar{x}^0} = \frac{\partial f_0}{\partial \bar{x}^0} - \frac{\partial f_0}{\partial \bar{y}^0} \left(\frac{\partial \bar{f}}{\partial \bar{y}^0} \right)^{-1} \left(\frac{\partial \bar{f}}{\partial \bar{x}^0} \right) \quad (1)$$

Step 1.2

The projected reduced gradient, \bar{p}^0 , is now calculated to satisfy the Kuhn-Tucker condition. For each component of the independent vector, \bar{x} ,

$$p_i^0 = \begin{cases} 0 & \text{if } x_i = \text{lower bound and } g_i^0 \leq 0 \\ 0 & \text{if } x_i = \text{upper bound and } g_i^0 \geq 0 \\ g_i^0 & \text{otherwise} \end{cases}$$

$$i = 1, 2, \dots, N-M$$

Step 1.3

At this point the projected reduced gradient may become the direction of movement with $\bar{h}^0 = \bar{p}^0$. It may also be modified in several ways. The rapidly convergent conjugate gradient methods of Davidon, Fletcher and Powell, and Fletcher and Reeves may be used to modify the final direction of movement. A restriction on the choice is that $\bar{h}^0 \cdot \bar{p}^0 > 0$.

Step 2

It is desired to remain within the feasible region at all times. Since this is not possible, a direction of movement is picked that will remain close to the feasible region. The step from \bar{x}^0 is represented by

$$\bar{x}^0 + \theta \bar{h}^0$$

Likewise, for \bar{y}^0 , the step is

$$\bar{y}^0 + \theta \bar{k}^0$$

The vector, \bar{k}^0 , determines the direction of movement for \bar{y}^0 .

Step 2.1

A desired movement would be along the surface of the constraints, that is tangent to them. This is accomplished by finding the tangent to

$\bar{f}(\bar{x}^0 + \theta\bar{h}^0, \bar{y}^0 + \theta\bar{k}^0) = 0$ at the point $[\bar{x}^0, \bar{y}^0]$. It becomes

$$\frac{\partial \bar{f}}{\partial \bar{x}^0} \bar{h}^0 + \frac{\partial \bar{f}}{\partial \bar{y}^0} \bar{k}^0 = 0$$

or

$$\bar{k}^0 = - \left(\frac{\partial \bar{f}}{\partial \bar{y}^0} \right)^{-1} \left(\frac{\partial \bar{f}}{\partial \bar{x}^0} \right) \bar{h}^0 \quad (5)$$

Step 2.2

The problem is now to maximize

$$f_0(\bar{x}^0 + \theta\bar{h}^0, \bar{y}^0 + \theta\bar{k}^0)$$

for θ . This may be done by an one-dimensional search such as the Golden section search, Fibonacci search, or by quadratic interpolation.

Step 3

Calculate

$$\tilde{x}^1 = \bar{x}^0 + \theta\bar{h}^0, \text{ and } \tilde{y}^1 = \bar{y}^0 + \theta\bar{k}^0$$

The values for the independent variables are projected into the bounds,

$\bar{a} \leq \bar{x} \leq \bar{b}$. The relation is

$$x_j^1 = \begin{cases} \text{lower bound, if } x_j^0 + \theta h_j^0 \leq \text{lower bound} \\ \text{upper bound, if } x_j^0 + \theta h_j^0 \geq \text{upper bound} \\ x_j^0 + \theta h_j^0, \text{ otherwise} \end{cases}$$

$$j = 1, \dots, N-M$$

In general the calculated dependent variables, \tilde{y}^1 , do not satisfy the feasibility conditions. It is hoped, though, that the values are "close" to feasibility. Indeed, if θ is small this condition exists. Usually the

vector, \bar{y}^1 , will be the starting point for finding \bar{y}^1 iteratively at step 4.

Step 4

A feasible solution is developed by solving the system

$$\bar{f}(\bar{x}^1, \bar{y}) = 0$$

This may be done by an iterative method. The existance of

$$\bar{y} = \bar{\phi}(\bar{x}^1)$$

is insured by the Implicit Function Theorem under the non-degeneracy conditions. If a component of \bar{y} violates a boundary condition, a degeneracy, then a change of basis occurs as described in section 7. There are two ends to this procedure.

Step 4.1

If the iterative procedure does not converge to a \bar{y}^1 , then \bar{x}^1 is out of the functional domain. This is alleviated by reducing θ (i.e., set $\theta = \frac{1}{2} \theta$) and returning to step 3. An example of this would be to find y given $x = 2$ for

$$x^2 + y^2 = 1$$

When using Newton's method, convergence of \bar{y}^1 depends on the shape of the vector function, $\bar{f}(\bar{x})$, even if \bar{x}^1 is within its domain. The problem is alleviated by restricting the procedure to a finite number of iterations, then reducing θ and returning to step 3.

Letting \bar{y}^1 be the solution obtained, it must be determined if the new solution, $\bar{x}^1 = (\bar{x}^1, \bar{y}^1)$, improves the objective function. Since \bar{y}^1 is not on the tangent which guarantees improvement then it will improve the objective function only if close to \bar{y}^1 . If the objective function is not improved,

decrease θ by one-half and return to Step 3.

Step 5

In this step it is normal to set $\bar{x}^0 = \bar{x}^1$ and repeat the algorithm.

But if for some reason a better value for θ may be determined by previous information, a return to step 3 is possible.

A flow diagram exemplifying the procedure is given in Fig. 3.1.

Normally, the starting point, \bar{x}^0 , is a feasible solution but this is not necessary. Any non-feasible \bar{x}^0 can be made feasible by adding artificial variables to satisfy the equality constraints, then forcing them to zero values by penalizing the objective function. Let an artificial variable, x_{N+i} , satisfy the condition

$$0 \leq x_{N+i} \leq \infty$$

If, at the starting point, \bar{x}^0 , the i^{th} constraint violates the equality condition and it is greater than zero, $f_i(\bar{x}^0) > 0$, it is made feasible by subtracting the artificial variable, x_{N+i} .

$$f_i(\bar{x}^0) - x_{N+i} = 0$$

For the less than zero condition, $f_i(\bar{x}^0) < 0$, the equation is formulated by adding the artificial variable.

$$f_i(\bar{x}^0) + x_{N+i} = 0$$

There are two methods for obtaining a feasible solution, the penalty method and the two-phase method. For the penalty method let M_i be a large number.

The objective function to be maximized becomes

$$f_0(\bar{x}) - \sum_{i \in NF} M_i x_{N+i}, \quad NF = \{i \mid f_i(\bar{x}) \neq 0\}$$

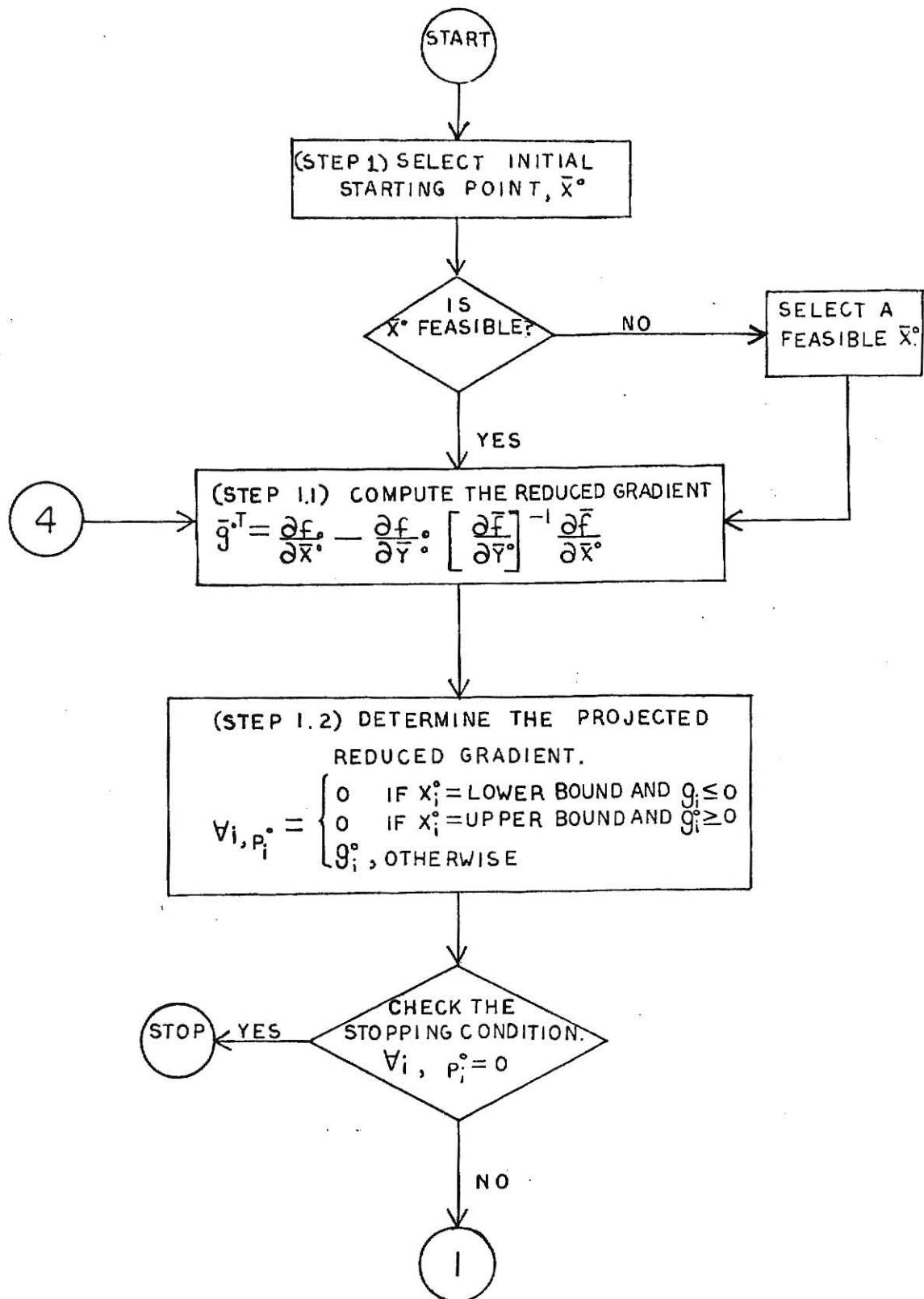


Fig. 3.1. Computer flow diagram for the GRG algorithm.

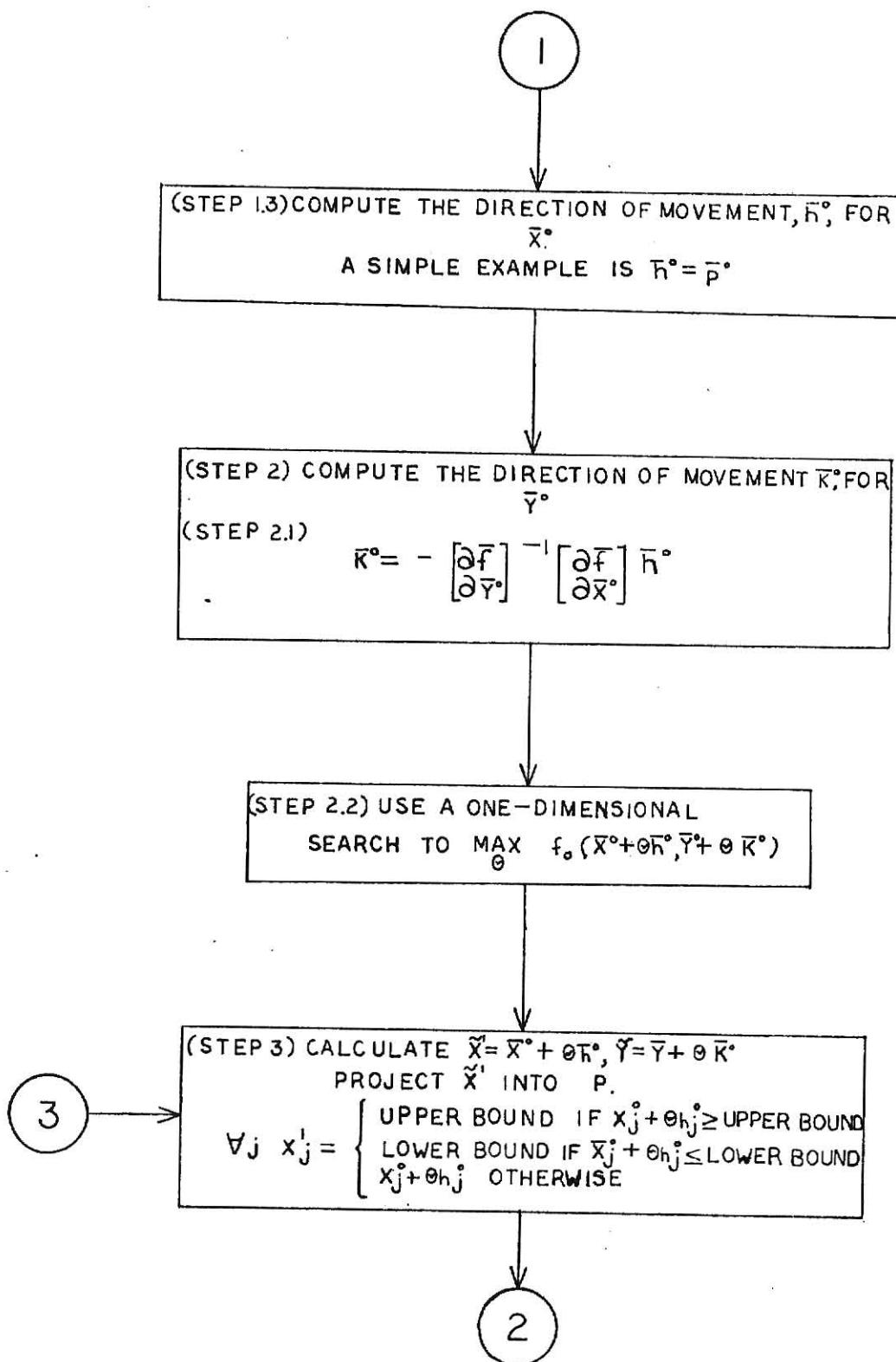


Fig. 3.1. (continued)

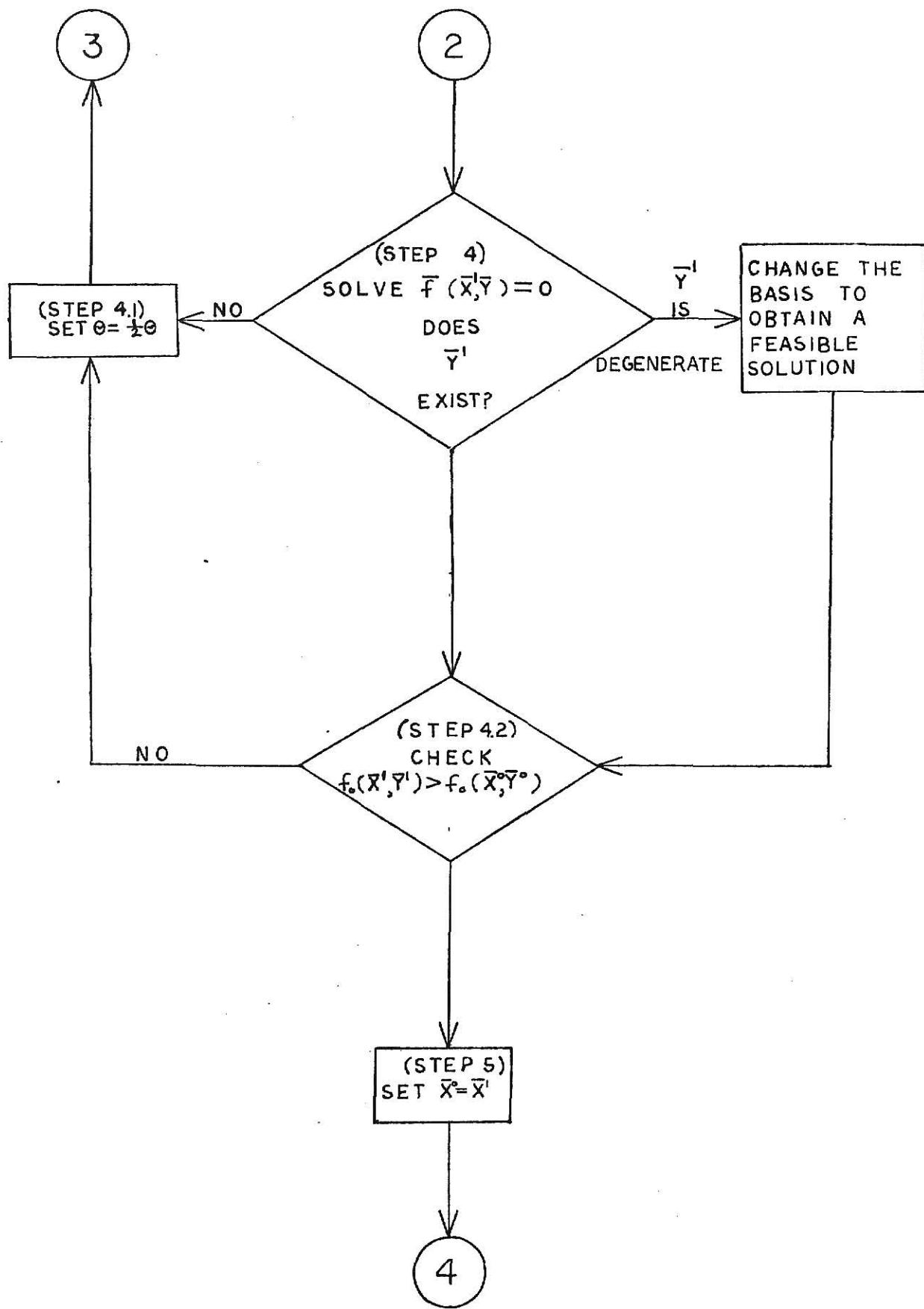


Fig. 3.1. (continued)

The artificial variables are forced to zero by the penalty, then disregarded.

The two-phase method occurs in two steps. The first step is to maximize

$$-\sum_{i \in NF} x_{N+i}, \quad NF = \{i \mid f_i(\bar{x}) \neq 0\}$$

When all of the artificial variables are zero the solution, \bar{x} , is feasible for $\bar{f}(\bar{x}) = 0$. In the second step the optimization problem continues normally. If it is impossible to force the artificial variables to zero values then a feasible solution may not exist. Both of the above techniques are similar to linear programming methods.

Theoretically, the stopping condition for the GRG algorithm is when

$p_i^0 = 0$, $i = 1, \dots, N-M$. This is not possible when using a finite number of digits for a real number. When convergence occurs, three appropriate stopping criteria are

$$1. \quad || \bar{p}^0 || = \sqrt{\sum_{i=1}^{N-M} (p_i^0)^2} < \epsilon_1$$

$$2. \quad p_i^0 < \epsilon_2, \quad i = 1, \dots, N-M$$

$$3. \quad |f_0(\bar{x}^1) - f_0(\bar{x}^0)| < \epsilon_3$$

3.6 A NUMERICAL EXAMPLE

A numerical example is solved for illustrating the GRG algorithm in detail.

Example 3

Maximize

$$f_0(\bar{x}) = (2x_1 - \frac{1}{2}x_1^2) + (3x_2 - \frac{1}{2}x_2^2)$$

subject to the constraints

$$f_1(\bar{x}) = x_1^2 + x_2^2 + x_3 - 1 = 0$$

$$x_i \geq 0, \quad i = 1, 2, 3$$

Note that the variable x_3 is essentially a slack variable for an inequality constraint.

The problem will be solved using the GRG algorithm with the exceptions that the value of the dependent variable will be determined analytically instead of using an iterative method, and the optimization of θ will be accomplished by direct differentiation. Table 3.1 summarizes the following numerical calculations and Figures 3.2 and 3.3 show the results graphically.

The independent variables, \bar{x} , and the dependent variable, \bar{y} , are chosen as

$$\bar{x} = \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix}, \quad \bar{y} = [x_2]$$

The reduced gradient for the given set of independent variables is (refer to equation (1))

$$\frac{\partial F}{\partial x} = \left(\frac{x_2(2-x_1) - x_1(3-x_2)}{x_2}, \frac{x_2 - 3}{2x_2} \right) \quad (6)$$

The formula for \bar{k} is (refer to equation (5))

$$\bar{k} = \left(-\frac{x_1}{x_2}, -\frac{1}{2x_2} \right) \bar{h}$$

Maximization of $f_0(\bar{x})$ is accomplished by direct differentiation. Let,

$$x_1 = x_1^0 + \theta d_1$$

		STEP 1	STEP 2	STEP 3	STEP 4	STEP 5
$\bar{X}^0 = \begin{bmatrix} 0.5 \\ 0.5 \\ 0.5 \end{bmatrix}$	$\bar{Y} = \begin{bmatrix} X_1 \\ X_2 \\ X_3 \end{bmatrix}$	$\frac{\partial F}{\partial \bar{X}}$	$\bar{h}^0 = \bar{p}^0$	\bar{k}^0	Θ	\bar{x}^1 $= \bar{x}^0 + \theta \bar{h}^0 = \bar{y}^0 + \theta \bar{k}^0$ $\bar{x}^1 \rightarrow \bar{x}^1$
1	$\begin{bmatrix} .5 \\ .5 \end{bmatrix}$	$\begin{bmatrix} 5 \\ 1 \end{bmatrix}$	$\begin{bmatrix} -1 \\ -2.5 \end{bmatrix}$	$\begin{bmatrix} 3.5 \\ 0 \end{bmatrix}$.547	$\begin{bmatrix} -.047 \\ -.867 \end{bmatrix}$ $\begin{bmatrix} 2.4 \\ 0 \end{bmatrix}$
2	$\begin{bmatrix} 0 \\ 0 \end{bmatrix}$	$\begin{bmatrix} 1 \\ 2, -1 \end{bmatrix}$	$\begin{bmatrix} 2 \\ 0 \end{bmatrix}$	$\begin{bmatrix} 0 \\ 1.0 \end{bmatrix}$	$\begin{bmatrix} 2 \\ 0 \end{bmatrix}$	$\begin{bmatrix} 1 \\ 0 \end{bmatrix}$
2.1				$.5$	$\begin{bmatrix} 1 \\ 0 \end{bmatrix}$	$\begin{bmatrix} 1 \\ 0 \end{bmatrix}$
2.2				$.25$	$\begin{bmatrix} .5 \\ 0 \end{bmatrix}$	$\begin{bmatrix} .5 \\ 0 \end{bmatrix}$
3	$\begin{bmatrix} .5 \\ 0 \end{bmatrix}$	$\begin{bmatrix} .816 \\ 2.77, -1.34 \end{bmatrix}$	$\begin{bmatrix} .277 \\ 0 \end{bmatrix}$	$\begin{bmatrix} .193 \\ 0 \end{bmatrix}$	$\begin{bmatrix} .554 \\ .849 \end{bmatrix}$ $\begin{bmatrix} .554 \\ 0 \end{bmatrix}$	$\begin{bmatrix} .834 \\ 0 \end{bmatrix}$

Table 3.1. Numerical results of Example 3.

$$x_2 = x_2^0 + \theta d_2$$

$$x_3 = x_3^0 + \theta d_3$$

where d_1 , d_2 , and d_3 are the directions of move for x_1^0 , x_2^0 , and x_3^0 .

$$\frac{df}{d\theta} = \frac{\partial f}{\partial x_1} \frac{dx_1}{d\theta} + \frac{\partial f}{\partial x_2} \frac{dx_2}{d\theta} = 0$$

or

$$\frac{df}{d\theta} = (2 - x_1^0) d_1 + (3 - x_2^0) d_2 = 0$$

$$\frac{df}{d\theta} = (2 - x_1^0 - \theta d_1) d_1 + (3 - x_2^0 - \theta d_2) d_2 = 0$$

This gives

$$\theta = \frac{(2 - x_1^0) d_1 + (3 - x_2^0) d_2}{(d_1^2 + d_2^2)} \quad (8)$$

Since x_2 is the dependent variable, $d_2 = k_1$, and for the independent variable, x_1 , $d_1 = h_1$.

Iteration 1

Step 1

The starting point is picked as $\bar{x}^0 = [0.5, 0.5, 0.5]$. Dividing this vector into independent and dependent variables gives

$$\bar{x}^0 = \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 0.5 \\ 0.5 \\ 0.5 \end{pmatrix}$$

$$\bar{y}^0 = [x_2] = [0.5]$$

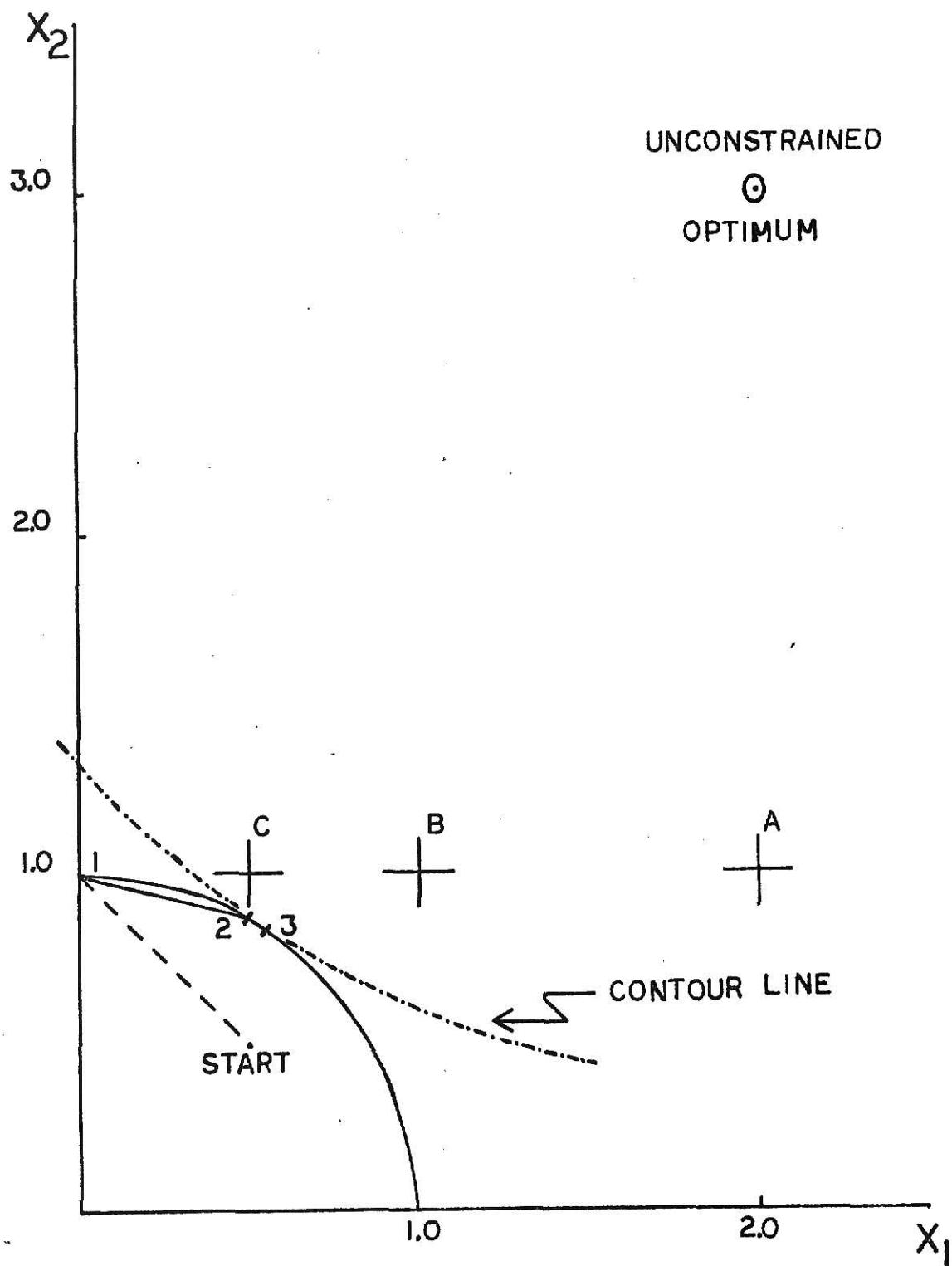


Fig. 3.2. Two-dimensional graph of the problem solving procedure of Example 3.

At this point the non-degeneracy of \bar{y}^1 is assumed.

Step 1.1

From equation (6)

$$\bar{g}^0 = \frac{\partial F}{\partial \bar{x}} = \left(\frac{0.5(2-0.5)-0.5(3-0.5)}{0.5}, \frac{0.5-3}{2(0.5)} \right) = [-1.0, -2.5].$$

Step 1.2

$$\bar{p}^0 = \begin{pmatrix} -1.0 \\ -2.5 \end{pmatrix}$$

Step 1.3

The direction of movement is along the reduced gradient in this example.

$$\bar{h}^0 = \bar{p}^0 = \begin{pmatrix} -1.0 \\ -2.5 \end{pmatrix}$$

Step 2

Step 2.1

The direction of movement for the dependent variable is in the direction of \bar{k} . From equation (7),

$$\bar{k}^0 = [-1.0, -1.0] \begin{pmatrix} -1.0 \\ -2.5 \end{pmatrix} = [3.5]$$

Step 2.2

The optimizing value for θ is calculated from equation (8).

$$\theta = \frac{(2.0-0.5)(-1) + (3.0-0.5)(3.5)}{(-1.0)^2 + (3.5)^2} = 0.547$$

Step 3

The values for \bar{x}^1 and \bar{y}^1 are

$$\bar{x}^1 = \begin{pmatrix} 0.5 \\ 0.5 \end{pmatrix} + 0.547 \begin{pmatrix} -1.0 \\ -2.5 \end{pmatrix} = \begin{pmatrix} -0.047 \\ -0.867 \end{pmatrix}$$

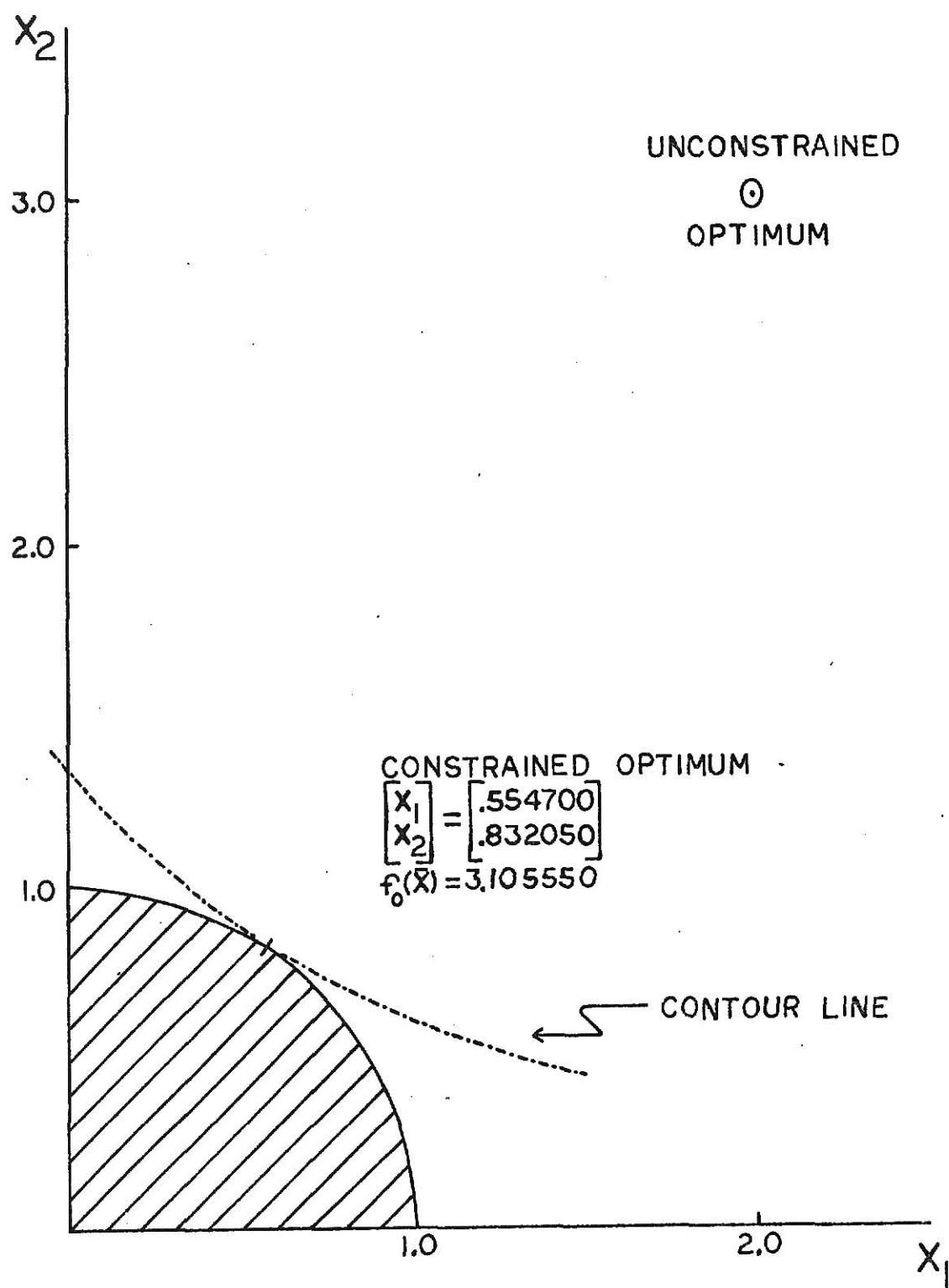


Fig. 3.3. Optimum solution of Example 3.

$$\bar{y}^1 = [0.5] + 0.547 [3.5] = [2.41]$$

The independent variables must be projected into the boundary set,
 $\bar{a} \leq \bar{x} \leq \bar{b}$. The projected independent vector is

$$\bar{x}^1 = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$$

Step 4

The dependent variable is solved for via the constraint, $f_1(\bar{x}) = 0$.

$$\bar{y}^1 = \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \left[\sqrt{1 - x_1^2 - x_3} \right]$$

$$\bar{y}^1 = \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = [1.0]$$

Step 4.1

This step is by-passed because convergence to \bar{y}^1 has occurred.

Step 4.2

In this step, improvement of the objective is checked.

$$f_0(\bar{x}^0) = 2.25 < 2.50 = f_0(\bar{x}^1)$$

Step 5

At this point, \bar{x}^1 becomes the starting point, \bar{x}^0 , for the second iteration.

$$\bar{x}^0 = \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}$$

Iteration 2

Since the slack variable, x_3 , is zero, the next problem step is shown graphically in Fig. 3.2.

$$\bar{x}^0 = \begin{pmatrix} 0 \\ 0 \end{pmatrix}, \quad \bar{y}^0 = [1]$$

Step 1.1

From equation (6),

$$\bar{g}^0 = \frac{\partial F}{\partial \bar{x}^0} = \begin{pmatrix} \frac{1(2-0)-0(3-1)}{1} & \frac{1-3}{2(1)} \end{pmatrix} = [2, 1].$$

Step 1.2

Since $\bar{x}_3^0 = 0$, the lower bound, and $\bar{g}_2^0 \leq 0$, the projected reduced gradient differs from \bar{g}^0 .

$$\bar{p}^0 = \begin{pmatrix} 2 \\ 0 \end{pmatrix}$$

Step 1.3

$$\bar{h}^0 = \bar{p}^0 = \begin{pmatrix} 2 \\ 0 \end{pmatrix}$$

Step 2

Step 2.1

From equation (7),

$$\bar{k} = [0, -1] \begin{pmatrix} 2 \\ 0 \end{pmatrix} = [0]$$

Step 2.2

From equation (8),

$$\theta = \frac{(2-0)(2) + (3-1)(0)}{(2^2 + 0)} = 1$$

Step 3

$$\bar{x}^1 = \begin{pmatrix} 0 \\ 0 \end{pmatrix} + 1 \begin{pmatrix} 2 \\ 0 \end{pmatrix} = \begin{pmatrix} 2 \\ 0 \end{pmatrix}$$

$$\tilde{y}^1 = [1] + 1 [0] = [1]$$

$$\bar{x}^1 = \tilde{x}^1$$

The boundary condition is satisfied with \tilde{x}^1 .

Step 4

A problem occurs at this point. A value for \bar{y}^1 does not exist. This is shown graphically in Fig. 3.2. Point A is located outside of the domain of $f_1(\bar{x})$.

Step 4.1

At this point, θ is reduced by one-half and the procedure returns to Step 3. This is referred to as Iteration 2.1 in Table 3.1.

Iteration 2.1

Step 3

$$\theta = 0.5$$

$$\tilde{x}^1 = \begin{pmatrix} 0 \\ 0 \end{pmatrix} + 0.5 \begin{pmatrix} 2 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$$\tilde{y}^1 = [1] + 0.5 [0] = [1]$$

$$\bar{x}^1 = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

The boundary condition is satisfied at this point. Refer to point B in Fig. 3.2.

Step 4

$$\bar{y}^1 = [x_2] = \sqrt{1 - x_1^2 - x_3^2}$$

$$\bar{y}^1 = [x_2] = [0]$$

Step 4.1

This step is by-passed because convergence to a feasible solution has occurred.

Step 4.2

The objective function has a value of 1.50 which is not an improvement over the previous solution, \bar{x}^0 . At this time reduce θ by one-half and return to Step 3. This will be referred to as Iteration 2.2 in Table 3.1.

Iteration 2.2Step 3

$$\theta = 0.25$$

$$\bar{x} = \begin{pmatrix} .5 \\ 0 \end{pmatrix}$$

Refer to point C in Fig. 3.2.

Step 4

Solving for the dependent variable yields

$$\bar{y}^1 = [0.816]$$

Step 4.1

This step is by-passed because convergence to a feasible solution has occurred.

Step 4.2

The value of the objective function improves to 3.10

Step 5

The next starting point is

$$\bar{x}^0 = \begin{pmatrix} 0.5 \\ 0.816 \\ 0 \end{pmatrix}$$

Iteration 3

This iteration continues smoothly in the same manner as the first iteration. Figure 3.2 indicates that convergence will be rapid. The solution is shown in Fig. 3.3 as

$$\bar{x} = \begin{pmatrix} \frac{2\sqrt{13}}{13} \\ \frac{3\sqrt{13}}{13} \\ 0 \end{pmatrix} = \begin{pmatrix} 0.554700 \\ 0.832050 \\ 0.0 \end{pmatrix}$$

The third iteration yields

$$\bar{x}^1 = \begin{pmatrix} 0.554 \\ 0.834 \\ 0.0 \end{pmatrix}$$

which is approaching the optimum. It is apparent that more than slide-rule precision is needed for a fourth iteration.

3.7 THE PSUEDO-NEWTON METHOD AND CHANGING THE BASIS

In general, numerical calculations are necessary for evaluating

$$\bar{y} = \bar{\phi}(\bar{x})$$

in step 4 of the GRG algorithm. It is desired that an iterative method produce a sequence, $\{\bar{y}^1, \dots, \bar{y}^t\}$, that satisfies the stopping condition

$$||\bar{f}(\bar{x}^1, \bar{y}^t)|| < \varepsilon$$

where the symbol, $||\bar{f}(\bar{x}^1, \bar{y}^t)||$, indicates the norm of the vector $\bar{f}(\bar{x}^1, \bar{y}^t)$. Upon convergence \bar{y}^t becomes \bar{y}^1 .

One of the most efficient procedures is the generalized Newton method.

The Taylor series expansion around the point \bar{y}^1 yields an approximate value for $\bar{f}(\bar{x}^1, \bar{y}^1)$ as

$$\bar{f}(\bar{x}^1, \bar{y}^1) \approx \bar{f}(\bar{x}^1, \bar{y}^1) + \frac{\partial \bar{f}}{\partial \bar{y}^1} (\bar{y}^1 - \bar{y}^1)$$

Since $\bar{f}(\bar{x}^1, \bar{y}^1) = 0$, the iterative formula becomes

$$\bar{y}^1 \approx \bar{y}^{i+1} = \bar{y}^1 - \left[\frac{\partial \bar{f}}{\partial \bar{y}^1} \right]^{-1} \bar{f}(\bar{x}^1, \bar{y}^1)$$

When the above approximation is good, it converges to the solution very rapidly. Sometimes it may not converge making it necessary to limit the maximum number of iterations. There are two reasons for nonconvergence. One is that the starting point is not close enough to the final solution. This is a problem caused by rapidly changing surface curvature, that is, the elements of the Jacobian, $\frac{\partial \bar{f}}{\partial \bar{y}^1}$, are "very" nonlinear. A second cause is that a real solution for \bar{y}^1 does not exist for the current values of the independent variables in \bar{x}^1 . Both of these problems are alleviated in step 4.1 of the GRG algorithm in which θ is reduced by 1/2. This causes the independent vector, \bar{x}^1 , and the starting point, \bar{y}^1 , to be closer to the feasible point $[\bar{x}^0, \bar{y}^0]$.

To determine \bar{y} , Abadie and Carpentier have proposed a modified Newton method and termed it the Pseudo-Newton method [2,5]. The method satisfies the boundary condition, P , by a logical procedure. When a degenerency occurs with a component, y_i^t , converging outside of P , the value of that component is set to its lower or upper bound. The degenerate, dependent component, y_i^t , is then interchanged with an independent component, x_j , thus changing the composition of the independent vector, \bar{x} , and the dependent vector, \bar{y} . The procedure has been termed by Abadie and Carpentier as changing the basis [2,5]. With the new basis the generalized Newton method is used to determine

a new value \tilde{y}^t .

Improvement of the objective function is not insured when a change of basis occurs during an iteration. Improvement is insured for values of θ less than a maximum value θ^* in the direction of the reduced gradient. But when a change of basis occurs the independent vector, \bar{x} , changes and is no longer related to the reduced gradient of the previous basis. For this reason improvement is not insured regardless of how many times θ has been reduced. Abadie has noted that the condition can cause the GRG 66 and GRG 69 codes to cycle [2]. He has studied the degeneracy problem extensively [1] and has implemented an anticycle procedure in the GREG code.

Example 4

Suppose a hypothetical objective function is subject to the constraints

$$x_1^2 + x_2^2 \leq 4$$

$$x_1^2 + x_3^2 \leq 1$$

$$x_i \geq 0 \quad i = 1, 2, 3$$

The constraints are pictured in Fig. 3.4 and, by using slack variables, have the form

$$x_1^2 + x_2^2 + x_4^2 = 4$$

$$x_1^2 + x_3^2 + x_5^2 = 1$$

$$x_i \geq 0, \quad i = 1, 2, 3, 4, 5.$$

Assume the i^{th} iteration yields (point 1 in Fig. 3.4)

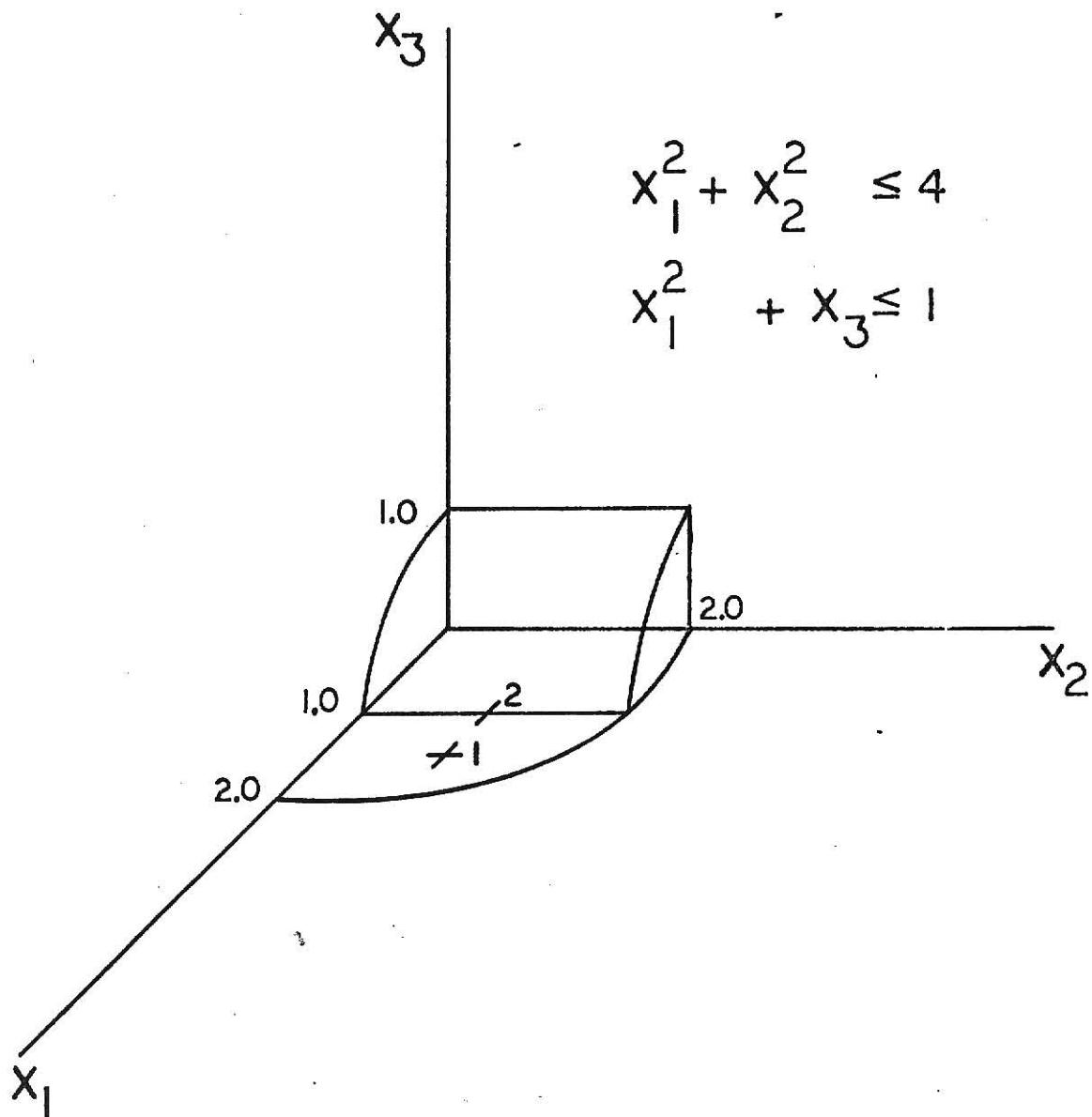


Fig. 3.4. Obtaining a feasible solution, point 2,
by changing the basis in Example 4.

$$\bar{x} = \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 3/2 \\ 1 \\ 0 \end{pmatrix}$$

and, by solving for \bar{y}

$$\bar{y} = \begin{pmatrix} x_4 \\ x_5 \end{pmatrix} = \begin{pmatrix} 3/4 \\ -5/4 \end{pmatrix}$$

The variable x_5 violates a boundary condition and this must be corrected before proceeding. Let $x_5 = 0$ and make it an independent variable in place of x_1 .

This yields

$$\bar{x} = \begin{pmatrix} x_2 \\ x_3 \\ x_5 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}$$

Solving for \bar{y} yields

$$\bar{y} = \begin{pmatrix} x_1 \\ x_4 \end{pmatrix} = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

a feasible solution (point 2 in Fig. 3.4). The new point becomes the next starting point only if the objective function is improved. This criterion depends on both the "shape" of the objective function and how "close" the point is to the point produced by the reduced gradient.

To prevent degenerate situations Abadie and Guigou [2,6] have recommended replacing a dependent variable, y_i , with an independent variable, x_j that maximizes the expression

$$\min \{ |\alpha_{ij}|(b_j - x_j), |\alpha_{ij}|(x_j - a_j) \}, \text{ for } \forall j$$

$$\alpha_{ij} = \left(\frac{\partial \bar{f}}{\partial y} \right)^{-1} \left[\frac{\partial \bar{f}}{\partial \bar{x}} \right]$$

The intention is to choose a value located far from its boundaries. This conjecture, though, has not been proven [2].

3.8 CONCLUDING REMARKS

The GRG algorithm is an attempt at the solution of the general continuous nonlinear programming problem. It is restricted to problems consisting of continuous constraints that have continuous first order partial derivatives. This restriction is a result of the Implicit Function Theorem cited earlier. The objective function is not as restricted in structure as the constraints. It is required to be continuous with the first order partial derivatives defined at each point. But, the first order partial derivatives do not necessarily have to be continuous if the objective function is monotonically increasing (or decreasing). This allows the use of a piecewise linear monotonically increasing (or decreasing) objective function.

The method is powerful because it possesses all of the convergence properties of the popular gradient techniques for unconstrained nonlinear optimization problems. Its main restriction is in step 4 of the algorithm. A method must exist for reentry into the feasible domain that converges rapidly, is reliable, and will control degenerate situations which are inherent in most problems. It is obvious that the re-entry procedure adds to the numerical difficulties that exist in large computer coded algorithms and that the GRG method is not easily coded and applied to large scale problems. These difficulties are largely technical, though, and can be alleviated with time and experience. The GRG method is theoretically a powerful technique and it appears that the only limitation to its application will be the computer code that represents it.

CHAPTER FOUR
INTRODUCTION TO THE GREG PROGRAM

4.1. INTRODUCTION

The generalized reduced gradient algorithm is coded in FORTRAN IV in the GREG program [30]. The program has been developed, by Abadie and his associates of Electricité de France. The GREG program consists of a main program, nine permanent or internal subroutines, and four user supplied, temporary or external subroutines. A typical source deck consists of approximately 2500 cards, requires approximately 150 K ($K = 1024$ bytes) of storage and 18 minutes of CPU time to compile by the IBM FORTRAN IV (G) compiler. For execution it requires, with the present dimensioning approximately 120 K of storage. The main program and the permanent subroutines have been compiled and stored in a partitioned data set. This simple step sharply reduces the CPU time, and consequently the cost. It also allows the program to be run in less than 128 K of storage.

As presently dimensioned the GREG program will handle problems involving up to 50 inequality and/or equality constraints.

The maximum number of variables a problem may have depends on the type of constraints involved. The program automatically provides slack and artificial variables. There is one slack variable added for each inequality constraint. For a given constraint, if the starting point, \bar{x}^0 , is not feasible, an artificial variable is supplied with an appropriate penalty attached to the objective function. Therefore, the dimension which represents the total number of variables a problem has depends on the following four numbers.

NV = the number of variables in the original problem.

NIN = the number of inequality constraints and, therefore, the number of slack variables added.

NIN1 = the number of inequality constraints not satisfied at the starting point and, therefore, the number of artificial variables added for this reason.

NIN4 = the number of equality constraints not satisfied at the starting point and, therefore, the number of artificial variables added for this reason.

The code is presently dimensioned with the following two constraints.

$$NV + NIN + NIN1 + NIN4 \leq 150$$

and

$$NV \leq 100$$

Use of the GREG program is approached in four steps. The first is developing the four user supplied, temporary subroutines. These subroutines, basically, define the problem. Next is organizing the input data. This step is extremely important when using the GREG program. The program is extremely sensitive to certain parameter values and they must be chosen carefully. The third step involves applying the necessary job control language. Along with the access procedures, the job control language is given that was used to load the program into the system. The title headings are in French and shall be translated as finally.

4.2. THE GREG USER SUPPLIED SUBROUTINES

The user supplied subroutines must define the optimization problem in the form of maximize

$$f_0(\bar{X}), \quad \bar{X} = (x_j \mid j = 1, \dots, NV)$$

subject to the constraints

$$f_i(\bar{X}) \leq \text{or } = 0, \quad i = 1, \dots, NC$$

$$a_j \leq x_j \leq b_j, \quad j = 1, \dots, NV$$

Any nonlinear programming problem may be put into this form. The four subroutines that describe this problem set-up to the GREG program are PHIX,

CPHI, JACOB, and GRADFI. PHIX defines the objective function, CPHI the constraint function, JACOB the gradients of the constraint functions, and GRADFI the gradient of the objective function. Each of the four subroutines performs a unique task. They are referred to many times during the execution of the program and warrant careful programming considerations.

Each of the user supplied subroutines must contain a set of "COMMON" statements that are commensurate to the internal GREG subroutines. Figure 4.1 shows a list of the GREG common block definition statements. An error in this portion of the programming can result in completely erroneous results.

The subroutines are called in the following order.

1. PHIX
2. CPHI
3. JACOB
4. GRADFI

With the variable, IT, which is the iteration counter, the subroutines can be used to initialize values used within them by "READ", "DATA", or arithmetic statements. When the subroutines are called for the first time, IT = -1. By checking the condition, if IT < 0, the initializations are performed on the first call to the subroutine. The programming for the following problem is shown in Figs. 4.2 through 4.5. The problem shall serve as an example in the description of the four external, user supplied subroutines.

Maximize

$$f_0(\bar{x}) = \sum_{j=1}^5 \ln (1 - \exp(B_j \cdot X_j))$$

subject to the constraints

```
DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)      GEGA  20
1,Y(150),C(150),VC(50),IBAS(50),IHB(100),IVC(50),IVA(100)      GEGA  30
COMMON   A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVA,IVB      GEGA  40
COMMON   NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NT0,NINI,NIN2,NIN3,NIGEGA  50
1N4,NVNINI,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCGEGA  60
2DB,KFIL,KLIN,KREN,KD,FI1,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEKA  70
COMMON   KFUNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEKA  80
IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21      GEGA  90
COMMON   IDIREC,DELTFI,ETA,JK0,LC,YSORT      GEGA 100
```

Fig. 4.1. GREG common block definition statements.

$$f_1(\bar{x}) = \sum_{j=1}^5 CC_{1j} \cdot x_j^2 \leq 110$$

$$f_2(\bar{x}) = \sum_{j=1}^5 CC_{2j} (x_j + \exp(x_j/4)) \leq 175$$

$$f_3(\bar{x}) = \sum_{j=1}^5 CC_{3j} \cdot x_j \exp(x_j/4) \leq 200$$

PHIX

The external, user-supplied subroutine PHIX defines the objective function to the GREG program. This value is stored in the FORTRAN variable PHI, and is described in terms of the FORTRAN vector array, XC(J). Only the original problem variables are used. That is, J ranges from one to NV. The penalties due to the artificial variables are added to PHI automatically in an internal subroutine.

If it is necessary to initialize constants used in defining PHI, the variable IT may be used as a logical indicator so constant values are initialized only once. Since PHIX is the first subroutine called, it may also be used to initialize constants used in CPHI, JACOB, and GRADFI as long as they have been declared in a common block definition statement. Figure 4.2 shows an example of PHIX.

CPHI

CPHI defines the constraint functions as previously defined (\leq or $= 0$). The values are stored in the vector array VC(I), I = 1, ..., NC, and in terms of the original problem variables, XC(J), J = 1, ..., NV. The constraints must be ordered with inequalities first and equalities second.

Let,

```

C001      SUBROUTINE PHIX
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)    GEGA 20
C003      1,Y(150),C(150),VC(50),IBAS(50),IB(1CC),IVC(5C),IVA(100)    GEGA 30
C004      CC*MON   A,PLFA,X,XC,XI,XS,Y,C,VC,IBAS,IBH,IVC,IVB,IVA,IVB    GEGA 40
C005      COMMON   NV,NC,NK,NEG,NIN,NTV,AV1,NEV,NEVL,NTO,NINI,NIN2,NIN3,NICEGA 50
C006      LN4,INVNIN1,INVNIN2,INVNIN3,INDEX,II,IR,IR1,IS,ISI,IT,IBP,ICCB,JCCB,KCGEGA 60
C007      2CB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSIL0,EPSIL2GEGA 70
C008      COMMON KFCNC,KGRAC,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEKA 80
C009      IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21                           GEGA 90
C010      COMMON ICIREC,DELIFI,ETA,JKO,LC,YSORT                           GEGA 100
C011      COMMON B(5),CR(3,5)
C012      IF(IIT)100,101,1C1
C013      100  B(1)=ALOG(0.2)
C014      B(2)=ALOG(0.15)
C015      B(3)=ALOG(0.1)
C016      B(4)=ALOG(0.35)
C017      B(5)=ALOG(0.25)
C018      1C1  PHI=C.C
C019      C01C2I=1,5
C020      102  PHI=PHI+ALOG(1.0-EXP(B(I1)*XC(I1)))
C021      RETURN
C022      END

```

Fig. 4.2. An example of PHIX.

NIN = the number of inequality constraints,

NEG = the number of equality constraints.

Then, VC(I), I = 1, ..., NIN, are inequality constraints and VC(I), I = NIN+1, ..., NIN+NEG, are equality constraints. The total number of constraints is

$$NC = NIN + NEG$$

Like PHIX, constants may be initialized in CPHI by using IT as a logical indicator. Since CPHI is the second procedure call by GREG, it may be used to initialize values in JACOB and GRADFI. An example of CPHI is shown in Fig. 4.3.

JACOB

The subroutine JACOB defines the gradients of the constraint functions. The partial derivative $\frac{\partial f_i}{\partial x_j}$ is stored in the matrix array A(i,j). The rows of the matrix represent each constraint function, $f_i(\bar{X})$, i = 1, ..., NC, in the same order as sequenced in CPHI. The partial derivatives are represented in terms of the FORTRAN variable XC(j), j = 1, ..., NV.

Constant values may also be initialized in JACOB. It is the third subroutine referred to and may be used to initialize values in GRADFI. An example of JACOB is shown in Fig. 4.4.

GRADFI

The fourth and final user-supplied subroutine is GRADFI. This subroutine defines the gradient of the objective function in terms of the array XC(J), J = 1, ..., NV. The component values are stored in the vector array C(J), J = 1, ..., NV.

Like the other subroutines, initialization may be accomplished in GRADFI, but only for this subroutine since it is the last one called for initialization purposes. Fig. 4.5 shows an example of GRADFI.

```

      SUBROUTINE CPHI
      DIMENSION A(50,100),ALFA(50,50),X(15C),XC(150),XI(150),XS(150)
      1,Y(150),C(150),VC(50),IBAS(5C),IMB(1CC),IVA(100),
      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IBB,IVC,IVA,IVB
      COMMON NV,NC,NK,NEG,NIN,NTV,NVI,NEVL,NTO,NINI,NINZ,NIN3,NICEGA
      1N4,NVNINI,NVNIN2,NVNINZ,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDR,KCGEGA
      2CB,KFIL,KLIN,KREN,KD,FI,PFI,PSI,PSI3,TB,TD,TC,EPSIL,EPSIL0,EPSIL2GEGA
      COMMON KFCNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJAC0,KMAX1,KMAX2,KGEWA
      IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21
      COMMON IDIREC,DELTF1,ETA,JKO,LC,YSORT
      COMMON B(5),CC(3,5)
      IF(IT)100,101,101
      100 CC(1,1)=1.0
      CC(2,1)=7.0
      CC(3,1)=7.0
      CC(1,2)=2.0
      CC(2,2)=7.0
      CC(3,2)=8.0
      CC(1,3)=3.0
      CC(2,3)=5.0
      CC(3,3)=8.0
      CC(1,4)=4.0
      CC(2,4)=9.0
      CC(3,4)=6.0
      CC(1,5)=2.0
      CC(2,5)=4.0
      CC(3,5)=9.0
      VC(1)=0.0
      VC(2)=0.0
      EC102J=1,5
      101 VC(1)=VC(1)+CR(1,J)*XC(J)**2
      VC(1)=VC(1)-110.0
      102 VC(1)=VC(1)+CC(2,J)*(XC(J)+EXP(XC(J)/4.0))
      VC(2)=VC(2)-175.0
      VC(3)=0.0
      EC104J=1,5
      103 VC(2)=VC(2)+CC(3,J)*(XC(J)+EXP(XC(J)/4.0))
      VC(2)=VC(2)-200.0
      RETURN
      END
      C001
      C002
      C003
      C004
      C005
      C006
      C007
      C008
      C009
      C010
      C011
      C012
      C013
      C014
      C015
      C016
      C017
      C018
      C019
      C020
      C021
      C022
      C023
      C024
      C025
      C026
      C027
      C028
      C029
      C030
      C031
      C032
      C033
      C034
      C035
      C036
      C037

```

Fig. 4.3. An example of CPHI.

```

C001      SUBROUTINE JACOB
C002      DIMENSION A(50,150),ALFA(50,50),X(150),XC(150),XS(150)      GFGA   20
C003      L,Y(150),C(150),VC(5C),IBAS(5C),IMB(1CC),IVC(5C),IV(10C)    GEGA   30
C004      COMMON A,ALFA,X,VC,XI,XS,Y,C,VC,IBAS,IMB,IVC,IVA,IVB      GEGA   40
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NIN1,NIN2,NIN3,NIGEGA 50
C006      LN4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,KCCEGA 60
C007      ZCB,KFIL,KLIN,KREN,KD,F1,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA 70
C008      COMMON KFCNC,KGRAC,KCONT,KINV1,KINV2,KCDBA,KJACQ,KMAX1,KMAX2,KGEGA 80
C009      IRENI,KREN2,KINV,KCDBA1,KREN1,KREN2,ICIREC,DELTFL,ETA,JKO,LC,YSORT 90
C010      COMMON B(5),CC(3,5)                                           GEGA   100
C011      COMMON E01C1J=1,5
C012      101 A(1,J)=2.0*CC(1,J)*XC(J)
C013      E0102J=1,5
C014      1C2 A(2,J)=CC(2,J)*(1.0+0.25*EXP(XC(J)/4.C))
C015      CC103J=1,5
C016      1C3 A(3,J)=CC(3,J)*EXP(XC(J)/4.0)*(1.0+0.25*XC(J))
C017      RETURN
C018      END

```

Fig. 4.4. An example of JACOB.

```

C001      SUBROUTINE GRADFI
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XI(150),XS(150)    GEGA   20
C003      1,Y(150),C(150),VC(150),IB(100),IVC(50),IVA(100)    GEGA   30
C004      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IB,IVC,IVA,IVB    GEGA   40
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEY,NEVL,NTC,NIN1,NIN2,NIN3,NICEGA 50
C006      1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IRI,IS,ISI,IT,IP,ICDB,JCDB,KCGEGA 60
C007      2CB,KFIL,KLIN,KREN,KD,F1,PHI,PSI,PSI3,TB,TC,EPSIL,EPSILO,EPSIL2GEGA 70
C008      COMMON KFNC,KGRAD,KCONT,KINV,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEKA 80
C009      COMMON KREN2,KINV,KCDBA1,KREN1,KREN21
C010      COMMON IGIREC,DELTFI,ETA,JKO,LC,YSORT    GEGA   90
C011      COMMON B(5),CC(3,5)
C012      E010 J=1,5
C013      AA=EXP(B(J)*XF(J))
C014      C(J)=B(J)*AA/(AA-1.C)
C015      RETURN
C016      END

```

Fig. 4.5. An example of GRADFI.

Since the constants used in PHIX will be related to those in GRADFI, and likewise for CPHI and JACOB, the initializations may usually be accomplished in PHIX and CPHI. To do this the constant variables are defined in "COMMON" statements so their values may be passed, but be cautious when doing this.

Note that the arrays A(I,J) and C(J) in the subroutines JACOB and GRADFI are initialized automatically at zero, thus eliminating the need for defining any zero values for them.

4.3. INPUT DATA FOR THE GREG PROGRAM

To use the GREG program, values for nineteen parameters, a starting point, a lower bound, and an upper bound must be established. The parameter input has been programmed in such a way that a minimum of two may be read in as input (NV and NIN or NEG) while the remaining ones take on default values. Besides this, the user may have designed some input for initialization purposes in PHIX, CPHI, JACOB, and GRADFI. It is the purpose of this section to discuss the above three types of input data.

The list of parameters and their definitions are given in Table 4.1. Each parameter is given a default value which is used if it is not changed in the parameter input list. Note that if NV is not given a value greater than zero the program will terminate without notice. The parameter input follows the rules of the IBM FORTRAN "NAMELIST" input procedure. This type of input is exemplified in Fig. 4.6. The GREG program is extremely sensitive to the values given to the parameters. It must be emphasized that the program is written in single precision arithmetic. This allows only seven significant digits and stopping criteria smaller than 1.0×10^{-7} can yield erroneous results.

Table 4.1. Parameters

<u>FORTRAN Program Symbols</u>	<u>Explanations</u>	<u>Default Value</u>
1. NV	the number of original problem variables.	0
2. NIN	the number of inequality constraints.	0
3. NEG	the number of equality constraints.	0
4. NEVL	the maximum number of iterations for Newton's Method.	20*
5. NTØ (NT zero)	the maximum number of bisections in the parabolic interpolation process when maximizing a function of a single variable.	6*
6. ITET	the number of previous iterations used to help determine a maximum value for Ø.	20*
7. ICONJ	equals 1 if conjugate directions are desired when maximizing, otherwise zero.	1
8. IDIAG	equals 1 if diagonal directions are desired when maximizing, otherwise zero.	0
9. ITMAX	the maximum number of iterations.	50
10. KFIL	equals 1 if the cost function is linear, otherwise, zero.	0
11. KLIN	equals 1 if <u>all</u> of the constraints are linear, otherwise, zero.	0
12. NCØ (NC zero)	equals zero for linear programming problems and is > the number of constraints for nonlinear programming problems.	10
13. ITSOR	the iteration which recording of the intermediate output starts. (1 ≤ ITSOR ≤ ITMAX)	1*

Table 4.1. (continued)

14.	ISOLSR	the solution is printed out every ISOLSR iterations. For small problems, ISOLSR = ITMAX. For large problems, ISOLSR < ITMAX.	50
15.	EPSIL	used as a criterion for the choice of a pivot in the changing and inversion of a basis. $(10^{-2} \leq EPSIL \leq 10^{-1})$	0.1E0*
16.	EPSIL0 (EPSIL zero)	is used as a stopping criteria for Newton's Method. ($EPSIL0 \geq 10^{-7}$)	0.1E-02
17.	EPSIL1	is used as a stopping criteria if the problem is declared convex. ($EPSIL1 \geq 10^{-7}$)	0.1E-02
18.	EPSIL2	is used as a stopping criterion by using the norm of the reduced gradient. ($EPSIL2 \geq 10^{-7}$)	0.1E-02
19.	PC	equals zero if the problem is non-convex; equals one if the problem is convex. This parameter affects only the stopping criteria.	0.0

* - it is recommended that these values are not changed.

```
C  
O  
L  
U  
M  
N  
123456789  
first card  
second card  
&PARM NV=10,NIN=3,NEG=2,PC=1.0,  
EPSILO=0.1E-4,EPSIL2=0.1E-03,&END
```

Fig. 4.6. An example of a FORTRAN namelist input for the GREG program. The title of the namelist is PARM.

After the parameter data, the starting point, the lower bound, and the upper bound are read in that order. The format is 8G10.4 which is equivalent to 8F10.4 and 8E10.4. This allows eight values per card. If a card is filled, skip to the next card and repeat the process until all NV values are listed. Each vector, the starting point, the lower bound, and the upper bound will start on a new card. The last data to enter the program is the optional, user defined initialization input. The data has to be ordered in the same calling sequence of the external subroutines PHIX, CPHI, JACOB, and GRADFI. Figure 4.7 shows the sequence of the input data.

4.4. JOB CONTROL LANGUAGE FOR UTILIZING THE GREG PROGRAM

The job control language given in this section is all that is necessary to utilize the GREG program at Kansas State University. Details can be found in the IBM FORTRAN IV (G and H) Programmers Guide. One can also consult the Kansas State University Computing Center User's Guide [12] for details on their procedures for preserving permanent data-sets.

There are two types of procedures of interest to the GREG user. They are the load and the access procedures. The two load procedures that have been used to install GREG into the system are given in Figs. 4.8, 4.9, and 4.10. Figures 4.11, 4.12, and 4.13 show the three access procedures for running a particular problem. It is recommended that object decks be used when several runs for the same problem are made. This will substantially reduce CPU time, and consequently, the cost.

4.5 INTERPRETING THE GREG OUTPUT

Output headings for the GREG program are in French. This section translates the headings and interprets the output. The language feature was not converted because of excessive costs to recompile the program.

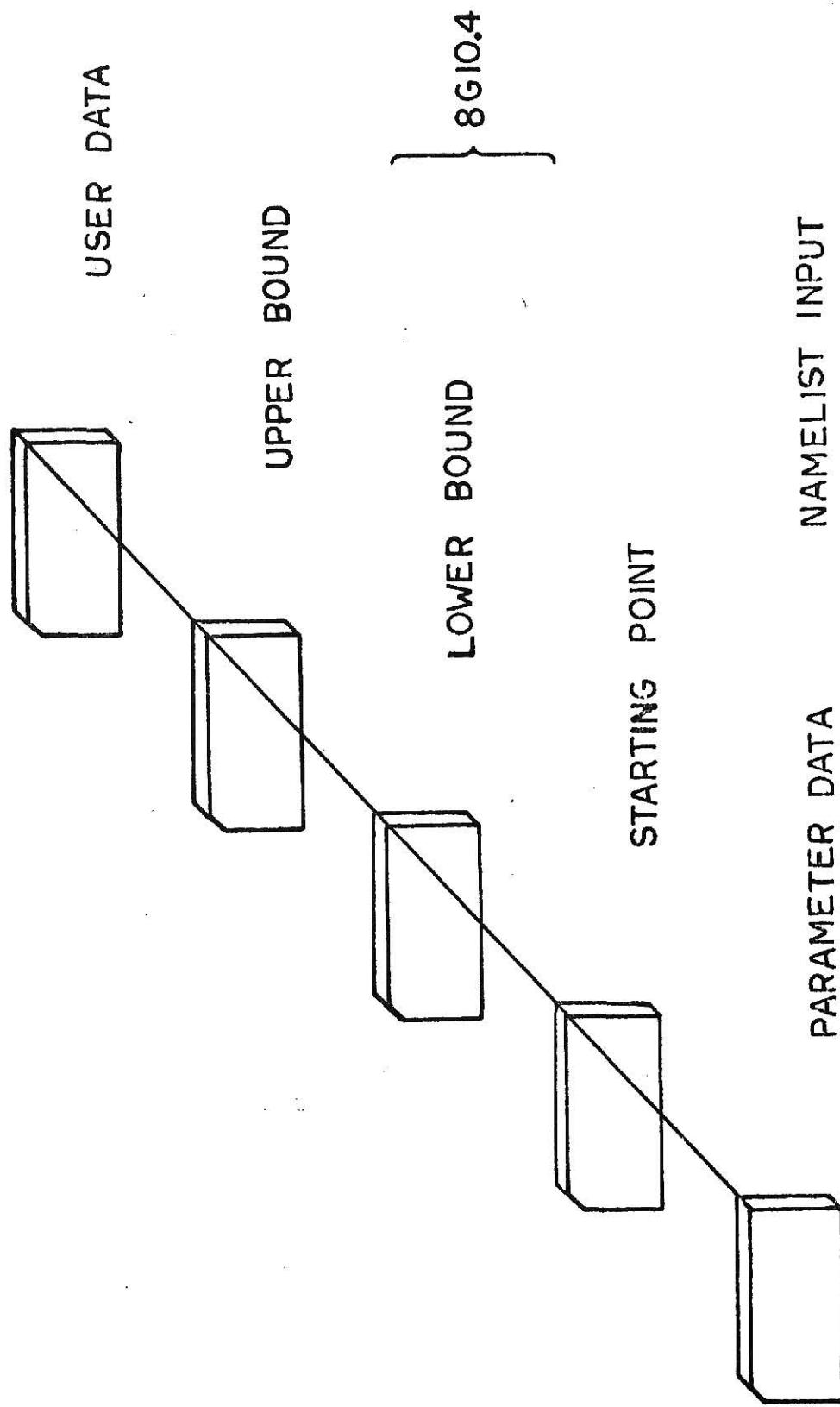


Fig. 4.7. The sequence of the input data deck.

```
//      JOB ( standard K.S.U. job card )
//      EXEC FORTGCLG
//FORT.SYSIN DD *

          complete GREG source deck including the subroutines
          PHIX, CPHI, JACOB, and GRADFI.

/*
//LKED.SYSLMOD DD DSN=yyyyyy.xxxxxxxx(name),DISP=(NEW,KEEP),UNIT=2314
//      VOL=SER=nnnnnn,SPACE=(TRK,(10,5,3),RLSE)
//GO.SYSIN DD *

        data

/*

```

Fig. 4.8. The necessary job control language for loading the GREG source deck into a partitioned data-set.

```
//      JOB ( standard K.S.U. job card )
//      EXEC FORTGLG
//LKED.SYSLMOD DD DSN=yyyyyy.xxxxxxxx(name),DISP=(NEW,KEEP),UNIT=2314
//      VOL=SER=nnnnnn,SPACE=(TRK,(10,5,3),RLSE)
//LKED.SYSIN DD *

      complete GREG object deck including the subroutines
      PHIX, CPHI, JACOB, and GRADFI.

/*
//GO.SYSIN DD *

      data

/*
```

Fig. 4.9. The necessary job control language for loading the GREG object deck into a partitioned data-set.

```
//          JOB ( standard K.S.U. job card )
//          EXEC FORTGCLG
//FORT.SYSIN DD *

        source deck

/*
//LKED.SYSLMOD DD DSN=yyyyyy.xxxxxxxx(name),DISP=(NEW,KEEP),UNIT=2314
//      VOL=SER=nnnnnn,SPACE=(TRK,(10,5,3),RLSE)
//LKED.SYSIN DD *

        object deck

/*
//GO.SYSIN DD *

        data

/*

```

Fig. 4.10. The necessary job control language for loading a combination of source and object decks into a partitioned data-set.

```
//          JOB ( standard K.S.U. job card )
//          EXEC FORTGCLG
//FORT.SYSIN DD *

                      source deck for PHIX, CPHI, JACOB, and GRADFI.

/*
//LKED.LIB DD DSN=yyyyyy.yxxxxxxxxx,DISP=SHR,UNIT=2314,VOL=SER=nnnnnn
//LKED.SYSIN DD *
      INCLUDE LIB(name)
      ENTRY MAIN
/*
//GO.SYSIN DD *

      data

/*
```

Fig. 4.11. Access to the stored program using a source deck for PHIX, CPHI, JACOB, and GRADFI.

```
//      JOB ( standard K.S.U. job card )
//      EXEC FORTGLG
//LKED.LIB DD DSN=yyyyyy.xxxxxxxx,DISP=SHR,UNIT=2314,VOL=SER=nnnnnn
//LKED.SYSIN DD *
                                         object deck for PHIX, CPHI, JACOB, and GRADFI.

INCLUDE LIB(name)
ENTRY MAIN
/*
//GO.SYSIN DD *

      data

/*

```

Fig. 4.12. Access to the stored program via an object deck.

```
//          JOB ( standard K.S.U. job card )
//          EXEC FORTGCLG
//FORT.SYSIN DD *

        source deck

/*
//LKED.LIB DD DSN=yyyyyy.xxxxxxxx,DISP=SHR,UNIT=2314,VOL=SER=nnnnnnn
//LKED.SYSIN DD *

        object deck

INCLUDE LIB(name)
ENTRY MAIN
/*
//GO.SYSIN DD *

        data

/*
```

Fig. 4.13. Access to the stored program via an object and a source deck for PHIX, CPHI, JACOB, and GRADFI.

Another feature of the output is that all real numbers are printed with fourteen significant digits because the program was developed on a machine with fifteen digit accuracy. Only the first seven digits are valid when using the IBM 360/50 computer. This feature was not changed because of the cost and plans to initiate a double precision version to handle the numerical difficulties inherent in some problems.

The output consists of five sections. They are

1. parameter echo check
2. starting conditions
3. intermediate output
4. final solution
5. summary of the results.

Table 4.2 translates the main French titles.

Of chief concern are the stopping conditions (NIVEAU DE SORTIE).

Numerically, the stopping condition codes are from zero to eight. The descriptions follow.

0. The program is terminated when the total number of iterations is greater than or equal to a specified maximum number of iterations. The number of iterations, IT, will exceed the maximum only during the course of a cycle of conjugate directions.
1. The program is terminated when the norm of the projected reduced gradient is less than or equal to ϵ_2 (EPSIL2) times the norm of the projected reduced gradient of the first iteration that has no artificial variables.

2. Let

$$n = \max. n_i, \forall i$$

$$n_i = |p_i (b_i - x_i)| \text{ if } p_i > 0$$

$$n_i = |p_i (a_i - x_i)| \text{ if } p_i < 0$$

b_i = the upper bound for an independent variable x_i .

a_i = the lower bound for an independent variable x_i .

The program is terminated when

$$n < \epsilon_3, \epsilon_3 = \frac{\epsilon_1}{n}, \epsilon_1 = \text{EPSILL}.$$

3. If the problem has been declared convex, then the program is terminated when

$$\Delta\phi < \epsilon_1, \epsilon_1 = \text{EPSILL}$$

$$\Delta\phi = \sum n_i, \forall i$$

4. Let $\bar{c}^T = \frac{\partial f_0}{\partial X}$, the gradient of the objective function and $\bar{y}^T =$ the direction of movement. This termination occurs when $c \cdot Y < 0$ (scalar product).

If \bar{Y} is a conjugate direction the iteration is repeated and \bar{Y} is set equal to \bar{C} . If the condition occurs again an inversion followed by a search for a new basis is performed. The program is terminated if $\bar{C} \cdot \bar{Y} < 0$ occurs another time.

5. If during the course of a re-entry into the feasible region the Psuedo-Newton method yields a point satisfying the constraints, but does not improve the function value, then the re-entry point is modified by choosing a θ^* such that

$$\theta^* = \max_{\text{int}} (\theta/10, \theta)$$

θ_{int} is an interpolated value between (zero and θ)

This condition is internal and is not printed. It is normally allowed to occur twice in succession after which an inversion with a new basis

is initiated. If this fails to correct the condition then the stopping condition 7 is enacted.

6. The program is terminated when

$$|f_0(\bar{x}^{\circ} + \theta\bar{h}^{\circ}, \bar{y}^{\circ} + \theta\bar{k}^{\circ}) - f_0(\bar{x}^{\circ}, \bar{y}^{\circ})| < \epsilon^* |f_0(\bar{x}^{\circ}, \bar{y}^{\circ})|$$

This stopping condition occurs only if the basis consists of slack variables.

7. The program is terminated when

$$|f_0(\bar{x}^{\circ} + \theta\bar{h}^{\circ}, \bar{y}^{\circ} + \theta\bar{k}^{\circ}) - f_0(\bar{x}^{\circ}, \bar{y}^{\circ})| < \epsilon^* |f_0(\bar{x}^{\circ}, \bar{y}^{\circ})|$$

but the basis does not consist of just slack variables as in condition 6. Before stopping, an inverse with a new basis is initiated after which the direction of movement is conjugate to the gradient. If this does not work then the direction is along the projected reduced gradient, after which, the program terminates. This condition is usually implemented by successive reductions of θ due to failure to re-enter the feasible region or improve the objective function.

8. The program terminates when

$$|f_0(\bar{x}^{\circ} + \theta\bar{h}^{\circ}, \bar{y}^{\circ} + \theta\bar{k}^{\circ}) - f_0(\bar{x}^{\circ}, \bar{y}^{\circ})| < \epsilon^* |f_0(\bar{x}^{\circ}, \bar{y}^{\circ})|$$

Unlike condition 7, no attempt to re-enter the feasible region has failed. An attempt to override the stopping condition is effected by selecting and inverting a new basis and taking the directions of movement as the projected reduced gradient.

The desired stopping conditions are one, two, and three. Conditions six, seven, and eight indicate that precision difficulties may exist. Attempting an improvement from these stopping conditions may be desired since their presence indicates premature stopping. There are two methods.

1. Try a new starting point since the route from the present starting point must be very "flat".

2. Subtract a constant from the objective function so $\epsilon^* | f(\bar{x}^*, \bar{y}^*)|$ will be smaller.

These methods are not guaranteed to work, but they have helped solve some problems (such as the Colville problem number 3 worked in Chapter 5). If the difficulty persists then the only alleviation is more precision.

$\epsilon^* = 1.0 \times 10^{-12}$ for GREG.

$\epsilon^* = 1.0 \times 10^{-6}$ for GREGR (a more realistic value for seven digit precision).

PARAMETRES	(parameters)
NV	5
KIN	6
NEG	0
NEVL	20
NTO	6
IYE1	20
ICONJ	1
ICIAIG	C
ITMAX	50
KFTL	0
KLIN	C
NCO	10
ITSQR	1
ISOLSR	10
EPSLL	0.1E 00
EPSTL0	0.1E-04
EPSTL1	0.1E-02
EPSTL2	0.1E-02
PC	0.C

Table 4.2. Translation of the five output sections of the GREG program.

GRADIENT REDUIT GENERALISE

		NCMBRE DE VARIABLES	NATURELLES	5 (number of primal variables)	
		NCMBRE TOTAL DE VARIABLES	11 (total number of variables)		
		NCMBRE DE CONTRAINTES	6 (number of constraints)		
EPSILON	CE	NEWTON	0.10CCE-C4 (stopping condition for Newton's method)		
EPSILON	TEST GRADIENT		0.10CCE-02 (stopping condition for gradient test)		
 FUNCTION ECONOMIQUE (objective function)					
0.1171875CCCCOCOE-C1					
BORNE INFERIEURE (lower bound)		VARIABLE NATURELLE (starting point)		BORNE SUPERIEURE (upper bound)	
x[1]	0.78000000000000E 02	x[1]	0.7861995117188E C2	x[1]	0.11290000000000E 03
x[2]	0.33000000000000E 02	x[2]	0.33439987182617E 02	x[2]	0.46000000000000E 02
x[3]	0.27000000000000E 02	x[3]	0.3106999206530E 02	x[3]	0.45000000000000E 02
x[4]	0.27000000000000E 02	x[4]	0.44099990844727E C2	x[4]	0.45000000000000E 02
x[5]	0.27000000000000E 02	x[5]	0.35219985961914E 02	x[5]	0.45000000000000E 02
 VALEUR DES CONTRAINTES (constraint values)					
C[1]	-C.21136474609275E 00				
C[2]	-C.917886352306E 02				
E[3]	-C.11C5728159414E 02				
C[4]	-C.86942718305859E 01				
C[5]	-C.4872711816406E 01				
C[6]	-C.12728881835938E 00				

Table 4.2. (continued)

IT 1	PHI	0.1181757E-125CCOE	0.3	DIR.GRAF.	NO	2	YN 0.35E 03	DELTAIFI 0.19E 04	ETA 0.14E 04	NCDB 0	NCN 2	NITN 2
IT 2	PHI	0.1182343150000E	0.3	DIR.CONJ.	LC	3				NCDB 0	NCN 2	
IT 3	PHI	0.1182890E-210000E	0.3	DIR.CONJ.	LC	3	YN 0.95E 03	DELTAIFI 0.57E 03	ETA 0.70E 03	NCDB 0	NCN 8	NITN 20
IT 4	PHI	0.12540035C625000E	0.3	DIR.GRAC.	NO	2	YN 0.83E 03	DELTAIFI 0.38E 02	ETA 0.38E 02	NCDB 0	NCN 2	NITN 7
IT 5	PHI	0.25632C31250.000E	0.3	DIR.GRAC.	NC	3	YN 0.83E 03	DELTAIFI 0.36E 02	ETA 0.36E 02	NCDB 0	NCN 3	NITN 12
IT 6	PHI	0.29201171875000E	0.3	DIR.GRAC.	NO	4	YN 0.83E 03	DELTAIFI 0.43E-02	ETA 0.43E-02	NCDB 0	NCN 2	
	PHI	0.29201171875000E	0.3				YN 0.81E 03					

CUREE CU CALCUL 165 CENTISECONDES
(execution time)

NO = number of variables at their lower or upper bounds.

LC = number of conjugate cycles when using conjugate directions.

YN = norm of the reduced gradient.

DELTAIFI = $\Delta\phi$ as defined by stopping condition 3.

ETA = n as defined by stopping condition 2.

NCDB = number of changes of basis during the iteration.

NCN = number of calls to the Newton method.

NITN = number of iterations of the Newton method during the course of NCN calls.

Table 4.2. (continued)

VARIABLE NATURELLE (primal variable)	DUALE ASSOCIEE (associated reduced gradient value)
$x_1 = C.7800000CCC000E 02$	$x_1(1) = V(1) = -0.48917602339063E 02$
$x_1(2) = 0.3300000CCC000E 02$	$x_1(2) = V(2) = -0.84309921264648E 02$
$x_1(3) = C.299938564E375E 02$	$x_1(3) = V(3) = 0.$
$x_1(4) = 0.45000000CC0000E 02$	$x_1(4) = V(4) = 0.26632446289063E 02$
$x_1(5) = C.36776123C46875E 02$	$x_1(5) = V(5) = 0.*$

CONTRAINTE (constraint)	DUALE ASSOCIEE (Lagrange multiplier)
CCNTRAINTE 1	$U(1) = -C.152587890625C0E-04$
CCNTRAINTE 2	$U(2) = -C.91999984741211E 02$
CCNTRAINTE 3	$U(3) = -C.11158248901367E 02$
CCNTRAINTE 4	$U(4) = -C.68417510986378E 01$
CCNTRAINTE 5	$U(5) = -C.499998474121C9E 01$
CCNTRAINTE 6	$U(6) = -C.152587890625C0E-04$

Table 4.2. (continued)

NIVEAU	DE	SORTIE	7	(stopping condition)
NCMBRE	C	ITERATIONS	6	(number of iterations)
NCMBRE	C APPELS	CU S/P	CALCULANT	LES CONTRAINTES 93 (number of times the constraints were calculated)
NCMBRE	D APPELS	DL S/P	CALCULANT	LA FCNCTION ECONOMIQUE 69 (number of times the objective function was calculated)
NCMBRE	C APPELS	DU S/P	CALCULANT	LE GRADIENT DE LA FONCTION ECONOMIQUE 14 (number of times the gradient was calculated)
NCMBRE	C APPELS	DL S/P	CALCULANT	L INVERSE AVEC BASE DONNEE 4 (number of inverse calculations with the given base)
NCMBRE	D APPELS	DL S/P	CALCULANT	L INVERSE AVEC RECHERCHE DE BASE 1 (number of inverse calculations with another base)
NCMBRE	D APPELS	DU S/P	EFFECTUANT LES CHANGEMENTS DE BASE 3 (number of base changes)	
NCMBRE	C APPELS	DU S/P	CALCULANT LE JACOBIEN DES CONTRAINTES 8 (number of times the Jacobien of the constraints was calculated)	
NCMBRE	D APPELS	DL S/P	EFFECTUANT UNE RECHERCHE DE MAX SANS DICHOTOMIE 14 (number of maximums found without bisection)	
NCMBRE	C APPELS	DU S/P	EFFECTUANT UNE RECHERCHE DE MAX AVEC DICHOTOMIE 0 (number of maximums found by bisection)	
NCMBRE	C APPELS	CU S/P	EFFECTUANT LA RENTREE DANS LE DOMAINE 26	NUMBER OF ITERATIONS 74 (total number of Newton's iterations)

Table 4.2. (end)

CHAPTER FIVE

NUMERICAL TESTING OF THE GREG PROGRAM

5.1 INTRODUCTION

Because this study was the first implementation of the GREG code at Kansas State University it was necessary to test the program and discover any limitations it might have. For this reason a series of numerical comparisons were performed. The problems involved were of three categories. There were simple example problems, small scale problems from recent publications that were solved by other techniques, and Colville test problems. The Colville study [10] involved both GRG 66 and GRG 69, which compared favorably with other existing nonlinear programming codes. The GREG version was derived from both of these codes and should be superior in all aspects.

The GREG code was developed on the CDC 6600 computer [30] which has fifteen digit single precision arithmetic in its FORTRAN compiler. It must be emphasized that this is a distinct advantage over the IBM system 360/50 with seven digit single precision arithmetic. This means that the code, which is very sensitive to numerical precision, can behave quite differently in the IBM system 360/50 .

In performing the numerical study it was hoped that the reliability, limitations, accuracy, and speed of the program could be judged. All of the pertinent programming material is given in Appendix 4. The speed of the program is measured as the time of execution of the algorithm. It does not include compiling of the subroutines and initialization of the program. Initialization involves, basically, the adding of slack and artificial variables.

5.2 THE EXAMPLE PROBLEM OF CHAPTER THREE

The first problem worked was Example 3 of Chapter 3. The problem definition is maximize

$$f_0(\bar{x}) = (2x_1 - \frac{1}{2}x_2^2) + (3x_2 - \frac{1}{2}x_2^2)$$

subject to the constraint

$$f_1(\bar{x}) = x_1^2 + x_2^2 - 1 \leq 0$$

$$x_i \geq 0, i = 1, 2, 3$$

The iterative process in the GREG program was different from the example because it used the slack variable as the dependent variable in the starting basis. The solution was obtained in 0.5 seconds as

$$f_0 = 3.105550$$

$$x_1 = 0.5545919$$

$$x_2 = 0.8321228$$

5.3 RELIABILITY OF A COMPLEX SYSTEM

The following model was proposed by Tillman, Hwang, Fan, and Lai [35,44] in the study of the reliability of spacecraft life support systems. Two optimization problems were considered from it. The first was to maximize the system reliability subject to a weight constraint. The second was to minimize the system weight subject to a reliability constraint. The problems were optimized by the sequential unconstrained minimization technique [25]. The codes employed were the RAC program [38] and a code employing the Hooke and Jeeves pattern search [33,35]. The second code was superior to the RAC program yielding a computational time for both problems of 90.4 seconds [35].

A schematic diagram of the system is shown in Fig. 5.1 [35,44]. Each of the numbered components is assumed to have a reliability of R_i , $i = 1, 2, 3, 4$. Assuming that component three does not fail the system unreliability is

$$Q_1 = R_3 \left((1-R_1) (1-R_4) \right)^2$$

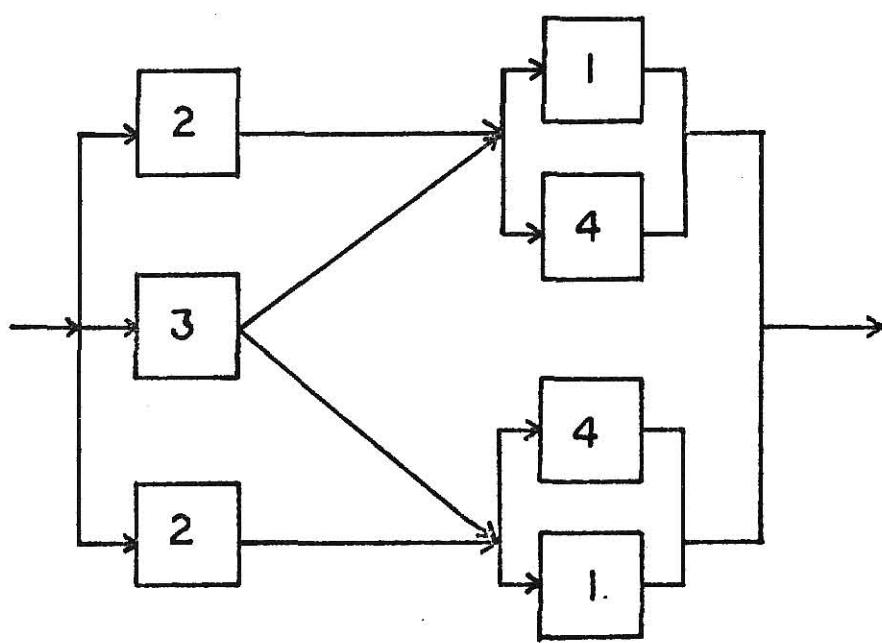


Fig. 5.1. Schematic diagram of a complex system.

If component three does fail then the probability of system failure becomes

$$Q_2 = (1 - R_3) (1 - R_2) (1 - (1 - R_1)(1 - R_4)))^2$$

Since the total reliability of the system is the complement of the total unreliability, the reliability is

$$R_s = 1 - Q_1 - Q_2$$

The weight of the system was given as a function of the component reliability, R_i , $i = 1, 2, 3, 4$.

$$\text{weight} = \sum_{i=1}^4 K_i R_i^{\alpha_i}$$

The resultant optimization problems are

maximize

$$R_s = 1 - R_3 ((1-R_1)(1-R_4))^2 - (1-R_3)(1-R_2) (1- (1 - R_1)(1 - R_4)))^2$$

subject to the constraint

$$K_1 R_1^{\alpha_1} + K_2 R_2^{\alpha_2} + K_3 R_3^{\alpha_3} + K_4 R_4^{\alpha_4} \leq C$$

and minimize

$$\text{weight} = K_1 R_1^{\alpha_1} + K_2 R_2^{\alpha_2} + K_3 R_3^{\alpha_3} + K_4 R_4^{\alpha_4}$$

subject to the constraints

$$1 - R_3 ((1-R_1)(1-R_4))^2 - (1-R_3)(1-R_2) (1- (1-R_1)(1-R_4)))^2 \geq R_s, \min$$

$$R_i \geq R_i, \min, i = 1, 2, 3, 4$$

Numerical values for the parameters were given as

$K_1 = 100, K_2 = 100, K_3 = 200, K_4 = 150$
 $C = 800, R_{s,\min} = 0.9$
 $\alpha_i = 0.6, R_{i,\min} = 0.5, i = 1, 2, 3, 4$

Numerical results for the GREG code were as follows.

First Problem

$R_s = 1.0$
 $R_1 = 1.0$
 $R_2 = 1.0$
 $R_3 = 0.5394473$
 $R_4 = 0.7974168$

Execution time = 0.35 seconds

Physically, the results may be interpreted as investing all effort into components one and two to make their reliability as high as possible. The best optimum solution by the previous methods was $R_s = 0.999997$.

Second Problem

weight = 641.8232
 $R_1 = 0.5$
 $R_2 = 0.8389202$
 $R_3 = 0.5$
 $R_4 = 0.5$

Execution time = 0.48 seconds

The best previous solution was

weight = 642.249

Total computational time for both problems was less than a second.

5.4 OPTIMUM RELIABILITY OF A MULTISTAGE PARALLEL SYSTEM

The problem in this section was optimized in Chapter Two using separable programming. It was obtained from a study by Tillman, Hwang, Fan, and Balbale [45]. It involves the optimization of a mixed system with N stages in series with components at each stage in parallel. The system is shown schematically in Fig. 5.2. The problem is to choose the number of parallel components at each stage that maximizes the systems reliability subject to a weight constraint. In reality, this problem should have an integer solution [43], but this point was neglected in the following solution since it is desired to use the example as a test problem only.

For a given stage, i , if R_i is the reliability of the components involved then

$$(1 - R_i)^{M^i}$$

is the unreliability of M^i components in parallel. The reliability is the complement of the unreliability and for N stages the total reliability is

$$R_s = \prod_{i=1}^N (1 - (1 - R_i)^{M^i})$$

A constraint involving weight and volume was given as

$$\sum_{i=1}^N p^i (M^i)^2 \leq P$$

The value of p^i is the product of weight per unit and volume per unit at stage i .

The cost constraint was given as

$$\sum_{i=1}^N c^i (M^i + \exp(M^i/4)) \leq C$$

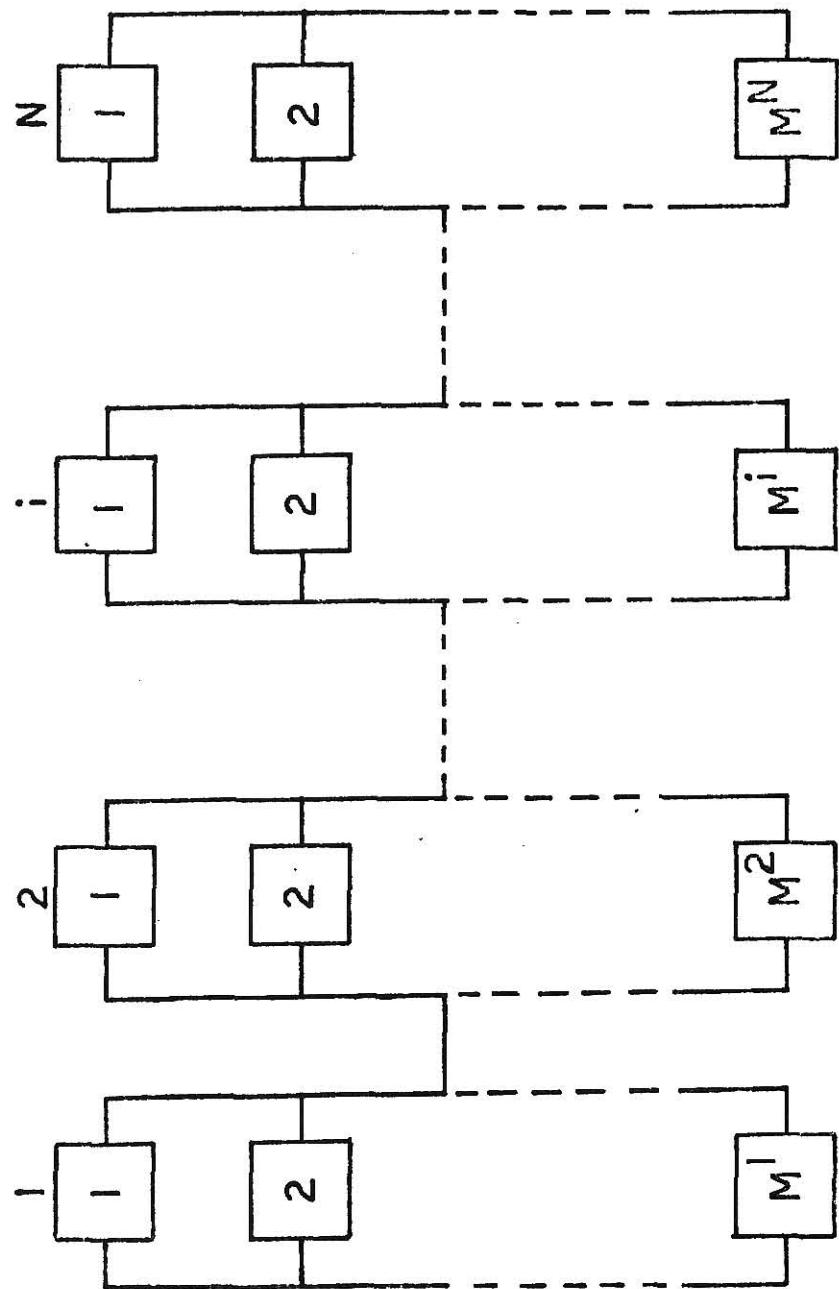


Fig. 5.2. Schematic diagram of a multistage parallel system.

The value $c^i M^i$ represents the cost of the units at stage i and $c^i \exp(M^i/4)$ is the additional cost for interconnecting parallel units.

The weight constraint was given as

$$\sum_{i=1}^N w^i M^i \exp(M^i/4) \leq W$$

The value $w^i M^i$ is the weight of the units at stage i and the factor $\exp(M^i/4)$ allows for the weight of interconnecting the units.

The problem was easier to work by maximizing $\ln(R_s)$.

This makes the problem separable. The parameter values used were

$$\begin{array}{llll} R_1 = 0.80, & p^1 = 1.0, & c^1 = 7.0, & w^1 = 7.0 \\ R_2 = 0.85, & p^2 = 2.0 & c^2 = 7.0 & w^2 = 8.0 \\ R_3 = 0.90, & p^3 = 3.0 & c^3 = 5.0 & w^3 = 8.0 \\ R_4 = 0.65 & p^4 = 4.0 & c^4 = 9.0 & w^4 = 6.0 \\ R_5 = 0.75 & p^5 = 2.0 & c^5 = 4.0 & w^5 = 9.0 \end{array}$$

The problem is (changing M^i to x_i)

maximize

$$\begin{aligned} f_0 = \ln(1-.2x_1) + \ln(1-.15x_2) + \ln(1-.10x_3) + \ln(1-.35x_4) \\ + \ln(1-.25x_5) \end{aligned}$$

subject to the constraints

$$x_1^2 + 2x_2^2 + 3x_3^2 + 4x_4^2 + 2x_5^2 \leq 110$$

$$7(x_1 + e^{x_1/4}) + 7(x_2 + e^{x_2/4}) + 5(x_3 + e^{x_3/4}) + 9(x_4 + e^{x_4/4})$$

$$\begin{aligned} &+ 4(x_5 + e^{x_5/4}) \leq 175 \\ 7x_1 e^{x_1/4} + 8x_2 e^{x_2/4} + 8x_3 e^{x_3/4} + 6x_4 e^{x_4/4} + 9x_5 e^{x_5/4} \leq 200 \end{aligned}$$

The solution was determined as

$$x_1 = 2.677995$$

$$x_2 = 2.352543$$

$$x_3 = 2.070096$$

$$x_4 = 3.530631$$

$$x_5 = 2.792024$$

Execution time = 2.86 seconds

This is very close to the results obtained by the separable programming method in Chapter Two. The results of the separable programming method (Chapter Two) with a grid size of 0.1 were

$$x_1 = 2.70000$$

$$x_2 = 2.32929$$

$$x_3 = 2.10000$$

$$x_4 = 3.50000$$

$$x_5 = 2.80000$$

and the results of solution by the discrete maximum principle were [45]

$$x_1 = 2.6000$$

$$x_2 = 2.2816$$

$$x_3 = 2.0075$$

$$x_4 = 2.6882$$

$$x_5 = 3.3981$$

5.5 TWO PROBLEMS FROM THE COLVILLE STUDY

The next two problems were taken from "A Comparative Study of Nonlinear Programming Codes", by A.R. Colville [10]. The report is often referred to as the Colville study. The purpose of the study was to compare and evaluate

existing nonlinear programming codes.

The first problem considered is a modified version of test problem number two in the study. It was modified by Guigou [30] to exemplify usage of the GREG code. It is presented in that form. The problem is to maximize

$$f_0(\bar{x}) = \sum_{i=1}^{10} b_i x_{i+5} - \sum_{i=1}^5 \sum_{j=1}^5 c_{ij} x_i x_j - 2 \sum_{j=1}^5 d_j x_j^3$$

subject to the quadratic constraints

$$\sum_{j=1}^{10} a_{i,j+5} x_{j+5} - e_i - 2 \sum_{j=1}^5 c_{ji} x_j - 3d_i x_i^2 \leq 0$$

$$i = 1, 2, 3$$

$$\sum_{j=1}^{10} a_{i,j+5} x_{j+5} - e_i - 2 \sum_{j=1}^5 c_{ji} x_j - 3d_i x_i^2 + x_{i+12} = 0$$

$$i = 4, 5$$

$$0 \leq x_k \leq 100, \quad k = 1, \dots, 17$$

Table 5.1 gives the parameter values. The problem was defined in the Colville study with five inequality constraints. The slack variables have been added to the last two constraints (x_{16} and x_{17}) to exemplify the GREG code when a nonfeasible starting point is encountered with both inequality and equality constraints.

The solution obtained was

$$f_0 = -32.34867$$

$$x_1 = 0.2994560$$

$$x_2 = 0.3332918$$

$$x_3 = 0.4006066$$

i	1	2	3	4	5
d_i	4	8	10	6	2
e_i	-15	-27	-36	-18	-12

i	1	2	3	4	5	6	7	8	9	10
b_i	-40	-2	-1/4	-4	-4	-1	-40	-60	5	1

Value of c_{ij} :

i \ j	1	2	3	4	5
1	30	-20	-10	32	-10
2	-20	39	-6	-31	32
3	-10	-6	10	-6	-10
4	32	-31	-6	39	-20
5	-10	32	-10	-20	30

Table 5.1 Parameter values for the modified Colville test problem two.

Value of a_{ij} :

$i \backslash j$	1	2	3	4	5
1	0	0	0	0	0
2	0	0	0	0	0
3	0	0	0	0	0
4	0	0	0	0	0
5	0	0	0	0	0
6	-16	2	0	1	0
7	0	-2	0	0.4	2
8	-3.5	0	2	0	0
9	0	-2	0	-4	-1
10	0	-9	-2	1	-2.8
11	2	0	-4	0	0
12	-1	-1	-1	-1	-1
13	-1	-2	-3	-2	-1
14	1	2	3	4	5
15	1	1	1	1	1

Table 5.1 (continued)

$$x_4 = 0.4287636$$

$$x_5 = 0.2243463$$

$$x_6 = x_7 = 0.0$$

$$x_8 = 0.5175596$$

$$x_9 = 0.0$$

$$x_{10} = 0.3081506$$

$$x_{11} = 0.1183264$$

$$x_{12} = x_{13} = 0.0$$

$$x_{14} = 0.1027765$$

$$x_{15} = x_{16} = x_{17} = 0.0$$

Execution time = 12.88 seconds

The best solution in the Colville study produced an objective function value of -32.349.

The second problem is test problem number three. It involves a quadratic objective function with six quadratic constraints. The problem is defined as minimize

$$f_0(\bar{x}) = 5.3578547 x_3^2 + 0.8356891 x_1 x_5 + 37.293239 x_1, - 40792.141$$

subject to

$$85.334407 + 0.0056858 x_2 x_5 + 0.0006262 x_1 x_4 - 0.0022053 x_3 x_5 \underset{92}{\leq}$$

$$80.51249 + 0.0071317 x_2 x_5 + 0.0029955 x_1 x_2 + 0.0021813 x_3^2 \underset{110}{\leq} 90$$

$$9.300961 + 0.0047026 x_3 x_5 + 0.0012547 x_1 x_3 + 0.0019085 x_3 x_4 \underset{25}{\leq} 20$$

$$78 \leq x_1, \leq 102$$

$$33 \leq x_2, \leq 45$$

$$27 \leq x_3, x_4, x_5 \leq 45$$

In this problem the optimum solution was not reached by the normal procedure. It was found that the code would stop too prematurely with stopping condition number seven. The problem was that the number of desired significant digits and the magnitude of the objective function required more precision than was available. To solve the problem a constant was subtracted from the objective function to scale it. The constant value was 30373.94. The final solution was then found to be

$$\begin{aligned}f_0 &= -(292.01 + 30373.94) = -30665.95 \\x_1 &= 78.00000 \\x_2 &= 33.00000 \\x_3 &= 29.99389 \\x_4 &= 45.00000 \\x_5 &= 36.77612\end{aligned}$$

Execution time = 1.65 seconds

The best solution given in the Colville study was -30665.5.

5.6 SOLUTION TO THE GEOMETRIC PROGRAMMING EXAMPLE OF CHAPTER TWO

For the purpose of comparison, the primal problem of the geometric programming example of Chapter Two was solved using the GREG code. The result was very close to the previous solution.

$$\begin{aligned}|f_0| &= 87.98774 \\x_1 &= 5.087201 \\x_2 &= 2.681470 \\x_3 &= 7.33074\end{aligned}$$

Execution time = 0.83 seconds

5.7 CONCLUDING REMARKS

The GREG code is very successful at solving nonlinear programming problems and appears to execute very efficiently. It does have some faults. One is that it is very sensitive to numerical error and behaves very erratically without any indication of what the problems are. Usage could be aided by the implementation of an error routine. FORTRAN programs of the magnitude of GREG are extremely difficult to control. For this reason the user should be proficient at FORTRAN programming.

The programming of the external subroutines could be aided by restructuring the GREG common block. Only a few variables and vectors are needed in the user supplied subroutines PHIX, CPHI, JACOB, and GRADFI. If these variables were placed in a named common block separate from the remaining GREG common variables, it would enhance programming by relieving the user of the possibility of creating conflicting variable names. Errors of that type are extremely difficult to locate since an error message would probably point to the permanent subroutines because a variable in common was unknowingly altered.

Despite the minor difficulties, the GREG code is extremely easy to use when considering the complexity of the problems it is designed to optimize. Being coded in FORTRAN, it may be used by a large number of programmers and could be altered to meet specific, specialized needs. The program is fast, efficient, and reasonably accurate in single precision. The code is obviously an enormous contribution to nonlinear optimization.

CHAPTER SIX

APPLICATION OF THE GRG METHOD TO A COMPLEX WATER QUALITY CONTROL SYSTEM

6.1 INTRODUCTION TO THE WATER QUALITY MANAGEMENT PROBLEM

Excessive degradation of water quality in streams has resulted from the rapid industrial growth of the past two decades. The realization has become apparent that even the largest rivers can no longer be treated as infinite sinks for human and industrial waste. Several factors are involved. The most obvious to the average person is the destruction of the esthetics in a normal ecological system. A more economically oriented problem is that the quality of water passed to downstream users may require excessive treatment to be used effectively, especially in the case of drinking water. Solutions to the water quality problems are in the realm of the political and social structure of our society. The problems must be compromised from all points of view so that the socio-ecological system serves for the mutual benefit of all members and not the destruction of any member.

The systems analysis approach to water quality management will by no means transpose water quality problems from the social and political system in our society, but as indicated by past experiences, can offer invaluable insights to the decision making processes. Much value can be derived from simplified mathematical models as long as the assumptions of the model have been effectively realized.

As early as 1925, the pioneering study of Streeter and Phelps [41] first attempted to apply the mathematical modeling approach to water quality control. Since that time, numerous problems have been defined and various linear and nonlinear models applied to them. A significant step has been the development of the dispersion models [20,22,23]. Traditionally the problems involved only one water quality standard, usually DO (dissolved oxygen), at a time. The usual form of pollution has been described as BOD (biochemical oxygen demand). Recently, models have been extended to include thermal pollution and water temperature

as a quality standard [16,17,18,23].

The evolution of water quality models has imposed two major problems on the systems analysis approach. The first, and the most severe, is the general lack of data to estimate the mathematical parameters. This problem is more acute as the model becomes more complex. A second problem is the complex array of mathematical equations that describe the system and difficulties in solving them. Of course, if enough data are available complex numerical methods to perform the mathematical analysis are worth the effort. The availability of data usually determines the model used and will ultimately affect the results of any study.

In this chapter a water quality management model that includes DO and temperature as quality standards, is presented. The model was originally proposed by Dysart and Hines [17]. Their approach was to define a multistage system with a treatment cost objective function and optimize it using a dynamic programming algorithm. The present chapter is concerned with reformulation of the model so that the generalized reduced gradient method via the GREG program can optimize it. The objective of the research was to demonstrate the effectiveness of the GRG method on water quality management problems and to demonstrate its flexibility for application to more involved models.

6.2 DO, BOD, AND TEMPERATURE RELATIONSHIPS IN FLOWING STREAMS

To develop the mathematical model, adequate transfer relationships of the state variables, temperature, BOD, and DO must be developed. The following transfer relations have been termed first order models and are extensions of the original Streeter-Phelps formulation [41].

The natural temperature of a stream is characterized by seasonable fluctuations and stochastic fluctuations within a given season. This makes it a

difficult variable to study in water quality control. One basic assumption is for any given time of the year the stream is characterized by an equilibrium temperature. When this temperature is raised by an unnatural process such as thermal pollution the stream temperature will tend to reduce to the equilibrium temperature by first order decay. The model is based, then, on the excess temperature of the stream.

$$\frac{d(T-E)}{dt} = -K(T-E)$$

T = water temperature in $^{\circ}\text{F}$

E = equilibrium temperature in $^{\circ}\text{F}$

t = time in days

K = temperature dissipation rate coefficient in days $^{-1}$

The analytical solution to this first order differential equation is

$$T_2 = E + (T_1 - E) \exp (-K (t_2 - t_1))$$

where

T_1 = initial water temperature at an upstream point in $^{\circ}\text{F}$

T_2 = water temperature at a downstream point in $^{\circ}\text{F}$

$t_2 - t_1$ = flow time from the initial starting point, t_1 , to the downstream point, t_2 , in days

The condition $T-E \geq 0$ is imposed on this model.

The development of the BOD transfer equation assumes that the normal BOD concentration is zero and that the dissipation of BOD follows a first order decay model.

$$\frac{dL}{dt} = -K_1 L$$

L = BOD concentration in mg/l

K_1 = deoxygenation rate coefficient in days⁻¹
 t = time in days

This differential equation has the simple solution

$$L_2 = L_1 \exp (-K_1(t_2 - t_1))$$

where

L_1 = initial BOD concentration in mg/l at time = t_1

L_2 = BOD concentration in mg/l at time = t_2

$t_2 - t_1$ = flow time in days

The dissolved oxygen transfer relation is derived from the classical Streeter-Phelps formulation [41]. It assumes that a saturation level exists in the stream at a given temperature and when the dissolved oxygen is depleted below this level the deficit approaches zero by a first order decay model. It also assumes the dissolved oxygen deficit grows with the influx of BOD by a first order growth model. The relation is

$$\frac{dD}{dt} = K_1 L - K_2 D$$

where

D = DO deficit in mg/l

K_2 = reaeration rate coefficient in days⁻¹

t = time in days

Knowing the BOD transfer equation the DO relationship becomes

$$\frac{dD}{dt} = K_1 L_1 \exp (-K_1(t-t_1)) - K_2 D$$

$$\frac{dD}{dt} + K_2 D = K_1 L_1 \exp(-K_1(t-t_1))$$

which is a first order linear differential equation. The solution is

$$D_2 = \frac{K_1 L_1}{K_2 - K_1} \left[\exp(-K_1(t_2 - t_1)) - \exp(-K_2(t_2 - t_1)) \right] \\ + D_1 \exp(-K_2(t_2 - t_1))$$

where

D_1 = initial DO deficit in mg/l

D_2 = DO deficit in mg/l at time t_2

$t_2 - t_1$ = flow time in days

D_2 is the DO deficit and it can be represented as

$$D_2 = DO_s - DO$$

where

DO_s = saturation DO level in mg/l at a given temperature at time = t_2

DO = DO level in mg/l at time = t_2

The DO transfer relation becomes

$$DO = DO_s - \frac{K_1 L_1}{K_2 - K_1} \left[\exp(-K_1(t_2 - t_1)) - \exp(-K_2(t_2 - t_1)) \right] \\ - D_1 \exp(-K_2(t_2 - t_1))$$

Temperature interacts with the BOD and the DO profiles and has drastic effects on the DO saturation level for a flowing stream. These interactions have been studied intensively. The DO saturation level of a stream has been

derived as a polynomial relation [42].

$$DO_s = 14.652 - 0.41022 \text{ TC} + 0.0079910 \text{ TC}^2 - 0.000077774 \text{ TC}^3$$

where

DO_s = DO saturation level in mg/l

TC = temperature in $^{\circ}\text{C}$

This curve was fitted to a set of data with a multiple correlation coefficient of 0.99980.

Likewise, the coefficients K_1 and K_2 are related to temperature by the relationships [19,11]

$$K_1 = K_{1,20} (1.047)^{(\text{TC}-20)}$$

$$K_2 = K_{2,20} (1.024)^{(\text{TC}-20)}$$

where

$K_{1,20}$ = deoxygenation rate coefficient at 20°C in days⁻¹

$K_{2,20}$ = reaeration rate coefficient at 20°C in days⁻¹

TC = temperature in $^{\circ}\text{C}$

A referral to the DO transfer relation indicates how dependent it is on the stream temperature.

The transfer relations represent a simplified model of a flowing stream and the study continued with this concept in mind. It must be realized that correct results are jeopardized by applying the parameter estimations of DO_s , K_1 , and K_2 to any existing stream. Regional variations in climate, geography, and soil composition justify doubt in their validity in application to streams other than the one for which they were developed.

The following assumptions about the transfer equations are implied:

1. uniform stream width
2. uniform stream depth
3. uniform and constant flow

6.3 DEFINITION OF THE N-STAGE RIVER BASIN MODEL

A river basin may now be modeled by dividing the system into stages, one at each point of discontinuity. A point of discontinuity consists of conditions that violate the assumptions of the transfer equations, changes in water quality standards, or boundaries determined by law that define a region under separate jurisdiction. Examples of discontinuities that require a new stage are

1. organic and/or thermal waste inputs
2. a state line
3. a change in the flow regime such as changes in stream width or depth
4. increases in the flow such as tributary inputs
5. changes in water quality standards

Figure 6.1 due to Dysart and Hines [17] exemplifies a generalized river basin model.

Consider a source at a stage, n, that has both organic and thermal input. The amount of treatment that must be supplied to its pollution input depends on two factors:

1. the quality of the stream water entering the stage from upstream sources
2. the water quality standards within the given stage.

Suppose that reliable data may be acquired to estimate the following values.

XT_n = thermal input in BTU/hr at the begining of stage n

XB_n = organic, BOD, input in lbs/day at the beginning of stage n

Q_n = uniform flow within stage n in ft^3/sec

TF_n = time of flow from the beginning to the end of stage n, in days ($t_2 - t_1$)

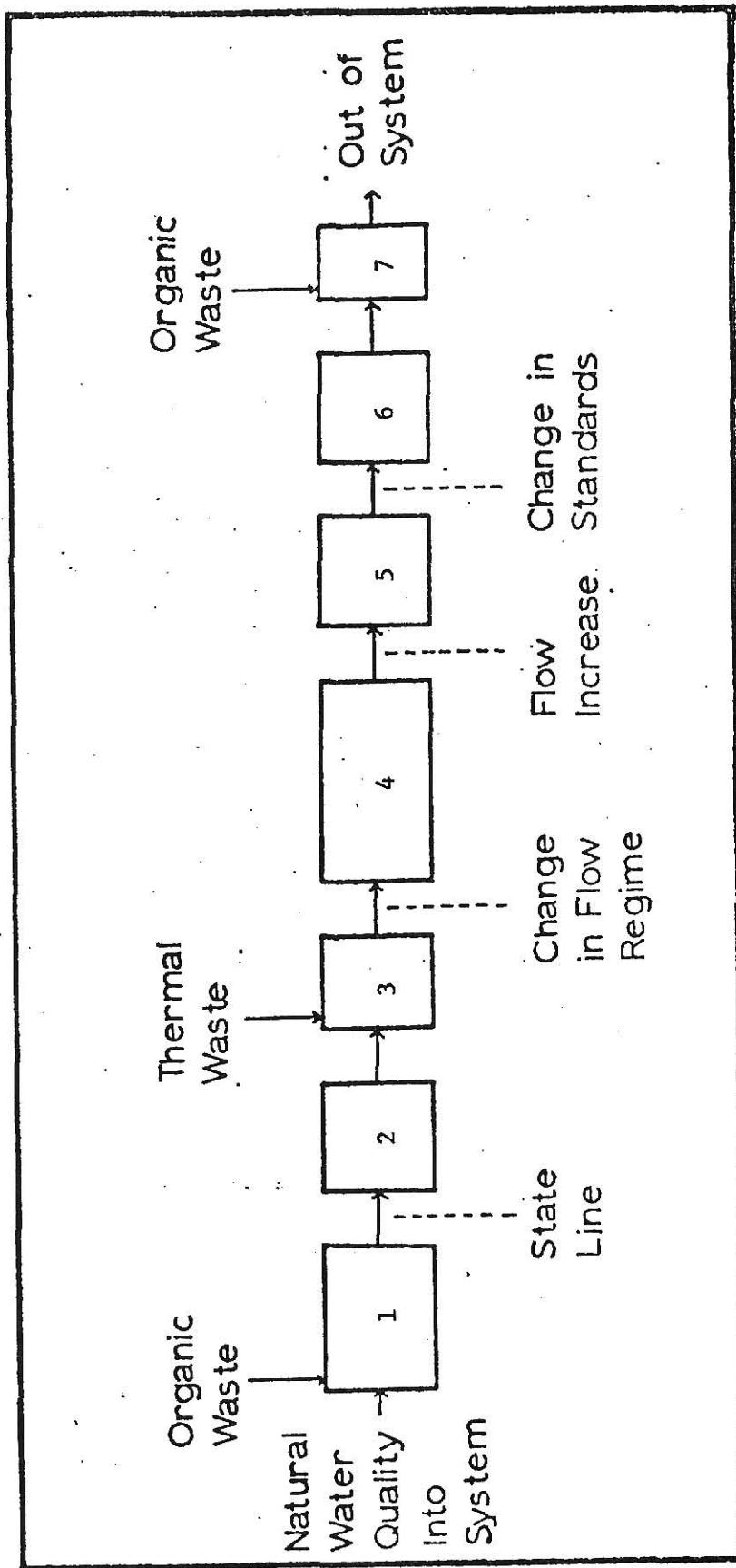


Fig. 6.1. An example of the discrete staging of the river basin model [17].

It is assumed that XT_n , XB_n , Q_n , and TF_n are constant. This is an unrealistic assumption because each value will vary daily, weekly, monthly, and seasonally. But, to formulate the optimization problem stochastic fluctuations are not considered and a static, steady state assumption is made. The rationalization is that if conditions that contribute to pollution such as low flow and high pollution input are studied then the system will behave favorably under normal conditions. It then becomes a question of how extreme should the extreme conditions be. This is a question that cannot be answered by the systems analyst. It becomes a philosophical question as to what risk of stream pollution is society willing to accept. One opinion is that as little risk as possible should be accepted. That is, very extreme conditions should be studied and treatment facilities should be designed so they will not be overloaded under the stress.

Thermal input in BTU/hr can be converted to the rise in stream temperature by the relation

$$\text{thermal input in } {}^{\circ}\text{F} = \frac{C_1 \cdot XT_n}{Q_n}$$

where

C_1 is a conversion constant

$$C_1 = 4.44970 \times 10^{-6} \quad (\frac{\text{hr} \cdot \text{ft}^3 \cdot {}^{\circ}\text{F}}{\text{sec} \cdot \text{BTU}})$$

The conversion of BTU to the rise in water temperature in ${}^{\circ}\text{F}$ and of water volume to mass is relative to the water temperature at 39.1 ${}^{\circ}\text{F}$. This is not an exact conversion at other temperatures, but the error is small when compared to the magnitude and significance level of a typical thermal pollution input value.

A similar expression is obtained when converting the BOD input to mg/ℓ .

$$\text{BOD input in } \text{mg}/\ell = \frac{C_2 \cdot XB_n}{Q_n}$$

$$C_2 = 0.185404 \left(\frac{\text{days} \cdot \text{ft}^3 \cdot \text{mg}}{\text{sec} \cdot \text{lbs} \cdot \ell} \right)$$

The values for deriving the conversion factors, C_1 and C_2 , were obtained from a handbook containing physical constants [32].

The temperature rise of the stream is now represented as a function of a treatment decision variable, $d(T,n)$, for the thermal input at stage n.

$$t\text{-rise}_n = \left[\frac{C_1 \cdot X T_n}{Q_n} \right] \left[\frac{100 - d(T,n)}{100} \right]$$

The expression represents the rise in stream temperature after a cooling treatment of $d(T,n)$ per cent.

$d(T,n) = \%$ of thermal input cooled to the equilibrium water temperature in stage n.

The maximum stream temperature occurs at the thermal pollution source and is represented as

$$\text{temp}_n = ST_n + \left[\frac{C_1 \cdot X T_n}{Q_n} \right] \left[\frac{100 - d(T,n)}{100} \right]$$

where ST_n is the incoming temperature for stage n.

Three water quality constraints for each stage are considered from the above temperature equations and the DO transfer equation.

$$t\text{-rise}_n \leq TRS_n$$

$$\text{temp}_n \leq TMAX_n$$

$$\min DO_n \geq MDO_n$$

The minimum DO constraint is basic because the normal ecological process of the stream is dependent on DO concentration.

Of extreme importance to the N-stage river basin model formulation are the boundary conditions of the state variables, ST_n , SB_n , and SO_n (temperature,

BOD, and DO, respectively). The output from stage n for a state variable must equal the input for stage n+1. The boundary relations are given in terms of the state variables.

Temperature

$$ST_{n+1} = E + (T_n - E) \exp (-K \cdot TF_n)$$

$$T_n = temp_n = ST_n + \left[\frac{C_1 \cdot XT_n}{Q_n} \right] \left[\frac{100 - d(T, n)}{100} \right]$$

where

ST_n = incoming temperature for stage n, in $^{\circ}\text{F}$

ST_{n+1} = outgoing temperature for stage n and the incoming temperature for stage n+1, in $^{\circ}\text{F}$

BOD

$$SB_{n+1} = L_n \exp (-K_1 \cdot TF_n)$$

$$L_n = SB_n + \left[\frac{C_2 \cdot XB_n}{Q_n} \right] \left[\frac{100 - d(B, n)}{100} \right]$$

where $d(B, n)$ = % treatment of BOD at the beginning of stage n

SB_n = incoming BOD for stage n, in mg/l

SB_{n+1} = outgoing BOD for stage n and the incoming BOD for stage n+1, in mg/l

DO

$$SO_{n+1} = DO_s - \frac{K_1 L_n}{K_2 - K_1} \left[\exp(-K_1 \cdot TF_n) - \exp(-K_2 \cdot TF_n) \right] - D_1 \exp(-K_2 \cdot TF_n)$$

D_1 = initial DO saturation level - SO_n

= initial DO deficit for stage n

SO_n = incoming DO for stage n in mg/l

S_{n+1} = outgoing DO for stage n and the incoming DO for stage n+1.

For all cases, $n = 1, \dots, N-1$.

In summary, the input of pollution into a stage is represented by the constant values XT_n and/or XB_n . If a source does not exist then the appropriate input value equals zero. Each stage is constrained by three water quality constraints based on the desired water quality standards. The BOD level is implicitly controlled by the minimum DO constraint. The boundary conditions for each stage are derived from the transfer relations and the initial state variable values, ST_1 , SB_1 , and SO_1 are considered constant in the steady state assumption. Figure 6.2 exemplifies a four stage system.

6.4 FORMULATION OF THE MATHEMATICAL OPTIMIZATION PROBLEM FOR THE GREG PROGRAM

This section describes the development of the water quality constraints for optimization by the GREG program. The optimization problem may be defined as minimize

$$\text{cost} = \sum_{n=1}^N C_n$$

subject to the constraints

$$t\text{-rise}_n \leq TRS_n$$

$$\text{temp}_n \leq TMAX_n$$

$$\min DO_n \geq MDO_n$$

$$n = 1, \dots, N$$

where C_n is the treatment cost as of function of the treatment level (decision variables) for the thermal and BOD inputs of stage n. The problem can be defined in terms of the decision variables. Let

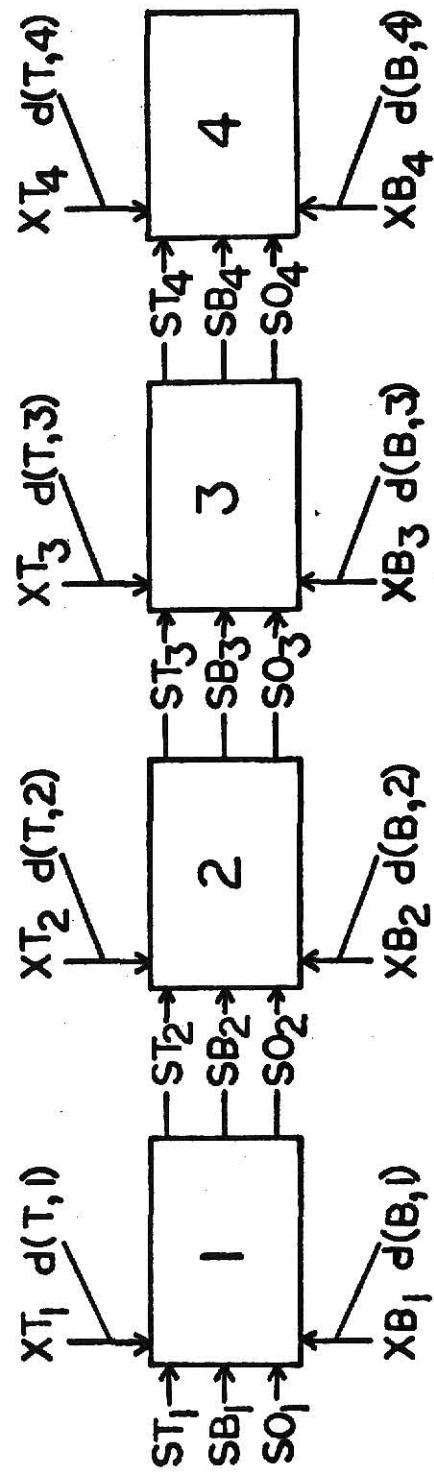


Fig. 6.2. A schematic diagram of the general river basin model with four stages. If an input does not exist in a stage then its value is appropriately set equal to zero.

$$x_i = d(T, i), \quad x_{N+i} = d(B, i)$$

$i = 1, \dots, N, \quad N = \text{number of stages}$

The total number of decision variables is $2 \times N$. The temperature rise constraint is then given as

$$f_i(\bar{x}) = \left[\frac{C_i \cdot x_{T,i}}{100 Q_i} \right] [100 - x_i] - TRS_i \leq 0 \quad i = 1, \dots, N$$

The first N constraints consist of the temperature rise constraint for each stage. The maximum temperature constraint is formulated as

$$f_{N+i}(\bar{x}) = ST_i + \left[\frac{C_i \cdot x_{T,i}}{100 Q_i} \right] [100 - x_i] - TMAX_i \leq 0 \quad i = 1, \dots, N$$

This set of N constraints contains the maximum temperature constraints for each stage. The final set of constraints is for the minimum DO restriction.

$$f_{2N+i}(\bar{x}) = -\min [DO] + MDO_i \leq 0 \quad i = 1, \dots, N$$

The total number of constraints is $3 \times N$.

Calculation of the constraints at a point, \bar{x} , consists of two operations. First, the boundary conditions are updated for each stage. That is, the values for ST_i , SB_i , and SO_i are calculated. Only ST_1 , SB_1 , and SO_1 remain constant. The values of the other state variables depend on the values of the previous decision variables. ST_i , SB_i , and SO_i are functions of x_j and x_{N+j} , $j = 1, \dots, i - 1$.

After the boundaries are calculated the values of the constraints are determined. The calculation of the temperature rise and maximum temperature constraints is direct, but the value of the minimum DO in a stage is more difficult to obtain.

It is known that the minimum DO is between $t = 0$ and $t = TF_i$ for stage i .

The problem is to determine a value $t = t_{\min}$ such that the DO concentration is minimum. A Fibonacci search is used to determine the minimum. The number of iterations is determined during the initialization of the program so that the desired precision is obtained in the calculation. The number of iterations is determined so the length of the final interval containing t_{\min} is 0.0001.

Calculation of the partial derivatives is the next task in preparing the problem for the GREG program. For the temperature rise constraints the partial derivatives are determined as

$$\frac{\partial f_i}{\partial x_i} = - \left[\frac{C_1 \cdot X T_i}{100 Q_1} \right]$$

$$\frac{\partial f_j}{\partial x_j} = 0, \quad j = 1, \dots, N \text{ and } j \neq i$$

$$\frac{\partial f}{\partial x_{N+i}} = 0, \quad j = 1, \dots, N$$

For this set of constraints the partial derivatives are all constant. Similar constant results occur with the maximum temperature constraint except the calculations involve previous decision variables.

$$\frac{\partial f_{N+i}}{\partial x_i} = - \left[\frac{C_1 \cdot X T_i}{100 Q_1} \right]$$

$$\frac{\partial f_{N+i}}{\partial x_j} = 0, \quad j > i$$

$$\frac{\partial f_{N+i}}{\partial x_j} = \frac{\partial S T_i}{\partial x_j}, \quad j < i$$

$$\frac{\partial f_{N+i}}{\partial X_{N+j}} = 0, \quad j = 1, \dots, N$$

The partial derivatives of the state variable, ST_i , are calculated inductively.

$$ST_i = XE_{i-1} + (T_i - XE_{i-1}) \exp (-K_{i-1} TF_{i-1})$$

$$T_i = ST_{i-1} + \frac{C_1 \cdot XT_{i-1}}{100 Q_{i-1}} [100 - X_i]$$

K_{i-1} = (K) the temperature dissipation rate coefficient specified for a given stage, $i - 1$. This subscripting has not been used in the previous discussion, but this, and all other constants, are generally different in each stage.

Thus,

$$\frac{\partial ST_i}{\partial X_{i-1}} = - \left[\frac{C_1 \cdot XT_{i-1}}{100 Q_{i-1}} \right] \exp (-K_{i-1} TF_{i-1})$$

$$\frac{\partial ST_i}{\partial X_{i-k}} = - \left[\frac{C_1 \cdot XT_{i-k}}{100 Q_{i-k}} \right] \exp \left(- \sum_{m=1}^k K_{i-m} TF_{i-m} \right), \quad 0 < k < i$$

By letting $j = i - k$,

$$\frac{\partial f_{N+i}}{\partial X_j} = - \left[\frac{C_1 \cdot XT_j}{100 Q_j} \right] \exp \left(\sum_{l=j}^{i-1} K_l \cdot TF_l \right), \quad j < i$$

Analytic partial differentiation of the minimum DO function is considered very impractical so it is accomplished numerically. The procedure is to calculate the minimum DO at the point, \bar{X} , and store the value, FMIN. Then, for each

variable X_j (and X_{N+j}) the calculation is

$$1. \quad FMIN1 = \min [DO] \text{ when } X_j \leftarrow X_j + \Delta$$

$$2. \quad \frac{\partial f_{2N+i}}{\partial X_j} \approx -\left[\frac{FMIN1 - FMIN}{\Delta} \right]$$

The minus sign comes from the minus sign in the constraint function

$$f_{2N+i} = -\min [DO_i] + MDO_i \leq 0$$

The calculations are performed only for the indices j such that $j \leq i$. Any decision, X_j and X_{N+j} , made in a downstream stage ($j > i$) has no effect on the previous stages and their constraints, thus indicating the partial derivatives for these cases are zero.

6.5 APPLICATION TO THE CHATTAHOOCHEE RIVER BASIN

To test the application of the GRG method to a N-stage water quality management model, data obtained from the study of Dysart and Hines [16,17,18] on the Chattahoochee River Basin below Atlanta, Georgia was utilized. A map of the system is shown in Fig. 6.3 [17]. The portion of the river under consideration in the study was bounded above at Atlanta, specifically at Clayton. The lower bound of the system was Yellowjacket Creek. There are numerous points of organic and thermal pollution input along this stretch of the river, but only the major contributors were considered in the study. The only major organic polluter is the Clayton plant. There are three major thermal polluters, McDonough, Atkinson, and Yates power plants. Four sources comprise the majority of both the organic and thermal waste in the system under consideration.

Figure 6.4 [17] summarizes the system and divides it into four steady-

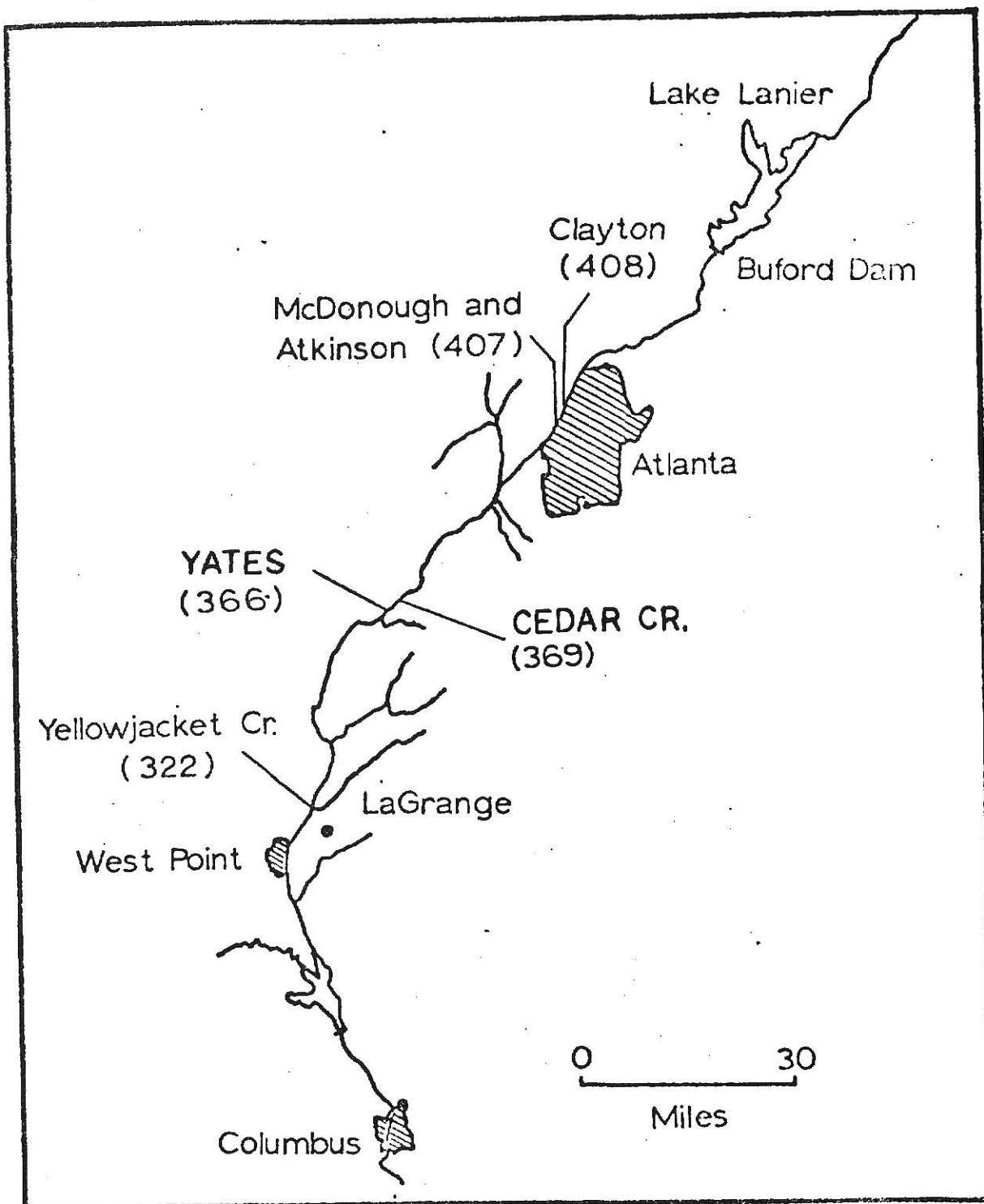


Fig. 6.3. A map of the Chattahoochee River basin [17].

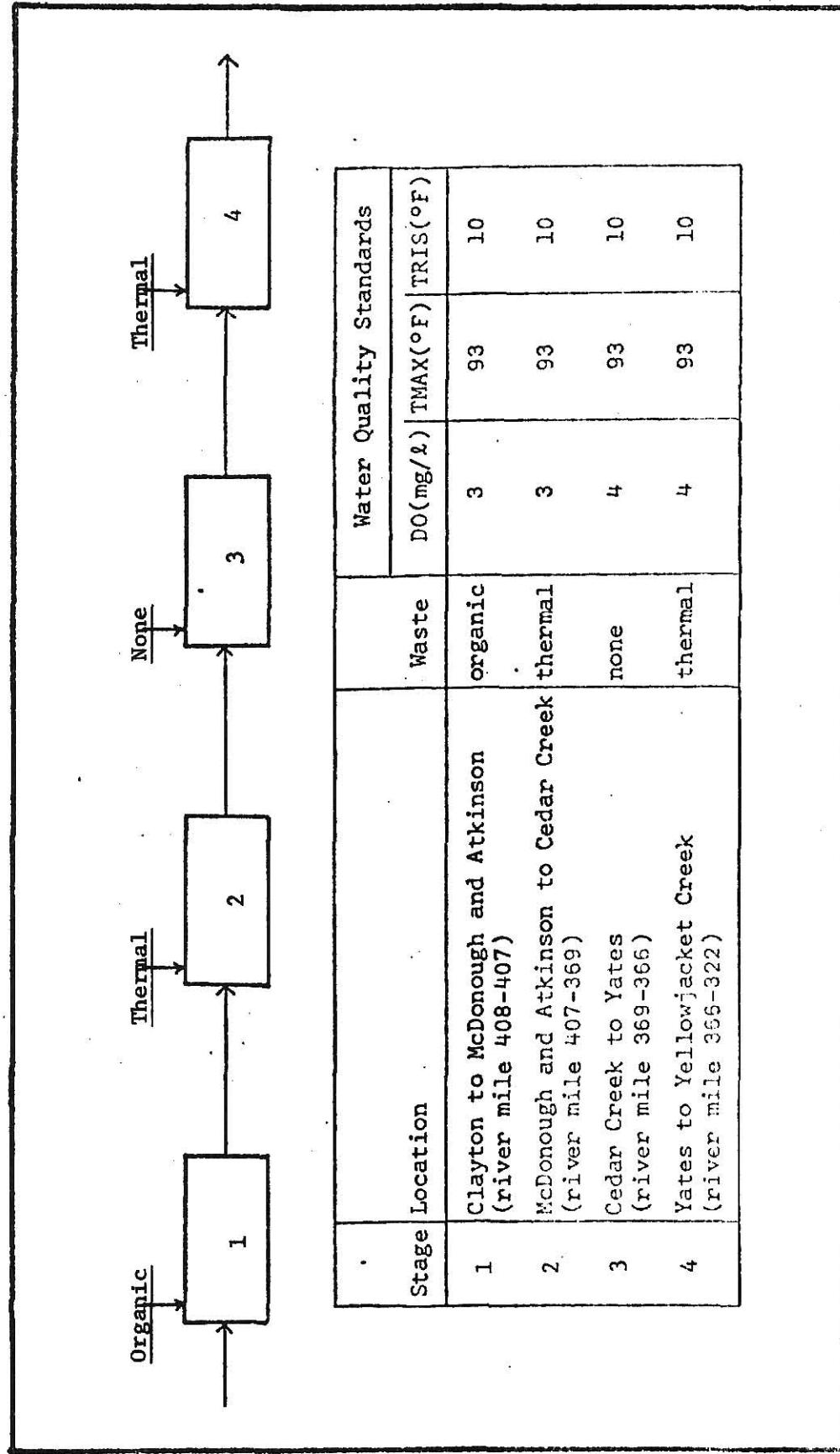


Fig. 6.4. The staging and the water quality standards of the river basin model [17].

state stages. Stage one describes the organic waste input and the resultant effect on the system for the Clayton plant. This stage extends to the McDonough and Atkinson plants where a discontinuity occurs with thermal discharge. These two plants are close together and can be treated as one source for stage two. The second stage extends to Cedar Creek where another discontinuity occurs. At Cedar Creek the DO standard changes as indicated in Fig. 6.4. The third stage is created and extends to the Yates plant, another thermal source. This becomes the upper bound for the fourth stage which extends to Yellowjacket Creek, the lower bound of the system. In this example, discontinuities resulting in new stages occurred, for all practical reasons, at pollution sources and because of a change in the DO standard. It was not considered necessary to divide the system into more stages because of changes in flow regime or for flow increases. The model described in Fig. 6.4 is believed to represent the system within the assumptions of the water quality management river basin model.

Input data for the model for the purpose of the present study are given in Table 6.1. The data were extracted from the "One-in-Ten September" case of Dysart and Hines [17]. This means that the lowest flow of the river occurring in September during the previous ten years was used as the basis for a conservative abatement, or treatment policy for the system. It was reasoned that if the extreme case was studied for a given time of the year then policy based on it would be satisfactory for more favorable conditions. Of course, this would not be true where significant BOD concentration is created by runoff (such as feedlot runoff) in high flow conditions.

The objective function of the system was given in terms of treatment costs for the three stages with pollution input. Figure 6.5 [17] shows the cost of organic waste treatment in millions of dollars as a function of the abatement level for the Clayton plant. This function, for the purpose of this

n	Location	XT_n (BTU/hr)	XB_n (lbs/day)	Q_n (ft^3 /sec)	XE_n (°F)	K (days $^{-1}$)	$K_{1,20}$ (days $^{-1}$)	$K_{2,20}$ (days $^{-1}$)	TF_n (days)	distance (miles)
1.	Clayton to McDonough and Atkinson	0	1.1×10^5	820	80.7	1.24	1.00	0.90	0.05	1
2.	McDonough and Atkinson to Cedar Creek	2.4×10^9	0	870	80.7	1.29	1.00	0.90	1.87	38
3.	Cedar Creek to Yates	0	0	885	80.7	1.10	1.00	0.90	0.17	3
4.	Yates to Yellowjacket Creek	1.6×10^9	0	920	80.7	1.34	1.00	0.90	2.40	44

Table 6.1. The input parameters for the "One-in-Ten September" case.

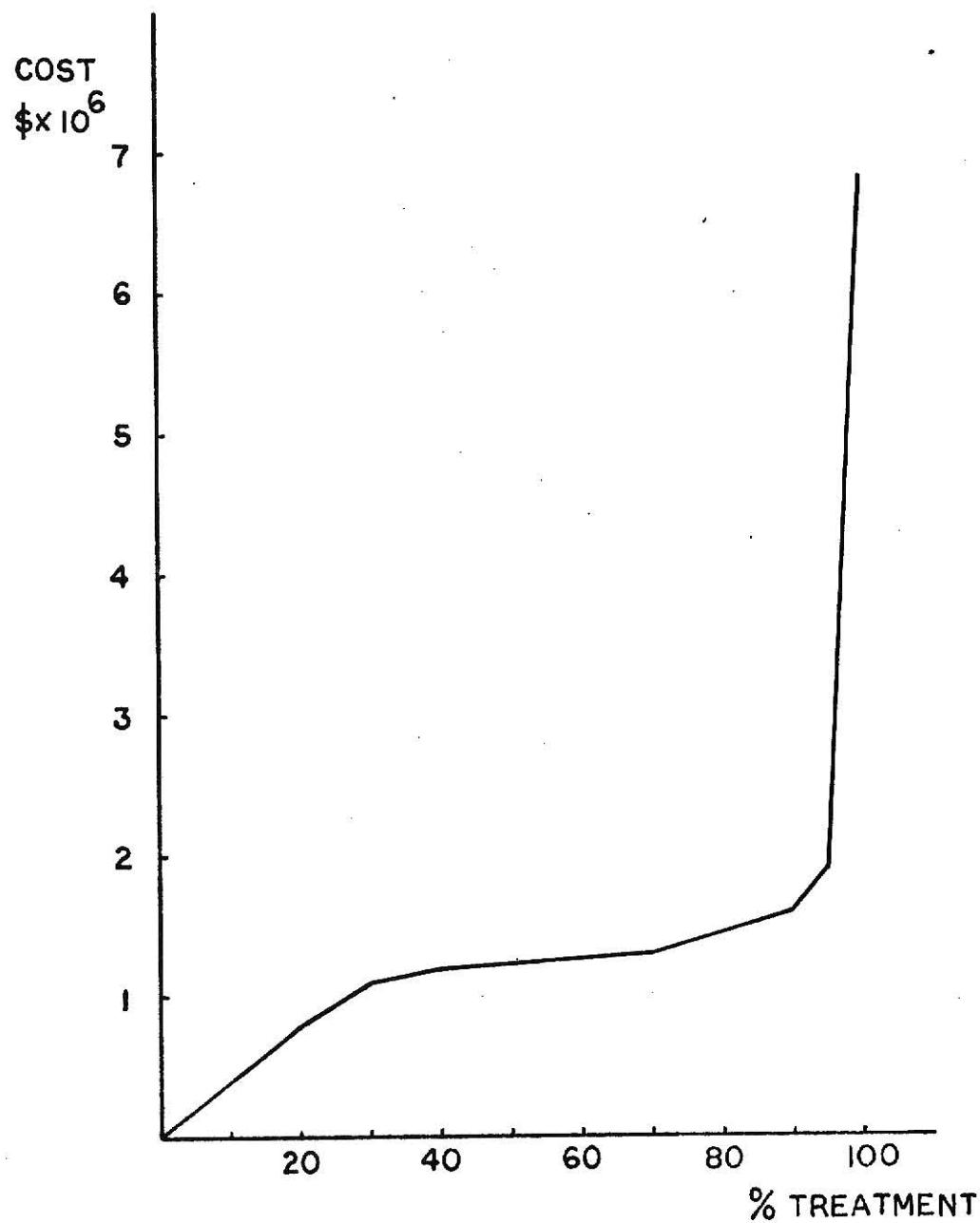


Fig. 6.5. The Clayton plant abatement cost of organic waste treatment as a function of per cent treatment.

study, is represented by a piecewise linear function.

$X = \% \text{ treatment}$

$$\text{cost} = 0.04X, \quad 0 \leq X < 20$$

$$\text{cost} = 0.03X + 0.2, \quad 20 \leq X < 30$$

$$\text{cost} = 0.01X + 0.8, \quad 30 \leq X < 40$$

$$\text{cost} = \left(\frac{0.01}{3} \right) X - \left(\frac{0.40}{3} - 1.2 \right), \quad 40 \leq X < 70$$

$$\text{cost} = 0.015X + 0.25, \quad 70 \leq X < 90$$

$$\text{cost} = 0.06X - 3.8, \quad 90 \leq X < 95$$

$$\text{cost} = 0.98X - 91.2, \quad 95 \leq X < 100$$

The cost for the treatment of the thermal waste was given as (in \$)

$X = \% \text{ treatment}$

$$\text{Cool Cost} = \left(\frac{817,000}{0.9} \right) [1 - \exp (- 0.023 X)]$$

(McDonough
and
Atkinson)

$$\text{Cool Cost} = \left(\frac{575,000}{0.9} \right) [1 - \exp (- 0.023 X)]$$

(Yates)

The objective function is separable and its gradient is easily defined for optimization with the GREG program.

The programming of the external GREG subroutines is given in Appendix 5. The results are encouraging since both accuracy and computation time are favorable when using the GRG method. The accuracy of the problem analysis is improved over the dynamic programming approach because the system is treated on a continuous basis. The main improvement, though, is that computation to an optimum solution is significantly faster than the dynamic programming approach.

The abatement policy was determined to be

% treatment = 66.9%
(Clayton)

% treatment = 18.5%
(McDonough
and
Atkinson)

% treatment = 0.0%
(Yates)

The total cost of the system abatement policy by the cost function used was \$1.60 million. The cost function used for the Clayton plant was a little different from the Dysart and Hines study [17] because values for it were extrapolated from a graph. The solution to this problem from their study was

% treatment = 70%
(Clayton)

% treatment = 20%
(McDonough
and
Atkinson)

% treatment = 0.0%
(Yates)

$$\text{Total cost} = \$1.576 \times 10^6$$

The execution time of the GREG program on the IBM 360/50 computer from a feasible starting point of 100% treatment for all inputs was 39.41 seconds. From the nonfeasible starting point of 0% treatment for all sources the execution time was 30.35 seconds. The execution times cannot be compared directly but the dynamic programming approach was attributed a typical 10-15 minute execution time on a Univac 1108 computer [18]. The program was written in the ALGOL language.

The resulting DO profile is shown in Fig. 6.6. A minimum DO concentration is reached in stage two. This minimum is caused by the organic waste of stage one. The temperature profile is shown in Fig. 6.7. The two jumps in the graph are caused by the thermal inputs. The variation of K_1 and K_2 , which depends on temperature, is shown in Fig. 6.8.

6.6 CONCLUDING REMARKS

The application of the GRG method to the water quality management model demonstrates its flexibility, speed, and accuracy. It serves as a pilot study to further research involving more complicated water quality models that are more realistic in describing the behavior of such a complex system. Two important results of this study were that the minimum DO and its partial derivatives were determined numerically with much success. This implies applications of the technique to an assortment of constraints. The most important application to water quality models would be the numerical integration of simultaneous differential equations to define the constraints. This makes it conceivable to apply the procedure to the more complex dispersion models involving BOD, DO, and temperature with less restrictive assumptions (see Fan et al. [23]).

This study also demonstrated the use of a monotonically increasing piecewise linear cost function. This demonstrated the flexibility in defining objective functions for the GRG method. It is the flexibility of the GRG method that is desirable. It has a wide variety of applications including water quality optimization models.

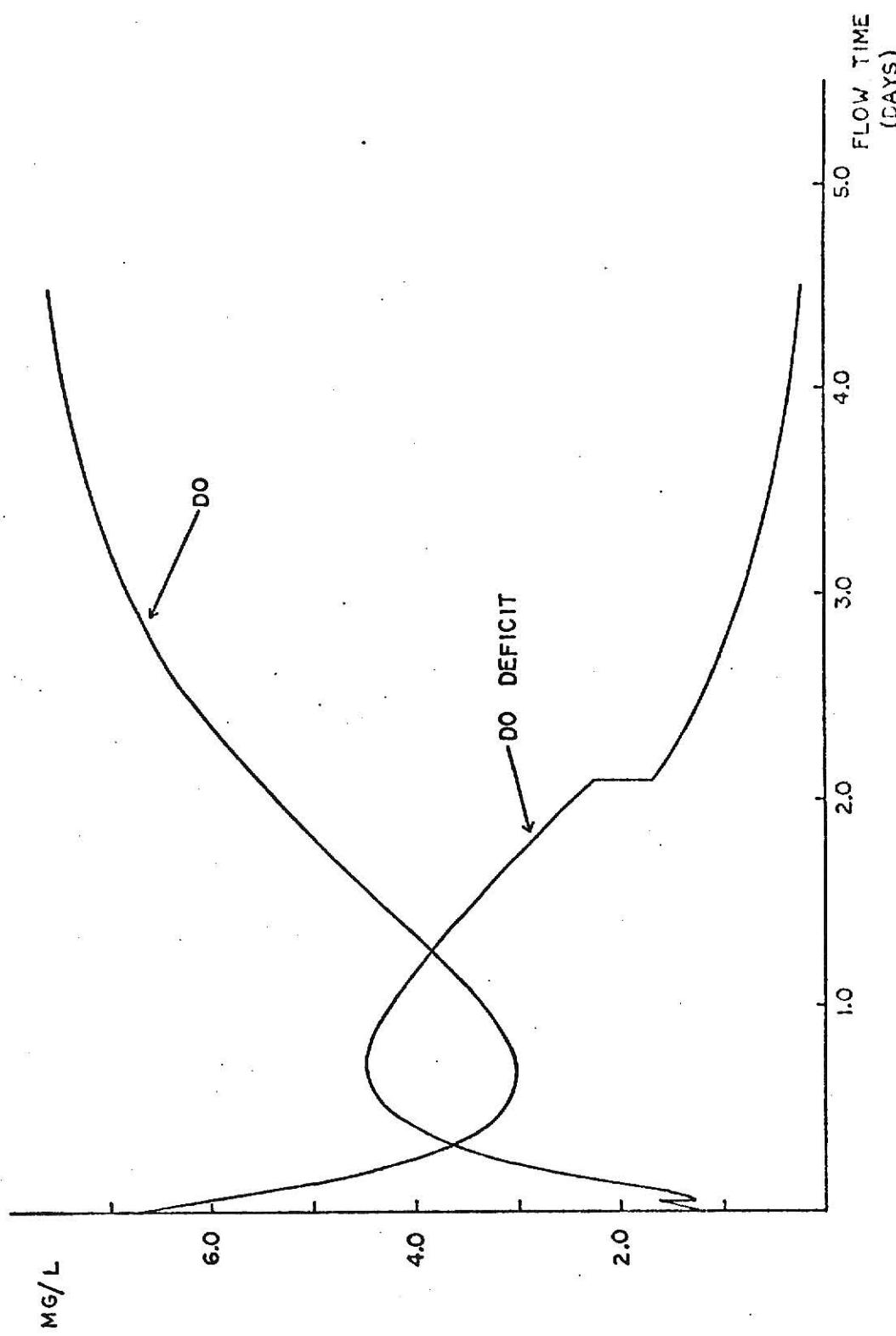


Fig. 6.6. The DO and DO deficit profiles for the system at the optimum solution. The discontinuities of the DO deficit curve represent thermal pollution input from the McDonough and Atkinson plants and the Yates plant.

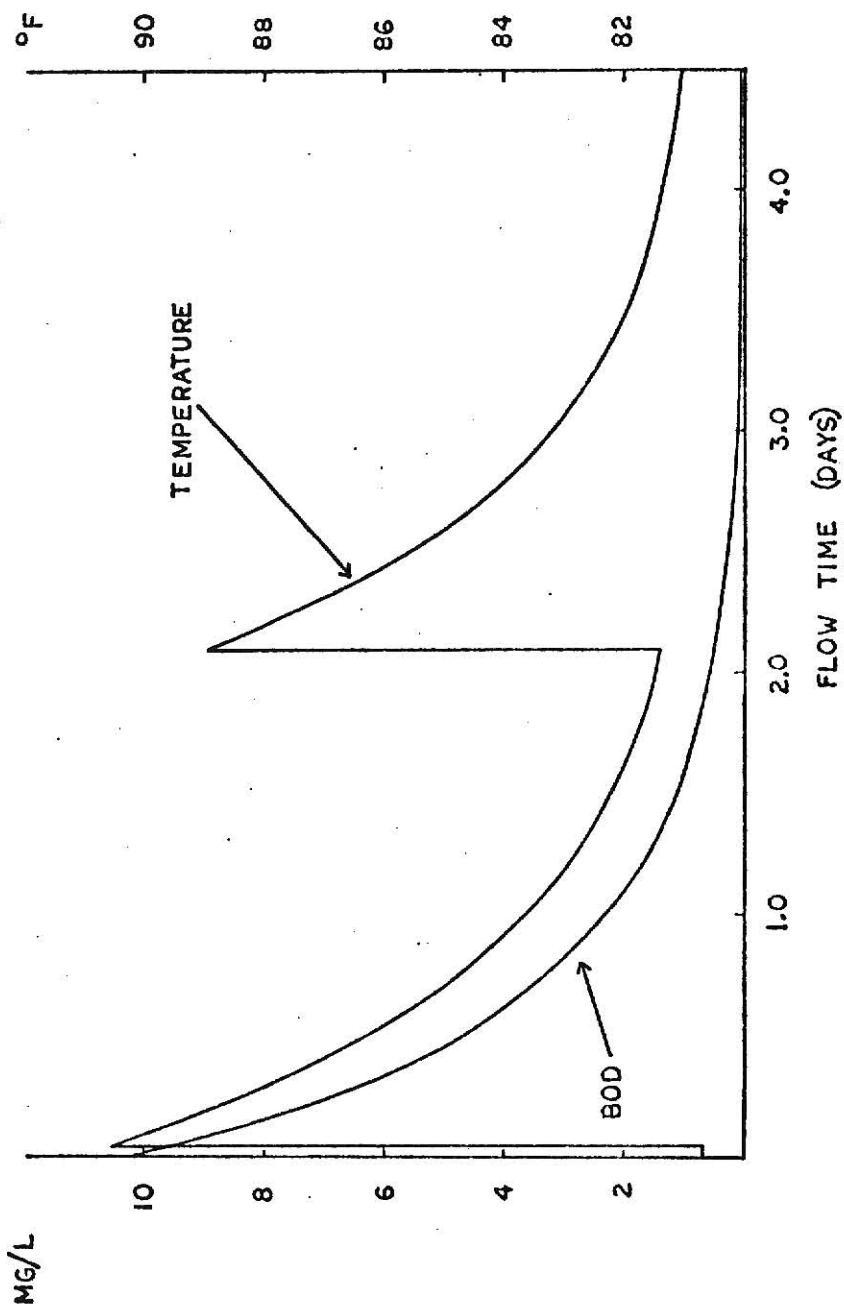


Fig. 6.7. The temperature and the BOD profiles for the system. The BOD input occurs at the Clayton plant (flow time = 0.0). The discontinuities of the temperature curve occur at the source of the two thermal inputs.

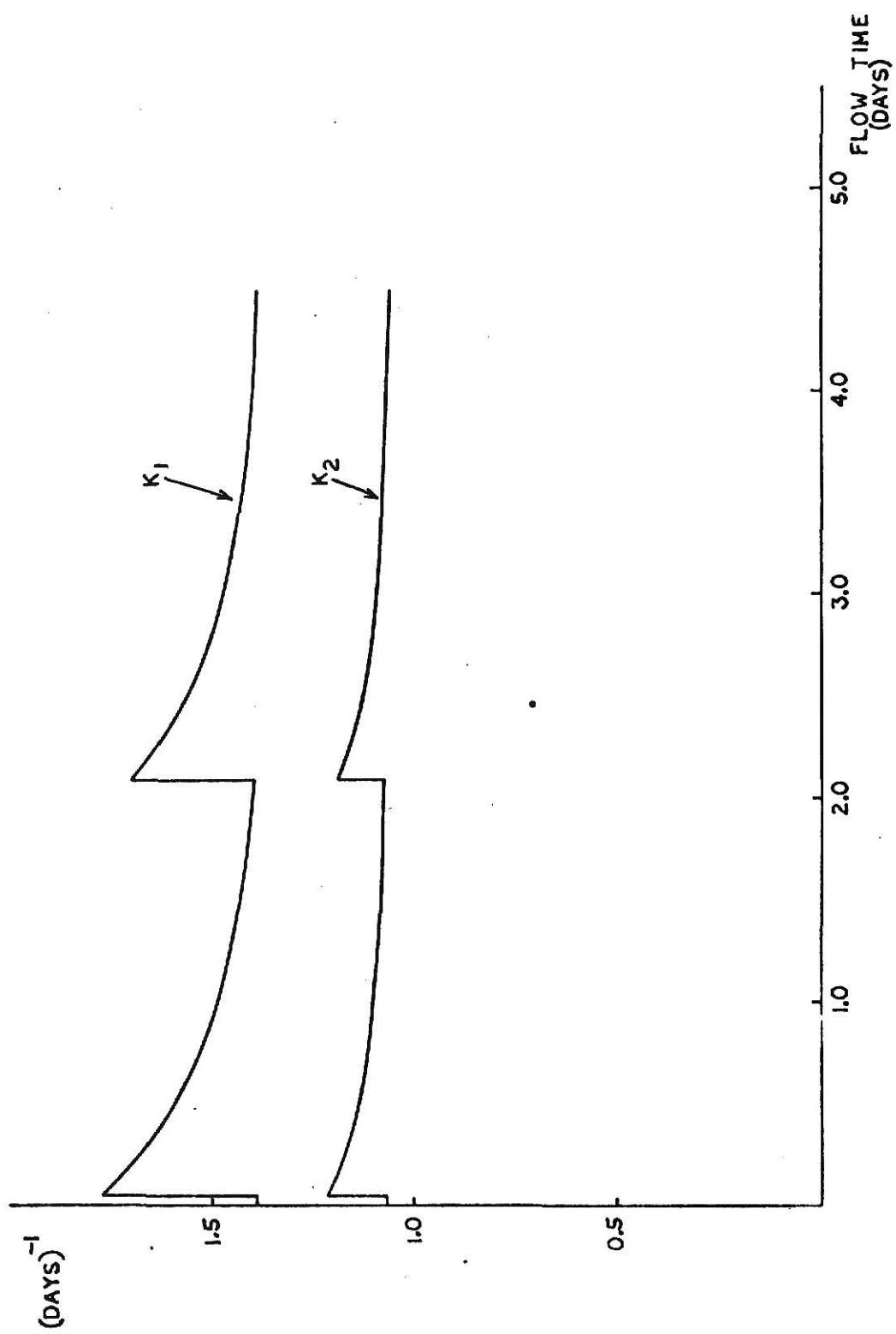


Fig. 6.8. The variation of the deoxygenation rate coefficient and the reaeration rate coefficient. Notice the close relation to the temperature curve of Fig. 6.7.

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APPENDIX 1. COMPUTER PROGRAM FOR GENERATING SEPARABLE PROGRAMMING DATA

This appendix contains a listing of the main program for generating the separable programming data for MPS/360. The user must supply the subroutine FUNC.

```

C.....JEREL WILLIAMS.....JULY,1971.....
C
C....THIS ROUTINE IS DESIGNED TO GENERATE THE LINEAR APPROXIMATION VALUES
C    AND TO PUNCH DATA CARDS WITHIN THE SPECIFICATIONS OF MPS/360
C    LINEAR AND SEPARABLE PROGRAMMING.
C
C....A USER DESIGNED SUBROUTINE FUNC MUST ACCOMPANY THIS ROUTINE.
C
C....THE USER MUST SUPPLY DATA FOR SIX VARIABLES.  ALL DATA IS ECHO CHECKED,
C    NFUN=THE NUMBER OF VARIABLES
C    NROWS=THE NUMBER OF CONSTRAINTS+1
C    START=THE INITIAL STARTING POINT FOR EACH VARIABLE
C    INTER=THE NUMBER OF UNIFORM PARTITIONS DESIRED BETWEEN ALOW & UPPER
C    ALOW=THE LOWER BOUND OF THE UNIFORM REGION
C    UPPER=THE UPPER BOUND OF THE UNIFORM REGION
C
C....THE DATA CARD FORMAT IS AS FOLLOWS.
C    FIRST CARD
C        READ NFUN & NROWS, FORMAT 215
C        THERE NOW OCCURS ONE DATA CARD FOR EACH VARIABLE CONTAINING
C        START, INTER, ALOW, & UPPER, FORMAT F10.0, I10, 2F10.0
C
C    900 FORMAT(5X,5HEPROR/5X,30HALOW EQUALS UPPER FOR VARIABLE,I5)
C    901 FORMAT(5X,5HERROR/5X,31HINVALID VALUE FOR NFUN OR NROWS)
C    902 FORMAT(5X,5HERROR/5X,27HINTER EQUALS 0 FOR VARIABLE,I5)
C    1000 FORMAT(215)
C    2000 FORMAT(4X,1HP,I4,5X,3IROW,I3,4X,F12.5,3X,3HROW,I3,4X,F12.5)
C    3000 FORMAT(4X,1HP,I4,5X,4HGRID,I3,3X,F12.5)
C    3001 FORMAT(4X,1HP,I4,5X,3IROW,I3,4X,F12.5,3X,4HGR10,I3,3X,F12.5)
C    4000 FFORMAT(4X,3HSFT,I3,4X,8H'MARKER',17X,8H'SEPORG')
C    5000 FORMAT(F10.0,I10,2F10.0)
C    6000 FORMAT(5X,23HFCHO CHECK FOR VARIABLE,I5/5X,6HSTART=,F10.2/
C                15X,6HINTER=,I10 /5X,6HALOW =,F10.2/5X,6HUPPER=,F10.2/)
C    7000 FORMAT(4X,6HENSET,4X,8H'MARKER',17X,8H'SEPEND'/80(1H*))
C    8000 FORMAT(1X,14HUP SEPBLND P,I4, 9X,3H1.0)
C    9000 FORMAT(5X,27HFCHO CHECK FOR NFUN & NROWS/
C                15X,6HNFUN =,[5/5X,6HNROWS=,I5/)

C
C....UNIT NUMBERS
C
C    NREAD=5
C    NPRNT=6
C    NPNCH=7
C
C....NREAD=CARD INPUT,NPRNT=PRINT OUTPUT,NPNCH=PUNCH OUTPUT
C
C    DIMENSION ROW(50),F150,50)
C    READ(NREAD,1000) NFUN,NROWS
C    WRITE(NPRNT,5000) NFUN,NROWS
C    IF(NFUN.LT.1.OR.NROWS.LT.1) GO TO 96
C    KKK=0
C    KKK=1000
C
C....DO CALCULATIONS FOR EACH VARIABLE
C
C    DO400KK=1,NFUN
C    IAX=1
C    READ(NREAD,5000) START,INTER,ALOW,UPPER

```

```

      WRITE(NPRNT,6000) KK,START,INTER,ALOW,UPPER
C
C.....CHECK DATA FOR ERRORS. PROGRAM DOES NOT ALLOW INTER .GT. 50.
C
      IF(INTER.GT.50) INTER=50
      IF(INTER.EQ.0) GO TO 97
      IF(ALOW.EQ.UPPER) GO TO 98
C
C.....DETERMINE UNIFORM PARTITION STEP-SIZE
C
      AN=INTER
      DELTA=(UPPER-ALOW)/AN
      LL=KK+100
      WRITE(NPNCH,4000) LL
      NN=NROWS+KK
C
C.....CHECK ALOW,NE,START. CREATE ANOTHER PARTITION IF TRUE
C
      IJ=0
      IF(ALOW-START)2,3,2
      2 X=START
      AX=ALOW
      IJ=1
      INTER=INTER+1
      IAX=2
      DO11IK=1,NROWS
      11 ROW(IK)=0.0
      GO TO 8
C
C.....CALCULATE PARTITION VALUES
C
      3 IJ=IJ+1
      DO10IK=1,NROWS
      10 ROW(IK)=0.0
      XK=IJ-IAX
      X=XK*DELTA+ALOW
      AX=X+DELTA
C
C.....CALCULATE FUNCTION DIFFERENCES
C
      8 DO100J=1,NROWS
      CALL FUNC(J,KK,AX,F,NFUN,NROWS)
      ROW(J)=F(J,KK)
      CALL FUNC(J,KK, X,F,NFUN,NROWS)
      100 ROW(J)=ROW(J)-F(J,KK)
C
C.....CREATE GRID VALUE
C
      GRID=X-AX
      II=IJ+KKK
C
C.....START PUNCH LOOP
C
      IR=0
      16 IR=IR+1
      IF(IR.GT.NROWS) GO TO 90
      IF(ROW(IR))17,200,17
      17 IRRR=IR+1

```

```
      GO TO 20
19 IRRR=IRR+1
20 IF(IRR>NRCWS) GO TO 89
   IF(ROW(IRR))75,19,75
89 IRR=IR+100
   WRITE(NPNCH,3C01) II,IRR,ROW(IR),LL,GRID
   IR=NN
   GO TO 200
90 IR=NN
   WRITE(NPNCH,3C001) II,LL,GRID
   GO TO 200
75 IRR=IR+100
   IMM=IRR+100
   WRITE(NPNCH,2C001) II,IRR,ROW(IR),IMM,ROW(IRR)
199 IR=IRR
200 IF(IR.LT.NN) GO TO 16
   IF(IJ.LT.INTER) GO TO 3
   KKK=KKK+INTER
   NKK=NKK+INTER
400 CONTINUE
   WRITE(NPNCH,7C001)

C
C.....CREATE BOUNDS SECTION
C
   C01I=1,NKK
   II=I+1000
1  WRITE(NPNCH,8C001) II
   STOP
96 WRITE(NPRNT,9C1)
   STOP
97 WRITE(NPRNT,9C2) KK
   STOP
98 WRITE(NPRNT,9C01) KK
   STOP
END
```

APPENDIX 2. COMPUTER OUTPUT FROM THE RELIABILITY PROBLEM

This appendix contains a listing of the subroutine FUNC and the MPS/360 control language program used to solve the reliability problem. The MPS/360 solution output is also given.

```
SUBROUTINE FUNC(J,KK,X,F,NFUN,NRCWS)
DIMENSION F(50,50)
MM=(KK-1)*NRCWS+J
GO TO(1,2,3,4,5,6,7,8,9,10,11,12,13,14,15,16,17,18,19,20),MM
1 F(1,1)= ALOG(1.-.2**X)
RETURN
2 F(2,1)=X**2
RETURN
3 F(3,1)=7.*(X+EXP(X/4.))
RETURN
4 F(4,1)=7.*X*EXP(X/4.)
RETURN
5 F(1,2)= ALOG(1.-.15**X)
RETURN
6 F(2,2)=2.*X**2
RETURN
7 F(3,2)=7.*(X+EXP(X/4.))
RETURN
8 F(4,2)=8.*X*EXP(X/4.)
RETURN
9 F(1,3)= ALOG(1.-.1**X)
RETURN
10 F(2,3)=3.*X**2
RETURN
11 F(3,3)=5.*(X+EXP(X/4.))
RETURN
12 F(4,3)=8.*X*EXP(X/4.)
RETURN
13 F(1,4)= ALOG(1.-.35**X)
RETURN
14 F(2,4)=4.*X**2
RETURN
15 F(3,4)=9.* (X+EXP(X/4.))
RETURN
16 F(4,4)=6.*X*EXP(X/4.)
RETURN
17 F(1,5)= ALOG(1.-.25**X)
RETURN
18 F(2,5)=2.*X**2
RETURN
19 F(3,5)=4.*(X+EXP(X/4.))
RETURN
20 F(4,5)=9.*X*EXP(X/4.)
RETURN
END
```

CONTROL PROGRAM COMPILER - MPS/360 V2-M8

```
C001      PROGRAM
C002      INITIALZ
C064      MOVE(XDATA,'DATA-SET')
C065      MOVE(XPBNAME,'SEPARABLE')
C066      CONVERT
C067      BCROUT
C068      MOVE(XOBJ,'ROW101')
C069      MOVE(XRHS,'LIMITS')
C070      TITLE('PRIMAL-LOCAL OPTIMUM')
C071      XSFTLB=-1    $* BOUNDED VARIABLES AT LOWER LIMIT
C073      SETUP('BOUND','SEPBOUND','MAX')
C074      PRIMAL
C075      SOLUTION    $* LOCAL OPTIMUM
C077      TITLE('PRIMAL-GLOBAL OPTIMUM')
C078      XSFTLB=+1    $* BOUNDED VARIABLES AT UPPER LIMIT
C080      SETUP('BOUND','SEPBOUND','MAX')
C081      PRIMAL
C082      SOLUTION    $* GLOBAL OPTIMUM
C084      TITLE('DUAL-GLOBAL OPTIMUM')
C085      SETUP('BOUND','SEPBOUND','MAX')
C086      DUAL
C087      SOLUTION    $* GLOBAL OPTIMUM-DUAL
C089      EXIT
C090      PEND
```

DUAL-LINEAL OPTIMUM

SECTION 1 - PICS

SOLUTION (OPTIMAL)

TYPE = C.21 X15. ITERATION NUMBER =

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.ACTV ***ACTIVITY***
 DUAL-LINEAL ESTIMATES AT 1.12976
 0.05514 1.12976
 0.05514 1.12976
 0.05514 1.12976
 0.05514 1.12976

.ACTV ***ACTIVITY***
 DEFINED AS
 ROW LIMITS
 SEPARATED

DUAL-LINEAL OPTIMUM

SECTION 1 - PICS

PAGE 4 - 72/048

NUMBER	ACTIVITY	AT	ACTIVITY	SLACK ACTIVITY	LOWER LIMIT	UPPER LIMIT	DUAL ACTIVITY
1	WORK1	BS	1.12976	1.12976-	NONE	NONE	1.00000
2	PROJ1C2	PS	84.05514	13.94486	NONE	96.00000	*
3	PROJ1C3	RS	79.69370	22.29840	NONE	101.99200	*
4	PROJ1C4	UL	151.21870	*	NONE	151.21870	-0.0098-
5	GRD1C1	EC	1.00000	1.00000	1.00000	1.00000	*
6	GRD1C2	EE	1.00000	1.00000	1.00000	1.00000	*
7	GRD1C3	EC	1.00000	1.00000	1.00000	1.00000	*
8	GRD1C4	EC	1.00000	1.00000	1.00000	1.00000	*
9	GRD1C5	EC	1.00000	1.00000	1.00000	1.00000	*

PRIMAL-LINEAL OPTIMUM
SECTION 2 - COLUMNS

NUMBER	COLUMN	AT	ACTIVITY	INPUT COST	LOWER LIMIT	UPPER LIMIT	REDUCE COST
10	x1	FS	2.7000				
11	x2	FS	2.3241				
12	x3	BS	2.1000				
13	x4	BS	3.5000				
14	x5	PS	2.8000				
15	p1c1	UL	1.0000	1.8232			
16	p1c2	UL	1.0000	0.6117			
17	p1c3	UL	1.0000	0.6443			
18	p1c4	UL	1.0000	0.6443			
19	p1c5	UL	1.0000	0.5375			
20	p1c6	UL	1.0000	0.0319			
21	p1c7	UL	1.0000	0.0270			
22	p1c8	UL	1.0000	0.0230			
23	p1c9	UL	1.0000	0.0195			
24	p1c10	UL	1.0000	0.0166			
25	p1c11	UL	1.0000	0.0141			
26	p1c12	UL	1.0000	1.3376			
27	p1c13	UL	1.0000	0.0397			
28	p1c14	UL	1.0000	0.0327			
29	p1c15	UL	1.0000	0.0270			
30	p1c16	BS	0.2928	0.0223			
31	p1c17	LL		0.0184			
32	p1c18	LL		0.0152			
33	p1c19	LL		0.0125			
34	p1c20	LL		0.0104			
35	p1c21	LL		0.0086			
36	p1c22	LL		0.0071			
37	p1c23	UL	1.0000	0.0323			
38	p1c24	UL	1.0000	0.0669			
39	p1c25	UL	1.0000	0.0529			
40	p1c26	UL	1.0000	0.0418			
41	p1c27	UL	1.0000	0.0331			
42	p1c28	UL	1.0000	0.0262			
43	p1c29	UL	1.0000	0.0108			
44	p1c30	UL		0.0165			
45	p1c31	UL		0.0131			
46	p1c32	LL		0.0104			
47	p1c33	LL		0.0082			
48	p1c34	UL	1.0000	0.3869			
49	p1c35	UL	1.0000	0.0445			
50	p1c36	UL	1.0000	0.0399			
51	p1c37	UL	1.0000	0.0358			
52	p1c38	UL	1.0000	0.0221			
53	p1c39	UL	1.0000	0.0233			
54	p1c40	UL		0.0239			
55	p1c41	UL		0.0233			
56	p1c42	UL		0.0159			
57	p1c43	UL		0.0169			
58	p1c44	UL		0.0169			

PRIMAL-LINEAL OPTIMUM

NUMBER	COLUMN	AT	ACTIVITY...	INPUT CUST.	LOWER LIMIT.	UPPER LIMIT.	REDUCED COST.
57	P1C45	LL	1.00000	*25593		1.00000	*23613
62	P1C46	UL	1.00000	*03417		1.00000	*00142
61	P1C47	UL	1.00000	*03461		1.00000	*00073
62	P1C48	LL	1.00000	*05314		1.00000	*00014
63	P1C49	LL		*0572		*00000	*00052-
64	P1C50	LL		*0036		1.00000	*00033-
65	P1C51	LL		*0055		1.00000	*00127-
66	P1C52	LL		*0078		1.00000	*00167-
67	P1C53	LL		*00155		1.00000	*00204-
68	P1C54	LL		*0035		1.00000	*00214-
69	P1C55	LL		*00117		1.00000	*00170-

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PRIMAL-GLOBAL OPTIMUM
 SOLUTION (OPTIMAL)
 TIME = 0.41 SECS. ITERATION NUMBER = 28
 * * * LINE * * * * ACTIVITY * * *
 PRIMAL DEFINED AS
 DUAL LIMITS
 SEPARATION
 BOUNDS * * *

PRIMAL-GLOBAL OPTIMUM
 TIME = 0.41 SECS. ITERATION NUMBER = 28
 * * * LINE * * * * ACTIVITY * * *
 PRIMAL DEFINED AS
 DUAL LIMITS
 SEPARATION
 BOUNDS * * *

PRIMAL-GLOBAL OPTIMUM

SECTION 1 - ROWS

NUMBER	ROW	ACTIVITY	SLACK ACTIVITY	LOWER LIMIT	UPPER LIMIT	DUAL ACTIVITY
1	PCK101	BS	1.12976	1.12976	NONE	NONE
2	GRD1C2	BS	84.05514	13.94486	98.00000	1.00000
3	GRD1C3	BS	79.69340	22.29840	101.99260	*
4	GRD104	UL	159.7180	*	151.21800	*.00048-
A	5	GRD1G1	EC	1.00000	1.00000	1.00000
A	6	GRD1C2	EC	1.00000	1.00000	1.00000
A	7	GRD1G3	EC	1.00000	1.00000	1.00000
A	8	GRD1G4	EC	1.00000	1.00000	1.00000
A	9	GRD1G5	EC	1.00000	1.00000	1.00000

PRIMAL-GLOBAL OPTIMUM
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PRIMAL-GLOBAL OPTIMUM
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SECTION 2 - COLUMNS

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ORIGINAL-GLOBAL OPTIMUM

NUMBER	COLUMN	AT	ACTIVITY	INPUT COST	LOWER LIMIT	UPPER LIMIT	REDUCED COST
12	x1	BS	2.70000	*	*	*	*
11	y2	BS	2.32979	*	*	*	*
12	x3	BS	2.10000	*	*	*	*
13	x4	BS	3.50000	*	*	*	*
14	y5	BS	2.86500	*	*	*	*
15	PICC1	UL	1.70000	18232	0.617	1.6457	*
16	PICC2	UL	1.00000	0.0523	0.00000	0.00000	*
17	PICC3	UL	1.00000	0.0523	0.00000	0.00000	*
18	PICC4	UL	1.00000	0.0443	0.00000	0.00000	*
19	PICC5	UL	1.00000	0.0375	1.00000	0.00180	*
20	PICC6	UL	1.00000	0.0319	1.00000	0.00116	*
21	PICC7	UL	1.00000	0.0270	1.00000	0.00038	*
22	PICC8	UL	1.00000	0.0230	1.00000	0.00110	*
23	PICC9	UL	*	0.0195	1.00000	0.00010	*
24	PIC10	LL	*	0.0166	1.00000	0.00014	*
25	PIC11	LL	*	0.0141	1.00000	0.00030	*
26	PIC12	UL	1.00000	1.9976	1.00000	1.00000	*
27	PIC13	UL	1.00000	0.0397	1.00000	0.00200	*
28	PIC14	UL	1.00000	0.0327	1.00000	0.00122	*
29	PIC15	UL	1.00000	0.0270	1.00000	0.00035	*
30	PIC16	BS	2.29258	0.0223	1.00000	0.00048	*
31	PIC17	LL	*	0.0184	1.00000	0.00090	*
32	PIC18	LL	*	0.0152	1.00000	0.00127	*
33	PIC19	LL	*	0.0125	1.00000	0.00158	*
34	PIC20	LL	*	0.0104	1.00000	0.00187	*
35	PIC21	LL	*	0.0086	1.00000	0.00213	*
36	PIC22	LL	*	0.0071	1.00000	0.00662	*
37	PIC23	UL	1.00000	0.0323	1.00000	0.00509	*
38	PIC24	UL	1.00000	0.0669	1.00000	0.0362	*
39	PIC25	UL	1.00000	0.0529	1.00000	0.0244	*
40	PIC26	UL	1.00000	0.0418	1.00000	0.0155	*
41	PIC27	UL	1.00000	0.0331	1.00000	0.0073	*
42	PIC28	UL	1.00000	0.0262	1.00000	0.00111	*
43	PIC29	UL	1.00000	0.0208	1.00000	0.00045	*
44	PIC30	LL	*	0.0165	1.00000	0.00031	*
45	PIC31	LL	*	0.0131	1.00000	0.00019	*
46	PIC32	LL	*	0.0104	1.00000	0.00015	*
47	PIC33	LL	*	0.0082	1.00000	0.00012	*
48	PIC34	UL	1.20000	3.0676	1.00000	0.00030	*
49	PIC35	UL	1.00000	0.6445	1.00000	0.00224	*
50	PIC36	UL	1.00000	0.5319	1.00000	0.0169	*
51	PIC37	UL	1.00000	0.3352	1.00000	0.0117	*
52	PIC38	UL	1.00000	0.2321	1.00000	0.0072	*
53	PIC39	UL	1.00000	0.2248	1.00000	0.0030	*
54	PIC40	LL	*	0.1533	1.00000	0.0029	*
55	PIC41	LL	*	0.2331	1.00000	0.0046	*
56	PIC42	LL	*	0.2459	1.00000	0.0041	*
57	PIC43	LL	*	0.1613	1.00000	0.0144	*
58	PIC44	LL	*	0.0169	1.00000	0.0060	*

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PROBABILISTIC OPTIMUM						PRICE	
NUMBER	•CUST. I.	AT	••ACTIVITY••	••INPUT CUST. •	••LOWER LIMIT. •	••UPPER LIMIT. •	•REDUCED CUST. •
59	P1C45	UL	1.00000	*25593	*	1.00000	*22619
60	P1C46	UL	1.00000	*00417	*	1.00000	*0145
61	P1C47	UL	1.00000	*03361	*	1.00000	*0078
62	P1C48	UL	1.00000	*0314	*	1.00000	*0019
63	P1C49	UL	*	*02272	*	1.00000	*0035-
64	P1C50	UL	*	*03246	*	1.00000	*0063-
65	P1C51	UL	*	*0205	*	1.00000	*0127-
66	P1C52	UL	*	*00178	*	1.00000	*0167-
67	P1C53	UL	*	*00155	*	1.00000	*0204-
68	P1C54	UL	*	*00135	*	1.00000	*0238-
69	P1C55	UL	*	*00117	*	1.00000	*0270-

LJL-LIL-1000

SOLUTION (OPTIMAL)

TIME = 0.000115. INITIATION NUMBER = 31

••• NAME •••	••• ACTIVITY •••	DEFINED AS
FUNCTIONAL DEPARTMENT	1.12976	RUNOUT LIMITS SEPBUAD
EC1403•••		

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DUAL-SLUGAL OPTIMUM

SECTION 1 - ROWS

NUMBER	••• ROW •••	AT	••• ACTIVITY •••	SLACK ACTIVITY	•• LOWER LIMIT •	•• UPPER LIMIT •	•• DUAL ACTIVITY •
1	RCW101	B5	1.12976	1.12976-	NONE	NONE	1.00000
2	RCW102	B5	84.35514	13.94485	NONE	9E-00000	*
3	RCW103	B5	79.69340	22.29843	NONE	101.99266	*
4	RCW104	UL	151.2140	*	NONE	151.21800	*.00048-
A	SP10101	EC	1.00000	*	1.00000	1.00000	*
A	SP10102	EC	1.00000	*	1.00000	1.00000	*
A	SP10103	EC	1.00000	*	1.00000	1.00000	*
A	SP10104	EC	1.00000	*	1.00000	1.00000	*
A	SP10105	EC	1.00000	*	1.00000	1.00000	*

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SECTION 2 - COLUMNS
DUAL-GLOBAL OPTIMUM

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NUMBER	COLUMN#	AT	ACTIVITY...	INPUT COST..	LOWER LIMIT.	UPPER LIMIT.	REDUCED COST.
16	X1	BS	2.70CRO	*	*	*	NONE
11	X2	BS	2.32979	*	*	*	NONE
12	X3	BS	2.190R0	*	*	*	NONE
13	X4	BS	3.50LCR0	*	*	*	NONE
14	X5	BS	2.80CRO	*	*	*	NONE
15	PIC01	UL	1.29CRO	*18232	1.60C90	*1.6857	
16	PIC02	UL	1.30CRO	*00417	1.00C60	*00442	
17	PIC03	UL	1.60CRO	*00523	1.00Q00	*00343	
13	PIC04	UL	1.60CRO	*00443	1.00000	*00256	
19	PIC05	UL	1.60CRO	*00375	1.00000	*00180	
20	PIC06	UL	1.60CRO	*00319	1.00000	*00116	
21	PIC07	UL	1.60CRO	*00270	1.00000	*00058	
22	PIC08	UL	1.60CRO	*00230	1.00000	*00010	
23	PIC09	UL	1.60CRO	*00195	1.00000	*00034-	
24	PIC10	LL	*	*00166	1.00000	*00073-	
25	PIC11	UL	*	*00141	1.00000	*00107-	
26	PIC12	UL	1.60CRO	*14976	1.00000	*12404	
27	PIC13	UL	1.00CC0	*00397	1.00000	*00260	
28	PIC14	UL	1.00GRO	*00327	1.00000	*00122	
29	PIC15	UL	1.00000	*00270	1.00000	*00056	
30	PIC16	BS	2.92P8	*00223	1.00000	*	
31	PIC17	LL	*	*00184	1.00000	*00049-	
32	PIC18	UL	*	*00152	1.00000	*00090-	
33	PIC19	LL	*	*00125	1.00000	*00127-	
34	PIC20	LL	*	*0004	1.00000	*00158-	
35	PIC21	UL	*	*00086	1.00000	*00167-	
36	PIC22	UL	*	*00071	1.00000	*00213-	
37	PIC23	UL	1.60CRO	*67363	1.00000	*06622	
38	PIC24	UL	1.60CRO	*00669	1.00000	*00509	
39	PIC25	UL	1.60CRO	*00529	1.00000	*00162	
40	PIC26	UL	1.60CRO	*00418	1.00000	*00244	
41	PIC27	UL	1.60CRO	*00331	1.00000	*00150	
42	PIC28	UL	1.60CRO	*00262	1.00000	*00073	
43	PIC29	UL	1.60CRO	*00208	1.00000	*00211	
44	PIC30	LL	*	*00165	1.00000	*00040-	
45	PIC31	LL	*	*00131	1.00000	*00093-	
46	PIC32	LL	*	*001C4	1.00000	*00119-	
47	PIC33	LL	*	*00082	1.00000	*00153-	
48	PIC34	UL	1.20G00	*36696	1.00000	*35770	
49	PIC35	UL	1.60CRO	*00445	1.00000	*00224	
50	PIC36	UL	1.60CRO	*0399	1.00000	*00169	
51	PIC37	UL	1.60G00	*0358	1.00000	*00119	
52	PIC38	UL	1.60CRO	*0321	1.00000	*00072	
53	PIC39	UL	1.60CRO	*0288	1.00000	*00330	
54	PIC40	UL	*	*0259	1.00000	*00330	
55	PIC41	UL	*	*0233	1.00000	*00346-	
56	PIC42	UL	*	*0209	1.00000	*00188	
57	PIC43	UL	*	*0188	1.00000	*00113-	
58	PIC44	UL	*	*0159	1.00000	*00143-	

DUAL-GLOBAL OPTIMUM

DUAL-GLOBAL OPTIMUM							PAGE 6 = 72/048
NUMBER	•CCUWY•	AT	••ACTIVITY••	••INPUT CUST••	••LOWER LIMIT••	••UPPER LIMIT••	* P-FNUCT: CUST,
59	PIC65	LL	1.00000	* 25593	-	-	* 22619
60	PIC66	UL	1.00000	* .00417	-	-	* .00145
61	PIC67	UL	1.00000	* .00361	-	-	* .00050
62	PIC68	UL	1.00000	* .00314	-	-	* .00045
63	PIC69	LL	-	* .00272	-	-	* .00035
64	PIC70	LL	-	* .00236	-	-	* .00033
65	PIC71	LL	-	* .00205	-	-	* .00127
66	PIC72	LL	-	* .00178	-	-	* .00167
67	PIC73	LL	-	* .00155	-	-	* .00144
68	PIC74	LL	-	* .00135	-	-	* .00130
69	PIC75	LL	-	* .00117	-	-	* .00120

PAGE	R LIMIT	P-FLUKE(CJST)
1	1.00000	*22619
1	1.00000	*00145
1	1.00000	*00376
1	1.00000	*00114
1	1.00000	*00235
1	1.00000	*00083
1	1.00000	*00127
1	1.00000	*00167
1	1.00000	*00134
1	1.00000	*00216
1	1.00000	*00270

APPENDIX 3. GEOMETRIC PROGRAMMING USING MPS/360

This appendix contains the necessary programming material to solve the example geometric programming dual problem. The contents are

1. the main FORTRAN program and its supporting subroutine
2. the job control language
3. sample input data to the FORTRAN program
4. an echo check output from the FORTRAN program
5. the MPS/360 control language program with the solution output

The separable data is read with the same formats as the program of Appendix one.

NFUN = the number of dual function variables including the transformed variables.

NROWS is replaced by NCON.

NCON = the number of primal constraints.

ALOW can never equal zero (because $\ln(0)$ does not exist) and UPPER ≤ 1 .

The data for the variables are ordered as

$w_{01}, w_{02}, \dots, w_{04}, w_{11}, w_{12}, w_{21}, A_1$, and A_2 .

The variable START is set equal to the coefficients

$c_{01}, c_{02}, \dots, c_{04}, c_{11}, c_{12}, c_{21}$

in the same order as the dual variables. For the transformed variables, A_1 and A_2 , it equals 1.0. This procedure is generalized and the FORTRAN program can be used to solve any posynomial geometric programming problem.

The separable data is not punched but is stored in a data-set where it is accessed by the MPS/360 control language program. This eliminates the cumbersome process of punching cards.

```

1000 FORMAT(2I5)
1001 FORMAT(20A4)
2000 FORMAT(4X,1H#,I4,5X,6HOBJECT,4X,F12.8,3X,4HGRID,I3,3X,F12.5)
40CC FORMAT(4X,3HSFT,I3,4X,8H'MARKER',17X,8H'SEPORG')
50CC FORMAT(F10.0,I10,2F10.0)
6000 FORMAT(5X,23HECHO CHECK FOR VARIABLE,I5/5X,6HSTART=,F10.5/
15X,6HINTER=,I10/5X,6HALOW =,F10.5/5X,6HUPPER=,F10.5/)
70CC FORMAT(4X,6HEMDSET,4X,8H'MARKER',17X,8H'SEPEND')
7001 FCFORMAT(6HBOUNFS)
80CC FORMAT(1X,14H'P SEPBOUND P,I4,9X,3H1.0)
900C FORMAT(5X,35HTHE NUMBER OF SEPARATED VARIABLES =,I5/
15X,35HTHE NUMBER OF DEPENDENT VARIABLES =,I5//)
NREAD=5
NPRNT=6
NDISK=9
EIMENSION A(20)
DOUBLE PRECISON OBJ,F
COMMON KK,START,NFC
DATA SEPA,ENDA/4HSEPA,4HENDA/
REWIND NDISK
20 READ(NREAD,1001) A
IF(A(1).EQ.SEPA) GO TO 21
WRITE(NDISK,1001) A
GO TO 20
21 READ(NREAD,1000) NFUN,NCON
WRITE(NPRNT,9000) NFUN,NCON
IF(NCON.GT.NFUN) STOP9999
NFC=NFUN-NCON
KKK=0
KKK=1000
C0400KK=1,NFUN
IAX=1
READ(NREAD,5000) START,INTER,ALOW,UPPER
WRITE(NPRNT,6000) KK,START,INTER,ALOW,UPPER
IFI(INTER.GT.50)INTER=50
IFI(ALOW.EQ.UPPER)STOP8888
IFI(INTER.EQ.0)STOP7777
AN=INTER
DELTA=(UPPER-ALOW)/AN
LL=KK+100
WRITE(NDISK,4000) LL
IJ=0
IFI(ALOW-START)2,3,2
2 X=START
AX=ALOW
IJ=1
INTER=INTER+1
IAX=2
GO TO 8
3 IJ=IJ+1
XK=IJ-IAX
X=XK*DELTA+ALOW
AX=X+DELTA
8 CALL FUNC(AX,OBJ)
CALL FUNC(X,F)
OBJ=OBJ-F
GRID=X-AX
II=IJ+KKK

```

```
      WRITE(NDISK,2C00) II,OBJ,LL,GRID
      IF(IJ-INTER)3,301,301
301 KKK=KKK+INTER
      NKK=NKK+INTER
400 CONTINUE
      WRITE(NDISK,7C00)
22 READ(NREAD,1C01) A
      IF(A(1).EQ.ENFA) GO TO 23
      WRITE(NDISK,1C01) A
      GO TO 22
23 WRITE(NDISK,7C01)
      C01I=1,NKK
      II=I+1000
1  WRITE(NDISK,8C00) II
      WRITE(NDISK,1C01) A
END FILE NDISK
STOP
END
```

```
SUBROUTINE FUNC(X,F)
COMMON KK,START,NFC
DOUBLE PRECISION SS,XX,F
XX=X
IF(KK.GT.NFC) GO TO 1
SS=START
F=XX*DLOG(SS/XX)
RETURN
1 F=XX*DLOG(XX)
RETURN
END
```

```
//          JOB ( standard K.S.U. job card )
//JOBLIB DD DSNAME=SYS1.USERLIB,DISP=SHR
//STEP1  EXEC FORTGCLG
//FORT.SYSIN DD *

        FORTRAN source deck

/*
//GO.FT09F001 DD DSN=&INPUT,UNIT=SYSDA,DISP=(NEW,PASS),
//           SPACE=(TRK,(5,5)),DCB=(RECFM=FB,LRECL=80,BLKSIZE=7280)
//GO.SYSIN DD *

        data

/*
//STEP2  EXEC KSLP
//CPC.SYSIN DD *

        MPS control language program

/*
//EXEC.SYSIN DD DSN=&INPUT,UNIT=SYSDA,DISP=(OLD,PASS)
//
```

NAME DATA-SET
ROWS

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N OBJECT
E GRID101
E GRID102
E GRID103
E GRID104
E GRID105
E GRID106
F GRID107
E GRID108
E GRID109
E NORM
E ORTH1
E ORTH2
E ORTH3
E ORTH4
E CONST1
E CONST2

COLUMNS

W1	GRID101	1.0	NORM	1.0
W1	ORTH1	1.0		
W2	GRID102	1.0	NORM	1.0
W2	ORTH2	1.0		
W3	GRID103	1.0	NORM	1.0
W3	ORTH3	1.0		
W4	GRID104	1.0	NORM	1.0
W4	ORTH4	1.0		
W5	GRID105	1.0	ORTH1	2.0
W5	ORTH4	-2.0	CONST1	-1.0
W6	GRID106	1.0	ORTH2	2.0
W6	ORTH4	-2.0	CONST1	-1.0
W7	GRID107	1.0	ORTH1	-1.0
W7	ORTH2	-1.0	ORTH3	-1.0
W7	CONST2	-1.0		
A1	GRID108	1.0	CONST1	1.0
A2	GRID109	1.0	CONST2	1.0

SEPARABLE DATA

9	2		
4.0	20	.228	.238
10.0	20	.298	.308
4.0	20	.328	.338
2.0	20	.125	.135
1.0	20	.045	.055
1.0	20	.010	.020
100.0	20	.328	.338
1.0	20	.060	.070
1.0	20	.328	.338

RHS

LIMITS	GRID101	4.0	GRID102	10.0
LIMITS	GRID103	4.00	GRID104	2.0
LIMITS	GRID105	1.0	GRID106	1.0
LIMITS	GRID107	100.0	GRID108	1.0
LIMITS	GRID109	1.0	NORM	1.0

ENDATA

THE NUMBER OF SEPARATED VARIABLES = 9
THE NUMBER OF DEPENDENT VARIABLES = 2

ECHO CHECK FOR VARIABLE 1

START= 4.00000
INTER= 20
ALCW = 0.22800
UPPER= 0.23800

ECHO CHECK FOR VARIABLE 2

START= 10.00000
INTER= 20
ALCW = 0.29800
UPPER= 0.30800

ECHO CHECK FOR VARIABLE 3

START= 4.00000
INTER= 20
ALOW = 0.32800
UPPER= 0.33800

ECHO CHECK FOR VARIABLE 4

START= 2.00000
INTER= 20
ALCW = 0.12500
UPPER= 0.13500

ECHO CHECK FOR VARIABLE 5

START= 1.00000
INTER= 20
ALOW = 0.04500
UPPER= 0.05500

ECHO CHECK FOR VARIABLE 6

START= 1.00000
INTER= 20
ALOW = 0.01000
UPPER= 0.02000

ECHO CHECK FOR VARIABLE 7

START= 100.00000
INTER= 20
ALCW = 0.32800
UPPER= 0.33800

ECHO CHECK FOR VARIABLE 8

START= 1.00000
INTER= 20
ALCW = 0.06000
UPPER= 0.07000

ECHO CHECK FOR VARIABLE 9

START= 1.00000
INTER= 20
ALOW = 0.32800
UPPER= 0.33800

CONTROL PROGRAM COMPILER - MPS/360 V2-M8

C001	PROGRAM
C002	INITIALZ
C064	MOVE(XDATA,'DATA-SET')
C065	MOVE(XPBNAME,'EXAMPLE')
C066	CONVERT
C067	BCPOUT
C068	MOVE(XOBJ,'OBJECT')
0069	MCVE(XRHS,'LIMITS')
C070	TITLE('GEOMETRIC PROGRAMMING')
C071	SETUP('BOUND','SEPBOUND','MAX','SCALE')
C072	PRIMAL
C073	SOLUTION
C074	EXIT
C075	PEND

GEOMETRIC PROGRAMMING

SOLUTION (OPTIMAL)

TIME = 0.84 MINS. ITERATION NUMBER = 125

...NAME...	...ACTIVITY...	DEFINED AS
FUNCTIONAL RESTRAINTS	4.47719	OBJECT LIMITS
BCLNCS...		SEP BOUND

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GEOMETRIC PROGRAMMING

SECTION 1 - ROWS

NUMBER	...ROW...	AT	...ACTIVITY...	SLACK ACTIVITY	...LOWER LIMIT...	...UPPER LIMIT...	DUAL ACTIVITY
1	OBJECT	BS	4.47719	4.47719-	NONE	1.00000	
2	GRID01	EC	6.00000	*	4.00000	4.00000	1.63504,
3	GRID02	EC	10.00000	*	10.00000	10.00000	2.49042
4	GRID03	EC	4.00000	*	4.00000	4.00000	1.48504
5	GRID04	EC	2.00000	*	2.00000	2.00000	1.72360
6	GRID05	EC	1.00000	*	1.00000	1.00000	1.97104
7	GRID06	EC	1.00000	*	1.00000	1.00000	3.25104
8	GRID07	EC	100.00000	*	100.00000	100.00000	4.70310
9	GRID08	EC	1.00000	*	1.00000	1.00000	1.72360-
10	GRID09	EC	1.00000	*	1.00000	1.00000	*0.9886-
11	NCMPH	EC	1.00000	*	1.00000	1.00000	3.4789-
12	CATH1	EC	*	*	*	1.62649	
13	CATH2	EC	*	*	*	*4F643	
14	GRTP3	EC	*	*	*	1.99137	
15	CRT4	EC	*	*	*	1.74731	
16	CCNST1	EC	*	*	*	1.72940	
17	CCNST2	EC	*	*	*	*0.9886	

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GEOMETRIC PROGRAMMING
SECTION 2 - COLUMNS

NUMBER	COLUMN	AT	ACTIVITY	INPUT COST	LOWER LIMIT	UPPER LIMIT	REDUCED COST
18	W1	ES	*23117	*	*	*	*
19	W2	ES	*30500	*	*	*	*
20	W3	ES	*33323	*	*	*	*
21	W4	ES	*13C50	*	*	*	*
22	W5	ES	*05108	*	*	*	*
23	W6	ES	*01417	*	*	*	*
24	W7	ES	*33233	*	*	*	*
25	A1	ES	*06525	*	*	*	*
26	A2	ES	*33323	*	*	*	*
27	PIC01	UL	*66315	*	*	*	*
28	PIC02	UL	*00CC0	*00093	*	*	*
29	PIC03	UL	*00CC0	*00093	*	*	*
30	PIC04	UL	*00CC0	*00093	*	*	*
31	PIC05	UL	*00CC0	*00093	*	*	*
32	PIC06	UL	*00000	*00093	*	*	*
33	PIC07	UL	*00000	*00093	*	*	*
34	PIC08	ES	*33323	*	*	*	*
35	PIC09	UL	UL	*00092	*	*	*
36	PIC10	UL	UL	*00092	*	*	*
37	P1011	UL	UL	*00092	*	*	*
38	P1012	UL	UL	*00092	*	*	*
39	P1013	UL	UL	*00092	*	*	*
40	P1014	UL	UL	*00092	*	*	*
41	P1015	UL	UL	*00092	*	*	*
42	P1C16	UL	UL	*00092	*	*	*
43	P1C17	UL	UL	*00092	*	*	*
44	P1018	UL	UL	*00091	*	*	*
45	P1C19	UL	UL	*00091	*	*	*
46	P1C20	UL	UL	*00091	*	*	*
47	P1C21	UL	UL	*CC091	*	*	*
48	P1C22	UL	UL	1.04695	1.00000	1.00000	25.20900
49	P1C23	UL	UL	*00126	*00001	*00001	*
50	P1C24	UL	UL	*00126	*00000	*00000	*
51	P1025	UL	UL	*00125	*00000	*00000	*
52	P1C26	UL	UL	*00125	*00000	*00000	*
53	P1C27	UL	UL	*00125	*00000	*00000	*
54	P1C28	UL	UL	*00125	*00000	*00000	*
55	P1C29	UL	UL	*00125	*00000	*00000	*
56	P1C30	UL	UL	*00125	*00000	*00000	*
57	P1C31	UL	UL	*00125	*00000	*00000	*
58	P1C32	UL	UL	*00125	*00000	*00000	*
59	P1C33	UL	UL	*00125	*00000	*00000	*
60	P1C34	UL	UL	*00125	*00000	*00000	*
61	P1C35	UL	UL	*00125	*00000	*00000	*
62	P1C36	UL	UL	*00125	*00000	*00000	*
63	P1037	UL	UL	*00124	*00000	*00000	*
64	P1C38	UL	UL	*00124	*00000	*00000	*
65	P1C39	UL	UL	*00124	*00000	*00000	*
66	P1C40	UL	UL	*00124	*00000	*00000	*

NUMBER	COLUMN	AT	ACTIVITY		INPUT COST	LOWER LIMIT	UPPER LIMIT	REDUCED COST
			•	•				
67	PIC41	LL	.	.	.00124	1.00000	1.00000	1.00000
68	PIC42	LL	1.00000	.	.00124	1.00000	1.00000	1.00000
69	PIC43	LL	1.00000	.	.82034	1.00000	1.00000	1.00000
70	PIC44	UL	1.00000	.	.CCC75	1.00000	1.00000	1.00000
71	PIC45	UL	1.00000	.	.00075	1.00000	1.00000	1.00000
72	PIC46	UL	1.00000	.	.00075	1.00000	1.00000	1.00000
73	PIC47	UL	1.00000	.	.CCC75	1.00000	1.00000	1.00000
74	PIC48	UL	1.00000	.	.00075	1.00000	1.00000	1.00000
75	PIC49	UL	1.00000	.	.00075	1.00000	1.00000	1.00000
76	PIC50	UL	1.00000	.	.CCC75	1.00000	1.00000	1.00000
77	PIC51	UL	1.00000	.	.00074	1.00000	1.00000	1.00000
78	PIC52	UL	1.00000	.	.CCC74	1.00000	1.00000	1.00000
79	PIC53	UL	1.00000	.	.CCC74	1.00000	1.00000	1.00000
80	PIC54	BS	.	66667	.00074	1.00000	1.00000	1.00000
81	PIC55	LL	.	.	.CCC74	1.00000	1.00000	1.00000
82	PIC56	LL	.	.	.CCC74	1.00000	1.00000	1.00000
83	PIC57	LL	.	.	.00074	1.00000	1.00000	1.00000
84	PIC58	LL	.	.	.CCC74	1.00000	1.00000	1.00000
B5	PIC59	LL	.	.	.CCC74	1.00000	1.00000	1.00000
86	PIC60	LL	.	.	.00074	1.00000	1.00000	1.00000
87	PIC61	LL	.	.	.CCC74	1.00000	1.00000	1.00000
88	PIC62	LL	.	.	.CCC74	1.00000	1.00000	1.00000
89	PIC63	LL	.	.	.00074	1.00000	1.00000	1.00000
90	PIC64	UL	1.00000	.	.34657	1.00000	1.00000	1.00000
91	PIC65	UL	1.00000	.	.CCC89	1.00000	1.00000	1.00000
92	PIC66	UL	1.00000	.	.CCC88	1.00000	1.00000	1.00000
93	PIC67	UL	1.00000	.	.00088	1.00000	1.00000	1.00000
94	PIC68	UL	1.00000	.	.0CC88	1.00000	1.00000	1.00000
95	PIC69	UL	1.00000	.	.00088	1.00000	1.00000	1.00000
96	PIC70	UL	1.00000	.	.00088	1.00000	1.00000	1.00000
97	PIC71	UL	1.00000	.	.CCC87	1.00000	1.00000	1.00000
98	PIC72	UL	1.00000	.	.CCC87	1.00000	1.00000	1.00000
99	PIC73	LL	1.00000	.	.00087	1.00000	1.00000	1.00000
100	PIC74	UL	1.00000	.	.CCC87	1.00000	1.00000	1.00000
101	PIC75	UL	1.00000	.	.CCC87	1.00000	1.00000	1.00000
102	PIC76	LL	.	.	.00086	1.00000	1.00000	1.00000
103	PIC77	LL	.	.	.00086	1.00000	1.00000	1.00000
104	PIC78	LL	.	.	.00086	1.00000	1.00000	1.00000
105	PIC79	LL	.	.	.00086	1.00000	1.00000	1.00000
106	PIC80	UL	.	.	.00086	1.00000	1.00000	1.00000
107	PIC81	UL	.	.	.CCC85	1.00000	1.00000	1.00000
108	PIC82	UL	.	.	.00085	1.00000	1.00000	1.00000
109	PIC83	UL	.	.	.00085	1.00000	1.00000	1.00000
110	PIC84	UL	.	.	.CCC85	1.00000	1.00000	1.00000
111	PIC85	UL	1.00000	.	.13955	1.00000	1.00000	1.00000
112	PIC86	UL	1.00000	.	.00105	1.00000	1.00000	1.00000
113	PIC87	UL	1.00000	.	.CCC14	1.00000	1.00000	1.00000
114	PIC88	UL	1.00000	.	.CCC14	1.00000	1.00000	1.00000
115	PIC89	UL	1.00000	.	.CCC13	1.00000	1.00000	1.00000
116	PIC90	UL	1.00000	.	.00103	1.00000	1.00000	1.00000
117	PIC91	UL	1.00000	.	.CCC12	1.00000	1.00000	1.00000

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GEOMETRIC PROGRAMMING		AT	ACTIVITY...	INPUT COST.	LOWER LIMIT.	UPPER LIMIT.	REDUCED COST.
NUMBER	COLUMN.	UL	1.00C00	1.00C00	1.00C00	1.00C00	.00003
118	PIC92	UL	1.00000	1.00000	1.00000	1.00000	.00003
119	PIC93	UL	1.00000	1.00000	1.00000	1.00000	.00003
120	PIC94	UL	1.00C00	1.00101	1.00000	1.00000	.00002
121	PIC95	UL	1.00000	1.00100	1.00000	1.00000	.00001
122	PIC96	UL	1.00000	1.00100	1.00000	1.00000	.00001
123	PIC97	UL	1.00000	1.00099	1.00000	1.00000	.00000
124	PIC98	ES	*16667	*00099	*00099	*00099	*00099
125	PIC99	LL	*	*	*	*	*
126	P11C0	UL	*	*	*	*	*
127	P1101	LL	*	*	*	*	*
128	P11C2	LL	*	*	*	*	*
129	P11C3	LL	*	*	*	*	*
130	P1104	LL	*	*	*	*	*
131	P1105	LL	*	*	*	*	*
132	P1106	LL	1.00000	1.00098	1.00000	1.00000	.00001
133	P1107	UL	1.00000	1.00097	1.00000	1.00000	.00000
134	P1108	UL	1.00000	1.00177	1.00000	1.00000	.00014
135	P1109	UL	1.00000	1.00174	1.00000	1.00000	.00012
136	P1110	UL	1.00C00	1.00172	1.00000	1.00000	.00013
137	P1111	UL	1.00C00	1.00170	1.00000	1.00000	.00003
138	P1112	UL	1.00C00	1.00168	1.00000	1.00000	.00006
139	P1113	UL	1.00000	1.00166	1.00000	1.00000	.00004
140	P1114	UL	1.00C00	1.00164	1.00000	1.00000	.00002
141	P1115	ES	*333333	*00163	*00163	*00163	*00163
142	P1116	LL	*	*00161	*00161	*00161	*00161
143	P1117	LL	*	*00159	*00159	*00159	*00159
144	P1118	LL	*	*00158	*00158	*00158	*00158
145	P1119	LL	*	*00156	*00156	*00156	*00156
146	P1120	LL	*	*00154	*00154	*00154	*00154
147	P1121	LL	*	*00153	*00153	*00153	*00153
148	P1122	LL	*	*00152	*00152	*00152	*00152
149	P1123	LL	*	*00150	*00150	*00150	*00150
150	P1124	LL	*	*00149	*00149	*00149	*00149
151	P1125	LL	*	*00148	*00148	*00148	*00148
152	P1126	LL	*	*00146	*00146	*00146	*00146
153	P1127	UL	1.00C00	1.67613	1.00000	1.00000	470.70332
154	P1128	UL	1.00C00	*00136	*00136	*00136	*00136
155	P1129	UL	1.00000	*00236	*00236	*00236	*00236
156	P1130	UL	1.00000	*00236	*00236	*00236	*00236
157	P1131	UL	1.00C00	*00236	*00236	*00236	*00236
158	P1132	UL	1.00000	*00236	*00236	*00236	*00236
159	P1133	UL	1.00000	*00236	*00236	*00236	*00236
160	P1134	UL	1.00C00	*00235	*00235	*00235	*00235
161	P1135	UL	1.00000	*00235	*00235	*00235	*00235
162	P1136	UL	1.00C00	*00235	*00235	*00235	*00235
163	P1137	UL	1.00C00	*66667	*66667	*66667	*66667
164	P1138	ES	*	*	*	*	*
165	P1139	LL	*	*	*	*	*
166	P1140	LL	*	*	*	*	*
167	P1141	LL	*	*	*	*	*
168	P1142	LL	*	*	*	*	*

NUMBER	GEOMETRIC PROGRAMMING		INPUT CONST.	LOWER LIMIT.	UPPER LIMIT.	REDUCED CONST.	TRUE	T ₂ /T ₄)
	COLUMN	AT						
169	P1143	LL	•00235	1.00000	•00000-	•00000-	•00000-	
170	P1144	LL	•C0235	1.00000	•00000-	•00000-	•00000-	
171	P1145	LL	•C0235	1.00000	•00000-	•00000-	•00000-	
172	P1146	LL	•00235	1.00000	•00000-	•00000-	•00000-	
173	P1147	LL	•00235	1.00000	•00000-	•00000-	•00000-	
174	P1148	UL	1.00000	1.6880-	1.00000	1.00000	1.7944-	
175	P1149	UL	1.00000	•00290-	1.00000	1.00000	•00004-	
176	P1150	UL	1.00000	•00090-	1.00000	1.00000	•00004-	
177	P1151	UL	1.00000	•00090-	1.00000	1.00000	•00003-	
178	P1152	UL	1.00000	•00089-	1.00000	1.00000	•00003-	
179	P1153	UL	1.00000	•00089-	1.00000	1.00000	•00003-	
180	P1154	UL	1.00000	•CC088-	1.00000	1.00000	•00002-	
181	P1155	UL	1.00000	•CCC88-	1.00000	1.00000	•00002-	
182	P1156	UL	1.00000	•00088-	1.00000	1.00000	•00001-	
183	P1157	UL	1.00000	•CC087-	1.00000	1.00000	•00001-	
184	P1158	UL	1.00000	•CCC87-	1.00000	1.00000	•00001-	
185	P1159	BS	•50000	•00088-	1.00000	1.00000	•00000-	
186	P1160	LL	•00086-	•CC086-	1.00000	1.00000	•00000-	
187	P1161	LL	•00086-	•CC086-	1.00000	1.00000	•00001	
188	P1162	LL	•00085-	•CC085-	1.00000	1.00000	•00001	
189	P1163	LL	•00085-	•CC085-	1.00000	1.00000	•00002	
190	P1164	LL	•CC085-	•CC085-	1.00000	1.00000	•00002	
191	P1165	LL	•CC084-	•CC084-	1.00000	1.00000	•00002	
192	P1166	LL	•00084-	•CC083-	1.00000	1.00000	•00003	
193	P1167	LL	•00084-	•CC083-	1.00000	1.00000	•00003	
194	P1168	LL	•00083-	•CC083-	1.00000	1.00000	•00003	
195	P1169	UL	1.00000	•36564-	1.00000	1.00000	•43207-	
196	P1170	UL	1.00000	•CCCC6-	1.00000	1.00000	•00002-	
197	P1171	UL	1.00000	•CCCC6-	1.00000	1.00000	•00001-	
198	P1172	UL	1.00000	•00006-	1.00000	1.00000	•00001-	
199	P1173	UL	1.00000	•00005-	1.00000	1.00000	•00001-	
200	P1174	UL	1.00000	•00005-	1.00000	1.00000	•00000-	
201	P1175	UL	1.00000	•00005-	1.00000	1.00000	•00000-	
202	P1176	UL	1.00000	•00005-	1.00000	1.00000	•00000-	
203	P1177	UL	1.00000	•00005-	1.00000	1.00000	•00000-	
204	P1178	UL	1.00000	•00005-	1.00000	1.00000	•00000-	
205	P1179	UL	1.00000	•00005-	1.00000	1.00000	•00000-	
206	P1180	BS	•66667	•00005-	1.00000	1.00000	•00000-	
207	P1181	LL	•	•00005-	1.00000	1.00000	•00000-	
208	P1182	LL	•	•00005-	1.00000	1.00000	•00000-	
209	P1183	LL	•	•00005-	1.00000	1.00000	•00000-	
210	P1184	LL	•	•00005-	1.00000	1.00000	•00000-	
211	P1185	LL	•	•00005-	1.00000	1.00000	•00000-	
212	P1186	LL	•	•00005-	1.00000	1.00000	•00000-	
213	P1187	LL	•	•00004-	1.00000	1.00000	•00001	
214	P1188	LL	•	•00004-	1.00000	1.00000	•00001	
215	P1189	LL	•	•00004-	1.00000	1.00000	•00001	

APPENDIX 4. GREG PROGRAMMING MATERIAL FOR CHAPTER FIVE

This appendix contains listings of the external subroutines and the resulting output for the example problems of Chapter five.

```

C001      SUBROUTINE PHIX
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)
C003      1,Y(150),C(150),VC(50),IBAS(50),IB(100),IVC(50),IVA(100)    GEGA 20
C004      COMMON A,ALFA,X,VC,XI,XS,Y,C,VC,IBAS,IB,IVC,IVA,IVB    GEGA 30
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGESA   GEGA 40
C006      LN4,NVNINI,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IP,ICDB,JCDB,KCGEGA 50
C007      2DB,KFIL,KLIN,KREN,KFNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA 60
C008      COMMON KFNC,KGRAD,KCDBA1,KREN11,KREN21,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA 70
C009      1REN1,KREN2,KINV,KCDBA1,KREN11,KREN21,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA 80
C010      COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT    GEGA 90
C011      PHI=2.0*X(1)-0.5*X(1)**2+3.0*X(2)-0.5*X(2)**2
C012      RETURN
C013      END

```

```

C001      SUBROUTINE CPHI
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)    GEGA 20
C003      1,Y(150),C(150),VC(50),IBAS(50),IB(100),IVC(50),IVA(100)    GEGA 30
C004      COMMON A,ALFA,X,VC,XI,XS,Y,C,VC,IBAS,IB,IVC,IVA,IVB    GEGA 40
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGEGA 50
C006      LN4,NVNINI,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IP,ICDB,JCDB,KCGEGA 60
C007      2DB,KFIL,KLIN,KREN,KFNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA 70
C008      COMMON KFNC,KGRAD,KCDBA1,KREN11,KREN21,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA 80
C009      1REN1,KREN2,KINV,KCDBA1,KREN11,KREN21,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA 90
C010      COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT    GEGA 100
C011      VC(1)=XC(1)**2+XC(2)**2-1.0
C012      RETURN
C013      END

```

```

C001
C002      SUBROUTINE JACOB
          DIMENSION A(50,100),ALFA(50,50),XC(150),XI(150),XS(150)   GEGA 20
          1,Y(150),C(150),VC(50),IBAS(50),IB(100),IVC(50),IVA(100)   GEGA 30
          COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IB,IH,B,IH,B,IH,B,IH,B
          COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGEGA 50
          1N4,NVNINI,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,KCGEGA 60
          2DB,KFIL,KLIN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA 70
          COMMON KFDNC,KGRAD,KCONT,KINV1,KINV2,KINV2,KINV2,KMAX1,KMAX2,KGEKA 80
          1RENI,KREN2,KINV,KCDBA1,KREN1,KREN21                           GEGA 90
          COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT
          A(1,1)=2.0*XC(1)
          A(1,2)=2.0*XC(2)
          RETURN
          END

C006
C007
C008
C009
C010

```

```

C001
C002      SUBROUTINE GRADFI
          DIMENSION A(50,100),ALFA(50,50),XC(150),XI(150),XS(150)   GEGA 20
          1,Y(150),C(150),VC(50),IBAS(50),IB(100),IVC(50),IVA(100)   GEGA 30
          COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IB,IH,B,IH,B,IH,B,IH,B
          COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGEGA 50
          1N4,NVNINI,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,KCGEGA 60
          2DB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA 70
          COMMON KFDNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEKA 80
          1RENI,KREN2,KINV,KCDBA1,KREN1,KREN21                           GEGA 90
          COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT
          C(1)=2.0-XC(1)
          C(2)=3.0-XC(2)
          RETURN
          END

```

**THE PAGE
NUMBER ON THE
FOLLOWING PAGE
IS ILLEGIBLE DUE
TO BEING CUT OFF.**

**THIS IS AS
RECEIVED FROM
CUSTOMER.**

PARAMÈTRES

NV	2
MIN	1
NFS	0
NEVL	20
NTJ	6
ITET	20
ICDJ	0
INDAG	C
ITPAX	20
KFIL	0
KLIN	C
NCO	10
ITSOR	1
ITSLSP	20
EPSIL	0.1E-00
EPSIL0	0.1E-06
EPSIL1	0.1E-04
EPSIL2	0.1E-04
PC	0.1E-01

GRADIENT REDUIT GÉNÉRALISÉ

Nombre de variables naturelles	2
Total de variables	3
Nombre de contraintes	1
EPSILON DE NEWTON	0.1000E-06
EPSILON TEST GRADIENT	0.1000E-04

FONCTION ÉCONOMIQUE $0.22500000000000E 01$

BORNÉ INFÉRIEUR VARIABLE NATURELLE BORNE SUPÉRIEURE

X(1)	0.0	X(1)	0.500000000000E 00	X(1)	0.100000000000E 02
X(2)	0.0	X(2)	0.500000000000E 00	X(2)	0.100000000000E 02

VALEUR DES CONTRAINTES

C(1) -0.500000000000E 00

1	PHI	0.2764642437503E-01	DIR.GRAD.	NO	0	YN	0.29E-01	DELTAIFI	G.38E-02	ETA	0.24E-02	NCDB	0	NCN	2	VITN	2
2	PHI	0.29301598649931E-01	DIR.GRAD.	NO	0	YN	0.27E-01	DELTAIFI	G.35E-02	ETA	0.22E-02	NCDB	0	NCN	2	VITN	2
3	PHI	0.3065533C7611788E-01	DIR.GRAD.	NO	0	YN	0.27E-01	DELTAIFI	G.34E-02	ETA	0.21E-02	NCDB	0	NCN	2	VITN	2
4	PHI	0.30774078365141E-01	DIR.GRAD.	NO	0	YN	0.6E-01	DELTAIFI	G.44E-02	ETA	0.21E-02	NCDB	0	NCN	3	VITN	11
5	PHI	0.3105645C4392453E-01	DIR.GRAD.	NO	1	YN	0.59E-00	DELTAIFI	G.39E-00	ETA	0.39E-00	NCDB	0	NCN	4	NITN	23
6	PHI	0.31305494R121696E-01	DIR.GRAD.	NO	1	YN	0.63E-02	DELTAIFI	G.59E-01	ETA	0.59E-01	NCDB	3	NCN	3	VITN	3
7	PHI	0.3105555C7659912E-01	DIR.GRAD.	NO	1	YN	0.19E-02	DELTAIFI	G.10E-02	ETA	0.10E-02	NCDB	0	NCN	3	VITN	3
7	PHI	0.31055507659912E-01	DIR.GRAD.	NO	1	YN	0.51E-03	DELTAIFI	G.10E-02	ETA	0.54E-02	NCDB	0	NCN	3	VITN	1

MÉTODES DE CALCUL 50 CENTISECONDES

VARIABLE NATURELLE

11 = 0.554591953754431E 00
21 = 0.83212280273438E 00

VARIABLE DUALE ASSOCIEE

V(1) = 0.56648254394531E-03
V(2) = 0.

CONTRATE

8-6

WILMUS G. TOLAI: A LITERATURE

```

C001
C002      SUBROUTINE PHI(X
C003          DIMENSION A(50,100),ALFA(50,50),X(150),XI(150),XS(150)    GEGA 20
C004          1,Y(150),C(150),IBAS(50),IB(100),IVB(100),IVA(100)    GEGA 30
C005          COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IBH,IVC,IVA,IVB    GEGA 40
C006          COMMON NV,NC,NK,NEG,NIN,NTV,NVL,NEV,NEVL,NT0,NIN1,NIN2,NIN3,NIGEGA 50
C007          NV,NVNL,NVN1,NVN2,NVN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCGEGA 60
C008          2EB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA 70
C009          COMMON KFCNG,KGRAD,KCONT,KINV,KINV2,KJACO,KCDBA,KINV2,KINV1,KMAX1,KMAX2,KGEGA 80
C010          IRENL,KREN2,KINV,KCDBA1,KREN1,KREN21                           GEGA 90
C011          COMMON IUTREC,DELTIFI,ETA,JK0,LC,YSORT                         GEGA 100
C012          XX1=1.0-XC(1)
C013          XX3=1.0-XC(3)
C014          XX4=1.0-XC(4)
C015          PHI=1.0-XC(3)*(XX1*XX4)**2-XX3*(1.0-XC(2)*(1.0-XX1*XX4))**2
C016          RETURN
C017          END

```

```

C001
C002      SUBROUTINE CPHI
C003          DIMENSION A(50,100),ALFA(50,50),X(150),XI(150),XS(150)    GEGA 20
C004          1,Y(150),C(150),VC(50),IBAS(50),IB(100),IVC(50),IVA(100)    GEGA 30
C005          COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IBH,IVC,IVA,IVB    GEGA 40
C006          COMMON NV,NC,NK,NEG,NIN,NTV,NVL,NEV,NEVL,NT0,NIN1,NIN2,NIN3,NIGEGA 50
C007          NV,NVNL,NVN1,NVN2,NVN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCGEGA 60
C008          2EB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA 70
C009          COMMON KFCNG,KGRAD,KCONT,KINV,KINV2,KJACO,KCDBA,KINV2,KINV1,KMAX1,KMAX2,KGEGA 80
C010          IRENL,KREN2,KINV,KCDBA1,KREN1,KREN21                           GEGA 90
C011          COMMON IUTREC,DELTIFI,ETA,JK0,LC,YSORT                         GEGA 100
C012          VC(1)=200.0*X(1)**0.6+200.0*X(2)**0.6+200.0*X(3)**0.6
C013          1+300.0*XC(4)**0.6-800.0
C014          RETURN
C015          END

```

```

0001      SUBROUTINE JACOB
0002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)
0003      1,Y(150),C(150),VC(50),IBAS(50),IB(100),IVC(50),IVA(100)
0004      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IB,IHC,IVC,IVA,IVB
0005      COMMON NV,NC,NK,NEG,NIN,NTV,NV,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGECA
0006      IN4,NVRINI,NV$IN3,INDEX,II,IR,IRI,IS,IS1,IT,IBP,ICDB,JCDB,KCGECA
0007      ZDB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GECA
0008      COMMON KFCMC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGECA
0009      IREN1,KREN2,KINV,KCDBA,KREN1,KREN21
0010      COMMON IIREC,DELTIF,ETA,JKO,LC,YSORT
0011      A(1,1)=120.0*X(1)**(-0.4)
0012      A(1,2)=120.0*X(2)**(-0.4)
0013      A(1,3)=120.0*X(3)**(-0.4)
0014      A(1,4)=180.0*X(4)**(-0.4)
0015      RETURN
0016      END

```

```

0001      SUBROUTINE GRAFI
0002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)
0003      1,Y(150),C(150),VC(50),IBAS(50),IB(100),IVC(50),IVA(100)
0004      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IB,IHC,IVC,IVA,IVB
0005      COMMON NV,NC,NK,NEG,NIN,NTV,NV,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGECA
0006      IN4,NVNINI,NVNTN2,NVNIN3,INDEX,II,IR,IRI,IS,IS1,IT,IBP,ICDB,JCDB,KCGECA
0007      ZDB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GECA
0008      COMMON KFCMC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGECA
0009      IREN1,KREN2,KINV,KCDBA,KREN1,KREN21
0010      COMMON IIREC,DELTIF,ETA,JKO,LC,YSORT
0011      X(1)=1.0-X(1)
0012      X(2)=1.0-X(2)
0013      X(3)=1.0-X(3)
0014      X(4)=1.0-X(4)
0015      C(1)=2.0*X(3)*XX1*XX4+2.0*X(3)*(1.0-XC(2))*(1.0-XC(2))*XC(2)
0016      C(2)=2.0*X(4)*XX1*XX4*(1.0-XC(2))*(1.0-XC(2))
0017      C(3)=-(XX1*XX4)*2+(1.0-XC(2))*(1.0-XX1*XX4)**2
0018      C(4)=2.0*X(3)*XX1*XX4*XX1+2.0*XX3*(1.0-XC(2))*(1.0-XC(2))*XC(2)
0019      1*X(1)
0020      END

```

SOMMAIRE DES

XV	4
YIV	1
YEV	6
YEVL	2C
YEV2	6
YFET	20
YGET	1
ZDLS	5
ZDLS2	5
ZMAX	5C
ZFIL	0
ZKIN	C
ZCC	1C
ZTSR2	1
ZTSLSe	50
ZSIL	0.1E-96
ZSIL0	0.1E-55
ZSIL1	0.1E-55
ZSIL2	0.1E-55
ZC	0.1E-01

GRADIENT REDUIT GENERALISE

Nombre de variables naturelles	4
Nombre total de variables	5
Nombre de contraintes	1
EPSILON DE DEFINITION	2.10E-05
EPSILON TEST GRADIENT	3.10GE-05

FONCTION ECONOMIQUE 0.6F623362779617E 60

BOUCLIER INFERRUE

VARIABLE	NATURELLE
X(1)	0.500000000000E 00
X(2)	0.500000000000E 00
X(3)	0.000000000000E 00
X(4)	0.000000000000E 00

BORNES SUPERIEURES

BORNE	SUPERIEURE
X(1)	0.100000000000E 01
X(2)	0.100000000000E 01
X(3)	0.100000000000E 01
X(4)	0.100000000000E 01

VALORS DES CONTRAINTES

C1 : 1 -7.175056540425E 03

IT 1	PHI	6.77469337174395E-25	DIRECTIONS	YH 0.47E 00	DELTAIFI 0.36E 00	ETA 0.13E 03	NCDB 3
IT 2	PHI	6.775924759235E-20	DIR-COJAS	YH 2	DELTAIFI 2	NCDB 2	
IT 3	PHI	5.33543639571457E-20	DIR-CIJAS	LC 4	DELTAIFI 2	NCDB 2	
IT 4	PHI	0.31571322547768E-20	ITERATION	BASE DEGENEREE	DELTAIFI 0.21E-01	NCDB 1	
IT 5	PHI	0.00000000000000E-21	DIREGRAD	NO 2	YH 0.81E-02	NCDB 0	
PHI	6.15000000000000E-01		YH 0.6	DELTAIFI 0.0	ETA 0.0	NCDB 1	

CONTRAINTE 1 C.0

VARIABLE NATURELLE

VARIABLE ASSOCIEE

X(1) =	0.165699000000000E 01 = X(1)	V(1) = 0.0
X(2) =	0.100000000000000E 01 = X(2)	V(2) = 0.0
X(3) =	0.5374473075647E 02	V(3) = 0.0
X(4) =	0.797416965P2565E 00	V(4) = 0.0

CONTRAINTE

VARIABLE ASSOCIEE

U(1) = 0.0

INVOC DE SCOTT:	1	
INVOC DE LIBERATION:	5	
INVOC DE APPELS DU SUB CALCULAT LES CONTRAINTE	17	
INVOC DE APPELS DU SUB CALCULAT LA FONCTION ECRITIQUE	19	
INVOC DE APPELS DU SUB CALCULAT L'UNIVERS BASE DEDUITE PAR LE SUBSEQUENT DE BASE.	20	
INVOC DE APPELS DU SUB CALCULAT LES CONTRAINTE BASE	1	
INVOC DE APPELS DU SUB CALCULAT LA FONCTION ECRITIQUE	1	
INVOC DE APPELS DU SUB CALCULAT LA FONCTION ECRITIQUE	1	
INVOC DE APPELS DU SUB CALCULAT LA FONCTION ECRITIQUE	1	
INVOC DE APPELS DU SUB CALCULAT LA FONCTION ECRITIQUE	1	

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C001      SUBROUTINE PHIX
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)   GEGA 26
C003      1,Y(150),C(150),VC(50),IBAS(50),IB(100),IVC(50),IVA(100)   GEGA 30
C004      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IB,IVC,IVA,IVB   GEGA 40
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NVI,NEV,NEVL,NTO,NIN1,NIN2,NIN3,NIGEGA 50
C006      1,N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCGEGA 60
C007      2EB,KFIL,KLIN,KREN,KD,F11,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA 70
C008      COMMON KFNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA 80
C009      IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21   GEGA 90
C010      COMMON IDIREC,DELTF1,ETA,JK0,LC,YSORT   GEGA 100
C011      PHI=200.+9.*XC(1)**G.6+200.+0.*XC(2)**0.6+200.+0.*XC(3)**0.6
C012      1+300.+0.*XC(4)**0.6
C013      PHI=-PHI
C014      RETURN.
C015      END

```

```

C001      SUBROUTINE CPHI
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)   GEGA 20
C003      1,Y(150),C(150),VC(50),IBAS(50),IB(100),IVC(50),IVA(100)   GEGA 30
C004      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IB,IVC,IVA,IVB   GEGA 40
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NVI,NEV,NEVL,NTO,NIN1,NIN2,NIN3,NIGEGA 50
C006      1,N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICUB,KCGBGA 60
C007      2EB,KFIL,KLIN,KREN,KD,F11,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA 70
C008      COMMON KFNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA 80
C009      IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21   GEGA 90
C010      COMMON IDIREC,DELFI,ETA,JK0,LC,YSORT   GEGA 100
C011      XX1=1.+0.-XC(1)
C012      XX3=1.+0.-XC(3)
C013      XX4=1.+0.-XC(4)
C014      VC(1)=1.+0.-XC(2)*(XX1*XX4)**2-XX3*(1.0-XX1*XX4)**2
C015      VC(1)=-VC(1)+C.9
C016      RETURN.
C017      END

```

```

SUBROUTINE JARIB
  DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150) 20
  I,Y(15C),C(15C),VC(50),IBAS(50),IB(100),IVC(50),IVA(10C)
  COMMON A,PLFA,X,XC,XI,XS,Y,C,VC,IBAS,IB,IH,VB,GEGA 30
  COMMON NV,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NIN2,NIN3,NIGEGA 40
  IN4,NVNIN1,NVNIN2,NVNIN3,INDEX,KR,IR,IS,ISI,T,I,BP,ICDB,KCCEGA 50
  ZDC,KFL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GE GA 60
  COMMON KFLNC,KGRAD,KCONT,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGE GA 70
  IRE,JL,KREN2,KINV,KCDBA1,KREN1,KREN21,KGE GA 80
  COMMON 10IREC,DELTFI,ETA,JKO,LC,YSORT 90
  GEGA 100

C001
  COMMON 10IREC,DELTFI,ETA,JKO,LC,YSORT
  XX1=1.0-XC(1)
  XX2=1.0-XC(2)
  XX3=1.0-XC(3)
  XX4=1.0-XC(4)
  A(:,1)=2.0*X C(3)*XX1*XX4+2.0*XX3*(1.0-XC(2)*(1.0-XC(1)*XX4))+XC( 12)*XX4
  A(1,1)=-A(1,1)
  A(1,2)=2.0*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4)
  A(1,2)=-A(1,2)
  A(1,3)=-(XX1*XX4)**2+(1.0-XC(2)*(1.0-XX1*XX4))**2
  A(1,3)=-A(1,3)
  A(1,4)=2.0*X C(3)*XX1*XX4+2.0*XX3*(1.0-XC(2)*(1.0-XC(1)*XX4))+XC( 12)*XX1
  A(1,4)=-A(1,4)
  RETURN
  END

C011
  COMMON 10IREC,DELTFI,ETA,JKO,LC,YSORT
  XX1=1.0-XC(1)
  XX2=1.0-XC(2)
  XX3=1.0-XC(3)
  XX4=1.0-XC(4)
  A(1,1)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  A(1,2)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  A(1,3)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  A(1,4)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  RETURN
  END

C021
  COMMON 10IREC,DELTFI,ETA,JKO,LC,YSORT
  XX1=1.0-XC(1)
  XX2=1.0-XC(2)
  XX3=1.0-XC(3)
  XX4=1.0-XC(4)
  A(1,1)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  A(1,2)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  A(1,3)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  A(1,4)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  RETURN
  END

C031
  COMMON 10IREC,DELTFI,ETA,JKO,LC,YSORT
  XX1=1.0-XC(1)
  XX2=1.0-XC(2)
  XX3=1.0-XC(3)
  XX4=1.0-XC(4)
  A(1,1)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  A(1,2)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  A(1,3)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  A(1,4)=-(XX1*XX2*(1.0-XC(2)*(1.0-XX1*XX4))+(1.0-XX1*XX4))
  RETURN
  END

```

PARAMETRES

IV	4
IV1	1
IV2	5
IVL	29
ITD	6
ITF-T	26
ITG-N	1
ITIAG	7
ITW-LX	50
KFIL	6
KLIN	6
LCJ	10
ITS-C2	1
ITS-LP	50
EPSIL	6.1E-4
EPSIL1	6.1E-4
EPSIL2	6.1E-6
PC	0.0

GRADIENT REBUILT GENERATIVE

fonction	ecoulement	variable naturelle
fonction	total	total

fonction = $-0.77660461914063 \pm 0.1$

CORRE SUPERIEURE

fonction	ecoulement	variable naturelle
x(1)	0.6939999407907E 00	x(1)
x(1)	0.6939999407907E 00	x(2)
x(1)	0.6939999407907E 00	x(3)
x(1)	0.6939999407907E 00	x(4)

fonction = $0.6939999407907E 00$

fonction = $0.6939999407907E 00$

IT 1 $x_1 = 0.41 - 0.5 \cdot 0.230375 \cdot 0.4385 \cdot 0.3$
 IT 2 $x_1 = 0.41 - 0.5 \cdot 0.230375 \cdot 0.4385 \cdot 0.3$
 IT 3 $x_1 = 0.41 - 0.5 \cdot 0.230375 \cdot 0.4385 \cdot 0.3$
 IT 4 $x_1 = 0.41 - 0.5 \cdot 0.230375 \cdot 0.4385 \cdot 0.3$
 IT 5 $x_1 = 0.41 - 0.5 \cdot 0.230375 \cdot 0.4385 \cdot 0.3$

DUREE DU CALCUL : 6.17 SECONDES

RAISONNEMENT

VARIABLE DUALE ASSOCIEE

$x_1(1) = 0.50652200000000 = x_{11} \quad 1$
 $x_1(2) = 0.83995235234509 \quad 0.9$
 $x_1(3) = 0.50652200000000 = x_{11} \quad 3$
 $x_1(4) = 0.50652200000000 = x_{11} \quad 4$

CONSTRAINTES

VARIABLE DUALE ASSOCIEE

CONTRAINTE 1 : $-3.11926728455678 \leq 0.6$

$U(1) = 0.46287292226563 \leq 0.3$

DETAILOU DE SPLITTING

VARIABLE DUALE ASSOCIEE

INVERSE D'APPELS	0.1	EFFECTIONNE	1	CONTRAINTE	1
INVERSE D'APPELS	0.1	CALCULENT	4	CONTRAINTE	4
INVERSE D'APPELS	0.1	CALCULENT	4	FUNCTION ECONOMIQUE	30
INVERSE D'APPELS	0.1	CALCULENT	4	GRADIENT DE LA FUNCTION ECONOMIQUE	12
INVERSE D'APPELS	0.1	CALCULENT	4	L'UNIVERS AVANT	445
INVERSE D'APPELS	0.1	CALCULENT	4	L'UNIVERS APRES	445
INVERSE D'APPELS	0.1	EFFECTUER	4	DECHERCHE DE BASE	0
INVERSE D'APPELS	0.1	CALCULENT	4	JACOBIEN DES CONTRAINTES	7
INVERSE D'APPELS	0.1	EFFECTUER	4	MAX STABILITE	4
INVERSE D'APPELS	0.1	EFFECTUER	4	MAX AVANT DECHERCHE	0
INVERSE D'APPELS	0.1	EFFECTUER	4	LA PREMIERE ETAT DE BASE	16

```

C001      SUBROUTINE PHI(X
C002      CIMENSION A(50,100),ALFA(50,50),X(150),XI(150),XS(150)   GEGA 20
C003      1,Y(15C),C(15C),VC(50),IBA(5C),IHB(1CC),IVC(50),IVA(100)  GEGA 30
C004      CC#MON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVA,IVB    GEGA 40
C005      COMMON NV,NC,NK,NEG,NIN,NTV,AVI,NEV,NEVL,NTO,NINI,NIN2,NIN3,NICEGA 50
C006      IN4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IRI,IS,IS1,IT,IBP,ICDB,JCDB,KCGEGA 60
C007      2FB,KFIL,KLIN,KREN,KD,F11,PH1,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA 70
C008      COMMON KFCNC,KGRAC,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEKA 80
C009      IREN1,KREN2,KINV,KCDBAI,KREN11,KREN21                           GEGA 90
C010      COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT                         GEGA 100
C011      COMMON B(5),C(3,5)
C012      IF(IIT)100,101,101
C013      100 B(1)=ALOG(0.2)
C014      B(2)=ALOG(0.15)
C015      B(3)=ALOG(0.1)
C016      B(4)=ALOG(0.35)
C017      B(5)=ALOG(0.25)
C018      101 PHI=0.0
C019      C0102I=1,5
C020      102 PHI=PHI+ALOG(1.0-EXP(B(I)*XC(I)))
C021      RETURN
C022      END

```

```

SUBROUTINE CPHI
DIMENSION A(50,100),ALFA(50,50),X(15C),XC(150),XI(150),XS(150)
1,Y(150),C(150),VC(50),IBAS(5C),IH(50),IV(100)
COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IH,B,IVB
COMMON NV,NC,NK,NEG,NIN,NIV,NV1,NEV,NEVL,NT0,NINI,NIN2,NIN3,NIGEGA
1N4,NVNIN1,NVNIN2,NVNIN2,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCGEGA
2CB,KFIL,KLIN,KREN,KD,FI1,P1,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA
COMMON KFCNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEKA
1RENI,KREN2,KINV,KCDBAI,KREN11,KREN21
COMMON IDIREC,DELTIFI,ETA,JKO,LC,YSORT
COMMON B(5),CC(13,5)
COMMON IF(IT)100,101,101
100 CC(1,1)=1.0
CC(2,1)=7.0
CC(3,1)=7.0
CC(1,2)=2.0
CC(2,2)=7.0
CC(3,2)=8.0
CC(1,3)=3.0
CC(2,3)=5.0
CC(3,3)=8.0
CC(1,4)=4.0
CC(2,4)=9.0
CC(3,4)=6.0
CC(1,5)=2.0
CC(2,5)=4.0
CC(3,5)=9.0
101 VC(1)=0.0
    CC(2)=1.5
102 VC(1)=VC(1)+CR(1,J)*XC(J)**2
    VC(1)=VC(1)-110.0
    VC(2)=0.0
    VC(3)=0.0
    E0103J=1,5
103 VC(2)=VC(2)+CC(2,J)*(XC(J)+EXP(XC(J)/4.0))
    VC(2)=VC(2)-175.0
    VC(3)=0.0
    E0104J=1,5
104 VC(3)=VC(3)+CC(3,J)*XC(J)*EXP(XC(J)/4.0)
    VC(3)=VC(3)-2C0.0
RETURN
END

```

```

C001      SUBROUTINE JACOB
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)
C003      1,Y(150),C(150),VC(50),IBAS(50),IB(50),IVC(50),IVA(100)
C004      COMMON   A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IB,IVC,IVA,IVB
C005      COMMON   NV,NC,NK,NEG,NIN,NTV,NVI,NEV,NEVL,NTC,NINI,NIN2,NIN3,NICEGA
C006      1A4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,ISI,IT,IBP,ICDB,JCDB,KCCEGA
C007      2CB,KFIL,KLIN,KREN,KD,FI1,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA
C008      COMMON KFCNC,KGRAC,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA
C009      IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21
C010      COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT
C011      COMMON B(5),CC(3,5)
C012      E0101J=1,5
C013      101 A(1,J)=2.0*CC(1,J)*XC(J)
C014      E0102J=1,5
C015      102 A(2,J)=CC(2,J)*(1.0+0.25*EXP(XC(J)/4.0))
C016      E0103J=1,5
C017      103 A(3,J)=CC(3,J)*EXP(XC(J)/4.0)*(1.0+0.25*XC(J))
C018      RETURN
C019      END

```

```

C001      SUBROUTINE GRADFI
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)
C003      1,Y(150),C(150),VC(50),IBAS(50),IB(50),IVC(50),IVA(100)
C004      COMMON   A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IB,IVC,IVA,IVB
C005      COMMON   NV,NC,NK,NEG,NIN,NTV,NVI,NEV,NEVL,NTC,NINI,NIN2,NIN3,NICEGA
C006      1A4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,ISI,IT,IBP,ICDB,JCDB,KCCEGA
C007      2CB,KFIL,KLIN,KREN,KD,FI1,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA
C008      COMMON KFCNC,KGRAC,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA
C009      IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21
C010      COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT
C011      COMMON B(5),CC(3,5)
C012      E0101J=1,5
C013      AA=EXP(B(J)*XC(J))
C014      C(J)=B(J)*AA/(AA-1.0)
C015      RETURN
C016      END

```

PARAMETRES

NV	5
NIV	3
NEG	0
NEVL	20
NTO	6
ITET	20
ICONJ	1
ICIAJ	0
ITMAX	5C
KFIL	0
KLIN	0
NCO	1C
ITSCR	1
ISOLSR	50
EPSIL	0.1E-00
EPSILO	C.1E-04
EPSIL1	C.1E-02
EPSIL2	C.1E-02
PC	C.C

GRADIENT REDUIT GENERALISE

NCHERE	DE VARIABLES NATURELLES
NCHERE	TOTAL DE VARIABLES
NCHERE	DE CONTRAINTE
EPSILCA	DE NEWTON
EPSILON	TEST GRADIENT

5
8
3
0.1000E-04
0.1000E-02

FUNCTION ECONOMIQUE -0.26884609460831E CC

VARIABLE INFERIEURE	VARIABLE NATURELLE	HORITE	SUPERIEURE
X(1)	0.1C0C000C0C0C0E 01	X(1)	0.2C0C000C0C0C0E 01
X(2)	0.1C0C000C0C0C0E 01	X(2)	0.2C0C000C0C0C0E 01
X(3)	0.100000000000E 01	X(3)	0.200000000000E 01
X(4)	0.1C0C000C0C0C0E 01	X(4)	0.2C0C000C0C0C0E 01
X(5)	0.1C0C000C0C0C0E 01	X(5)	0.2C0C000C0C0C0E 01

VALEUR DES CONTRAINTES

C1	1)	-C.620CCCCCCCC0C0CE G2
C1	2)	-C.1B240336279297E 02
C1	3)	-C.74697174772766 E 02

IT 1	PHI=0.12854999303818E 00	DIR.GRAC.	NO 0	YN 0.19E 00	DELTAFI 0.11E 01	ETA 0.44E 00	NCDF 0	NC4 2	NITV 2
IT 2	PHI=0.105256498461E 00	DIR.CONJ.	LC 5	DIR.CONJ.	LC 5	DIR.CONJ.	NCDF 0	NC4 2	NITV 2
IT 3	PHI=0.91195169489C0E-01	DIR.CONJ.	LC 5	DIR.CONJ.	LC 5	DIR.CONJ.	NCDF 1	NC4 3	NITV 3
IT 4	PHI=0.218144C83D0307E-01	DIR.CONJ.	LC 5	DIR.CONJ.	LC 5	DIR.CONJ.	NCDF 0	NC4 4	NITV 4
IT 5	PHI=0.809409C2231124E-01	DIR.GRAC.	NO 1	YN 0.15E-01	DELTAFI 0.38E-01	ETA 0.13E-01	NCDF 0	NC4 5	NITV 5
IT 6	PHI=0.7961871212066E-01	DIR.GRAC.	NO 1	YN C.11E-01	DELTAFI 0.27E-01	ETA 0.12E-01	NCDF 0	NC4 6	NITV 6
IT 7	PHI=0.79601645469666E-01	DIR.CONJ.	LC 4	DIR.GRAC.	NO 1	DIR.GRAC.	NCDF 0	NC4 7	NITV 7
IT 8	PHI=0.79601526260376E-01	DIR.GRAC.	NO 1	YN 0.82E-03	DELTAFI 0.27E-02	ETA 0.10E-02	NCDF 0	NC4 8	NITV 8
IT 9	PHI=0.79594455702454E-01	DIR.GRAC.	NC 1	YN 0.82E-03	DELTAFI 0.27E-02	ETA 0.10E-02	NCDF 0	NC4 9	NITV 9
	PHI=0.79594459702454E-01			YN C.30E-03	DELTAFI 0.12E-02	ETA 0.55E-03			

DUREE DU CALCUL 286 CENTISECONDES

VARIABLE NATURELLE

VARIABLE DUALE ASSOCIEE

x(1) =	0.26779956817627E 01	v(1) =	-0.22783875465351E-04
x(2) =	C.25524381C8716E 01	v(2) =	0.18665520490875E-03
x(3) =	0.2070096C159302E 01	v(3) =	0.18708407878816E-03
x(4) =	0.3530631653687E 01	v(4) =	0.19506737589836E-03
x(5) =	C.2792024C124268E 01	v(5) =	0.

CONTRAINTE

CONTRAINTE 1	-C.134213397216E 02
CONTRAINTE 2	-C.220563964438E 02
CONTRAINTE 3	-C.152287890625r0E-04

VARIABLE DUALE ASSOCIEE

U(1) =	0.0
U(2) =	0.0
U(3) =	0.96090566427827E-03

NIVEAU	DE SORTIE	7	9	147
NOMBRE	C. ITERATIONS			
NOMBRE	D APPELS	CU S/P CALCULANT LES CONTRAINTEs		
NOMBRE	D APPELS	CL S/P CALCULANT LA FONCTION ECONOMIQUE	110	
NOMBRE	C APPELS	DU S/P CALCULANT LE GRADIENT DE LA FONCTION ECONOMIQUE		15
NOMBRE	D APPELS	CL S/P CALCULANT L INVERSE AVEC BASE DONNEE		
NOMBRE	D APPELS	CL S/P CALCULANT L INVERSE AVEC RECHERCHE DE BASE	6	1
NOMBRE	D APPELS	DU S/P EFFECTUANT LES CHANGEMENTS DE BASE	i	
NOMBRE	C APPELS	CU S/P CALCULANT LE JACOBIEN DES CONTRAINTEs	11	
NOMBRE	D APPELS	CL S/P EFFECTUANT UNE RECHERCHE DE MAX SANS DICHOTOMIE		12
NOMBRE	C APPELS	DU S/P EFFECTUANT UNE RECHERCHE DE MAX AVEC DICHOTOMIE		3
NOMBRE	C APPELS	CL S/P EFFECTUANT LA RENTREE DANS LE DOMAINE	31	
				TOTAL DE MAPPINGS 130

```

SUBROUTINE PHIX
DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)
1,Y(150),C(150),VC(50),IBAS(50),IHB(100),IVC(50),IVA(100) GEGK 10
COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVA,IVB GEGK 20
COMMON NV,NC,NK,NEG,NIN,NTV,NVI,NEV,NEVL,NTG,NINI,NIN2,NIN3,NIGEGK 30
1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IPB,ICDB,JCDB,KCGEGK 40
2D8,KFIL,KLIN,KREN,KD,FI1,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGK 50
COMMON KFCNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEKG 60
IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21 GEGK 70
COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT GEGK 80
DIMENSION B(10) GEGK 90
COMMON B GEGK 100
IF (IT.GE.0) GO TO 100 GEGK 110
READ (5,1001) (B(I),I=1,10) GEGK 120
1001 FORMAT (10F6.0) GEGK 130
100 CONTINUE GEGK 140
PHI = -8.*XC(1)**3-16.*XC(2)**3-20.*XC(3)**3-12.*XC(4)**3-4.*XC(5) GEGK 150
1**3-30.*XC(1)**2-39.*XC(2)**2-10.*XC(3)**2-39.*XC(4)**2-30.*XC(5)* GEGK 160
2*2+40.*XC(1)*XC(2)+20.*XC(1)*XC(3)-64.*XC(1)*XC(4)+20.*XC(1)*XC(5) GEGK 170
3+12.*XC(2)*XC(3)+62.*XC(2)*XC(4)-64.*XC(2)*XC(5)+12.*XC(3)*XC(4)+2 GEGK 180
40.*XC(3)*XC(5)+40.*XC(4)*XC(5) GEGK 190
DO 20 I=6,15 GEGK 200
PHI = PHI+B(I-5)*XC(I) GEGK 210
20 CONTINUE GEGK 220
RETURN GEGK 230
END GEGK 240

```

```

SUBROUTINE CPHI
DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)
1,Y(150),C(150),VC(50),IBAS(50),IHBA(100),IVC(50),IVA(100)
COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVA,IVB
COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NIN1,NIN2,NIN3,VIC
1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IHP,ICDB,JCDB,KCGEGL
2DB,KFIL,KLIN,KREN,KD,FI1,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGL
COMMON KFONC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGL
IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21
COMMON IDIREC,DELTFI,ETA,JK0,LC,YSORT
IF (IT.GE.0) GO TO 103
READ (5,1001) (I,J,A(I,J),L=1,37)
1001 FORMAT (6(2I3,F6.0)/6(2I3,F6.0)/6(2I3,F6.0)/6(2I3,F6.0)/6(2I3,F6.0))
1F6.0)/6(2I3,F6.0)/6(2I3,F6.0))
A(4,16) = 1.
A(5,17) = 1.
A(1,2) = 40.
A(1,3) = 20.
A(1,4) = -64.
A(1,5) = 20.
A(2,1) = 40.
A(2,3) = 12.
A(2,4) = 62.
A(2,5) = -64.
A(3,1) = 20.
A(3,2) = 12.
A(3,4) = 12.
A(3,5) = 20.
A(4,1) = -64.
A(4,2) = 62.
A(4,3) = 12.
A(4,5) = 40.
A(5,1) = 20.
A(5,2) = -64.
A(5,3) = 20.
A(5,4) = 40.
103 CONTINUE
VC(1) = 15.-6C.*XC(1)+40.*XC(2)+20.*XC(3)-64.*XC(4)+20.*XC(5)-12.*GEGL
1XC(1)**2
VC(2) = 27.+40.*XC(1)-78.*XC(2)+12.*XC(3)+62.*XC(4)-64.*XC(5)-24.*GEGL
1XC(2)**2
VC(3) = 36.+20.*XC(1)+12.*XC(2)-20.*XC(3)+12.*XC(4)+20.*XC(5)-30.*GEGL
1XC(3)**2
VC(4) = 18.-64.*XC(1)+62.*XC(2)+12.*XC(3)-78.*XC(4)+40.*XC(5)-18.*GEGL
1XC(4)**2
VC(5) = 12.+2C.*XC(1)-64.*XC(2)+20.*XC(3)+40.*XC(4)-60.*XC(5)-6.*XGEGL
1C(5)**2
DO 30 I=1,5
VC1 = 0.
DO 25 J=6,15
VC1 = VC1+A(I,J)*XC(J)
25 CONTINUE
VC(I) = VC(I)+VC1
30 CONTINUE
VC(4) = VC(4)+XC(16)
VC(5) = VC(5)+XC(17)
RETURN
END

```

```

SUBROUTINE JACOB                                GEGM 10
DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150) GEGM 20
1,Y(150),C(150),VC(50),IBAS(50),IHB(100),IVC(50),IVA(100) GEGM 30
COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVA,IVB GEGM 40
COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTD,NIN1,NIN2,NIN3,NIGEGM 50
1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,Jcdb,Kcgegm 60
2db,Kfil,Klin,Kren,Kd,FI1,Phi,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGM 70
COMMON KFNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGM 80
1REN1,KREN2,KINV,KCDBA1,KREN11,KREN21 GEGM 90
COMMON IDIREC,DELTFI,ETA,JK0,LC,YSORT GEGM 100
DIMENSION B(10) GEGM 110
COMMON B GEGM 120
A(1,1) = -24.*XC(1)-60. GEGM 130
A(2,2) = -48.*XC(2)-78. GEGM 140
A(3,3) = -60.*XC(3)-20. GEGM 150
A(4,4) = -36.*XC(4)-78. GEGM 160
A(5,5) = -12.*XC(5)-60. GEGM 170
RETURN GEGM 180
END GEGM 190

```

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SUBROUTINE GRADFI                               GEGN 10
DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)   GEGN 20
1,Y(150),C(150),VC(50),IBAS(50),IHB(100),IVC(50),IVA(100)      GEGN 30
COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVA,IVB            GEGN 40
COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NIN1,NIN2,NIN3,NIGEGN 50
1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCD,B,KCGEGN 60
2DB,KFIL,KLIN,KREN,KD,FI1,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEUN 70
COMMON KFDNC,KGRAC,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGN 80
1REN1,KREN2,KINV,KCDBA1,KREN11,KREN21                         GEGN 90
COMMON IDIREC,DELTFI,ETA,JK0,LC,YSORT                           GEGN 100
DIMENSION B(10)                                         GEGN 110
COMMON B                                              GEGN 120
IF (IT.GE.0) GO TO 100                                     GEGN 130
DO 10 I=6,15                                              GEGN 140
C(I) = B(I-5)                                            GEGN 150
10 CONTINUE                                              GEGN 160
100 CONTINUE                                             GEGN 170
C(1) = -24.*XC(1)**2-60.*XC(1)+40.*XC(2)+20.*XC(3)-64.*XC(4)+20.*XGEGN 180
1C(5)                                              GEGN 190
C(2) = -48.*XC(2)**2+40.*XC(1)-78.*XC(2)+12.*XC(3)+62.*XC(4)-64.*XGEGN 200
1C(5)                                              GEGN 210
C(3) = -60.*XC(3)**2+20.*XC(1)+12.*XC(2)-20.*XC(3)+12.*XC(4)+20.*XGEGN 220
1C(5)                                              GEGN 230
C(4) = -36.*XC(4)**2-64.*XC(1)+62.*XC(2)+12.*XC(3)-78.*XC(4)+40.*XGEGN 240
1C(5)                                              GEGN 250
C(5) = -12.*XC(5)**2+20.*XC(1)-64.*XC(2)+20.*XC(3)+40.*XC(4)-60.*XGEGN 260
1C(5)                                              GEGN 270
RETURN                                              GEGN 280
END                                                 GEGN 290

```

PARAMETRES

NV	17
NIN	3
NEG	2
NEVL	20
NTO	6
ITET	20
ICONJ	1
IDIAG	C
ITMAX	50
KFIL	C
KLIN	C
NCO	10
ITSCR	1
ISOLSR	50
EPSIL	0.1E 00
EPSILO	0.1E-03
EPSIL1	0.1E-03
EPSIL2	0.1E-03
PC	0.C

GRADIENT REDUIT GENERALISE

NCMBRE	LE VARIABLES NATURELLES
NCMBRE	TOTAL DE VARIABLES
NCMBRE	DE CONTRAINTES
EPSILON	DE NEWTON
EPSILON	TEST GRADIENT

17
23
5
0.100E-03
0.100E-03

FUNCTION ECONOMIQUE -0.890000000000E 01

BORNE INFÉRIEUR

	VARIABLE	NATURELLE	BORNE SUPERIEURE
x(1)	0.0	x(1)	0.100000000000E 03
x(2)	0.0	x(2)	0.100000000000E 03
x(3)	0.0	x(3)	-0.100000000000E 03
x(4)	0.0	x(4)	0.100000000000E 03
x(5)	0.0	x(5)	0.100000000000E 03
x(6)	0.0	x(6)	0.100000000000E 03
x(7)	0.0	x(7)	0.100000000000E 03
x(8)	0.0	x(8)	0.100000000000E 03
x(9)	0.0	x(9)	0.100000000000E 03
x(10)	0.0	x(10)	0.100000000000E 03
x(11)	0.0	x(11)	0.100000000000E 03
x(12)	0.0	x(12)	0.100000000000E 03
x(13)	0.0	x(13)	0.100000000000E 03
x(14)	0.0	x(14)	0.100000000000E 03
x(15)	0.0	x(15)	0.100000000000E 03
x(16)	0.0	x(16)	0.100000000000E 03
x(17)	0.0	x(17)	0.100000000000E 03

VALUEUR DES CONTRAINTES

c(1)	0.410000000000E 02
c(2)	-C.260000000000E 02
c(3)	C.720000000000E 02
c(4)	C.790000000000E 02
c(5)	C.0

IT 1 PHI-0.20075135690000E 11 DIR.GRAD- NO 6 YN 0.22E 11 DELTAFI 0.32E 13 ETA 0.16E 13 NCDB 0 NCN 2
 IT 2 PHI-C.1685716582400E 11 DIR.CONJ. LC 12 NO 7 YN 0.22E 11 DELTAFI 0.32E 13 ETA 0.16E 13 NCDB 0 NCN 2
 IT 3 PHI-0.1685C477536000E 11 DIR.CONJ. LC 12 NO 8 YN 0.22E 11 DELTAFI 0.32E 13 ETA 0.16E 13 NCDB 0 NCN 2
 IT 4 PHI-0.14910701568000E 11 ITERATION SPECIALE BASE DEGENEREE NO 9 YN 0.11E 11 DELTAFI 0.21E 13 ETA 0.83E 12 NCDB 1 NCN 1
 IT 5 PHI-0.14100354049000E 11 DIR.GRAD- NO 10 YN 0.41E 10 DELTAFI 0.21E 13 ETA 0.70E 18 NCDB 1 NCN 1
 IT 6 PHI-0.1495215744C0000E 10 DIR.GRAD- NO 11 YN 0.12E 11 DELTAFI 0.21E 13 ETA 0.85E 12 NCDB 1 NCN 1
 IT 7 PHI-0.52067491B40000E 10 DIR.GRAD- NO 12 YN 0.14E 11 DELTAFI 0.19E 13 ETA 0.11E 13 NCDB 2 NCN 5
 IT 8 PHI-0.3124048849C000E 10 DIR.GRAD- NO 13 YN 0.12E 09 DELTAFI 0.31E 10 ETA 0.31E 10 NCDB 0 NCN 3
 IT 9 PHI-G.1130372C09217E 03 DIR.GRAD- NO 14 YN 0.12E 09 DELTAFI 0.31E 10 ETA 0.31E 10 NCDB 0 NCN 3
 LES VARIABLES ARTIFICIELLES SONT TOUTES ANNULEES NO 15 YN 0.92E 02 DELTAFI 0.12E 05 ETA 0.56E 04 NCDB 1 NCN 23
 IT 10 PHI-0.652564920898E 02 DIR.GRAD- NO 16 YN 0.42E 02 DELTAFI 0.58E 04 ETA 0.39E 04 NCDB 1 NCN 4
 IT 11 PHI-0.5215751171828E 02 DIR.GRAD- NO 17 YN 0.93E 01 DELTAFI 0.97E 03 ETA 0.60E 03 NCDB 1 NCN 5
 IT 12 PHI-0.4240324401855E 02 DIR.GRAD- NO 18 YN 0.40E 01 DELTAFI 0.25E 03 ETA 0.18E 03 NCDB 0 NCN 3
 IT 13 PHI-0.39398252868652E 02 DIR.GRAD- NO 19 YN 0.68E 01 DELTAFI 0.18E 03 ETA 0.18E 03 NCDB 0 NCN 2
 IT 14 PHI-0.393982358624023E 02 DIR.CONJ. LC 2 YN 0.29E 01 DELTAFI 0.18E 03 ETA 0.18E 03 NCDB 0 NCN 3
 IT 15 PHI-0.37773712158203E 02 DIR.CONJ. LC 3 YN 0.29E 01 DELTAFI 0.18E 03 ETA 0.13E 03 NCDB 0 NCN 1
 IT 16 PHI-0.37434371948247E 02 DIR.GRAD- NO 20 YN 0.28E 01 DELTAFI 0.12E 03 ETA 0.12E 03 NCDB 0 NCN 2
 IT 17 PHI-0.3704394C053711E 02 DIR.GRAD- NO 21 YN 0.25E 01 DELTAFI 0.12E 03 ETA 0.11E 03 NCDB 0 NCN 1
 IT 18 PHI-0.3677922532226E 02 DIR.GRAD- NO 22 YN 0.23E 01 DELTAFI 0.12E 03 ETA 0.10E 03 NCDB 0 NCN 2
 IT 19 PHI-0.34977713999258E 02 DIR.GRAD- NO 23 YN 0.23E 01 DELTAFI 0.12E 03 ETA 0.10E 03 NCDB 0 NCN 2
 IT 20 PHI-0.32819366455078E 02 DIR.CONJ. LC 4 YN 0.31E 01 DELTAFI 0.12E 03 ETA 0.13E 03 NCDB 0 NCN 1
 IT 21 PHI-0.32717651367118E 02 DIR.CONJ. LC 5 YN 0.31E 01 DELTAFI 0.12E 03 ETA 0.13E 03 NCDB 0 NCN 1
 IT 22 PHI-0.32351867124660154F 02 DIR.GRAD- NO 24 YN 0.19E 01 DELTAFI 0.12E 03 ETA 0.18E 03 NCDB 0 NCN 2
 IT 23 PHI-0.323722CC277347E 02 DIR.CONJ. LC 6 YN 0.53E 00 DELTAFI 0.18E 02 ETA 0.11E 02 NCDB 0 NCN 2
 IT 24 PHI-0.3236625203472E 02 DIR.GRAD- NO 25 YN 0.55E 00 DELTAFI 0.18E 02 ETA 0.11E 02 NCDB 0 NCN 2
 IT 25 PHI-0.32366222509765E 02 DIR.CONJ. LC 7 YN 0.53E 00 DELTAFI 0.10E 02 ETA 0.10E 02 NCDB 0 NCN 3
 IT 26 PHI-G.3236021421339E 02 DIR.CONJ. LC 8 YN 0.53E 00 DELTAFI 0.10E 02 ETA 0.99E 01 NCDB 0 NCN 8
 IT 27 PHI-0.32351867124660154F 02 DIR.GRAD- NO 27 YN 0.16E 00 DELTAFI 0.41E 01 ETA 0.36E 01 NCDB 0 NCN 8
 IT 28 PHI-0.32351364135742E 02 DIR.GRAD- NO 28 YN 0.12E 00 DELTAFI 0.41E 01 ETA 0.36E 01 NCDB 0 NCN 8
 IT 29 PHI-0.323506492749023E 02 DIR.CONJ. LC 9 YN 0.12E 00 DELTAFI 0.41E 01 ETA 0.11E 02 NCDB 0 NCN 2
 IT 30 PHI-0.323503817573242E 02 DIR.GRAD- NO 31 YN 0.47E-01 DELTAFI 0.34E 01 ETA 0.18E 01 NCDB 0 NCN 3
 IT 31 PHI-0.323502246985352E 02 DIR.GRAD- NO 32 YN 0.59E-01 DELTAFI 0.34E 01 ETA 0.43E 01 NCDB 0 NCN 4
 IT 32 PHI-0.3235015669140E 02 DIR.GRAD- NO 33 YN 0.37E-01 DELTAFI 0.33E 01 ETA 0.18E 01 NCDB 0 NCN 2
 IT 33 PHI-0.3234938649316E 02 DIR.GRAD- NO 34 YN 0.37E-01 DELTAFI 0.33E 01 ETA 0.18E 01 NCDB 0 NCN 2
 IT 34 PHI-0.32348738624023E 02 DIR.CONJ. LC 10 YN 0.67E-02 DELTAFI 0.33E 00 ETA 0.31E 00 NCDB 0 NCN 3
 IT 35 PHI-0.32348709106445E 02 DIR.CONJ. LC 11 YN 0.67E-02 DELTAFI 0.33E 00 ETA 0.31E 00 NCDB 0 NCN 0
 IT 36 PHI-0.32348678598867E 02 DIR.CONJ. LC 12 YN 0.67E-02 DELTAFI 0.33E 00 ETA 0.31E 00 NCDB 0 NCN 0

VARIABLE	NATURELLE	VARIABLE	DUALE	ASSOCIEE
X(1)	= C.2994460599371E 00	V(1)	= 0.	
X(2)	= C.33329182863735E 00	V(2)	= -0.69427490234375E- 03	
X(3)	= C.40660669183731E 00	V(3)	= 0.	
X(4)	= C.428763628C0598E 00	V(4)	= 0.	
X(5)	= C.22434639930125E 00	V(5)	= 0.	
X(6)	= C.0	V(6)	= -0.36304092407227E 02	
X(7)	= C.0	V(7)	= -0.1957429809570E 01	
X(8)	= C.51755962371826E 01	V(8)	= -0.3147244453302E- 02	
X(9)	= C.0	V(9)	= -0.13940868377686E 01	
X(10)	= C.3061566271323E 01	V(10)	= 0.	
X(11)	= C.11632645416260E 02	V(11)	= 0.352478027734375E-02	
X(12)	= C.0	V(12)	= -0.3813537597656E 02	
X(13)	= C.0	V(13)	= -0.56750320434570E 02	
X(14)	= C.102776587C0243E 00	V(14)	= -0.46758651733398E-02	
X(15)	= C.0	V(15)	= -0.68643802404404E 00	
X(16)	= C.0	V(16)	= -0.42875647544861E 00	
X(17)	= C.0	V(17)	= -0.22437608242035E 00	

VARIABLE	DUALE	ASSOCIEE	
CONTRAINTE			
CONTRAINTE 1	-0.11444091796875E-03	U(1)	= 0.29944628477097E 00
CONTRAINTE 2	-0.61035156220000E-04	U(2)	= 0.3332568737559E 00
CONTRAINTE 3	-0.9155273437500E-04	U(3)	= -0.40061430765152E 00
CONTRAINTE 4	0.57220458994375E-05	U(4)	= 0.4287564754481E 00
CONTRAINTE 5	0.17166137695313E-04	U(5)	= 0.22437608242035E 00

NIVEAU	DE SORTIE	1	36
NOMBRE	D ITERATIONS		
NOMBRE	D APPELS	DU S/P	CALCULANT LES CONTRAINTES 466
NOMBRE	D APPELS	DU S/P	CALCULANT LA FONCTION ECONOMIQUE 298
NOMBRE	D APPELS	DU S/P	CALCULANT LE GRADIENT DE LA FONCTION ECONOMIQUE 104
NOMBRE	D APPELS	DU S/P	CALCULANT L INVERSE AVEC BASE DONNEE RECHERCHE DE BASE 39
NOMBRE	D APPELS	DU S/P	CALCULANT L INVERSE AVEC RECHERCHE DE BASE 0
NOMBRE	D APPELS	DU S/P	EFFECTUANT LES CHANGEMENTS DE BASE 9
NOMBRE	D APPELS	DU S/P	CALCULANT LE JACOBIEN DES CONTRAINTES 43
NOMBRE	D APPELS	DU S/P	EFFECTUANT UNE RECHERCHE DE MAX SANS DICHOTOMIE 33
NOMBRE	D APPELS	DU S/P	EFFECTUANT UNE RECHERCHE DE MAX AVEC DICHOTOMIE 3
NOMBRE	D APPELS	DU S/P	EFFECTUANT LA RENTREE DANS LE DOMAINE 115
			NOMBRE TOTAL D ITERATIONS 357

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C001      SUBROUTINE PHI(X
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)
C003      1,Y(15C),C(15C),VC(5C),IBA(5C),IB((5C),IVB((10C),
C004      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IBB,IVC,IVA,IVB
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NT0,NINI,NIN2,NIN3,NICEGA
C006      LN4,NVNINI,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCCEGA
C007      2EB,KFIL,KLIN,KREN,KD,FI1,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA
C008      COMMON KFNC,KGRAC,KCONT,KINV1,KINV2,KMAX1,KMAX2,KCEGA
C009      IREN1,KREN2,KINV,KCDBA,KJACO,KJAC1,KREN11,KREN21
C010      COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT
C011      COMMON B(5),CC(3,4)
C012      IF(ITT)100,11C,110
C013      100 E(1)=-40792.14
C014      E(2)=37.29324
C015      E(3)=0.835689
C016      E(4)=5.357855
C017      11C  PHI=B(1)+B(2)*XC(1)+B(3)*XC(1)*XC(5)+B(4)*XC(3)**2
C018      PHI=PHI+30373.94
C019      PHI=-PHI
C020      RETURN
C021      END

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      SUBROUTINE CPH'I
      DIMENSION A(50,100),ALFA(50,50),X(150),XI(150),XS(150)
      1,Y(150),C(150),VC(50),IBAS(50),IHBI(50),IVC(50),IVA(100)
      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVA,IVB
      COMMON NV,NC,NK,NEG,NIN,NTV,NVI,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGE
      1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IPB,ICDB,JCDB,KCCEGA
      2CB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEA
      COMMON KFCNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGE
      1REN1,KREN2,KINV,KCDBA1,KREN11,KREN21
      COMMON ICIREC,DELTFI,ETA,JKO,LC,YSORT
      COMMON B(5),CC(3,4)
      IF(IT)100,11C,11C
      100 CC(1,1)=85.33441
      CC(1,2)=0.005686
      CC(1,3)=0.000626
      CC(1,4)=-0.002205
      CC(2,1)=80.51249
      CC(2,2)=0.007132
      CC(2,3)=0.002996
      CC(2,4)=0.002181
      CC(3,1)=9.300961
      CC(3,2)=0.004703
      CC(3,3)=0.001255
      CC(3,4)=0.00108
      CC(1)=CC(1,1)+CC(1,2)*XC(2)*XC(5)+CC(1,3)*XC(1)*XC(4)+CC(1,4)
      1*XC(3)*XC(5)
      VC(2)=-VC(1)
      VC(1)=VC(1)-92.0
      VC(3)=CC(2,1)+CC(2,2)*XC(2)*XC(5)+CC(2,3)*XC(1)*XC(2)+CC(2,4)
      1*XC(3)*#2
      VC(4)=-VC(3)+90.0
      VC(3)=VC(3)-110.0
      VC(5)=CC(3,1)+CC(3,2)*XC(3)*XC(5)+CC(3,3)*XC(1)*XC(3)+CC(3,4)
      1*XC(3)*XC(4)
      VC(6)=-VC(5)+20.0
      VC(5)=VC(5)-25.0
      RETURN
      END

      CC01
      CC02
      CC03
      CC04
      CC05
      CC06
      CC07
      CC08
      CC09
      CC10
      CC11
      CC12
      CC13
      CC14
      CC15
      CC16
      CC17
      CC18
      CC19
      CC20
      CC21
      CC22
      CC23
      CC24
      CC25
      CC26
      CC27
      CC28
      CC29
      CC30
      CC31

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C001      SUBROUTINE JAF08
C002      DIMENSION A(50,100),ALFA(50,50),XC(150),XI(150),XS(150)    GEGA 20
C003      1,Y(150),C(150),VC(50),IBAS(50),IVC(50),IVA(100)    GEGA 30
C004      COMMON A,ALFA,X,VC,XI,XS,Y,C,VC,IBAS,IBB,IVC,IVA,IVB    GEGA 40
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NFV,NEVL,NTO,NIN1,NIN2,NIN3,NIEGA 50
C006      IN4,NVNINI,NVNIN2,NVNIN2,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCCEGA 60
C007      2TB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSILC,EPSIL2,SEGAG 70
C008      COMMON KFNC,KGRAC,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA 80
C009      IREN1,KRENZ,KIRV,KCDBAI,KREN11,KREN21    GEGA 90
C010      COMMON IDIREC,DELTIF,ETA,JKO,LC,YSORT    GEGA 100
C011      COMMON B(5),CC(3,4)
C012      A(1,1)=CC(1,3)*XC(4)
C013      A(1,2)=CC(1,2)*XC(5)
C014      A(1,3)=CC(1,4)*XC(5)
C015      A(1,4)=CC(1,3)*XC(1)
C016      A(1,5)=CC(1,2)*XC(2)+CC(1,4)*XC(3)
C017      A(3,1)=CC(2,3)*XC(2)
C018      A(3,2)=CC(2,2)*XC(5)+CC(2,3)*XC(1)
C019      A(3,3)=2.0*CC(2,4)*XC(3)
C020      A(3,5)=CC(2,2)*XC(2)
C021      A(5,1)=CC(3,3)*XC(3)
C022      A(5,3)=CC(3,2)*XC(5)+CC(3,4)*XC(4)+CC(3,3)*XC(1)
C023      A(5,4)=CC(3,4)*XC(3)
C024      A(5,5)=CC(3,2)*XC(3)
C025      C011=2,6,2
C026      EC1J=1,5
C027      1 A(I,J)=-A(I-1,J)
C028      RETURN
C029      END

```

```

SUBROUTINE GRADFI
C001   DIMENSION A(50,100),ALFA(50,50),X(150),XI(150),XS(150)    GEGA 20
C002   1,Y(150),C(15C),VC(5C),IBA(5C),IB(1CC),IVC(50),IVA(100)  GEGA 30
C003   COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IRAS,IHB,IVC,IVA,IVB  GEGA 40
C004   COMMON NV,NC,NK,NEG,NIN,NTV,KV1,NEV,NEVL,NTO,NIN1,NIN2,NIN3,NICEGA 50
C005   IN4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,Jcdb,KCGEGA 60
2LB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSTL,EPSTL2,CGEGA 70
COMMON KFCNC,KGRAC,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA 80
IREN1,KREN2,KIRY,KCOBA1,KREN11,KREN21  GEGA 90
COMMON IDIREC,DELTIFI,ETA,JK0,LC,YSORT  GEGA 100
COMMON B(5),CC(3,4)
C(1)=-B(2)-B(3)*XC(5)
C(3)=-2*0*B(4)*XC(3)
C(5)=-B(3)*XC(1)
RETURN
END
C006
C007
C008
C009
C010
C011
C012

```

PARAMETRES

NV	5
NIN	6
NEG	0
NEVL	20
NTO	6
ITET	20
ICONJ	1
ICLAG	0
ITMAX	50
KFIL	0
KLIN	0
NCO	10
ITSCR	1
ISOLSR	10
EPSIL	0.1E 00
EPSILO	0.1E-04
EPILL1	0.1E-02
EPILL2	0.1E-02
PC	0.C

GRADIENT RÉSULTAT GÉNÉRALISÉ

Nombre de variables naturelles 5
 Nombre total de variables 11
 Nombre de contraintes 6
 EPSILON DE NEWTON 0.1000E-04
 EPSILON TEST GRADIENT 0.1000E-02

FONCTION ÉCONOMIQUE 0.11718750000000E-01

	BORNÉ INFÉRIEUR	VARIABLE NATURELLE	BORNÉ SUPÉRIEUR
X1(1)	0.780000000000E 02	X(1) 0.78619995117188E 02	X5(1) 0.102000000000E 03
X1(2)	0.320000000000E 02	X(2) 0.33439987182517E 02	X5(2) 0.420000000000E 02
X1(3)	0.270000000000E 02	X(3) 0.3106992065630E 02	X5(3) 0.450000000000E 02
X1(4)	0.270000000000E 02	X(4) 0.44099990844727E 02	X5(4) 0.450000000000E 02
X1(5)	0.270000000000E 02	X(5) 0.35219985961914E 02	X5(5) 0.450000000000E 02

VALEUR DES CONTRAINTES

C(1) -C.21136474609375E 00
 C(2) -C.91788635251906E 02
 C(3) -C.111C5728149414E 02
 C(4) -C.88942718503859E 01
 C(5) -C.4872111816406E 01
 C(6) -C.1272881835938E 00

IT 1	PHI 0.11817576125000E 03	DIR.GRAC.	NO 2	YN 0.35E 03	DELTAFI 0.19E 04	ETA 0.14E 04	NCDF 3	NCN 2
IT 2	PHI C.118234315000C0E 03	DIR.CONJ.	LC 3	DIR.CONJ.	DIR.CONJ.	DIR.CONJ.	NCIS 0	NCY 2
IT 3	PHI 0.118289C6250C0E 03	DIR.CONJ.	LC 3	DIR.CONJ.	DIR.CONJ.	DIR.CONJ.	NCIS 0	NCY 2
IT 4	PHI 0.25400339625000E 03	DIR.GRAC.	NO 2	YN 0.95E 03	DELTAFI 0.57E 03	ETA 0.70E 03	NCDS 2	NCY 20
IT 5	PHI 0.*25632031250000E 03	DIR.GRAC.	NO 3	YN 0.83E 03	DELTAFI 0.38E 02	ETA 0.38E 02	NCDB 0	NCY 2
IT 6	PHI 0.*29011718750000E 03	DIR.GRAU.	NC 4	YN 0.83E 03	DELTAFI 0.26E 02	ETA 0.36E 02	NCIR 0	NCY 7
	PHI 0.29201171875000E 03			YN 0.81E 03	DELTAFI 0.43E-02	ETA 0.43E-02		NCY .2

DUREE EU CALCUL 165 CENTISECONDES

VARIABLE NATURELLE VARIABLE DUALE ASSOCIEE

x(1) =	0.78000000000000E 02	= x(1)	y(1) =	-0.48917602539063E 02
x(2) =	0.33000000000000E 02	= x(2)	y(2) =	-0.8309921264648E 02
x(3) =	0.299993896484375E 02	= x(3)	y(3) =	0.
x(4) =	0.*45C000CCCC000E 02	= x(4)	y(4) =	0.*26632446289063E 02
x(5) =	0.367776123046875E 02	= x(5)	y(5) =	0.

CONTRAINTE VARIABLE DUALE ASSOCIEE

CCNTRAINTE 1	-C.15258789062500E-04	U(1) =	0.40318652343750E 03
CCNTRAINTE 2	-C.9199998474121E 02	U(2) =	0.0
CCNTRAINTE 3	-C.1115824590134E 02	U(3) =	0.0
CCNTRAINTE 4	-0.*8841751098632E 01	U(4) =	0.0
CCNTRAINTE 5	-C.*999984741219E 01	U(5) =	0.0
CCNTRAINTE 6	-C.15258789062500E-04	U(6) =	0.80937573242188E 03

NIVEAU DE SORTIE	7	NOMBRE D'ITERATIONS	6	NOMBRE	C APPELS DU S/P CALCULANT LES CONTRAINTE	93	NOMBRE	C APPELS DU S/P CALCULANT LA FONCTION ECONOMIQUE	69	NOMBRE	D APPELS DU S/P CALCULANT LE GRADIENT DE LA FONCTION ECONOMIQUE	14
				NOMBRE	C APPELS DU S/P CALCULANT L INVERSE AVEC DONNEE RECHERCHE DE BASE	4	NOMBRE	C APPELS DU S/P CALCULANT L INVERSE AVEC RECHERCHE DE BASE	1			
				NOMBRE	D APPELS DU S/P EFFECTUANT LES CHANGEMENTS DE BASE	3						
				NOMBRE	D APPELS DU S/P CALCULANT LE JACOBIEN DES CONTRAINTES	8						
				NOMBRE	C APPELS DU S/P EFFECTUANT UNE RECHERCHE DE MAX SANS DICHOTOMIE	14	NOMBRE	C APPELS DU S/P EFFECTUANT UNE RECHERCHE DE MAX AVEC DICHOTOMIE	0	NOMBRE	TOTAL D'ITERATIONS	74
				NOMBRE	C APPELS DU S/P EFFECTUANT LA RENTREE DANS LE DOMAINE	26						

```

C001      SUBROUTINE PHIX
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XI(150),XS(150)
C003      1,Y(150),C(150),VC(50),IBAS(50),IHB(100),IVC(50),IVA(100)
C004      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVI,IVB
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGEGA
C006      IN4,NVNINI,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCGEGA
C007      2DB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSSIL,EPSSIL0,EPSSIL2GEGA
C008      COMMON KFCNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA
C009      IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21
C001      COMMON IDTREC,DELTFI,ETA,JK0,LC,YSORT
C002      PHI=-4.0*XC(1)-10.0*XC(2)-4.0*XC(3)-2.0*SQRT(XC(1)*#2+XC(2)*#2)
C003      RETURN
C004      END

```

```

C001      SUBROUTINE CPHI
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XI(150),XS(150)
C003      1,Y(150),C(150),VC(50),IBAS(50),IHB(100),IVC(50),IVA(100)
C004      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVI,IVB
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGEGA
C006      IN4,NVNINI,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCGEGA
C007      2DB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSSIL,EPSSIL0,EPSSIL2GEGA
C008      COMMON KFCNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEGA
C009      IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21
C001      COMMON IDTREC,DELTFI,ETA,JK0,LC,YSORT
C002      VC(1)=-XC(1)*XC(2)*XC(3)+100.C
C003      RETURN
C004      END

```

```

C001      SUBROUTINE JACOB
C002      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)    GEGA 20
C003      1,Y(150),C(150),VC(50),IBAS(50),IH(50),IVC(50),IVA(100)    GEGA 30
C004      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IH,IVC,IVA,IVB    GEGA 40
C005      COMMON NV,NC,NK,NEG,NIN,NTV,NW1,NEV,NEVL,NT0,NINI,NIN2,NIN3,NIGEGA 50
C006      1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCGEGA 60
C007      2CFB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2,GECA 70
C008      COMMON KFNC,KGRAD,KCDBA,KINV2,KINV1,KJACO,KMAX1,KMAX2,KGEGA 80
C009      1RENI,KREN2,KINV,KCDBA1,KREN11,KREN21    GEGA 90
C010      COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT    GEGA 100
C011      END

C006      SUBROUTINE GRADFI
C007      DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)    GEGA 20
C008      1,Y(150),C(150),VC(50),IBAS(50),IH(50),IVC(50),IVA(100)    GEGA 30
C009      COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IH,IVC,IVA,IVB    GEGA 40
C010      COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NT0,NINI,NIN2,NIN3,NIGEGA 50
C011      1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCDB,KCGEGA 60
C012      2CFB,KFIL,KLIN,KREN,KD,FIL,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2,GECA 70
C013      COMMON KFNC,KGRAD,KCDBA,KINV2,KINV1,KJACU,KMAX1,KMAX2,KGEGA 80
C014      1RENI,KREN2,KINV,KCDBA1,KREN11,KREN21    GEGA 90
C015      COMMON IDIREC,DELTFI,ETA,JKO,LC,YSORT    GEGA 100
C016      SS=SGRT(XC(1)**2+XC(2)**2)
C017      C(1)=-4.0-2.0*XC(1)/SS
C018      C(2)=-10.0-2.0*XC(2)/SS
C019      C(3)=-4.0
C020      RETURN
C021      END

```

PARAMETRES

NV	3
NIV	1
NE5	0
NFVL	20
NT0	6
ISTT	20
ICD1J	1
IDIAG	0
ITMAX	50
KFIL	0
KLIN	0
NC0	10
ITSCR	1
ISLSR	50
EPSIL	0.1E 00
EPSIL0	0.1E-04
EPSIL1	0.1E-02
EPSIL2	0.1E-02
PC	0.1E 01

GRADIENT REDUIT GENERALISE

NCBRE DE VARIABLES NATURELLES 3
 NCBRE TOTAL DE VARIABLES 4
 NCBRE DE CONTRAINTES 1
 EPSILON DE NEWTON 0.1000E-04
 EPSILON TEST GRADIENT 0.1000E-02

FONCTION ECONOMIQUE -0.2E28427124023E 03

	BORNE INFERIEURE	VARIABLE NATURELLE	BORNE SUPERIEURE
X(1)	0.0	X(1)	0.10000000000000E 02
X(2)	0.0	X(2)	0.10000000000000E 02
X(3)	0.0	X(3)	0.10000000000000E 02

VALEUR DES CONTRAINTES

C(1) -0.90000000000000E 03

IT 1	PHI-0.1327776487578E 03	DIR.GRAD.	NO C	YN 0.13E 02	DELTAFI 0.21E 03	ETA 0.14E 03	NCLH 2	NCV 2
IT 2	PHI-0.92377822875977E 02	DIR.CONJ.	LC 3	DIR.CONJ.	LC 3	DIR.CONJ.	NCBL 2	NCV 2
IT 3	PHI-0.89144851C742189E 02	DIR.CONJ.	LC 3	DIR.CONJ.	LC 3	DIR.CONJ.	NCBL 2	NCV 2
IT 4	PHI-C.69113655467963E 02	DIR.CONJ.	LC 3	DIR.CONJ.	LC 3	DIR.CONJ.	NCBL 3	NCV 3
IT 5	PHI-0.867149658203E 02	DIR.CONJ.	NC 1	YN 0.28E 01	DELTAFI 0.21E 03	ETA 0.24E 03	NCBL 3	NCV 3
IT 6	PHI-0.B735727319063E 02	DIR.CONJ.	LC 2	DIR.CONJ.	LC 2	DIR.CONJ.	NCBL 3	NCV 3
IT 7	PHI-0.87990325927734F 02	DIR.CONJ.	LC 2	DIR.CONJ.	LC 2	DIR.CONJ.	NCBL 3	NCV 3
IT 8	PHI-0.87877747192383E 02	DIR.GRAD.	NO Q	YN 0.29E 00	DELTAFI 0.39E 00	ETA 0.24E 00	NCBL 0	NCV 10
IT 9	PHI-0.87787747192383E 02	DIR.GRAD.	NO 1	YN 0.18E-01	DELTAFI 0.39E 00	ETA 0.17E 01	NCBL 0	NCV 10
IT q	PHI-0.87987747192383E 02	DIR.GRAD.	YN 0.68E-02	DELTAFI 0.42E-01	ETA 0.34E-01	ETA 0.34E-01	NCBL 0	NCV 6

DUREE DU CALCUL	83 CENTISECONDES
VARIABLE NATURELLE	VARIABLE DUALE ASSOCIEE

$$\begin{aligned}
 x_1 &= 0.50872011164692E 01 & v_1 &= -0.66871643066406E-02 \\
 x_2 &= 0.26847987C717E 01 & v_2 &= 0. \\
 x_3 &= 0.73304378896729E 01 & v_3 &= -0.10337829989844E-02
 \end{aligned}$$

CONTRAINTE

CONTRAINTE 1 C.0

$U(1) = 0.29315400123596E 00$

NIVEAU DE SORTIE	1	NOMBRE D ITERATIONS	9
NOMBRE D APPELS DU S/P	CALCULANT LES CONTRAINTES	167	
NOMBRE C APPELS DU S/P	CALCULANT LA FONCTION ECONOMIQUE	95	
NOMBRE D APPELS CL S/P	CALCULANT LE GRADIENT DE LA FONCTION ECONOMIQUE	19	
NOMBRE D APPELS CL S/P	CALCULANT L'INVERSE AVEC BASE DONNÉE	7	
NOMBRE C APPELS CU S/P	CALCULANT L'INVERSE AVEC RECHERCHE DE BASE	1	
NOMBRE D APPELS DU S/P	EFFECTUANT LES CHANGEMENTS DE BASE	1	
NOMBRE D APPELS DU S/P	CALCULANT LE JACOBIEN DES CONTRAINTES	12	
NOMBRE D APPELS DU S/P	EFFECTUANT UNE RECHERCHE DE MAX SANS DICHOTOMIE	8	
NOMBRE D APPELS DU S/P	EFFECTUANT UNE RECHERCHE DE MAX AVEC DICHOTOMIE	2	
NOMBRE D APPELS DU S/P	EFFECTUANT LA RENTREE DANS LE DOMAINE	31	
TOTAL DITERATIONS		144	

APPENDIX 5. GREG PROGRAMMING MATERIAL FOR CHAPTER SIX

This appendix contains the GREG external subroutines and four supporting subroutines used to optimize the water quality control model of Chapter six.

The four supporting subroutines are

1. MINDO - minimizes the dissolved oxygen function in a given stage by a Fibonacci search.
2. INPUT - reads in the supplementary parameter values, initializes variables, and determines the number of iterations the Fibonacci search must perform in each stage to guarantee a final search interval length = 0.0001.
3. FUNCT - a function that calculates the DO profile in each stage.
4. CHANGE- calculates the boundary values.

Definition of the Variables*

XT(I)	thermal input in BTU/hr
XQ(I)	flow in ft^3/sec
XB(I)	organic waste in lbs/day
XE(I)	equilibrium temperature in $^{\circ}\text{F}$
AK(I)	temperature dissipation rate coefficient in (days) $^{-1}$
AK1(I)	deoxygenation rate coefficient in (days) $^{-1}$ at $20\ ^{\circ}\text{C}$
AK2(I)	reaeration rate coefficient in (days) $^{-1}$ at $20\ ^{\circ}\text{C}$
ST(I)	incoming water temperature in $^{\circ}\text{F}$
SB(I)	incoming BOD in mg/l
SO(I)	incoming DO in mg/l
TF(I)	flow in days
TRS(I)	maximum allowable temperature rise in $^{\circ}\text{F}$
TMAX(I)	maximum allowable temperature in $^{\circ}\text{F}$
MDO(I)	minimum allowable DO in mg/l
TMIN	the flow time when DO is minimum, $0 \leq \text{TMIN} \leq \text{TF(I)}$
FMIN	minimum DO
NSTAGE	the number of stages
CC1	the conversion factor C_1
CC2	the conversion factor C_2

* - each subscripted variable ranges from $I = 1, \dots, \text{NSTAGE}$.

```

SUBROUTINE PHIX
DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)    GEGA  20
1,Y(150),C(150),IBAS(50),IFB(100),IVC(50),IVA(100)      GEGA  30
COMMON   A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IFB,IVC,IVA,IVB      GEGA  40
COMMON   NV,NC,NK,NEG,AIN,NTV,NV1,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGEGA  50
1N4,NVNINI,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCD8,KCGEGA  60
2DP,KFIL,KLIN,KREN,KD,FI1,PHI,PSI,PSI3,TB,TC,TC,EPSIL,EPSILO,EPSIL2GEGA  70
COMMON   KFCNC,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACO,KMAX1,KMAX2,KGEKA  80
1REN1,KREN2,KINV,KCDBA1,KREN11,KREN21      GEGA  90
COMMON   ICTREC,DELTIF1,ETA,JKO,LC,YSCRT      GEGA 100
COMMON/AREA1/XT(4),XQ(4),XB(4),XE(4),AK(4),AK1(4),AK2(4),ST(4)
1,SB(4),SO(4),TF(4),TRS(4),TMAX(4),MDC(4),NSTAGE
2,CC1,CC2,DELTA
REAL MDO
IF(IT.LT.0) CALL INPUT
PHI=0.0
PHI=PHI+(0.817/0.9)*(1.0-EXP(-0.023*XC(2)))
PHI=PHI+(0.575/0.9)*(1.0-EXP(-0.023*XC(4)))
IF(XC(5).LT.+2E-3) GO TO 10
IF(XC(5).LT.+3E-3) GO TO 11
IF(XC(5).LT.+4E-3) GO TO 12
IF(XC(5).LT.+7E-3) GO TO 13
IF(XC(5).LT.+9E-3) GO TO 14
IF(XC(5).LT.+9E-3) GO TO 15
PHI=PHI+0.98*YC(5)-91.2
GO TO 30
10 PHI=PHI+0.04*YC(5)
GO TO 30
11 PHI=PHI+0.03*YC(5)+0.2
GO TO 30
12 PHI=PHI+0.01*YC(5)+0.8
GO TO 30
13 PHI=PHI+(0.01/3.0)*XC(5)-(0.4/3.0)+1.2
GO TO 30
14 PHI=PHI+0.015*XC(5)+0.25
GO TO 30
15 PHI=PHI+0.06*YC(5)-3.8
3C PHI=-PHI
RETURN
END

```

```

SUBROUTINE GRDFI
DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)    GEGA  20
1,Y(150),C(150),VC(50),IBAS(50),IH8(100),IVC(50),IVA(100)    GEGA  30
COMMON   A,/LFA,X,XC,XI,XS,Y,C,VC,IBAS,IH8,IVC,IVA,IVB          GEGA  40
COMMON   NV,NC,NK,NEG,NIN,NTV,NVI,NEV,NEVL,NTO,NIN1,NIN2,NIN3,NIGEA  50
1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IPB,ICDB,JCDP,KCGEA  60
2DB,KFIL,KLIN,KREN,KC,FI1,PFI,PSI,PSI3,TB,TC,EPSIL,EPSIL0,EPSIL2GEA  70
COMMON   KFNC,KGRAD,KCONT,KINV1,KINV2,KCCBA,KJACC,KMAX1,KMAX2,KGEA  80
1REN1,KREN2,KINV,KCCBA1,KREN1I,KREN2I                         GEGA  90
COMMON   ICTREC,DELTFI,ETA,JKD,LC,YSCRT                         GEGA 100
COMMON/AREA1/XT(4),XQ(4),XB(4),XE(4),AK(4),AK1(4),AK2(4),ST(4)
1,SB(4),SD(4),TF(4),TRS(4),TMAX(4),MDC(4),NSTAGE
2,CC1,CC2,DELTA
REAL MDO
C(2)=-(0.817/F.9)*EXP(-0.023*XC(2))*C.023
C(4)=-(0.575/F.9)*EXP(-0.023*XC(4))*C.023
IF(XC(5).LT.20.0) GO TO 10
IF(XC(5).LT.30.0) GO TO 11
IF(XC(5).LT.40.0) GO TO 12
IF(XC(5).LT.70.0) GO TO 13
IF(XC(5).LT.90.0) GO TO 14
IF(XC(5).LT.95.0) GO TO 15
C(5)=-0.98
RETURN
10 C(5)=-0.04
RETURN
11 C(5)=-0.03
RETURN
12 C(5)=-0.01
RETURN
13 C(5)=-(0.01/3.0)
RETURN
14 C(5)=-0.015
RETURN
15 C(5)=-0.06
RETURN
END

```

```

SUBROUTINE CPNI
DIMENSION A(FO,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150) GEGA 20
1,Y(150),C(150),VC(50),IBAS(50),IH(100),IVC(50),IVAI(100) GEGA 30
COMMON A,PLFA,X,XC,XI,XS,Y,C,VC,IBAS,IH,IVC,IVA,IVB GEGA 40
COMMON NV,NC,NK,NEG,NIN,NTV,NVI,NEV,NEVL,NTC,NINI,NIN2,NIN3,NIGEGA 50
1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IP,ICDP,JCDB,KCGEGA 60
2CB,KFIL,KLIN,KREN,KD,FI1,PHI,PSI,PSI3,TB,TD,TC,EPSIL,EPSILO,EPSIL2GEGA 70
COMMON KFDNC,KGRAC,KCONT,KINV1,KINV2,KCCBA,KJACO,KMAX1,KMAX2,KGEGA 80
IREN1,KREN2,KIFV,KLDBA1,KREN11,KREN21 GEGA 90
COMMON IDIREC,DELTFI,ETA,JKO,LC,YSCRT GEGA 100
COMMON/AREA1/YT(4),XQ(4),XB(4),XE(4),AK(4),AK1(4),AK2(4),ST(4)
1,SB(4),SO(4),TF(4),TRS(4),TMAX(4),MCC(4),NSTAGE
2,CC1,CC2,DELTA
REAL MDO

C.....INITIALIZE
IF(IT.GE.0) GO TO 100
DO51=1,NSTAGE
  XB(I)=(CC2*XE(I))/(100.0*XC(I))
  XT(I)=(CC1*XT(I))/(100.0*XC(I))
5 CONTINUE
C.....END OF INITIALIZATION
C.....SET UP BOUNDARY CONDITIONS
100 CALL CHANGE(1,NSTAGE-1)
C.....END OF BOUNDARY CONDITIONS
C.....INEQUALITY CONSTRAINTS
C.....I SUBSCRIPTS...MAX TEMP RISE
C.....NSTAGE+I SUBSCRIPTS...MAX TEMP
C.....2*NSTAGE+I SUBSCRIPTS...MIN CC
DO101=1,NSTAGE
  VC(I)=XT(I)*(100.0-XC(I))-TRS(I)
  VC(NSTAGE+I)=ST(I)+XT(I)*(100.0-XC(I))-TMAX(I)
  CALL MINDO(I,TMIN,FMIN)
  VC(2*NSTAGE+I)=-FMIN+MDO(I)
10 CONTINUE
RETURN
END

```

```

SUBROUTINE JACOB
DIMENSION A(FO,130),ALFA(50,50),X(150),XC(150),XI(150),XS(150)      GEGA  20
1,Y(150),C(150),VC(50),IBAS(50),IHB(100),IVC(50),IVA(100)      GEGA  30
COMMON   A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVA,IVB      GEGA  40
COMMON   NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NINI,NIN2,NIN3,NICEGA  50
1N4,NVAIN1,NVNTN2,NVNIN3,INDEX,II,IR,IP1,IS,IS1,IT,IBP,ICDB,JCDP,KCGEGA  60
2DB,KFIL,KLIN,KREN,KE,FII,PHI,PSI,PSI3,TB,TC,TC,EPSIL,EPSILO,EPSIL2GEGA  70
COMMON   KFCNG,KGRAC,KCCT,KINV1,KINV2,KCDBA,KJACC,KMAX1,KMAX2,KGEGA  80
1RE41,KREN2,KINV,KCDBA1,KREN11,KRE421      GEGA  90
COMMON   ICTREC,DELTIF1,ETA,JKO,LC,YSCRT      GEGA 100
COMMON/AREA1/XT(4),XQ(4),XB(4),XE(4),AK(4),AK1(4),AK2(4),ST(4)
1,SB(4),SO(4),TF(4),TRS(4),TMAX(4),MDC(4),NSTAGE
2,CC1,CC2,DELTA
REAL MDO
C.....INITIALIZE CONSTANTS
IF(IT.GE.0) GO TO 100
C0100J=1,NSTAGE
AXXX=0.0
K=J
50 A(NSTAGE+J,K)=-XT(K)*EXP(AXXX)
K=K-1
IF(K.LE.0) GO TO 55
AXXX=AXXX-AK(K)*TF(K)
GO TO 50
55 A(J,J)=-XT(J)
60 CONTINUE
C.....END OF INITIALIZATION
100 CALL CHANGE(1,NSTAGE-1)
C.....NUMERICAL PARTIAL DERIVATIVES FOR THE MIN DO CONSTRAINT
C030I=1,NSTAGE
CALL MINDO(I,TMIN,FMIN)
IF(XT(I).EQ.0.0) GO TO 5
XXCC=XC(I)
XC(I)=XC(I)+DELT
CALL MINDO(I,TMIN1,FMIN1)
A(2*NSTAGE+I,I)=(FMIN-FMIN1)/DELTA
XC(I)=XXCC
5 IF(XB(I).EQ.0.0) GO TO 6
XXCC=XC(NSTAGE+I)
XC(NSTAGE+I)=YC(NSTAGE+I)+DELTA
CALL MINDO(I,TMIN1,FMIN1)
A(2*NSTAGE+I,NSTAGE+I)=(FMIN-FMIN1)/DELTA
XC(NSTAGE+I)=XXCC
6 IF(I.EC.1) GO TO 30
III=I-1
C020J=1,III
IF(XT(J).EQ.0.0) GO TO 7
XXCC=XC(J)
XC(J)=XC(J)+DELT
CALL CHANGE(J,III)
CALL MINDO(I,TMIN1,FMIN1)
A(2*NSTAGE+I,J)=(FMIN-FMIN1)/DELTA
XC(J)=XXCC
7 IF(XB(J).EQ.0.0) GO TO 8
XXCC=XC(NSTAGE+J)
XC(NSTAGE+J)=YC(NSTAGE+J)+DELTA
CALL CHANGE(J,III)
CALL MINDO(I,TMIN1,FMIN1)

```

```
A(2*NSTAGE+I,NSTAGE+J)=(FMIN-FMIN1)/DELTA
XC(NSTAGE+J)=XXCC
E IF(XT(J).EQ.0.0.AND.XB(J).EQ.0.0) GC TC 20
CALL CHANGE(J,III)
2C CONTINUE
3C CONTINUE
RETURN
END
```

```

SUBROUTINE CHANGE(J,N)
DIMENSION A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150)    GEGA  20
1,Y(150),C(150),VC(50),IBAS(50),IHB(100),IVC(50),IVA(100)      GEGA  30
COMMON   A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IHB,IVC,IVA,IVB      GEGA  40
COMMON   NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NINI,NIN2,NIN3,NIGEGA  50
IN4,NVNIN1,NVNIN2,NVNIN3,INCEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,Jcdb,KCGEGA  60
2DB,KFIL,KLIN,YREN,KC,FI1,PHI,PSI,PSI3,TB,TC,EPSIL,EPSILO,EPSIL2GEGA  70
COMMON   KFCNC,KGRAC,KCONT,KINV1,KINV2,KCDBA,KJACC,KMAX1,KMAX2,KGEA  80
IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21                           GEGA  90
COMMON   ICIREC,DELTIF,ETA,JKO,LC,YSORT                           GEGA 100
COMMON/AREA1/XT(4),XQ(4),XB(4),XE(4),AK(4),AK1(4),AK2(4),ST(4)
1,SB(4),SO(4),TF(4),TRS(4),TMAX(4),MDC(4),NSTAGE
2,CC1,CC2,DELTA
REAL MDC
C.....FUNCTION DEFINITIONS
TEMPC(T)=(5.0/9.0)*(T-32.0)
CS(T)=14.652-0.41022*TEMPC(T)+0.79910E-2*TEMPC(T)**2-0.77774E-4
1*TEMPC(T)**3
C.....END OF FUNCTION DEFINITIONS
C.....SET UP BOUNDS FOR STAGES J+1 TO N+1
DO20I=J,N
TMIX=ST(I)+XT(I)*(100.0-XC(I))
ST(I+1)=XE(I)+(TMIX-XE(I))*EXP(-AK(I)*TF(I))
BDC=SB(I)+XB(I)*(100.0-XC(NSTAGE+I))
AAK1=AK1(I)*(1.047)**(TEMPc(ST(I+1))-2.0)
SB(I+1)=BDC*EXP(-AAK1*TF(I))
AAK2=AK2(I)*(1.024)**(TEMPc(ST(I+1))-2.0)
DODEF=CS(TMIX)-SO(I)
IF(DODEF<0.0) DODEF=0.0
SO(I+1)=CS(ST(I+1))-(AAK1*BDC/(AAK2-AAK1))*(EXP(-AAK1*TF(I))
1-EXP(-AAK2*TF(I)))-DODEF*EXP(-AAK2*TF(I))
20 CONTINUE
RETURN
END

```

```

SUBROUTINE MINDC(I,TMIN,FMIN)
DIMENSION A(50,150),ALFA(50,50),X(150),XC(150),XI(150),XS(150)    GEGA  20
1,Y(150),C(150),VC(50),IBAS(50),IH8(150),IVC(50),IVA(150)      GEGA  30
COMMON A,ALFA,X,XC,XI,XS,Y,C,VC,IBAS,IH8,IVC,IVA,IVR      GEGA  40
COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTO,NIN1,NIN2,NIN3,NICEGA 50
1N4,NVNIN1,NVNIN2,NVNIN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDB,JCCB,KCGEGA 60
2DB,KFIL,KLIN,KREN,KD,FI1,PPI,PSI,PSI3,TB,TC,TC,EPSIL,PSILO,EPSIL2GEGA 70
COMMON KFCNC,KGRAC,KCONT,KINV1,KINV2,KCDBA,KJACC,KMAX1,KMAX2,KCGEGA 80
1RENI,KREN2,KINV,KCDBA1,KREN11,KREN21      GEGA  90
COMMON IDIREC,DELTIF,ETA,JKO,LC,YSORT      GEGA 100
COMMON/AREA1/XT(4),XQ(4),XB(4),XE(4),AK(4),AK1(4),AK2(4),ST(4)
1,SB(4),SO(4),TF(4),TRS(4),TMAX(4),MDC(4),NSTAGE
2,CC1,CC2,DELTI
COMMON/AREA2/NTIMES(4),DEL(4),NF(30)
REAL NDC
C.....FIBONNACI SEARCH FOR THE MIN EC IN STAGE I
XA=0.0
XBR=TF(I)
X1=XA+DEL(I)*(XBB-XA)
X2=XA+XBB-X1
Y1=FUNCT(I,X1)
Y2=FUNCT(I,X2)
YA=FUNCT(I,XA)
YE=FUNCT(I,XE)
C.....CHECK FOR MINIMUM AT ECUNES
FMIN=4MIN1(Y1,Y2,YA,YE)
IF(FMIN.EQ.YA) GO TO 100
IF(FMIN.EQ.YE) GO TO 200
N=NTIMES(I)
DO300J=1,N
IF(Y1.LT.Y2) GO TO 10
XBR=X1
X1=X2
Y1=Y2
X2=XA+XBB-X1
Y2=FUNCT(I,X2)
GO TO 300
10 XA=X2
X2=X1
Y2=Y1
X1=XA+XBR-X2
Y1=FUNCT(I,X1)
300 CONTINUE
IF(Y1.LT.Y2) GO TO 20
TMIN=X2
FMIN=Y2
RETURN
20 TMIN=X1
FMIN=Y1
RETURN
100 TMIN=0.0
RETURN
200 TMIN=TF(I)
RETURN
END

```

```

FUNCTION FUNCT(I,TTT)
CIMENSI0N A(50,100),ALFA(50,50),X(150),XC(150),XI(150),XS(150) GEGA 20
1,Y(150),C(150),VC(50),IRAS(50),IHB(100),IVC(50),IVA(100) GEGA 30
COMMON A,PLFA,X,XC,XI,XS,Y,C,VC,IPAS,IHB,IVC,IVA,IVB GEGA 40
COMMON NV,NC,NK,NEG,NIN,NTV,NV1,NEV,NEVL,NTC,NIN1,NIN2,NIN3,NIGEGA 50
1N4,NVAIN1,NVNTN2,NVNTN3,INDEX,II,IR,IR1,IS,IS1,IT,IBP,ICDR,JCDR,KCGEGA 60
2DB,KFIL,KLIN,KREN,KD,FI1,PFI,PSI,PSI3,TB,TC,EP0L,EP0L0,EP0L2GEGA 70
COMMON KFCNG,KGRAD,KCONT,KINV1,KINV2,KCDBA,KJACC,KMAX1,KMAX2,KGEWA 80
IREN1,KREN2,KINV,KCDBA1,KREN11,KREN21 GEGA 90
COMMON IC7REC,DELTFL,ETA,JKD,LC,YSCRT GEGA 100
COMMON/AREAI/XT(4),XQ(4),XB(4),XE(4),AK(4),AK1(4),AK2(4),ST(4)
1,SR(4),SD(4),TF(4),TRS(4),TMAX(4),M00(4),NSTAGE
2,CC1,CC2,DELTA
REAL MDO
C.....FUNCTION DEFINITIONS
TEMPC(T)=(5.0/9.0)*(T-32.0)
CS(T)=14.652-0.41022*TEMPc(T)+0.79910E-2*TEMPc(T)**2-0.77774E-4
1*TEMPc(T)**3
C.....END OF FUNCTION DEFINITIONS
TMIX=ST(I)+XT(I)*(100.C-XC(I))
STI1=XE(I)+(TMIX-XE(I))*EXP(-AK(I)*TTT)
B0D=SD(I)+XB(I)*(100.C-XC(NSTAGE+I))
AAK1=AK1(I)*(1.047)**(TEMPc(STI1)-20.C)
AAK2=AK2(I)*(1.024)**(TEMPc(STI1)-20.C)
DODEF=CS(TMIX)-SD(I)
IF(DODEF.LT.0.0) DODEF=0.0
FUNCT=CS(STI1)-(AAK1*B0D/(AAK2-AAK1))*(EXP(-AAK1*TTT)
1-EXP(-AAK2*TTT))-DODEF*EXP(-AAK2*TTT)
RETURN
END

```

```

SUBROUTINE INPUT
COMMON/AREA1/XT(4),XQ(4),XB(4),XE(4),AK(4),AK1(4),AK2(4),ST(4)
1,SR(4),SO(4),TF(4),TRS(4),TMAX(4),MDO(4),NSTAGE
2,CC1,CC2,DELTA
COMMON/AREA2/NTIMES(4),DEL(4),NF(30)
REAL MDO
C.....NUMBER OF STAGES
NSTAGE=4
READ(5,1000) XT,XQ,XB,XE,AK,AK1,AK2,ST,SB,SC,TF,TRS,TMAX,MDO
READ(5,2000) DELTA
C.....CONVERSION FACTORS
CC1=1.0/(3.6E7*62.4262)
CC2=0.4535924E6/(8.64E4*28.31605)
C.....DETERMINE NUMBER OF ITERATIONS FOR FIBONACCI SEARCH
NF(1)=1
NF(2)=1
DO70I=1,NSTAGE
DO50J=3,30
NTIMES(I)=J-1
NF(J)=NF(J-1)+NF(J-2)
AF1=NF(J-1)
AF=NF(J)
AF=1.0/AF
IF((AF*TF(I)).LE.0.0001) GO TO 70
50 CONTINUE
70 CEL(I)=AF1*AF
WRITE(6,3000) XT,XQ,XB,XE,AK,AK1,AK2,ST,SB,SO,TF,TRS,TMAX,MDO
WRITE(6,4000) NSTAGE,DELTA,CC1,CC2
WRITE(6,5000) NTIMES,DEL
RETURN
1000 FORMAT(4G10.6)
2000 FORMAT(F10.0)
3000 FORMAT(4(5X,E13.6))
4000 FORMAT(5X,I5,?(5X,E13.6))
5000 FORMAT(4(5X,I5)/4(5X,E13.6))
END

```

PARAMETRES

NV	8
NIN	12
NEG	0
NEVL	20
NTO	6
ITET	20
ICONJ	1
IDIAG	C
ITMAX	50
KFIT	0
KLIN	C
NCO	20
ITSCR	1
ISOLSR	50
EPSIL	0.1E 00
EPSILO	0.1E-02
EPSIL1	0.1E-02
EPSIL2	0.1E-02
PC	0.0
0*C	0.24C000E 10
0*820C00E 03	0.870060E C3
0*110C00E 06	1*0
0*807C00E 02	0*8C7000E C2
0*124C00E 01	0*129000E C1
0*100C00E 01	0*13C030E C1
0*900C00E C0	0*900C00E C0
0*807000E 02	0*0
0*20000CE 01	0*0
0*670000E C1	0*0
0*500000E-01	0*187C00E C1
0*100C00E 02	0*1C0000E C2
0*93CC00E 02	0*93CC00E C2
0*300000E 01	0*3CC000E C1
4	0*100000E-03
14	22
0*618033E 00	0*618034E C0
	0.444970E-C5
	22
	0.616034E 00
	0.618034E C0

GRADIENT RÉDUIT GÉNÉRALISÉ

NCMBRE DE VARIABLES NATURELLES	6
NCMBRE TOTAL DE VARIABLES	24
NCMBRE DE CONTRAINTES	12
EPSILON DE NEWTON	0.1000E-02
EPSILON TEST GRADIENT	0.1000E-02

FONCTION ÉCONOMIQUE 0.0

ÉCRAN INFÉRIEUR

	VARIABLE	NATURELLE	BCRNE SUPERIEURE
XI(1)	x(1)	0.0	x(1)
XI(2)	x(2)	0.0	x(2)
XI(3)	x(3)	0.0	x(3)
XI(4)	x(4)	0.0	x(4)
XI(5)	x(5)	0.0	x(5)
XI(6)	x(6)	0.0	x(6)
XI(7)	x(7)	0.0	x(7)
XI(8)	x(8)	0.0	x(8)

VALEUR DES CONTRAINTES

C1 1)	-C.100000000000E+00	02
C1 2)	C.22750282287E98E	01
C1 3)	-C.100000000000E+00	02
C1 4)	-C.22613983154797E	01
C1 5)	-C.12357093051758E	02
C1 6)	-C.24978637695213E-01	
C1 7)	-C.057982798E	02
C1 8)	-C.36490631103516E	01
C1 9)	-C.20126600265603E	01
C1 10)	C.71467208862305E	01
C1 11)	C.30198316574C97E	01
C1 12)	C.221C1259231567E	01

IT 1	PHI-0.775471E75000CE 04	DIR.GRAD.	NC 1C	YN C.34E 03	DELTAFI 0.37E 05	ETA 0.33E 05	NCD B 0	NCN 2
IT 2	PHI-J.58594921875CC0E 04	DIR.CONJ.	LC 2				NCD B 0	NCN 2
IT 3	PHI-C.46764521875CC0E C4	DIR.CCNJ.	LC 2				NCD B 0	NCN 2
IT 4	PHI-U.4325335937500CE 04	DIR.GRAD.	NC 1C	YN 0.29E 03	DELTAFI 0.26E 05	ETA 0.16E 05	NCD B 0	NCN 2
IT 5	PHI-O.294J1230566406E 04	DIR.GRAD.	NC 11	YN C.27E 04	DELTAFI 0.26E 05	ETA 0.27E 14	NCD B 1	NCN 2
IT 6	PHI-G.96953G27343750E 03	DIR.GRAD.	NO 11	YN C.16E 04	DELTAFI 0.26E 05	ETA 0.16E 14	NCD B 1	NCN 4
IT 7	PHI-O.16048212U51392E 01	DIR.GRAD.	NC 11	YN 0.12E 03	DELTAFI 0.11E 05	ETA 0.11E 05	NCD B 1	NCN 3
LES VARIABLES ARTIFICIELLES SONT TOUTES ANNULÉES	PHI-G.16048212C51392E 01		YN 0.0	DELTAFI 0.0		ETA 0.0	NCD B 1	NCN 2

DURÉE CU CALCUL 3035 CENTISECONDES

VARIABLE NATURELLE	VARIABLE DUALE ASSOCIEE
$x_1 = 0.0$	$v_1 = 0.0$
$x_2 = 0.8541931152344E-02$	$v_2 = 0.0$
$x_3 = 0.0$	$v_3 = 0.0$
$x_4 = 0.0$	$v_4 = -0.146944411C9856E-01$
$x_5 = 0.66895965576172E-02$	$v_5 = 0.0$
$x_6 = 0.0$	$v_6 = 0.0$
$x_7 = 0.0$	$v_7 = 0.0$
$x_8 = 0.0$	$v_8 = 0.0$

CONTRAINTE	VARIABLE DUALE ASSOCIEE
C ₁	$U_1 = 0.0$
C ₂	$U_2 = 0.1111C3844642639E+00$
C ₃	$U_3 = 0.0$
C ₄	$U_4 = 0.0$
C ₅	$U_5 = 0.0$
C ₆	$U_6 = 0.0$
C ₇	$U_7 = 0.0$
C ₈	$U_8 = 0.0$
C ₉	$U_9 = 0.0$
C ₁₀	$U_{10} = 0.34952521324158E-01$
C ₁₁	$U_{11} = 0.0$
C ₁₂	$U_{12} = 0.0$

NIVEAU	DÉ SORTIE	1	NOMBRE D'ITERATIONS	7	NOMBRE D'appels	DL	S/P	CALCULANT LES CONTRAINTES	75
NCMBRE	D APPELS	DL	S/P	CALCULANT LA FONCTION ECONOMIQUE	39	NCMBRE	D APPELS	DL	S/P
NCMBRE	D APPELS	DL	S/P	GRADIENT DE LA FONCTION ECONOMIQUE	15	NCMBRE	D APPELS	DL	S/P
NCMBRE	D APPELS	DL	S/P	L'INVERSE AVEC BASE DONNÉE	4	NCMBRE	D APPELS	DL	S/P
NCMBRE	D APPELS	DL	S/P	L'INVERSE AVEC RECHERCHE DE BASE	0	NCMBRE	D APPELS	DL	S/P
NCMBRE	D APPELS	DL	S/P	EFFECTUANT LES CHANGEMENTS DE BASE	4	NCMBRE	D APPELS	DL	S/P
NCMBRE	D APPELS	DL	S/P	CALCULANT LE JACOBIEN DES CONTRAINTES	9	NCMBRE	D APPELS	DL	S/P
NCMBRE	D APPELS	DL	S/P	EFFECTUANT UNE RECHERCHE DÉ MAX SANS DICHOTOMIE	7	NCMBRE	D APPELS	DL	S/P
NCMBRE	D APPELS	DL	S/P	EFFECTUANT UNE RECHERCHE DE MAX AVEC DICHOTOMIE	0	NCMBRE	D APPELS	DL	S/P
				DENTREE DANS LE DOMAINE	17				NOMBRE TOTAL D'ITERATIONS 60

PACAMETTES

GRADIENT REDUIT GENERALISE

NUMBER DE VARIABLES NATURELLES	8
NUMBER TOTAL DE VARIABLES	20
NUMBER DE CONTRAINTES	12
EPSILON DE NEWTON	0.100E-02
EPSILON TEST GRADIENT	0.100E-02

FUNCTION ECONOMIQUE -G.81915740966797E 01

	BORNE INFERIEURE		VARIABLE NATURELLE	BORNE SUPERIEURE
	X(1)	X(2)		
X(1)	0.0	x(1)	0.0	x(1)
X(2)	0.0	x(2)	0.100000000000E 03	x(2)
X(3)	0.0	x(3)	0.0	x(3)
X(4)	0.0	x(4)	0.100000000000E 03	x(4)
X(5)	0.0	x(5)	0.100000000000E 03	x(5)
X(6)	0.0	x(6)	0.0	x(6)
X(7)	0.0	x(7)	0.0	x(7)
X(8)	0.0	x(8)	0.0	x(8)

VALEUR DES CONTRAINTES

C(1)	-C.100000000000E 02
C(2)	-C.100000000000E 02
C(3)	-C.100000000000E 02
C(4)	-C.100000000000E 02
C(5)	-C.12300003051758E 02
C(6)	-C.12300003051758E 02
C(7)	-C.12300003051758E 02
C(8)	-C.12300003051758E 02
C(9)	-C.36501507949E 29E 01
C(10)	-C.34234899875488E 01
C(11)	-C.3196732521C571E 01
C(12)	-C.32800188064E75E 01

IT 1	PHI-C.29049C5E914J85E 01	DIR-GRAD.	NO 5	YN C.98E CO	CELTAF1 0.98E 02	ETA C.98E 02	NCDB 0	NCN 2	NITN 2
IT 2	PHI-C.2676637649E361E 01	DIR-CONJ.	LC 2				NCDB 0	NCN 6	NITN 6
IT 3	PHI-C.26742123167725E 01	DIR-CONJ.	LC 2				NCDB 0	NCN 2	NITN 2
IT 4	PHI-C.22251167297362E 01	DIR-CONJ.	LC 2				NCDB 2	NCN 10	NITN 45
IT 5	PHI-D.16048C78536987E 01	DIR-GRAD.	NC 6	YN C.11E-C1	DELTAF1 0.55E 00	ETA 0.32E 00	NCDB 1	NCN 7	NITN 26
	PHI-D.16048C78536987E 01			YN C.C	DELTAF1 C.0	ETA 0.0			

VARIABLE	NATURELLE	VARIABLE	DUALE	ASSOCIEE
X(1)	= C*J	= X(1)	= V(1)	= 0.0
X(2)	= C*18541931152344E 02	= X(2)	= V(2)	= 0.
X(3)	= C*0	= X(3)	= V(3)	= 0.0
X(4)	= C*2	= X(4)	= V(4)	= -0.14694441109896E-01
X(5)	= C*66891937255859E 02	= X(5)	= V(5)	= 0.
X(6)	= C*0	= X(6)	= V(6)	= 0.0
X(7)	= C*0	= X(7)	= V(7)	= 0.0
X(8)	= C*0	= X(8)	= V(8)	= 0.0

CONTRAINTE				
CC	CONTRAINTE	VARIABLE	DUALE	ASSOCIEE
CC	CONTRAINTE 1	-C*19000000000000E 02	U(1)	= 0.0
CC	CONTRAINTE 2	-C*95945C68359575E-03	U(2)	= 0.1103844642639E 00
CC	CONTRAINTE 3	-C*10000000000000E 02	U(3)	= 0.0
CC	CONTRAINTE 4	-C*2261398315427E 01	U(4)	= 0.0
CC	CONTRAINTE 5	-C*123011030517F8E 02	U(5)	= 0.0
CC	CONTRAINTE 6	-C*230101C1318259E C1	U(6)	= 0.0
CC	CONTRAINTE 7	-C*11404306958C5BE 02	U(7)	= 0.0
CC	CONTRAINTE 8	-C*381823730466875E 01	U(8)	= 0.0
CC	CONTRAINTE 9	-C*309462119506P4_ C1	U(9)	= 0.0
CC	CONTRAINTE 10	-C*57029724121C94E-03	U(10)	= 0.34952521324158E-01
CC	CONTRAINTE 11	-C*11567068099976E 01	U(11)	= 0.0
CC	CONTRAINTE 12	-C*14771C514066f0E 01	U(12)	= 0.0

NIVEAU	DE	SCRIPTÉ	?	?				
NOMBRE	DE	ITERATIONS	5					
NOMBRE	D APPELS	CL	S/P	CALCULANT	LES CONTRAINTES	105		
NOMBRE	D APPELS	CL	S/P	CALCULANT	LA FONCTION ÉCONOMIQUE	57		
NOMBRE	D APPELS	CL	S/P	CALCULANT	LE GRADIENT DE LA FONCTION ÉCONOMIQUE	21		
NOMBRE	D APPELS	CL	S/P	CALCULANT	L'INVERSE AVEC BASE DONNÉE	4		
NOMBRE	D APPELS	CL	S/P	CALCULANT	L'INVERSE AVEC RECHERCHE DE BASE	0		
NOMBRE	D APPELS	CL	S/P	EFFECTUANT	LES CHANGEMENTS DE BASE	3		
NOMBRE	D APPELS	CL	S/P	CALCULANT	LE JACOBIEN DES CONTRAINTES	9		
NOMBRE	D APPELS	CL	S/P	EFFECTUANT	UNE RECHERCHE DE MAX SANS DICHOTOMIE	5		
NOMBRE	D APPELS	CL	S/P	EFFECTUANT	UNE RECHERCHE DE MAX AVEC DICHOTOMIE	0		
NOMBRE	D APPELS	CL	S/P	EFFECTUANT LA RENTREE DANS LE DOMAINE	27			
				NOMBRE TOTAL D'ITERATIONS	81			

OPTIMIZATION OF INDUSTRIAL SYSTEMS WITH
THE SEPARABLE PROGRAMMING AND
THE GENERALIZED REDUCED GRADIENT METHODS

by

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B.A., KANSAS STATE TEACHERS COLLEGE, EMPORIA, 1970

AN ABSTRACT OF A MASTER'S THESIS

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requirements for the degree

MASTER OF SCIENCE

Department of Industrial Engineering

KANSAS STATE UNIVERSITY
Manhattan, Kansas

1972

This thesis reviews two nonlinear programming methods and their application to industrial systems. Separable programming is described and exemplified. The same procedure is presented that is used in the separable programming subroutine of the Mathematical Programming System/360 (MPS/360). A supplemental FORTRAN program for assisting the usage of MPS/360 when solving separable programming problems is presented. An interesting application of separable programming to solve the geometric programming dual problem with N-degrees of difficulty is also exhibited. The second technique considered is the generalized reduced gradient method. The mathematical theory is presented and numerical examples are used to exemplify it. Its application via the GREG program is evaluated, and numerous computer examples are worked. The study is concluded with the application of the technique to optimize a water quality control model.